

# Algorithm for calculating HNF of a matrix M∈Z<sup>n×m</sup> in C++

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#### Pseudo code

- You create a matrix L to change the leftmost non lower triangular columns lowest (positionally) non zero element into zero
- Proceed upwards until the column is in lower triangular form
- Move onto the next column
- Continue until the matrix is lower triangular form

#### Algorithm 1: HNF(M)

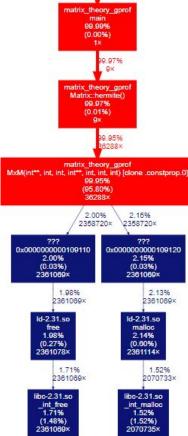
```
[n,m] \leftarrow size(M)
U \leftarrow IMat(n)
L \leftarrow IMat(n)
H \leftarrow M
for j \leftarrow 1...m do
    for i \leftarrow n...j + 1 do
         if H_{i,j} = 0
              Nothing
         else if H_{i-1,j} = 0
              Swap rows i-1 and i in L
         else
             a \leftarrow H_{i-1,j}, b \leftarrow H_{i,j}
             q \leftarrow qcd(a,b)
              [x, y] = DiophantesSolver(a, b, g)
             \alpha \leftarrow \frac{a}{a}, \beta \leftarrow \frac{b}{a}
              L_{i,i} \stackrel{g}{\leftarrow} \alpha, \ L_{i,i-1} \stackrel{g}{\leftarrow} -\beta, \ L_{i-1,i} \stackrel{g}{\leftarrow} y, \ L_{i-1,i-1} \stackrel{g}{\leftarrow} x
         end
         U \leftarrow LU, H \leftarrow LH
         L \leftarrow IMat(n)
    end
end
return [U, H]
```

#### Pseudo code complexity

- Current complexity is O(n<sup>5</sup>)
  - Since Matrix multiplication is O(n³)
- Everything else is negligible time
  - O(log(max(A, B))) time complexity for Diophante solver
  - O(n) time for swapping rows of L
  - O(log(min(a, b)) time for gcd

# **Optimizing the algorithm**

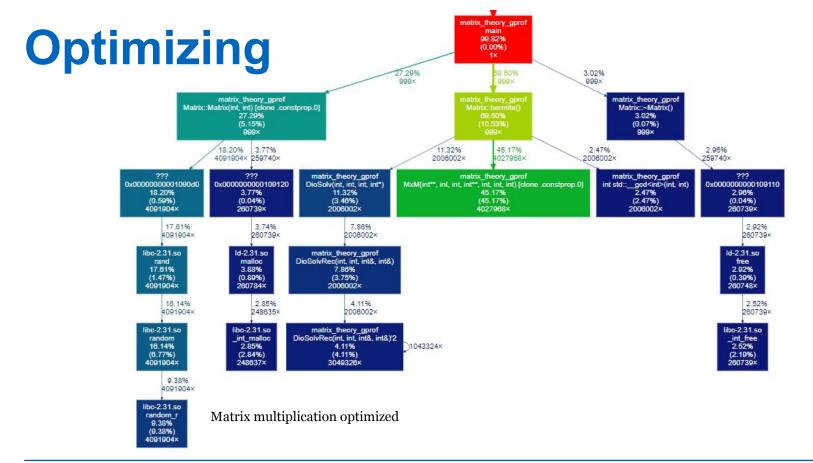
- Running profiling we can detect bottlenecks
  - Clearly only relevant bottleneck is matrix multiplication



Naive implementation

#### **Optimizing the algorithm**

- Since L is a sparse matrix, with the form  $I \oplus G \oplus I = L$ , were  $G \in \mathbb{Z}^{2 \times 2}$ , we see that in LU, L only acts on the rows where G is located
- This means that  $LU = [U_1, GU_2, U_3]^T$  were  $U_2 \in \mathbb{Z}^{2 \times m}$
- Hence we can change the matrix multiplication used in the algorithm to a specific case with complexity O(m)
- Also resetting L to a identity matrix can be done by replacing G with a identity matrix  $I \in \mathbb{Z}^{2 \times 2}$

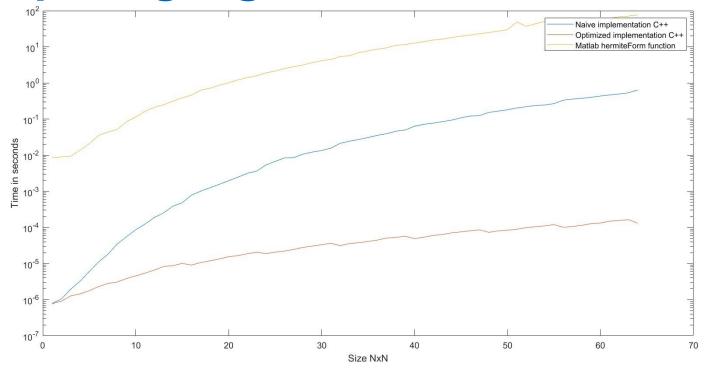




## **Optimized algorithm**

- Still over half of the time used in calculating the HNF is taken by matrix multiplication
- Diophantine equation solver takes around 15% of the time
- And around 15% is spent inside the HNF algorithm itself
  - Mainly updating variables and if statements

#### **Comparing algorithms to Matlab**





#### **Comparing algorithms to Matlab**

- Clearly the optimized version is a lot more efficient in runtime
- Compared to the matlab version both were vastly superior, why?
  - Possibly since Matlab automatically stores values as double precision floats
  - Matlab is used for mostly matrices over real and complex numbers
    - HNF implementation might be of little importance
  - Overflow checking, see next slide

## **Shortcomings of algorithm in C++**

- There is no integer overflow guard in the algorithm
  - Since the values, especially at the bottom right grow large extremely quickly overflow errors happen and the result is not correct
  - This happens with elements in  $M \in \mathbb{Z}^{n \times n}$  being between -10 and 10 and n > 5
- Sometimes gives a U, H such that det(U)=-1 and UM=H
  - Bug for some unknown reason, fairly rare
    - 1 in 10
  - Fix is to just let  $U \leftarrow -U$  and  $H \leftarrow -H$

#### Sources

- <a href="https://github.com/KristianWasas/HNF-calc">https://github.com/KristianWasas/HNF-calc</a>
- Chapter 2 of lecture notes
- https://github.com/google/benchmark
  - For benchmarking
- https://github.com/jrfonseca/gprof2dot
  - For visualization of profiling data
- Matlab R2022b for plotting