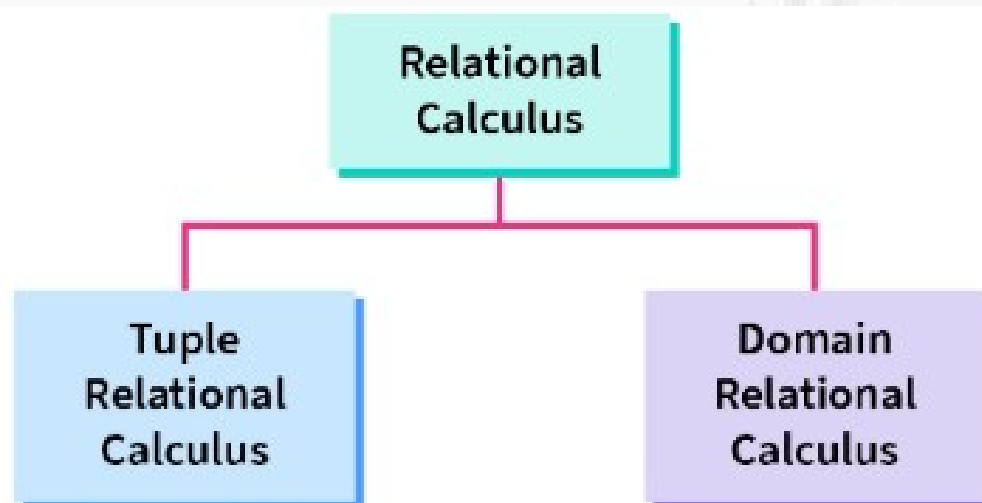


# Relational Calculus

- Relational Calculus in database management system (DBMS) is all about ***"What you want ?"***.
- Relational calculus does not tell us how to get the results from the Database, but it just cares about what we want.
- The theory of Relational calculus was introduced by computer scientist and mathematician Edgar Codd.



# Tuple Relational Calculus

- Tuple Relational Calculus is a **non-procedural query language** unlike relational algebra.
- Tuple Calculus provides only the description of the query but it does not provide the methods to solve it.
- Thus, it explains what to do but not how to do.
- In Tuple Calculus, a query is expressed as :

**$\{t \mid P(t)\}$**

where  $t$  = resulting tuples,

$P(t)$  = known as Predicate and these are the conditions that are used to fetch  $t$

- Thus, it generates set of all tuples  $t$ , such that Predicate  $P(t)$  is true for  $t$ .
- $P(t)$  may have various conditions logically combined with OR ( $\vee$ ), AND ( $\wedge$ ), NOT( $\neg$ ).  
It also uses quantifiers:  
 $\exists t \in r (Q(t))$  = "there exists" a tuple in  $t$  in relation  $r$  such that predicate  $Q(t)$  is true.  
 $\forall t \in r (Q(t))$  =  $Q(t)$  is true "for all" tuples in relation  $r$ .

# Example

- Let's say we have a table called “Employees” with the following attributes:
  - EmployeeID
  - Name
  - Salary
  - DepartmentID

To retrieve the names of all employees who earn more than \$50,000 per year, we can use the following TRC query:

$$\{ t \mid \text{Employees}(t) \wedge t.\text{Salary} > 50000 \}$$

In this query, the “Employees(t)” expression specifies that the tuple variable  $t$  represents a row in the “Employees” table. The “ $\wedge$ ” symbol is the logical AND operator, which is used to combine the condition “ $t.\text{Salary} > 50000$ ” with the table selection.

# Example contd..

Table-1: Customer

Customer name	Street	City
Saurabh	A7	Patiala
Mehak	B6	Jalandhar
Sumiti	D9	Ludhiana
Ria	A5	Patiala

Table-2: Branch

Branch name	Branch city
ABC	Patiala
DEF	Ludhiana
GHI	Jalandhar

Table-3: Account

Account number	Branch name	Balance
1111	ABC	50000
1112	DEF	10000
1113	GHI	9000
1114	ABC	7000

Table-4: Loan

Loan number	Branch name	Amount
L33	ABC	10000
L35	DEF	15000
L49	GHI	9000
L98	DEF	65000

Table-5: Borrower

Customer name	Loan number
Saurabh	L33
Mehak	L49
Ria	L98

Table-6: Depositor

Customer name	Account number
Saurabh	1111
Mehak	1113
Sumiti	1114



# Exercise

- 1:** Find the loan number, branch, amount of loans of greater than or equal to 10000 amount.
- 2:** Find the loan number for each loan of an amount greater or equal to 10000.
- 3:** Find the names of all customers who have a loan and an account at the bank.
- 4:** Find the names of all customers having a loan at the “ABC” branch.

# Solution

- **1:** Find the loan number, branch, amount of loans of greater than or equal to 10000 amount.

$$\{t \mid t \in \text{loan} \wedge t[\text{amount}] \geq 10000\}$$

- **2:** Find the loan number for each loan of an amount greater or equal to 10000.

$$\{t \mid \exists s \in \text{loan}(t[\text{loan number}] = s[\text{loan number}] \wedge s[\text{amount}] \geq 10000)\}$$

- **3:** Find the names of all customers who have a loan and an account at the bank.

$$\{t \mid \exists s \in \text{borrower}(t[\text{customer-name}] = s[\text{customer-name}]) \\ \wedge \exists u \in \text{depositor}(t[\text{customer-name}] = u[\text{customer-name}])\}$$

- **4:** Find the names of all customers having a loan at the “ABC” branch.

$$\{t \mid \exists s \in \text{borrower}(t[\text{customer-name}] = s[\text{customer-name}] \\ \wedge \exists u \in \text{loan}(u[\text{branch-name}] = \text{“ABC”} \wedge u[\text{loan-number}] = s[\text{loan-number}]))\}$$

# Domain Relational Calculus

- **Domain Relational Calculus** is a non-procedural query language equivalent in power to Tuple Relational Calculus.
- Domain Relational Calculus provides only the description of the query but it does not provide the methods to solve it.
- Domain Relational Calculus uses domain Variables to get the column values required from the database based on the predicate expression or condition.

- In Domain Relational Calculus, a query is expressed as,

$$\{ \langle x_1, x_2, x_3, \dots, x_n \rangle \mid P(x_1, x_2, x_3, \dots, x_n) \}$$

where,  $\langle x_1, x_2, x_3, \dots, x_n \rangle$  represents resulting domains variables and  $P(x_1, x_2, x_3, \dots, x_n)$  represents the condition or formula equivalent to the Predicate calculus.

- **Note:** The domain variables those will be in resulting relation must appear before  $\mid$  within  $\langle$  and  $\rangle$  and all the domain variables must appear in which order they are in original relation or table.

# Example

Table-1: Customer

Customer name	Street	City
Debomit	Kadamtala	Alipurduar
Sayantan	Udaypur	Balurghat
Soumya	Nutanchati	Bankura
Ritu	Juhu	Mumbai

Table-2: Loan

Loan number	Branch name	Amount
L01	Main	200
L03	Main	150
L10	Sub	90
L08	Main	60

Table-3: Borrower

Customer name	Loan number
Ritu	L01
Debomit	L08
Soumya	L03



# Queries

**1:** Find the loan number, branch, amount of loans of greater than or equal to 100 amount.

$$\{ \langle l, b, a \rangle \mid \langle l, b, a \rangle \in \text{loan} \wedge (a \geq 100) \}$$

**2:** Find the loan number for each loan of an amount greater or equal to 150.

$$\{ \langle l \rangle \mid \exists b, a (\langle l, b, a \rangle \in \text{loan} \wedge (a \geq 150)) \}$$

**3:** Find the names of all customers having a loan at the “Main” branch and find the loan amount.

$$\{ \langle c, a \rangle \mid \exists l (\langle c, l \rangle \in \text{borrower} \wedge \exists b (\langle l, b, a \rangle \in \text{loan} \wedge (b = \text{“Main”}))) \}$$