## Number Theory - 1

## **Properties of Modulus**

```
int x = -8\%5; // x = -3
-> -3,-8,-13,2
x+=m
```

#### **Modulus:**

```
Non-negative remainder when x is divided by m. It is also said as "x mod m".

Range of "x mod m": [0,m-1]

x%m -> modulus

If x is negative -> modulus = (x%m + m) % m
```

# Q. N is a given number Find the factorial N % m

```
Eg. n=5 m = 20
5! = 120
120%m = 0;
```

```
long long a=1e18,b=2e18;
int m = 1e9+7; // -> 1000000007
int x = (a+b)%m;
cout<<x;
a>0,b>0
```

#### **Properties of Modulus:**

```
1. (a+b)\%m = ((a\%m)+(b\%m))\%m
```

```
2. (a-b)\%m = ((a\%m)-(b\%m)+m)\%m
3. (a*b)\%m = ((a\%m)*(b\%m))\%m
4. (a/b)%m = ((a%m)*(b^-1%m))%m
m=5,a=9,b=14;
a\%m = 4.b\%m = 4
m=5,a=8,b=4;
(a-b)\%m = 4
a\%m = 3
b\%m = 4
-1\%m = -1
Q. Find (N!)%m. Where (N! = 1*2*3*4*5*.....*n)
N -> input
M = 1e9+7;
(n!\%m)?
N!\%m = (n*(n-1)*(n-2)*(n-3)....1)\%m = ((n\%m)*((n-1)*(n-2)...1))\%m))\%m
1<=n<=1000
int ans=1;
for(int i=1;i<=n;i++){</pre>
      ans = (ans\%m)*(i\%m);
      ans%=m;
```

### Q. Find $(x^n)$ . (x raised to the power n or $x^n$ )

#### Method 1: (Using a for loop)

return ans;

```
int ans=1;
for(int i=1;i<=n;i++){</pre>
```

```
ans*=x;
}
```

#### Method 2: (Using a recursive function)

```
int fun(int x,int n) //fun(x,n) -> x^n;
{
    if(n==0){
        return 1;
    }
    return x*fun(x,n-1);
}
```

 $fun(2,5) \rightarrow fun(2,4) \rightarrow fun(2,3) \rightarrow fun(2,2) \rightarrow fun(2,1) \rightarrow fun(2,0)$ Time complexity of both methods is O(n)

## **Binary Exponentiation**

```
2^8 = (2^2)^4 = 4^4 = (4^2)^2 = 16^2 = (16^2)^1 = 256 -> logn

x^n -> logn

2^5 = 2^2^4

x^n
```

#### Q. Find (x^n) mod m

#### Method 1: Using a recursive function

```
int binaryExponentiation(int x,int n,int m) // x^n O(logn)
{
    if(n==0){
        return 1;
    }
    if(n%2==0){
        return binaryExponentiation(((x%m)*(x%m))%m,n/2,m);
    }
}
```

```
return
((x%m)*binaryExponentiation(((x%m)*(x%m))%m,(n-1)/2,m)%m)%m;
}
```

Iterative (Using a loop) :-

```
int binaryExponentiation(int x,int n,int m)  // O(logn)
{
    int res=1;
    while(n!=0){
        if(n%2==1){
            res = ((res%m)*(x%m))%m;
        }
        x = ((x%m)*(x%m))%m;
        n/=2;
    }
    return res;
}
```

Time complexity of both methods is O(log N) here, which works very fast.

Even, for large numbers like N=10<sup>20</sup>, log N has a very small value.

#### **Prime Number**

#### **Definition**

Prime numbers are those numbers that are divisible by only 1 and the number itself. i.e the number of divisors should be 2.

Q. 2,3,4,6,7,8,9 PPNNPNN Q. Write a C++ code to check whether the given number is prime or not.

```
Time Complexity - O(n)
[ This was a slow method ]
```

#### **Important Key Point**

Consider a natural number N
If i is a divisor of N.
Then, (N/i) is also a divisor of N.

```
e.g
N=6;
2 is divisor of N bcz (6%2==0);
6/2 = 3 is also divisor of N bcz (6%3==0);
N<=10^10 [Worst case of (N)]
```

i is divisor of N then (N/i) is also a divisor of N.

```
N=12 -> 1 2 3 4 6 12
1 to <=sqrt(N)
```

Property: If you have got a divisor > sqrt(N), Then there must be a divisor that is less than sqrt(N)

```
N=a*b;
a<=sqrt(N);
b>=sqrt(N);
```

## Fast method to check if n is prime

```
int n;
cin>>n;
int divisors=0;
for(int i=1;i<=sqrt(n);i++){
    if(n%i==0) {
        // n is divisible by i
        int first_Divisor=i;
        int second_Divisor=(N/i);
        if(first_Divisor!=second_Divisor) divisors+=2;
        else divisors++;
    }
}
if(divisors==2) cout<<"The given number is a prime number"<<endl;
else cout<<"The given number is not a prime number"<<endl;</pre>
```

```
Time complexity of this method: O ( sqrt(N) ) [ Faster than previous method ] sqrt(16)=4 16-> 1 2 4 8 16 i=4 i*i=16 10^6 -> n/2 == 5*10^5 sqrt(N) == 10^3
```

## Some important in-built functions

- 1.  $pow(n, x) => Finds n^x in O(logn)$
- 2. sqrt(n) => Finds square root on n.

**Caution:** If the numbers are small then only use pow() function, otherwise use the Binary Exponentiation method to calculate power.