



香港中文大學 (深圳)
The Chinese University of Hong Kong

CSC3100 Data Structures

Course information

Li Jiang
School of Data Science (SDS)
The Chinese University of Hong Kong, Shenzhen



About instructor (Li Jiang)

► Background

- PhD: 2017-2021, CUHK
- Postdoc: 2021-2023, MPI-INF
- Assistant Professor: 2024-present, CUHK-SZ



Prof. Li Jiang

► Contact

- Email: jiangli@cuhk.edu.cn
- Office: Teaching Complex C610

► Research

- General area: computer vision, artificial intelligence
- Topics: 3D scene understanding, motion prediction, representation learning, autonomous driving, VLM, AIGC, etc.



Lecture location and time

- ▶ Two sessions in total
 - Leading instructor: Prof. Wenye Li



Prof. Wenye Li

时段	L01-LEC(3948)	注册人数限制	120	状态	●
期次	正常	注册合计	69		
星期和时间		教室	讲师	会议日期	
星期一-星期三 10:30 - 11:50		Teaching Complex B206	Wenye LI	2025/01/06 - 2025/05/09	

时段	L02-LEC(3949)	注册人数限制	120	状态	●
期次	正常	注册合计	99		
星期和时间		教室	讲师	会议日期	
星期二-星期四 13:30 - 14:50		Teaching Complex C303	Li JIANG	2025/01/06 - 2025/05/09	



Lecture location and time

▶ Session by Li Jiang

- Lecture location: Room 303, Teaching Complex C
- Lecture time (3 hours/week)
 - 1:30 pm ~ 2:50 pm on Tuesday
 - 1:30 pm ~ 2:50 pm on Thursday
- Office hour
 - Time: 4:00 pm ~ 5:00 pm on every Tuesday
 - Location: Room 610, Teaching Complex C

A break of 5~10
mins



TA information

- ▶ 4 TA and 4 USTF
 - Leading TA: [Mr. Shunlin Lu](#)
 - Host tutorials; assignments; exams; office hour

TAs



Shunlin Lu



Yaomin Wang



Chen Shi



Jingjing Qian

USTFs

Zeyuan He

Diyuan Deng

Hanjun Zheng

Wangmeiyu Zhang



Office hours of TAs

- ▶ Office hours of TAs
 - Monday 19:00-20:00, **To be determined**

ID	Name	PGTA/USTF	Email address
1	Shunlin Lu (leading TA)	PGTA	223040243@link.cuhk.edu.cn
2	Yaomin Wang	PGTA	222042013@link.cuhk.edu.cn
3	Chen Shi	PGTA	224040349@link.cuhk.edu.cn
4	Jingjing Qian	PGTA	224040366@link.cuhk.edu.cn
5	Zeyuan He	USTF	123090168@link.cuhk.edu.cn
6	Diyuan Deng	USTF	123090079@link.cuhk.edu.cn
7	Hanjun Zheng	USTF	122090797@link.cuhk.edu.cn
8	Wangmeiyu Zhang	USTF	123090825@link.cuhk.edu.cn



Tutorials



群聊: CSC3100 Student 2025

► Tutorials

- Check announcements in your e-mail box or Blackboard
- Check Wechat group messages



该二维码7天内(1月13日前)有效, 重新进入将更新

Tutorial time	Monday	Tuesday	Wednesday	Thursday	Friday
19:00 - 19:50		TC414		TC407	
20:00 - 20:50		TC414			



Tutorials

► Tutorial arrangement

Week	Content	PGTF/USTF
1 (Jan 6 - Jan 10)	No tutorial	-
2 (Jan 13 - Jan 17)	Java & OJ	Diyuan & Hanjun
3 (Feb 10 - Feb 14)	Data Structures & Examples	Zeyuan He
4 (Feb 17 - Feb 21)		Jingjing Qian
5 (Feb 24 - Feb 28)		Yaoming Wang
6 (Mar 3 - Mar 7)		Yaoming Wang
7 (Mar 10 - Mar 14)		Yaoming Wang
8 (Mar 17 - Mar 21)	Midterm, no tutorial	-
9 (Mar 24 - Mar 28)	Data Structures & Examples	Chen Shi
10 (Mar 31 - Apr 3)		Chen Shi
11 (Apr 7 - Apr 11)		Chen Shi
12 (Apr 14 - Apr 18)		Jingjing Qian
13 (Apr 21 - Apr 25)		Jingjing Qian



Assessment



	Release date (Tentative)	Due date (Tentative)	Weight
Programming assignment 1	Jan 17	Feb 14	10%
Programming assignment 2	Feb 21	Mar 14	10%
Midterm exam	Mar 21 ?	--	20%
Programming assignment 3	Mar 14	Apr 3	10%
Programming assignment 4	Apr 3	Apr 25	10%
Final exam	May 10 - 17 ?		40%

* For assignments, students can use Python/C/Java to answer the questions in Online Judge (OJ) system

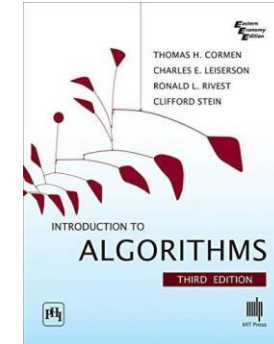
* OJ system url: <http://oj.cuhk.edu.cn/>



Textbook and references

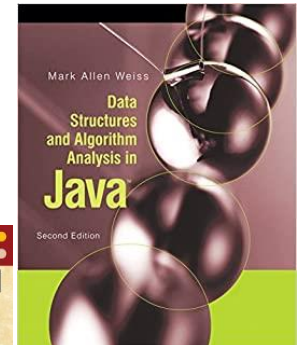
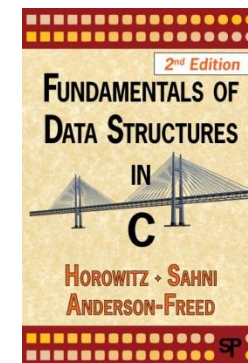
▶ Textbook

- T.H. Cormen, C.E. Leiserson, R.L. Rivest, C. Stein, Introduction to Algorithms, 3rd Edition, The MIT Press, 2009.
- [https://edutechlearners.com/download/Introduction to algorithms-3rd%20Edition.pdf](https://edutechlearners.com/download/Introduction%20to%20algorithms-3rd%20Edition.pdf)



▶ References

- M.A. Weiss, Data Structures and Algorithm Analysis in Java, 2nd Edition, Addison-Wesley, 2007.
- Ellis Horowitz, Sartaj Sahni, Susan Anderson-Freed, Fundamentals of Data Structures in C.





Teaching methods

- ▶ Lectures (3 hours/week)
 - Important materials from the textbook will be covered
 - Please ensure you stay up with the materials
 - Java will be used for illustrating the data structures
 - Feel free to interrupt to ask questions

- ▶ Tutorials (1 hour/week)
 - Illustrate more on the difficult parts
 - Show students more examples
 - Discuss assignment/exercise



How to do well in this course?

▶ Common suggestions

- Attend the lectures physically
- Slides of lectures/tutorials will be uploaded to Blackboard before lectures/tutorials; learn them in advance
- Use examples to facilitate learning
- **Data structures are not just reading materials; you need to write more codes!!!**

▶ Special suggestions

- If you feel difficult,
 - Try to focus on the content of slides and seek for help (**TAs and instructors are ready to answer your questions!**)
- If you feel easy,
 - Read more details (e.g., theoretical analysis) in the textbook
 - Use the learned techniques to solve ACM-ICPC problems



How to do well in this course?

- ▶ 100 Data Structure Problems in OJ
 - */**/** refer to easy/normal/difficult
 - 8 categories in total

Category	Number
Array/Linked List	12
Stack/Queue	12
Tree (BST, Heap)	16
Graph (BFS, DFS, MST, SP)	16
Sorting	10
Hashing	8
Recursion/Divide-and-conquer	16
Strings	10

Acknowledgement:

Tianci Hou, Yohandi, Ziyi Zhao,
Letian Cheng, Yuzhou Lin, Yutian Wang
Yiyu Ren, Frederick Khasanto

Maximum Weight | ★★ - Chir x +

Not secure 10.26.200.13/problem/c... ☆

Yixiang Fang - Go... 2023年度国家自然... 费控系统 >> All Bookmarks

Maximum Weight (★★)

[Submit solution](#)

[All submissions](#)

[Best submissions](#)

✓ Points: 100 (partial)

⌚ Time limit: 1.0s

📄 Memory limit: 256M

✎ Author: USACO

▼ Problem type: 100 Data Structure Problems | Graph

DESCRIPTION

Given a connected weighted simple graph with n vertices and m edges. Consider an edge with id i . Let's determine for this edge the maximum integer weight we can give to it so that it is considered in all minimum spanning trees of the graph if we don't change the other weights.

INPUT

The first line contains two integers n and m ($2 \leq n \leq 2 \times 10^5, n - 1 \leq m \leq 2 \times 10^5$), where n and m are the number of vertices and the number of edges in the graph, respectively.

Each of the next m lines contains three integers u, v and c ($1 \leq u, v \leq n, v \neq u, 1 \leq c \leq 10^9$) meaning that there is an edge between vertices u and v with weight c .

Finally there is a number, indicating the edge id i .



Learn more and practice more



- ▶ **ACM-ICPC**
 - ACM: Association for Computing Machinery
 - ICPC: International Collegiate Programming Contest

- ▶ **Online Judge (OJ) systems**
 - Many programming problems
 - CUHKSZ OJ (campus only)
 - <http://oj.cuhk.edu.cn/>
 - SJTU OJ
 - <https://acm.sjtu.edu.cn/OnlineJudge/>
 - PKU OJ
 - <http://poj.org/>



CUHK-SZ programming team

- ▶ Coaches: Yixiang Fang, Chenhao Ma
- ▶ Duration (hours): 4
- ▶ Questions: 12
- ▶ The first CUHK-SZ programming contest
 - Time: April 2, 2023
 - Participates: 235
 - Sponsor: Bopu
- ▶ The second CUHK-SZ programming contest
 - Time: April 13, 2024
 - Participates: 380
 - Sponsor: Bopu / Xinyoudui / Huawei / Lianqiang



News:

<https://sds.cuhk.edu.cn/article/991>

<https://sds.cuhk.edu.cn/event/1675>



Course policy

- ▶ Your work **MUST** be your own
 - Cheating is against “fair-play” and will not be tolerated under any circumstances

- ▶ Assignments
 - Penalty will be imposed on copying assignments; minimum penalty is zero mark for the assignments
 - **There will be penalty for late submission:**
 - 0~24 hours after deadline: **final score = your score x 0.8**
 - 24~72 hours after deadline: **final score = your score x 0.5**
 - 72+ hours after deadline: **final score = your score x 0**



Examples of code plagiarism

← 返回 | 代码对比

1.cpp 与 1.cpp 相似度为 100%

```
#include<bits/stdc++.h>
using namespace std;
int n,a[100100],r[100100];
long long ans=0;
void mergesort(int s,int t)
{
    if(s==t) return ;
    int mid=(s+t)/2;
    mergesort(s,mid);
    mergesort(mid+1,t);
    int i=s,j=mid+1,k=s;
    while(i<=mid&&j<=t)
    {
        if(a[i]<a[j]) r[k++]=a[i++];
        else r[k++]=a[j++],ans+=mid-i+1;
    }
    while(i<=mid) r[k++]=a[i++];
    while(j<=t) r[k++]=a[j++];
    for(int i=s;i<=t;++i)
        a[i]=r[i];
}

int main()
{
    scanf("%d",&n);
    for(int i=1;i<=n;++i)
        scanf("%d",&a[i]);
    mergesort(1,n);
    printf("%lld\n",ans);
}
```

```
#include<bits/stdc++.h>
using namespace std;
int n,a[500500],r[500500];
long long ans=0;
void msort(int s,int t)
{
    if(s==t) return ;
    int mid=(s+t)/2;
    msort(s,mid);msort(mid+1,t);
    int i=s,j=mid+1,k=s;
    while(i<=mid&&j<=t)
    {
        if(a[i]<a[j]) r[k++]=a[i++];
        else r[k++]=a[j++],ans+=mid-i+1;
    }
    while(i<=mid) r[k++]=a[i++];
    while(j<=t) r[k++]=a[j++];
    for(int i=s;i<=t;++i)
        a[i]=r[i];
}

int main()
{
    scanf("%d",&n);
    for(int i=1;i<=n;++i)
        scanf("%d",&a[i]);
    msort(1,n);
    printf("%lld\n",ans);
}
```



Examples of code plagiarism

Sample1.cpp 与 Sample2.cpp 相似度为 100%

```
#include<iostream>
using namespace std;

void fool() {
    // an easy function
    int a = 10;
    cout<<a;
}

void foo2() {
    /* I
       am
       comments*/
    int b = 20;
    cout<<b;
}

int main() {
    fool();
    foo2();
}
```

```
#include<iostream>
using namespace std;

void fun2() {
    int c=20; cout<<c;
}

void fun1() {
    int d=10; cout<<d;
}

int main() {
    fun1();
    fun2();
}
```

Changing the variable and function names, and re-organizing function orders are typical kinds of code plagiarism... We will use a software tool to detect code plagiarism automatically!

Don't send your source codes to others, which may make you be the one who is copied from!



Students with diverse backgrounds

Year	Percentage
1 st year	0%
2 nd year	57.3%
3 rd year	24.2%
4 th year	18.5%
Overall	100%

▶ Student distribution

- **SME: 35.4 %**
 - Financial engineering
 - Finance
 - Economics
 - Global Business Studies
- **SSE: 14.0%**
 - Electrical and Computer Engine
 - Mathematics and Applied Math
 - Electronic Info Engineering
- **SDS: 48.9%**
 - Data Science and Big Data Tech
 - Computer Science and Engineering
 - Statistics
- **Others: 1.7%**
 - Bioinformatics
 - English Studies

Please try to consider your classmates before complaining



Feedback & acknowledgements

- ▶ Feedback is important and also welcome!
 - Talk to course instructors and TAs, or send us emails
 - Please talk to us before complaining to others

- ▶ Acknowledgements
 - Lecture slides
 - Prof. Yixiang Fang (CUHK-SZ), Prof. Tianshu Yu (CUHK-SZ), Prof. Xiang Wan (CUHK-SZ), Prof. Kaiming Shen (CUHK-SZ), Prof. Wenye Li (CUHK-SZ), Prof. Minming Li (CityU, HK), Prof. Zengfeng Huang (FDU), Prof. Jane You (PolyU, HK), Prof. Sibow Wang (CUHK)
 - Tutorial slides
 - Mr. Xingchao Wang (CUHK-SZ), Mr. Panwen Hu (CUHK-SZ), Mr. Ziteng Weng (CUHK-SZ)



香港中文大學 (深圳)
The Chinese University of Hong Kong

CSC3100 Data Structures

Lecture 1: Introduction

Li Jiang
School of Data Science (SDS)
The Chinese University of Hong Kong, Shenzhen



Outline

- ▶ Introduction
 - Why take this course?
 - Basic concepts, e.g., abstract data type (ADT)
 - Relationship of ADT, data structures, and algorithms
 - Topics in this course
 - Tentative teaching plan



Why take this course?

- ▶ Required course
 - Also very important and very useful
 - A fundamental course in computer science
- ▶ Learn to save the data (**data structure**) and manipulate the data (**algorithms**) effectively and efficiently
- ▶ No single data structure fits all scenarios
 - Array: friendly to search (if sorted), not friendly to updates
 - List: friendly to updates, not friendly to search
- ▶ How useful in practice?
 - E.g., validate one Chinese ID in 1.4 billion people
 - E.g., find the best driving route
 - E.g., find webpages in Google



Real examples

▶ Route planning

- Find the shortest path between two specific locations
- Input:
 - A road network
 - A source node (location)
 - A destination node (location)
- Output
 - A path, or a sequence of edges, with the shortest total distance





Real examples

► Google search

- Find the documents matching your query keywords
- Use sophisticated algorithms to create an index structure which is just a data structure

30 trillion (30×10^{12})
webpages



data structure



Books

Images

Videos

More

Tools

About 3,290,000,000 results (0.65 seconds)

A data structure is a **specialized format for organizing, processing, retrieving and storing data**. There are several basic and advanced types of data structures, all designed to arrange data to suit a specific purpose. Data structures make it easy for users to access and work with the data they need in appropriate ways.

<https://www.techtarget.com/definition/data-structure>

[What are Data Structures? - Definition from WhatIs.com](#)

<https://www.geeksforgeeks.org/data-structures>

[Data Structures - GeeksforGeeks](#)

Aug 10, 2022 — A **data structure** is a storage that is used to store and organize data. It is a way of arranging data on a computer so that it can be accessed ...

[Introduction to Data Structures](#) · [Data Structure Alignment](#) · [Data Structures \[129\]](#)

People also ask

What is data structure and types?



What is data structure and its example?



Where is data structures used?



What are the four basic data structures?



Feedback



One sentence about this course

This course is about how to use computing resources and tools to solve practical problems **correctly** and **efficiently**

↓
Logical thinking and math

↓
Good at data structures,
algorithms, and programming



Data abstraction

- ▶ A clear **separation** between
 - the **abstract properties** of a data type and
 - the concrete **details of its implementation**

- ▶ Example: smartphone
 - Users do not know
 - How calls are made (e.g., CPU, memory, battery, and electric wire)
 - How the phone accesses the Internet
 - How the data is stored in the phone

 - Users do know
 - To make a call: input a person's phone number
 - To access the Internet: open the browser



Abstract Data Type (ADT)

- ▶ An ADT is for encapsulation (information hiding)
 - The implementation of an ADT and its operations can be localized to one section of the program
 - Procedures that make use of the ADT can safely ignore its implementation details

- ▶ Benefit of ADT
 - **User-friendly:** users do not need to know the mechanisms of how to connect to the Internet
 - **Designer-friendly:** designer can change mechanisms without affecting users
 - **Protection:** others cannot know your secrets!



How to separate?

- ▶ An ADT only provides the definition of operations
 - It consists of names of every operation (function), the type of its arguments, and the type of its result

1	ADT IntegerSet
2	IntegerSet createEmptySet();
3	IntegerSet addElementToSet(integer, SetA);
4	Boolean search(integer, SetA);
5	IntegerSet intersection(setA, setB);
6	IntegerSet union(setA, setB);
7	IntegerSet difference(setA, setB);

- Does not reveal the internal implementation details
 - How the set is represented? Array, List, Tree, etc.?
 - How the operations are implemented? Many different ways



Relationships: ADT, DS, algorithms

- ▶ A data structure (DS): **implementation** of an ADT
 - List ADT
 - Implementation: ArrayList or LinkedList
 - Set ADT
 - Implementation: Hash Table or Red-Black Tree
- ▶ An algorithm: **implementation of operations** in ADT





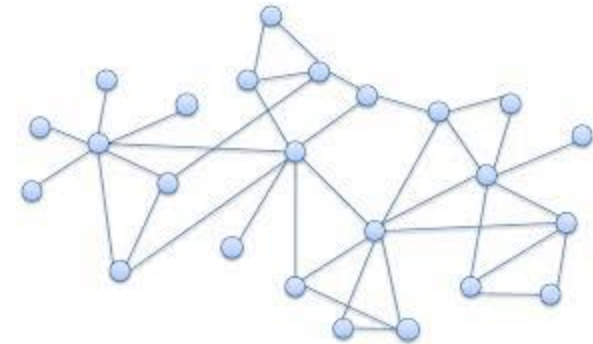
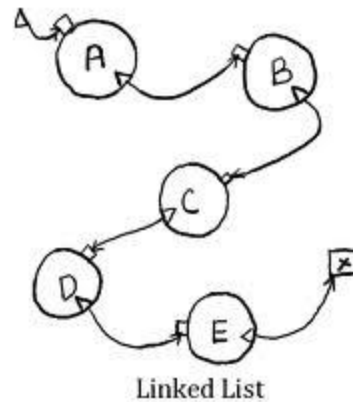
What are the common ADTs?

- ▶ Arrays, lists, stacks, queues, trees, graphs, etc.

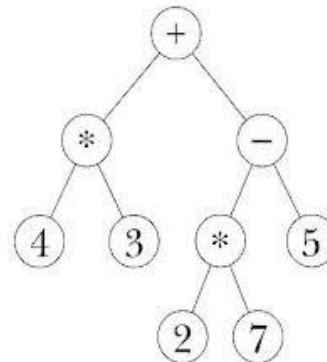
Array:

23	4	6	15	5	7
0	1	2	3	4	5

↑
Array index



Tree





Topics covered

- ▶ How to define "efficient"?
 - Time complexity analysis
- ▶ No single data structure fits all scenarios
 - **Array**: Efficient for sorting and search, not for updates
 - **List**: Efficient for updates, but not in searching/sorting
 - **Stack/queues**: Fast for certain types of updates, not in searching/sorting
 - **Trees**: Heap & binary search trees
 - **Hash tables**: Hash functions
 - **Graphs**: Graph algorithms
- ▶ Efficient operations
 - **Sorting**
 - Searching/Updates on different data structures (Discussed when learning the specific data structure)



Tentative teaching plan

*This tentative teaching plan is subject to change as the instructor sees fit.

Week	Content
1	Course overview, Java basics
2	Arrays, insertion/merge sort
3	Time complexity
4	List
5	Stack, queues
6	More sorting algorithms
7	Trees

Week	Content
8	Midterm exam
9	Trees
10	Hashing
11	Graphs
12	Graphs
13	Graphs, other data structures
14	Course review

Throughout this course, we will see:

- How to estimate the time cost of a program
- How to select proper data structure(s) to solve real problems
- Techniques to improve the speed of a program



Example: selection problem

► Problem

- Given N numbers, determine the k^{th} largest, where $N > k$

► Solution 1:

- 1) read N number into an array
- 2) sort the array in descending order
- 3) return the element in position k

► Solution 2:

- 1) read the first k elements into an array and sort them in descending order
- 2) each remaining element is read one by one,
 - 2.1) it is ignored if it is smaller than or equal to the k^{th} element in the array
 - 2.2) otherwise, it is placed in its correct spot in the array, bumping one element out of the array
- 3) the element in the k^{th} position is returned as the answer



Example: selection problem

- ▶ Two natural questions:
 - Which solution is better ?
 - By simulation
 - By theoretical analysis
 - Is either algorithm good enough (particularly when N is very large)?
 - A simulation using 1 million elements and $k = 500,000$ will show that NEITHER algorithm finishes in a reasonable amount of time
 - Is there a better algorithm?
- ▶ What is the conclusion?
 - Writing a working program is not good enough!
 - The optimum solution: a correct algorithm with minimum resources and minimum running time



Exercise

- ▶ Is it possible to find a better algorithm to select the k^{th} largest number?



Mathematics review

► Exponents

$$X^A X^B = X^{A+B}$$

$$\frac{X^A}{X^B} = X^{A-B}$$

$$(X^A)^B = X^{AB}$$

$$X^N + X^N = 2X^N \neq X^{2N}$$

$$2^N + 2^N = 2^{N+1}$$



Mathematics review

- ▶ Logarithm definition:

$$X^A = B \text{ if and only if } \log_x B = A$$

- ▶ All log are to be **base 2** unless specified otherwise

– *Useful equalities*

$$\log_A B = \frac{\log_c B}{\log_c A}; C > 0$$

$$\log AB = \log A + \log B$$

$$\log\left(\frac{A}{B}\right) = \log A - \log B$$

$$\log x < x, \quad \forall x > 0$$

$$\log 1 = 0, \log 2 = 1, \log 1024 = 10, \log 65536 = 16$$



Mathematics review

- Series: arithmetic series

$$\sum_{i=1}^N i = \frac{N(N+1)}{2}$$

$$\sum_{i=1}^N i^2 = \frac{N(N+1)(2N+1)}{6}$$

Example: To find the sum
 $2+5+8+\dots+(3k-1)$

$$= 3(1+2+3+\dots+k) \\ - (1+1+1+\dots+1)$$

$$= \frac{3k(k+1)}{2} - k$$



Mathematics review

- ▶ Series: geometric series

$$\sum_{i=0}^N A^i = \frac{1 - A^{N+1}}{1 - A}$$

If $0 < A < 1$, then

$$\sum_{i=0}^N A^i \leq \frac{1}{1 - A}$$

- ▶ Derivation

$$\text{Let } S = 1 + A + A^2 + \dots \quad (1)$$

where, $0 < A < 1$

$$\text{then } AS = A + A^2 + A^3 + \dots \quad (2)$$

Subtracting (1) and (2), we get $S - AS \leq 1$, i.e.

$$S \leq \frac{1}{1 - A}$$



Mathematics review

► To prove a false statement:

- proof by counter example

e.g., Fibonacci number:

$$F_0 = 1, F_1 = 1, F_{k+1} = F_k + F_{k-1}$$

To show the statement $F_k \leq k^2$ is false, we can compute a concrete counter example, e.g.,

$$F_{11} = 144 > 11^2.$$

► To prove a correct statement

- proof by induction
 - (1) proving a base case
 - (2) inductive hypothesis
- proof by contradiction
 - (1) assume it is false
 - (2) show that this assumption is false