

ESC201: Lecture 15



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Goal of Simplification

In the SOP expression:

1. Minimize number of product terms
2. Minimize number of literals in each term

Simplification \Rightarrow Minimization

Minimization

$$y = \overline{x_1} \cdot \overline{x_2} \cdot x_3 + \overline{x_1} \cdot x_2 \cdot x_3 + x_1 \cdot \overline{x_2} \cdot x_3 + x_1 \cdot x_2 \cdot x_3$$

$$y = \overline{x_1} \cdot x_3 \cdot (\overline{x_2} + x_2) + x_1 \cdot x_3 \cdot (\overline{x_2} + x_2)$$

$$y = \overline{x_1} \cdot x_3 + x_1 \cdot x_3$$

$$y = (\overline{x_1} + x_1) \cdot x_3$$

$$y = x_3$$

Principle used: $x + \overline{x} = 1$

$$f = \bar{x} \cdot \bar{y} + \bar{x} \cdot y + x \cdot \bar{y}$$

Apply the Principle: $x + \bar{x} = 1$ to simplify

$$f = \bar{x} \cdot (\bar{y} + y) + x \cdot \bar{y}$$

$$f = \bar{x} + x \cdot \bar{y}$$

How do we simplify further?

$$f = \bar{x} \cdot \bar{y} + \bar{x} \cdot y + x \cdot \bar{y} = \bar{x} \cdot \bar{y} + \bar{x} \cdot \bar{y} + \bar{x} \cdot y + x \cdot \bar{y}$$

Principle used : $x + x = x$

$$\begin{aligned} f &= \bar{x} \cdot \bar{y} + \bar{x} \cdot y + \bar{x} \cdot \bar{y} + x \cdot \bar{y} \\ &= \bar{x} \cdot (\bar{y} + y) + (\bar{x} + x) \cdot \bar{y} = \bar{x} + \bar{y} \end{aligned}$$

Simplify

$$f = \overline{x_1} \cdot \overline{x_2} \cdot \overline{x_3} \cdot x_4 + \overline{x_1} \cdot x_2 \cdot \overline{x_3} \cdot x_4 + x_1 \cdot \overline{x_2} \cdot \overline{x_3} \cdot x_4 + x_1 \cdot x_2 \cdot \overline{x_3} \cdot x_4 + \\ \overline{x_1} \cdot \overline{x_2} \cdot x_3 \cdot x_4 + x_1 \cdot \overline{x_2} \cdot x_3 \cdot x_4$$

Principle: $x + \overline{x} = 1$ and $x + x = x$

Need a systematic and simpler method for applying these two principles

Karnaugh Map (K map) is a popular technique for carrying out simplification

It represents the information in problem in such a way that the two principles become easy to apply

K-map representation of truth table

x	y	min term
0	0	$\overline{x} \cdot \overline{y}$ m0
0	1	$\overline{x} \cdot \underline{y}$ m1
1	0	$\underline{x} \cdot \overline{y}$ m2
1	1	$\underline{x} \cdot \underline{y}$ m3

		y	
		0	1
x	0	m ₀	m ₁
	1	m ₂	m ₃

x	y	f ₁
0	0	0
0	1	1
1	0	1
1	1	0



		y	
		0	1
x	0	0	1
	1	1	0

$$f_2 = \sum (0, 2, 3)$$



		y	
		0	1
x	0	0 1	1 0
	1	1	1

		y	
		0	1
x	0	1	0
	1	0	1



$$f = \bar{x}.\bar{y} + x.y$$

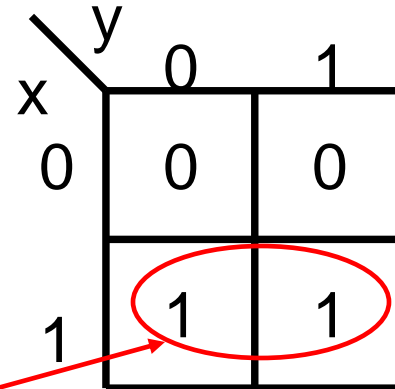
Minimization using Kmap

$$f_2 = \sum (2, 3)$$

$$f = x.\bar{y} + x.y$$

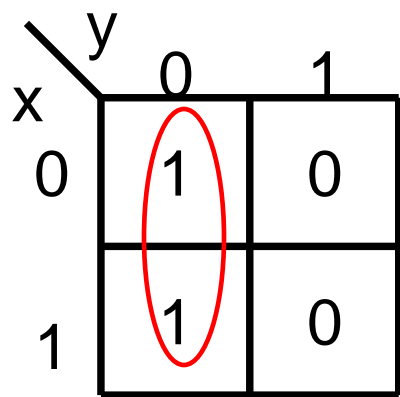
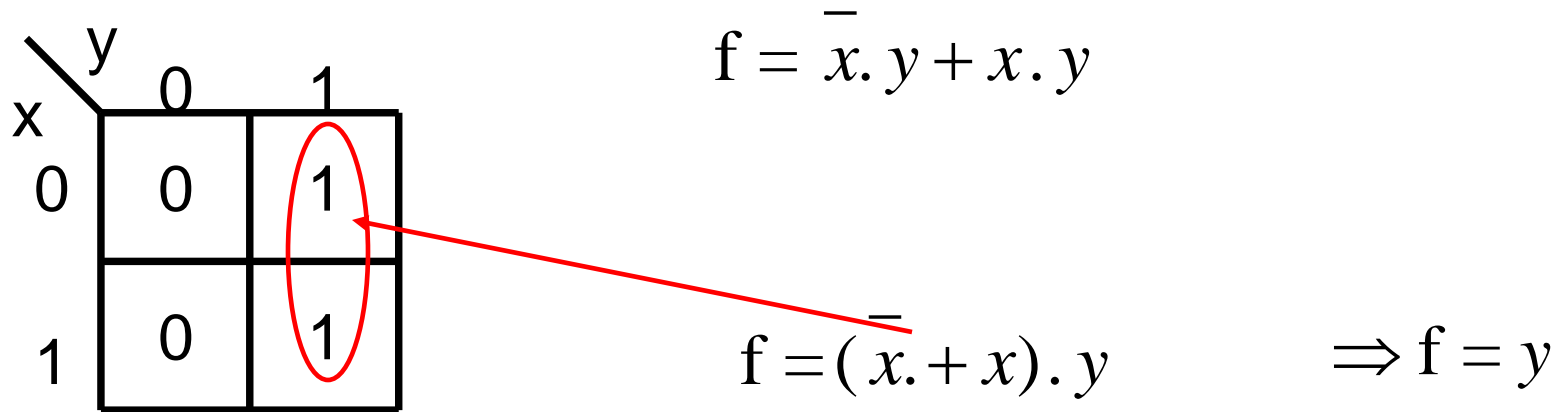
$$f = x.(\bar{y} + y)$$

$$f = x$$

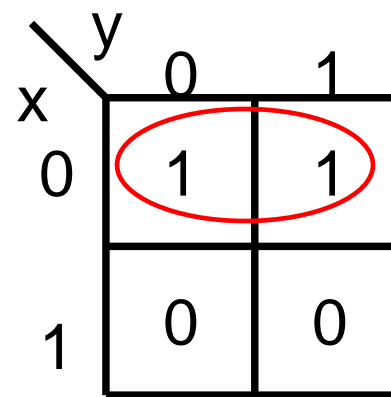


		y	
		0	1
x	0	0	0
	1	1	1

Combine terms which differ in only one bit position. As a result, whatever is common remains.



$$\Rightarrow f = \bar{y}$$



$$\Rightarrow f = \bar{x}$$

$$f_2 = \sum (0, \overset{1}{\cancel{2}}, 3)$$

		y	
		0	1
x	0	0	1
	1	1	1

$$A + A = A$$

$$f = x.\bar{y} + x.y + \bar{x}.y$$

$$\begin{aligned} f &= x.(\bar{y} + y) + \bar{x}.y \\ &= x + \bar{x}.y \end{aligned}$$

$$\begin{aligned} f &= x + \bar{x}.y + x.y \\ &= x + (\bar{x} + x).y \\ &= x + y \end{aligned}$$

The idea is to cover all the 1's with as few and as simple terms as possible

3-variable K-map representation

x	y	z	min terms	
0	0	0	$\bar{x} \cdot \bar{y} \cdot \bar{z}$	m0
0	0	1	$\bar{x} \cdot \bar{y} \cdot z$	m1
0	1	0	$\bar{x} \cdot y \cdot \bar{z}$	m2
0	1	1	$\bar{x} \cdot y \cdot z$	m3
1	0	0	$x \cdot \bar{y} \cdot \bar{z}$	m4
1	0	1	$x \cdot \bar{y} \cdot z$	m5
1	1	0	$x \cdot y \cdot \bar{z}$	m6
1	1	1	$x \cdot y \cdot z$	m7

		yz			
		00	01	11	10
x	0	m ₀	m ₁	m ₃	m ₂
	1	m ₄	m ₅	m ₇	m ₆

x	y	z	f
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1



		yz			
		00	01	11	10
x	0	0	1	1	0
	1	0	1	1	0

yz		00	01	11	10
x	0	1	0	1	0
	1	0	1	1	0

$$f = \bar{x}.\bar{y}.\bar{z} + \bar{x}.y.z + x.\bar{y}.z + x.y.z$$

3-variable minimization

x \ yz	00	01	11	10
0	1	0	1	0
1	0	1	1	0

$x.z$

$$f = \bar{x}.\bar{y}.\bar{z} + \bar{x}.y.z + x.y.z + x.\bar{y}.z$$

$y.z$

$$f = \bar{x}.\bar{y}.\bar{z} + y.z + x.z$$

3-variable minimization

x \ yz	00	01	11	10
0	1	0	0	1
1	0	1	1	0

$x.z$

$$f = \bar{x}.\bar{y}.\bar{z} + \bar{x}.y.\bar{z} + x.y.z + x.\bar{y}.z$$

$\bar{x}.\bar{z}$

$$f = \bar{x}.\bar{z} + x.z$$

3-variable minimization

x \ yz	00	01	11	10
	0	0	0	0
1	1	1	1	1

$x.\bar{y}$

$x.y$

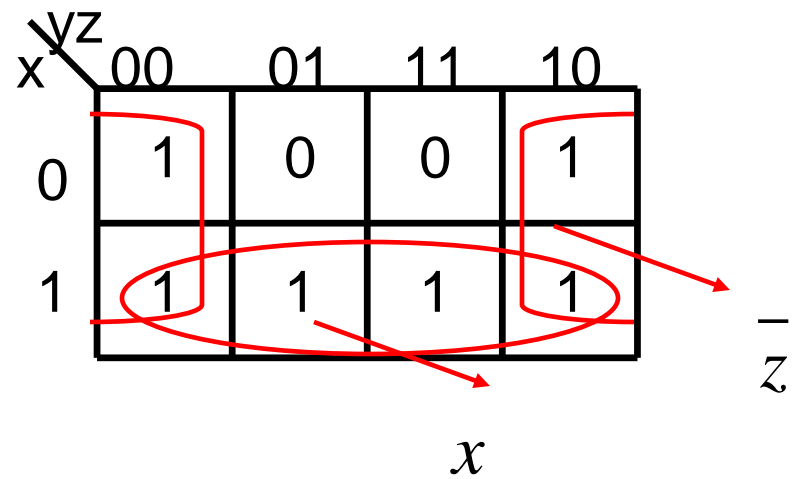
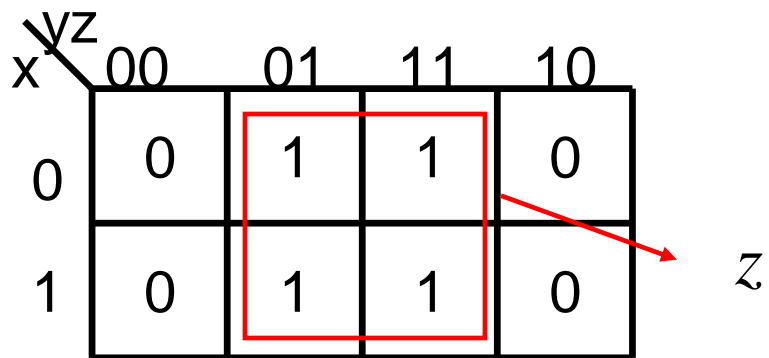
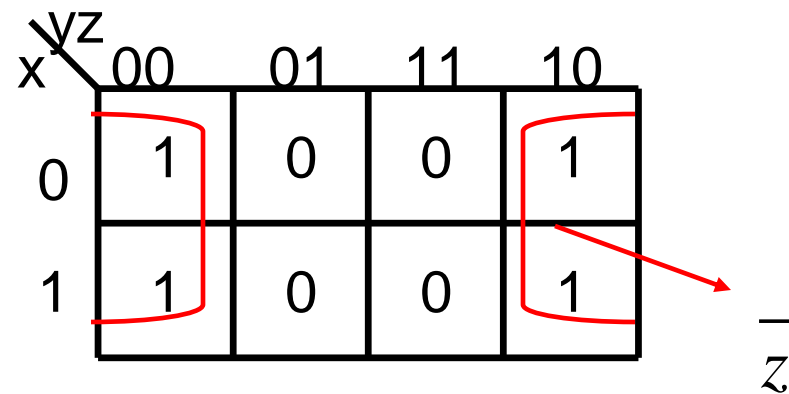
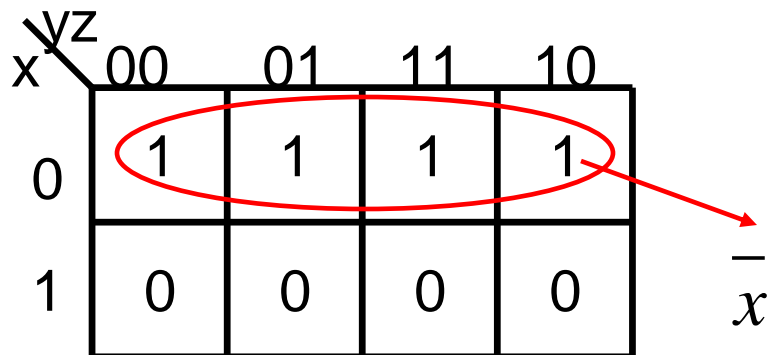
$$f = x.\bar{y}.\bar{z} + x.\bar{y}.z + x.y.z + x.y.\bar{z}$$

$$f = x.\bar{y} + x.y$$

x \ yz	00	01	11	10
	0	0	0	0
1	1	1	1	1

x

$$f = x.(\bar{y} + y) = x$$



$$f = x + \bar{z}$$

Can we do this ?

$\begin{array}{c} yz \\ x \end{array}$		00	01	11	10
		0	0	0	0
0	0	0	0	0	0
1	1	1	1	1	0

Note that each encirclement should represent a single product term. In this case it does not.

$$\begin{aligned} f &= x.\bar{y}.\bar{z} + x.\bar{y}.z + x.y.\bar{z} \\ &= x.\bar{y} + x.z \end{aligned}$$

We do not get a single product term. In general we cannot make groups of 3 terms.

Can we use kmap with the following ordering of variables?

x \ yz	00	01	10	11
0	0	0	0	0
1	0	1	1	0

Can we combine these two terms into a single term ?

$$\begin{aligned}f &= x.\bar{y}.z + x.y.\bar{z} \\ &= x.(\bar{y}.z + y.\bar{z})\end{aligned}$$

Note that no simplification is possible. Kmap requires information to be represented

$x \backslash yz$		00	01	10	11
0	0	1	0	1	
1	0	0	0	0	

These two terms can be combined into a single term but it is not easy to show that on the diagram.

$$\begin{aligned}
 f &= \bar{x}.\bar{y}.z + \bar{x}.y.z \\
 &= \bar{x}.(\bar{y} + y).z = \bar{x}.z
 \end{aligned}$$

Kmap requires information to be represented in such a way that it is easy to apply the principle $x + \bar{x} = 1$

4-variable K-map representation

w	x	y	z	min terms
0	0	0	0	m_0
0	0	0	1	m_1
0	0	1	0	m_2
0	0	1	1	m_3
⋮	⋮	⋮	⋮	⋮
1	1	1	0	m_{14}
1	1	1	1	m_{15}

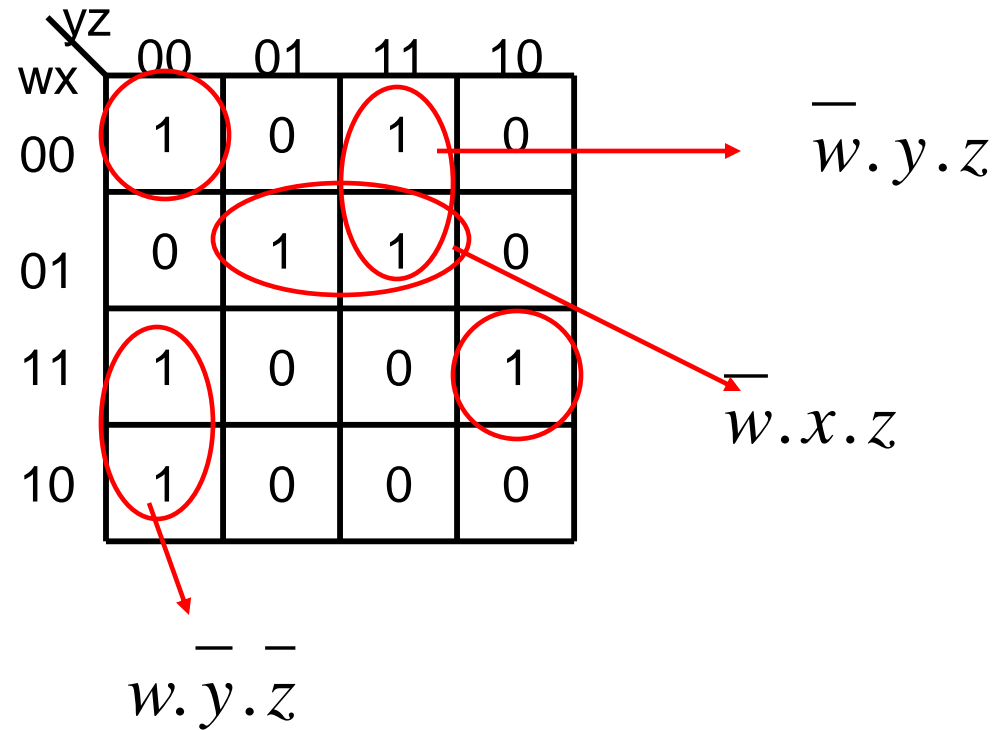


wx \ yz	00	01	11	10
00	0	1	3	2
01	4	5	7	6
11	12	13	15	14
10	8	9	11	10

wx \ yz	00	01	11	10
00	1	0	1	0
01	0	1	1	0
11	1	0	0	1
10	1	0	0	0

$$\begin{aligned}
 f = & \overline{w}. \overline{x}. \overline{y}. \overline{z} + \overline{w}. \overline{x}. y. \overline{z} + \overline{w}. x. \overline{y}. \overline{z} + \overline{w}. x. y. \overline{z} \\
 & + w. \overline{x}. \overline{y}. \overline{z} + w. \overline{x}. y. \overline{z} + w. x. \overline{y}. \overline{z}
 \end{aligned}$$

4-variable minimization



$$f = \overline{w}.y.z + \overline{w}.x.z + w.\overline{y}.\overline{z} + \overline{w}.x.y.\overline{z} + w.x.y.\overline{z}$$

But is this the simplest expression ?

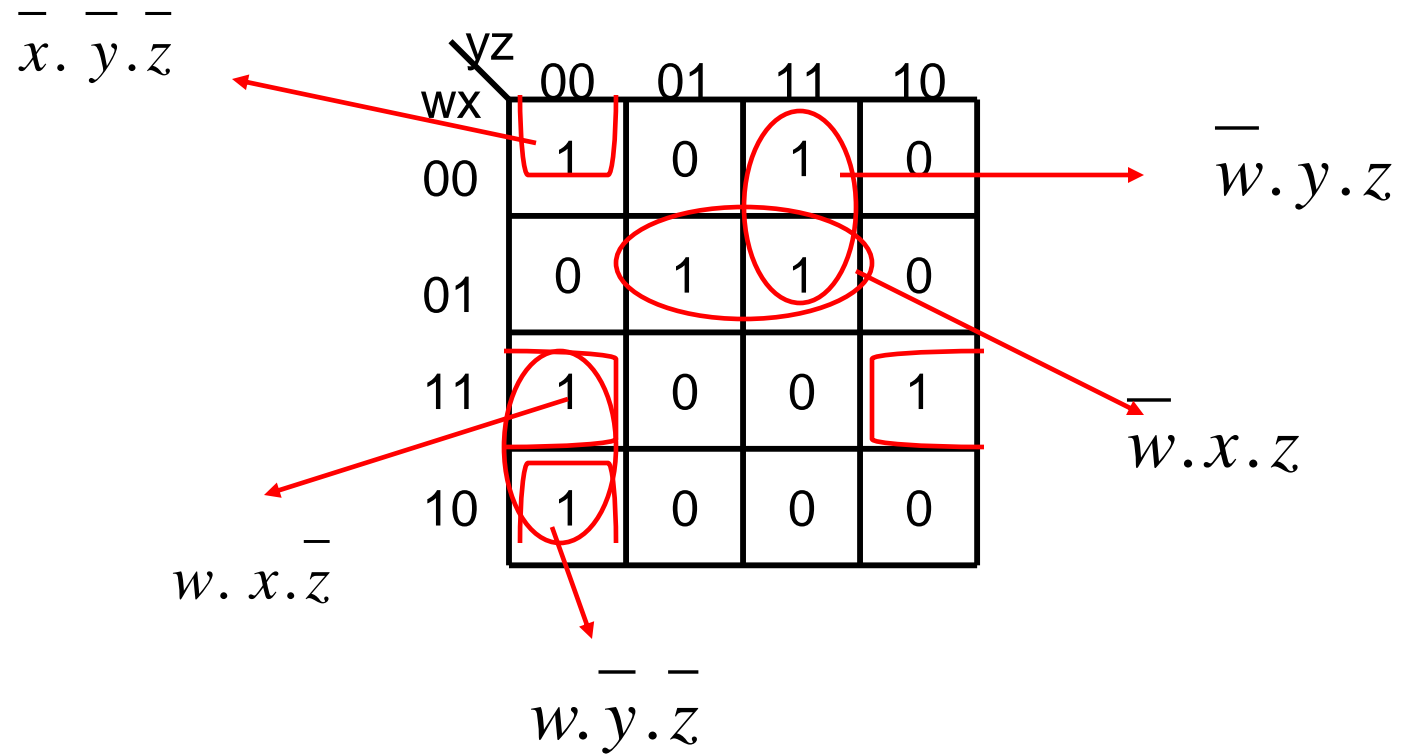
wx \ yz	00	01	11	10
00	1	0	1	0
01	0	1	1	0
11	1	0	0	1
10	1	0	0	0

$$w \cdot x \cdot \bar{y} \cdot \bar{z} + w \cdot x \cdot y \cdot \bar{z} = w \cdot x \cdot \bar{z}$$

wx \ yz	00	01	11	10
00	1	0	1	0
01	0	1	1	0
11	1	0	0	1
10	1	0	0	0

$$w \cdot \bar{x} \cdot \bar{y} \cdot \bar{z} + w \cdot \bar{x} \cdot y \cdot \bar{z} = \bar{x} \cdot \bar{y} \cdot \bar{z}$$

4-variable minimization



$$f = \overline{w}.y.z + \overline{w}.x.z + w.\overline{y}.\overline{z} + w.x.\overline{z} + \overline{x}.\overline{y}.\overline{z}$$

Is this the best that we can do ?

Cover the 1's with minimum number of terms

wx \ yz	00	01	11	10
00	1	0	1	0
01	0	1	1	0
11	1	0	0	1
10	1	0	0	0

$$f = \overline{w}.y.z + \overline{w}.x.z +$$

$$w.\overline{y}.\overline{z} + w.x.\overline{z} + \overline{x}.\overline{y}.\overline{z}$$

wx \ yz	00	01	11	10
00	1	0	1	0
01	0	1	1	0
11	1	0	0	1
10	1	0	0	0

$$f = \overline{w}.y.z + \overline{w}.x.z +$$

$$w.x.\overline{z} + \overline{x}.\overline{y}.\overline{z}$$

4-variable minimization

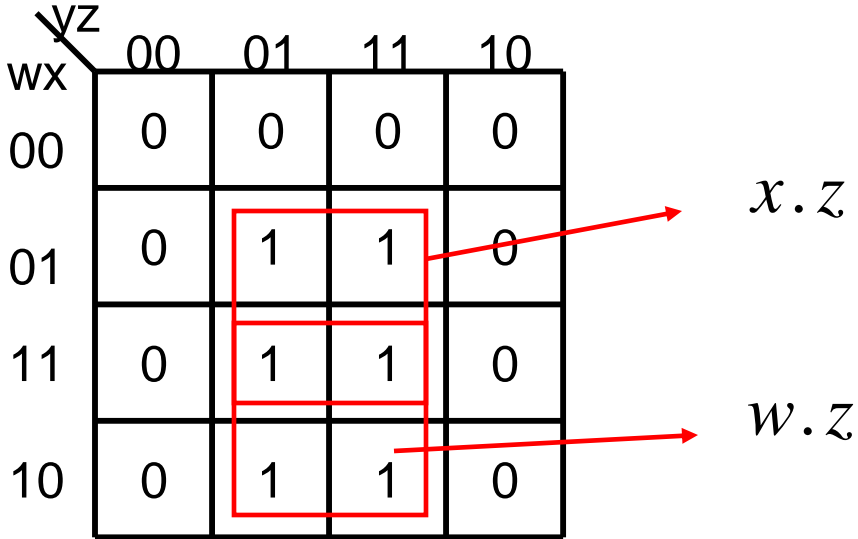
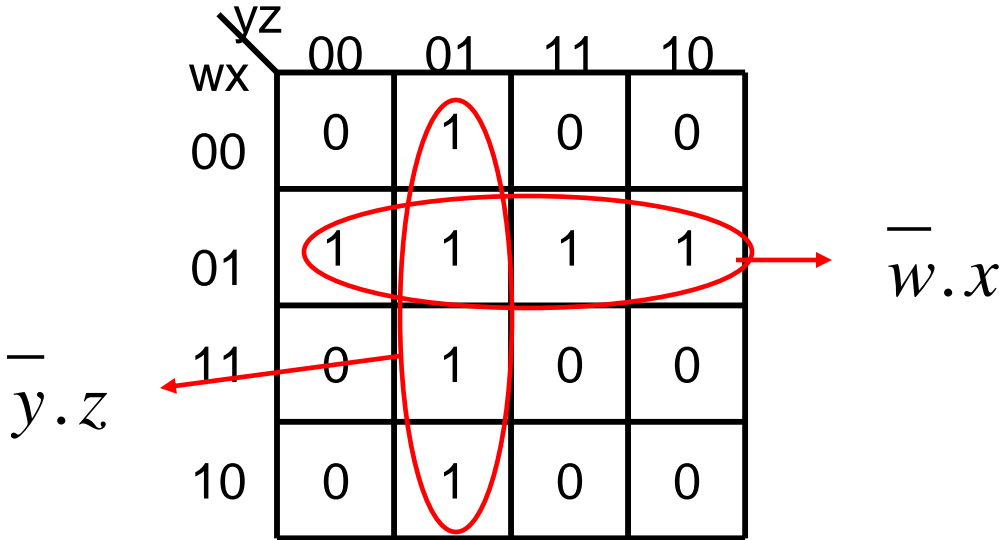
wx \ yz	00	01	11	10
00	1	0	0	0
01	1	1	0	0
11	0	0	0	0
10	1	0	0	1

$$f = \overline{w}.x.\overline{y} + \overline{w}.x.\overline{z} + \overline{w}.\overline{y}.\overline{z}$$

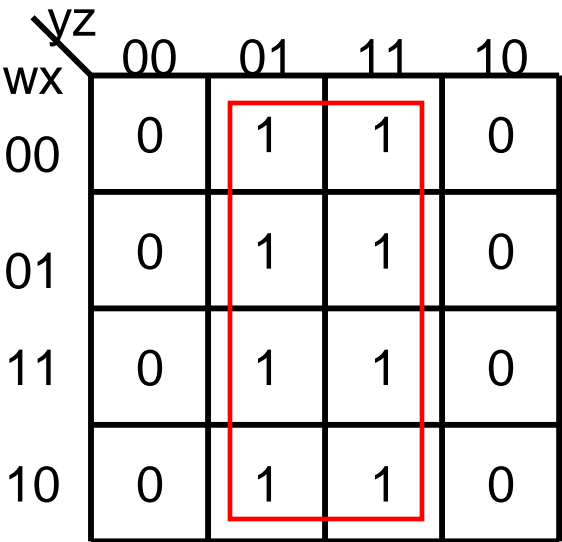
wx \ yz	00	01	11	10
00	1	0	0	0
01	1	1	0	0
11	0	0	0	0
10	1	0	0	1

$$f = \overline{w}.x.\overline{y} + \overline{w}.x.\overline{z} + x.\overline{y}.\overline{z}$$

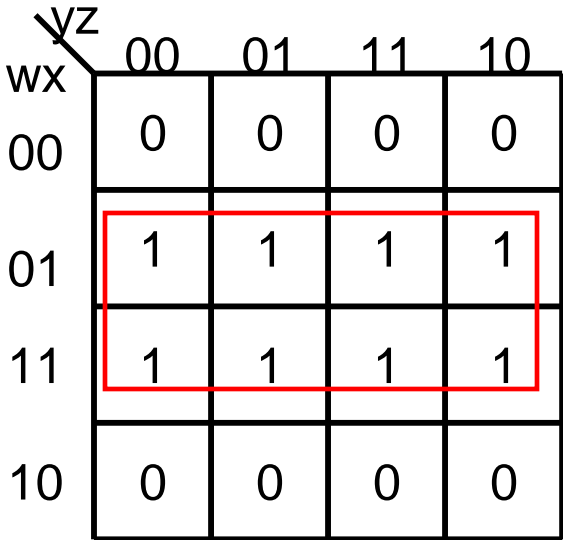
Groups of 4



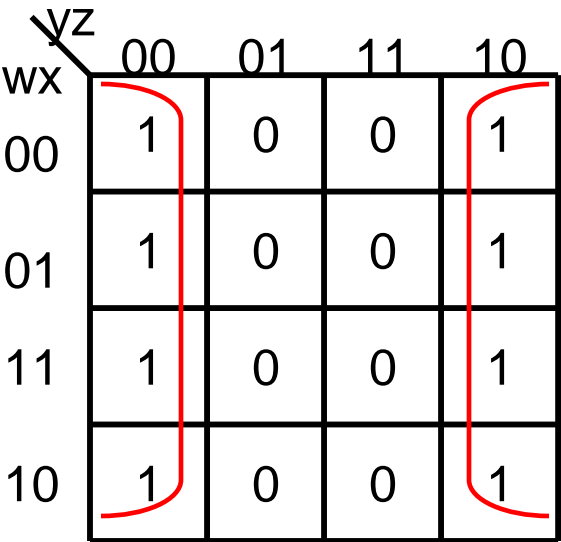
Groups of 8



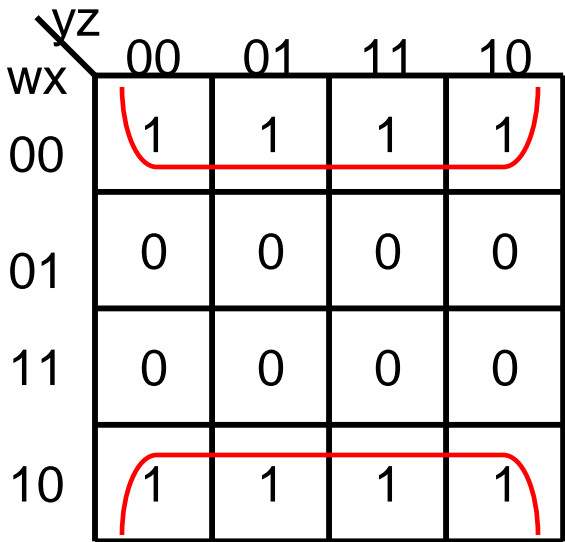
z



x

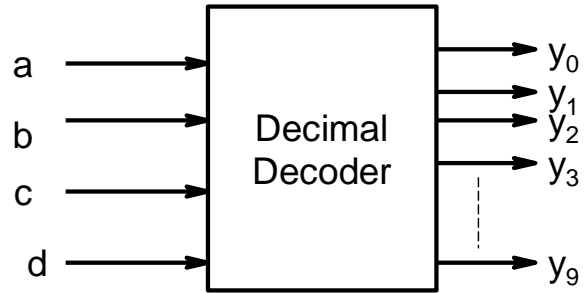


\bar{z}



\bar{x}

Don't care terms



Y_3

	cd	00	01	11	10
ab					
00		0	0	1	0
01		0	0	0	0
11		x	x	x	x
10		0	0	x	x

$$y_3 = \bar{a}.\bar{b}.c.d$$

$$y_3 = \bar{b}.c.d$$

Don't care
conditions

a	b	c	d	y_0	y_1	y_2	y_3	y_4	y_5	y_6	y_7	y_8	y_9
0	0	0	0	1	0	0	0	0	0	0	0	0	0
0	0	0	1	0	1	0	0	0	0	0	0	0	0
0	0	1	0	0	0	1	0	0	0	0	0	0	0
0	0	1	1	0	0	0	1	0	0	0	0	0	0
0	1	0	0	0	0	0	0	1	0	0	0	0	0
0	1	0	1	0	0	0	0	0	1	0	0	0	0
0	1	1	0	0	0	0	0	0	0	1	0	0	0
0	1	1	1	0	0	0	0	0	0	0	1	0	0
1	0	0	0	0	0	0	0	0	0	0	0	1	0
1	0	0	1	0	0	0	0	0	0	0	0	0	1

Don't care terms can be chosen as 0 or 1. Depending on the problem, we can choose the don't care term as 1 and use it to obtain a simpler Boolean expression

Decimal Encoding using
Binary Coded Decimal (BCD)

Y_3

cd \ ab	00	01	11	10
00	0	0	1	0
01	0	0	0	0
11	x	x	x	x
10	0	0	x	x

$$y_3 = \bar{b}.c.d$$



Don't care terms should only be included in encirclements if it helps in obtaining a larger grouping or smaller number of groups.

Don't Care Terms

- Don't care terms can be chosen as 0 or 1
 - Sometimes choosing them as 1 yields simpler expressions in SOP minimization

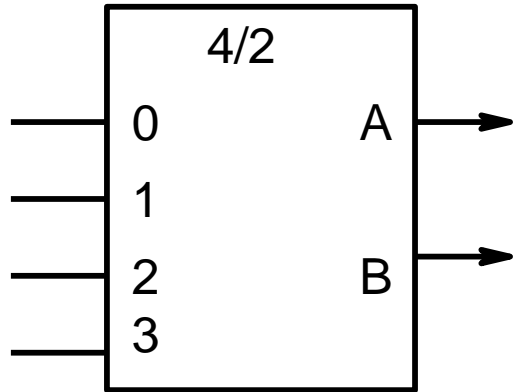
wx \ yz	00	01	11	10
00	1	x	0	1
01	1	0	1	1
11	0	x	1	1
10	1	x	1	x

$$f = \overline{w} \cdot \overline{z} + x \cdot y + w \cdot \overline{x}$$

Don't care terms should only be included in encirclements if it helps in obtaining a larger grouping or smaller number of groups.

Encoder

- Inverse of decoders



d_3	d_2	d_1	d_0	B	A
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	0	1	0
1	0	0	0	1	1

$$A = \overline{d_2} \overline{d_0}$$

$$B = \overline{d_1} \overline{d_0}$$

$$A = \overline{d_2} \overline{d_0}$$

$d_3 \backslash d_1 d_0$	00	01	11	10
00	x	0	x	1
01	0	x	x	x
11	x	x	x	x
10	1	x	x	x

A

$d_3 \backslash d_1 d_0$	00	01	11	10
00	x	0	x	0
01	1	x	x	x
11	x	x	x	x
10	1	x	x	x

B

$$B = \overline{d_1} \overline{d_0}$$

You know what to do 😊

