ESC201: Lecture 18

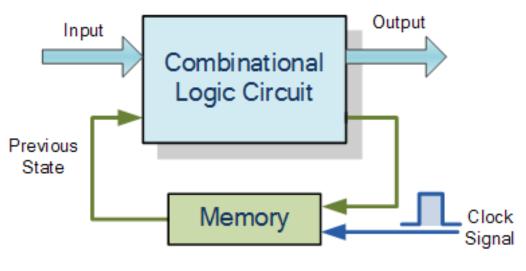


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2024-25 SEM-I | ESC201 INTRODUCTION TO ELECTRONICS

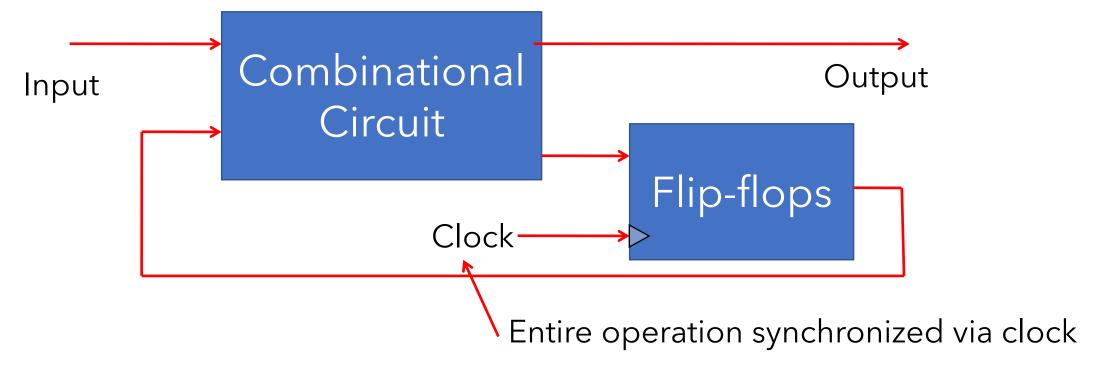
Sequential Circuits



- Calculation divided into steps
- Each step is triggered by a clock
- At each step,
 - output is based on the current values of inputs and past values of inputs/outputs.

• Requires memory

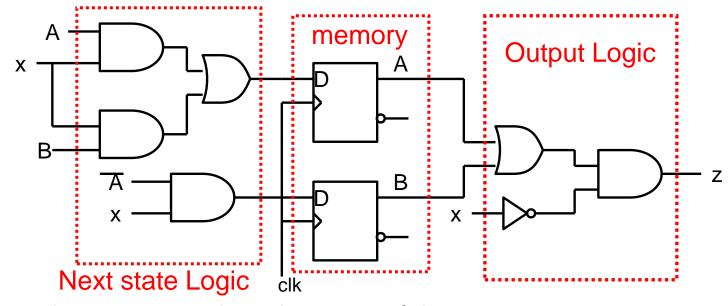
Synchronous Clocked Sequential Circuits



Employs signals that affect the stored value only at discrete instants of time.

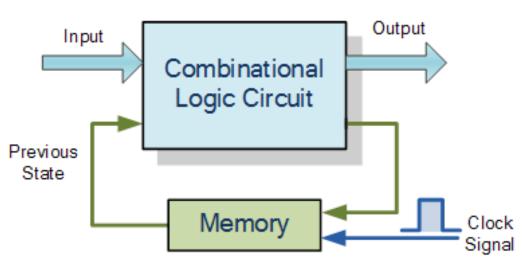
Synchronization is achieved via the *clock pulses*.

Analyzing sequential circuits



- Output z depends on the input x and on the state of the memory (A,B)
- The memory has 2 FFs and each FF can be in state 0 or 1.
- Thus there are four possible states: AB: 00,01,10,11
- To describe the behavior of a sequential circuit, we need to show
 - How the system goes from one memory state to the next as the input changes
 - How the output responds to input in each state

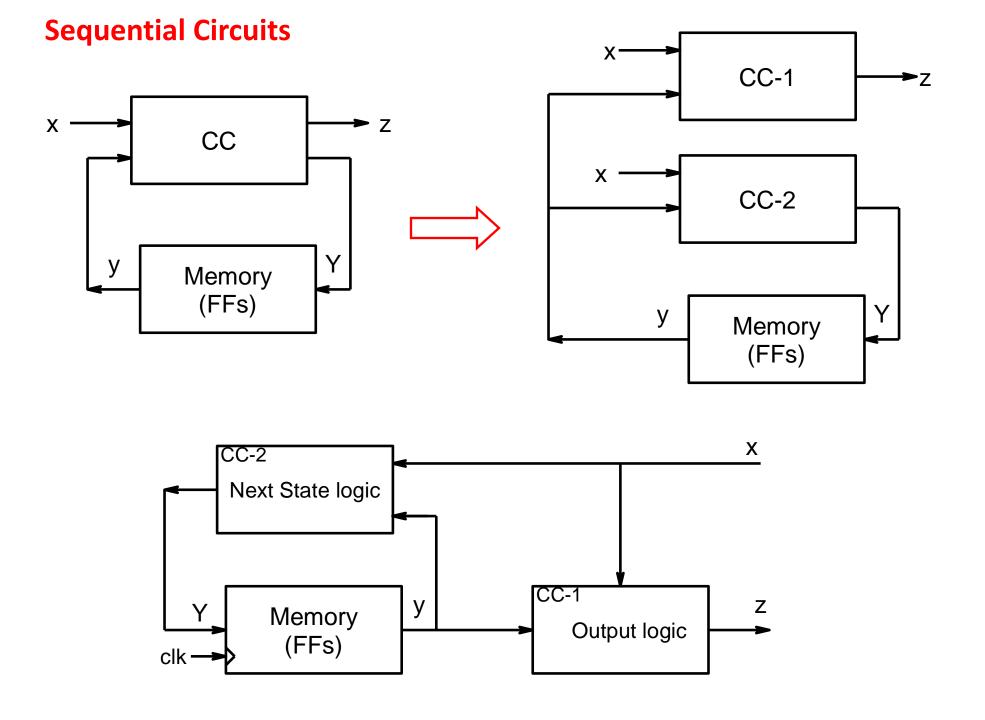
Sequential Circuits



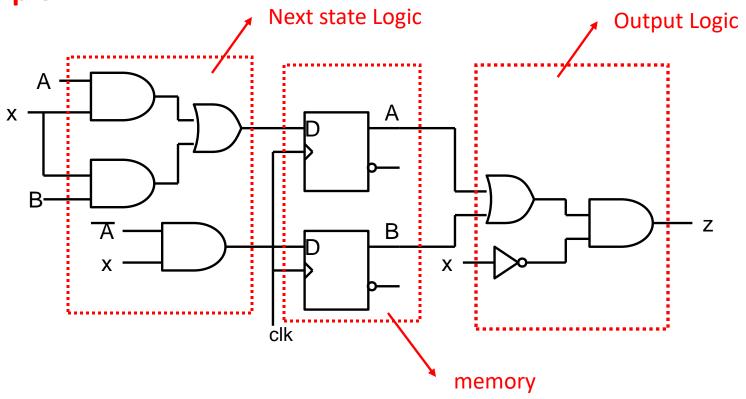
The binary information stored in the storage elements at any given time defines the **state** of the sequential circuit at that time

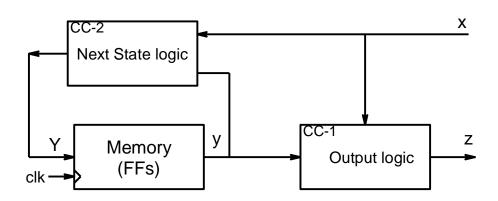
Output is a function of input as well as the present state (the stored value).

Next state is also a function of the present state and inputs.

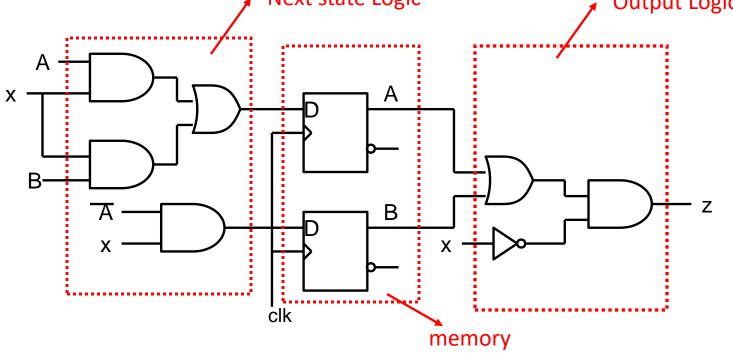


Example





Analysis Next state Logic **Output Logic**



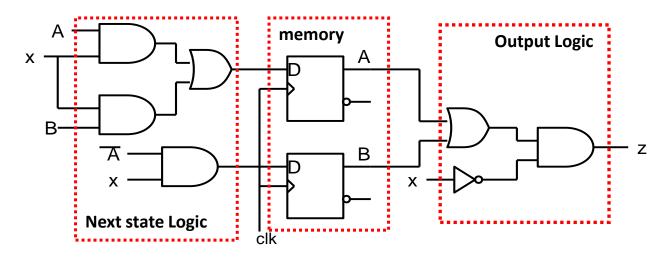
The dependence of output z on input x depends on the state of the memory (A,B)

The memory has 2 FFs and each FF can be in state 0 or 1. Thus there are four possible states: AB: 00,01,10,11.

To describe the behavior of a sequential circuit, we need to show

- how the system goes from one memory state to the next as the input changes
- How the output responds to input in each state

Analysis of Sequential Circuits



$$D_A = A.x + B.x$$
; $D_B = \overline{A}.x$; $z = (A + B).\overline{x}$

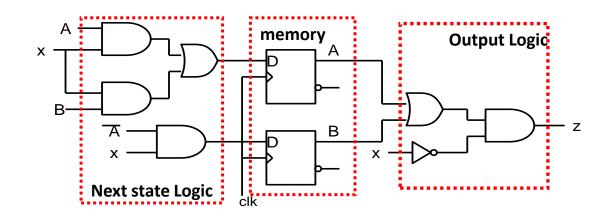
$$A(t+1) = A(t).x + B(t).x$$

$$B(t+1) = \overline{A(t)}.x$$

$$z = (A + B).\overline{x}$$

State Transition Table

		Input	Next State		Output
A(t)	B(t)	x(t)	A(t+2	L) B(t+1)	z(t)
0	0	0	0	0	0
0	0	1	0	1	0
0	1	0	0	0	1
0	1	1	1	1	0
1	0	0	0	0	1
1	0	1	1	0	0
1	1	0	0	0	1
1	1	1	1	0	0

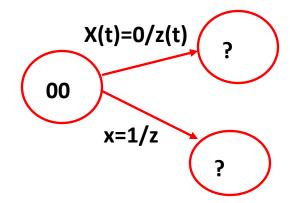


State Transition Table

Present State Inp		Input	Next State		Output
A(t) B(t)	x(t)	A(t+:	1 B(t+1)	z(t)
0	0	0	0	0	0
0	0	1	0	1	0
0	1	0	0	0	1
0	1	1	1	1	0
1	0	0	0	0	1
1	0	1	1	0	0
1	1	0	0	0	1
1	1	1	1	0	0

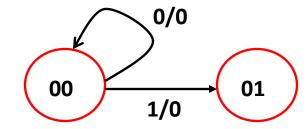
00)——

Memory state in which FF A& B have output values 00

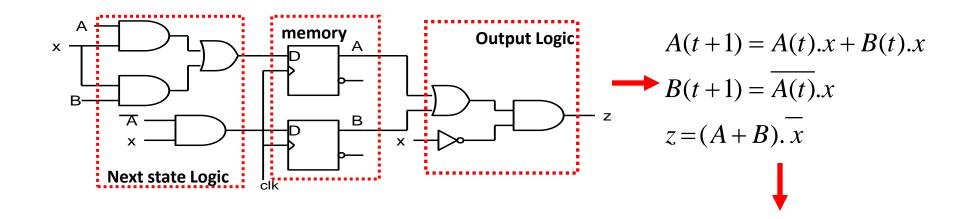


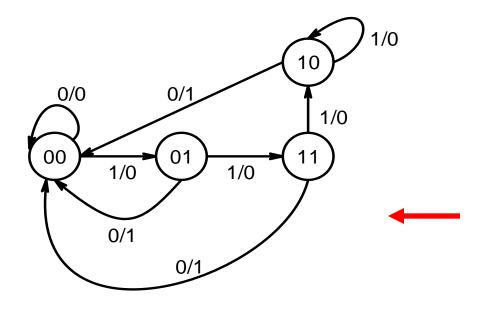
If x = 0 then z = 0, When the clock edge comes the system would stay in 00 state.

If x = 1 then z = 0. When the clock edge comes the system would go to 01 state.



Analysis of Sequential Circuits





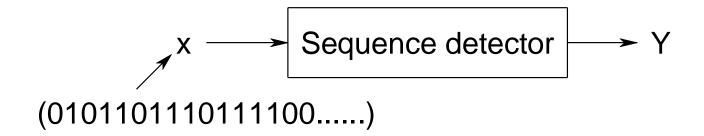
State transition Graph

State Transition Table

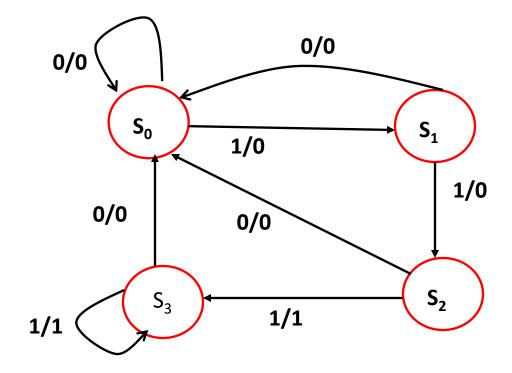
Þr	esent	State	Input	Next State		Output
	A(t)	B(t)	x(t)	A(t+:	1) B(t+1)	z(t)
	0	0	0	0	0	0
	0	0	1	0	1	0
	0	1	0	0	0	1
	0	1	1	1	1	0
	1	0	0	0	0	1
	1	0	1	1	0	0
	1	1	0	0	0	1
	1	1	1	1	0	0

Design of Sequential Circuits Specifications State Diagram State Transition Table Present State Input **Next State** Output В Z Choose FF 0 State Encoding State Table Synthsize Combinational Circuit

System specification to State diagram

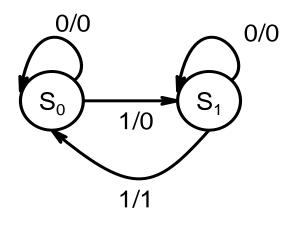


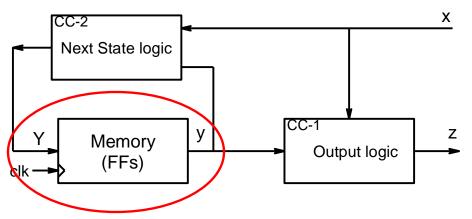
Detect 3 or more consecutive 1's in the input stream



Conversion of State transition graph to a circuit

Example-1





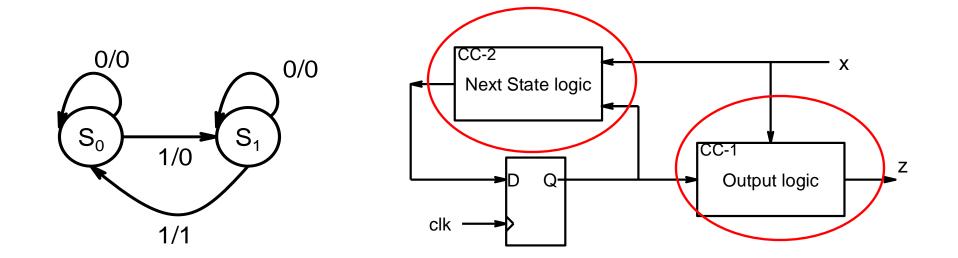
3 blocks need to be designed

- 1. How many FFs do we need?
- N FFS can represent 2^N states so Minimum is 1

2. Which FF do we choose?

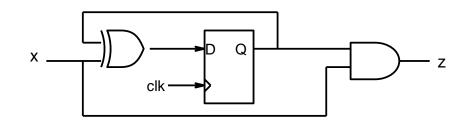
Say D FF

- 3. How are the states encoded?
- Say FF output Q=0 represents S₀ and Q=1 represents S₁ state



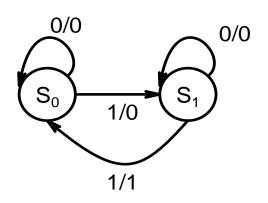
State Transition Table

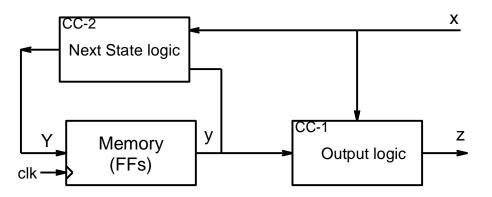
Present State	Input	Next State	D	Output
Q(t)	Х	Q(t+1)		Z
0	0 1	0 1	0 1	0
1	0	1 0	1 0	0 1



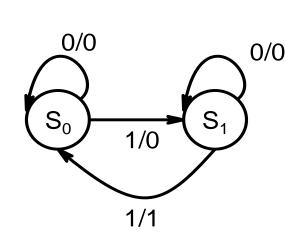
$$D = \overline{Q}.x + Q.\overline{x} \quad ; \quad z = Q.x$$

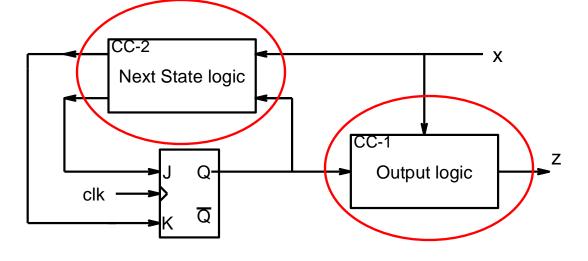
Example-2





- 1. How many FFs do we need?
- 2. Which FF do we choose? Say JK FF
- 3. How are the states encoded? Say FF output Q=0 represents S_0 and Q=1 represents S_1 state

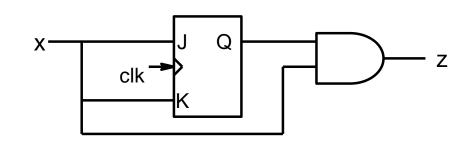




State Transition Table

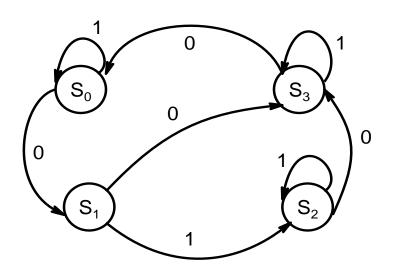
Present State	Input	Next State	J K	Output
Q(t)	Х	Q(t+1)		z
0	0 1	0 1	0 X 1 X	0
1	0	1 0	X 0 X 1	0

Q(t)	Q(t+1)	J K
0	0	0 X
0	1	1 X
1	0	X 1
1	1	X 0



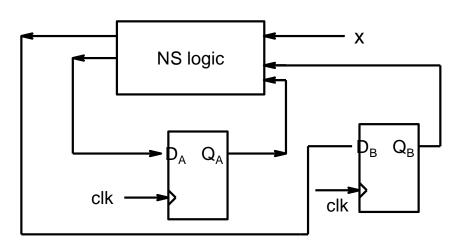
$$J = x$$
; $K = x$; $z = Q.x$

Example-3



	FF	O/P
State	Α	В
S _o	0	0
S ₁	0	1
S_2	1	0
S_3	1	1

For 4 states a minimum of two FFs will be required. Let us choose 2 D FFs A &B



Present State		Input	Next	State		
Α	В	Х	Α	В	D_A	D_B
0 0	0	0 1	0 0	1 0	0	1
0	1	0	1 1	1 0	1	1
1	0	0 1	1 1	1 0	1	1
1	1	0	0 1	0	0	0

Preser	nt State	Input	Next	State		
Α	В	Х	Α	В	D_A	D _B
0	0	0 1	0 0	1 0	0	1 0
0	1	0 1	1 1	1 0	1	1
1 1	0	0 1	1 1	1 0	1	1
1	1	0	0	0	0	0

ΛF	2	D_A			
x AB 00		01	11	10	
0	0	:17	0	:17	
1	0	1.	1:	.1	

$$D_A = \overline{A}B + xB + A\overline{B}$$
$$= A \oplus B + x.B$$

$$D_{B} = \overline{x}.\overline{A} + \overline{x}.\overline{B} + x.A.B$$

$$= \overline{x}.(\overline{A} + \overline{B}) + x.A.B$$

$$= \overline{x}.\overline{AB} + x.AB = \overline{x} \oplus \overline{AB}$$

$$D_A = A \oplus B + x.B$$
 $D_B = \overline{x \oplus AB}$

