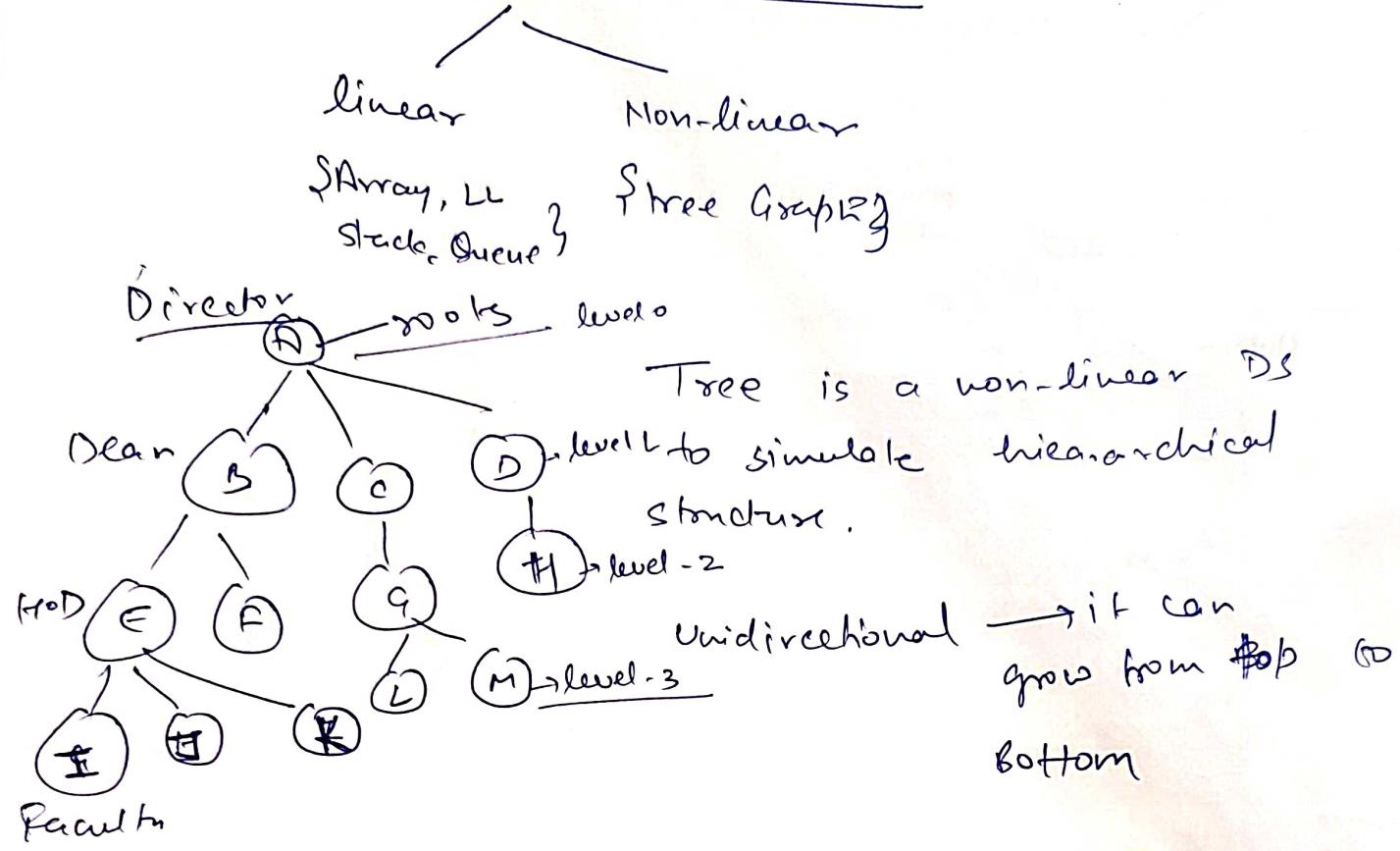


## Introduction to Trees

0



Nodes - | stores info + link to other node

“Tree is a non-linear DS, defined as collection of nodes linked together to simulate a hierarchy.”

Root → node with no parent (A)

order → A, B, C, D, E, F, G, H, I, J, K, L, M

Pastent → A, B, C, D, E, F, G

Rebuild  $\Rightarrow$  L, M are children of G  
child / children

Leaf nodes  $\rightarrow$  nodes with no child/children  
Ex I, J, K, L, M (External node)

Path :- Sequence of consecutive edges from source node to destination node.

(2)

Ancestor

any predecessor node on the path from root to that node.

Exancestor of L  $\rightarrow$  A, C, Gancestor of H  $\rightarrow$  D, ADecendant  $\rightarrow$ 

Any successor node on the path from that node to ~~root~~ leaf node.

Exdecendant of C  $\rightarrow$  G, E, Mdecendant of B  $\rightarrow$  E, F, I, J, KDegree  $\rightarrow$ 

Max. degree of tree = Max. no. of children a node can have in that tree.

degree of tree  $\rightarrow$  Max. degree among all nodes.

Depth of node :-

length of path from root to that node

Ex

Depth of J = 4      Depth of root node = 0

height of node  $\Rightarrow$ 

no. of edges in the longest path from that node to leaf.

Ex

height of A  $\Rightarrow$  4

height of G  $\Rightarrow$  1

height of B  $\Rightarrow$  2

Note:- height of tree = height of root node

Siblings

$\rightarrow$  children of same parent

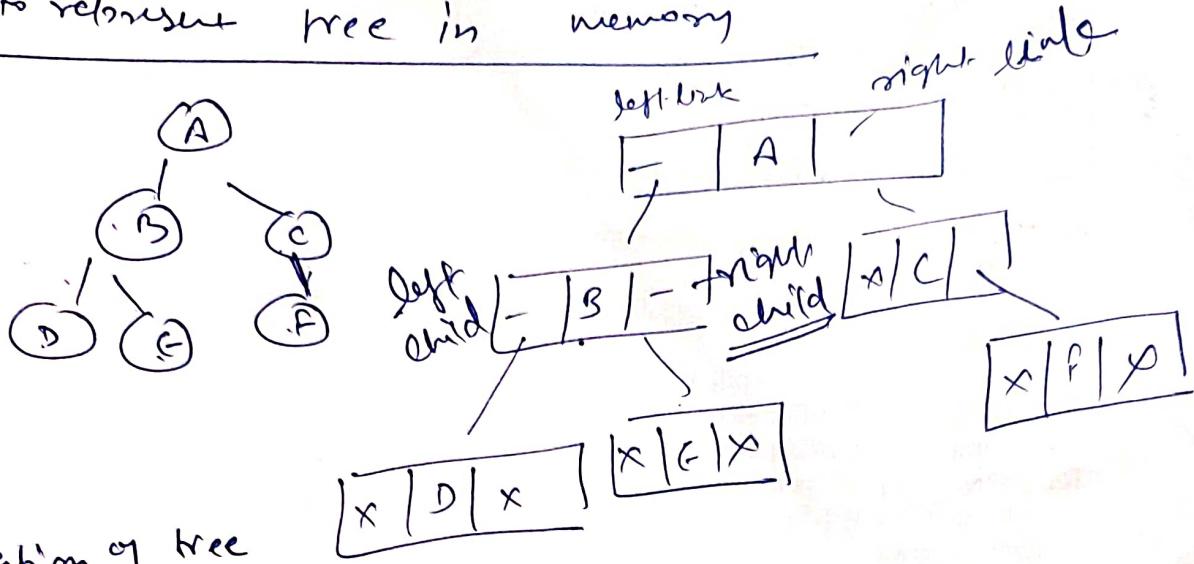
B, C, D are siblings

level of node :-

the tree given has total 4 levels starting from 0 to 3.

(3) In any tree if there are  $n$  nodes then  
no. of edges =  $(n-1)$

How to represent tree in memory



Representation of tree node in C

Struct node

```
of
struct data;
```

```
struct node *left;
struct node *right;
```

```
}
```

Applications of Tree Data

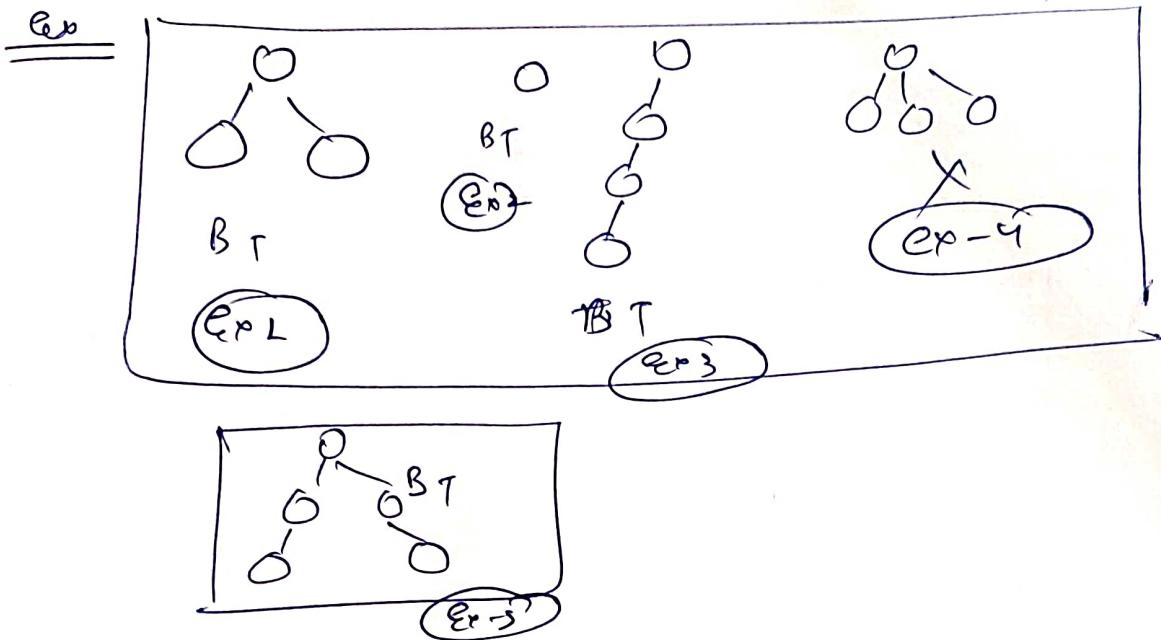
File System

Routing Protocols  
Quick Searcher.

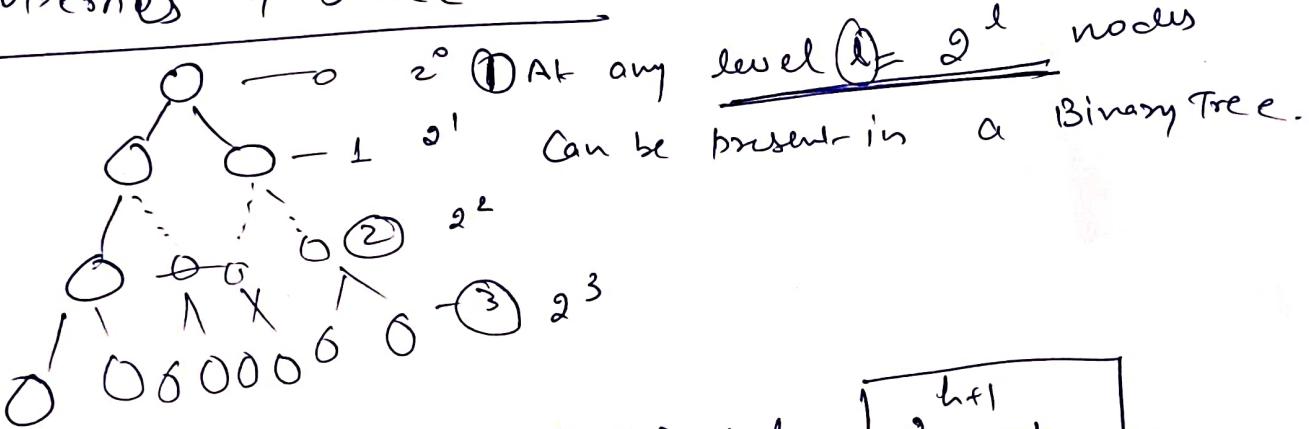
## (2) 2

Binary Tree

Each node can have atmost 2 children (either 1 or 2 or no child)



## Properties of B-Tree



(2) max. no. of nodes of height  $h = 2^{h+1} - 1$

Proof: at level 0  $\rightarrow 0$

at level 1  $\rightarrow 2$

at level 2  $\rightarrow 4$

$$\begin{aligned} \text{Total nodes} &= 2^0 + 2^1 + 2^2 + \dots + 2^h \\ &= (2^{h+1} - 1) \end{aligned}$$

(3) min. no. of nodes of ~~the~~ at height  $h = (h+1)$

Proof: total nodes =  $(h+1)$

height 1 =  $\frac{2}{3}$  nodes

height 2 =  $\frac{4}{3}$  nodes

height 3 = 4 nodes

(5)

- (4) Total no. of leaf nodes in binary Tree  
 $=$  total no. of degree 2 nodes + 1

- (5) Max. no. of binary Tree with  $n$  nodes  
 $= \frac{2^n C_n}{n+1} \Rightarrow$

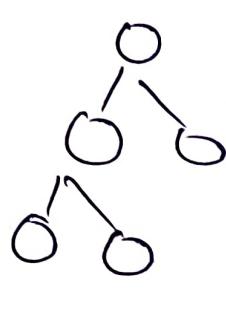
$${}^n C_r = \frac{n!}{r!(n-r)!}$$

- (6) Height of a Binary Tree = Total levels in BT - 1  
 $\boxed{h = L - 1}$

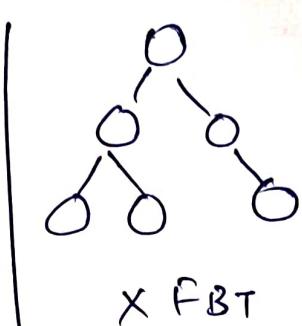
(6)

## Types of Binary Tree

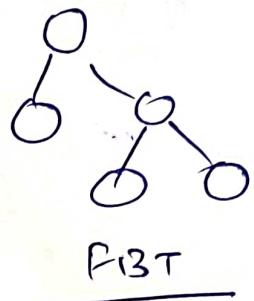
Full Binary Tree/strict BT with either 0 or 2 children.

Ex:-

FBT

o  
FBT

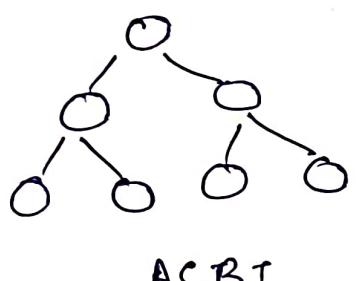
X FBT



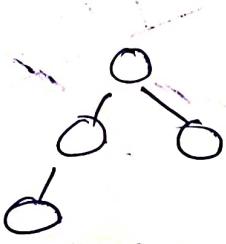
FBT

## Almost Complete Binary Tree

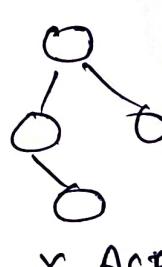
- A BT which always fill from left to right
- only after filling one level you will go to next level



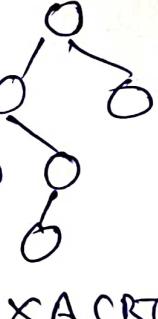
ACBT



ACBT



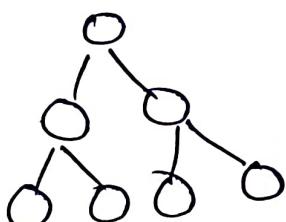
X ACBT



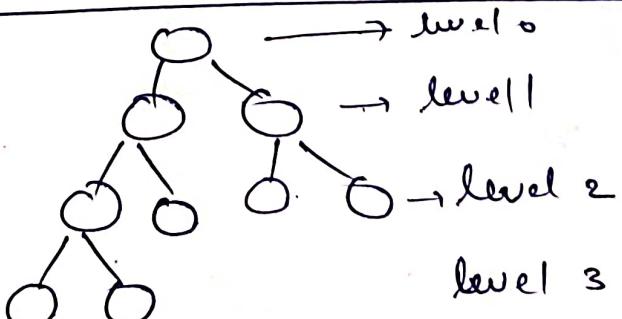
X ACBT

## Complete Binary Tree

A BT where internal nodes always have 2 children except leaves, and no holes are allowed



complete BT tree

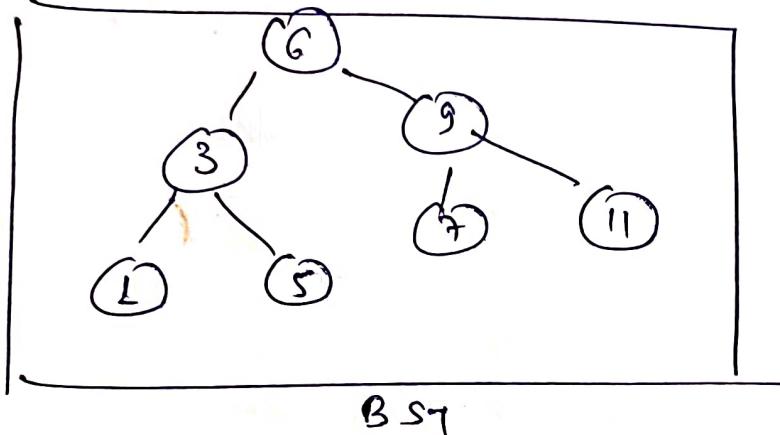


level 3

\* complete BT tree (No holes allowed)

## Binary Search Tree

- \* Is a special kind of Binary Tree, which overcome the drawback of Binary Tree where nodes are inserted with no order.
- \* In BST, a node's key is greater than or equal to its left child's key but less than or equal to right child's key.

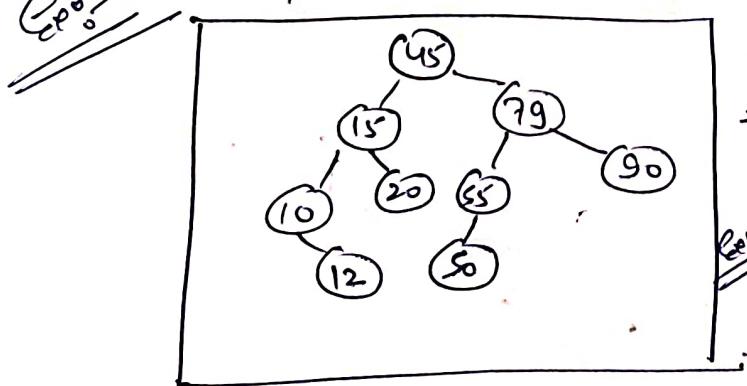


- \* BST inorder will always give an sorted array  
in above tree inorder

[ 1, 3, 5, 6, 7, 9, 11 ]

- \* How to create a BST, given an array

45, 15, 79, 90, 10, 55, 12, 20, 50



Q. Generate a BST by inserting following keys

~~45, 15, 79, 90, 10, 55, 12, 20, 50~~

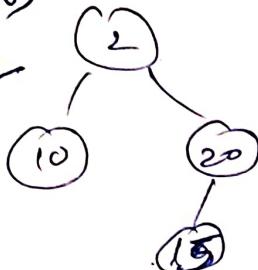
Process:- Select first value as node, & than add every key value in the left if smaller or right if greater than root.

## Deletion in BST

Deletion of leaf node

Deletion of Non-leaf node

Ex:



delete 15

(No change in BST after Deletion of leaf node)

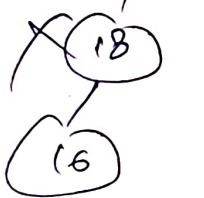
non-leaf node with one child

non-leaf child with 2 children

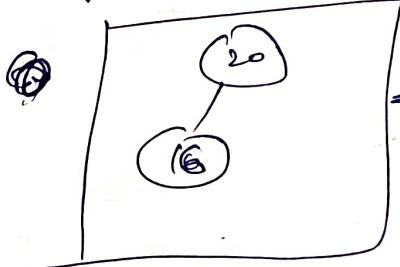
Removing Non-leaf node with one child

Process: Join the leaf node with grand Parent node

Ex:

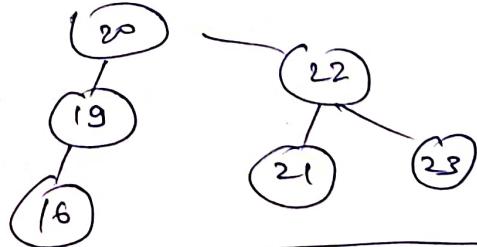
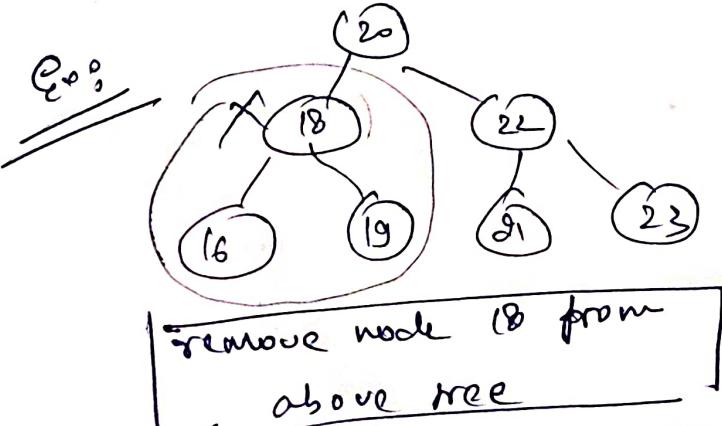


Solution



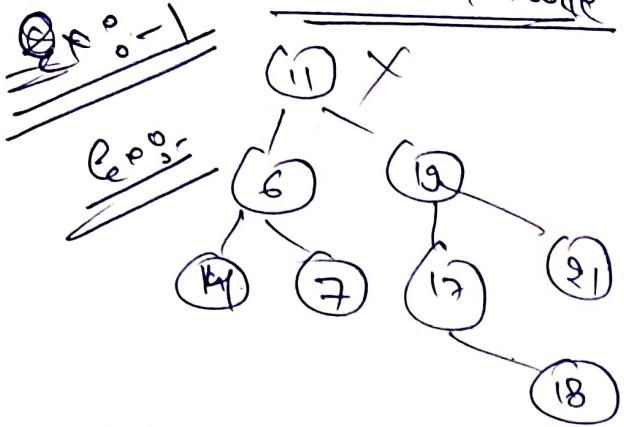
Removing Non-leaf node with 2 children

Method 1: Pick the largest element from the left subtree of the node which you want to delete

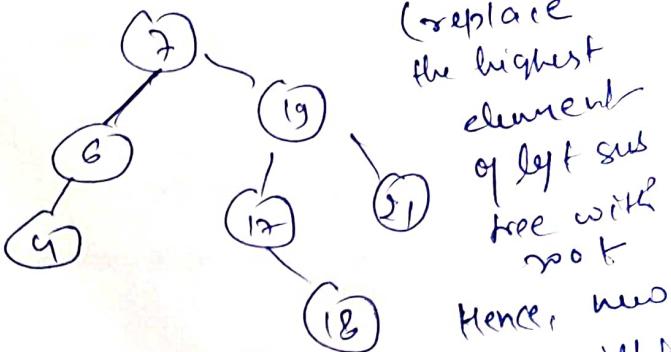


tree after removal of 18

Remove root node



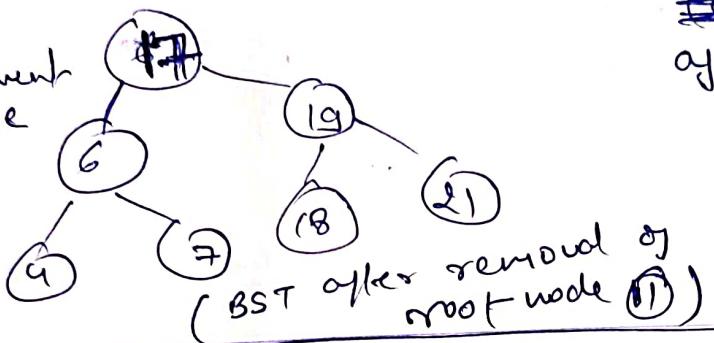
from left subtree (Method I)



(replace the highest element of left sub tree with root)  
Hence, new root will be 7.

Method-II

Replace the smallest element of right subtree with node that is being deleted

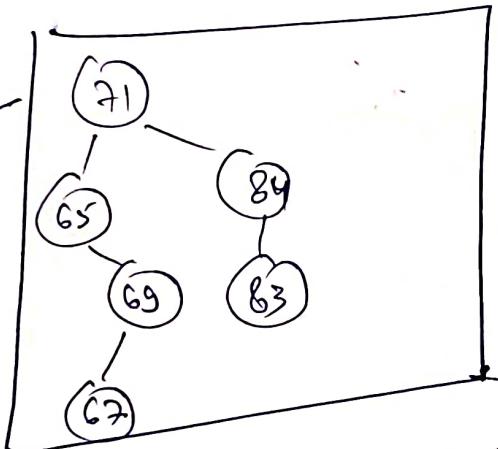


~~bst~~ (BST after removal of 11)

(Insert the given keys to a BST)

71, 65, 84, 69, 67, 83

Solution

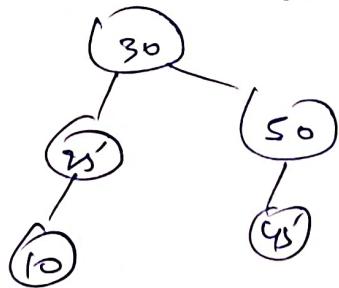
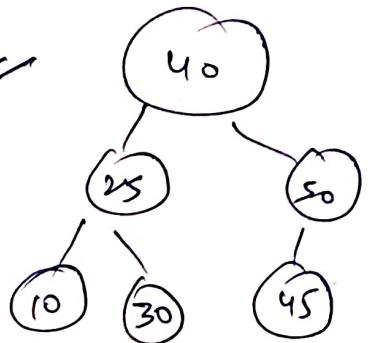


Method - 2

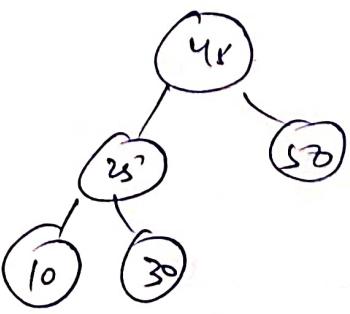
removing 40, and

~~left subtree~~ replacing it with largest node of left subtree i.e. 30

Q 10



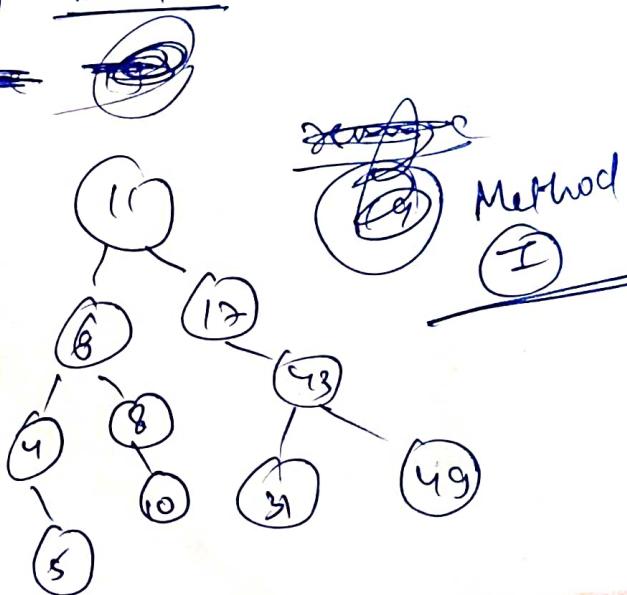
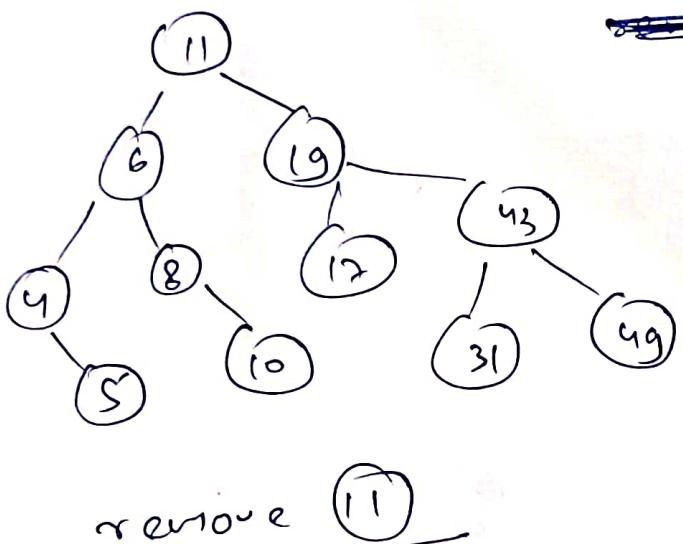
(Remove root node 40 from above BST)



~~replacing~~ ~~left subtree~~

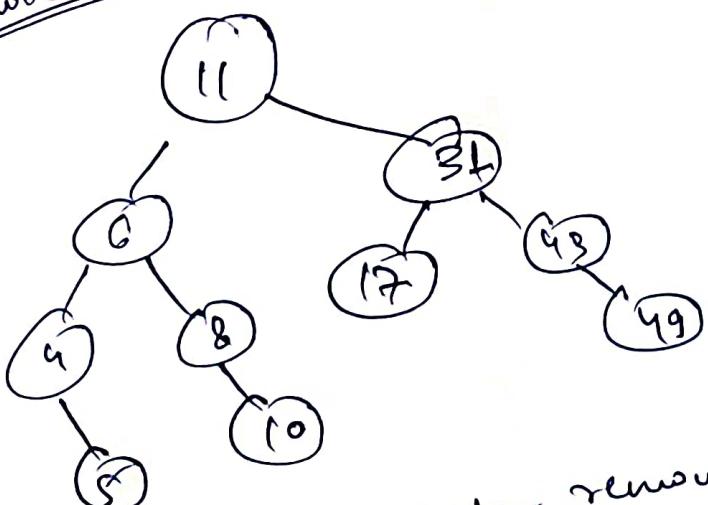
removing 40, and replacing it with smallest node of right subtree i.e., 45

Remove (19) from the following BST.



remove 11

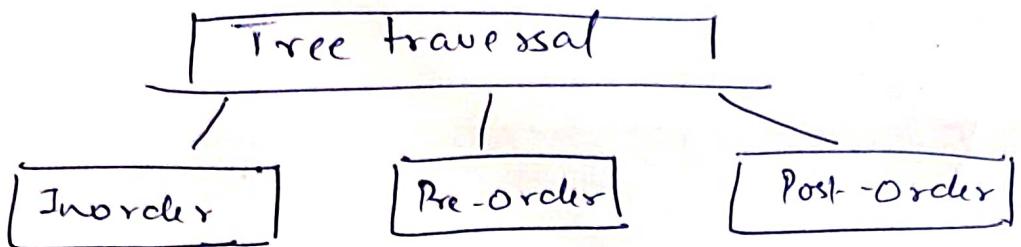
Method 2



Solution

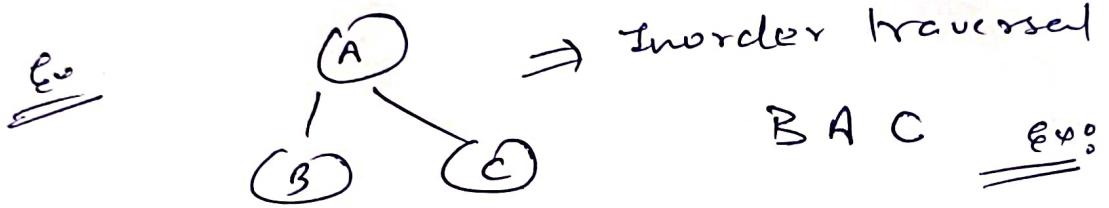
(BST after removal of  
~~node 19~~ by  
replacing it with smallest  
node of right subtree  
node 3).

## BINARY TREE TRAVERSAL



### INORDER TRAVERSAL

LEFT - ROOT - RIGHT



(Here we recursively visit the left subtree, then root & then explore the right subtree)  
used to create a sorted list

### PRE-ORDER TRAVERSAL

ROOT - LEFT - RIGHT

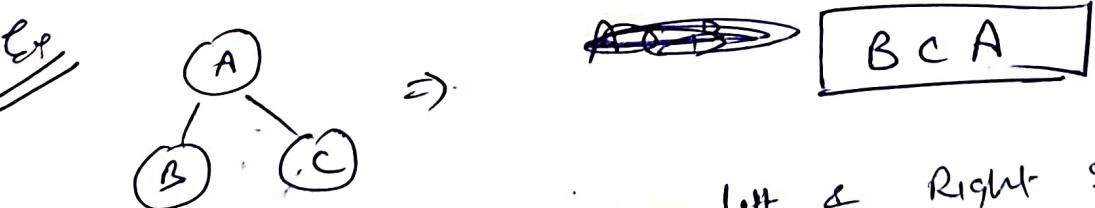


[the order of traversed

### POST-ORDER TRAVERSAL

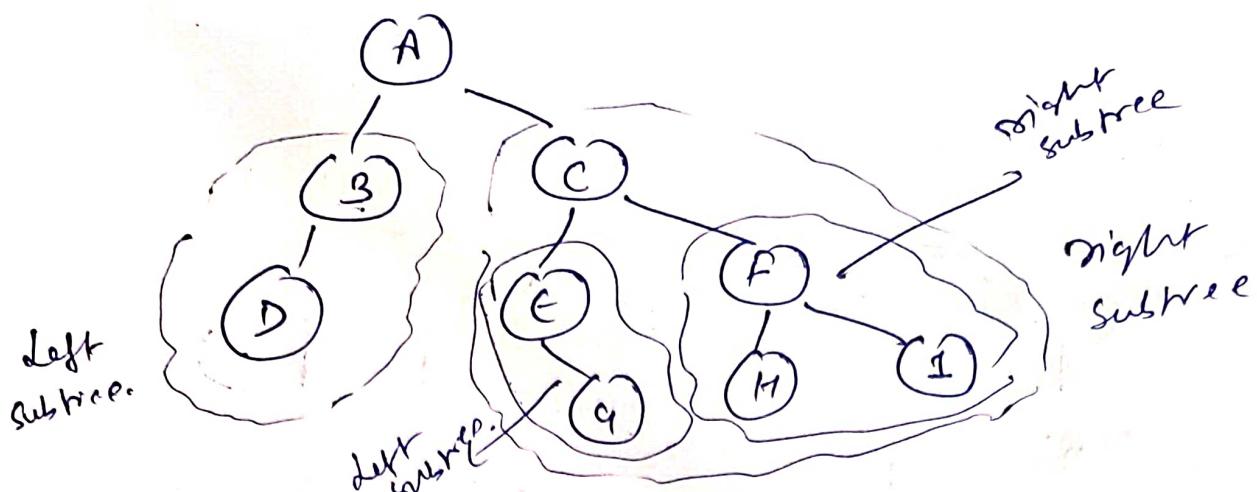
~~ROOT - RIGHT - LEFT~~

~~LEFT - RIGHT - ROOT~~ [Order of traversal]



### POST ORDER

visits the left & right subtree first,  
then ends up at root.  
This is often used to delete trees.



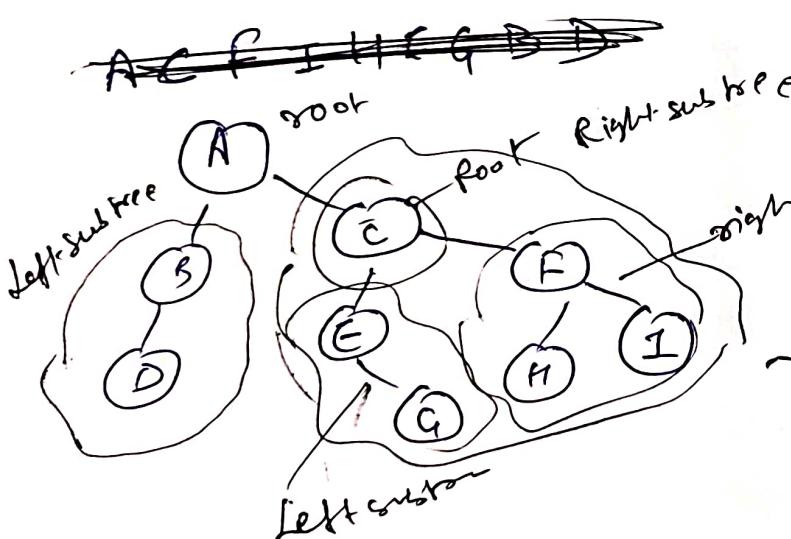
IN-ORDER (Left - Root - Right)

[DB A E G C H F I]

PRE-ORDER (Root - Left - Right)

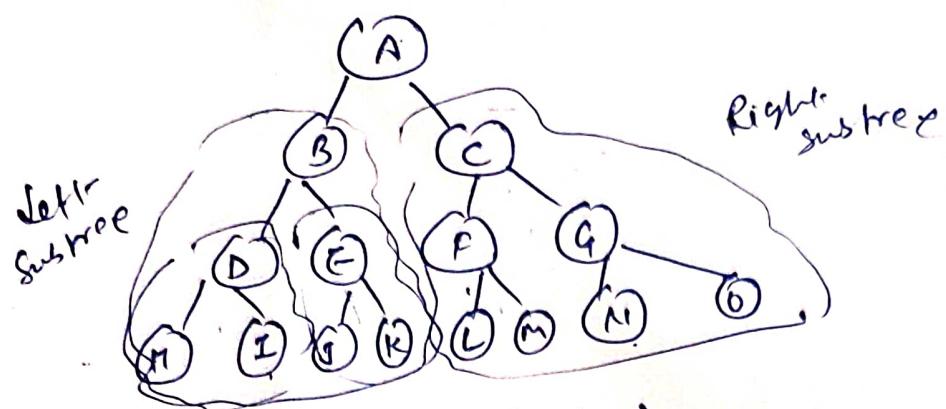
[A B D C E G F H I]

POST-ORDER (Root - Right - Left) ~~(Root - Right - Left)~~ left right Root



[DB G E H I F C A]

Ex-2  
Find out the Pre-order, Postorder, Inorder traversals of Binary Tree.



INORDER: (Left - Root - Right)

H D I B J E K A L F . M C N G O

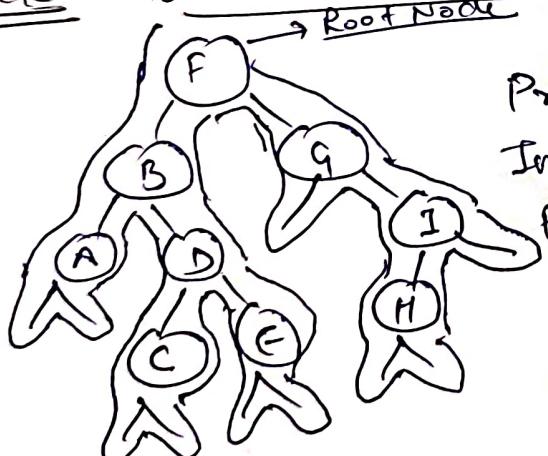
PREORDER: (Root - Left - Right)

A B D H I E J K C F L M G N O

POSTORDER: (Left - Right - Root)

H I D J K E B L M F N O G C A

TRICK to traverse a Binary Tree  
→ Root Node



Pre-Order - F B A D C E G I H

In-Order - A B C D E F G H I

Post-Order - A C E D B H I G F

Trick :- if node is visited      once put it in Pre-order  
              "                          twice                          " - In-order  
              "                          thrice                          " - Post-order