

Lab Problems

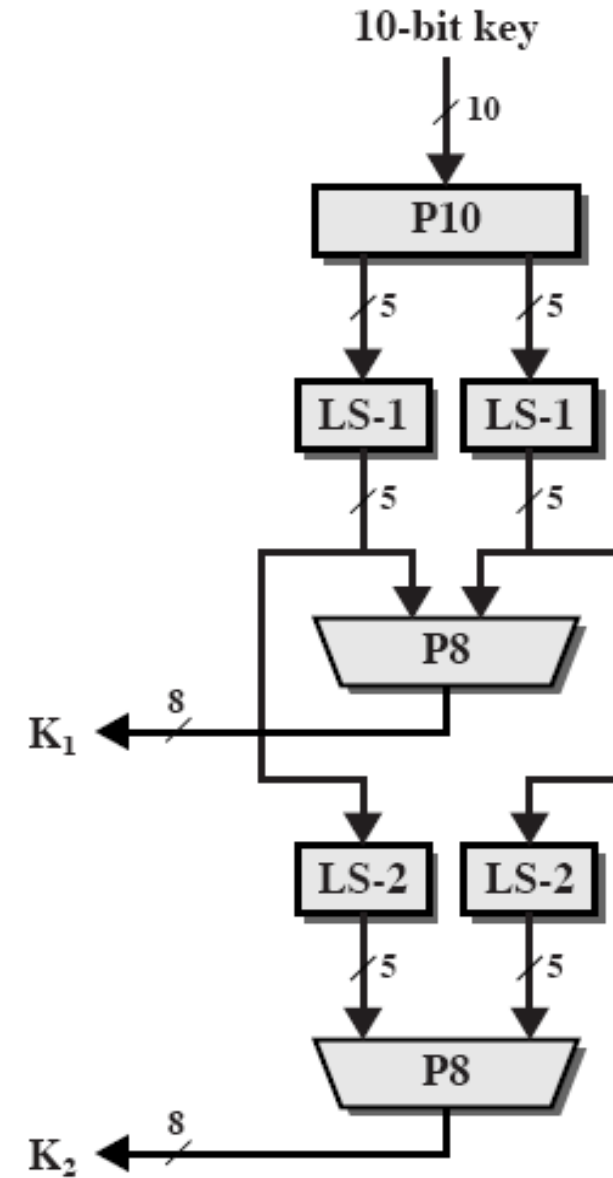
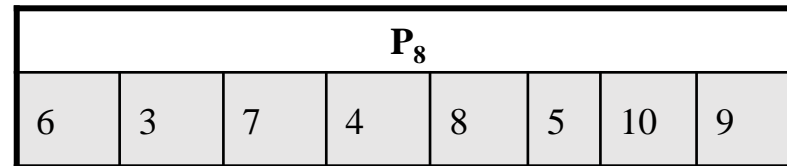
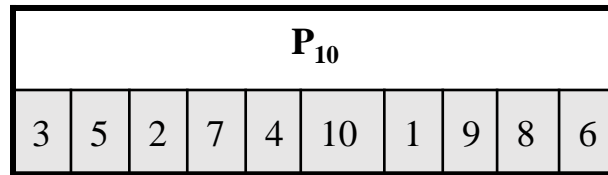
CSE-312

Lab Problem-12

Implement the Key Generation for S-DES

Input: 10 bit key

Output: K1 and K2



Lab Problem-13

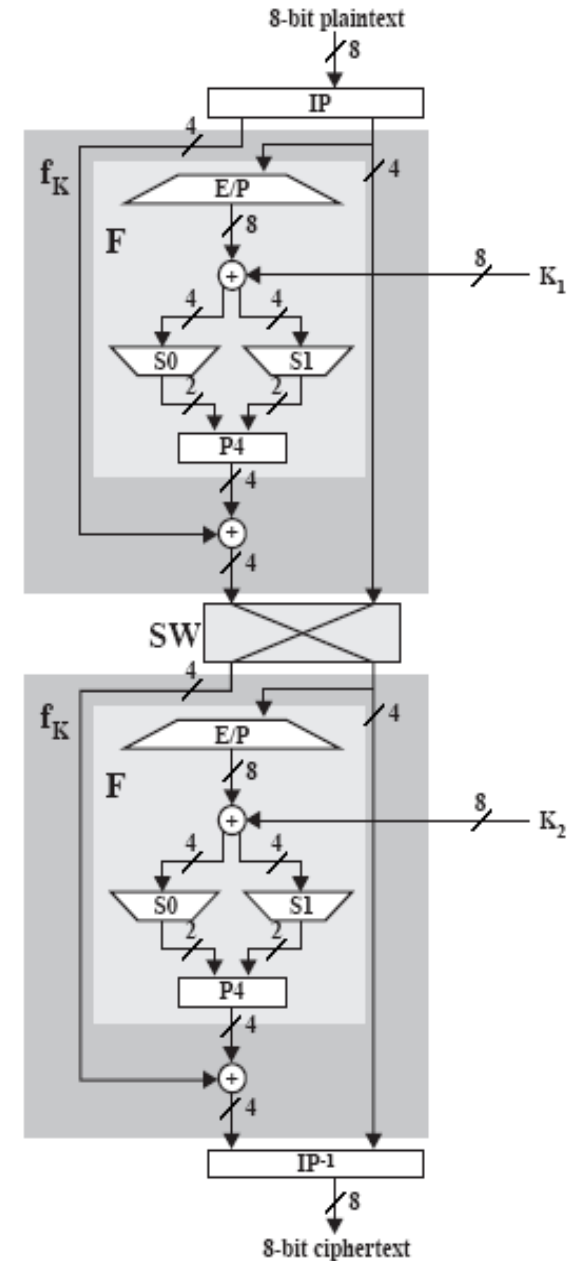
Implement the S-DES Encryption

IP							
2	6	3	1	4	8	5	7

$$S0 = \begin{bmatrix} 1 & 0 & 3 & 2 \\ 3 & 2 & 1 & 0 \\ 0 & 2 & 1 & 3 \\ 3 & 1 & 3 & 2 \end{bmatrix}, \quad P4 (2431)$$

$$S1 = \begin{bmatrix} 0 & 1 & 2 & 3 \\ 2 & 0 & 1 & 3 \\ 3 & 0 & 1 & 0 \\ 2 & 1 & 0 & 3 \end{bmatrix}$$

IP ⁻¹							
4	1	3	5	7	2	8	6



Lab Problem-14

Implement the S-DES Decryption

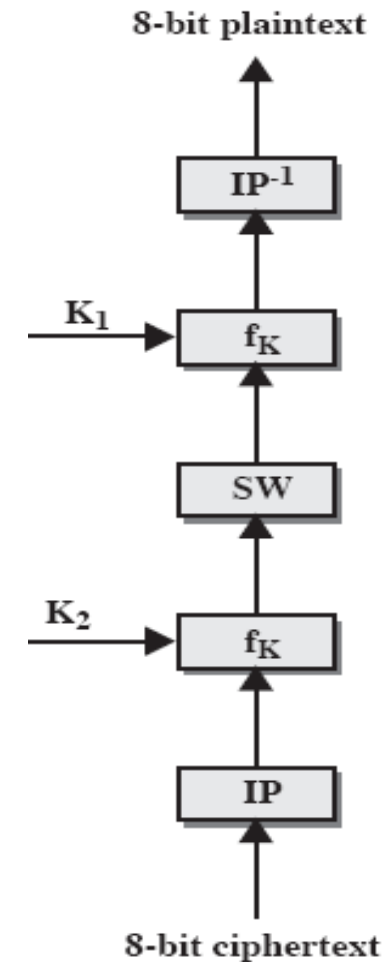
IP							
2	6	3	1	4	8	5	7

$$S0 = \begin{bmatrix} 1 & 0 & 3 & 2 \\ 3 & 2 & 1 & 0 \\ 0 & 2 & 1 & 3 \\ 3 & 1 & 3 & 2 \end{bmatrix}, \quad \text{P4 (2431)}$$

$$S1 = \begin{bmatrix} 0 & 1 & 2 & 3 \\ 2 & 0 & 1 & 3 \\ 3 & 0 & 1 & 0 \\ 2 & 1 & 0 & 3 \end{bmatrix}$$

IP ⁻¹							
4	1	3	5	7	2	8	6

DECRYPTION



Lab Problem-15

1. Implement the Euclidean Algorithm for integers and polynomials

EXTENDED EUCLID(m, b)

1. (A_1, A_2, A_3) = ($1, 0, m$);

(B_1, B_2, B_3) = ($0, 1, b$)

2. **if** $B_3 = 0$

return $A_3 = \gcd(m, b)$; no inverse

3. **if** $B_3 = 1$

return $B_3 = \gcd(m, b)$; $B_2 = b^{-1} \bmod m$

4. $Q = A_3 \text{ div } B_3$

5. (T_1, T_2, T_3) = ($A_1 - Q B_1, A_2 - Q B_2, A_3 - Q B_3$)

6. (A_1, A_2, A_3) = (B_1, B_2, B_3)

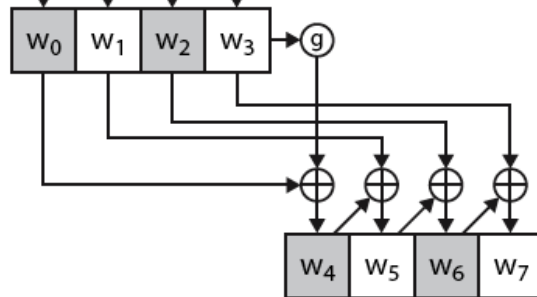
7. (B_1, B_2, B_3) = (T_1, T_2, T_3)

8. **goto** 2

Lab Problem-16

Implement AES Key Expansion

k_0	k_4	k_8	k_{12}
k_1	k_5	k_9	k_{13}
k_2	k_6	k_{10}	k_{14}
k_3	k_7	k_{11}	k_{15}



• $w[0] = (54, 68, 61, 74)$, $w[1] = (73, 20, 6D, 79)$, $w[2] = (20, 4B, 75, 6E)$, $w[3] = (67, 20, 46, 75)$

• $g(w[3])$:

- circular byte left shift of $w[3]$: $(20, 46, 75, 67)$
- Byte Substitution (S-Box): $(B7, 5A, 9D, 85)$
- Adding round constant $(01, 00, 00, 00)$ gives: $g(w[3]) = (B6, 5A, 9D, 85)$

• $w[4] = w[0] \oplus g(w[3]) = (E2, 32, FC, F1)$:

0101 0100	0110 1000	0110 0001	0111 0100
1011 0110	0101 1010	1001 1101	1000 0101
1110 0010	0011 0010	1111 1100	1111 0001
E2	32	FC	F1

• $w[5] = w[4] \oplus w[1] = (91, 12, 91, 88)$, $w[6] = w[5] \oplus w[2] = (B1, 59, E4, E6)$,
 $w[7] = w[6] \oplus w[3] = (D6, 79, A2, 93)$

• first roundkey: **E2 32 FC F1 91 12 91 88 B1 59 E4 E6 D6 79 A2 93**

Lab Problem-17

Implementation of RC4

```
/* Initialization */  
for i = 0 to 255 do  
  S[i] = i;  
  T[i] = K[i mod keylen];
```

```
/* Initial Permutation of S */
```

```
  j = 0;  
  for i = 0 to 255 do  
    j = (j + S[i] + T[i]) mod 256;  
    Swap (S[i], S[j]);
```

```
/* Stream Generation */
```

```
i, j = 0;  
while (true)  
  i = (i + 1) mod 256;  
  j = (j + S[i]) mod 256;  
  Swap (S[i], S[j]);  
  t = (S[i] + S[j]) mod 256;  
  k = S[t];
```

Lab Problem-18

Implement **RSA** Algorithm encryption :

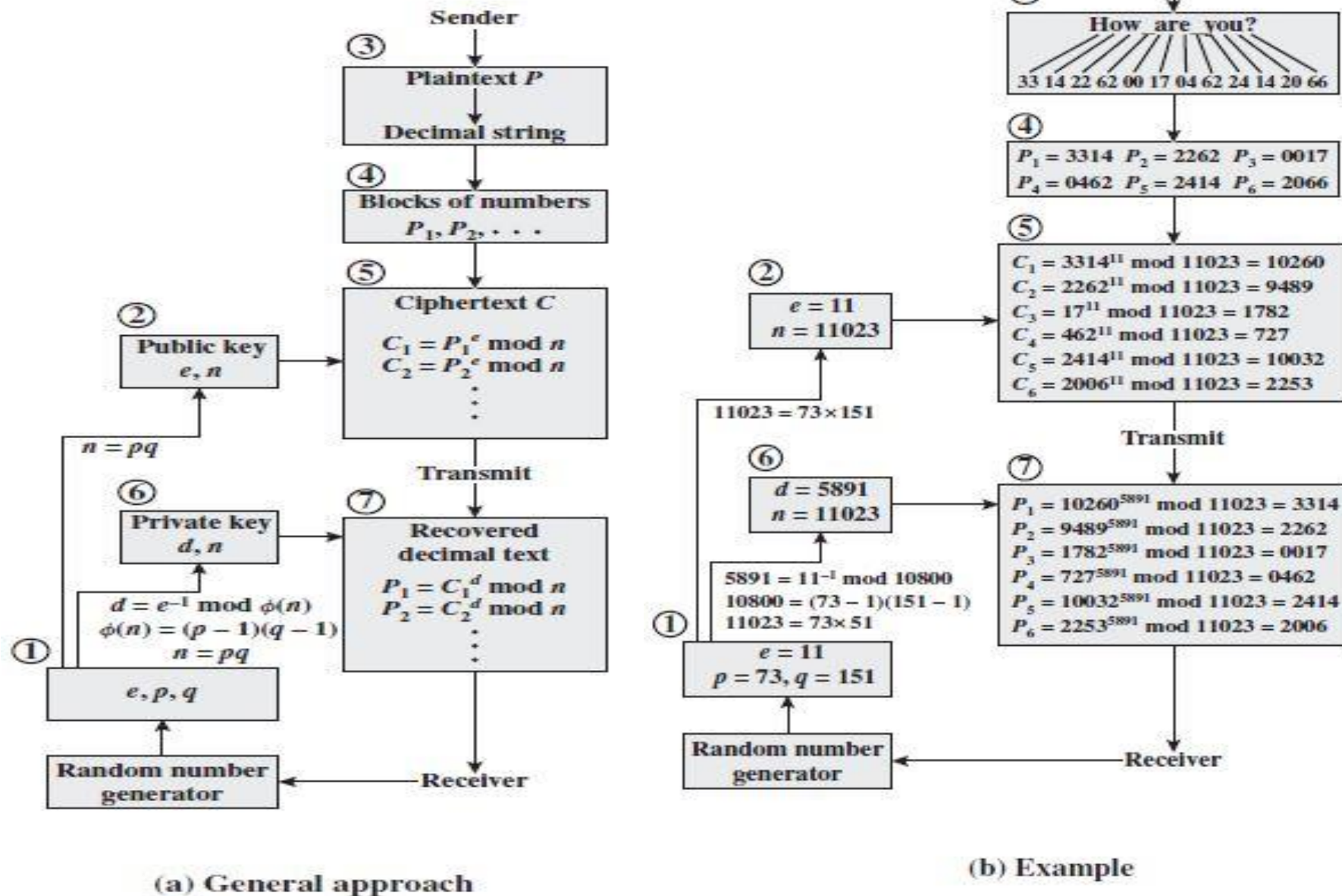


Figure 9.7 RSA Processing of Multiple Blocks