VASAVI COLLEGE OF ENGINEERING (Affiliated to Osmania University)

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9	PRELAB QUESTIONS-5 HOW Goreedy approach did			le &
0	Divide and Conques	Goreed	y Algorithm	Λ ,
>	Finds the sol? but does not aim to find optimal sol? divides the problem into small sub problems and each problem is solved independently. Recursive in nature and, slower and inefficient	optimal feasible I this I feasible shall described	sol from sol solved to le sol the	get In We optimal
2)	Write the control abstrapproach:			A
	Algorithm Goveredy (a, f Solution: = 0	n) lo		

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	{ æ:=select(a);	55
	is feasible (solution, a) then	
	solution: = union (solution, a)	, Ze
3)	Differentiate between feasible sol & optin	nal
>	feasible sol satisfies all constraints	
)	optimal soin which is a feasible soin which optimizes objective fn.	ch
)	What is the time complexity of Prim's	
	algorithm?	
¥.	Time complexity of Prim's alanyithm in	
	Time complexity of Prim's algorithm in $O((E+V) \log V)$ where $E \rightarrow no \cdot ob edges$,	
	V -> no. of vertices.	
	Woute the significance of union and find operation in Kruskal's algorithm.	

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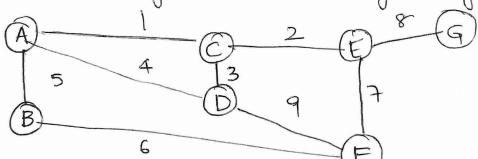
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It merges a different subsets into a single subset and the suppresentative of one set becomes suppresentative of another. The disjoint-set also supports one other important operation called Make Set, which creates a set containing only a given element in it.

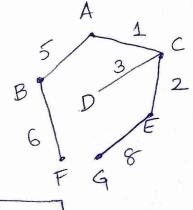
Fund minimum spanning tree by using Prim's and Kruskal's algorithms for the given graph.



* Kruskal's Algorithm

$$\langle A,C \rangle = 1$$

 $\langle C,E \rangle = 2$ $\langle F,E \rangle = 7$
 $\langle C,D \rangle = 3$ $\langle E,G \rangle = 8$
 $\langle A,D \rangle = 4$ $\langle D,F \rangle = 9$



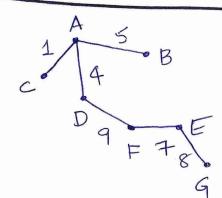
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Cost = 34

PRELAB PROGRAMS - 5:

Implement Kruskal's algorithm using C:

#include < Stdio.h >

#includexstalib.h>

int comparator (const void * p1, const void * p2)

{ const int (*x)[3] = p1;const int (*y)[3] = p2;

netwin (*x*)[2] - (* y)[2],}

void makeset (int parent[], int rank[], int n)

{ for (int i=0; kn; i++)

{ pagent[i]=i; rank[i]=0; }}

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int findparent (int parent[], int component)
 { is (parent [component] == component)
          return component;
     return parent [component] = findParent (parent)
                                       parent (component);
Void unionSet (int u, int v, int parent[],
                  int rank[], int n)
{ u = findParent (parent, u);
   V = find Parient (parient, v);
   ib (rank[u] < rank[v]) { parent[u]=v;}
   else 'ib (rank[u] > rank[v]) { parent[v] = u°, }
   else { parent[v]=u; rank[u]++;}}
void Kruskal(int n, int edge[n][3])
{ quort (edge, n, sizeof (edge(o)), comparator);
  int parent[n], rank[n];
   makeSet (parent, rank, n);
```

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int minCost = 0;

print ("Following are the edges in the spanning tree: \n");

for (int i = 0; i < n; i++)

{ int v1 = find Parent (parent, edge [i][i]);

int v2 = find Parent (parent, edge [i][i]);

int wt = edge[i][2]; if (VI! = V2)

f unionSet (VI, V2, parent, rank, n); minCost += Wt;

mincost+=wt)
printf("%d--%d==%d\n", edge[i][o],

edge[i][i], wt);}

printf ("Minimum Cost Spanning Tree: %d/n")
min Cost); }

int main() { int edge [5][3] = { {0,1,10}, {0,2,6}, {0,3,5}, {1,3,15}, {2,3,4}};

Kruskal (s,edge);

Following are the edges in the 1000 spanning

$$2 - - 3 = = 4$$
 $0 - - 3 = = 5$

Minimum Cost Spanning Tree: 19.

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[i]

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2) Implement Prim's algorithm in C.

#include x limits.h>

include Kstd bool.h>

#include (Stdio.h> # define V 5

int minkey (int key[], bool, mstSet[])

{ int min = INT_MAX, min_index;

for (int v=0; v < V; v++)

{ ib (mtset[v] == false && key[v] < min)

{ min = key[v], min_index = v ; }}

return min-index; }

int printMsT (int parent [], int graph[v][v])

{ printf ("Edge It Weight In");

for (int i=1; ix V; i++)

{ printf ("%d-%d\t %d\n", parent[i], i,

graph[i][parent(i]]);}

void primMST (int graph[V][V])

int parent[V]; int key[V]; book mstSet[M];

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NAME OF THE LABORATORY : DAA Name K.S.T. SIVANI Roll No. -052 for(int i=0:1xV:i++) { key[i] = INT_MAX, mstSet[i] = false; 'y key[0] = 0; parent[0] = -1; for (int count = 0; count < V-1; count ++) { int u=minKey(key, mstSet); mstSet[u] = true; for (int v = 0; v < V; v++) { ib (graph[u][v] & 2 mst Set[v] == false & 2 graph[u][v] key[v]) { parent[v] = u's key[v] = graph[u][v]; } print MST (parient, graph); } int main () { int graph [v][V] = {{0,2,0,6,0}, {2,0,3,8,5}, {0,3,0,0,4}, \$6,8,0,0,9}, 90,5,7,9,033; primmst (graph); gretwino; }

Edge Weight
0-1 2
1-2 3
0-3 6