

first 64 bits → Network ID  
Second 64 bits → Interface ID

14/12/2023

→ Header size

- └ TCP: 28 bytes
- └ UDP: 8 bytes

\* → TCP: end to end processing connection

\* → Establishing a network takes 3 steps.

\* → Removing a network connection takes 4 steps.

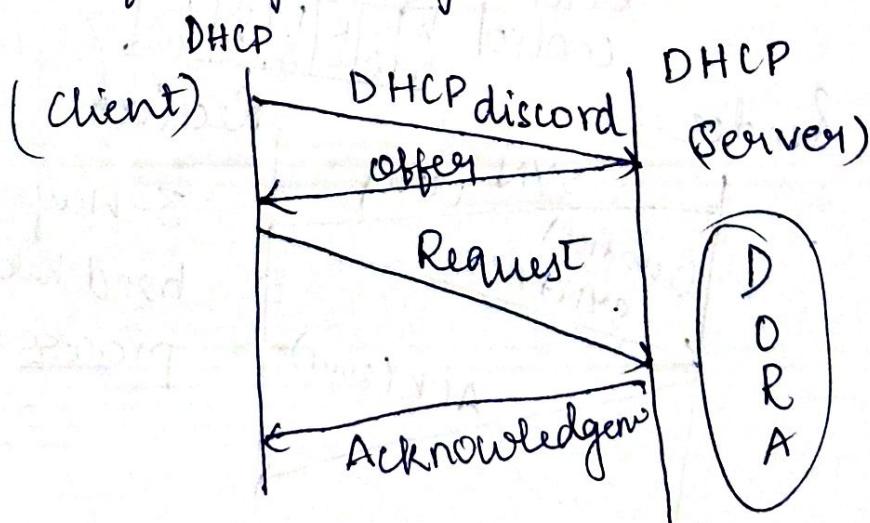
21/12/2023

→ ARP: Address Resolution Protocol

└ resolves IP to MAC  
length of MAC = 48 bits

→ DHCP: Dynamic Host Configuration Protocol

└ provides: IP address, Subnet Mask,  
default gateway, DNS server address.

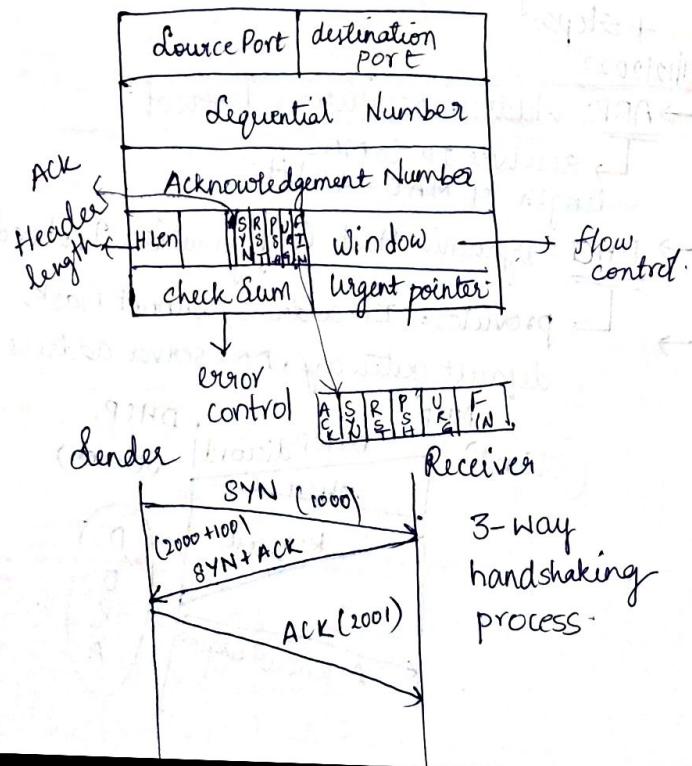


→ Elements of Transport protocol: Transport/TCP protocol

- Addressing
- Flow control & Error Control.
- Connection Establishment / Release
- Multiplexing & Demultiplexing
- Crash / Error recovery.

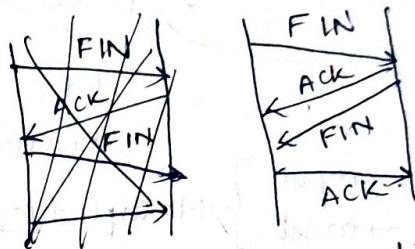
→ Flow and Error control are not possible in UDP protocol.

\* TCP header: (20 byte) Max(60 bytes) ⇒ Can be extended

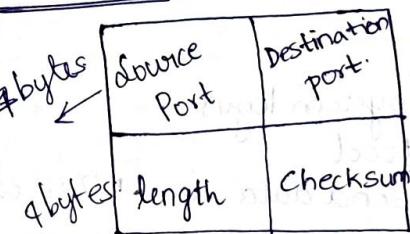


→ RST is set; if the connection is lost suddenly, the connection can be RESET to 1.

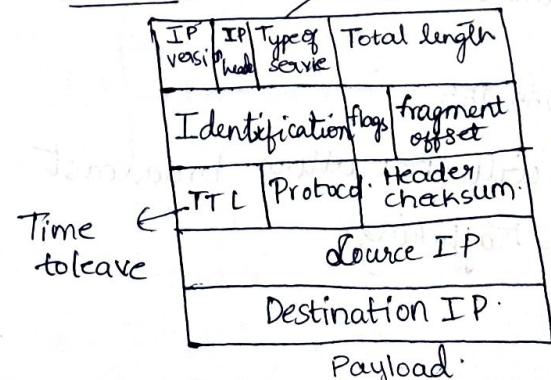
→ PSH, URG.



\* UDP header: → videos & audio streaming.



\* IP header: → gives the priorities of requests.



X10 → tells the receiver that more packets are arriving.

X01 → tells the receiver to close the connection.

\* Default port no. of ftp: 20 & 21

functions

data control connection connection

{ put | mput } Multiple files  
{ get | mget } Single file

\* telnet: (23)

remote system login

TCP protocol

Do not send data in encrypted form and data is not secure.

\* Limited and direct broadcast addresses

255.255.255.255

→ routers will not allow broadcast address trafficking.

→ What are the fields required to reassemble the fragments at receiver in IP address?

→ flags & ~~questions~~ fragments.

→ Application Layer Protocols

\* ICMP: Internet Control Msg protocol

→ Network layer

NAT

\* DNS: Domain Name System

→ resolves domain name to IP address

resolver → 2 ways: iterative & recursive

used to find

the IP address.

to resolve domain to IP in systems.

\* 5 tuple of DNS

→ Domain Name

→ Time to leave

→ Class (IPv4 | IPv6)

→ Type

→ Value

\* Telnet & FTP uses TCP protocol.

\*  $200 \cdot 40 \cdot 50 \cdot 60 / 2^4$

$255 \cdot 255 \cdot 255 \cdot 0$

$200 \cdot 40 \cdot 50 \cdot 0 \rightarrow \text{Network}$

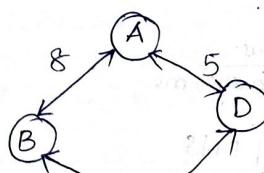
$200 \cdot 40 \cdot 50 \cdot 255 \rightarrow \text{Broadcast}$

\* TCP connections needed in FTP: 3-way hand shaking

\* 3 Main divisions of DNS:

- Generic Domains
- Country domains
- Inverse domains

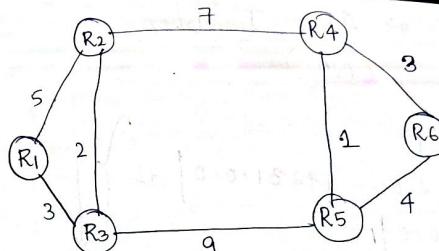
\* Problem:



A's routing table:

	destination	next hop	distance
	A		0
A	B	B	8
B	C	D	8
C	D		5
D			

$B \leftrightarrow C \Rightarrow \text{unused}$



$R_1 \rightarrow R_6$

$R_1 - R_3 - R_2 \rightarrow R_4 - R_6$

→ Take the unused links and reduce the distances by 2 (or) any value

\*  $172 \cdot 16 \cdot 0 \cdot 0 / 20 \Rightarrow$  \* subnet into 8 subnets

$\downarrow$  \* subnet mask

255.

for network \* ~~subnet weight~~

→ \* calculate range of each

→ 20 bits are there

for network

$\rightarrow 111111 \cdot 111111 \cdot 11110000 \cdot 00000000 / 20$

$\rightarrow 111111 \cdot 111111 \cdot 111100000000 / 20$

→ subnetmasking

## \* Network Address Translation \* (NAT)

10.0.0.0/8  
172.16.0.0 - 172.31.0.0/12 }  
192.16.0.0/16  
private address

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→ Given IP address: 10.15.20.60

Address class:

class A: 10.0.0.0 to 126.255.255.255

class B: 128.0.0.0 to 191.255.255.255.

class C: 192.0.0.0 to 223.255.255.255.

∴ 10 is given. IP address will be  
in class A.

In class A; first octet represents network  
and remaining 3 represent host

⇒ Network IP address: 10.0.0.0

### Direct Broadcast address:

→ In class A; direct broadcast address is obtained by setting all host bits to 1 in host portion of the address.

$$\Rightarrow 10.255.255.255$$

### Limited broadcast address:

→ is always 255.255.255.255.  
regardless of IP address class.

### Key features of UDP:

- Transport layer protocol.
- connectionless.
- provides fast and efficient way to transmit data.
- Doesn't have the overhead of establishing and maintaining connections.
- No flow control.
- supports broadcast & multicast communication.
- simple Header structure.
- Check Sum for error detection.

\*→ Real-time applications:

- ↳ Voice over Internet Protocol
- ↳ Online Gaming
- ↳ Streaming media
- ↳ DNS queries
- ↳ Internet of things.

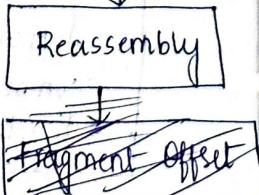
→ IP header - Reassembly of fragments:

↳ IP protocol allows for the fragmentation of packets when the payload size exceeds the Maximum Transmission Unit (MTU) of the underlying network.

\*→ Reassembly process:



- \*→ Fragment offset: (8 byte unit)  
gives the position of packet in original msg.
- \*→ More fragments flag:  
If this is set, additional fragments are expected



→ Time out mechanism is often employed to manage the reassembly process.

\*→ Complete packet is formed when all fragments are received & reassembled.

\*→ Identification field is used to uniquely identify a set of fragments belonging to same original packet

→ FTP: 3-way hand shaking

→ File Transfer Protocol

→ transfer files b/w client & server on TCP-based network

→ SYN packet: Synchronize signal is sent by the client. Packet also includes the client's initial sequence number (ISN).

→ SYN-ACK packet: On receiving the packet, server acknowledges the request by sending a TCP packet with both SYN and ACK flags set to 1.

→ ACK packet: FTP client acknowledges the server's response by sending a TCP packet with the ACK flag set.

\*→ 3 way handshake ensures that both the client and the server agree on initial sequence number & establishes reliable connection before data transfer.

→ FTP uses 2 separate channels for communication:

- ↳ Command channel (Control connection)
- ↳ Data channel (Additional connections)

\*→ Default port numbers of FTP:

- ↳ Data Channel : Port no - 20
- ↳ Control channel : Port no - 21

\*→ Default port number for TELNET : port 23

Telnet uses TCP for reliable connection.

TCP SYN → TCP SYN-ACK → ACK :

\*→ Once TCP connection is established, the user is prompted to log into the remote system.

\*→ After successful connection, user enters an interactive session with the remote host.

\*→ TELNET connection allows the user to execute commands & receive responses from the remote host.

\*→ IP address : 220.100.1.1

↳ to send a packet to all hosts in same network; we use broadcast address.

↳ Specific broadcast address depends on network's subnet mask.

Network address : 220.100.1.0

Broadcast address : 220.100.1.255

⇒ Source IP address : 220.100.1.1

Destination IP address : 220.100.1.255

\*→ Working of following protocols:

→ DHCP: Dynamic Host Configuration Protocol

↳ dynamically assigns IP addresses & other network configuration parameters to devices in a network. diagram

→ SMTP: Simple Mail Transfer Protocol

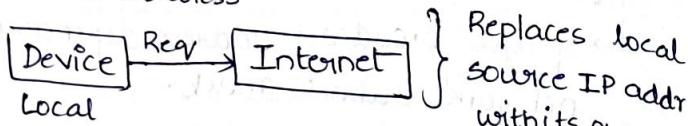
↳ sending emails b/w servers.

Client SMTP email server.

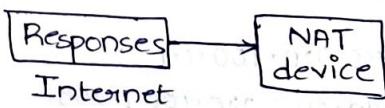
→ Server  $\xrightarrow{\text{SMTP}}$  recipient's email server.  
→ Recipient's email server stores the message until the recipient retrieves it.

### → NAT: Network Address Translation:

Map private IP addresses within a local network to a single public IP address.



NAT device maintains a table to track connections.



### → TCP: Transmission Control Protocol:

Reliable, connection-oriented, ordered delivery  
Establishes connection through 3-way handshake  
Manages flow control, error checking, congestion control.

### → ICMP: Internet Control Message Control:

used for Network error report & diagnostic

→ provides error messages such as ping replies  
→ used for network testing and troubleshooting

### → FTP: File transfer protocol:

→ transferring files b/w client and server.  
→ Opens a separate data connection for actual file transfer.  
→ Supports both active & passive models for data transfer.

### → TELNET:

→ provides remote access to the system  
→ provides terminal emulation over the network.  
→ Allows user to log into remote host & interact directly.

### \* 3 Main divisions of DNS:

→ translates the domain name to IP address  
→ Generic Domains:  
→ registered hosts  
→ 3 character labels; describes organization type.  
→ Ex: aero, biz, com, gov etc.

### Country Domain:

↳ 2 character country abbreviations.

### Inverse Domains:

↳ Mapping address to a name.

## \* Address Resolution Protocol:

→ Maps IP addr to MAC addr.

→ ARP request is a message sent by a device to request the MAC addr. associated with a specific IP addr.

### ARP Request Packet

→ If required mapping is not found in cache; the device sends an ARP request.

↳ Includes: Sender's MAC, IP addresses & target IP & MAC address.

### Broadcasting ARP request:

↳ Sends this request to all devices.

### Receiving and Responding:

↳ Device with IP addr mentioned in ARP request responds to request.

↳ response; ARP reply: contains MAC addr.

### Updating the ARP cache

→ Requesting device updates its ARP cache with newly obtained MAC addr.

→ Subsequent data packets destined for that IP addr are encapsulated in Ethernet frames using correct MAC addr.

## \* Role of Domain Resource Records: (DNS)

→ In DNS, Resource Records are fundamental elements that provide various types of information about domain names, mappings, server info.

→ Each RR has a specific format and is associated with a particular type.

### Commonly used RR:

#### Address Record (A):

→ Maps domain name to IPv4 address.

↳ Ex:

example.com. IN A 192.168.1.1

#### AAAA (IPv6 add) Record:

→ Maps domain name to IPv6 address.

↳ example.com IN AAAA 2001:

→ (MX) Mail Exchange Record:

→ Specifies mail servers responsible for receiving email on behalf of the domain.

→ Ex:

example.com IN MX 10 mail.example.com

→ (CNAME) Canonical Name Record:

→ creates an alias for existing domain.

→ Ex: www.example.com IN CNAME example.com

→ TXT (text) record:

→ holds human-readable text associated with domain.

→ Ex:

example.com IN TXT "This is a text record."

→ NS (Name Server) record:

→ Specifies authoritative DNS servers for domain.

→ Ex:

example.com IN NS ns1.example.com

\*→ WAN link = 2 Mbps

RTT b/w source & destination = 300 msec.

optimal TCP window size

needed to fully utilize the line?

Ans: Bandwidth delay product:

$$= \text{Link bandwidth} \times \text{Round Trip Time (RTT)}$$

$$= 2 \text{ Mbps} \times 0.3 \text{ sec.}$$

↓

$$2 \text{ Mbps} = \frac{2 \text{ Mbps}}{8} = 0.25 \text{ MBps}$$

$$\therefore \text{BDP} = 0.075 \text{ MB}$$

↓

convert to bytes

$$= 0.075 \text{ MB} \times 1024 \text{ KB/MB} \times 1024 \text{ bytes/KB}$$

$$\text{TCP window size} = 78,643 \text{ bytes}$$

\*→ Transport layer protocol - transmitting data

using FTP/TELNET:

→ Transport layer protocols: TCP/UDP.

→ When transmitting data using FTP/TELNET, TCP is primary protocol employed.

### \* FTP: Role of TCP:

- Reliable and ordered delivery.
- Establishes connection through 3-way handshake before data transfer.
- While data transfer process ; commands & responses are exchanged over the control channel ; while actual file data is transferred over data channel using TCP.

### \* TELNET: Role of TCP:

- Reliable; bidirectional communication b/w client & server.
- Ensures commands and responses are delivered in correct order & handles retransmission of lost or corrupted data.
- User IPs and OIDs are transmitted over the TCP connection in text-based format.

\* Given subnet mask: 255.255.255.124

Class C Subnet mask: 255.255.255.0.

→ Hence ; given subnet mask falls b/w 2 class C Subnets

→ Binary form:

1111111.1111111.1111111.0111100

In this 5 bits are for host addresses.

No. of hosts per subnet =  $2^n - 2$   
 $n$  = no. of bits for host addresses

$$\Rightarrow 2^5 - 2 = 30$$

⇒ Each subnet can have 30 hosts.

⇒ No. of bits borrowed for subnetting = 3 bits.

$$2^3 = 8 \text{ Subnets}$$

→ Leaky Bucket Algorithm:

→ flow control mechanism often used in networking to control the rate at which data is transmitted.

→ Primary purpose:

→ smooth outbursts of traffic and network operates at a steady rate.

## \* Leaky Bucket Algorithm

- Incoming data is poured into a bucket with a leak.
- bucket has a fixed capacity; if incoming data is too high; it spills out.
- prevents bursty traffic and policing to control the rate of data transfer.
- prevents network congestion.
- Quality of service in TCP typically involves factors like latency, packet loss.

\* windows send size = 65538 B.

1 Gbps channel

10 ms. one way delay.

max. throughput = ?

$BDP = \text{Link Bandwidth} * \text{Round trip time}$

$\text{Throughput} = \frac{\text{TCP window size}}{\text{RTT}}$

window size =  $65538 \times 8 = 524304 \text{ bits}$

$$RTT: 1 \text{ Gbps} = 10^9 \text{ bits per sec}$$

$$BDP = 10^9 \text{ bps} \times (20 \text{ msec}) = 20 \times 10^6 \text{ bits}$$

$$\text{Throughput} = \frac{524304}{20 \times 10^6} = 0.026/\text{sec}$$

$$= 38,044,948 \text{ bps.}$$

$$= 38.04 \text{ Mbps.}$$

\* CIDR block: 245.248.128.0/20.

→ subnet size =  $2^{12}$  addresses.

→ Divide the CIDR block into 2 parts for organisation A and organisation B.

→ Allocate addresses:

→ Half of the addresses =  $2^8 = 2048$  addresses.

Subnet: 245.248.128.0/21

Range: ~~245.248.128.0 to 245.248.135.255~~

245.248.128.0 to 245.248.135.255

→ Quarter of the addresses =  $2^10 = 1024$  addresses.

Subnet: 245.248.136.0/22

Range: 245.248.136.0 to 245.248.137.255

ISP:  
 Subnet: 245.248.140.0/21  
 Range: 245.248.140.0 to  
 245.248.147.255.

\* Header fields in TCP that contribute to reliable and ordered data.

- Sequence Number
- Acknowledgement Number
- TCP flags
- Window size
- Checksum

\* Minimum Window size: channel utilization  $> 80\%$ .

$$\text{Throughput} = \frac{\text{Window size}}{\text{RTT}} \times \text{Link BW}$$

$$U(\text{channel utilization}) = \frac{\text{Throughput}}{\text{Link BW}}$$

$$\begin{aligned} \text{Link BW} &= 1 \times 10^9 \text{ bps} \\ \text{RTT} &= 40 \text{ msec} = 0.04 \text{ sec} \\ \text{Packet size} &= 1000 \text{ bytes} \\ U &> 80\% \end{aligned}$$

$$\text{Throughput} = U \times \text{Link BW}$$

$$\text{window size} = \frac{\text{Throughput}}{\frac{1}{\text{RTT}}}$$

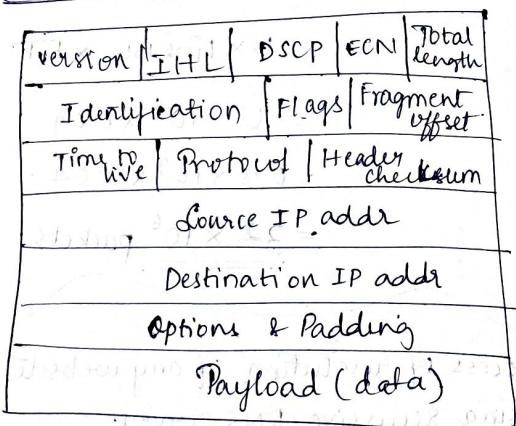
$$\begin{aligned} U \times \text{Link BW} \times \text{RTT} &= U \times 10^9 \times 0.04 \\ &= 0.8 \times 10^9 \times 0.04 \\ &= 32 \times 10^6 \text{ packets} \end{aligned}$$

\* Process of resolution of any website using recursive DNS server.

- Understand ethical hacking.
- Recursive resolution is a process where the DNS server will continue to search for an answer in lower-level domains.

- Recursive DNS server receives the DNS query from the user's device
- Cache check; to determine if it has the mapping of the requested domain name to an IP address.

### \* IPv4 format:



- Version: (4 bits)
  - ↳ Version of protocol (IPv4 / IPv6)
- Internet Header length (IHL): (4 bits)
  - ↳ 32 bit words
  - ↳ starting point of the data

- Total length (16 bits)
  - ↳ Header + Payload
  - ↳ Assists routers in processing & forwarding packets.
- Protocol: (8 bits)
  - ↳ protocol used
  - ↳ enables routers to ~~determine~~ determine how to process payload.

### \* SMTP: Sending and Receiving emails:

#### → Sending Email:

- User composes Email
- Email client initiates connections
- SMTP handshake
- Sender authentication
- Email transmission
- SMTP server routing
- SMTP relay
- Email delivery

#### → Receiving Email:

- Email retrieval
- Authentication
- Email download
- Display & Management

## Characteristics:

- Text-based protocol
- Reliability
- Connection-oriented

\* IP addr: 10.0.0.0/16.

500 hosts per subnet.

$$2^x \geq 500 \Rightarrow x = 9.$$

9 bits for host.

$$\rightarrow \text{No. of bits for subnet} = 32 - \text{No. of bits for hosts}$$

$$= 32 - 9 = 23$$

→ subnet mask for the new subnets will be 1/23.

$$\rightarrow \text{subnet range} = \text{Network Addr.} + \text{Incr} \times \text{Subnet No.}$$

$$\text{Incr} = 2^9 = 512$$

## \* Subnet 1

$$\hookrightarrow \text{NA} = 10.0.0.0$$

Range = 10.0.0.1 to 10.0.1.254  
 $\text{BA} = 10.0.1.255$

## \* Subnet 2

$$\hookrightarrow \text{NA} = 10.0.2.0$$

Range = 10.0.2.1 to 10.0.3.254  
 $\text{BA} = 10.0.3.255$

## \* Key Components of DNS:

- Server, Resolver.
- Record.
- Root DNS Servers.
- Cache.
- Top-level domain servers.

## \* Process of domain Name resolution:

- User initiates request
- DNS query sent to resolver
- Resolver checks cache
- Recursive DNS query
- Root DNS servers
- Response sent to resolver
- Response sent to user

### \* Significance :

- Human-readable names
- Distributed network
- Navigation
- Email delivery
- Dynamic IP address assignment
- Global accessibility

### \* Link State Routing Algorithm:

Optimal path for data transmission b/w nodes.

#### Steps

→ Router creates a Link State Advertisement (LSA) containing info about its directly connected links and states.

→ Router floods its LSA to all other routers in the network.

Each router collects LSAs from all other routers in the network and builds its link state database.

→ Using LSDB ~~represent~~ each router runs Dijkstra's algorithm to calculate the shortest path to every other router in network.

→ Shortest path tree is used to determine the next-hop router for each destination in the network.

→ Routing Table.