

Assignment-5

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Course – B.Tech CSE 4th sem

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1. **WAP to show the implementation of Kruskal's algorithm for computing MST, calculate its time complexity, and show the working of the algorithm.**
 2. **WAP for disjoint sets and show their working.**
-

Sol.

1) Kruskal's algorithm:

```
#include <stdio.h>
#include <stdlib.h>
#define MAX 100

typedef struct{
    int src, dest, weight;
} Edge;

typedef struct{
    int V, E;
    Edge* edges;
}Graph;

typedef struct Node{
    int parent;
    int rank;
} Node;

void printMST(Edge *result, int e);
Graph* createGraph(int V, int E);
int find(Node *set, int i);
void Union(Node *set, int x, int y);
int compareEdges(const void* a, const void* b);
void KruskalMST(Graph* graph);
```

```
int main() {
    int V = 6;
    int E = 8;
    Graph* graph = createGraph(V, E);

    graph->edges[0].src = 0;
    graph->edges[0].dest = 1;
    graph->edges[0].weight = 5;

    graph->edges[1].src = 0;
    graph->edges[1].dest = 3;
    graph->edges[1].weight = 11;

    graph->edges[2].src = 1;
    graph->edges[2].dest = 4;
    graph->edges[2].weight = 3;

    graph->edges[3].src = 1;
    graph->edges[3].dest = 2;
    graph->edges[3].weight = 7;

    graph->edges[4].src = 2;
    graph->edges[4].dest = 4;
    graph->edges[4].weight = 1;

    graph->edges[5].src = 2;
    graph->edges[5].dest = 5;
    graph->edges[5].weight = -3;

    graph->edges[6].src = 3;
    graph->edges[6].dest = 4;
    graph->edges[6].weight = 0;

    graph->edges[7].src = 5;
    graph->edges[7].dest = 4;
    graph->edges[7].weight = 2;

    KruskalMST(graph);

    free(graph->edges);
    free(graph);

    return 0;
}
```

```
//-----
```

```
void printMST(Edge *result, int e){  
    printf("Constructed MST:\n");  
    int totalWeight = 0;  
    for (int i = 0; i < e; ++i) {  
        printf("%d--> %d, weight = %d\n", result[i].src, result[i].dest, result[i].weight);  
        totalWeight += result[i].weight;  
    }  
    printf("MST weight: %d\n", totalWeight);  
}
```

```
//-----
```

```
Graph* createGraph(int V, int E) {  
    Graph* graph = (Graph*)malloc(sizeof(Graph));  
    graph->V = V;  
    graph->E = E;  
    graph->edges = (Edge*)malloc(graph->E * sizeof(Edge));  
    return graph;  
}
```

```
//-----
```

```
Node *InitSet(int V){  
    Node *set = (Node*)calloc(V, sizeof(Node));  
    for(int i = 0; i < V; i++){  
        set[i].parent = i;  
        set[i].rank = 0;  
    }  
    return set;  
}
```

```
//-----
```

```
int find(Node *set, int v){  
    if(set[v].parent != v) set[v].parent = find(set, set[v].parent);  
  
    return set[v].parent;  
}
```

```
//-----
```

```

void Union(Node*set, int x, int y){
    int rootX = find(set, x);
    int rootY = find(set, y);

    if (rootX == rootY) return;

    if(set[rootX].rank > set[rootY].rank){
        set[rootY].parent = rootX;
    }else if(set[rootX].rank < set[rootY].rank){
        set[rootX].parent = rootY;
    }else{
        set[rootY].parent = rootX;
        set[rootX].rank++;
    }
}

//-----

int compareEdges(const void* a, const void* b) {
    Edge* edgeA = (Edge*)a;
    Edge* edgeB = (Edge*)b;
    return edgeA->weight > edgeB->weight;
}

//-----

void KruskalMST(Graph* graph) {
    int V = graph->V;
    Edge result[MAX];
    int e = 0; // index of result
    int i = 0; // index of sorted edges
    //Sort edges in non-decreasing order
    qsort(graph->edges, graph->E, sizeof(graph->edges[0]), compareEdges);
    Node* set = InitSet(V);
    while (e < V- 1 && i < graph->E) {
        Edge nextEdge = graph->edges[i++];
        int x = find(set, nextEdge.src);
        int y = find(set, nextEdge.dest);
        if (x != y) {
            result[e++] = nextEdge;
            Union(set, x, y);
        }
    }
    printMST(result, e);
    free(set);
}

```

Output:

```
PS C:\Users\Ujjwal\Desktop\C\Graph_algo> cd "C:\Users\Ujjwal\Desktop\C\Graph_algo"
• Constructed MST:
2 --> 5, weight = -3
3 --> 4, weight = 0
2 --> 4, weight = 1
1 --> 4, weight = 3
0 --> 1, weight = 5
MST weight: 6
```

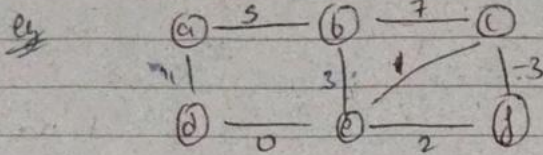
• Time complexity:

* Kruskal's Algo:

- 1) $A \leftarrow \phi$ \leftarrow solution set
- 2) $\forall v \in V: \left. \begin{array}{l} \text{find-set}(v) \\ \text{mark-set}(v) \end{array} \right\} \rightarrow O(V \alpha(V))$
- 3) $L =$ Arrange edges in non-decreasing order of their weight $\left. \right\} O(E \log E)$
- 4) $\forall (u, v) \in L: \left. \begin{array}{l} \text{if find-set}(u) \neq \text{find-set}(v) \\ A \leftarrow A \cup \{u, v\} \\ \text{Union}(u, v) \end{array} \right\} O(E \alpha(V))$
- 5) Return A

Time $\rightarrow O(V \alpha(V)) + O(E \alpha(V)) + O(E \log E)$
 $\approx O((V+E) \cdot \alpha(V)) + O(E \log E)$
 \rightarrow Since graph is connected $|E| \geq |V|-1$
hence we can replace V with E
 $O(E \alpha(V)) + O(E \log E)$
 $\rightarrow \alpha(V) = O(\log(V)) = O(\log(E))$
 $\boxed{O(E \log E)}$

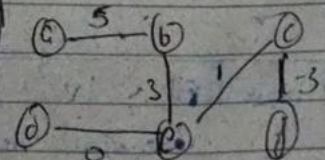
• Working:



After step ② & ③ {a} {b} {c} {d} {e} {f}

$$L = \{(c,d)^{-3}, (d,e)^0, (e,c)^1, (e,f)^2, (b,c)^3, (a,b)^5, (b,f)^7, (a,d)^1\}$$

Iteration	{a} {b} {c} {d} {e} {f}	A = ϕ
① (c,d)	$\text{findset}(c) \neq \text{findset}(d) \rightarrow \text{true} \rightarrow$ {a} {b} {c,d} {e} {f}	{(c,d)}
② (d,e)	$\text{findset}(d) \neq \text{findset}(e) \rightarrow \text{true} \rightarrow$ {c,d} {b} {d,e} {e} {f}	{(c,d), (d,e)}
③ (e,c)	$\text{findset}(e) \neq \text{findset}(c) \rightarrow \text{true} \rightarrow$ {a} {b} {c,d,e} {f}	{(c,d), (d,e), (e,c)}
④ (e,f)	$\text{findset}(e) \neq \text{findset}(f) \rightarrow \text{false}$ No change	—
⑤ (b,c)	$\text{findset}(b) \neq \text{findset}(c) \rightarrow \text{true} \rightarrow$ {a} {b,c,d,e} {f}	{(c,d), (d,e), (e,c), (b,c)}
⑥ (a,d) ⑥ (a,b)	$\text{findset}(a) \neq \text{findset}(b) \rightarrow \text{true} \rightarrow$ {a,b,c,d,e} {f}	{(c,d), (d,e), (e,c), (b,c), (a,b)}
⑦ (b,c)	false No change	
⑧ (a,d)	false No change	



2) WAP for disjoint sets and show their working.

```
#include<stdio.h>
#include<stdlib.h>

typedef struct Node{
    int parent;
    int rank;
} Node;
Node *InitSet(int V){
    Node *set = (Node*)calloc(V, sizeof(Node));
    for(int i = 0; i < V; i++){
        set[i].parent = i;
        set[i].rank = 0;
    }
    return set;
}
int find(Node *set, int v){
    if(set[v].parent != v) set[v].parent = find(set, set[v].parent);
    return set[v].parent;
}
void Union(Node*set, int x, int y){
    int rootX = find(set, x);
    int rootY = find(set, y);
    if (rootX == rootY) return;
    if(set[rootX].rank > set[rootY].rank){
        set[rootY].parent = rootX;
    }else if(set[rootX].rank < set[rootY].rank){
        set[rootX].parent = rootY;
    }else{
        set[rootY].parent = rootX;
        set[rootX].rank++;
    }
}
int main() {
    int V = 5;
    Node *set = InitSet(V);
    Union(set, 0, 1);
    Union(set, 1, 2);
    Union(set, 3, 4);
    printf("Find(2): %d\n", find(set, 2));
    printf("Find(4): %d\n", find(set, 4));
    Union(set, 2, 3);
    printf("Find(4) after merging: %d\n", find(set, 4));
    free(set);
    return 0;
}
```


Output:

```
PS C:\Users\Ujjwal\Desktop\C\Disj  
Find(2): 0  
Find(4): 3  
Find(4) after merging: 0  
PS C:\Users\Ujjwal\Desktop\C\Disj
```

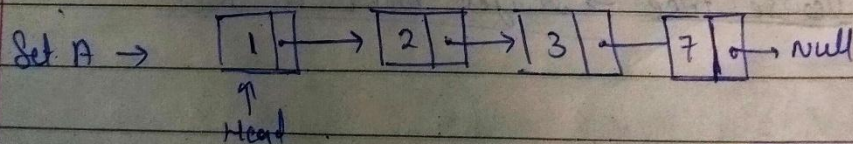
• Time complexity and Working:

* Disjoint sets: Refers to two or more sets that don't have any common element.
 $A \cap B = \emptyset$
eg: $A = \{1, 2, 3\}$ A & B are disjoint sets
 $B = \{4, 5, 6\}$

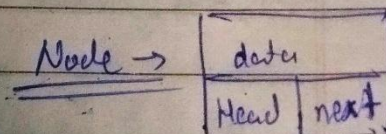
=> Operations

- FindSet(u): returns the Representative element of the set containing element 'u'.
- MakeSet(u): returns a singleton set containing only 'u' element $\rightarrow \{u\}$
- Union(u, v): if 'u' & 'v' belongs to different sets (RE of $u \neq RE$ of v) then merge the two sets
eg $\{1\}$ $\{2\}$
 $union(1, 2) \rightarrow \{1, 2\}$

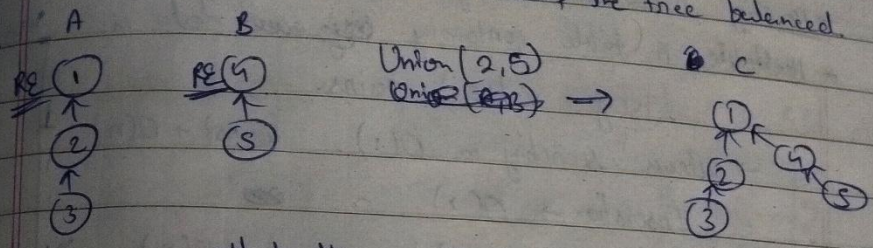
→ Implementation using linked list



- each node stores Head pointer which is RE of that set.



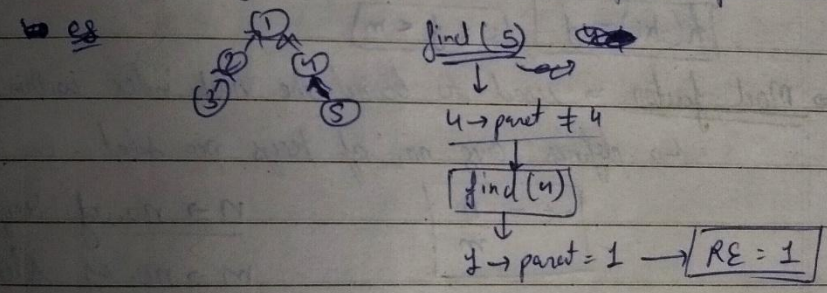
- Rooted tree → set is represented as tree
- each element points to the ~~root~~ parent
- rank is used to keep the tree balanced.



attach the smaller tree to root of larger tree

⇒ findset → uses path compression

recursively call findset(u) until ~~parent~~ u → parent = u



⇒ Time Complexity of Rooted tree method → $O(\alpha(v))$

~~α is Ackermann function~~

α → Inverse Ackermann function

(grows very slow) $\approx \underline{\underline{O(1)}}$