

Computer Security Lecture 5



Simplified Advanced Encryption Standard

Dr. Mohamed Loey

Lecturer, Faculty of Computers and Information Benha University

Egypt

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Simplified Advanced Encryption Standard

□ Simplified AES (S-AES) was developed by Professor Edward Schaefer of Santa Clara University in 2003

- ☐ its purpose is educational, since its key and block size are very small 16bits
- it is possible for students to encrypt or decrypt a block doing all operations by hand
- it easier for students to understand the structure AES

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Simplified Advanced Encryption Standard

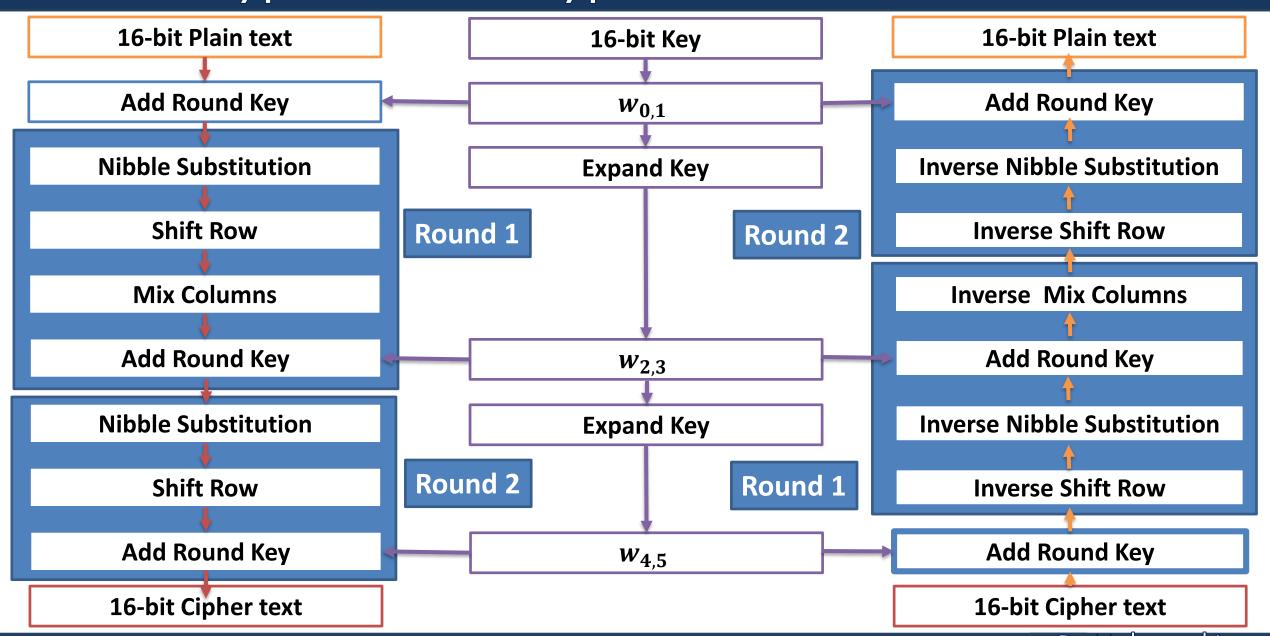
S-AES Encryption and Decryption

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S-AES Encryption and Decryption



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S-AES Encryption Example

□ 16-bit Plaintext, P= D7 28

=1101 0111 0010 1000

□ 16-bit Key, K= 4A F5

=0100 1010 1111 0101

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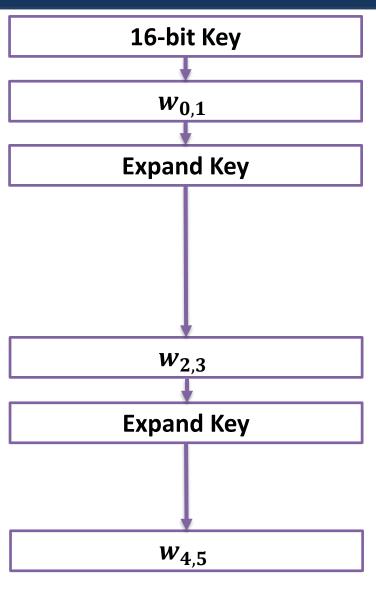
Simplified Advanced Encryption Standard

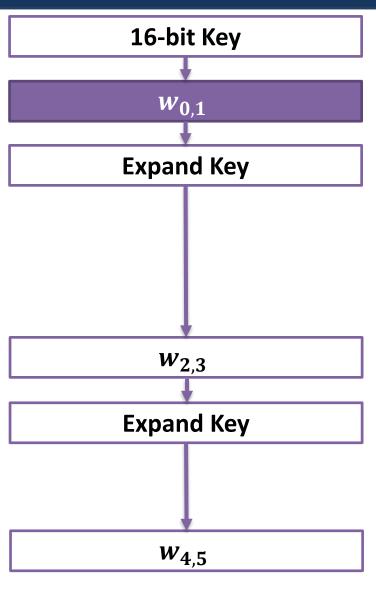
S-AES Encryption and Decryption

S-AES Key Generation

S-AES Encryption

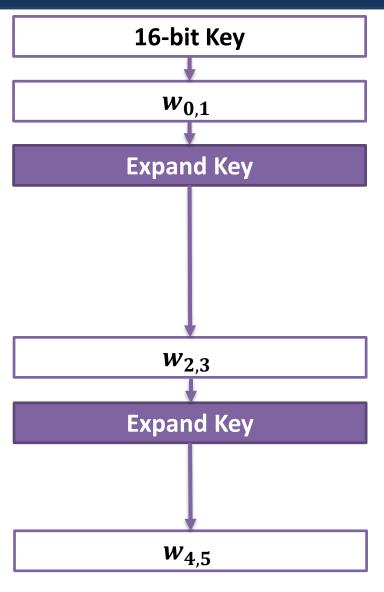
S-AES Decryption



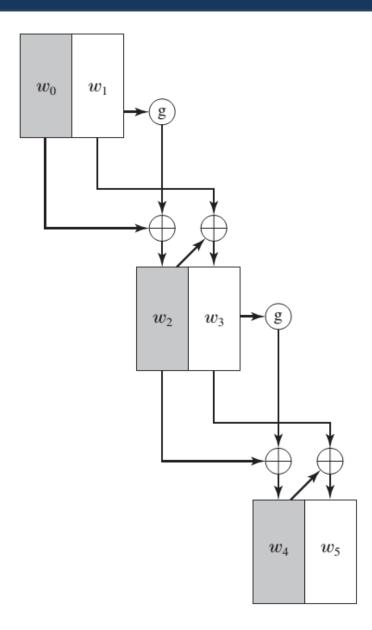


- □ K= 4A F5
 - = 0100 1010 1111 0101

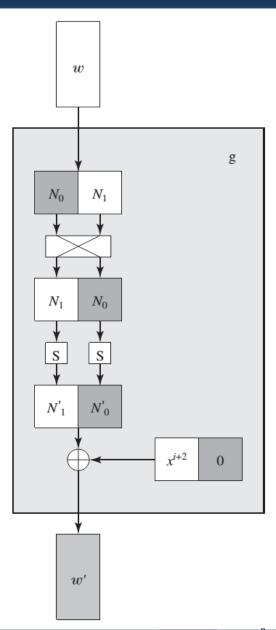
- ☐ The input key, K, is split into 2 words, w0 and w1:
- \square w0 = 0100 1010
- \square w1 = 1111 0101

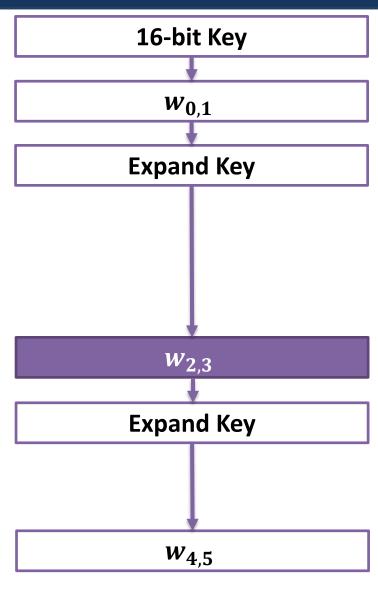


☐ S-AES Key Expansion



☐ Function g

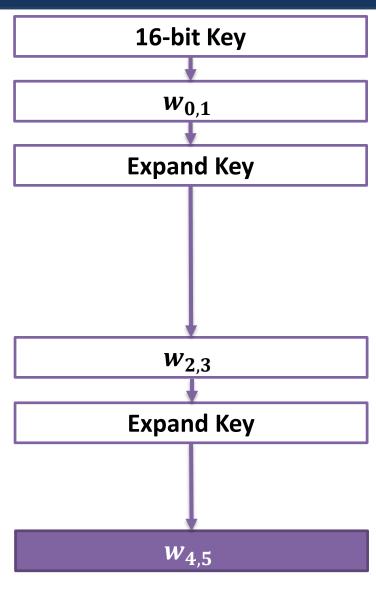




- \square w0 = 0100 1010, w1 = 1111 0101
- \square w2 = w0 \oplus Rcon(1) \oplus SubNib(RotNib(w1))
- □ RotNib() is "rotate the nibbles", which is equivalent to swapping the nibbles, Rcon is a round constant
- □ SubNib() is "apply S-Box substitution on nibbles using encryption S-Box"
- \square RotNib(w1) = 0101 1111
- □ SubNib(0101 11111) = 0001 0111
- \square Rcon(1) = 10000000

C D	20V	j						
S-Box		00	01	10	11			
i	00	9	4	A	В			
	01	D	1	8	5			
	10	6	2	0	3			
	11	С	Е	F	7			

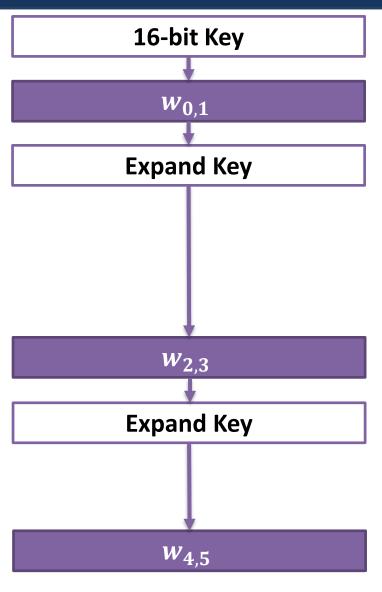
- \square w0 = 0100 1010, w1 = 1111 0101
- \square w2 = w0 \oplus Rcon(1) \oplus SubNib(RotNib(w1))
 - = 0100 1010 \(\oplus \) 1000 0000 \(\oplus \) 0001 0111
 - = 0100 1010 \(\oplus \) **1001 0111**=1101 1101
- \square w2 =1101 1101
- \square w3 = w2 \oplus w1 = 1101 1101 \oplus 1111 0101= 0010 1000
- \square w3 = 0010 1000



```
\square w2 = 1101 1101, w3 = 0010 1000
\square w4 = w2 \oplus Rcon(2) \oplus SubNib(RotNib(w3))
      = 1101 1101 \oplus 0011 0000 \oplus SubNib( 1000 0010 )
      = 1110 1101 \oplus 0011 0000 \oplus 0110 1010
      = 1110 1101 \oplus 0101 1010
      = 1011 \ 0111
```

$w5 = w4 \oplus w3$		
=1011 0111	⊕ 0010	1000
=1001 1111		

S-B	Юх	j							
		00	01	10	11				
	00	9	4	A	В				
i	01	D	1	8	5				
	10	6	2	0	3				
	11	С	Е	F	7				



- ☐ Key
- \square Key0 = w0w1
 - = 0100 1010 1111 0101
- \square Key1 = w2w3
 - = 1101 1101 0010 1000
- \Box Key2 = w4w5
 - = 1011 0111 1001 1111

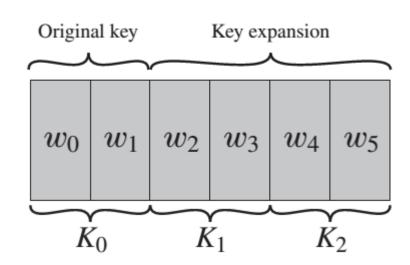


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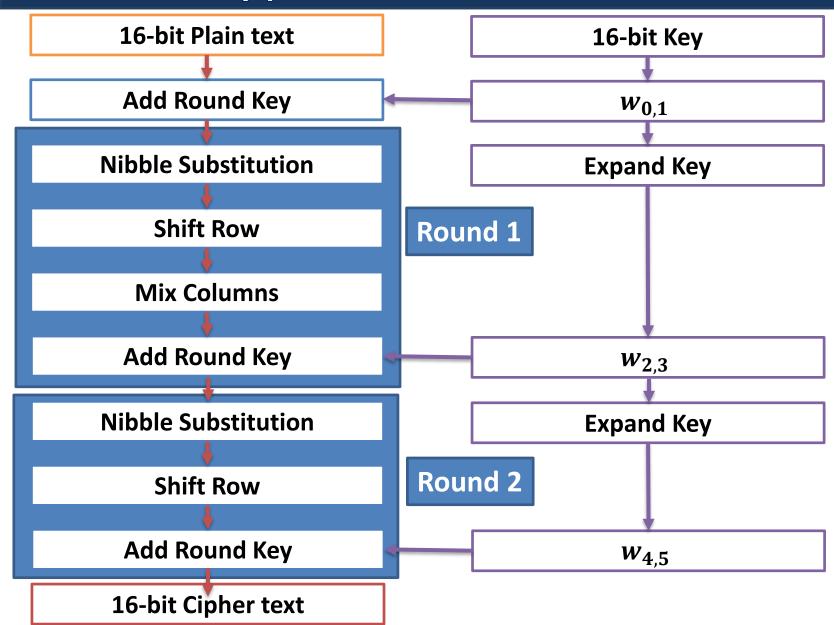
Simplified Advanced Encryption Standard

S-AES Encryption and Decryption

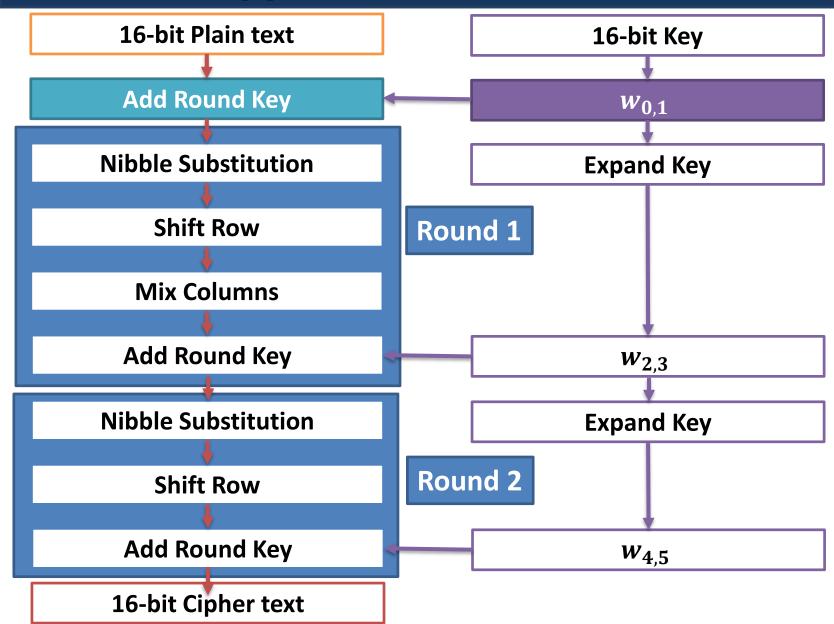
S-AES Key Generation

S-AES Encryption

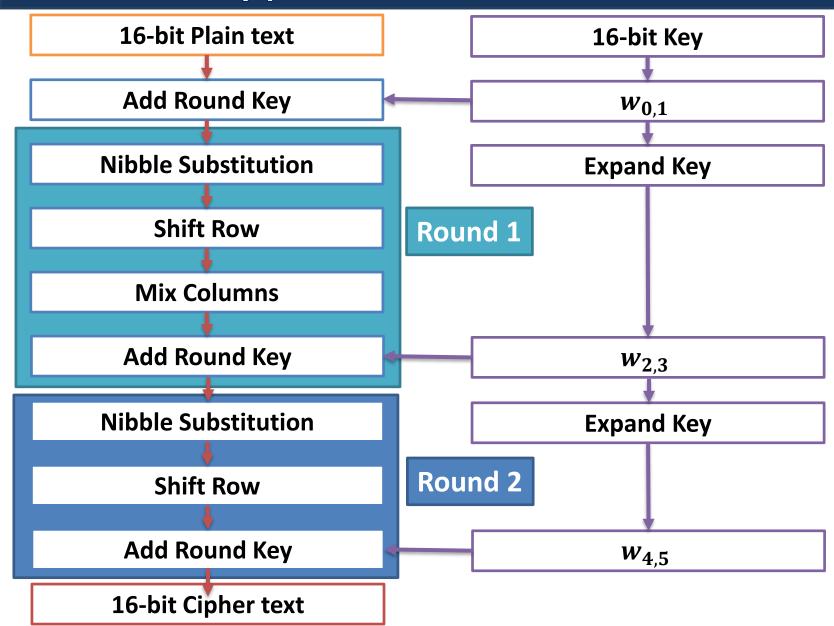
S-AES Decryption



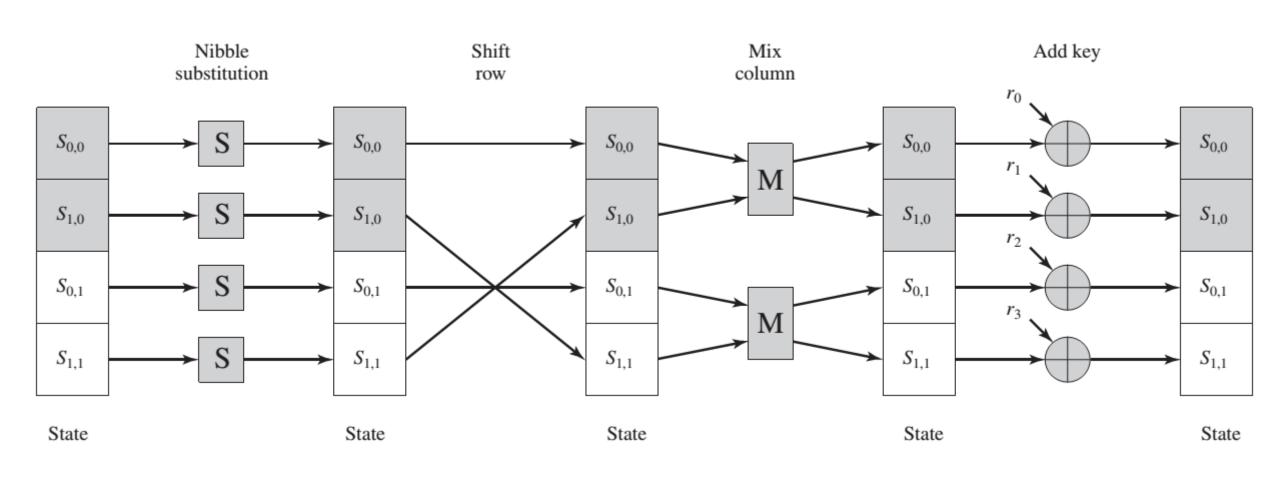
- ☐ Assume: P= 1101 0111 0010 1000
- \square Key0 = w0w1
 - = 0100 1010 1111 0101
- \square Key1 = w2w3
 - = 1101 1101 0010 1000
- \square Key2 = w4w5
 - = 1000 0111 1010 1111



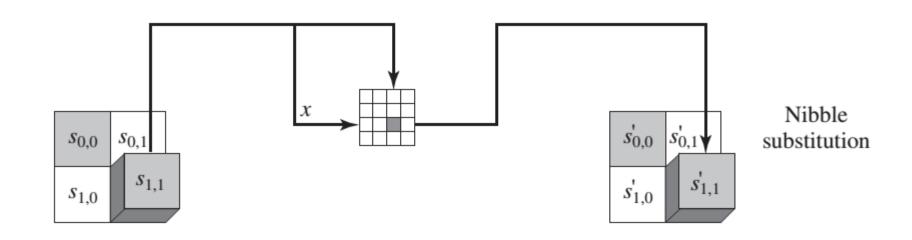
- □ Round 0
- □ P= 1101 0111 0010 1000
- ☐ Key0 = 0100 1010 1111 0101
- \square R0= P \oplus Key0
 - = 1101 0111 0010 1000 ⊕
 - 0100 1010 1111 0101
 - = 1001 1101 1101 1101

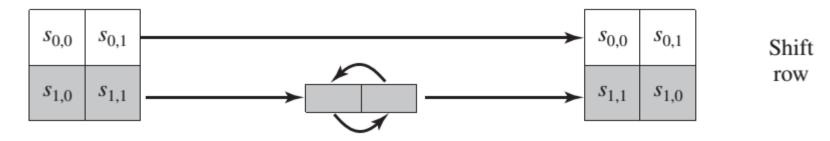


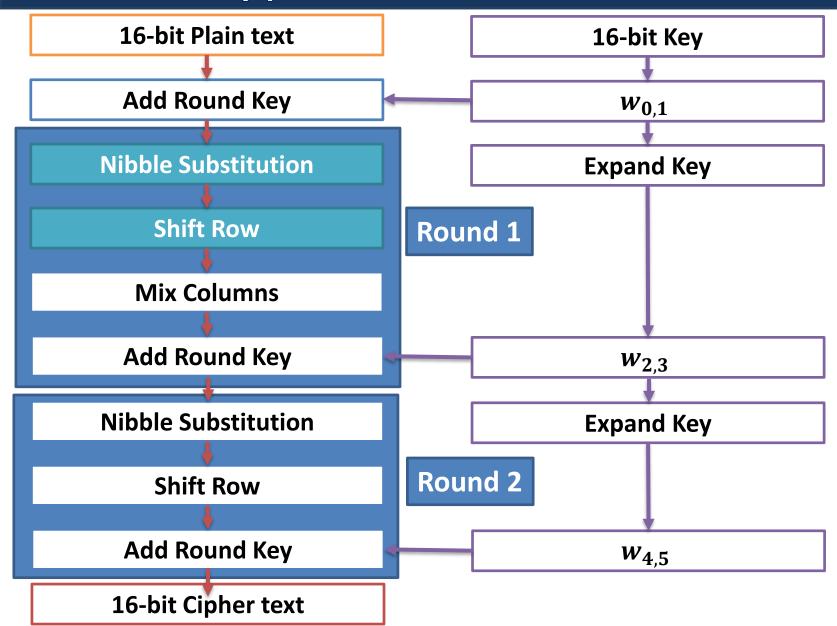
☐ S-AES Encryption Round



□ S-AES Transformation (Substitution and Shift row)







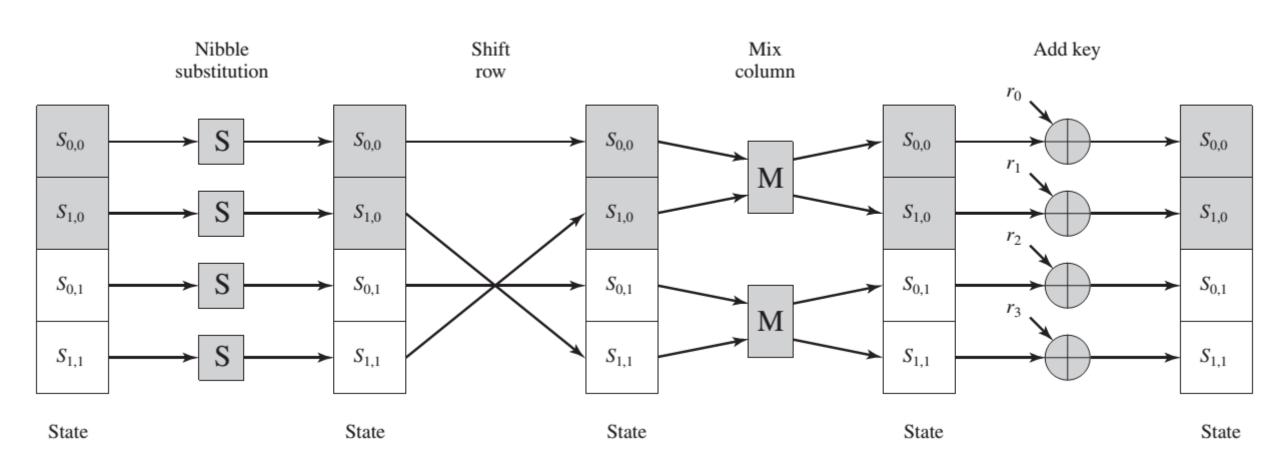
- ☐ Round 1
- 1) Nibble Substitution:
- 2) Shift Row:
- ☐ Swap 2nd nibble and 4th nibble
- ☐ ShRow(0010 1110 1110)

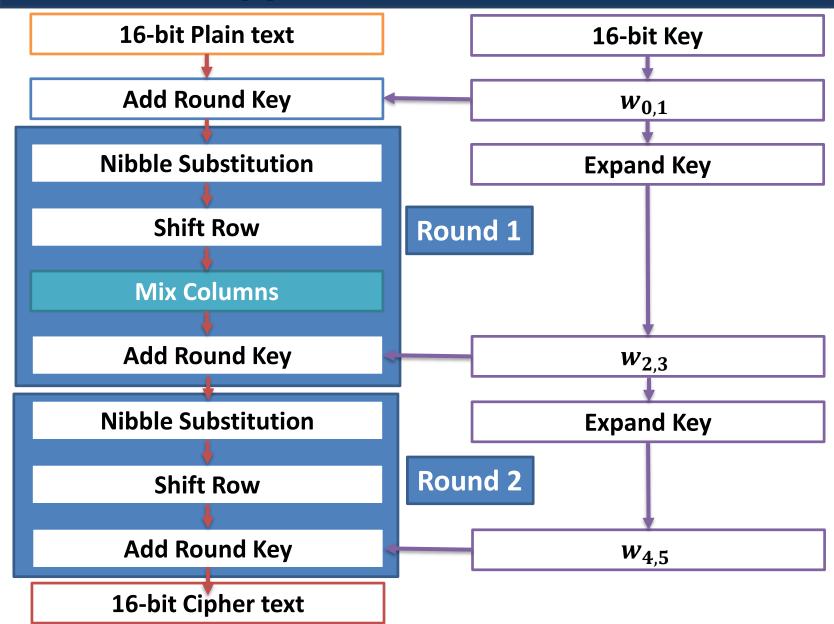
= 0010 1110 1110 1110

S-Box		j							
		00	01	10	11				
i	00	9	4	A	В				
	01	D	1	8	5				
	10	6	2	0	3				
	11	С	Е	F	7				

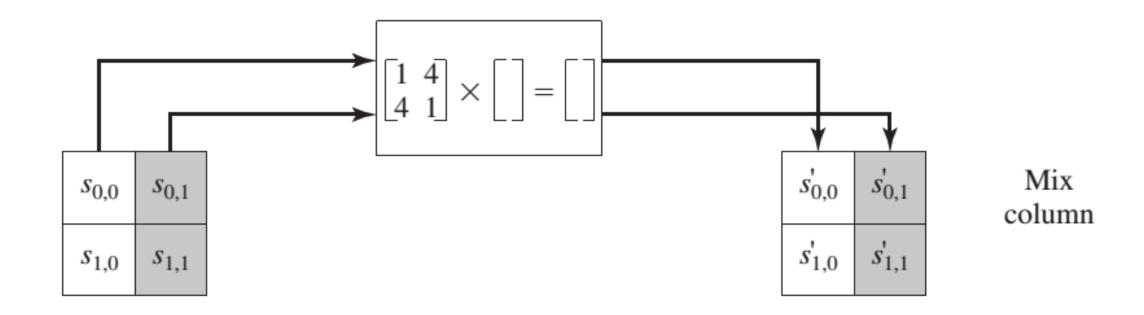
Key Generation

☐ S-AES Encryption Round





□ S-AES Transformation (Mix Column)



☐ Mix Column Table

*	1	2	3	4	5	6	7	8	9	Α	В	C	D	Ε	F
2	2	4	6	8	Α	С	E	3	1	7	5	В	9	F	D
4	4	8	С	3	7	В	F	6	2	E	Α	5	1	D	9
9	9	1	8	2	В	3	Α	4	D	5	С	6	F	7	Ε

- □ Round 1
- 3) Mix Columns:

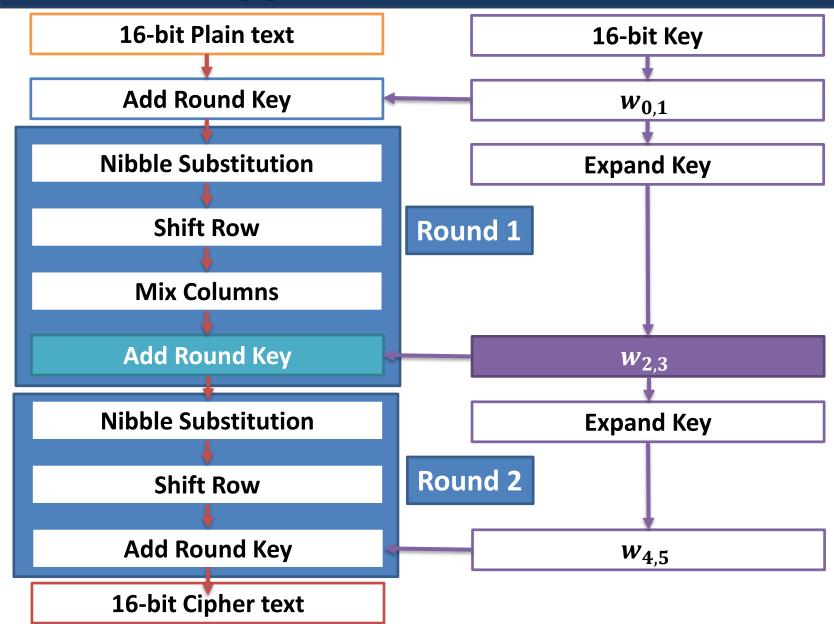
$$\square$$
 MixCol (0010 1110 1110 1110) = $\begin{pmatrix} 0010 & 1110 \\ 1110 & 1110 \end{pmatrix} * \begin{pmatrix} 1 & 4 \\ 4 & 1 \end{pmatrix} =$

$$\Box = \begin{pmatrix} 2 & E \\ E & E \end{pmatrix} * \begin{pmatrix} 1 & 4 \\ 4 & 1 \end{pmatrix} = \begin{pmatrix} (2*1 \oplus E*4) & (E*1 \oplus E*4) \\ (2*4 \oplus E*1) & (E*4 \oplus E*1) \end{pmatrix}$$

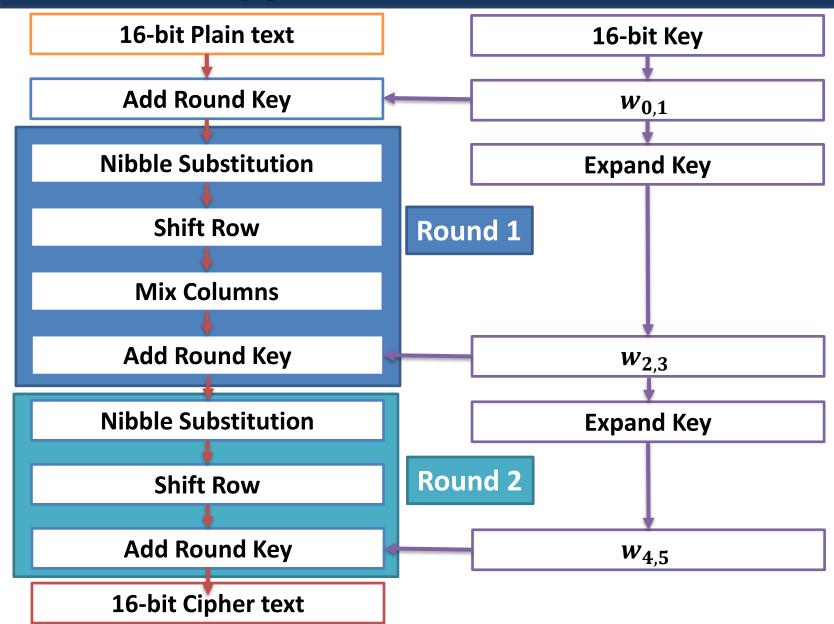
$$\Box = \begin{pmatrix} (2 \oplus D) & (E \oplus D) \\ (8 \oplus E) & (D \oplus E) \end{pmatrix} = \begin{pmatrix} (0010 \oplus 1101) & (1110 \oplus 1101) \\ (1000 \oplus 1110) & (1101 \oplus 1110) \end{pmatrix}$$

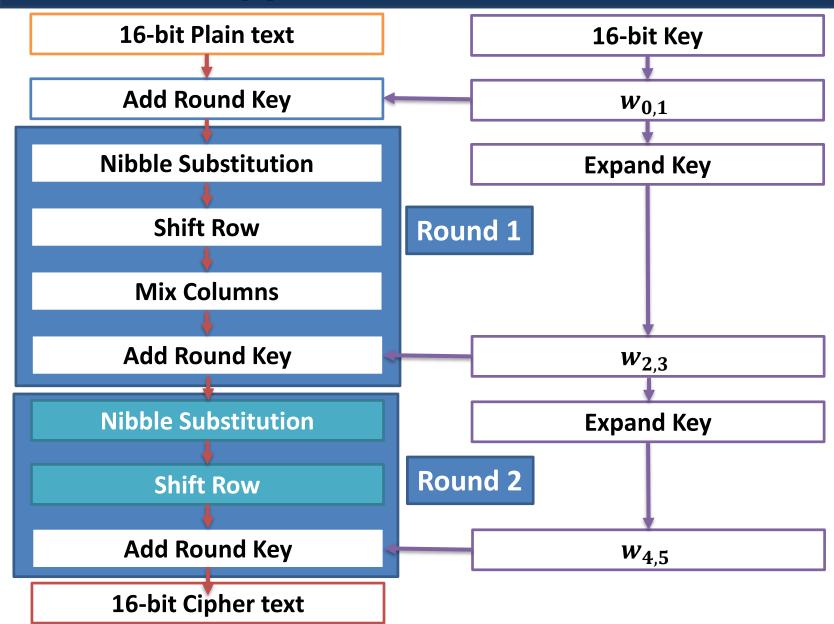
$$\square \begin{pmatrix} (0010 \oplus 1101) & (11110 \oplus 1101) \\ (1000 \oplus 1110) & (1101 \oplus 1110) \end{pmatrix} = \begin{pmatrix} 1111 & 0011 \\ 0110 & 0011 \end{pmatrix}$$

= 1111 0110 0011 0011

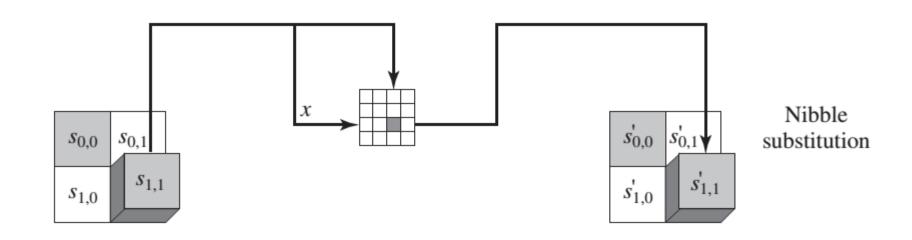


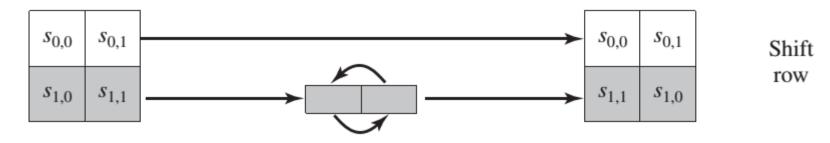
- ☐ Round 1
- 4) Add round Key1
- ☐ Key1 = 1101 1101 0010 1000
- \square R1= Key1 \oplus MixCol(ShRow(SubNib(R0)))
 - $= 1101 \ 1101 \ 0010 \ 1000 \oplus 1111 \ 0110 \ 0011 \ 0011$
 - = 0010 1011 0001 1011





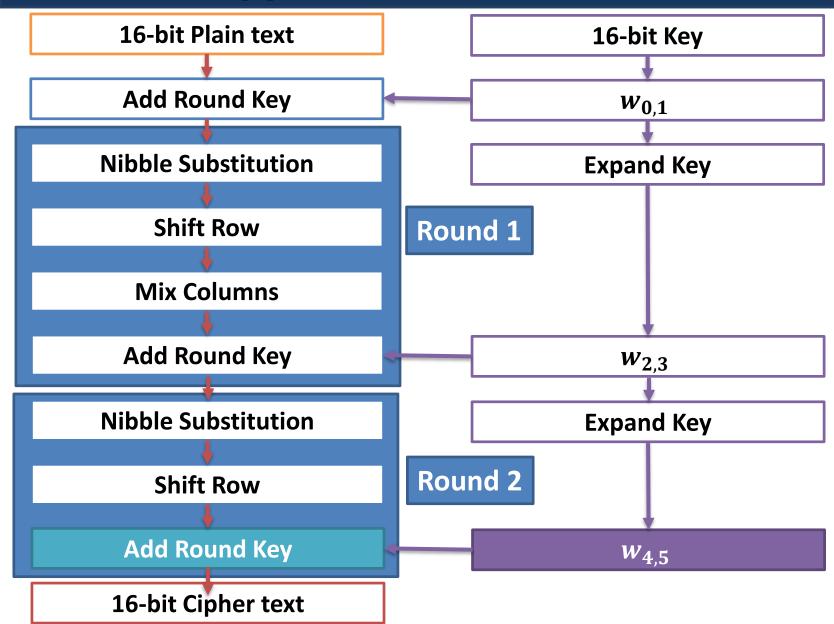
□ S-AES Transformation (Substitution and Shift row)





- □ Round 2
- 1) Nibble Substitution:
- □ SubNib(0010 1011 0001 1011)= 1010 0011 0100 0011
- 2) Shift Row:
- ☐ Swap 2nd nibble and 4th nibble
- ☐ ShRow(1010 0011 0100 0011)
 - = 1010 0011 0100 0011

S-Box		j			
		00	01	10	11
i	00	9	4	A	В
	01	D	1	8	5
	10	6	2	0	3
	11	С	Е	F	7



- □ Round 2
- 4) Add round Key2
- ☐ Key2 = 1000 0111 1010 1111
- \square R2= Key2 \oplus ShRow(SubNib(R1))
 - $= 1101 \ 1101 \ 0010 \ 1000 \oplus 1010 \ 0011 \ 0100 \ 0011$
 - = 0010 0100 1110 1100
 - Ciphertext = 0010 0100 1110 1100

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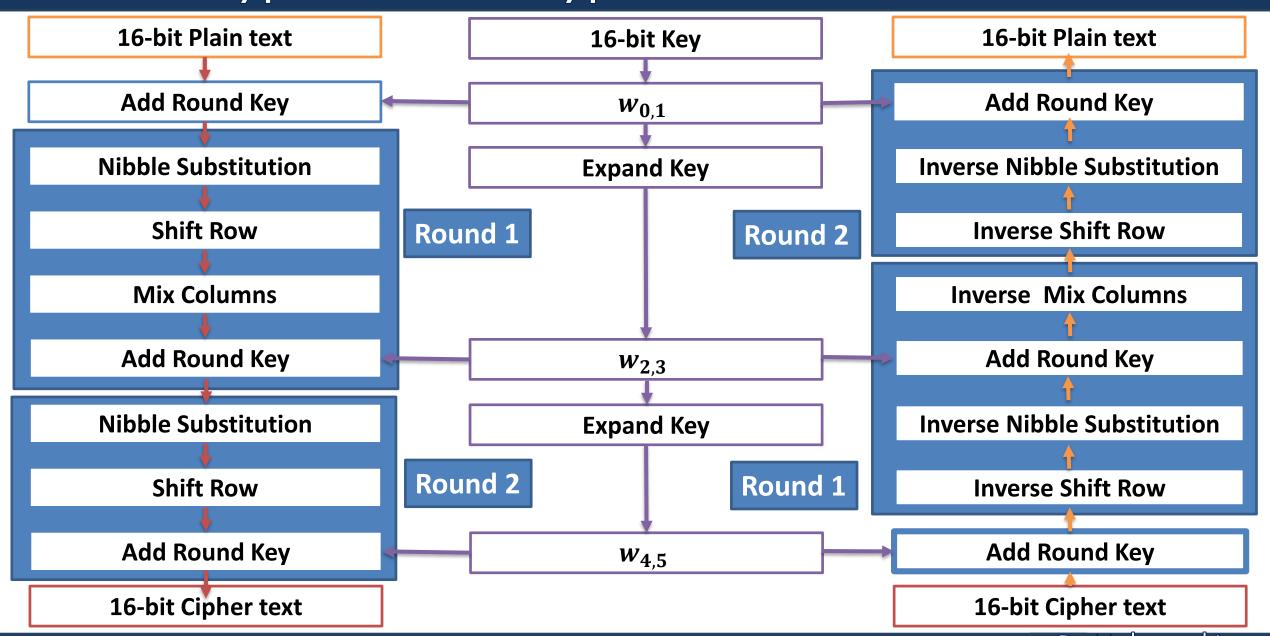
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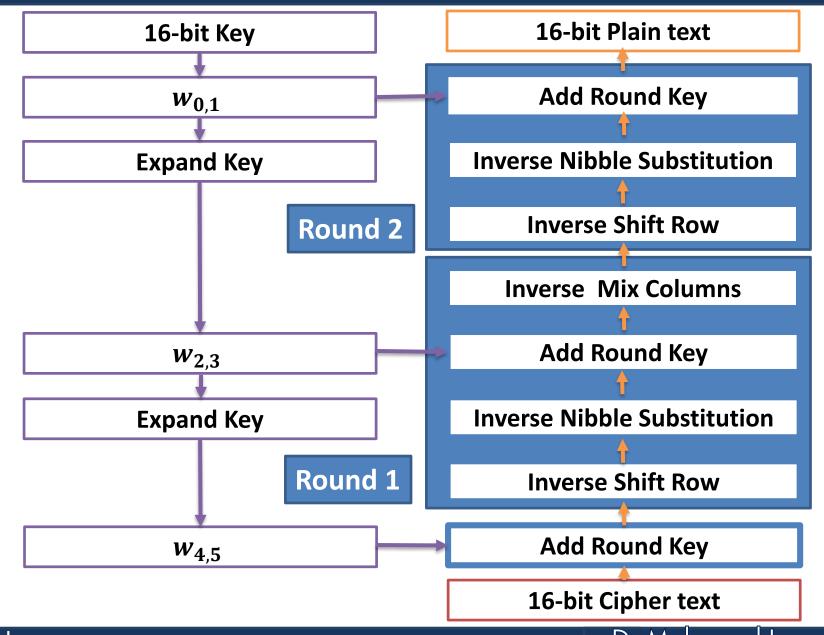
S-AES Key Generation

S-AES Encryption

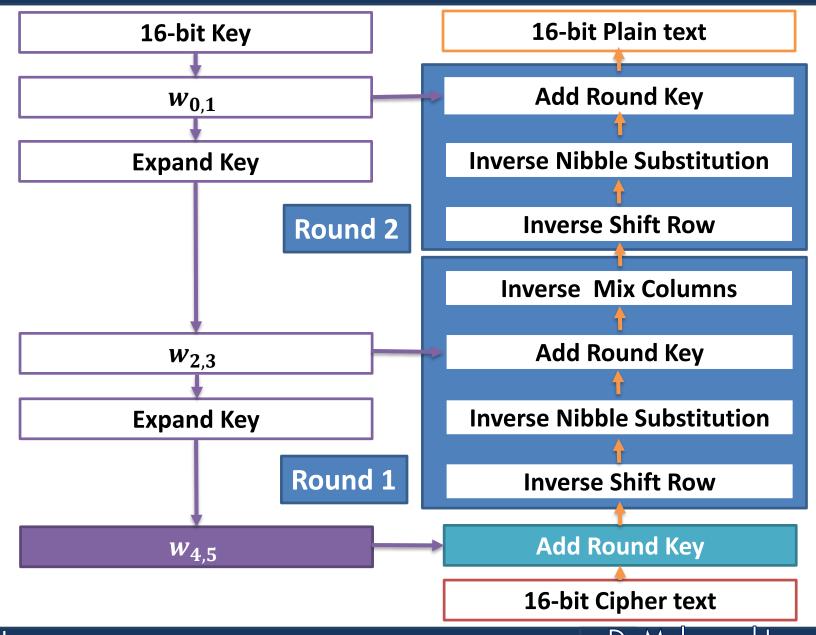
S-AES Encryption and Decryption



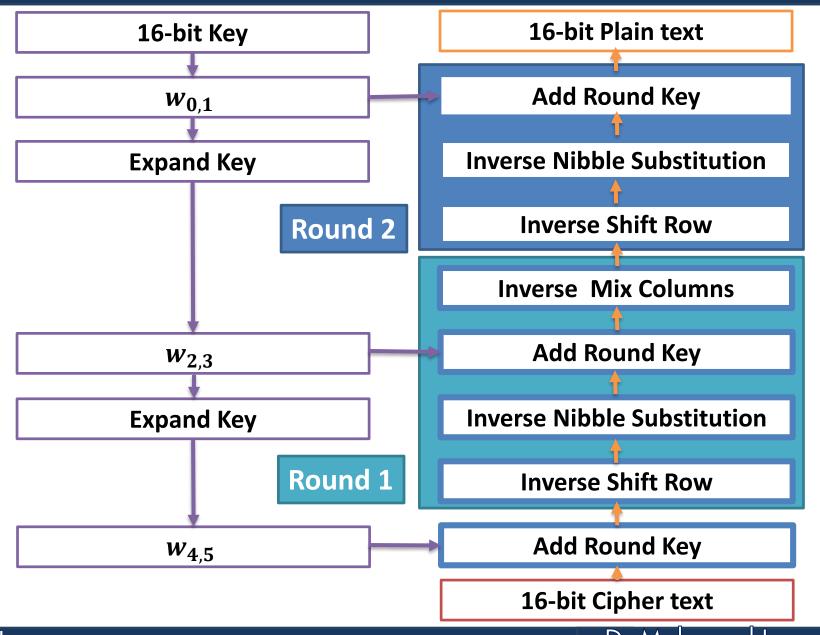
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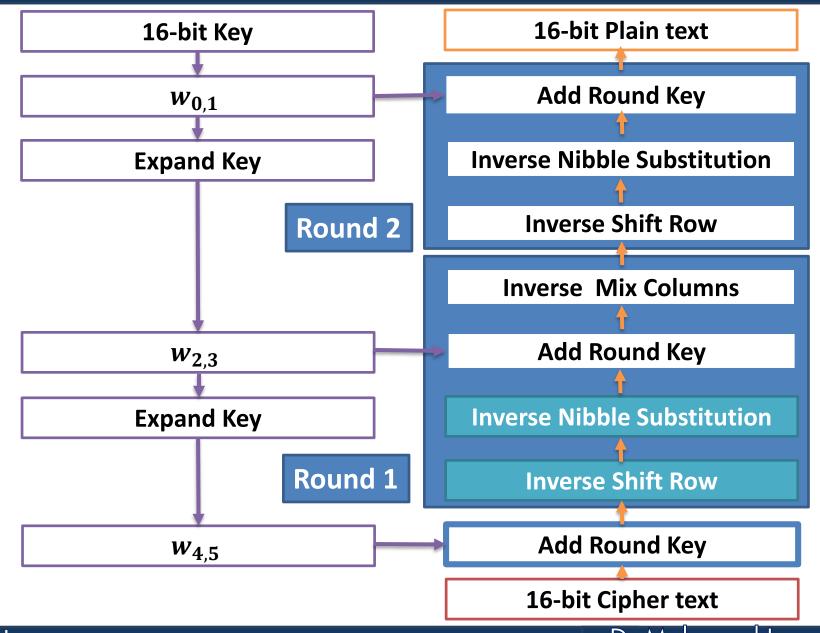


- ☐ Assume: C= 0010 0100 1110 1100
- \square Key0 = w0w1
 - = 0100 1010 1111 0101
- \square Key1 = w2w3
 - = 1101 1101 0010 1000
- \square Key2 = w4w5
 - = 1000 0111 1010 1111



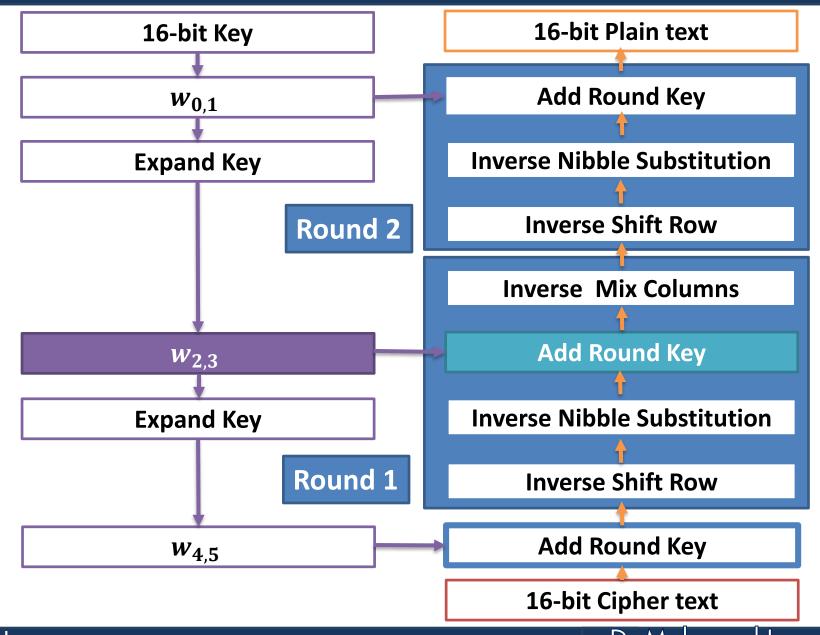
- \Box Key2 = w4w5
 - = 1000 0111 1010 1111
- □ C= 0010 0100 1110 1100
- □ R0= C ⊕ Key2 = 0010 0100 1110 1100 ⊕ 1000 0111 1010
 - 1111
- $\square = 1010\ 0011\ 0100\ 0011$



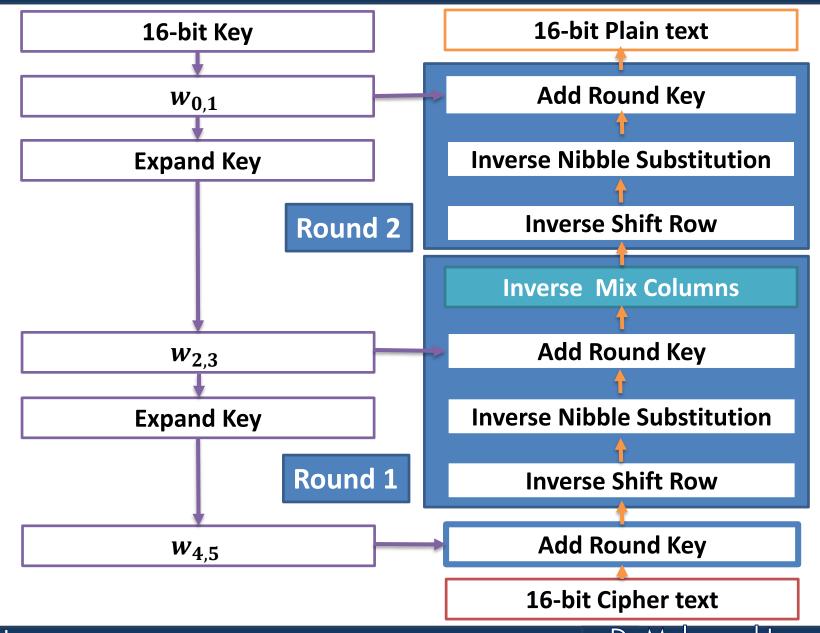


- ☐ Round 1
- 1) Inverse Shift Row
- \square IShRow(RO)=IShRow(1010 0011 0100 0011)=
 - = 1010 0011 0100 0011
- 2) Inverse Nibble Sub
- □ ISubNib(1010 0011 0100 0011)=
 - = 0010 1011 0001 1011

Inv S-Box		j				
		00	01	10	11	
i	00	A	5	9	В	
	01	1	7	8	F	
	10	6	0	2	3	
	11	С	4	D	Е	



- 3) Add Round 1 Key
- □ 0010 1011 0001 1011 ⊕ Key1
 - $= 0010 \ 1011 \ 0001 \ 1011 \oplus 1101 \ 1101 \ 0010 \ 1000$
 - =1111 0110 0011 0011



- □ Round 1
- 4) Inverse Mix Columns:
- \square MixCol (11111 0110 0011 0011) = $\begin{pmatrix} 11111 & 0011 \\ 0110 & 0011 \end{pmatrix} * \begin{pmatrix} 9 & 2 \\ 2 & 9 \end{pmatrix} =$

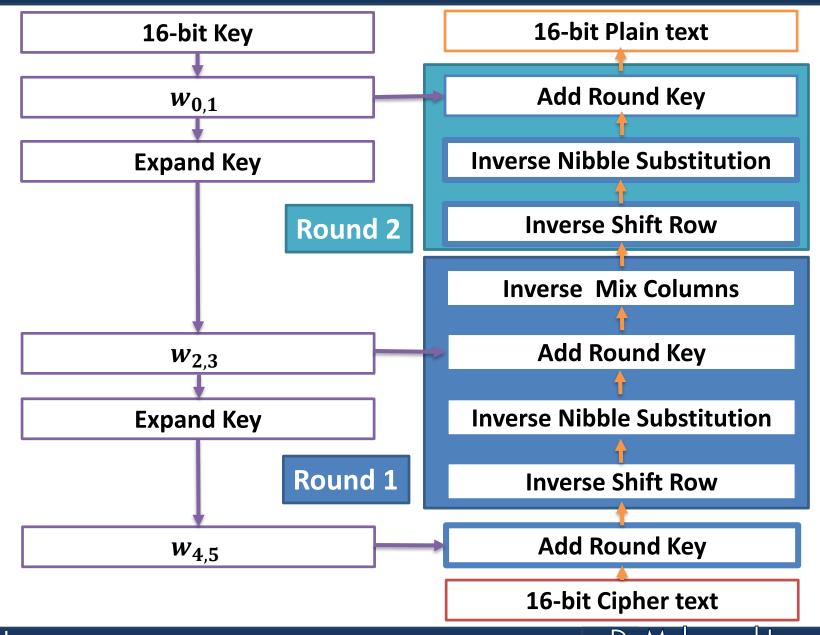
$$\Box = \begin{pmatrix} F & 3 \\ 6 & 3 \end{pmatrix} * \begin{pmatrix} 9 & 2 \\ 2 & 9 \end{pmatrix} = \begin{pmatrix} (F*9 \oplus 6*2) & (3*9 \oplus 3*2) \\ (F*2 \oplus 6*9) & (3*2 \oplus 3*9) \end{pmatrix}$$

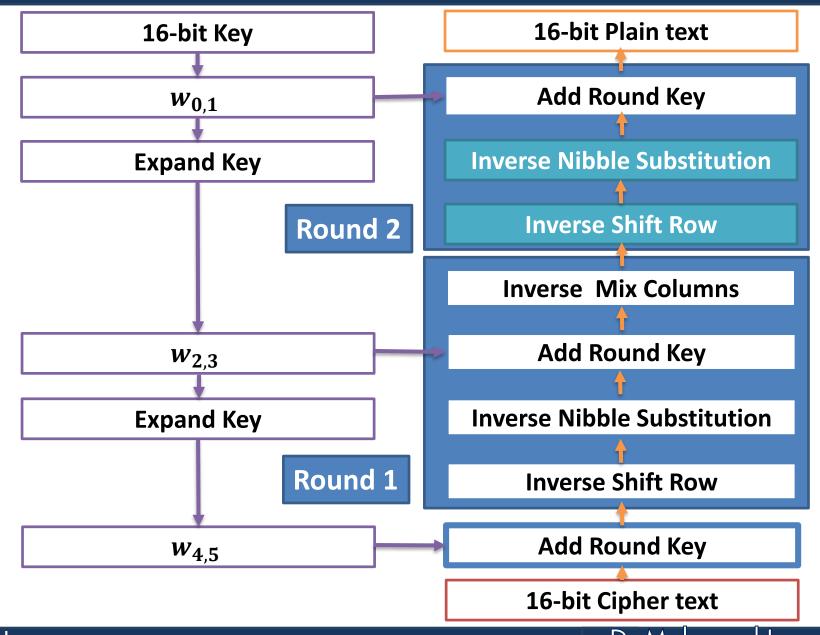
$$\Box = \begin{pmatrix} (E \oplus C) & (8 \oplus 6) \\ (D \oplus 3) & (6 \oplus 8) \end{pmatrix}$$

$$\Box = \begin{pmatrix} (1110 \oplus 1100) & (1000 \oplus 0110) \\ (1101 \oplus 0011) & (0110 \oplus 1000) \end{pmatrix}$$

$$\Box \begin{pmatrix} (1110 \oplus 1100) & (1000 \oplus 0110) \\ (1101 \oplus 0011) & (0110 \oplus 1000) \end{pmatrix} = \begin{pmatrix} 0010 & 1110 \\ 1110 & 1110 \end{pmatrix}$$

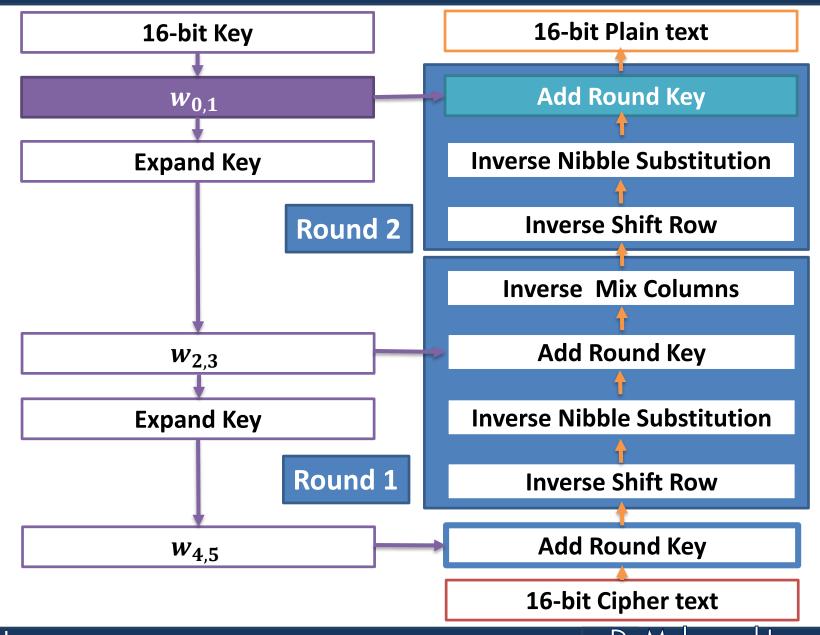
$$R1 = 0010 1110 1110 1110$$





- □ Round 2
- 1) Inverse Shift Row
- ☐ IShRow(R1)=IShRow(0010 1110 1110)=
 - = 0010 1110 1110 1110
- 2) Inverse Nibble Sub
- □ ISubNib(0010 1110 1110 1=
 - = 1001 1101 1101 1101

		j			
Inv S-Box		00	01	10	11
i	00	A	5	9	В
	01	1	7	8	F
	10	6	0	2	3
	11	С	4	D	Е



- 3) Add Round 2 Key
- □ R2= 1001 1101 1101 1101 ⊕ Key0
 - $= 1001 \ 1101 \ 1101 \ 1101 \oplus 0100 \ 1010 \ 1111 \ 0101$
 - =1101 0111 0010 1000

Plaintext = 1101 0111 0010 1000 = D7 28

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