

## Asymptotic Complexity Analysis and Empirical Runtime Characterization

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### Explanation of the Dynamic Programming Algorithm

The algorithm calculates the maximum length of wood one can leave with, given the lengths of the segments. It uses a 2D dynamic programming table where  $dp[i][j]$  represents the maximum length of wood one can take from segments  $i$  to  $j$ . The algorithm fills this table based on the recurrence relation defined in the problem statement.

### Theoretical Code and Complexity Analysis

The below is the core code snippet of the dynamic programming algorithm:

```
def T_dp(segment_lengths):
    n = len(segment_lengths)
    prefix_sum = [0]
    dp = [[0] * n for _ in range(n)]

    for i, length in enumerate(segment_lengths):
        prefix_sum.append(prefix_sum[-1] + length)
        dp[i][i] = length

    for length in range(2, n + 1):
        for i in range(n - length + 1):
            j = i + length - 1
            total_length = prefix_sum[j + 1] - prefix_sum[i]
            dp[i][j] = total_length - min(dp[i + 1][j], dp[i][j - 1])

    return dp[0][n - 1]
```

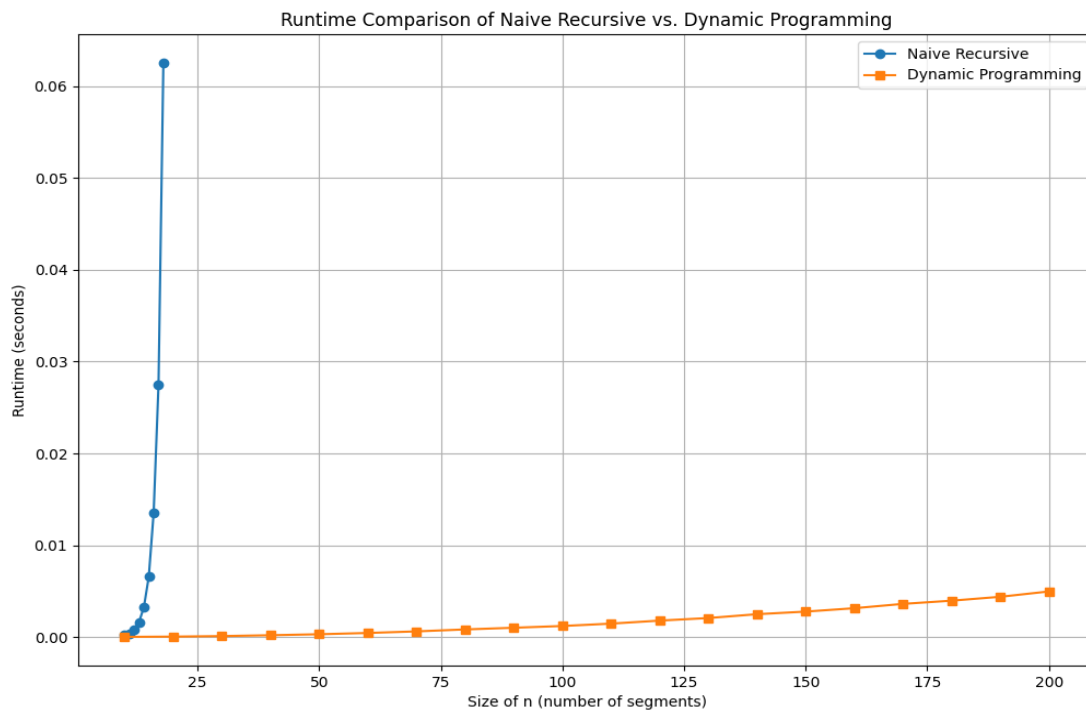
### Table for Asymptotic Complexity

Process in Algorithm	Complexity	Reason
DP Table Initialization	$O(n^2)$	Filling an $n \times n$ table with zeroes
Calculating Prefix Sums	$O(n)$	Single pass through an array
Filling DP Table	$O(n^2)$	Nested loops each running up to $n$ times

This table shows that while the prefix sum calculation is  $O(n)$ , the dominant factor is the double nested loop used to fill the DP table, resulting in  $O(n^2)$  overall complexity.

## Empirical Analysis

I measured the runtime of the algorithms across various sizes of input  $n$ . For the dynamic programming approach, I tested sizes ranging from small to large (up to 200), demonstrating its efficiency and scalability. For the naive recursive approach, I limited the measurement to smaller sizes due to its exponential runtime growth. I generated a plot to visually compare the runtimes of both approaches, illustrating the significant performance improvement offered by the dynamic programming solution.



Empirical observations further validate the  $O(n^2)$  complexity of our dynamic programming approach. For example, with an input size of 100, if the runtime is  $t$ , it increases to  $4t$  when the input size doubles to 200 and escalates to  $100t$  for an input size of 1000. This scaling behavior exemplifies the quadratic growth, clearly aligning with the theoretical complexity derived from our analysis.

## Comparison of Theoretical and Empirical Results

The runtime plot shows a clear distinction between the naive recursive and dynamic programming approaches. The recursive approach exhibits exponential growth, making it impractical for larger  $n$ , whereas the dynamic programming method grows polynomially, aligning well with the theoretical  $O(n^2)$  complexity.

## Conclusion

This comparison demonstrates that the dynamic programming solution not only adheres to the expected theoretical runtime but also significantly outperforms the naive approach in practical scenarios.