

CSE 310 Recitation 4

Objectives:

1. Heaps (array representation, tree representation, max heap, min heap), Max-heapify (heap function).

Instruction

1. For all recitation: **the solution should be clearly typed or written and must be saved in .pdf or .jpg format. Note: unreadable answer receives no credits!**
2. All recitation must be submitted through the link posted on Blackboard, we do NOT accept any hand-in submissions or submissions sent through emails!

Question

1. For the BUILD-MAX-HEAP(A) procedure we discussed in class (pseudo code given below), can we change the code in line 2, to increase the loop index i from 1 to $\text{length}[A]/2$? Why or why not?

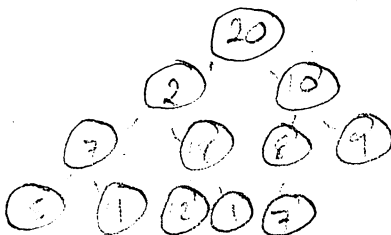
BUILD-MAX-HEAP(A)

- 1 heap-size[A] = length[A]
- 2 for $i = \text{length}[A]/2$ to 1 do
- 3 HEAPIFY(A, i)

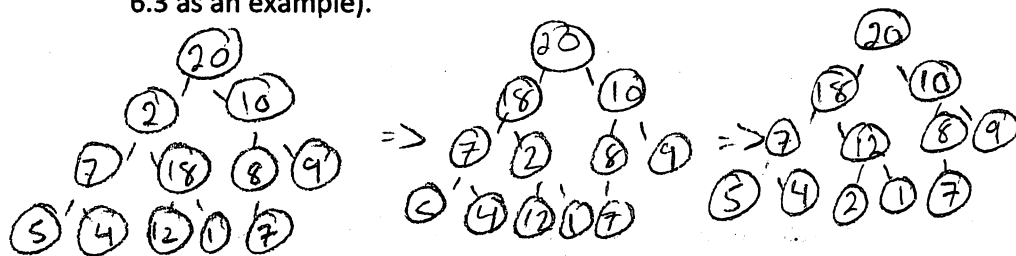
No, in order to build max heap, you need to start in the middle. Because its a recursive call of heapify, and heapify swaps the children of a parent node to let the largest be the parent. And if the bottom of the tree is not a max heap, then you won't be able to make one if you start at the top.

2. Consider the integer array $A = [20, 2, 10, 7, 18, 8, 9, 5, 4, 12, 1, 7]$

A. Give the complete binary tree representation for the array A



B. You can see that A is not a max heap. Use the pseudo code BUILD-MAX-HEAP(A) (textbook pp. 157) to draw the intermediate complete binary trees until you get a max heap (see pp.158 Figure 6.3 as an example).



C. HEAP-INCREASE-KEY(A, i, key) is an operation used in priority queues and we discussed it in class. Given the max-heap you get in above step B, show the operations of HEAP-INCREASE-KEY(A, 10, 19) by drawing the intermedia binary tree step-by-step (see pp.165 Figure 6.5 as one example)

