



Dr. Vishwanath Karad

**MIT WORLD PEACE  
UNIVERSITY** | PUNE

TECHNOLOGY, RESEARCH, SOCIAL INNOVATION & PARTNERSHIPS

## Scheduling

---

# Syllabus

---

**Scheduling:** Types of Scheduling, Scheduling Algorithms: FCFS, SJF, Priority, Round Robin.

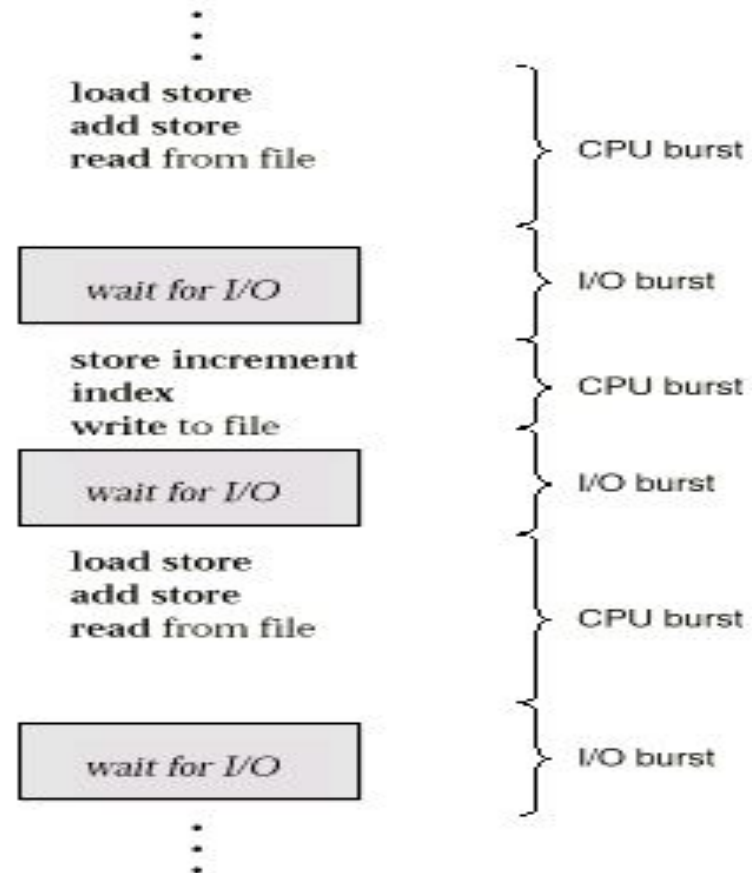
# Scheduling

---

- Maximum CPU utilization obtained with multiprogramming
- CPU–I/O Burst Cycle – Process execution consists of a *cycle* of CPU execution and I/O wait.

# Alternating Sequence of CPU And I/O Bursts

---



# CPU / Process Scheduler

---

- Selects from among the processes in memory that are ready to execute, and allocates the CPU to one of them.
- Scheduling may be *preemptive* or *nonpreemptive*
  - Nonpreemptive*** – If CPU allocated to a process, it keeps the CPU until it releases it by terminating or switching to waiting state
  - Preemptive*** - If CPU allocated to a process, it may be released if high priority process needs the CPU

# Scheduler

---

- CPU scheduling decisions may take place when a process:
  1. Switches from running to waiting state
  2. Switches from running to ready state
  3. Switches from waiting to ready state
  4. Terminates
- Scheduling under 1 and 4 is *nonpreemptive*
- All other scheduling is *preemptive*

# Scheduling Algorithms

---

They deal with the problem of deciding which of the processes in ready queue is to be allocated the CPU

Algo. compared based on following criteria

**CPU utilization** – keep the CPU as busy as possible

**Throughput** – # of processes completed or amount of work done per unit time

**Turnaround time** – time of submission of a process to the time of completion

**Waiting time** – amount of time a process has been waiting in the ready queue

**Response time** – amount of time it takes from when a request was submitted until the first response is produced

# Optimization Criteria

---

Max CPU utilization

Max throughput

Min turnaround time

Min waiting time

Min response time



# Types of process schedulers / Scheduling categories

---

- Scheduling is broken down into three categories:

## 1. Long term scheduling:

- Is performed when a new process is created.
- Long-term scheduler (or job scheduler) – selects which processes should be brought into the ready queue.
- Long-term scheduler is invoked very infrequently (seconds, minutes)  $\Rightarrow$  (may be slow)
- The long-term scheduler controls the *degree of multiprogramming*.

# Types of process schedulers / Scheduling categories

---

## 2. Medium term scheduling:

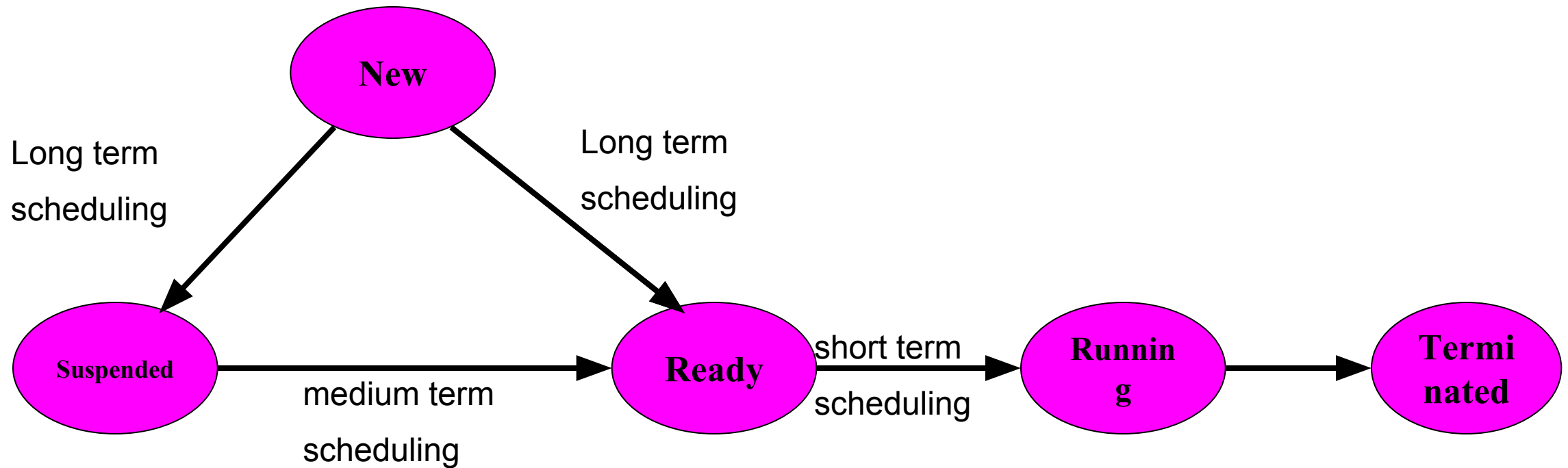
- Part of the swapping function
- Swapping-in decisions are taken by medium term scheduler
- Based on the need to manage the degree of multiprogramming

## 3.Short term scheduling:

- Determines which ready process will be assigned the CPU when it next becomes available and actually assign the CPU to this process .
- Short-term scheduler (or CPU scheduler) – selects which process should be executed next and allocates CPU.
- Short-term scheduler is invoked very frequently (milliseconds)⇒ (must be fast).

# Types of process schedulers / Scheduling categories

---



# Scheduling Algorithms

---

- FCFS (First Come First Serve)
- SJF (Shortest Job First)
- Priority scheduling
- Round Robin scheduling

# FCFS Scheduling: Characteristics

---

**Selection Function:**  $\max(w)$ , selects the process which is waiting in the ready queue for maximum time.

**Decision Mode :** Non\_preemptive

**Throughput:** Not emphasized

**Response Time:** May be high, especially if there is a large variance in the process execution times.

**Overhead:** Minimum

**Effect on Processes:** Penalizes short processes

**Starvation:** No

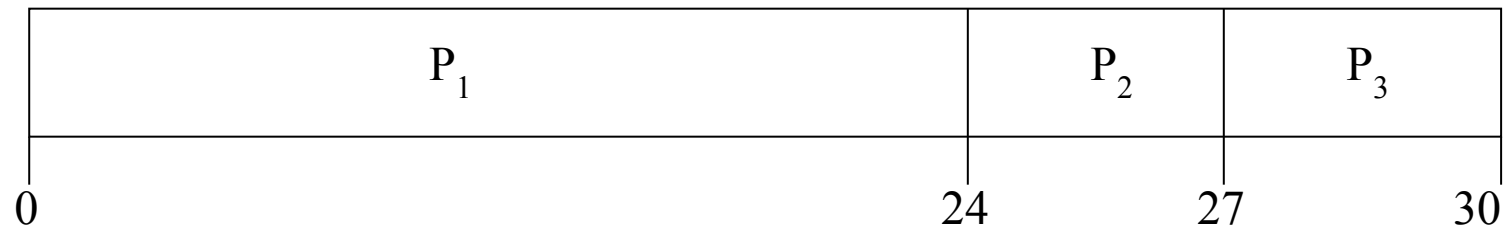
---

		<u>Process</u>	<u>Burst Time</u>
$P1$	24		
		$P2$	3
		$P3$	3

- Suppose that the processes arrive in the order:  $P1$  ,  $P2$  ,  $P3$

---

The Gantt Chart for the schedule is:



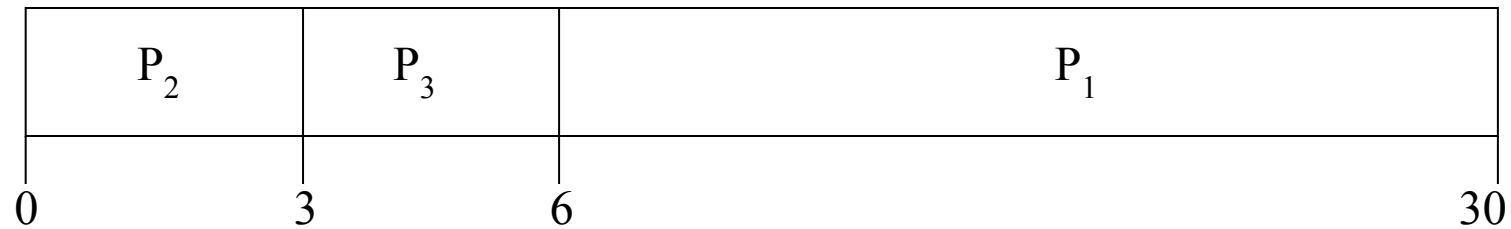
Waiting time for  $P_1 = 0$ ;  $P_2 = 24$ ;  $P_3 = 27$

Average waiting time:  $(0 + 24 + 27)/3 = 17$

## FCFS Scheduling (Cont.)

---

- Suppose that the processes arrive in the order :  $P_2, P_3, P_1$



- Waiting time for  $P_1 = 6$ ;  $P_2 = 0$ ;  $P_3 = 3$
- Average waiting time:  $(6 + 0 + 3)/3 = 3$



## Example-2 of FCFS

---

<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
----------------	---------------------	-------------------

<i>P1</i>	0.0	7
-----------	-----	---

<i>P2</i>	2.0	4
-----------	-----	---

<i>P3</i>	4.0	1
-----------	-----	---

<i>P4</i>	5.0	4
-----------	-----	---

## Example-2 of FCFS

---

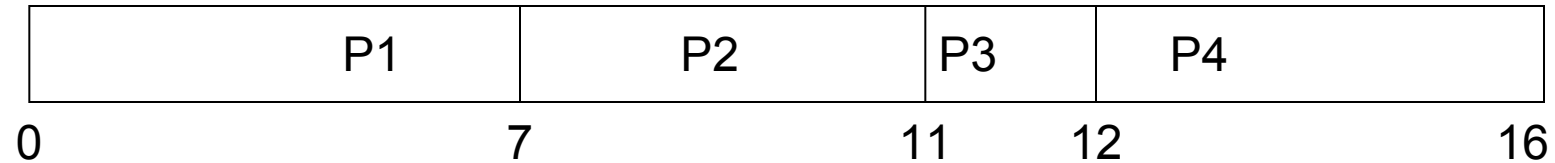
**Process   Arrival Time   Burst Time**

***P1*   0.0   7**

***P2*   2.0   4**

***P3*   4.0   1**

***P4*   5.0   4**



**Average waiting time =  $(0 + 5 + 7 + 7)/4 = 4.75$**

**Average Turnaround Time =  $(7+9+8+11)/4=8.75$**

# Shortest-Job-First (SJF) Scheduling

---

- Associate with each process the length of its next CPU burst. Use these lengths to schedule the process with the shortest time
- Two schemes:
  - **nonpreemptive** – once CPU given to the process it cannot be preempted until completes its CPU burst
  - **preemptive** – if a new process arrives with CPU burst length less than remaining time of current executing process, preempt. This scheme is known as the **Shortest-Remaining-Time-First (SRTF)**
- **SJF is optimal** – gives minimum average waiting time for a given set of processes

## Example of Non-Preemptive SJF

---

<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
<i>P1</i>	0.0	7
<i>P2</i>	2.0	4
<i>P3</i>	4.0	1
<i>P4</i>	5.0	4

## Example of Non-Preemptive SJF

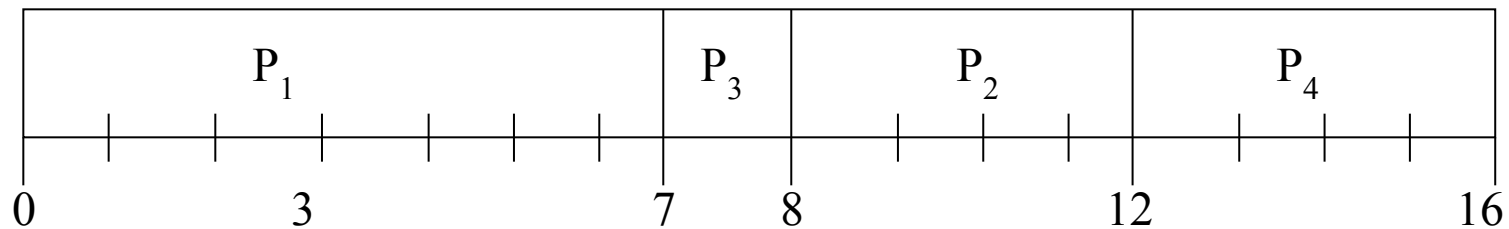
<u>Process</u>	<u>Arrival Time</u>	<u>Burst Time</u>
----------------	---------------------	-------------------

<i>P1</i>	0.0	7
-----------	-----	---

<i>P2</i>	2.0	4
-----------	-----	---

<i>P3</i>	4.0	1
-----------	-----	---

<i>P4</i>	5.0	4
-----------	-----	---



■ Average waiting time =  $(0 + 6 + 3 + 7)/4 = 4$

# Shortest Job First Preemptive or Shortest Remaining Time

---

- It is a preemptive version of SJF. In this policy, scheduler always chooses the process that has the shortest expected remaining processing time.
- When a new process arrives in the ready queue, it may in fact have a shorter remaining time than the currently running process.
- Accordingly, the scheduler may preempt whenever a new process becomes ready.
- Scheduler must have an estimate of processing time to perform the selection function.

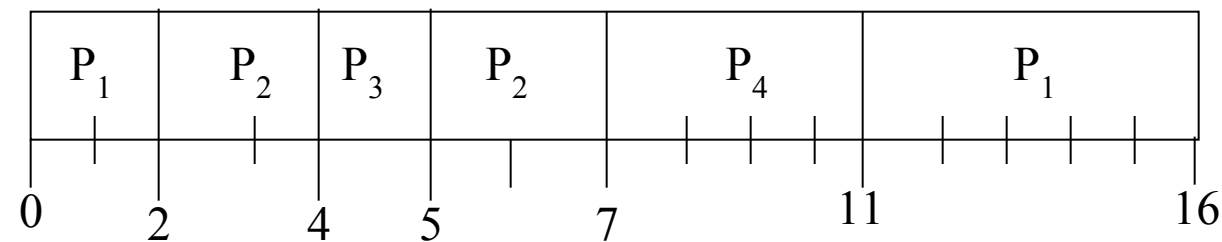
## Shortest Job First Preemptive or Shortest Remaining Time: characteristics

---

- **Selection Function:** minimum total service time required by the process, minus time spent in execution so far.
- **Decision Mode :** Preemptive ( At arrival time)
- **Throughput:** High
- **Response Time:** Provides good response time
- **Overhead:** Can be high
- **Effect on Processes:** Penalizes long processes.
- **Starvation:** Possible

# Example of Preemptive SJF

Process	Arrival Time	Burst Time
$P_1$	0.0	7
$P_2$	2.0	4
$P_3$	4.0	1
$P_4$	5.0	4

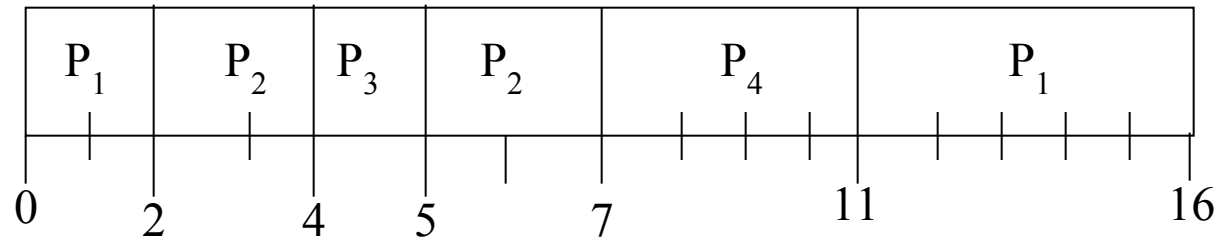




# Example of Preemptive SJF

---

Process	Arrival Time	Burst Time
$P_1$	0.0	7
$P_2$	2.0	4
$P_3$	4.0	1
$P_4$	5.0	4



- Average waiting time =  $(9 + 1 + 0 + 2)/4 = 3$
- Average turnaround time =  $(16 + 5 + 1 + 6)/4 = 7$

# Priority Scheduling

---

- A priority number (integer) is associated with each process
- The CPU is allocated to the process with the highest priority (smallest integer  $\equiv$  highest priority)
  - Preemptive
  - Nonpreemptive
- SJF is a priority scheduling where priority is the predicted next CPU burst time
- Problem  $\equiv$  Starvation – low priority processes may never execute
- Solution  $\equiv$  Aging – as time progresses increase the priority of the process

# Round Robin (RR)

---

- Each process gets a small unit of CPU time (*time quantum*)
- After this time has elapsed, the process is preempted and added to the end of the ready queue.
- If there are  $n$  processes in the ready queue and the time quantum is  $q$ , then each process gets  $1/n$  of the CPU time in chunks of at most  $q$  time units at once.
- No process waits more than  $(n-1)q$  time units.

## RR: characteristics

---

- **Selection Function:** constant
- **Decision Mode :** Preemptive ( At time quantum)
- **Throughput:** May be low if time quantum is too small
- **Response Time:** Provides good response time for short processes
- **Overhead:** Minimum
- **Effect on Processes:** Fair treatment
- **Starvation:** No

## Example of RR with Time Quantum = 20

---

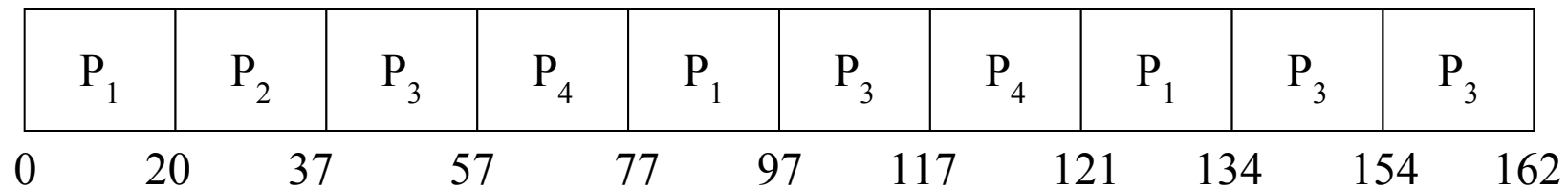
Process   Burst Time

*P1*   53

*P2*   17

*P3*   68

*P4*   24



- Typically, higher average turnaround than SJF, but better *response*

Example: Round Robin (By Default preemptive: Time Quantum 2)

---

Process	Arrival Time	Burst Time
$P_1$	0.0	7
$P_2$	2.0	4
$P_3$	4.0	1
$P_4$	5.0	4

Example: Round Robin (By Default preemptive: Time Quantum 2)

Process	Arrival Time	Burst Time
$P_1$	0.0	7
$P_2$	2.0	4
$P_3$	4.0	1
$P_4$	5.0	4

P1	P2	P1	P3	P2	P4	P1	P4	P1	
0	2	4	6	7	9	11	13	15	16

Average waiting time =  $(9 + 3 + 2 + 6)/4 = 5$

Average turnaround time =  $(16 + 7 + 3 + 10)/4 = 9$

# References

---

1. William Stallings, Operating System: Internals and Design Principles, Prentice Hall, ISBN-10: 0-13-380591-3, ISBN-13: 978-0-13-380591-8, 8th Edition
2. Abraham Silberschatz, Peter Baer Galvin and Greg Gagne, Operating System Concepts, WILEY, ISBN 978-1-118-06333-0, 9th Edition