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PERFECT SCORE OF **150/150** AS A TESTAMENT TO EXECEPTIONAL E-LEARNING METHODS

#University Category

# Unit 2: Process and Thread Management

Lecture 6
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- 1. Threads: Threads and its benefits
- 2. Multi-threading models
- 3. Types of thread
  - 1. Kernel Level thread
  - 2. User level thread
  - 3. Hybrid threads



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#### **Learning & Course Outcomes**

LO1: To understand Threads and Their Benefits

LO2: To evaluate Multi-threading Models

LO3:To differentiate Types of Threads

CO2: Evaluate and analyze process and thread scheduling techniques, discerning their benefits and challenges.



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#### **Motivation**

- Most modern applications are multithreaded
- Threads run within application
- Multiple tasks with the application can be implemented by separate threads
  - Update display
  - Fetch data
  - Spell checking
  - Answer a network request
- Process creation is heavy-weight while thread creation is light-weight
- Can simplify code, increase efficiency
- Kernels are generally multithreaded



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#### **Multithreaded Server Architecture**

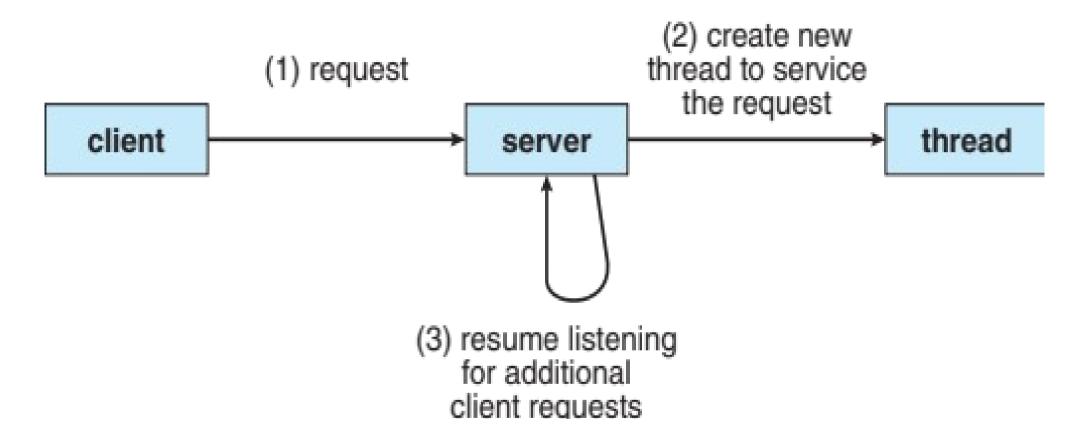


Fig 1. Multithreaded Server Architecture [1]



#### **Benefits**

- **Responsiveness** may allow continued execution if part of process is blocked, especially important for user interfaces
- **Resource Sharing** threads share resources of process, easier than shared memory or message passing
- Economy cheaper than process creation, thread switching lower overhead than context switching
- **Scalability** process can take advantage of multiprocessor architectures



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#### **Multicore Programming**

- **Multicore or multiprocessor** systems putting pressure on programmers, challenges include:
  - Dividing activities
  - Balance
  - Data splitting
  - Data dependency
  - Testing and debugging
- **Parallelism** implies a system can perform more than one task simultaneously
- Concurrency supports more than one task making progress
  - Single processor / core, scheduler providing concurrency



#### **Multicore Programming (Cont.)**

- Types of parallelism
  - Data parallelism distributes subsets of the same data across multiple cores, same operation on each
  - **Task parallelism** distributing threads across cores, each thread performing unique operation
- As # of threads grows, so does architectural support for threading
  - CPUs have cores as well as hardware threads
  - Consider Oracle SPARC T4 with 8 cores, and 8 hardware threads per core



### Concurrency vs. Parallelism

Concurrent execution on single-core system:

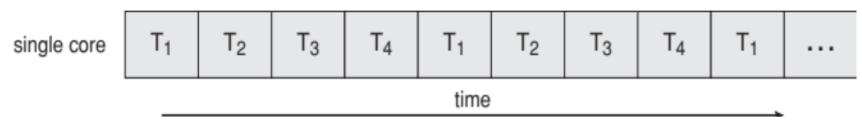


Fig 2. Concurrency [1]

Parallelism on a multi-core system:

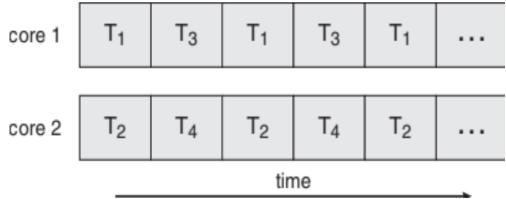


Fig 3. Parallelism [1]



#### Single and Multithreaded Processes

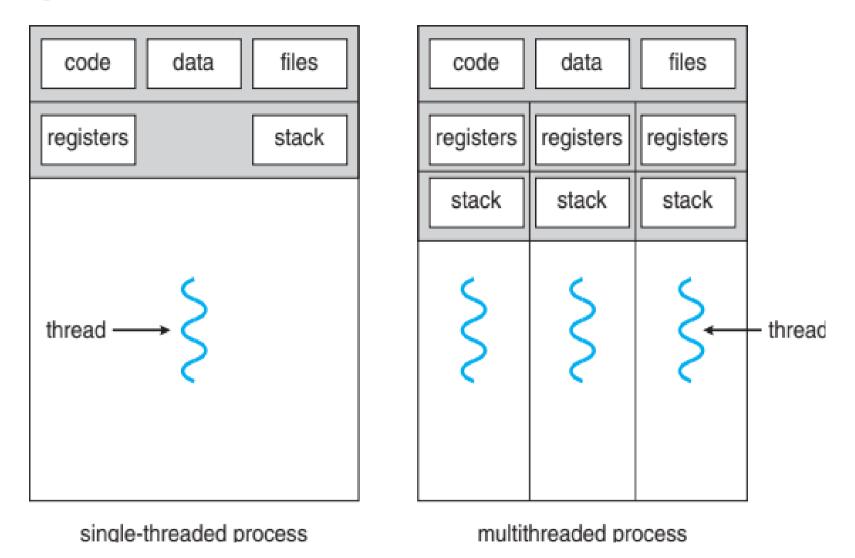


Fig 4. Single and Multithreaded Processes [1]



#### Concurrency vs. Parallelism

- Identifies performance gains from adding additional cores to an application that has both serial and parallel components
  - S is serial portion
  - N processing cores

$$speedup \le \frac{1}{S + \frac{(1-S)}{N}}$$

- That is, if application is 75% parallel / 25% serial, moving from 1 to 2 cores results in speedup of 1.6 times
- As N approaches infinity, speedup approaches 1 / S
- Serial portion of an application has disproportionate effect on performance gained by adding additional cores



#### Numerical of Amdahl's Lab

#### Exercise and solution on Amdahl's Law

- 1. What is the overall speedup if you make 10% of a program 90 times faster?
- 2. What is the overall speedup if you make 90% of a program 10 times faster?

· Amdahl's law

$$OverallSpeedup = \frac{1}{(1-f) + \frac{f}{s}}$$

 What is the overall speedup if you make 10% of a program 90 times faster?

$$\frac{1}{(1-0.1) + \frac{0.1}{90}} \approx \frac{1}{0.9011} \approx 1.11$$

 What is the overall speedup if you make 90% of a program 10 times faster

$$\frac{1}{(1-0.9)+\frac{0.9}{10}} = \frac{1}{0.19} \approx 5.26$$



#### **User Threads, Kernel Threads and Hybrid Threads**

- **User threads** management done by user-level threads library
- Three primary thread libraries:
  - POSIX Pthreads
  - Windows threads
  - Java threads
- **Kernel threads** Supported by the Kernel
- Examples virtually all general-purpose operating systems, including:
  - Windows
  - Solaris
  - Linux
  - Tru64 UNIX
  - Mac OS X
- **Hybrid thread** Combination of both



#### **Multithreading Models**

- Many-to-One
- One-to-One
- Many-to-Many



#### Many-to-One Model

- Many user-level threads mapped to single kernel thread
- One thread blocking causes all to block
- Multiple threads may not run in parallel on muticore system because only one may be in kernel at a time
- Few systems currently use this model
- Examples:
  - Solaris Green Threads
  - GNU Portable Threads

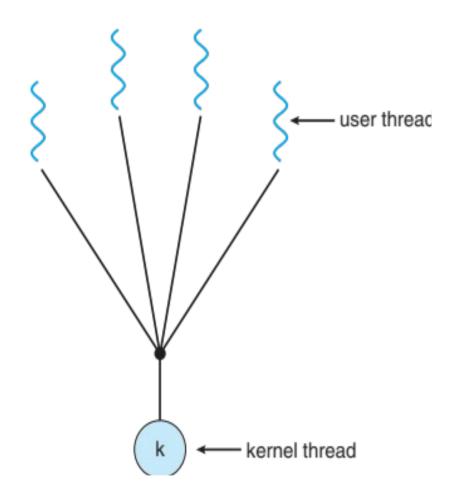


Fig 5. Many-to-One Model [1]



#### **One-to-One Model**

- Each user-level thread maps to kernel thread
- Creating a user-level thread creates a kernel thread
- More concurrency than many-to-one
- Number of threads per process sometimes restricted due to overhead
- Examples
  - Windows
  - Linux
  - Solaris 9 and later

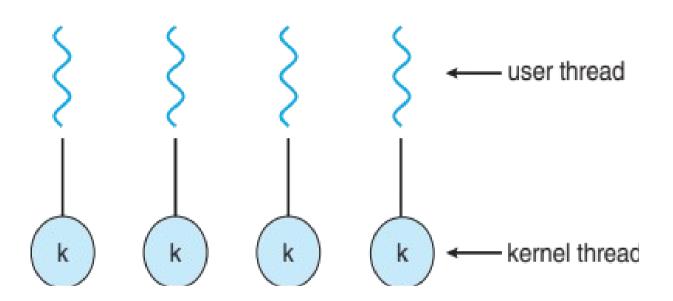


Fig 6. One-to-One Model [1]



#### **Many-to-many Model**

- Allows many user level threads to be mapped to many kernel threads
- Allows the operating system to create a sufficient number of kernel threads
- Solaris prior to version 9
- Windows with the ThreadFiber package

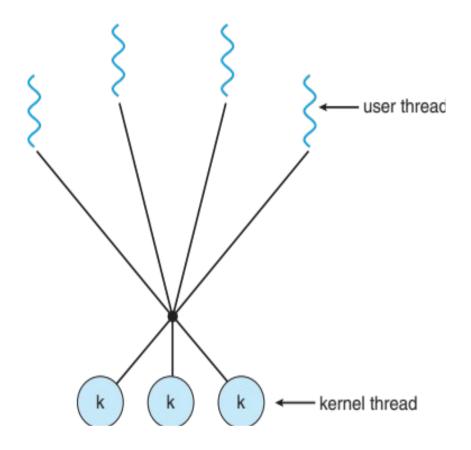


Fig 7. Many-to-Many Model [1]



#### **Two-Level Model**

- Similar to M:M, except that it allows a user thread to be bound to kernel thread
- Examples
  - IRIX
  - HP-UX
  - Tru64 UNIX
  - Solaris 8 and earlier

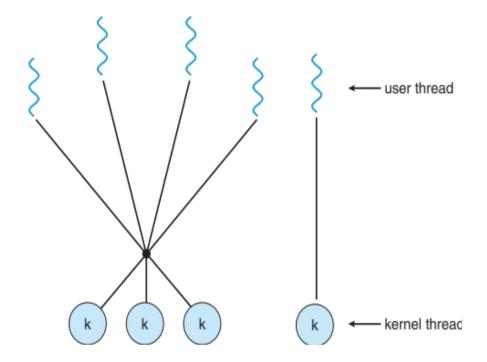


Fig 8. Two-Level Model [1]



#### Difference between multithreading models

| Aspect        | Many-to-One Model   | One-to-One Model  | Many-to-Many Model  |
|---------------|---|---|---|
| Description   | Maps many user-level threads to one kernel thread.                    | Maps each user-level thread to a kernel thread.   | Maps many user-level threads to many kernel threads.  |
| Concurrency   | Limited concurrency; only one thread can access the kernel at a time. | High concurrency; multiple threads can run in parallel on multiple processors.          | Flexible concurrency; can handle more threads than kernel threads.                          |
| Advantages    | Simple implementation; minimal overhead.                              | Allows true parallelism on multiprocessors; each thread can be independently scheduled. | Avoids the limitations of both many-to-one and one-to-one models; ent thread management.    |
| Disadvantages | Blocking of one thread blocks all threads; no true parallelism.       | Significant overhead for creating and managing many kernel threads.                     | More complex to implement than other models.  |
| Suitability   | Suitable for applications that do not require high concurrency.       | Suitable for applications that require high concurrency and true parallelism.           | Suitable for applications that need a balance between concurrency and resource utilization. |
| Examples      | Green threads in early versions of Java.                              | Windows operating system, Linux pthreads.   | Solaris operating system.   |



#### **Summary**

- **Threads and Their Benefits:** Explains what threads are and the advantages they offer in concurrent execution.
- **Multi-threading Models:** Discusses different models used for implementing multi-threading.
- **Types of Threads:** Describes various thread types, including kernel-level, user-level, and hybrid threads.
- **Kernel Level Threads:** Provides details about threads managed by the operating system kernel.
- **User Level Threads:** Covers threads managed by user-level libraries.
- **Hybrid Threads:** Combines aspects of both kernel and user-level threading for improved performance and flexibility.



#### **Reference Material**

- [1]. Silberschatz, A. & Galvin, P. (2009) Operating System Concepts. 8th ed. NJ: John Wiley & Sons, Inc.
  - Download Link- https://www.mbit.edu.in/wp-content/uploads/2020/05/Operating\_System\_Concepts\_8th\_EditionA4.pdf
- [2]. NPTEL Video Lecture: <a href="https://onlinecourses.nptel.ac.in/noc24">https://onlinecourses.nptel.ac.in/noc24</a> cs80/preview



#### MCQ's

- 1. Which of the following best defines a thread?
  - a) An independent program
  - b) A lightweight process within a process
  - c) A system call for executing code
  - d) A subroutine for handling exceptions
- 2. What is one of the primary benefits of using threads in a program?
  - a) Decreased parallelism
  - b) Increased memory usage
  - c) Improved responsiveness
  - d) Higher code complexity



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#### MCQ's

- 3. Which of the following is not a multi-threading model?
  - a) Kernel Level thread
  - b) User Level thread
  - c) Hybrid thread
  - d) Sequential thread
- 4. Which type of thread is managed entirely by the kernel of the operating system?
  - a) Kernel Level thread
  - b) User Level thread
  - c) Hybrid thread
  - d) Sequential thread



#### MCQ's

- 5. What is a characteristic of User Level threads?
  - a) They require less context-switching overhead.
  - b) They offer better OS support for thread management.
  - c) They are faster in context-switching than Kernel Level threads.
  - d) They are not portable across different operating systems.
- 6. Hybrid threads combine features of which two types of threads?
  - a) Kernel Level and User Level threads
  - b) Sequential and Concurrent threads
  - c) Synchronous and Asynchronous threads
  - d) Critical and Non-Critical threads



#### MCQ's Answers

#### Answers:

- 1.b) A lightweight process within a process
- 2.c) Improved responsiveness
- 3.d) Sequential thread
- 4.a) Kernel Level thread
- 5.a) They require less context-switching overhead.
- 6.a) Kernel Level and User Level threads



**What's Next** 

- 1. Thread Scheduling
  - 1. Content Scope
  - 2. Pthread Scheduling
  - 3. Threading Issues





## Thank You

