Measuring Runtime Performance

Complexity Notation

n = input size
 f(n) = max number of steps when input has size n
 O(f(n)) = asymptotic upper bound

```
void f(int *out, const int *in, int size) {
  int product = 1;
  for (int i = 0; i < size; ++i)
    product *= in[i];
  for(int i = 0; i < size; ++i)
    out[i] = product / in[i];
  } // f()

f(n) = 1 + 1 + 1 + 3n + 1 + 1 + 3n = 6n + 5 = O(n)</pre>
```

Ways to measure complexity

Analytically

- Analysis of the code itself
- Recognizing common algorithms/patterns
- Based on a recurrence relation

Empirically

- Measure runtime programmatically
- Measure runtime using external tools
- Test on inputs of a variety of sizes

Measuring Runtime Programmatically

- "Programmatically" measurements are taken from inside the code itself
- Varies greatly depending on the language
- Many different ways to do it even just in C/C++!

Measuring Time In C++11+

```
class Timer {
     std::chrono::time point<std::chrono::system clock> cur;
     std::chrono::duration<double> elap;
   public:
                                                                        Example
     Timer() : cur(), elap(std::chrono::duration<double>::zero()) {}
                                                                 int main() {
     void start() {
                                                              26
                                                                   Timer t;
10
       cur = std::chrono::system clock::now();
                                                              27
                                                                   t.start();
11
     } // start()
                                                              28
12
                                                                   t.stop();
13
     void stop() {
                                                              30
       elap += std::chrono::system clock::now() - cur;
14
                                                              31
15
     } // stop()
                                                              32
16
                                                              33
                                                                   t.reset();
17
     void reset() {
                                                              34
                                                                   t.start();
18
       elap = std::chrono::duration<double>::zero();
     } // reset()
19
                                                                   t.stop();
20
21
     double seconds() {
       return elap.count();
                                                                   return 0;
     } // seconds()
                                                                 } // main()
   }; // Timer{}
```

Note: Checking time too often will slow down your program!

doStuff1(); cout << "1: " << t.seconds() << "s" << endl: doStuff2(); cout << "2: " << t.seconds() << "s" << endl;

#include <chrono>

Let's try it!

Save the file to a folder you can access from a *NIX shell, and/or upload to CAEN

Browser Download

https://eecs281staff.github.io/search-demo/search.cpp

- Command Line Download
 - \$ mkdir search-demo && cd search-demo
 - \$ wget https://eecs281staff.github.io/search-demo/search.cpp



EECS 281: Search Demo

After Downloading

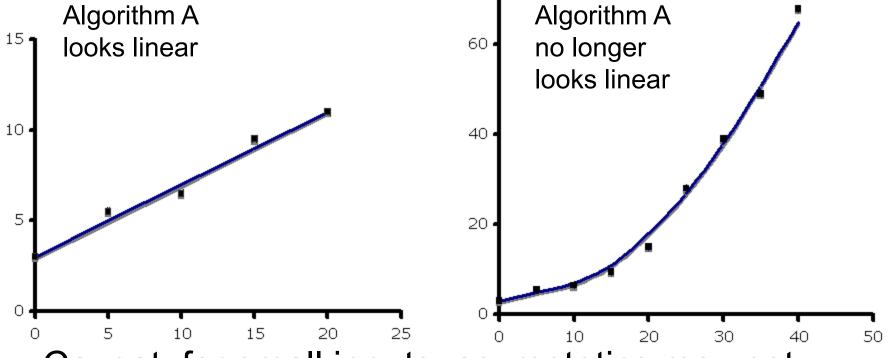
Compile:

```
$ g++ -std=c++17 -03 search.cpp -o search
```

- Run a binary search, 1M items:
 - \$./search b 1000000
- Run a linear search, 1M items:
 - \$./search 1 1000000
- Try with larger numbers!

Empirical Results

- Plot actual run time versus varying input sizes
- Include a large range to accurately display trend



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Caveat: for small inputs, asymptotics may not play out yet

Prediction versus Experiment

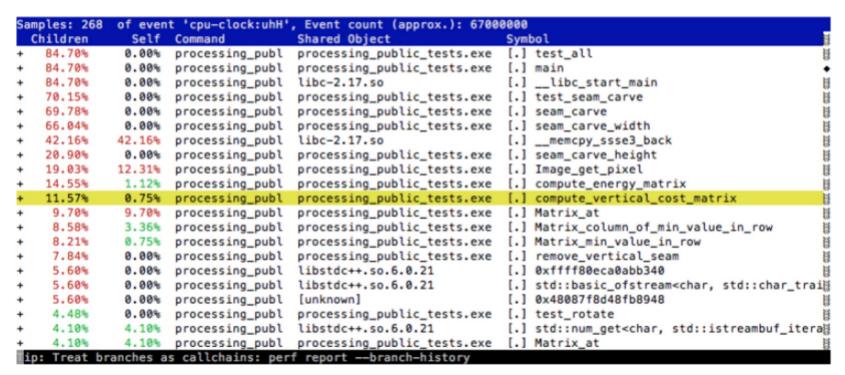
- What if experimental results are worse than predictions?
 - Example: results are exponential when analysis is linear
 - Error in complexity analysis
 - Error in coding (check for extra loops, unintended operations, etc.)
- What if experimental results are better than predictions?
 - Example: results are linear when analysis is exponential
 - Experiment may not have fit worst case scenario
 - Error in complexity analysis
 - Error in analytical measurements
 - Incomplete algorithm implementation
 - Algorithm implemented is better than the one analyzed
- What if experimental data match asymptotic prediction but runs are too slow?
 - Performance bug?
 - Check compile options (e.g. use -03)
 - Look for optimizations to improve the constant factors

Measuring Runtime Performance

Runtime Analysis Tools

Using a Profiling Tool

- This won't tell you the complexity of an algorithm, but it tells you where you program spends its time.
- Many different tools exist you'll learn to use perf in lab.



A snapshot of a perf report generated for EECS 280 Project 2. Image credit: Alexandra Brown

Measuring Runtime on Linux

If you are launching a program using command

% progName -options args

Then

% /usr/bin/time progName -options args
will produce a runtime report

0.84user 0.00system 0:00.85elapsed 99%CPU

If you're timing a program in the current folder, use ./

% /usr/bin/time ./progName -options args

Often, you can just type time rather than /usr/bin/time.

Measuring Runtime on Linux

 Example: this command just wastes time by copying zeros to the "bit bucket"

% time dd if=/dev/zero of=/dev/null

kill it with control-C

```
3151764+0 records in
3151764+0 records out
1613703168 bytes (1.6 GB) copied, 0.925958 s, 1.7 GB/s
Command terminated by signal 2
0.26user 0.65system 0:00.92elapsed 99%CPU
(0avgtext+0avgdata 3712maxresident)k
0inputs+0outputs (0major+285minor)pagefaults 0swaps
```

Measuring Runtime on Linux

- 0.26user 0.65system 0:00.92elapsed 99%CPU
- user time is spent by your program
- system time is spent by the OS on behalf of your program
- elapsed is wall clock time time from start to finish of the call, including any time slices used by other processes
- %CPU Percentage of the CPU that this job got. This is just (user + system) / elapsed
- man time for more information

Using valgrind

Suppose we want to check for memory leaks:

```
valgrind ./search b 1000000
```

- Force a leak!
 - Replace return 0 with exit(0), run valgrind using
 flags --leak-check=full --show-leak-kinds=all
- Who leaked that memory?
 - The memory address isn't very useful, we just know that main() called operator new
 - Recompile with -g3 instead of -03 and run valgrind one more time

Runtime Analysis Tools

Analyzing Recursion

Job Interview Question

Implement this function

```
// returns x^n
int power(int x, uint32_t n);
```

- The obvious solution using n-1 multiplications is O(n)
 - $-2^8 = 2*2*...*2$
- Less obvious: O(log n) multiplications
 - Hint: $2^8 = ((2^2)^2)^2$
 - How does it work for 2^7 ?
- Write both solutions iteratively and recursively

Computing xⁿ

```
int power(int x, uint32_t n) {
     if (n == 0) {
      return 1;
  } // if
     int result = x;
     for (int i = 1; i < n; ++i) {
8
      result = result * x;
     } // for
10
  return result;
12 } // power()
```

Analyzing Solutions

- Iterative functions use loops
- A function is recursive if it calls itself
- What is the time complexity of each function?

Iterative

```
int power(int x, int n) {
   int result = 1;
   for (int i = 0; i < n; ++i) {
     result = result * x;
   } // for()

return result;
} // power()

@(n)
It's just a regular loop.</pre>
```

Recursive

```
int power(int x, int n) {
if (n == 0) {
   return 1;
} // if()

return x * power(x, n - 1);

// power()

???
We need another tool to analyze this.
```

Recurrence Relations

- A recurrence relation describes the way a problem depends on a subproblem.
 - A recurrence can be written for a computation:

$$x^n = \begin{cases} 1 & n == 0 \\ x * x^{n-1} & n > 0 \end{cases}$$

– A recurrence can be written for the time taken:

$$T(n) = \begin{cases} c_0 & n == 0 \\ T(n-1) + c_1 & n > 0 \end{cases}$$

– A recurrence can be written for the amount of memory used*:

$$M(n) = \begin{cases} c_0 & n == 0\\ M(n-1) + c_1 & n > 0 \end{cases}$$

*Non-tail recursive

Solving Recurrences

- Substitution method
 - 1. Write out T(n), T(n 1), T(n 2)
 - 2. Substitute T(n-1), T(n-2) into T(n)
 - 3. Look for a pattern
 - 4. Use a summation formula
- Another way to solve recurrence equations is the Master Method (AKA Master Theorem)

Solving Recurrences: Linear

$$T(n) = \begin{cases} c_0 & n == 0 \\ T(n-1) + c_1 & n > 0 \end{cases}$$
 int power(int 2 if (n == 0) 3 return 1;

```
1  int power(int x, int n) {
2   if (n == 0)
3    return 1;
4
5   return x * power(x, n - 1);
6  } // power()
```

Recurrence: T(n) = T(n - 1) + cComplexity: $\Theta(n)$

Solving Recurrences: Logarithmic

$$T(n) = \begin{cases} c_0 & n = 0 \\ T(\frac{n}{2}) + c_1 & n > 0 \end{cases}$$

```
1  int power(int x, int n) {
2    if (n == 0)
3     return 1;
4
5    int result = power(x, n / 2);
6    result *= result;
7    if (n % 2 != 0) // n is odd
8     result *= x;
9
10    return result;
11 } // power()
```

Recurrence: T(n) = T(n/2) + cComplexity: $\Theta(\log n)$

A Logarithmic Recurrence Relation

$$T(n) = \begin{cases} c_0 & n == 0 \\ T(\frac{n}{2}) + c_1 & n > 0 \end{cases} \rightarrow \mathcal{O}(\log n)$$

- Fits the logarithmic recursive implementation of power()
 - The power to be calculated is divided into two halves and combined with a single multiplication
- Also fits Binary Search
 - The search space is cut in half each time, and the function recurses into only one half

Recurrence Thought Exercises

- What if a recurrence cuts a problem into two subproblems, and both subproblems were recursively processed?
- What if a recurrence cuts a problem into three subproblems and...
 - Processes one piece
 - Processes two pieces
 - Processes three pieces
- What if there was additional, non-constant work after the recursion?

Binomial Coefficient

$$\binom{n}{k} = \frac{n!}{k! (n-k)!}$$

- Binomial Coefficient "n choose k"
- Write this function with pen and paper
- Compile and test what you've written
- Options
 - Iterative
 - Recursive
 - Tail recursive
- Analyze

Analyzing Recursion