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SUBJECT	DAA							
EXPERIMENT NO:	03							
AIM:	To perform matrix multiplication using strassen's algorithm.							
PROBLEM STATEMENT 1:								
	The main reason for high time complexity in naive divide and conquer is because of 8 recursive calls. The strassen matrix reduces these recursive calls to 7. The idea is similar to naive divide and conquer i.e divide the matrix in sub-matrices of N/2xN/2 dimension until we get 2x2 matrix. We compute this matrix using formulae provided by strassens algorithm.							
	p1 = a(f - h) $p2 = (a + b)hp3 = (c + d)e$ $p4 = d(g - e)p5 = (a + d)(e + h)$ $p6 = (b - d)(g + h)p7 = (a - c)(e + f)$							
	The A x B can be calculated using above seven multiplications.  Following are values of four sub-matrices of result C							
	$\begin{bmatrix} a & b \\ c & d \end{bmatrix} \mathbf{X} \begin{bmatrix} e & f \\ g & h \end{bmatrix} = \begin{bmatrix} p5 + p4 - p2 + p6 & p1 + p2 \\ \hline p3 + p4 & p1 + p5 - p3 - p7 \end{bmatrix}$							
	$\begin{bmatrix} c & d & 1 \\ \end{bmatrix}$ $\begin{bmatrix} g & h \\ \end{bmatrix}$ $\begin{bmatrix} p3 + p4 \\ p1 + p5 - p3 - p7 \end{bmatrix}$							
	A B C  A, B and C are square metrices of size N x N  a, b, c and d are submatrices of A, of size N/2 x N/2  e, f, g and h are submatrices of B, of size N/2 x N/2  p1, p2, p3, p4, p5, p6 and p7 are submatrices of size N/2 x N/2							

As there are 7 recursive calls we get following recurrence-

 $T(N) = 7T(N/2) + O(N^2)$ 

Solving above recurrence we get time complexity of  $O(N^2.81)$  which is less than naive approach.

Disadvantages of strassens algorithm-

- 1.)The constants used in strassens multiplication are high and naive approach is used for typical application.
- 2.)For sparse matrix there are better algorithms.
- 3.)The sub-matrices in recursion take extra space.
- 4.)Because of limited percision of computer arithemetic on non integer values, large errors accumulate in strassens algorithm.

## **ALGORITHM**

1.)Divide the input matrices A and B into N/2xN/2 dimension. After dividing we get sub-matrices A11,A12,A21,A22 and B11,B12,B21,B22.

2.)Computing sub-matrices using following formulae-

M1=(A11+A22)\*(B11+B22)

M2=(A21+A22)\*B11

M3=A11\*(B12-B22)

M4=A22\*(B21-B11)

M5=(A11+A12)\*B22

M6=(A21-A11)\*(B11+B12)

M7=(A12-A22)\*(B21+B22)

3.)Compute the output using intermediate matrices

C11=M1+M4-M5+M7

C12=M3+M5

C21=M2+M4

C22=M1-M2+M3+M6

4.)Combine the output matrices for final matrix.

## **PROGRAM:**

```
#include<iostream>
     using namespace std;
     int main(){
         int A[2][2],B[2][2];
cout << "Enter 4 elements for matrix A: ";
for(int i=0;i<2;i++)</pre>
              for(int j=0;j<2;j++)
                  cin \gg A[i][j];
         cout << "Enter 4 elements for matrix B: ";</pre>
         for(int i=0;i<2;i++)
              for(int j=0;j<2;j++)
                  cin \gg B[i][j];
         int p[7];
         p[0] = A[0][0]*(B[0][1] - B[1][1]);
         p[1] = (A[0][0] + A[0][1])*B[1][1];
         p[2] = (A[1][0] + A[1][1])*B[0][0];
          p[3] = A[1][1]^*(B[1][0] - B[0][0]); 
 p[4] = (A[0][0] + A[1][1])^*(B[0][0] + B[1][1]); 
         p[5] = (A[0][1] - A[1][1])*(B[1][0] + B[1][1]);
         p[6] = (A[0][0] - A[1][0])*(B[0][0] + B[0][1]);
          int C[2][2];
         C[0][0] = p[3]+p[4]-p[1]+p[5];
         C[0][1] = p[0] + p[1];
         C[1][0] = p[2] + p[3];
         C[1][1] = p[0] + p[4] - p[2] - p[6];
         cout << "Multiplication of A and B using strassen's algorithm: " << endl;</pre>
          for(int i=0;i<2;i++){
              for(int j=0;j<2;j++){
                  cout << C[i][j] << "\t";
              cout << endl;</pre>
```

## **RESULT ( SNAPSHOT)**

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Through this experiment I understood how to compute matrices product using strassens algorithm. The complexity of this alogrithm was found out to be  $O(N^2.81)$  which is slightly better than naive approach whose time complexity is  $O(N^3)$ . This can have huge impact on matrices with high dimensions.