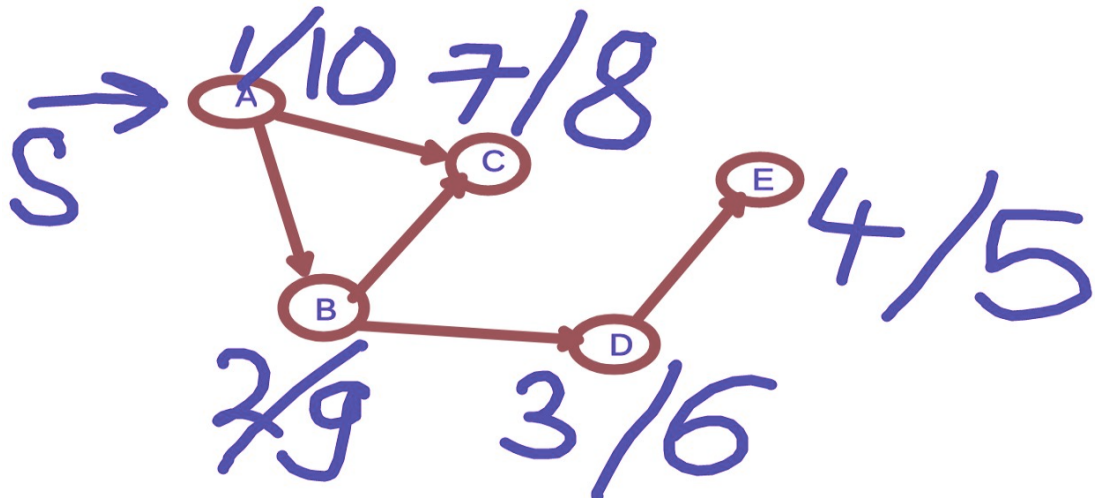


## DFS (Depth first search)



Visiting Time

Finishing Time

### Colour codes in DFS

**White:** Unvisited node

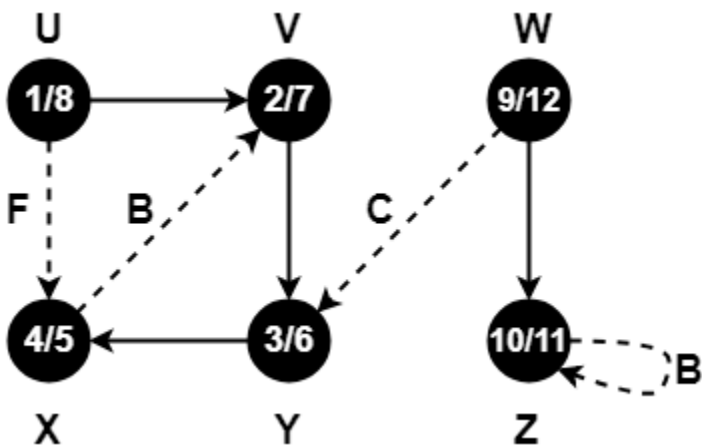
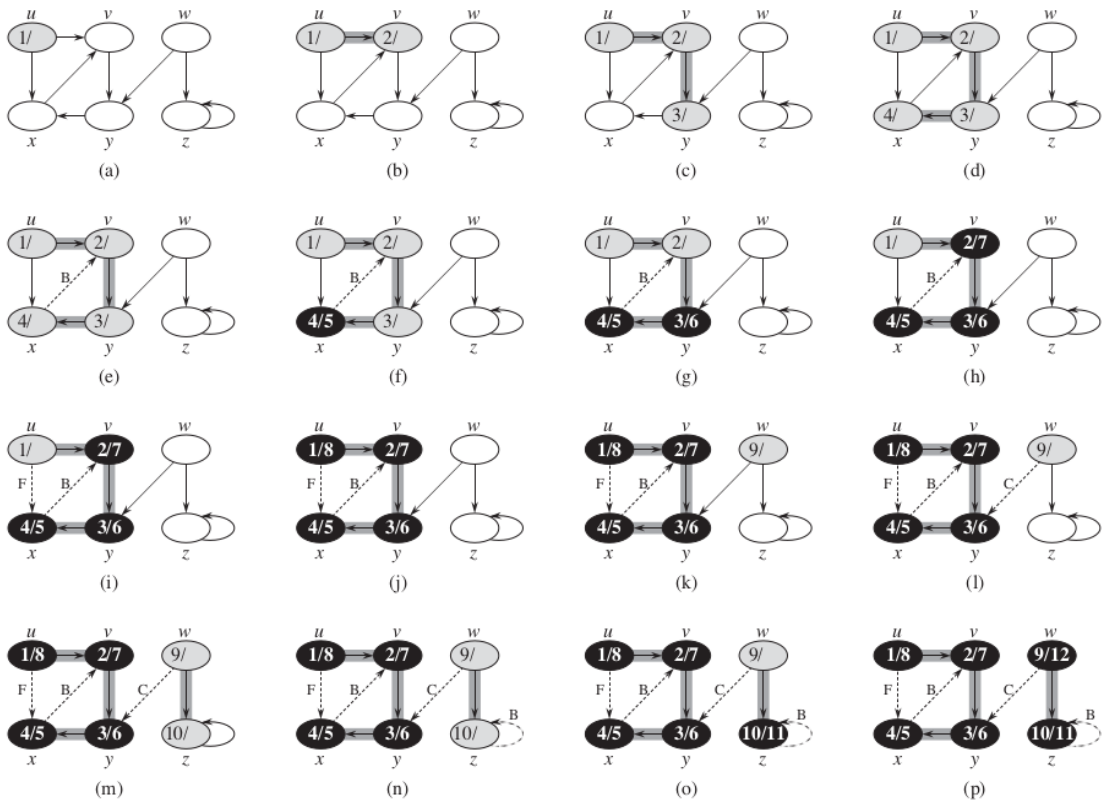
**Gray:** Visited but not finished node

(Visited but all its neighbours are not visited)

**Black:** Finished node

(Visited and all its neighbours are also visited)

**Try running DFS on below graph and calculating visiting and finishing time of each node:**



## Taking graph input as adjacency list

```
int n; // no. of nodes
vector<vector<int> > adj; // adjacency list
```

```

int main()
{
    cin>>n;
    int e; // no. of edges;
    cin>>e;
    adj.resize(n);
    while (e--)
    {
        int u,v;
        cin>>u>>v;
        adj[u].push_back(v);
        adj[v].push_back(u); // for undirected graph
    }
}

```

## Code for normal DFS

```

int n; // no. of nodes
vector<vector<int> > adj; // adjacency list
vector<bool> vis; // initialise all values
with false

void dfs(int node)
{
    vis[node]=true;
}

```

```

for(auto child: adj[node])
{
    if ( ! vis[child] )
    {
        dfs(child);
    }
}
}

```

Code for DFS when you need to calculate colour of each node at every instant

```

vector<int> colour;
// 0 - white (Use constants or #define for colours)
// 1 - gray
// 2 - black

void dfs(int node)
{
    colour[node]=gray;
    for(auto child: adj[node])
    {
        if ( colour[child] ==0 )
        {

```

```
        dfs(child);  
    }  
}  
colour[node]=black;  
}
```

## Code for DFS when you need to calculate visiting time and finishing time of nodes

```
int timerCode=1;
vector<int> visTime;
vector<int> finishTime;
void dfs(int node)
{
    vis[node]=true;
    visTime[node]=timerCode;
    timerCode++;
    for(auto child: adj[node])
    {
        if ( ! vis[child] )
        {
            dfs(child);
        }
    }
    finishTime[node]=timerCode;
    timerCode++;
}
```

## Calling DFS in main()

```
for(int i=0; i<n; i++)
{
```

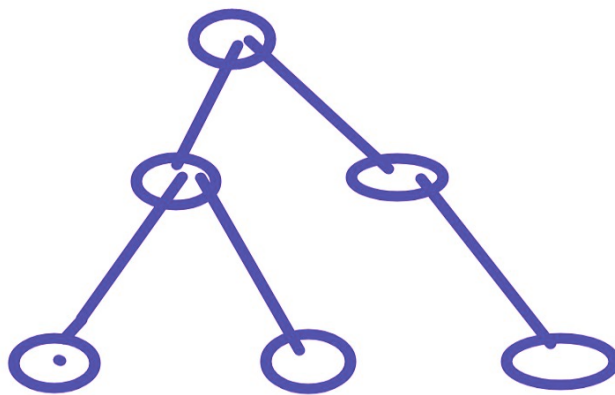
```
if (!vis[i])  
{  
    dfs(i);  
}  
}
```

## Tree

Tree is an undirected connected graph without any cycle

Properties:

- If there are  $n$  nodes in a tree, then it will have  $n-1$  edges



- If you are given a connected graph with  $n$  nodes and  $n-1$  edges, then the given graph is a tree

- If you are given a connected graph with  $n$  nodes and  $n$  edges, then the given graph will contain exactly 1 cycle
- Each node has a single parent in tree

**You have been given a tree of  $n$  nodes with root at 0.  
Find level of each node using DFS**

```
vector<int> level;
// level[0] =0;
void dfs(int node)
{
    vis[node]=true;
    for(auto child: adj[node])
    {
        if ( ! vis[child] )
        {
            level[child]=level[node]+1;
            dfs(child);
        }
    }
}
// in main, call dfs(0); // Because 0 is root
// of tree
```

**DFS on trees can also be implemented as:**

```
// call dfs(0, -1) in main()
void dfs(int node, int par)
```



```

{
for(auto child: adj[node])
{
    if (child != par)
    {
        level[child]=level[node]+1;
        dfs(child);
    }
}
}

```

**Q. Checking whether given graph is bipartite graph**

Link: <https://cses.fi/problemset/task/1668>

It can be done using DFS

```

vector<int> teamNum;
bool dfs(int node, int currTeam)
{
    vis[node]=true;
    teamNum[node]=currTeam;
    for(auto child: adj[node])
    {
        if ( teamNum[child] == teamNum[node])
            return false;
        if ( ! vis[child] )

```

```
        {
            bool temp=dfs(child, 3 - currTeam);
            if(temp==false)
                return false;
        }
    }
    return true;
}

int main()
{

for(int i=0; i<n; i++)
{
    if(!vis[i])
    {
        bool temp=dfs(i,1);
        if(!temp)
        {
            cout<<"IMPOSSIBLE";
            return 0;
        }
    }
}

// Print the team values if division is possible
```

```
}
```

## Detecting cycle in a graph

no of nodes=n

```
int colour[n];
```

**Return true if cycle present else false.**

```
bool dfs(int node, int parent)
{
    colour[node]=1;
    for(auto child: adj[node])
    {
        if ( child!=parent && colour[child] == 1 )
        {
            // You can remove the condition for
            child!=parent, when you want to detect even the
            2 node cycles in directed graph like 1->2, 2->1
            return true;
        }
        else if( child!=parent && colour[child] == 0 )
```

```

    {
        bool temp = dfs(child, node);
        if( temp == true) return true;
    }
}
colour[node]=2;
return false;
}

```

```

int main()
{
    // Take graph input
    // .....
    for(int i=0;i<n;i++)
        colour[i]=0;
    // Mark colour of all nodes as WHITE
    (unvisited)

    for(int i=0; i<n; i++)
    {
        if(colour[i] == 0)
        {
            bool cycle = dfs(i,-1);
            if(cycle)
            {

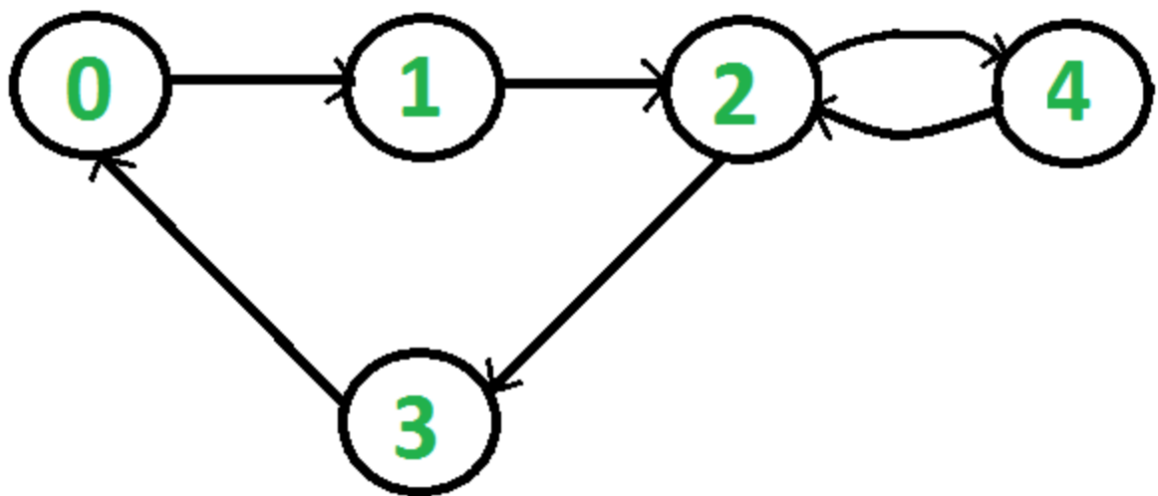
```

```

        cout<<"Cycle Present";
        return 0;
    }
}
}
cout<<"No cycle found";
return 0;
}

```

You can try running this code over below example:



```

*dfs(0)
col[0]=1
0->child==1
**dfs(1,0(parent))
col[1]=1
***dfs(2)
col[2]=1
****dfs(3)

```

col[3]==0

col[3]=1

3->

child=0

color of child==1-->cycle

col[0]==1

return true;

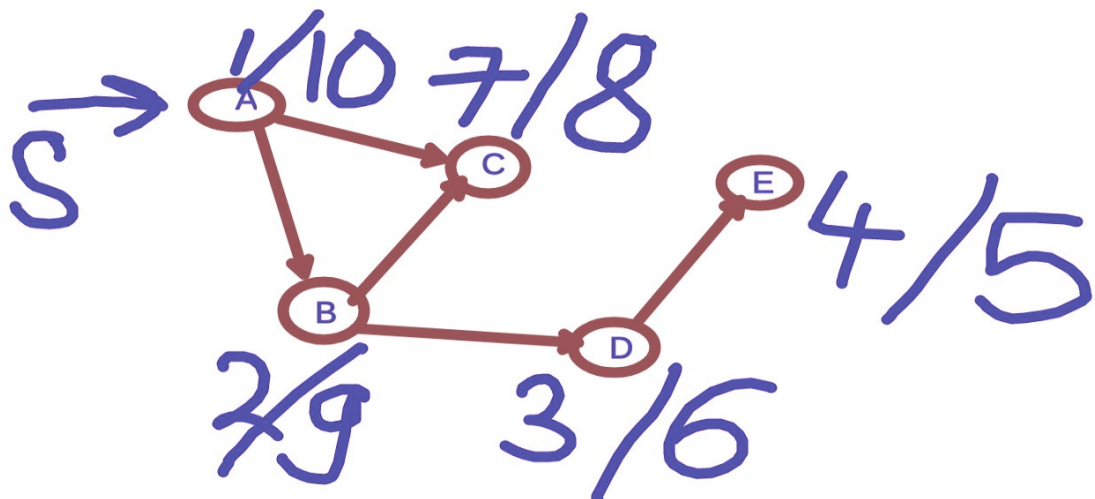
\*\*\*\*dfs(4)

col[3]=2

col[2]=2

col[1]=2

col[0]=2



dfs(A)

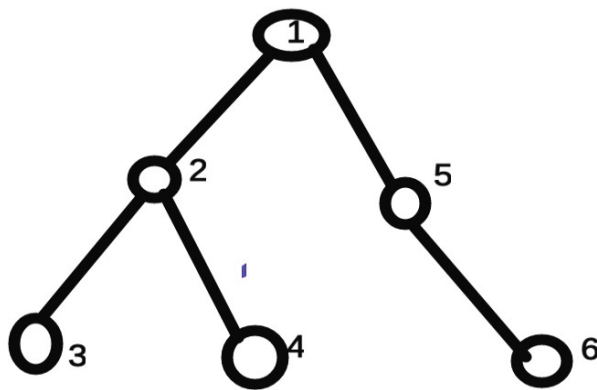
col[A]=1

A->child:B,C

```
*dfs(B)
  child:C,D
  **dfs(C)
    col[C]=1;
    child C-NULL
    col[C]=2
    false ret
```

```
**dfs(D)
  col[D]=1
  D->CHILD:E
  ***dfs(E)
    col[E]=1
    col [E]=2
    ret false;
  col[D]=2
  ret false
```

```
*dfs(C)
```



**Given a tree, you need to find the gcd of all the values of nodes in subtree of every node. 0 is root node**

Example, in above graph,  $\text{ans}[5] = \text{gcd}(5,6)$

$\text{ans}[2] = \text{gcd}(2,3,4)$

No. of nodes,  $n \leq 10^5$

**Try to solve this problem using DFS**

`vector<int> ans(n);`

```
// call dfs(0, -1) in main()
void dfs(int node, int par)
{
    ans[node]=val[node];
    for(auto child: adj[node])
    {
```



```

        if (child != par)
        {
            dfs(child);
            ans[node] = __gcd(ans[node], ans[child]);
        }
    }
}

```

## Topological Sorting: (Kanh's Algo)

```

int indeg[n+1]={0};

for(int i=0;i<m;i++){
    int u,v;
    cin>>u>>v;
    // u -> v
    indeg[v]++;
}

vector<int>topl;

queue<int>q;

for(int i=1;i<=n;i++)
if(indeg[i]==0)
q.push(i);

while(!q.empty()){
    auto fst=q.front();
    q.pop();

    topl.push_back(fst);
    for(auto child:adj[fst])
    {

```

```

        indeg[child]--;

        if (indeg[child]==0)
            q.push(child);
    }
}

for(auto e:topl)
    cout<<e<<" ";

```

DFS Method:

Arrange in decreasing order of exit times!

```

int n; // number of vertices
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> ans;

void dfs(int v) {
    visited[v] = true;
    for (int u : adj[v]) {
        if (!visited[u])
            dfs(u);
    }
    ans.push_back(v);
}

void topological_sort() {
    visited.assign(n, false);
    ans.clear();
    for (int i = 0; i < n; ++i) {
        if (!visited[i])
            dfs(i);
    }
    reverse(ans.begin(), ans.end());
}

```

Practice:

<https://cses.fi/problemset/task/1679>

[Link to Rough](#)