

# Coin change

total -> 11

coins-> 1,2,3,4,5 (infinite no. of each)

11-> (5,5,1), (5,4,2),.....

find the total number of ways?

base case->

ways to form 0 -> 1

-> we have formed amount 4 in 5 ways

-> we have one more coin '2'

-> how many ways will be there to form 6?

way 1+ 2-> 6

way 2+ 2-> 6

way 3+ 2-> 6

way 4+ 2-> 6

way 5+ 2-> 6

5 ways to form 6

//without permutation

n=4 and coins -> 3,2,1

0-> 1

1-> 0+1

2-> 1+1

3-> 1+0+2

4-> 0+1+3

```
//with permutation  
n=4 and coins -> 1,2,3
```

```
0-> 1=1  
1-> 0+1=1  
2-> 0+1+1=2  
3-> 0+1+1+2=4  
4-> 0+1+2+4=7
```

## Code for 'without permutation'

```
#include<bits/stdc++.h>  
using namespace std;  
  
typedef long long ll;  
  
int main(){  
    ll k,n; cin>>k>>n;  
    ll coins[n];  
    for(ll i=0;i<n;i++){  
        cin>>coins[i];  
    }  
    ll ways[k+1]={1};  
    //1,3,2 only one way to iterate  
    for(ll i=0;i<n;i++){  
        for(ll j=0;j<=k-coins[i];j++){  
            ways[j+coins[i]]+=ways[j];  
        }  
    }  
    cout<<ways[k];  
}  
  
// coins array size * amount (n *k)
```

## Code for with permutation

```
#include<bits/stdc++.h>
using namespace std;

typedef long long ll;

int main(){
    ll k,n; cin>>k>>n;
    ll coins[n];
    for(ll i=0;i<n;i++){
        cin>>coins[i];
    }
    ll ways[k+1]={1};
    for(ll j=0;j<=k;j++){
        for(ll i=0;i<n;i++){
            if(j+coins[i]<=k){
                ways[j+coins[i]]+=ways[j];
            }
        }
    }
    cout<<ways[k];
}

// complexity -> coins array size * amount (n * k)
```

<https://www.spoj.com/problems/TRT/>

```
#include <bits/stdc++.h>
using namespace std;

//dp[i][j] -> from start and end we are at ith and jth index
// indexes -> 1,2,3.....i.....j,j+1,j+2.....n

int dp[2005][2005];
int func(int i, int j, int a[],int n)
{
    if(i==0 && j==n+1)
    {
        return 0;
    }
    if(dp[i][j]!=-1)
        return dp[i][j];
    int ans1=0;
    int ans2 = 0;
    if(i) // When we have come from i-1,j and we are taking the ith
element
        ans1 = func(i-1,j,a,n) + a[i] * (i+n-j+1);
    if(j<=n) // When we have come from i,j+1 and we are taking the jth
element
        ans2 = func(i,j+1,a,n) + a[j] * (i+n-j+1);
    dp[i][j] = max(ans1,ans2);
    return dp[i][j];
}

void solve()
{
    int n;
    cin>>n;
    int a[n+1];
    for(int i=1;i<=n;i++)
        cin>>a[i];
```

```

int ans = 0;
memset(dp, -1, sizeof(dp));
// For the ending condition we must cover the cases when we have
covered all the elements.
// i.e. i + (n-j+1) = n or j - i = 1 or j = i+1
for(int i=0; i<=n; i++)
ans = max(ans, func(i, i+1, a, n));
cout<<ans<<"\n";
}

```

### Matrix Chain Multiplication

A(n,m)

B(m,k)

C(k,p)

A\*B -> n\*m\*k operations to multiply

Matrix Chain Multiplication

ABC -> (AB)C or A(BC)

(AB)C -> n\*m\*k+n\*k\*p  
(n\*k)

A(BC)-> m\*k\*p+n\*m\*p  
m\*p

AB -> n\*m\*k

ABCDE ->

0 1 2 3 4  
p[]={1,2,3,4,5}

$dp[i][j]$  = Minimum cost of multiplying matrices from  $i$  to  $j$

$dp[0][1]=0$

$dp[i][i+1]=0$  for every possible  $i$

$i \dots j$

$k$

$M_1(i,k) * M_2(k,j)$

ll  $dp[n][m]$ ;

if( $dp[i][j] \neq -1$ ) // memoization

$dp[i][j] = \min (dp[i][k]+dp[k][j]+p[i]*p[k]*p[j])$  for  $k>i$  &&  $k<j$

Code:

```
#include<bits/stdc++.h>
#define INF INT_MAX
using namespace std;

int dp[1001][1001];

int find(int l,int r,int p[]){
    if(l+1==r)
        return 0;

    if(dp[l][r]!=-1)
        return dp[l][r];

    dp[l][r]=INF;

    for(int k=l+1;k<r;k++){
        dp[l][r]=min(dp[l][r],find(l,k,p)+find(k,r,p)+p[l]*p[k]*p[r]);
    }
    return dp[l][r];
}

int main(){
```

```
memset(dp, -1, sizeof(dp));  
// input p[]  
int n;  
cin >> n;  
  
int p[n];  
  
for(int i=0; i<n; i++)  
    cin >> p[i];  
  
cout << find(0, n-1, p);  
  
}
```

Try this: [https://atcoder.jp/contests/dp/tasks/dp\\_n](https://atcoder.jp/contests/dp/tasks/dp_n)