

GARLAND

```
int dp[105][105][105][2];
```

```
int a[105];
```

```
int n;
```

```
/*
```

```
5 36 7 3 2 2 5 4 6 2 5 7 4 7 8 4
```

```
*/
```

```
int f(int i,int odd,int even,int prev){
```

```
    if(i==n)
        return 0LL;
```

```
    else if(dp[i][odd][even][prev]!=-1)
        return dp[i][odd][even][prev];
```

```
    //two cases
```

```
    //attached or removed bulb
```

```
    if(a[i]){
```

```
        int parity=(a[i]&1LL);
```

```
        return dp[i][odd][even][prev]=f(i+1,odd,even,parity)+(prev!=parity);
```

```
    }else{
```

```

int op1=INF,op2=INF;

if(odd){

    op1=f(i+1,odd-1,even,1)+(1-prev);

}

if(even){

    op2=f(i+1,odd,even-1,0)+prev;

}

return dp[i][odd][even][prev]=min(op1,op2);
}
}

```

```

inline void solve(){

    // int n;
    cin>>n;

    // vector<int>a(n);

```

```

int odd=0,even=0;
for(int i=0;i<n;++i){
    cin>>a[i];
    if(a[i]==0)
        continue;
    if(a[i]&1LL)
        odd++;
    else
        even++;
}

```

```

//odd,even represents number of odds and evens in given array
odd=(n+1)/2-odd;
even=n/2-even;

```

```

memset(dp,-1LL,sizeof(dp));
//now odd,even represents number of odds /evens we can take
//f(0,odd,even,prev)

```

```

//attached

```

```

if(a[0]){

```

```

    int parity=(a[0]&1ll);
    out(f(1,odd,even,parity));

```

```

}else{

```

```

    int op1=INF,op2=INF;

```

```

    if(odd){
        op1=f(1,odd-1,even,1);
    }

```

```

    if(even)
        op2=f(1,odd,even-1,0);

    out(min(op1,op2));//answer

}

}

```

Bit Manipulation

Decimal System -> 10 digits 0-9

Binary System -> 2 digits 0 and 1 -> bits

3 -> 11 -> $2^1 + 2^0$

Decimal -> 97 -> $9 \cdot (10^1) + 7 \cdot (10^0)$

Set -> Bit is 1

Unset -> Bit is 0

10110 -> $2^1 + 2^2 + 2^4 = 2 + 4 + 16 = 22$

Operators -> Bitwise AND, OR, NOT, XOR

1 AND 1 -> 1 (Both the operands should be 1 to give 1 as output)

10110

01101

00100 -> Bitwise AND

1 OR 0 -> 1 (Result will be 1 if either of the operands are 1)

10010

11000

11010 -> Bitwise OR

NOT 1 -> 0 (Inverts the operands)

NOT 0 -> 1

11101

00010 -> Bitwise NOT

1 XOR 1 -> 0 (Result will be 1 if both operands are different, and 0 if both are same)

1 XOR 0 -> 1

0 XOR 1 -> 1

0 XOR 0 -> 0

10010

11000

01010 -> Bitwise XOR

Shifts -> Left Shift, Right Shift

000110 left shift 2 -> 00011000

110101 right shift 2 -> 001101

Bitwise AND -> &

Bitwise OR -> |
Bitwise XOR -> ^
Bitwise NOT -> ~

Left shift -> <<
Right Shift -> >>

$2 \& 1 \rightarrow 10 \& 01 = 0$
 $2 | 1 \rightarrow 10 | 01 = 11 = 3$
 $2 \wedge 1 \rightarrow 10 \wedge 01 = 11 = 3$
 $\sim 2 \rightarrow \sim(10) = 01 = 1$

int -> 32 bits -> $2^0 + 2^1 + \dots 2^9 + 2^{11} + \dots 2^{30}$ (2^{33} not possible)
111000011010101011010101 -> maximum allowed width is 32
long long -> 64 bits
1010100101100011101000010101001101 -> maximum width is 64
 $2^{40} + 2^{41}$

int x = 3; // 000000...0011 -> 30 zeroes and 2 ones
int y = ~x; // 1111111...1100 -> 30 ones and 2 zeroes

int z = 3<<1; // 00000..0011 << 1 -> 000..00110 -> $6 = 3*2$
to multiply by 2^i -> perform $x<<i$
Similarly
 $z>>5$ -> equivalent to $z / 2^5$

int z -> 1110000000..0000 width 32
 $z<<5$ -> 11100 000000000..00 width 32 -> overflow

int w = 000...1110 width 32 = $2+4+8 = 14$
 $w>>6$ -> 00...0000 001110
//14/64 -> answer is not incorrect but set bits are being lost

long long z = 1110000..00 width 32 -> 0000..00 11000..0000 width 64
z<<5 -> 000.. 00 11100 0000000...000 -> no loss

$1 \ll 3 = 8$ equivalent to $1 * (2^3)$

$1 \ll i = 2^i$ -> 1 is being treated as int

$1LL \ll i$ -> 1 is being treated as long long, safe to use $i = 30, 40, 50$ etc.

checking whether a particular bit is set or not

long long x = 111000011 -> find out whether 3rd bit is set or not, with expression involving bitwise operators. (bits start from 0)

$(x \gg 3) \& 1$

111000011 -> 11100001 -> 1110000 -> 111000

111000 & 1 = 111000 & 000001 -> 0

$x \& (1 \ll 3)$ -> 111000011 & 000001000 -> 0

for checking ith bit is set or not -> use $x \& (1LL \ll i) == 0$

setting the ith bit -> $x | (1LL \ll i)$

unsetting the ith bit -> $x \& (\sim(1LL \ll i))$

11100011 unset the 1st bit

$1 \ll 1$ -> 00000010

$\sim(1 \ll 1)$ -> 11111101

11100011
&11111101
=11100001

toggling the i th bit (0 to 1 and vice versa) $\rightarrow x \oplus (1 \ll i)$

toggle 2nd bit of 11100011

11100011
 $\oplus 00000100$
=11100111 \rightarrow toggled