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CSC3100 Data Structures

Tutorial 6: Queue

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- Queue Concept
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Queue

First-In, First-Out (FIFO)

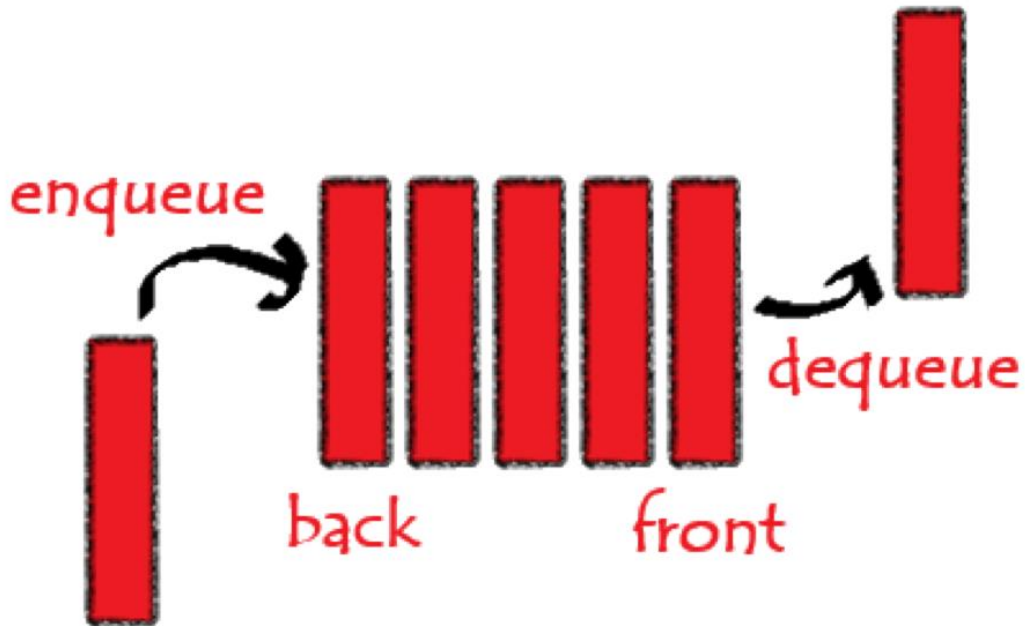


Figure: Queue analogy

Basic operations

- enqueue ()
 - Add elements to the back of the queue.
- dequeue ()
 - Remove elements from the front of the queue.



Double-ended Queue

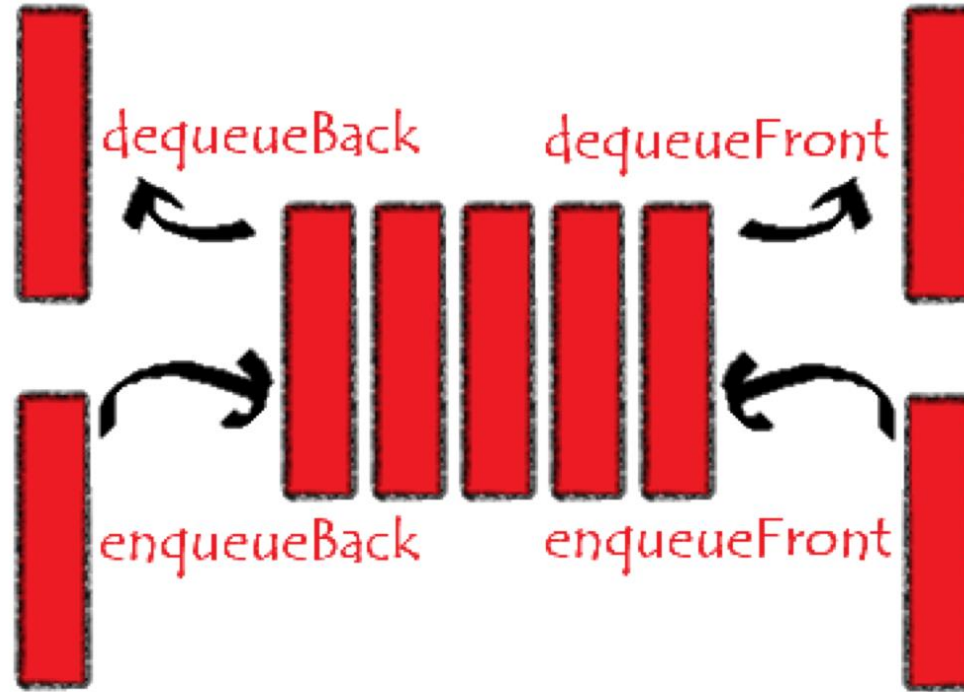


Figure: Double-ended Queue analogy

Double-ended queue is a more flexible queue that allows enqueue and dequeue at both ends.



Circular Queue

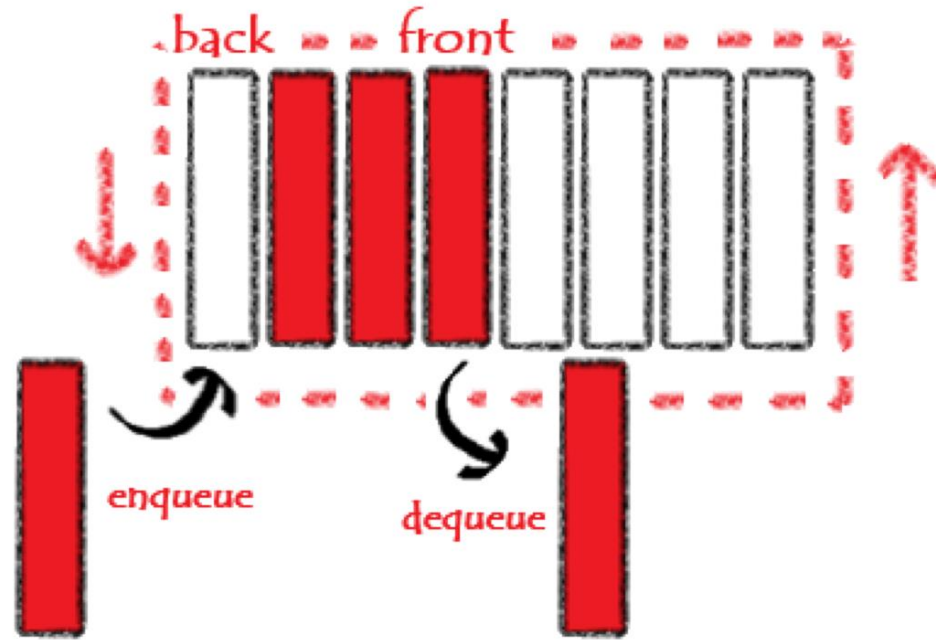
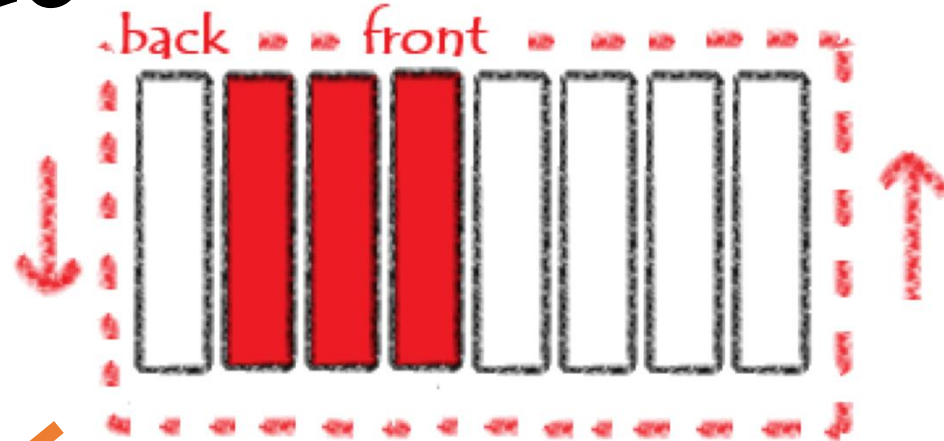


Figure: Circular Queue analogy

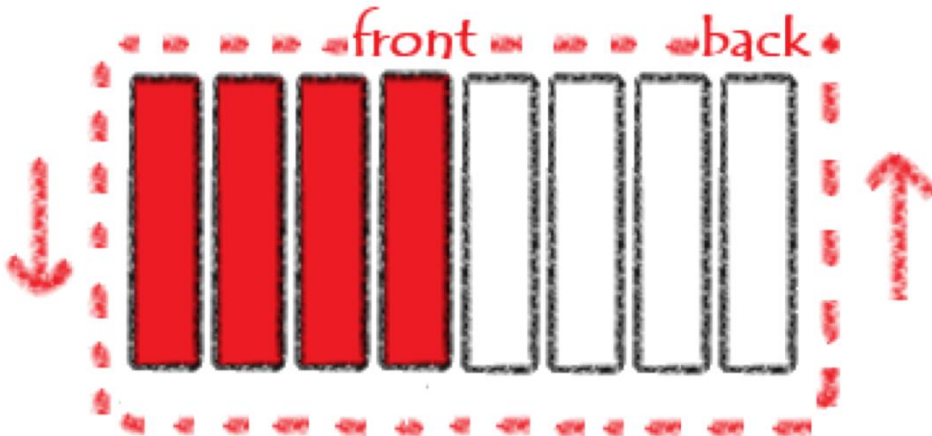
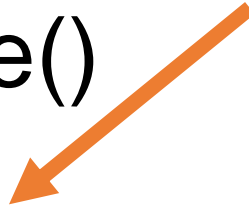
A queue in a fixed-size circular buffer, making it as if it were connected end-to-end. With fixed capacity but it is memory efficient.



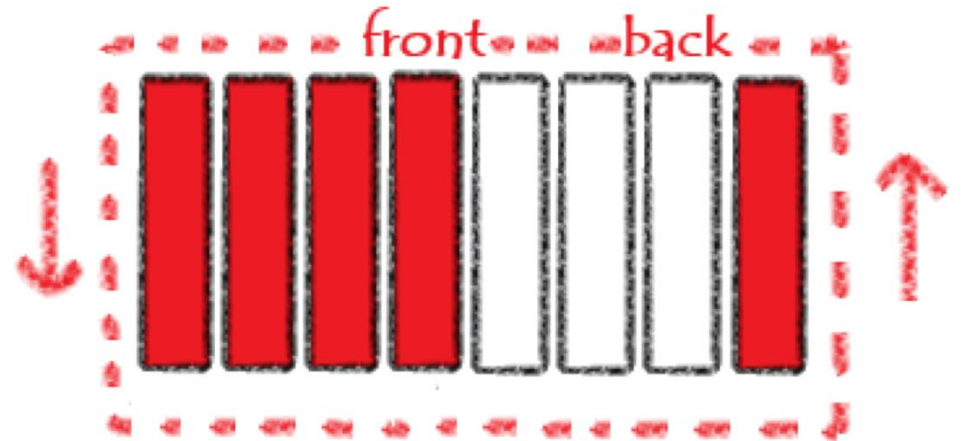
Circular Queue



enqueue()



enqueue()





Priority Queue

A queue where elements are processed based on priority.
Unlike the standard queue, it does not necessarily
implements a First-In, First-Out, policy.

Will be covered in more detail in Tutorial about heap.



Exercise 1: Implement stack using queues(LeetCode P225)

Implement a last-in-first-out (LIFO) stack using only two queues.



Exercise 2: Sliding Window Maximum(LeetCode P239)

Given an array of integers, there is a sliding window of size k which is moving from the very left of the array to the very right. You can only see the k numbers in the window. Each time the sliding window moves right by one position.

Example 1:

Input: `nums = [1,3,-1,-3,5,3,6,7]`, `k = 3`

Output: `[3,3,5,5,6,7]`

Explanation:

| Window position | Max |
|---------------------|-------|
| ----- | ----- |
| [1 3 -1] -3 5 3 6 7 | 3 |
| 1 [3 -1 -3] 5 3 6 7 | 3 |
| 1 3 [-1 -3 5] 3 6 7 | 5 |
| 1 3 -1 [-3 5 3] 6 7 | 5 |
| 1 3 -1 -3 [5 3 6] 7 | 6 |
| 1 3 -1 -3 5 [3 6 7] | 7 |



Solution 2: Sliding Window Maximum(LeetCode P239)

Method 1: Brute-Force

For a sequence with length n , there exists $n-k+1$ windows, the time complexity to find the maximum value in a window is $O(k)$. The overall time complexity is $O(nk)$.

We could save the time by **double-end queue** to reduce the time complexity of finding the maximum value in a window to $O(1)$.



Solution 2: Sliding Window Maximum(LeetCode P239)

Method 2: Monotonic queue

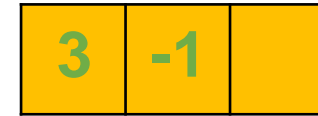
Key observations:

- If the **current** element is **greater than or equal to** the element at the **end** of the data structure, then the latter no longer influences the maximum of future windows; hence, it **can be discarded**.
- If the element at the **beginning** of the data structure and the **current** element is k **distance away**, then the former no longer influences the maximum of future windows; hence, it **can also be discarded**.
- From the beginning to the end of the data structure, the elements are **sorted in descending order**; hence, the **maximum lies** at the **beginning** of the data structure.



Solution 2: Sliding Window Maximum(LeetCode P239)

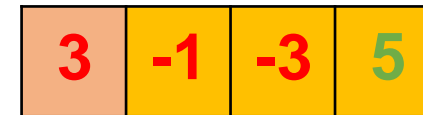
Method 2: Monotonic queue



3



3



5



5



6



7



Solution 2: Sliding Window Maximum(LeetCode P239)

Method 2: Monotonic queue

Initialization:

Define a double-ended queue “deque”, a result list “res”, and the length of the array “n”.

Sliding Window:

The left boundary ranges from $i = 1 - k$ to $i = n - k$;

The right boundary ranges from $j = 0$ to $j = n - 1$.

- (a) If $i > 0$ and the front element of the `deque` is equal to the deleted element `nums[i - 1]`, then dequeue the front element.
- (b) Remove all elements in the `deque` that are less than $\text{nums}[j]$
- (c) Add $\text{nums}[j]$ to the end of the `deque`.
- (d) If a window is formed (i.e., $i \geq 0$): Add the maximum value of the window (i.e., the front element of the `deque`) to the result list `res`.

Return value: Return the result list `res`.



Thanks for your Attention!

Q&A