

[https://atcoder.jp/contests/abc190/tasks/abc190\\_d](https://atcoder.jp/contests/abc190/tasks/abc190_d)

Q.) How many arithmetic progressions consisting of integers with a common difference of 1 have a sum of n?

Approach:-

Series :-  $a, a+1, a+2, \dots, a+(x-1)$  [ $a \rightarrow$  first term,  $x \rightarrow$  length]

$$\text{Sum} = x \cdot (2a + (x-1)) / 2 = n$$

$$x \cdot (2a + x - 1) = 2n$$

$$x \cdot (2a + x - 1) = m$$

Solution:-

$$x \cdot [2a + (x-1)] = m$$

$$x$$

$$k_1 = \frac{m}{x} \rightarrow \text{1st case}$$

$$\frac{m}{x}$$

$$k_2 = x \rightarrow \text{2nd case}$$

```
#define ll long long int
```

```
ll n;
```

```
cin >> n;
```

```
ll m = 2 * n, ans = 0;
```

```
for (ll x = 1; x * x <= m; x++) {
```

```
    if (m % x == 0) {
```

```
        ll k1 = (m / x), k2 = x;
```

```
        if ((k1 + 1 - x) % 2 == 0) ans++;
```

```
        if (k2 != k1 && (k2 + 1 - (m / x)) % 2 == 0) ans++;
```

```
        // If true then you got a series whose first term(a) and length=x
```

```
    }
```

```
}
```

```
cout << ans << endl;
```

Q:

Universe is defined by 2 elements -> s and t.

N universe are given .

Degree of connectivity is defined as

$$D(U_i, U_j) = S_i + S_j + |t_i - t_j|, \text{ if } |t_i - t_j| \leq m$$
$$0 \quad \text{Else}$$

3 3 -> n and m

0 0 -> 1st universe

3 0 -> 2nd universe

9 2 -> 3rd universe

$$D(1, 2) = 3$$

$$2 \leq n \leq 10^5$$

$$0 \leq m \leq 10^{16}$$

$$0 \leq s_i, t_i \leq 10^{16}$$

$$S_i + S_j + |t_i - t_j| \rightarrow S_i + S_j + t_i - t_j, \text{ if } t_i > t_j$$

$$S_i + S_j + t_j - t_i, \text{ if } t_j > t_i$$

$$S_i + S_j + t_i - t_j = (s_i + t_i) + (s_j - t_j)$$

Deque -> possible universes for that i that satisfy  $t_i - t \leq m$

F.....b

B -> push i s-t =9

4,5,6

<https://codeforces.com/problemset/problem/582/A>

```
int main()
{
    int n;
    cin>>n;
    map<int,int,greater<int> > freq;
    // store frequency of all elements
    for(int i=0 ; i<n*n; i++)
    {
        int num;
        cin>>num;
        freq[num]++;
    }
    vector<int> ans;
    for(auto it=freq.begin(); it!=freq.end(); )
    {
        if(it->second>0)
        {
            freq[(it->first)]--;
            for(int i=0; i<ans.size(); i++)
            {
                freq[__gcd(ans[i],(it->first))]-=2;
            }
            ans.push_back(it->first);
        }
        else
        {
            it++;
        }
    }
}
```

```
    }  
}  
for(int i=0; i<n; i++)  
{  
    cout<<ans[i]<<" ";  
}  
}
```

Time complexity:  $O(n^2 \log(n))$