

MIT WORLD PEACE UNIVERSITY

Computer Networks  
Second Year B.Tech Semester 3  
Academic Year 2022-23

---

---

OPERATING SYSTEMS

---

---

NOTES FROM TANANBAUM AND CLASSES

Prepared By

P34. Krishnaraj Thadesar

Batch A2

September 6, 2022

## **Contents**

<b>1</b>	<b>Concurrency Control - Waht to learn</b>	<b>2</b>
<b>2</b>	<b>Design Issues in concurrency</b>	<b>2</b>
<b>3</b>	<b>Contexts of Concurrency</b>	<b>2</b>
<b>4</b>	<b>Key terms related to concurrency</b>	<b>2</b>
4.1	Race Condition . . . . .	3
<b>5</b>	<b>Difficulties due to concurrency</b>	<b>3</b>
5.1	Example . . . . .	3

## 1 Concurrency Control - What to learn

**P**rocess Synchronization: Principles of Concurrency, Requirements for Mutual Exclusion: Hardware Support, OS Support

**C**lassical Synchronization Problems: Readers Writers Problem, Producer and Consumer Problem

**P**riniples of Deadlock, Deadlock modelling, prevention, avoidance and stuff.

**When 2 processes are running at the same time, when they are interdependent on each other through inputs and outputs, then you would call that concurrency. We aren't doing multiprocessing by choice, here we just gotta do it at the same time somehow.**

## 2 Design Issues in concurrency

1. Communication among processes
2. Sharing and Competition for resources
3. Synchronization of activities of multiple processes
4. Allocation of processor time to processes.

## 3 Contexts of Concurrency

1. Multiple Applications like Multiprogramming(diff programs on a single processor) and Multiprocessing(on diff processors)
2. Structured Applications: Some applications can be effectively programmed as a set of concurrent processes. (Principles of modular design and structured programming)
3. OS Structure: OS often implemented as a set of processes or threads.

## 4 Key terms related to concurrency

1. Atomic Operation: A sequence of one or more statements that appear to be indivisible that is no other process can see an intermediate state or interrupt the operations
2. Critical Section: A section of code within a process that requires access to shared resources and that must not be executed while another process is in corresponding section of code.
3. Deadlock: A situation in which two or more processes are unable to proceed because each is waiting for one of the others to do something.
4. Mutual Exclusion: The requirement that when one process is in critical section that access shared resources, no other processes may be in critical section that accesses any of those shared resources.
5. Race Condition: A situation in which multiple threads/Processes read and write a shared data item and final result depends on relative timing of their execution.
6. Starvation: A situation in which a runnable process is overlooked indefinitely by the scheduler, although it is able to proceed, it is never chosen.

### 4.1 Race Condition

A race Condition occurs when multiple competing processes or threads read and write data items so that final result depends on the order of execution of instructions in multiple processes. So 2 processes p1 and p2, say they share global variable 'a'. P1 updates a to 1, and p updates it to 2. Thus two tasks are in a race to write variable 'a'. Loser of race is the one that determines the value of 'a'.

## 5 Difficulties due to concurrency

1. Sharing of global resources: Eg. Two processes both make use of global variable and both perform read and write on that variable, in which read and write are done is critical.
2. Management of resources optimally. Eg. Process has gained the ownership of IO devices but is suspended before using it, thus locking IO device and preventing its use by other processes.
3. Error locating in Program: Results are not deterministic are reproducible.

### 5.1 Example

```
void echo()
{
    chin = getchar();
    chout = chin;
    putchar(chout);
}
```

1. Uniprocessor multiprogramming, single user environment
2. Many applications can call this procedure repeatedly to accept user input and display on screen.
3. User can jump from one application to other.
4. Each application needs or uses procedure echo.
5. Echo is made shared procedure and loaded into a portion of memory global to all applications
6. Single copy of echo procedure is used, saving space.
7. When echo procedure is invoked by some process, and if the process gets suspended for any reason before completing it, then no other process can invoke echo till process that was suspended gets resumed, and completes echo. Thus other processes are not allowed to access it.
8. It would make more sense if you visualize this in terms of a multi user operating system. So if many people are using a server kind of thing at the same time.