



Aalto University

Algorithm for calculating HNF of a matrix $M \in \mathbb{Z}^{n \times m}$ in C++

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Pseudo code

- You create a matrix L to change the leftmost non lower triangular columns lowest (positionally) non zero element into zero such that $\det(L)=1$
- Proceed upwards until the column is in lower triangular form
- Move onto the next column
- Continue until the matrix is lower triangular form
- Use the same idea to transform the pivots into strictly increasing order

Algorithm 1: HNF(M)

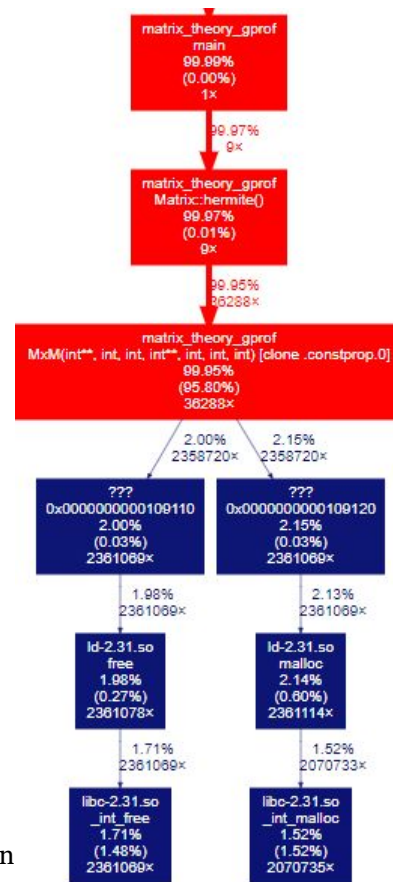
```
[n,m] ← size(M)
U ← IMat(n)
L ← IMat(n)
H ← M
for j ← 1...m do
  for i ← n...j + 1 do
    if  $H_{i,j} = 0$ 
      Nothing
    else if  $H_{i-1,j} = 0$ 
      Swap rows  $i - 1$  and  $i$  in  $L$ 
    else
       $a \leftarrow H_{i-1,j}, b \leftarrow H_{i,j}$ 
       $g \leftarrow \gcd(a, b)$ 
       $[x, y] = \text{DiophantesSolver}(a, b, g)$ 
       $\alpha \leftarrow \frac{a}{g}, \beta \leftarrow \frac{b}{g}$ 
       $L_{i,i} \leftarrow \alpha, L_{i,i-1} \leftarrow -\beta, L_{i-1,i} \leftarrow y, L_{i-1,i-1} \leftarrow x,$ 
    end
    U ← LU, H ← LH
    L ← IMat(n)
  end
end
Permute rows of H st. pivot locations are increasing
Store permutation matrix in L
U ← LU
L ← IMat(n)
while H pivot locations not strictly increasing at  $H_{i,j}, H_{i-1,j}$ 
   $a \leftarrow H_{i-1,j}, b \leftarrow H_{i,j}$ 
   $g \leftarrow \gcd(a, b)$ 
   $[x, y] = \text{DiophantesSolver}(a, b, g)$ 
   $\alpha \leftarrow \frac{a}{g}, \beta \leftarrow \frac{b}{g}$ 
   $L_{i,i} \leftarrow \alpha, L_{i,i-1} \leftarrow -\beta, L_{i-1,i} \leftarrow y, L_{i-1,i-1} \leftarrow x,$ 
  U ← LU, H ← LH
  L ← IMat(n)
end
return [U, H]
```

Pseudo code complexity

- Current complexity is $O(n^5)$
 - Since Matrix multiplication is $O(n^3)$
- Everything else is negligible time
 - $O(\log(\max(A, B)))$ time complexity for Diophante solver
 - $O(n)$ time for swapping rows of L
 - $O(\log(\min(a, b)))$ time for gcd

Optimizing the algorithm

- Running profiling we can detect bottlenecks
 - Clearly only relevant bottleneck is matrix multiplication

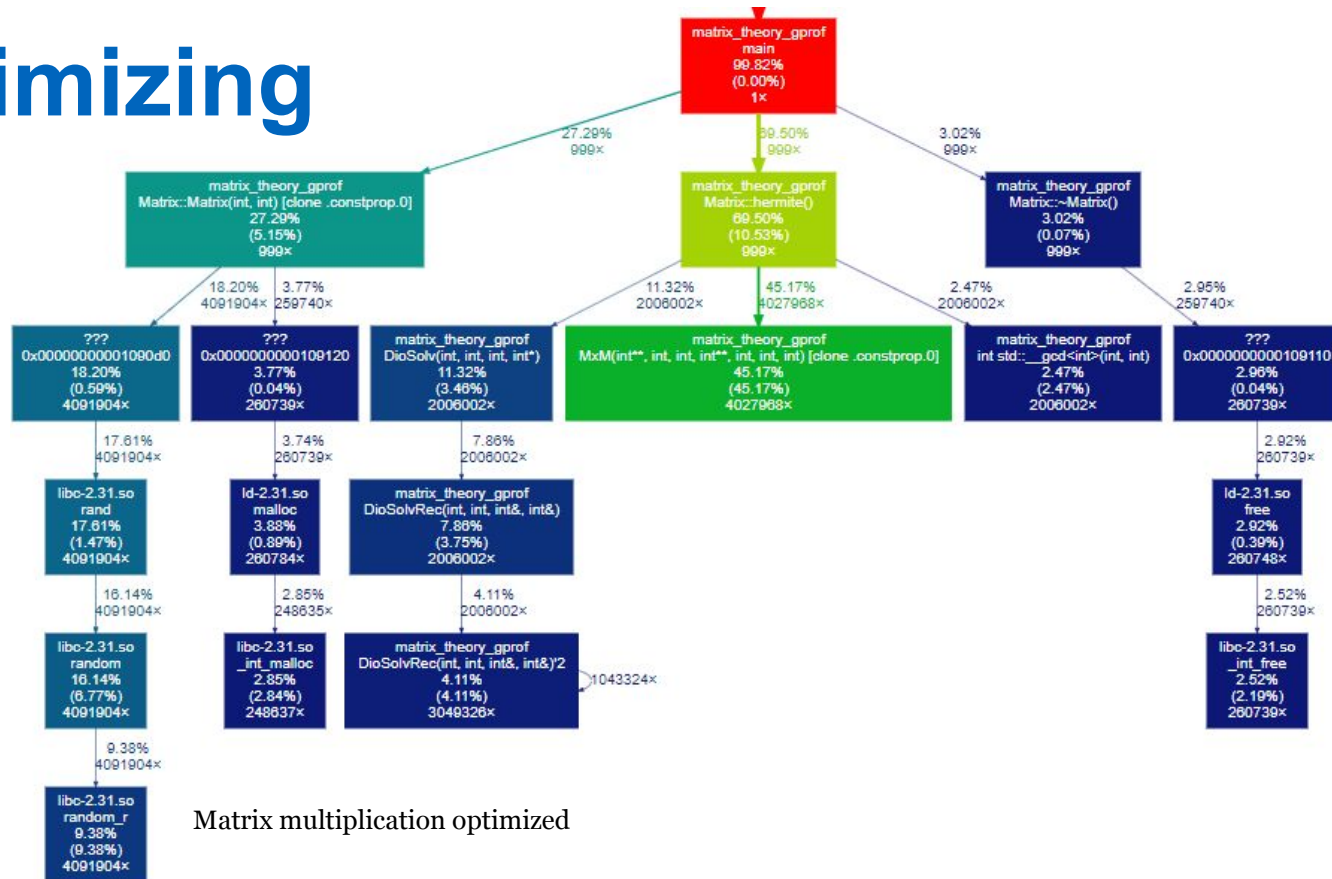


Naive implementation

Optimizing the algorithm

- Since L is a sparse matrix, with the form $I \oplus G \oplus I = L$, where $G \in \mathbb{Z}^{2 \times 2}$, we see that in LU , L only acts on the rows where G is located
- This means that $LU = [U_1, GU_2, U_3]^T$ where $U_2 \in \mathbb{Z}^{2 \times m}$
- Hence we can change the matrix multiplication used in the algorithm to a specific case with complexity $O(m)$
- Also resetting L to a identity matrix can be done by replacing G with a identity matrix $I \in \mathbb{Z}^{2 \times 2}$

Optimizing

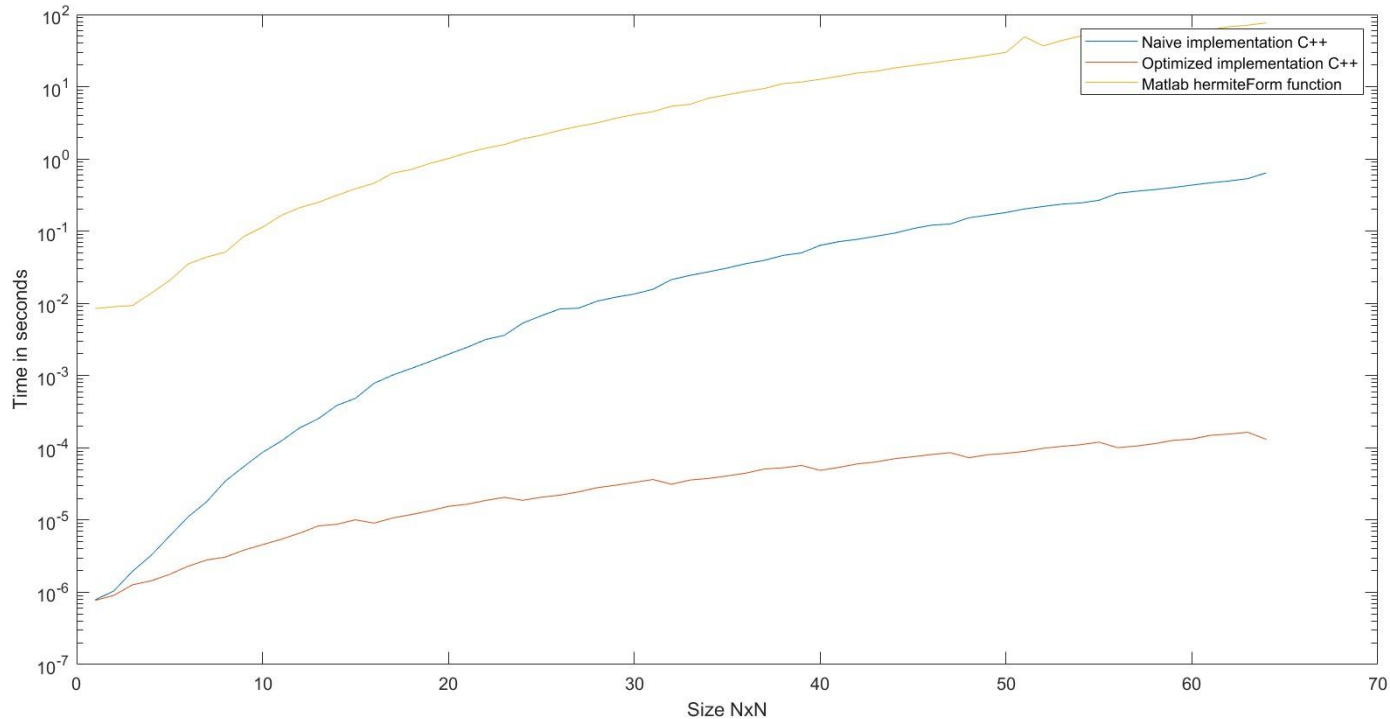


Matrix multiplication optimized

Optimized algorithm

- Still over half of the time used in calculating the HNF is taken by matrix multiplication
- Diophantine equation solver takes around 15% of the time
- And around 15% is spent inside the HNF algorithm itself
 - Mainly updating variables and if statements

Comparing algorithms to Matlab



Comparing algorithms to Matlab

- Clearly the optimized version is a lot more efficient in runtime
- Compared to the matlab version both were vastly superior, why?
 - Possibly since Matlab automatically stores values as double precision floats
 - Matlab is used for mostly matrices over real and complex numbers
 - HNF implementation might be of little importance
 - Overflow checking, see next slide

Shortcomings of algorithm in C++

- There is no integer overflow guard in the algorithm
 - Since the values, especially at the bottom right grow large extremely quickly overflow errors happen and the result is not correct
 - This starts happening when elements in $M \in \mathbb{Z}^{n \times n}$ being between -10 and 10 and $n > 5$

Sources

- <https://github.com/KristianWasas/HNF-calc>
- Chapter 2 of lecture notes
- <https://github.com/google/benchmark>
 - For benchmarking
- <https://github.com/jrfonseca/gprof2dot>
 - For visualization of profiling data
- Matlab R2022b for plotting