

Q1, Solve using Master's Theorem:-

$$1) T(n) = 3T(n/2) + n^2$$

$$T(n) \Rightarrow aT(n/b) + f(n) \text{ where } a=3, b=2$$

On comparing, $a=3, b=2, f(n)=n^2$

$$\text{Now, } c = \log_b a \text{ i.e., } \log_2 3 \text{ i.e., } 1.584$$

$$\text{and } n^c \text{ i.e., } n^{1.584} < n^2$$

$$\therefore f(n) > n^c \therefore \boxed{T(n) = O(n^2)}$$

$$2) T(n) = 4T(n/2) + n^2$$

$$a \geq 1, b > 1 \text{ and } a < 4, b = 2, f(n) = n^2$$

$$c = \log_2 4 \text{ i.e., } 2$$

$$n^c = n^2 \Rightarrow f(n) = n^2$$

$$\therefore \boxed{T(n) = O(n^2 \log n)}$$

$$3) T(n) = T(n/2) + 2^n$$

$$a=1, b=2, f(n)=2^n$$

$$c = \log_b a \Rightarrow \log_2 1 \text{ i.e., } 0$$

$$\text{Now, } n^c \text{ i.e., } n^0 \Rightarrow 1 \text{ and } f(n) > n^c$$

$$\boxed{T(n) = O(2^n)}$$

$$4) T(n) = 2^h T(n/2) + n^k$$

$$a = 2^h, h = 2, f(n) = n^k$$

$$c = \log_2 2^h \Rightarrow h \text{ and } n^c = n^h$$

$$f(n) = n^c$$

$$T(n) = O(n^2 \log n)$$

$$Q5, T(n) = 16T(n/4) + n$$

$$a = 16, h = 4, f(n) = n$$

$$c = \log_4 16 \Rightarrow \log_4 (4)^2 \text{ i.e., } 2 \log_4 4 \text{ i.e., } 2$$

$$\text{Now, } n^c = n^2 \text{ and } f(n) < n^c$$

$$\therefore T(n) = O(n^2)$$

$$Q6, T(n) = 2T(n/2) + n \log n$$

$$a = 2, h = 2, f(n) = n \log n$$

$$c = \log_2 2 \text{ i.e., } 1 \text{ and } n^c \Rightarrow n' \Rightarrow n$$

$$\text{Since, } n \log n > n \text{ i.e., } f(n) > n^c$$

$$\text{So, } T(n) = O(n \log n)$$

$$Q7, T(n) = 2T(n/2) + \frac{n}{\log n}$$

$$\text{So } a = 2, h = 2, f(n) = \frac{n}{\log n}$$

$$c = \log_2 2 \text{ i.e., } 1 \text{ and } n^c \Rightarrow n'$$

Since $\frac{n}{\log n} < n$ $\therefore f(n) < n^c$

$\therefore \boxed{T(n) = \theta(n)}$

8) $T(n) = 2T(n/4) + n^{0.51}$

$a=2, b=4, f(n) = n^{0.51}$

$c = \log_b a$ i.e. $\log_4 2$ i.e. 0.5

as $n^c = n^{0.5}$

So, $n^{0.5} < n^{0.51}$ i.e. $f(n) > n^c$

$\therefore \boxed{T(n) = \theta(n^{0.51})}$

9) $T(n) = 0.5 T(n/2) + 1/n$

$a=0.5, b=2, f(n) = 1/n$

Since, according to Master's theorem, $a \geq 1$ but here a is 0.5 so we can't apply Master's theorem.

10) $T(n) = 16T(n/4) + n!$

$a=16, b=4, f(n) = n!$

$c = \log_4 16$ i.e. 2

Now, $n^c = n^2$

as $n! > n^2$

$\therefore \boxed{T(n) = \theta(n!)}$

$$11) T(n/2) + \log n$$

$$a=4, b=2, f(n) = \log n$$

$$c = \log_2 4 \text{ i.e., } 2$$

$$\text{Now, } n^c = n^2$$

$$\text{Since } \log n < n^2$$

$$\therefore f(n) < n^c$$

$$\therefore \boxed{f(n) = O(n^2)}$$

$$(12) T(n) = \sqrt{n} T(n/2) + \log n$$

$$a = \sqrt{n}, b=2, f(n) = \log n$$

$$\text{Now, } c = \log_2 \sqrt{n} \text{ i.e., } \frac{1}{2} \log_2 n$$

$$\text{Now, } \frac{1}{2} \log_2 n < n \log n \quad \therefore f(n) < n^c$$

$$\text{So, } T(n) = O(f(n)) \text{ i.e., } O(\log(n))$$

$$(13) T(n) = 3T(n/2) + n$$

$$a=3, b=2, f(n) = n$$

Now,

$$c = \log_2 3 \text{ i.e., } 1.5849$$

$$\text{Now, } n^c = n^{1.5849}$$

$$\text{So, } n < n^{1.5849}$$

$$\text{i.e., } f(n) < n^c$$

$$\therefore \boxed{T(n) = O(n^{1.5849})}$$

~~$$T(n) = 3T(n/2) + n$$~~

~~$$a=3, b=2, f$$~~

$$14) T(n) = 3T(n/3) + \sqrt{n}$$

$$a=3, b=3, f(n) = \sqrt{n}$$

$$\text{Now, } c = \log_3 3 \text{ i.e., } 1 \quad \therefore n^c = n^1$$

$$\text{as } \sqrt{n} < n \quad \text{so, } f(n) < n$$

$$\therefore T(n) = \theta(n)$$

$$15) T(n) = 4T(n/2) + cn$$

$$a=4, b=2, f(n) = cn$$

$$\text{Now, } c = \log_2 4 \text{ i.e., } 2 \quad \text{so, } n^c = n^2$$

$$\therefore cn < n^2 \text{ (for any constant)}$$

$$\therefore f(n) < n^c$$

$$\therefore \boxed{T(n) = \theta(n^2)}$$

$$16) T(n) = 3T(n/4) + n \log n$$

$$a=3, b=4, f(n) = n \log n$$

$$\text{Now, } c = \log_4 3 \text{ i.e., } 0.792, \quad n^c \Rightarrow n^{0.792}$$

$$\text{Now, } n^{0.792} < n \log n$$

$$\therefore \boxed{T(n) = \theta(n \log n)}$$

$$17, T(n) = 6T(n/3) + n^2 \log n$$

$$a=6, b=3, f(n) = n^2 \log n$$

$$\text{Now, } c = \log_3 6 \Rightarrow 1.63903$$

$$\text{Now, } n^c \Rightarrow n^{1.63}$$

$$\text{So, } n^{1.63} < n^2 \log n$$

$$\text{So, } \boxed{T(n) = \theta(n^2 \log n)}$$

$$18, T(n) = 3T(n/3) + n/2$$

$$a=3, b=3, f(n) = n/2$$

$$\text{So, } c = \log_3 3 \text{ i.e., } 1,$$

$$\text{Now, } n^c \Rightarrow n^1$$

$$\therefore n^c > n/2 \text{ i.e., } f(n) < n^c$$

$$\text{So, } T(n) = \theta(n)$$

$$19, T(n) = 4T(n/2) + \frac{n}{\log n}$$

$$\rightarrow a=4, b=2, f(n) = \frac{n}{\log n}$$

$$c = \log_2 a = \log_2 4 = 2$$

$$n^c = n^2$$

$$\frac{n}{\log n} < n^2$$

$$\boxed{T(n) = \theta(n^2)}$$

$$1. T(n) = 64T(n/8) - n^2 \log n$$

$$a=64 \quad b=8$$

$$c = \log_b a = \log_8 64 = \log_8 (8)^2$$

$$c=2$$

$$n^c = n^2$$

$$\therefore n^2 \log n > n^2$$

$$\boxed{T(n) = \Theta(n^2 \log n)}$$

$$21. T(n) = 7T(n/3) + n^2$$

$$a=7, b=3, f(n) = n^2$$

$$c = \log_b a = \log_3 7 = 1.7712$$

$$n^2 < n^{1.7712}$$

$$n^{1.7712} < n^2$$

$$\boxed{T(n) = \Theta(n^2)}$$

$$22. T(n) = T(n/2) + n(2 - \cos n)$$

$$a=1, b=2$$

$$c = \log_b a = \log_2 1 = 0$$

$$n^c = n^0 = 1$$

$$n(2 - \cos n) > n^c$$

$$T(n) = \Theta(n(2 - \cos n))$$