

HW) You are given  $n$  objects the  $i$ th object has weight  $W_i$  and a target weight  $K$ . You have to tell the maximum value possible for any weight  $\leq K$

$n \leq 1e5$ ,  $k \leq 1e5$

Sum of weights  $\leq 1e5$

Q) Longest Increasing Subsequence

You are given an array of  $n$  integers find the longest strictly increasing subsequence of the array?

$n \leq 2000$

$A[i] \leq 1e9$

Input:

1 2 6 1 3 4

Output:

4

1 1

2  $dp[2] = \max(1, dp[1] + 1) = 2$

6 3

1 1

3 3

4 4

```
Int dp[n+1], ans=0;
```

```
for(int i=1; i<=n; i++)
```

```
{
```

```
    dp[i]=1;
```

```
    for(int j=1; j<i; j++)
```

```
    {
```

```
        if(a[j]<a[i])
```

```
        dp[i]=max(dp[i], dp[j]+1);
```

```
    }
```

```
    ans=max(ans, dp[i]);
```

```
}
```

```
Return ans;
```

Q) Longest common subsequence

Given 2 strings a and b find the longest common subsequence between the 2.

$|A| \leq 2000$

$|B| \leq 2000$

Input

abcdef	dp[4][4]	abcd	ahyc
	Abc	ahyc	abcd ahy abc ahy
Ahyczyerf	i-1, j	i, j-1	i-1, j-1

Output:

4 (acef)

$Dp[i][j]$  longest common subsequence of prefix of a consisting of i characters and prefix of b consisting of j characters.

$dp[3][4]=2$

abc ahyc

ab ahyc

abc ahy

$dp[2][3]=1$

Ab ahy

$dp[3][4]=\max(dp[3][4], dp[2][3]+1)=2$

$dp[i][j]=\max(dp[i-1][j], dp[i][j-1]);$

abc ahyc

abcd ahyc

$Dp[i-1][j-1] \quad a[i] == b[j]$

$dp[i][j]=\max(dp[i][j], dp[i-1][j-1]+1);$

$dp[i][j]=dp[i-1][j-1]+1;$

```

Int dp[n+1][m+1];
memset(dp, 0, sizeof(dp));
a='#'+a;
b='$'+b;
for(int i=1;i<=n;i++)
{
    for(int j=1;j<=m;j++)
    {
        dp[i][j]=max(dp[i-1][j], dp[i][j-1]);
        if(a[i]==b[j])
            dp[i][j]=max(dp[i][j], dp[i-1][j-1]+1);
    }
}
Int ans

```

## Coin Change Problem

You have infinite coins of certain values available. Find the minimum number of coins required to get a sum of N. If it is impossible to get a sum of N, print -1.

1, 2, 5

$9 = 5 + 2 + 2 = 3$  coins

2, 4, 5

$8 = 4 + 4$

$8 = 5 + 2 + 1$

$1 = -1$

$2 = 1$

$3 = -1$

$4 = 1$

$5 = 1$

$6 = 2$

$7 = 2$

$8 = 2$

States -  $dp[i]$  = Minimum number of coins required to get a sum equal to i.

Transitions - for(int j=0;j<n;j++)

$dp[i] = \min(dp[i], dp[i-a[j]]+1)$

Base Case -  $dp[0] = 0$

Goal -  $dp[N]$

```

int dp[N+1];
for(int i=0;i<=N;i++)
{
    dp[i] = INF;
    for(int j=0;j<n;j++)
    {
        if(i-a[j]>=0)
            dp[i] = min(dp[i],dp[i-a[j]]+1);
    }
}
if(dp[N]>=INF)
    cout << -1;
else
    cout << dp[N];

```

[https://atcoder.jp/contests/dp/tasks/dp\\_n](https://atcoder.jp/contests/dp/tasks/dp_n)

States -  $dp[i][j]$  = Minimum cost of merging all slimes from  $a[i]$  to  $a[j]$  into a single slime

Transitions - for(int  $k=i$ ;  $k<j$ ;  $k++$ )

$dp[i][j] = \min(dp[i][j], dp[i][k] + dp[k+1][j] + \text{sum}[i][j]);$

Base Case - if( $i==j$ )  $dp[i][j]=0$

Goal -  $dp[1][n]$

```

7 6 8 6 1 1
6 8 6 1
(6 8 6) + (1)
(6 8) + (6 1)
(6) + (8 6 1)

```

```

int mincost(int i,int j)
{
    if(i==j)
        return 0;
    if(dp[i][j]!=-1)
        return dp[i][j];
    dp[i][j]=INF;
    for(int k=i;k<j;k++)
        dp[i][j] = min(dp[i][j],mincost(i,k)+mincost(k+1,j)+sum[i][j]);
    return dp[i][j];
}

```

```
int main()
{
    int a[n];
    //Input
    memset(dp,-1,sizeof(dp))
    for(int i=0;i<n;i++)
    {
        sum[i][i]=a[i];
        for(int j=i+1;j<n;j++)
            sum[i][j] = sum[i][j-1]+a[j];
    }
    cout << mincost(0,n-1);
}
```

Time Complexity -  $O(N^3)$