SCALER

Array - Carry forward & Subarrays

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Notes

Count 'a-g' pairs

< *Question* > : Given a string s of lowercase characters, return the count of pairs (i, j) such that i < j and s[i] is 'a' and s[j] is 'g'.

s = "a b e g a g"
0 1 2 3 4 5

Ans = 3

i	j
0	3
0	5
4	5

s = "a c g d g a g"
0 1 2 3 4 5 6

Ans = 4

i	j
0	2
0	4
0	6
5	6

s = "b c a g g a a g"
0 1 2 3 4 5 6 7

Ans = 5

i	j
2	3
2	4
2	7
5	7
6	7

BF Idea

```
cnt = 0
for i → 0 to (N-1) {
    for j → 0 to (N-1) {
        if (i < j && s[i] == 'a' && s[j] == 'g')
            cnt++
    }
}
} return cnt
```

TC = O(N²) SC = O(1)

```
cnt = 0
for i → 0 to (N-2) {
    if (s[i] == 'a') {
        for j → i+1 to (N-1) {
            if (s[j] == 'g')
                cnt++
        }
    }
}
} return cnt
```

count # 'g' from (i+1) to (N-1).

ans += cnt_g[i+1]

TC = O(N²) → O(N) SC = O(1) → O(N)

---

s = " b c a g g a a g "
0 1 2 3 4 5 6 7

cnt_g → [3 3 3 3 2 1 1 1]

L ← R

ans = 0 + 3 + 1 + 1 = 5

SCALER

```
∀i, if (s[i]=='g')
      cnt_g[i] = cnt_g[i+1]+1
    else  // s[i] ≠ 'g'
      cnt_g[i] = cnt_g[i+1]
```

Idea Carry Forward → calculate & use

	0	1	2	3	4	5	6	7
str →	b	c	a	g	g	a	a	g

L ⟵ R

'g' → cnt 3 3 3 3 2 1 1 1

'a' → ans = 0 + 1 + 1 + 3 = 5

```
cnt = 0
ans = 0
for i → (N-1) to 0 {
    if (s[i] == 'g') cnt++ // calculate
    if (s[i] == 'a') ans += cnt // use
}
```

TC = O(N) SC = O(1)

i < j s[i] = 'a' s[j] = 'g'

count #'a' from left to right

0 —— i

cnt = 0 ↗

ans = 0

c a a g x g

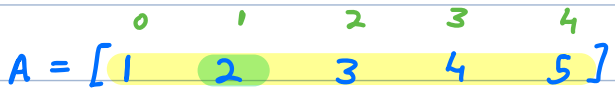
cnt = + 2

ans = 2 + 2 = 4

```
for i → 0 to N-1 {
    if (s[i] == 'a') cnt++ // calculate
    if (s[i] == 'g') ans += cnt // use
}
```

TC = O(N) SC = O(1)

Subarrays → continuous part of array



single element ✓

complete array ✓

Example: arr[ ] → [2 4 1 6 -3 7 8 4]

a. [1, 6, 8]

b. [1, 4]

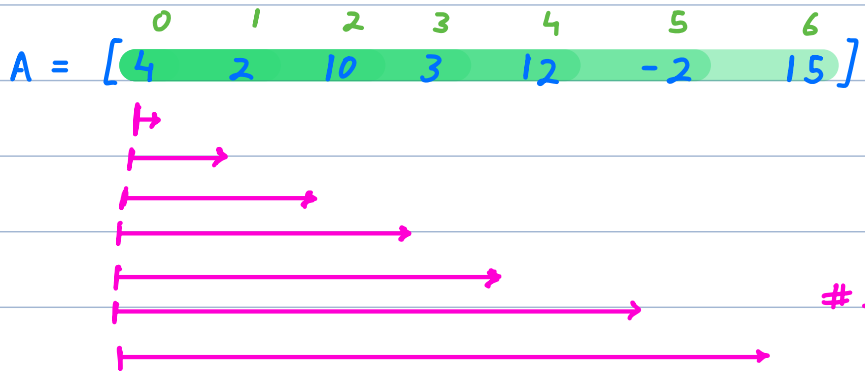
c. [6, 1, 4, 2]

d. [7, 8, 4,] ✓

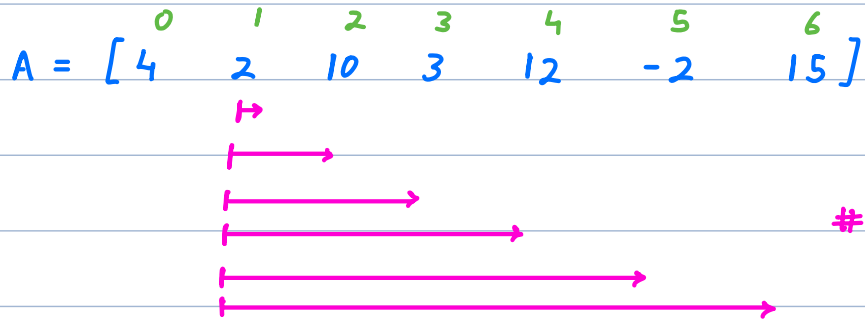
Representation of a subarray

1) L R — start & end

2) L (length) — start & length



subarrays starting from index 0 = N



Total number of subarrays

subarrays starting from ↓ = ↓

0	N
1	N-1
2	N-2
⋮	
(N-1)	1

N * (N+1) / 2

< Question > : Given an array, si and ei. Print from si to ei. $si \le ei$

arr → [4 2 10 3 12 -2 15] si = 2, ei = 5

0 1 2 3 4 5 6

o/p → 10 3 12 -2

- void printSubarray(arr, si, ei) {

```
for i → si to ei {
    print ( arr[i] )
}
```

}

print 1 subarray → T.C O(N)

SC = O(1)

< Question > : Print all the possible sub-arrays of the given array.

[5, 7, 3, 2]

0 1 2 3

$\frac{N * (N+1)}{2} * N \rightarrow O(N^3)$

8:30 AM

O/P - [5]

[5, 7]

[5, 7, 3]

[5, 7, 3, 2]

[7]

[7, 3]

[7, 3, 2]

[3]

[3, 2]

[2]

Idea Consider all the subarrays & print Subarray( )

```
for st → 0 to (N-1) {
    for end → st to (N-1) {
        for i → st to end {
            print (A[i])
        } print ('\n')
    }
}
```

TC = $O(N^3)$ SC = $O(1)$

Min Max

< *Question* > : Given an array of N integers, return the length of smallest subarray which contains both maximum and minimum elements of the array.

$1 \le N \le 10^6$

arr[ ] → [2 2 6 4 5 1 5 2 6 4 1]

0 1 2 3 4 5 6 7 8 9 10

Ans = 3

arr[ ] → [1 2 3 1 3 4 6 4 6 3]

Ans = 4

0 1 2 3 4 5 6 7 8 9

arr[ ] → [8 8 8 8 8 8]

0 1 2 3 4 5

min = 8 max = 8 Ans = 1

Bruteforce → Find min & max of the array → O(N)

∀ subarrays, check if it contains min & max and take smallest length subarray as answer.

$\frac{N*(N+1)}{2}$ → O(N)

TC = O(N + N^3) = $O(N^3)$ SC = O(1)

Observation

1. There must be exactly one occurrence of min & max element.

[min min - - - - max]

2. Min and max elements should be the end point of subarray.

3. ✓ ∀max check closest min on left

case-1: [min - - - - max] → ∀min check closest max on right

case-2: [max - - - - min]

✓ ∀min check closest max on left

∀max check closest min on right

	0	1	2	3	4	5	6	7	8	9	10
arr[ ] → [	2	2	6	4	5	1	5	2	6	4	1]
	i	i	i	i	i	i	i	i	i	i	i

min = 1

max = 6

latest_min_index = ~~-1~~ 10

latest_max_index = ~~-1~~ 8

length = 10 - 8 + 1 = 3

ans = ~~4~~ 3

< / > Code

```
minA = A[0]        maxA = A[0]
for i → 1 to (N-1) {
    minA = min (minA, A[i])
    maxA = max (maxA, A[i])
}
l_min = -1         l_max = -1
ans = N
for i → 0 to (N-1) {
    if (A[i] == minA) {
        l_min = i          // calculating
        if (l_max != -1) {
            ans = min (ans, i - l_max + 1)
        }                          // use
    }                    [l_max    i]
    if (A[i] == maxA) {
        l_max = i          min ——— max
        if (l_min != -1) {  [l_min ——— i]
            ans = min (ans, i - l_min + 1)
        }
    }
}
return ans
```

TC = O(N) SC = O(1)