

# # LC - Biweekly Contest - 151:-

A] 3467. Transform Array By Parity:- ✓

B] 3468. Find number of copy arrays:-

original = [1, 2, 3, 4]

bounds = [1, 10], [2, 9], [3, 8], [4, 7]

copy

x

y

z

q

(no. of ele. satisfy the condit<sup>n</sup> given)

ans = xxyxzxy

here 1 ele.

copy = [x, , , ]

org[i] - org[i-1] = copy[i] - copy[i-1]

d<sub>1</sub>

x + d<sub>1</sub>

x

(x + d<sub>1</sub>) + d<sub>2</sub>

(x + d<sub>1</sub> + d<sub>2</sub>) + d<sub>3</sub>

d<sub>1</sub> = d<sub>2</sub> = d<sub>3</sub> = d<sub>4</sub> ⇒ (here)

copy = [x, x+1, x+2, x+3]

bounds:

[1, 10]

[2, 9]

[3, 8]

[4, 7]

1 ≤ x ≤ 10

1 ≤ x ≤ 8

1 ≤ x ≤ 7

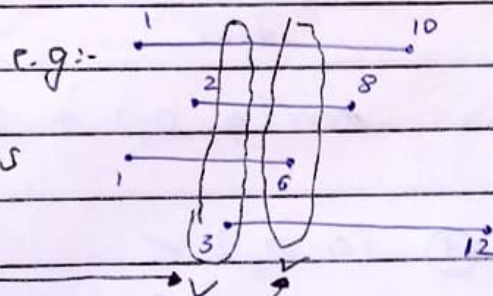
1 ≤ x ≤ 4

⇒ 1 ≤ x ≤ 4 final

max of min. side values (lower bounds)

min. of max. side values (upper bounds)

ans



C] 3469. Find min. cost to remove array elements:-

nums[] = [a<sub>1</sub>, a<sub>2</sub>, a<sub>3</sub>, a<sub>4</sub>] no. of op. = 2

For 2<sup>nd</sup> op.:-

nums[] = [a<sub>1</sub>, a<sub>2</sub>, a<sub>3</sub>, a<sub>4</sub>, a<sub>5</sub>] = 3



→ last op. can have either 2 or 1 element.

1 op.  $a_1, a_2, a_3$

3 possibilities

- $a_1, a_2$
- $a_2, a_3$
- $a_3, a_1$

whenever we have to try for all the possibilities → DP

→ so try for recu. for all possibilities.

→ then merge <sup>for</sup> our answer.

num[] =  $a_1, a_2, a_3, a_4, a_5$   
↑    ↑    ↑  
idx-1 idx idx+1  
possibilities

• We have not used  $idx_1, idx_2$  &  $idx_3$

∴  $dp[i_1][i_2][i_3]$

•  $O(n^3)$

• We have to memorize it also.

• We can only go up to  $O(n^2 \log n)$  max.

①  $(a_1, a_2)$ , cost =  $\max(a_1, a_2)$

Next stage:-

num =  $a_1, a_2, a_3, a_4, a_5$   
↑    ↑    ↑    ↑    ↑  
idx-1 idx idx+1 idx+2 idx+3  
previdx

cost =  $\max(a_1, a_2) + \text{solve}(\text{previdx}, \text{idx}+2, \text{idx}+3)$

previdx = idx+1

② Do we need previdx?

$(a_2, a_3)$  ✓, cost =  $\max(a_2, a_3)$

num  $a_1, a_2, a_3, a_4, a_5$   
↑    ↑    ↑    ↑    ↑  
idx-1 idx idx+1 idx+2 idx+3  
previdx

previdx ↑    ↑  
idx-1    idx    idx+1

$\max(a_2, a_3) + \text{solve}(\text{previdx}, \text{idx}+2, \text{idx}+3)$

→ previdx = idx-1

③  $(a_1, a_3)$  ✓

$a_1, a_2, a_3, a_4, a_5$   
↑    ↑    ↑    ↑    ↑  
idx-1 idx idx+1 idx+2 idx+3  
previdx

$\min(a_1, a_3) + \text{solve}(\text{previdx}, \text{idx}+2, \text{idx}+3)$

→ previdx = idx

•  $O(n^3)$

→ only previdx is variable

• So, we only need  $(prevIdx, idx)$  to derive minCost.

• code:- Base case

memoization case

↳  $dp[idx][prevIdx] \neq -1 \Rightarrow$  return  $dp[idx][prevIdx]$

• From 3 cases:-

① opt-1 :-  $\max(nums[prevIdx], nums[idx]) + \text{solve}(idx+1, idx+2)$

② opt-2 :-  $\max(nums[idx], nums[idx+1]) + \text{solve}(prevIdx, idx+2)$

③ opt-3 :-  $\max(nums[prevIdx], nums[idx+1]) + \text{solve}(idx, idx+2)$

So  $\text{minCost}(opt1, opt2, opt3)$

return  $dp[idx][prevIdx]$  minCost  $\Rightarrow$  memoize

• Base case:-

↳ 1 ele Rem. or 2  $\Rightarrow$  discussed initially

↳ if  $(idx == n)$  return  $\text{num}(prevIdx)$

if  $(idx == n-1)$  return  $\max(\text{num}(prevIdx), \text{num}(idx))$