DYNAMIC PROGRAMMING

Day 4

Q)https://www.geeksforgeeks.org/maximum-sum-path-in-a-matrix-from-top-left-to-bottom-right/?ref=rp

Given a matrix: n X m
eg.
A 2X3 matrix -
1 2 3
4 5 6

dp[2][2]-->1,2,5
-->1,4,5
dp[2][2]=10
dp[1][2]=1,2-->3
dp[1][3]=6
dp[2][1]=5
dp[2][3]=max(dp[1][3],dp[2][2])+6
=max(6,10)+6==16
movement:right or down
dest:bottom right(n,m)

start:top left(1,1)

1-2-3-6-->6+3+2+1=12
1-2-5-6-->1+2+5+6=14
1-4-5-6-->16
max val find?

(i,j)-->(i-1,j)(down)
→(i,j-1)(right)
(n,m)-->(n-1,m),(n,m-1)
dp[i][j]-->store max value we can obtain
(1,1)-->(i,j)

dp[n][m]-->ans
int dp[n+1][m+1]
i,j
dp[i][j]=max(dp[i-1][j],dp[i][j-1])+arr[i][j]

base case:(i==0||j==0)return 0;

```
for(int i=0;i<=n;i++)
 dp[i][0]=0;
for(int j=0;j<=m;j++)
 dp[0][j]=0;
```

```
for(int i=1;i<=n;i++)
{
 for(int j=1;j<=m;j++)
 {

dp[i][j]=max(dp[i-1][j],dp[i][j-1])+
arr[i][j];
 }
}
ans==dp[n][m]
```

Q)https://codeforces.com/problemset/problem/1286/A

N bulbs--numbering 1 to n
prev bulb odd numbered--current odd numbered
prev even--current even
missing bulbs--0

odd->even,even->odd→++1
complexity-->min

5
0 5 0 2 3
odd=1
even=1
missing:1,4
1-->odd 5 even 2 3
complexity=2(5,4),(2,3)
2-->even 5 odd 2 3
complexity=(4,5),(1,2),(2,3)==3
min complexity achieve?
odd-->fill,even-->fill

ith index-->upto n
odds rem,even rem

missing=8
odd=3
even=5
odd--,odd=2
even=8-2=6
odd+even=missing
even-->missing

-->ith index
-->odds rem
-->even rem
ith index==odd,prev??
-->prev(odd--1/even--0)
prev==2(i==0)

ith→(i+1)th
ith→(i-1)th
(i+1)th-->ith
ith fill
ith index-->assume garland upto i

i+1th
ith-->odd→(i+1)th/even->(i+1)th

ith-->odd/even

if(arr[i]==0)//filling req

else//already filled

```
int fun(int i,int oc,int ec,int
prev)
{
 int res=INT_MAX;

 if(arr[i]!=0)
{
 if(arr[i]%2==prev||i==0)
  res= fun(i+1,oc,ec,arr[i]%2);
 else
 res= 1+fun(i+1,oc,ec,arr[i]%2);
 }
else
{
 if(oc>0)
if(i==0)
{
 if(arr[i]==0)
  {
  }
 el
 }
res=min(res,fun(i+1,oc-1,ec,1)+(prev
```

```
==0)//slot odd fill
if(ec>0)

res=min(res,fun(i+1,oc,ec-1,0)+(prev
==1)//slot even fill
return res;
}
```

0 5 0 2 3

```
int arr[105];
int n;
int dp[105][55][55][3];

int fun(int i,int oc,int
ec,int prev)
{

  if(i==n)
   return 0;

if(dp[i][oc][ec][prev]!=-1)
```

```
    return
dp[i][oc][ec][prev];
    int res=1e9;
   if(arr[i]!=0)
   {

if((prev==2)||((arr[i]%2)==p
rev))

res=fun(i+1,oc,ec,arr[i]%2);
       else

res=1+fun(i+1,oc,ec,arr[i]%2
);
   }
   else
   {
       if(prev==2)
       {
          if(oc>0)
```

```
res=min(res,fun(i+1,oc-1,ec,
(int)1));
            if(ec>0)

res=min(res,fun(i+1,oc,ec-1,
(int)0));
        }
        else
        {
            if(oc>0)
            {

res=min(res,fun(i+1,oc-1,ec,
1)+(prev==0));
            }
            if(ec>0)

res=min(res,fun(i+1,oc,ec-1,
0)+(prev==1));
        }
    }
```

```
   return
dp[i][oc][ec][prev]=res;
}

int main()
{
    memset(dp,-1,sizeof(dp));
    cin>>n;
    int oc=0,ec=0;
    for(int i=0;i<n;i++)
    {
        cin>>arr[i];
        if(arr[i]%2)
         oc++;
        else if(arr[i]!=0)
         ec++;
    }
    oc=(n+1)/2-oc;
    ec=(n/2)-ec;
    //cout<<oc<<"
"<<ec<<"\n";
    int res=fun(0,oc,ec,2);
```

```
    cout<<res<<"\n";
    return 0;
}
```

Matrix Chain Multiplication
(MCM-type problems in DP)

A(BC) or (AB)C

A -> i*j, B->j*k
A*B -> i*k * j

ABCDEF -> minimize the number of addition operations
A -> 10*30, B -> 30*5, C->5*60
(AB)C ->10*30*5 + 10*5*60= 4500
A(BC) -> 30*5*60 + 10*30*60=27000

27000/4500 = 6 times

A[0….N]
Size of ith matrix is A[i-1]*A[i]

Hint : (AB)(CDEF) -> placing brackets lets us solve independent problems
Dimension of (AB) = first dimension of A * second dimension of B
Dimension of (CDEF) = first dimension of C * second dimension of F

DP state -> i,j ->index of first matrix and index of last matrix
M1,M2,M3,....Mn

Dp[i][j] -> min number of addition operations required if you multiply matrices Mi...Mj optimally

Dp[i][j] = min(dp[i][k]+dp[k+1][j] + A[i-1]*A[k]*A[j]) , k from i to j-1

(ABCDE) -> (A)(BCDE) or (AB)(CDE) or (ABC)(DE) or (ABCD)(E)

(A)(BCDE) -> M1=A and M2=BCDE
M1 -> first dim of A * second dim of A
M2 -> first dim of B * second Dim of E

```
int dp[n+1][n+1];
memset(dp,-1,sizeof dp);

int f(int l, int r)
{
if(l==r) return 0;
if(dp[l][r]!=-1) return dp[l][r];
dp[l][r]=1e9;
for(int k=i;k<r;k++)
dp[l][r]=min(dp[l][r],
dp[l][k]+dp[k+1][r]+
    a[l-1]*a[k]*a[r]);
return dp[l][r];
}
```

Now, try these problems:

1. https://www.spoj.com/problems/MIXTURES/
2. https://atcoder.jp/contests/dp/tasks/dp_n
3. https://leetcode.com/problems/minimum-cost-to-cut-a-stick/

Principle of Inclusion and Exclusion (PIE)

|A U B| = |A| + |B| - |A inter B|

|U Ai| = sum(|Ai|) - sum(|Ai inter Aj|)
+sum(|Ai inter Aj inter Ak|) - ….

D(n) -> derangement of length n
Number of permutation P where Pi notequal to i for all i.

N! - F
F -> number of permutations where there is atleast one index i such that Pi = i.

Ai -> number of permutations where ith element is fixed, i.e. Pi=i
F = |U Ai| = nC1 * (n-1)! - nC2*(n-2)! +...(-1)^(n+1)*nCn*(n-n)!

Homework:
https://codeforces.com/problemset/problem/559/C