

CSC3100 Data Structures Mid-term Exam

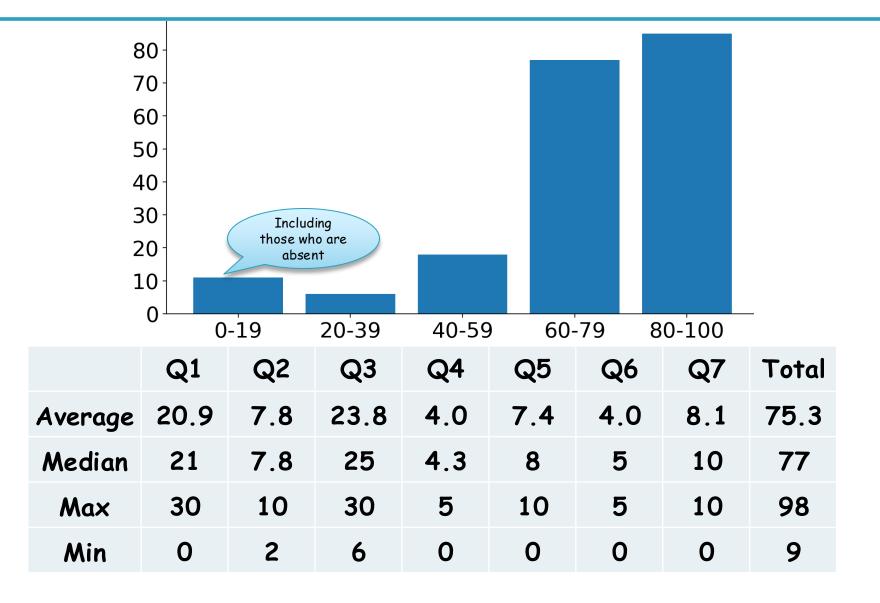
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#1 Multiple Choice Questions	(30)
> #2 True or False	(10)
#3 Short Answer Questions	(30)
#4 Doubly Linked List	(5)
> #5 Read Program	(10)
#6 Find Second Largest Value	(5)
#7 Reverse First K Elements in a Queue	(10)



Grade Distribution





- (1). Which of the following is a linear data structure?
 - A. Array
 - B. Linked List
 - C. Stack
 - D. All of the above
- (2). What is the worst-case time complexity of accessing an element in an array with n elements?
 - A. O(1)
 - B. $O(\log n)$
 - C. O(n)
 - D. $O(n^2)$
- (3). Which of the following is NOT a valid operation on an array?
 - A. Accessing an element at a given index.
 - B. Appending an element to the end.
 - C. Deleting an element at an arbitrary position without shifting.
 - D. Finding the maximum element in O(1) time (unsorted array).



- (4). In a singly linked list with **n** elements, what is the time complexity of inserting a new node at the beginning?
 - A. O(1)
 - B. $O(\log n)$
 - C. O(n)
 - D. $O(n^2)$
- (5). What is the time complexity of the following recursive function?

```
public void func(int n) {
    if (n <= 1) return;
    func(n/2);
    func(n/2);
}</pre>
```

- A. O(n)
- B. $O(n^2)$
- C. $O(n \log n)$
- D. $O(\log n)$.

- $g(n) = 2g(\frac{n}{2}) + O(1)$ and let g(1) = b
- When $n = 2^x$, we have that: $g(n) \le 2g\left(\frac{n}{2}\right) + c \le 4g\left(\frac{n}{4}\right) + c + 2c \le 8g\left(\frac{n}{8}\right) + c + 2c + 4c \dots \le 2^x \cdot g(1) + c + 2c + 4c + \dots + 2^{x-1}c \le bn + cn c$
- When $n \neq 2^x$, we can extend the array to the nearest power of two n' and show that $g(n) \leq g(n') \leq bn' + n' c \leq 2bn + 2cn c$.
- Thus, g(n) = O(n)



(6). What is the time complexity of the following code?

```
for (int i = 1; i <= n; i *= 2) {
   for (int j = 0; j < i; j++) {
        System.out.println(j);
   }
}</pre>
```

```
A. O(n \log n)
```

- B. O(n)
- C. $O(n^2)$
- D. $O(\log n)$.
- (7). What is the time complexity of the following Java method solve() for an input String s with n characters? (Note: A String is commonly stored as an array of characters. At line 4, s + i will create a new String. For example, "ABCD" + 1 will return "ABCD1".)

```
public void solve(String s) {
   int n = s.length();
   for (int i = 0; i < n; i++) s = s + i;
   System.out.println(s);
}</pre>
```

- A. O(n)
- B. $O(n^2)$
- C. O(1)
- D. $O(\log n)$.



- (8). Which of the following data structures allows insertion and deletion from both ends?
 - A. Stack
 - B. Dequeue
 - C. Queue
 - D. None of the above
- (9). One difference between a queue and a stack is:
 - A. Queues require linked lists, but stacks do not.
 - B. Stacks require linked lists, but queues do not.
 - C. Queues use two ends of the structure; stacks use only one.
 - D. Stacks use two ends of the structure, queues use only one.



(10). Here is an INCORRECT pseudocode for the algorithm which is supposed to determine whether a sequence of parentheses is balanced:

```
declare a character stack

while ( more input is available) {
    read a character
    if ( the character is a '('))
        push it on the stack
    else if ( the character is a ')' and the stack is not empty )
        pop a character off the stack
    else
        print "unbalanced" and exit
}

print "balanced"
```

Which of these unbalanced sequences does the above code think is balanced?

- A. ((())
- B. ())(()
- C. (()()))
- D. (()))()



- (1) (True/False) One can implement a stack based on a linked list so that each individual push/pop operation is in time O(1).
- (2) (True/False) One can reverse the order of the elements in a singly linked list with \mathbf{n} elements in time O(n).
- (3) (True/False) For any two given singly linked lists of arbitrary sizes, it is possible to concatenate (merge) them in time O(1).
- (4) (True/False) The worst-case time complexity of searching an element in a sorted array of n elements using binary search is $O(\log n)$.
- (5) (True/False) A queue can be implemented using two stacks.



(2) Reverse a linked list

```
// an iterative solution to reverse a linked list
public void reverseLinkedList() {
   Node currentNode = head;
   // for first node, previous node will be null
   Node previousNode = null;
   Node nextNode;
   while(currentNode != null) {
        nextNode = currentNode.next;
        // reverse the link
        currentNode.next = previousNode;
        // move current node and previous node by 1 node
        previousNode = currentNode;
        currentNode = nextNode;
   }
   head = previousNode;
}
```



- (1) (True/False) One can implement a stack based on a linked list so that each individual push/pop operation is in time O(1).
- (2) (True/False) One can reverse the order of the elements in a singly linked list with \mathbf{n} elements in time O(n).
- (3) (True False) For any two given singly linked lists of arbitrary sizes, it is possible to concatenate (merge) them in time O(1).
- (4) (True/False) The worst-case time complexity of searching an element in a sorted array of n elements using binary search is $O(\log n)$.
- (5) (True False) A queue can be implemented using two stacks.



#2 True or False

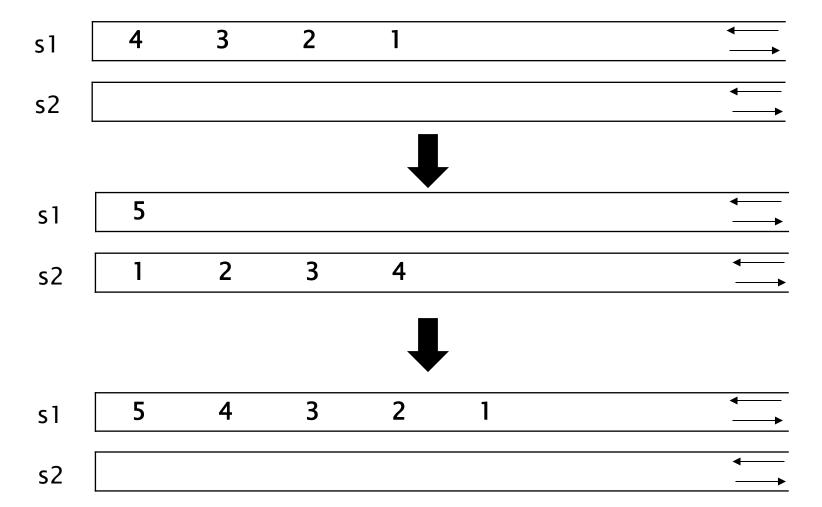
▶ (5) Implement a queue using two stacks

```
// Implement a queue using two stacks
   class Queue<T> {
        Queue() {
            s1 = new Stack<>();
            s2 = new Stack<>();
        // Add an item to the queue
        public void enqueue(T data) {
            while (!s1.isEmpty()) {
                s2.push(s1.pop());
            // push item into the first stack
            s1.push(data);
            // Move all elements back to the first stack from the second stack
            while (!s2.isEmpty()) {
                s1.push(s2.pop());
        public T dequeue() {
            // if the first stack is empty
            if (s1.isEmpty())
                System.out.println("Underflow!!");
            // return the top item from the first stack
            return s1.pop();
38 class Main {
        public static void main(String[] args) {
            int[] keys = { 1, 2, 3, 4, 5 };
            Queue<Integer> q = new Queue<Integer>();
            for (int key: keys) {
                q.enqueue(key);
            System.out.println(q.dequeue());
                                                // print 1
            System.out.println(q.dequeue());
                                                // print 2
```



#2 True or False

▶ (5) Implement a queue using two stacks





- 3. [6*5] Concisely answer the following short questions.
- (a) List five types of common data structures.

Answer:

Array, Linked list, Stack, Queue, Tree, ...

A typical error:

int, double, char, long, bool



- 3. [6*5] Concisely answer the following short questions.
- (b) What is the main advantage of using a linked list over an array?

Answer:

The main advantage of linked lists is dynamic memory allocation and efficient insertion/deletion operations.

A typical error:

Only answer one of the two key advantages.



- 3. [6*5] Concisely answer the following short questions.
- (c) How does the time complexity of inserting an element at the beginning of an array compare to that of a linked list?

Answer:

Array: O(n); Linked List: O(1).

Therefore, the time complexity for array insertion is greater than that of a linked list.

A typical error:

The time complexity for array insertion equals that of a linked list.



- 3. [6*5] Concisely answer the following short questions.
- (d) What is the advantage of using a circular queue over a linear queue?

Answer:

The main advantage of a circular queue is efficient space utilization.

A typical error:

The circular queue has lower time complexity.



- 3. [6*5] Concisely answer the following short questions.
- (e) Why is a stack more suitable for implementing undo operations, which are commonly seen in text editors, than a queue?

Answer:

A stack is ideal for undo operations as its LIFO behavior matches the expected undo sequence, while a queue's FIFO approach would inappropriately undo the oldest action first.

A typical error:

Stacks use LIFO without explaining why this property specifically benefits undo operations.



- 3. [6*5] Concisely answer the following short questions.
- (f) Asymptotic analysis of 6 functions.

Use "<" and/or "=" to order the following functions by their (asymptotic) growth rate with n: (a). $4n \log(n) + \sqrt{n}$; (b). $2^{\log(n^2)}$; (c). $3n + 20(\log(n))^2$; (d). $\log(n^5)$; (e). $n \log(n)$; (f). 2^{200} .

Answer:

$$f < d < c < a = e < b$$

A typical error:

$$n < (a), n < (b), n = (c), (d) < n, n < (e), (f) < n$$



#4 Doubly Linked List

4. [5*1] To implement a doubly linked list, maintaining a reference to the first and last node in the list, along with its size as follows.

```
public class LinkedList<Item> {
    private Node first; // the first node in the linked list
    private Node last; // the last node in the linked list
    private int N; // number of items in the linked list
    private class Node {
        private Item item; // the item
        private Node next, prev; // next and previous nodes
    }
    ...
}
```

What are the best estimates of the worst-case running time of the following operations in big-O notation? (1 point each, choose among O(1), $O(\sqrt{N})$, $O(\log N)$, O(N), $O(N \log N)$, $O(N^2)$.)



#4 Doubly Linked List

4. [5*1] To implement a doubly linked list, maintaining a reference to the first and last node in the list, along with its size as follows.

Answer:

(1) addFirst(Item item): Add an item to the beginning of the list.

O(1)

(2) get(int i): Return the item at position i of the list.

O(N)

(3) set(int i, Item item): Replace position i of the list with the item.

O(N)

(4) removeLast(): Delete and return the item at the end of the list.

O(1)

(5) contains(Item item): Is the item in the list?

O(N)



5. [5*2] Read the following Java programs, and answer questions.

```
public class Q5_1 {
      public static long fib1(int N) {
          if (N == 0) return 0;
          if (N == 1) return 1;
          return fib1(N-1) + fib1(N-2);
      public static void main(String[] args) {
          int N = 10;
          System.out.println(fib1(N));
  public class Q5_2 {
      public static long [] memo = new long[60];
      public static long fib2(int N) {
          if (memo[N] > 0) return memo[N];
         if (N == 0) return 0;
          if (N == 1) return 1;
          memo[N] = fib2(N-1) + fib2(N-2);
          return memo[N];
      public static void main(String[] args) {
          int N = 10;
          System.out.println(fib2(N));
24
```



5. [5*2] Read the following Java programs, and answer questions.

(1) Write down the outputs of the two programs, i.e., the returned values of fib1(10) and fib2(10) respectively.

Answer:

Fib1(10) = 55, Fib2(10) = 55

(2) For methods fib1(int N) and fib2(int N), estimate the orders of growth of operations (time complexity) as a function of the input N, respectively.

Answer:

Typical Errors:

For Fib1, O(2^N)

 $O(N^2), O(N)$

For Fib2, O(N)

O(logN)



#6 Find Second Largest Value

6. [5] Given an unsorted array of size n, write an algorithm in O(n) to find the second largest element in the array.

```
int findSecondLargest(int arr[], int n){
      ... // add your operations
      return secondLargest;
   }
}
```



6 Find Second Largest Value

6. [5] Given an unsorted array of size n, write an algorithm in O(n) to find the second largest element in the array.

```
Answer:
if (arr.length < 2) { return Integer.MIN_VALUE; }.
if (arr.length == 2) { return Math.min(arr[0], arr[1]); }
int largest = Math.max(arr[0], arr[1]);
int second_largest = Math.min(arr[0], arr[1]);
for (int idx = 2; idx < arr.length; idx++) {
  if (arr[idx] > largest) {
     second_largest = largest;
     largest = arr[idx]; }
  else if (arr[idx] > second_largest) {
     second_largest = arr[idx]; }
return second_largest;
```



6 Find Second Largest Value

6. [5] Given an unsorted array of size n, write an algorithm in O(n) to find the second largest element in the array.

```
Typical Error:
if (arr.length < 2) { return Integer.MIN_VALUE; }.
if (arr.length == 2) { return Math.min(arr[0], arr[1]); }
int largest = Math.max(arr[0], arr[1]);
int second_largest = Math.min(arr[0], arr[1]);
for (int idx = 2; idx < arr.length; idx++) {
  if (arr[idx] > largest) {
     second_largest = largest;
     largest = arr[idx]; }
  else if (arr[idx] > second_largest) { // Forget to consider this situation
     second_largest = arr[idx]; }
return second_largest;
```



#7 Reverse First K elements in a Queue

7. [10] Given an integer k and a queue Q with at least k elements, reverse the first k elements while keeping the rest of the queue unchanged.

Answer:

```
void reverseKELements(queue Q, int k){
 queue Q aux;
                                                               void reverseKELements(queue Q, int k){
 stack S aux:
                                                                 stack 5;
 int n = Q.size();
                                                                 int n = Q.size();
 for(int i=0;i<k;i++){
                                                                 for(int i=0;i<k;i++){
  S_aux.push(Q.front());
                                                                  S.push(Q.front());
  Q.pop();
                                                                  Q.pop();
 for(int i=k;i<n;i++){
                                                                 for(int i=0:i<k;i++){
  Q_aux.push(Q.front());
                                                                  Q.push(S.top());
  Q.pop();
                                                                  S.pop();
 for(int i=0;i<k;i++){
                                                                 for(int i=k;i<n;i++){
  Q.push(S_aux.top());
                                                                  Q.push(Q.front());
  S_aux.pop();
                                                                  Q.pop();
 for(int i=k;i<n;i++){
                                                                 return Q;
  Q.push(Q_aux.front());
  Q_aux.pop();
 return Q;
```



#7 Reverse First K elements in a Queue

7. [10] Given an integer k and a queue Q with at least k elements, reverse the first k elements while keeping the rest of the queue unchanged.

A typical error:

```
void reverseKELements(queue Q, int k){
  queue Q_aux;
  stack S_aux;
  int n = Q.size();
  for(int i=0;i<k;i++){
      S_aux.push(Q.front());
      Q.pop();
}
for(int i=0;i<k;i++){
      Q_aux.push(Q.front());
      Q.pop();
}
for(int i=0;i<k;i++){
      Q.push(S_aux.top());
      S_aux.pop();
}
for(int i=0;i<k;i++){
      Q.push(Q_aux.front());
      Q.aux.pop();
}
return Q;
}</pre>
```