

18.997 Probabilistic Method Problem Set #3

Holden Lee

3/22/11

Problem 1 ($\mathfrak{3.1}$, $R(k, k)$)

Lemma 1.1:

$$\binom{n}{k} \leq \frac{1}{e} \left(\frac{en}{k} \right)^k.$$

Proof. By integral estimation,

$$\begin{aligned} \ln k! &= \sum_{m=1}^k \ln m \\ &\geq \int_1^k \ln x \, dx \\ &= k \ln k - k + 1. \end{aligned}$$

Exponentiating gives $k! > e \left(\frac{k}{e} \right)^k$. Hence

$$\binom{n}{k} = \frac{n(n-1) \cdots (n-k+1)}{k!} \leq \frac{n^k}{k!} \leq \frac{1}{e} \left(\frac{en}{k} \right)^k.$$

□

Let $a = \left\lfloor \frac{k}{e} 2^{\frac{k}{2}} \right\rfloor / \frac{k}{e} 2^{\frac{k}{2}}$, and put $n = \left\lfloor \frac{k}{e} 2^{\frac{k}{2}} \right\rfloor = a \frac{k}{e} 2^{\frac{k}{2}}$ in

$$R(k, k) > n - \binom{n}{k} 2^{1-\binom{k}{2}}$$

and use the above estimate to get

$$\begin{aligned} R(k, k) &> a \frac{k}{e} 2^{\frac{k}{2}} - \frac{1}{e} \left(\frac{en}{k} \right)^k 2^{1-\binom{k}{2}} \\ &= \left(a - \frac{1}{k} \left(\frac{en}{k} \right)^k 2^{1-\frac{k^2}{2}} \right) \frac{k}{e} 2^{\frac{k}{2}}. \\ &= \left(a - \frac{2a^k}{k} \right) \frac{k}{e} 2^{\frac{k}{2}}. \end{aligned}$$

For $k \geq 3$, we have $1 - \frac{1}{2^{\frac{k}{2}}} < a \leq 1$, and $\frac{1}{2^{\frac{k}{2}}} = o(1)$. Since $\frac{2a^k}{k} \leq \frac{2}{k} = o(1)$ as well, $R(k, k) \geq (1 - o(1)) \frac{k}{e} 2^{\frac{k}{2}}$.

Problem 2 (*3.2, $R(4, k)$*)

Consider a complete graph with n vertices. Call the first color red and the second blue. Color each edge in the graph red with probability p and blue with probability $1 - p$. The probability that a given set of 4 vertices forms a red K_4 is $p^{\binom{4}{2}}$ and the probability that a given set of k vertices forms a blue K_k is $(1 - p)^{\binom{k}{2}}$.

Let X be the total number of red K_4 's and blue K_k 's. By linearity of expectation, since there are $\binom{n}{4}$ groups of 4 vertices and $\binom{n}{k}$ groups of k vertices.

$$\mathbb{E}(X) = \binom{n}{4} p^6 + \binom{n}{k} (1 - p)^{\binom{k}{2}}.$$

There exists a coloring with at most $\mathbb{E}(X)$ red K_4 's and blue K_k 's. Pick a vertex from each red K_4 and blue K_k and delete it. We obtain a graph with at least $n - \mathbb{E}(X)$ vertices and no red K_4 or blue K_k . This shows that for any $n \in \mathbb{N}$ and any $p \in [0, 1]$,

$$R(4, k) > n - \binom{n}{4} p^6 - \binom{n}{k} (1 - p)^{\binom{k}{2}}.$$

Assume $k \geq 3$. Now pick $n = a \left(\frac{k}{\ln k}\right)^2$ and $p = \frac{2 \ln k}{k}$, where a is to be chosen (depending on k to make n an integer, but close to a constant). Using Lemma 1.1,

$$\begin{aligned} R(4, k) &\geq n - \frac{n^4}{24} p^6 - \left(\frac{en}{k}\right)^k (1 - p)^{\binom{k}{2}} \\ &\geq n - \frac{n^4}{24} p^6 - \left(\frac{en}{k}\right)^k e^{-p \binom{k}{2}} \\ &= a \left(\frac{k}{\ln k}\right)^2 - \frac{1}{24} a^4 \left(\frac{k}{\ln k}\right)^8 2^6 \left(\frac{\ln k}{k}\right)^6 - \left(\frac{eak}{(\ln k)^2}\right)^k e^{-\frac{2 \ln k}{k} \binom{k}{2}} \\ &= a \left(\frac{k}{\ln k}\right)^2 - \frac{2^6 a^4}{24} \left(\frac{k}{\ln k}\right)^2 - \left(\frac{eak}{(\ln k)^2}\right)^k k^{-(k-1)} \\ &= \left(a - \frac{8}{3} a^4\right) \left(\frac{k}{\ln k}\right)^2 - \frac{k(ea)^k}{(\ln k)^{2k}}. \end{aligned}$$

Fix $a' \in \left(0, \sqrt[3]{\frac{3}{8}}\right)$, and choose a to be as close to a' as possible, so that n is an integer. Since $\left(\frac{k}{\ln k}\right)^2 \rightarrow \infty$, we have $a \rightarrow a'$ as $k \rightarrow \infty$.

Note the last term above goes to 0 as $k \rightarrow \infty$ because $\frac{k}{(\ln k)^k} \rightarrow 0$ and $\left(\frac{ea}{\ln k}\right)^k \rightarrow 0$. Since a converges to a' as $k \rightarrow \infty$, for large k , $a - \frac{8}{3} a^4$ is bounded below by some $c > 0$. Hence $R(4, k) = \Omega\left(\left(\frac{k}{\ln k}\right)^2\right)$.

Problem 3 (*3.3, Independent set in 3-uniform hypergraph*)

Let G be a 3-uniform hypergraph with n vertices and $m \geq \frac{n}{3}$ edges. Take a random subset A by placing each vertex of G in A independently with probability p . Let $X = |A|$; then

$\mathbb{E}(X) = np$. Let Y be the number of edges in the subgraph induced by A . The probability that a given edge is in A is p^3 , since each of its vertices, independently, has probability p of being in A . Since there are m edges, by linearity of expectation, $\mathbb{E}(Y) = mp^3$.

Now $\mathbb{E}(X - Y) = np - mp^3$. There exists a subset A such that $X - Y \geq np - mp^3$. For each edge in the subgraph induced by A , choose one of its vertices. Upon removing these vertices, we get a set of at least $X - Y \geq np - mp^3$ vertices with no edges between them, i.e. an independent set of size at least $np - mp^3$.

Now take $p = \left(\frac{n}{3m}\right)^{\frac{1}{2}}$ (legal since $m \geq \frac{n}{3}$). Then we get an independent set of size at least

$$np - mp^3 = n \left(\frac{n}{3m}\right)^{\frac{1}{2}} - m \left(\frac{n}{3m}\right)^{\frac{3}{2}} = \frac{2n^{\frac{3}{2}}}{3\sqrt{3}\sqrt{m}}.$$

Problem 4 (3.4, Even directed cycle)

We show that in fact, the statement holds when each outdegree is at least $\log_2 n - \alpha \log_2 \log_2 n$ where $\alpha \in [0, \frac{1}{2})$.

Lemma 4.1: Let G be a directed graph, whose vertices are colored in two colors such that for every vertex v , there exists a vertex w such that there is an edge from v to w , and w is colored oppositely from v . Then G has a directed even cycle.

Proof. Choose any vertex v_1 . Once v_k is chosen, choose v_{k+1} to be adjacent to v_k along an outgoing edge, of the opposite color as v_k . At some point, a vertex will be repeated. Say that the first repeated vertex is v_k , and $v_k = v_j$, $j < k$. Then v_j, v_{j+1}, \dots, v_k is a simple cycle, since v_k is the first repeated vertex. Since the colors of vertices in the cycle alternate, it must have even length. \square

The following is Corollary 3.5.2 in the text.

Theorem 4.2:

$$m(d) = \Omega \left(2^d \left(\frac{d}{\ln d} \right)^{\frac{1}{2}} \right).$$

In other words, there exists C such that for every $d \geq 2$, any d -uniform hypergraph with at most $C2^d \left(\frac{d}{\ln d}\right)^2$ edges can be colored with two colors, so that no edge is monochromatic.

We will only need the weaker bound

$$m(d) = \Omega \left(2^d d^{\alpha'} \right), \text{ for any } \alpha' \in \left[0, \frac{1}{2} \right).$$

Given a directed graph G all of whose vertices have outdegree at least $\delta = \log_2 n - \alpha \log_2 \log_2 n$, for each vertex v let S_v be a set of vertices consisting of v and $\lceil \delta \rceil - 1$ vertices adjacent along an outgoing edge. Choose α' so that $\alpha < \alpha' < \frac{1}{2}$. Take C such that

$m(d) > C2^d d^{\alpha'}$ for $d \geq 2$. Let D be a positive constant less than 1. For large enough n ,

$$\begin{aligned} m(\lceil \delta \rceil) &\geq C2^{\lceil \delta \rceil} \lceil \delta \rceil^{\alpha'} \\ &\geq C2^\delta \delta^{\alpha'} \\ &= Cn(\log_2 n)^{-\alpha} (\log_2 n - \alpha \log_2 \log_2 n)^{\alpha'} \\ &\geq Cn(\log_2 n)^{-\alpha} D(\log_2 n)^{\alpha'} \\ &= CDn(\log_2 n)^{\alpha' - \alpha} \geq n. \end{aligned}$$

Consider the $\lceil \delta \rceil$ -uniform hypergraph whose vertices are the vertices of G and whose edges are the n sets S_v . By the above calculations, (for large enough n) there exists a coloring so that none of the S_v are monochromatic, i.e. so that each vertex leads to a vertex of a different color. By the lemma, G has an even cycle.

Problem 5 (4.1, $P(X = 0)$)

Let $p_k = P(X = k)$. By the Cauchy-Schwarz inequality,

$$\left(\sum_{k \geq 0} kp_k \right)^2 \leq \left(\sum_{k \geq 0} p_k \right) \left(\sum_{k \geq 0} k^2 p_k \right).$$

(Note the $k = 0$ terms for kp_k and $k^2 p_k$ are 0.) We rewrite this as

$$\begin{aligned} \mathbb{E}(X)^2 &\leq (1 - P(X = 0))\mathbb{E}(X^2) \\ \implies P(X = 0)\mathbb{E}(X^2) &\leq \mathbb{E}(X^2) - \mathbb{E}(X)^2 \\ \implies P(X = 0)\mathbb{E}(X^2) &\leq \text{Var}(X) \\ \implies P(X = 0) &\leq \frac{\text{Var}(X)}{\mathbb{E}(X^2)}. \end{aligned}$$

Problem 6 (4.2)

We show the inequality with $c = \frac{\sqrt{5}-2}{2}$.

Lemma 6.1: Let a_1, \dots, a_k be any real numbers, and $\varepsilon_1, \dots, \varepsilon_k$ independent random variables taking the values -1 and 1 each with probability $\frac{1}{2}$. Let $X = \varepsilon_1 a_1 + \dots + \varepsilon_k a_k$. Then

$$\text{Var}(X) = a_1^2 + \dots + a_k^2$$

and

$$P(|X| \leq 1) \geq 1 - (a_1^2 + \dots + a_k^2).$$

Proof. Let $X_i = \varepsilon_i a_i$. Note $\text{Var}(X_i) = \mathbb{E}((\varepsilon_i a_i)^2) = a_i^2$. Since the X_i are independent, we have

$$\text{Var}(X) = \text{Var}(X_1) + \dots + \text{Var}(X_k) = a_1^2 + \dots + a_k^2.$$

Clearly, $\mathbb{E}(X_i) = 0$ so $\mathbb{E}(X) = 0$. By Chebyshev's inequality with $\lambda = \frac{1}{\sqrt{a_1^2 + \dots + a_k^2}}$ and $\sigma = \sqrt{a_1^2 + \dots + a_k^2}$, we get

$$P(|X| \geq 1) = P(|X - \mathbb{E}(X)| \geq \lambda\sigma) \leq \frac{1}{\lambda^2} = a_1^2 + \dots + a_k^2.$$

Since $P(|X| \leq 1) \geq 1 - P(|X| \geq 1)$, this proves the second part. \square

Let $\lambda = \sqrt{5} - 2$. Note $\frac{1-\lambda^2}{8} = \frac{\lambda}{2}$ and $c = \frac{\lambda}{2}$. Consider two cases.

1. There exists a_i with $a_i^2 \geq \lambda$. Without loss of generality, $a_1^2 \geq \lambda$. Then by the lemma,

$$P(|\varepsilon_2 a_2 + \dots + \varepsilon_k a_k| \leq 1) \geq 1 - (a_2^2 + \dots + a_n^2) = a_1^2 \geq \lambda.$$

Since $|a_1| \leq 1$, given that $|\varepsilon_2 a_2 + \dots + \varepsilon_k a_k| \leq 1$, if ε_1 is such that $\varepsilon_1 a_1$ has opposite sign from $\varepsilon_2 a_2 + \dots + \varepsilon_k a_k$, then we also have $|X| \leq 1$. Thus

$$P(|X| \leq 1) \geq \frac{1}{2} P(|\varepsilon_2 a_2 + \dots + \varepsilon_k a_k| \leq 1) \geq \frac{\lambda}{2}.$$

2. There does not exist a_i with $a_i^2 \geq \lambda$. Let k be the greatest index so that

$$a_1^2 + a_2^2 + \dots + a_k^2 \leq \frac{1+\lambda}{2}.$$

(Note $k \geq 1$ since $a_1^2 < \lambda < \frac{1+\lambda}{2}$.) Let $A = a_1^2 + \dots + a_k^2$. By the maximality assumption, $a_1^2 + \dots + a_k^2 + a_{k+1}^2 > \frac{1+\lambda}{2}$. Since $a_{k+1}^2 < \lambda$, we conclude $A > \frac{1-\lambda}{2}$. Thus by the lemma,

$$\begin{aligned} P(|\varepsilon_1 a_1 + \dots + \varepsilon_k a_k| \leq 1) &\geq 1 - (a_1^2 + a_2^2 + \dots + a_k^2) = 1 - A. \\ P(|\varepsilon_{k+1} a_{k+1} + \dots + \varepsilon_n a_n| \leq 1) &\geq 1 - (a_{k+1}^2 + \dots + a_n^2) = A. \end{aligned}$$

By symmetry,

$$P(\varepsilon_{k+1} a_{k+1} + \dots + \varepsilon_n a_n \in [0, 1]) = P(\varepsilon_{k+1} a_{k+1} + \dots + \varepsilon_n a_n \in [-1, 0]) \geq \frac{A}{2}.$$

Now, noting $A \in [\frac{1-\lambda}{2}, \frac{1+\lambda}{2}]$ implies $A(1-A) \geq (\frac{1-\lambda}{2})(\frac{1+\lambda}{2}) = \frac{1-\lambda^2}{4}$,

$$\begin{aligned} P(|X| \leq 1) &\geq P(\varepsilon_1 a_1 + \dots + \varepsilon_k a_k \in [0, 1]) P(\varepsilon_{k+1} a_{k+1} + \dots + \varepsilon_n a_n \in [-1, 0]) \\ &\quad + P(\varepsilon_1 a_1 + \dots + \varepsilon_k a_k \in [-1, 0]) P(\varepsilon_{k+1} a_{k+1} + \dots + \varepsilon_n a_n \in [0, 1]) \\ &\geq P(\varepsilon_1 a_1 + \dots + \varepsilon_k a_k \in [0, 1]) \left(\frac{A}{2}\right) + P(\varepsilon_1 a_1 + \dots + \varepsilon_k a_k \in [-1, 0]) \left(\frac{A}{2}\right) \\ &= P(|\varepsilon_1 a_1 + \dots + \varepsilon_k a_k| \leq 1) \left(\frac{A}{2}\right) \\ &\geq (1-A) \left(\frac{A}{2}\right) \geq \frac{1-\lambda^2}{8} = \frac{\lambda}{2} \end{aligned}$$

as needed.

Problem 7 (4.3)

We need the following estimate. The proof is similar to Chebyshev's inequality.

Lemma 7.1: Let a_1, \dots, a_n be vectors in \mathbb{R}^2 , and $\varepsilon_1, \dots, \varepsilon_n$ be independently chosen to be ± 1 with probability $\frac{1}{2}$. Then

$$P\left(\left\|\sum_{k=1}^n \varepsilon_k a_k\right\| \geq R\right) \leq \frac{\sum_{k=1}^n \|a_k\|^2}{R^2}.$$

Proof. First we calculate $\mathbb{E}\left(\left\|\sum_{k=1}^n \varepsilon_k a_k\right\|^2\right)$. Let $a_k = (x_k, y_k)$. Now

$$\mathbb{E}\left(\left\|\sum_{k=1}^n \varepsilon_k a_k\right\|^2\right) = \mathbb{E}((\varepsilon_1 x_1 + \dots + \varepsilon_n x_n)^2 + (\varepsilon_1 y_1 + \dots + \varepsilon_n y_n)^2)$$

Expanding, noting that $\mathbb{E}(\varepsilon_i \varepsilon_j x_i x_j) = \mathbb{E}(\varepsilon_i \varepsilon_j y_i y_j) = 0$ for $i \neq j$ (since $\varepsilon_i \varepsilon_j$ has equal probability of being ± 1), the expected value equals

$$\mathbb{E}(\varepsilon_1^2 x_1^2) + \dots + \mathbb{E}(\varepsilon_n^2 x_n^2) + \mathbb{E}(\varepsilon_1^2 y_1^2) + \dots + \mathbb{E}(\varepsilon_n^2 y_n^2) = x_1^2 + \dots + x_n^2 + y_1^2 + \dots + y_n^2 = \sum_{k=1}^n \|a_k\|^2.$$

By Markov's inequality,

$$\begin{aligned} P\left(\left\|\sum_{k=1}^n \varepsilon_k a_k\right\| \geq R\right) &= P\left(\left\|\sum_{k=1}^n \varepsilon_k a_k\right\|^2 \geq R^2\right) \\ &\leq \frac{\mathbb{E}\left(\left\|\sum_{k=1}^n \varepsilon_k a_k\right\|^2\right)}{R^2} \\ &= \frac{\sum_{k=1}^n \|a_k\|^2}{R^2}. \end{aligned}$$

□

Lemma 7.2: Let a_1, \dots, a_n be vectors in \mathbb{R}^2 , all of length at most r . Then there exist $\varepsilon_1, \dots, \varepsilon_n \in \{-1, 1\}$ so that

$$\|\varepsilon_1 a_1 + \dots + \varepsilon_n a_n\| \leq \sqrt{2}r.$$

Proof. For $\varepsilon = (\varepsilon_1, \dots, \varepsilon_n) \in \{-1, 1\}^n$ and $i \neq j$, let $\mathcal{P}_{i,j}(\varepsilon)$ be the (possibly degenerate) parallelogram bounded by the 4 vertices $v \pm a_i \pm a_j$, where $v = \sum_{k \neq i,j; 1 \leq k \leq n} \varepsilon_k a_k$. Let $\mathcal{P} = \bigcup_{1 \leq i < j \leq n, \varepsilon \in \{-1, 1\}^n} \mathcal{P}_{i,j}(\varepsilon)$. Note that if $Q_1 Q_2 Q_3 Q_4$ is one of these parallelograms, then we have $Q_2 = Q_1 \pm 2a_i$ for some i , and similarly for the other adjacent pairs of vertices.

We claim that \mathcal{P} contains the origin. First we show that \mathcal{P} is convex. Let Q be a vertex on the boundary of \mathcal{P} , and QQ_1 and QQ_2 be edges of \mathcal{P} , with $Q_1 \neq Q_2$ (i.e. QQ_1, QQ_2 are edges of some $\mathcal{P}_{i,j}(\varepsilon)$.) As mentioned, $Q_1 = Q \pm 2a_i$ for some i and $Q_2 = Q \pm 2a_j$ for some j , for $i \neq j$. Suppose the directed angle $\angle Q_1 Q Q_2$ is in the range $[0^\circ, 180^\circ]$. Let

$Q' = Q \pm 2a_i \pm 2a_j$, where the two signs match the signs in Q_1 and Q_2 , respectively. Then Q_1Q_2Q' is one of the parallelograms, in particular, P contains the angle $\angle Q_1QQ_2$. This shows that \mathcal{P} has no reflex angle on the boundary. \mathcal{P} has a well-defined outer boundary that traces a convex polygon, and has no “holes” (because holes would cause reflex angles as well). Hence \mathcal{P} is convex. Since \mathcal{P} is clearly symmetric around the origin, it must contain the origin.

Hence we can take a parallelogram $\mathcal{P}_{i,j}(\varepsilon)$ containing the origin. Suppose by way of contradiction that all its vertices P_1, P_2, P_3, P_4 are at a distance greater than $\sqrt{2}r$ from the origin O . One of the angles $\angle P_1OP_2, \angle P_2OP_3, \angle P_3OP_4, \angle P_4OP_1$ is at least 90° , say WLOG $\angle P_1OP_2$. Then by Pythagorean's inequality $|P_1P_2|^2 \geq |OP_1|^2 + |OP_2|^2 > 2(\sqrt{2}r)^2$ so $|P_1P_2| > 2r$. But $|P_1P_2| = 2a_i$ for some i , and $a_i > r$, contradiction. Thus one of P_1, P_2, P_3, P_4 is at most a distance of $\sqrt{2}r$ from O , proving the lemma. \square

Back to the problem, let $i_0 = 0$, let i_1 be the largest integer so that $\|a_1\|^2 + \dots + \|a_{i_1}\|^2 \leq \frac{1}{20}$, let i_2 be the largest integer so that $\|a_{i_1+1}\|^2 + \dots + \|a_{i_2}\|^2 \leq \frac{1}{20}$, and so on. (Note $i_{j+1} > i_j$ because $\|a_i\|^2 \leq \frac{1}{100}$ for all i .) Suppose this divides the a_i into t groups. For $0 \leq j < t-1$, by the maximality assumption on i_{j+1} , $\|a_{i_j+1}\|^2 + \dots + \|a_{i_{j+1}}\|^2 + \|a_{i_{j+1}+1}\|^2 > \frac{1}{20}$; since $\|a_{i_{j+1}+1}\|^2 \leq \frac{1}{100}$, we conclude $\|a_{i_j+1}\|^2 + \dots + \|a_{i_{j+1}}\|^2 > \frac{1}{25}$. Thus we've divided the a_i into t groups, and in all of them except the last, the sum of the squares of the absolute values is in the interval $(\frac{1}{25}, \frac{1}{20}]$. This shows $t \leq 25$.

By Lemma 7.1,

$$P \left(\left\| \sum_{k=i_j+1}^{i_{j+1}} \varepsilon_k a_k \right\| \geq \frac{1}{\sqrt{18}} \right) \leq 18 \sum_{k=i_j+1}^{i_{j+1}} \|a_k\|^2 \leq \frac{18}{20}$$

so

$$P \left(\left\| \sum_{k=i_j+1}^{i_{j+1}} \varepsilon_k a_k \right\| \leq \frac{1}{\sqrt{18}} \right) \geq \frac{1}{10}.$$

Thus the probability that $\left\| \sum_{k=i_j+1}^{i_{j+1}} \varepsilon_k a_k \right\| \leq \frac{1}{\sqrt{18}}$ for each $0 \leq j < t$ is at least $\frac{1}{10^t} \geq \frac{1}{10^{25}}$.

Let $v_j = \sum_{k=i_j+1}^{i_{j+1}} \varepsilon_k a_k$. Let S be the set of $\varepsilon = (\varepsilon_1, \dots, \varepsilon_n)$ such that $\|v_j\| \leq \frac{1}{\sqrt{18}}$ for all j ; call $v = \sum_{k=1}^n \varepsilon_k a_k$ the vector associated to ε . We say two vectors $\varepsilon, \varepsilon'$ are equivalent if

$$(\varepsilon_{i_j+1}, \dots, \varepsilon_{i_{j+1}}) = \pm(\varepsilon'_{i_j+1}, \dots, \varepsilon'_{i_{j+1}})$$

for each j . This divides S into equivalence classes, each containing 2^t elements. Note that the vectors associated to the 2^t elements in the equivalence class of ε are in the form $\omega_0 v_0 + \dots + \omega_{t-1} v_{t-1}$ where $\omega_i = \pm 1$. Since $\|v_i\| \leq \frac{1}{\sqrt{18}}$, by Lemma 7.2 there exists a choice of $\omega_0, \dots, \omega_{t-1}$ so that $\|\omega_0 v_0 + \dots + \omega_{t-1} v_{t-1}\| \leq \frac{\sqrt{2}}{\sqrt{18}} = \frac{1}{3}$ (in fact, there exist two choices, by symmetry). Hence in each equivalence class C of S , at least 2 of the 2^t elements of C have

associated vectors with absolute value at most $\frac{1}{3}$. Thus

$$\begin{aligned} P\left(\left\|\sum_{k=1}^n \varepsilon_k a_k\right\| \leq \frac{1}{3}\right) &\geq P\left(\bigwedge_{j=0}^{t-1} \left(\left\|\sum_{k=i_j+1}^{i_{j+1}} \varepsilon_k a_k\right\| \leq \frac{1}{\sqrt{18}}\right)\right) \cdot \frac{2}{2^t} \\ &\geq \frac{1}{10^{25}} \cdot \frac{1}{2^{24}}. \end{aligned}$$