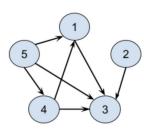
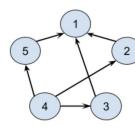


The trust relationships form a graph. Each trust pair, [a, b] represents a **directed** edge going from [a, b] to [b, b].

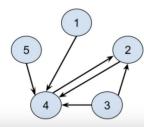
For example, with N=5 and trust = [[1,3],[2,3],[4,1],[5,3],[5,1],[5,4]] , we get the following graph. Who is the town judge?

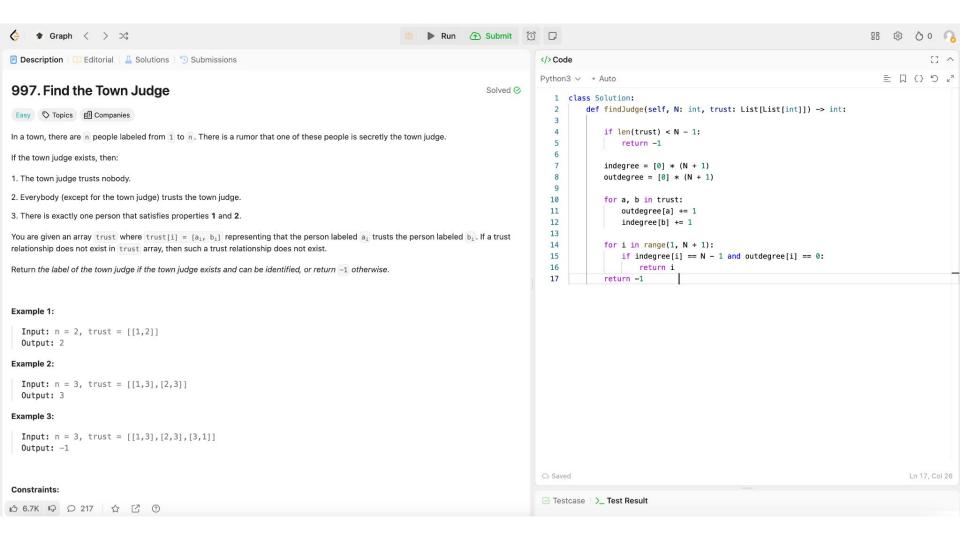


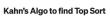
What about this example, with trust = [[2,1],[3,1],[4,2],[4,3],[4,5],[5,1]] ?



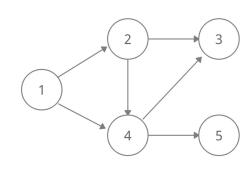
And what about this example, with trust = [[1,4],[2,4],[3,2],[3,4],[4,2],[5,4]] ?



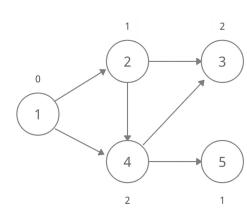




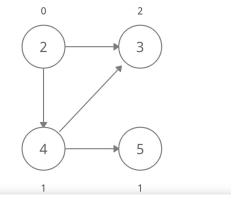
Consider the following graph.



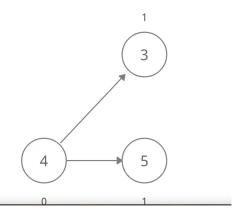
According to Kahn's Algo, First, we need to find the indegree of all the vertices. $\,$



The graph will look like this after removing the vertex $\overline{\mathbbm{1}}$ and its outgoing edges.



Delete the vertex $\,_2$ from the graph and all of its outgoing edges. Recalculate the indegree of the remaining vertices, the graph will look like this.

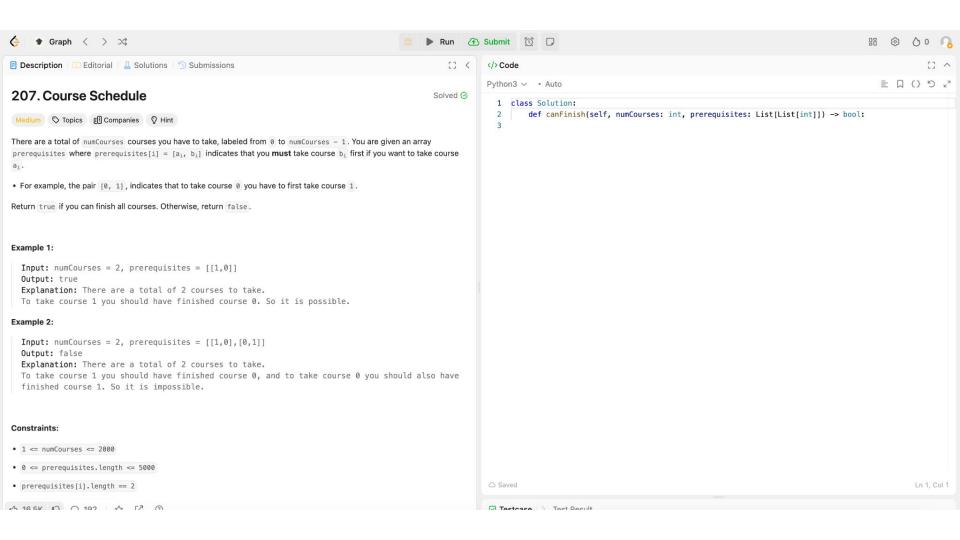


3

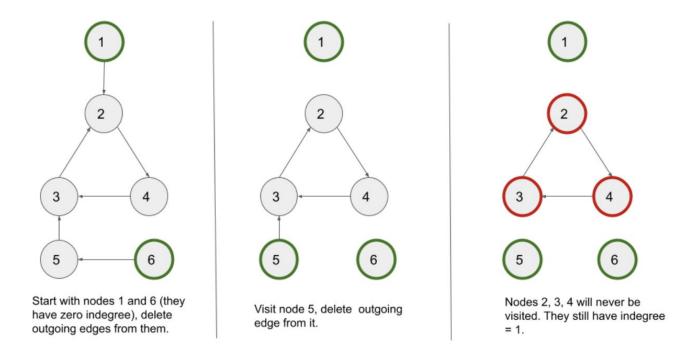


If we consider vertex $\ \ _3$ first, the topological order will become -

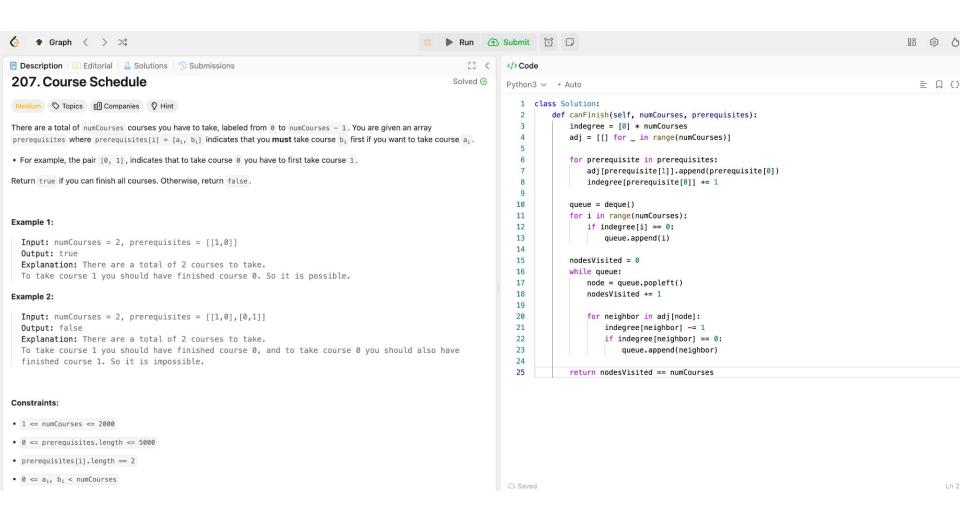
Topological order - 1 2 4 3 5

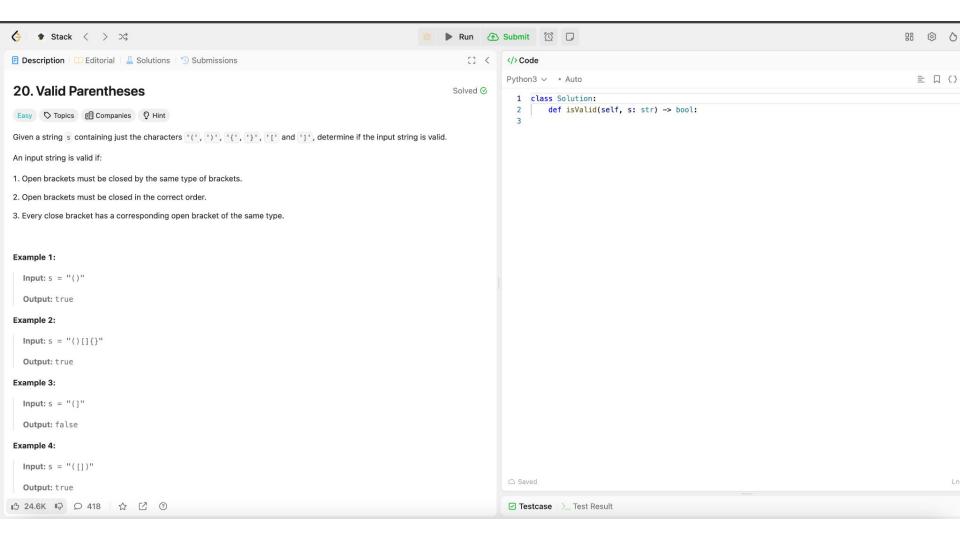


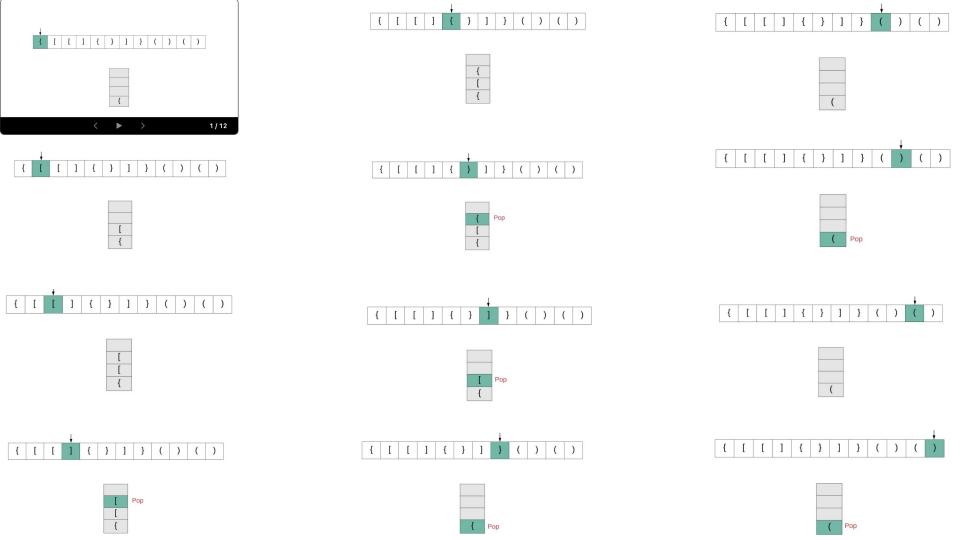
Let's perform Kahn's algorithm on a directed graph having a cycle. Here's a visual step-bystep representation of how it would work:

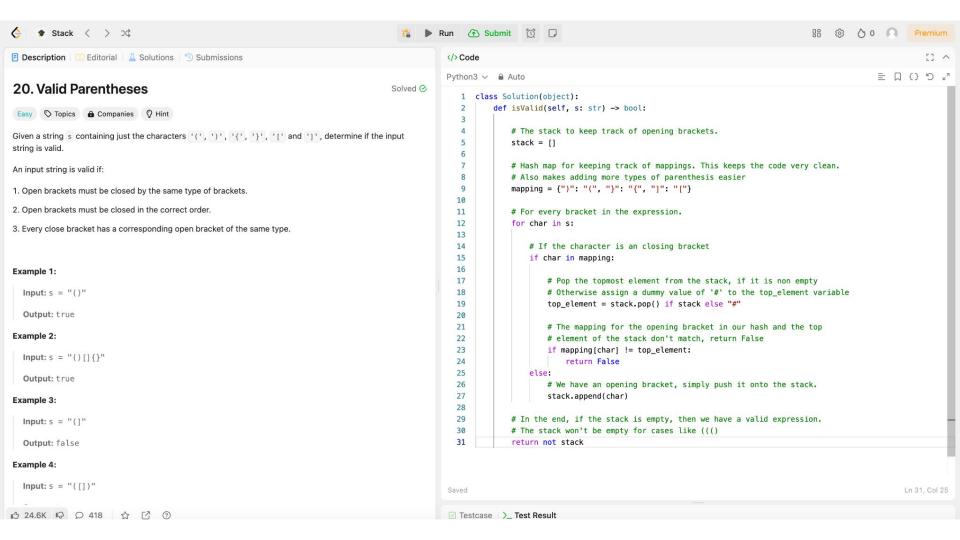


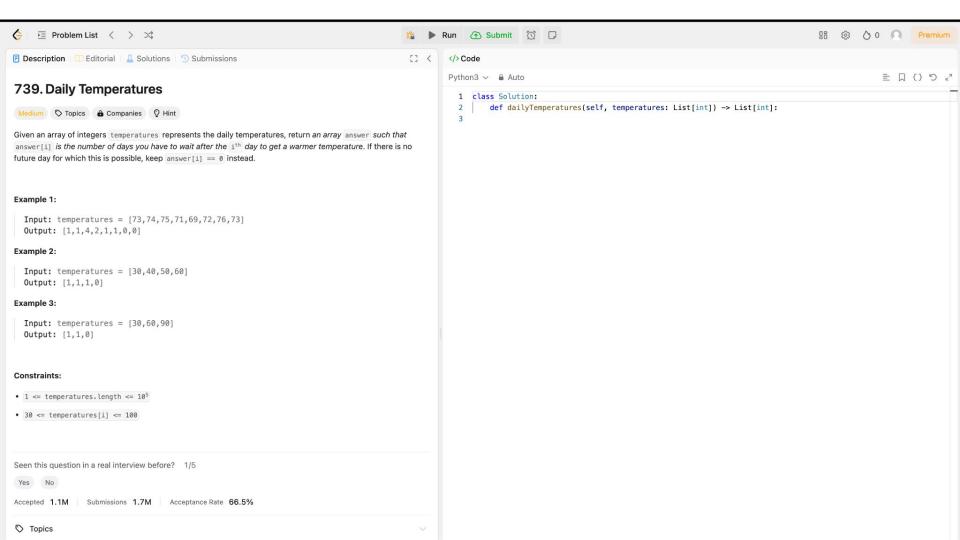
We can see that if there is a cycle, the indegree of nodes in the cycle cannot be set to 0 due to cyclic dependency. We are unable to visit the cycle's nodes. So, if the number of visited nodes is less than the total number of nodes in the graph, we have a cycle.

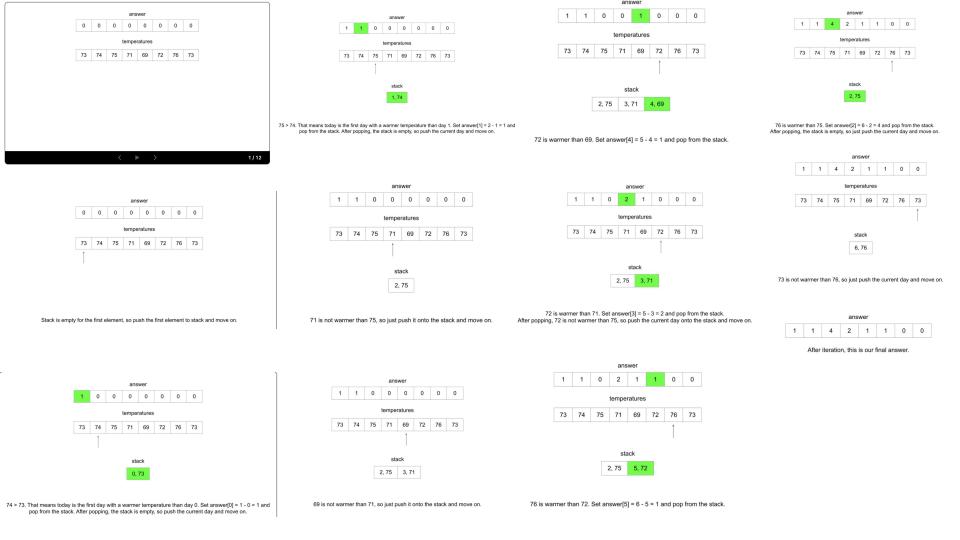


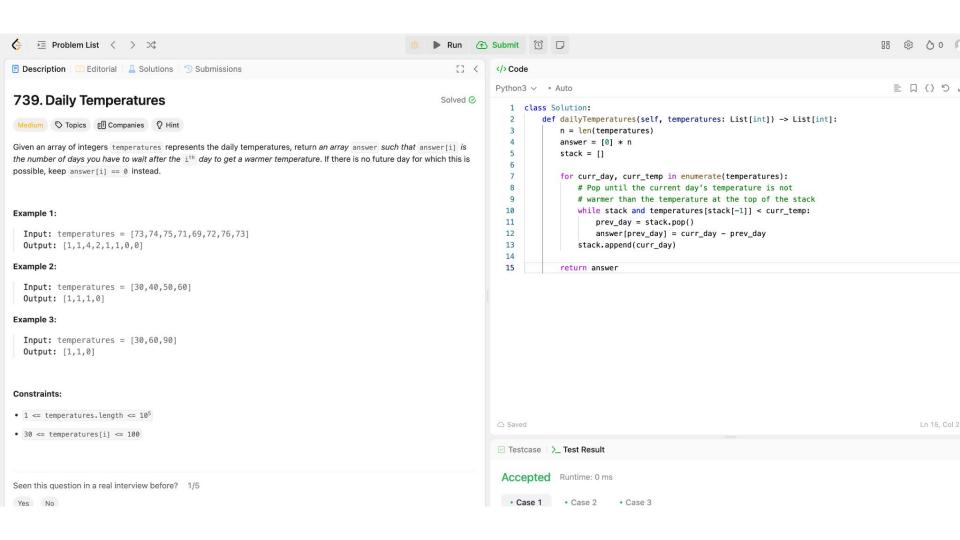


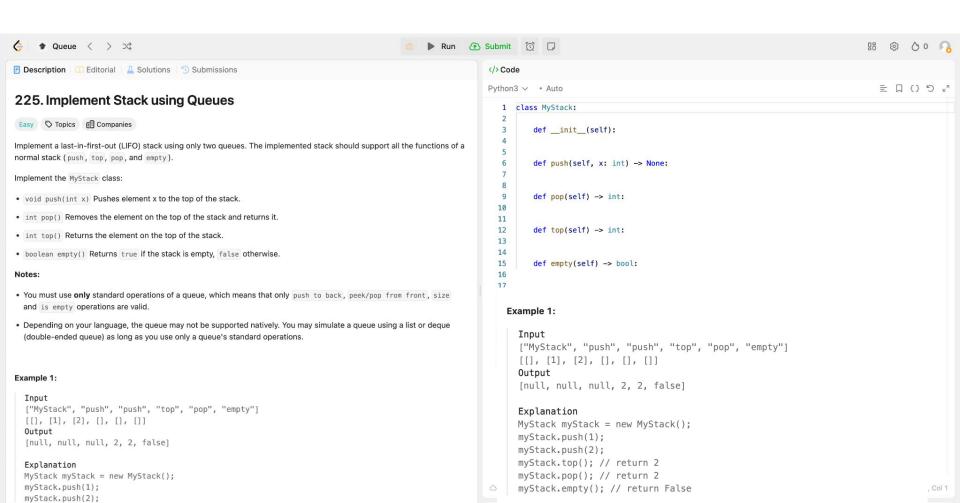














One-Queue Approach

- How it Works: Use a single queue and reorganize (rotate) its elements when pushing a new element.
- Time Complexity: O(n) for push, O(1) for pop, top, and empty.
- When to Use: Choose this approach when you expect pop and top operations to be more frequent than push.

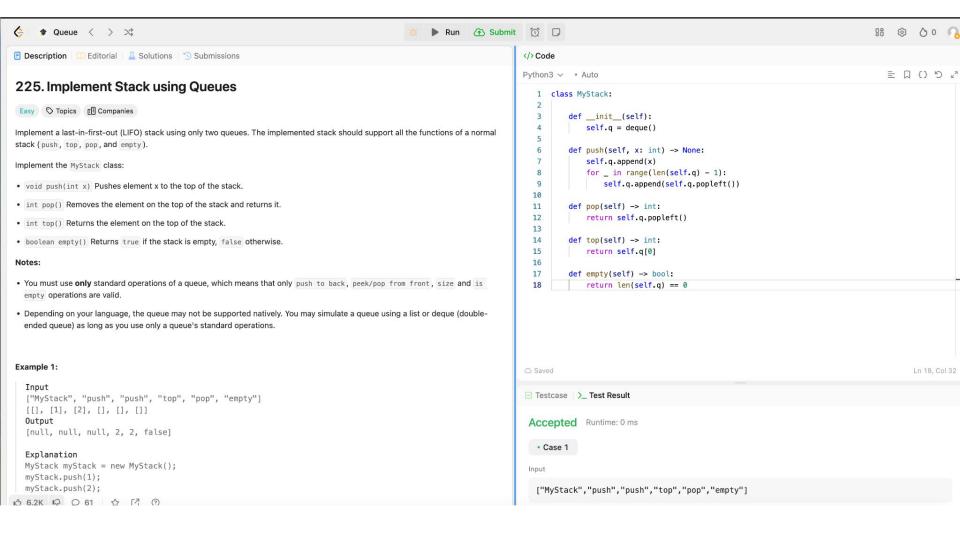
Push Operation

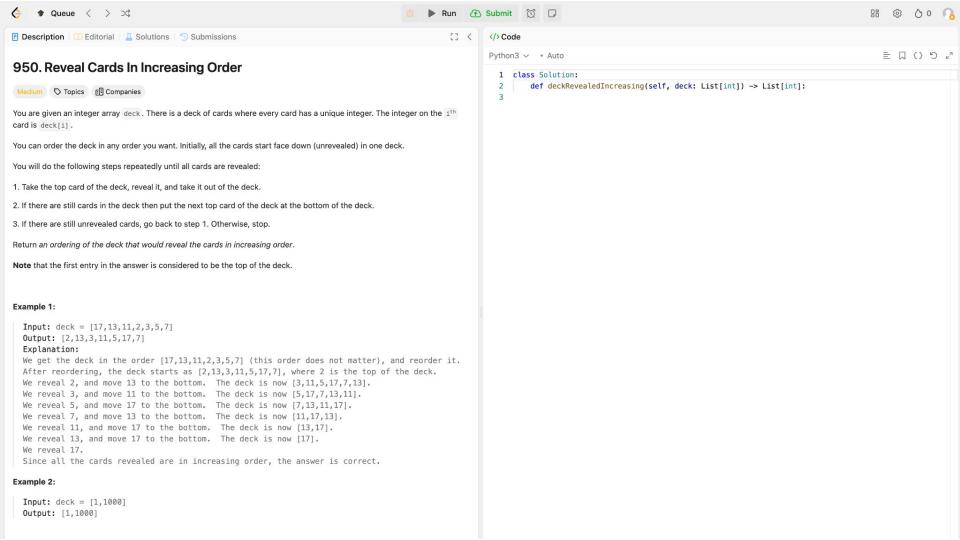
- Enqueue the new element at the back of q.
- To ensure that the last element can be accessed from the front, rotate the queue by dequeueing and enqueuing each element except the newly pushed element.

Pop and Top Operations

Simply dequeue from the front for pop and peek at the front for top

Col





		deck 2 3 5 7 11 13 17	deck sorted 2 3 5 7 11 13 17	deck sorted 2 3 5 7 11 13 17
deck	17 13 11 2 3 5 7	queue 2 3 4 5 6 1	queue 6 1 3 5	queue 1 5
deck sorted	2 3 5 7 11 13 17	0 1 2 3 4 5 6 result 2 0 3 0 0 0 0	0 1 2 3 4 5 6 result 2 0 3 0 5 0 7	0 1 2 3 4 5 6 result 2 13 3 11 5 0 7
deck sorted	2 3 5 7 11 13 17	deck 2 3 5 7 11 13 17	deck 2 3 5 7 11 13 17	deck sorted 2 3 5 7 11 13 17
queue	0 1 2 3 4 5 6	queue 4 5 6 1 3	queue 3 5 1	queue 5
result	0 1 2 3 4 5 6	0 1 2 3 4 5 6 result 2 0 3 0 5 0 0	result 2 0 3 11 5 0 7	0 1 2 3 4 5 6 result 2 13 3 11 5 17 7
		leadin 2 0 5 0 0 0		
deck				
deck sorted	2 3 5 7 11 13 17	deck sorted 2 3 5 7 11 13 17	deck sorted 2 3 5 7 11 13 17	return
	0 1 2 3 4 5 6	deck sorted 2 3 5 7 11 13 17 queue 4 5 6 1 3		return 0 1 2 3 4 5 6 result 2 13 3 11 5 17 7
sorted		sorted 2 3 3 7 11 13 17	sorted 2 3 5 7 11 13 17	0 1 2 3 4 5 6
gueue	0 1 2 3 4 5 6	queue 4 5 6 1 3 0 1 5 6	queue 3 5 1 0 1 2 3 4 5 6	0 1 2 3 4 5 6
sorted queue result	0 1 2 3 4 5 6 0 1 2 3 4 5 6 2 0 0 0 0 0 0	queue 4 5 6 1 3 0 1 2 3 4 5 6 result 2 0 3 0 5 0 0	queue 3 5 1 0 1 2 3 4 5 6 result 2 0 3 11 5 0 7	0 1 2 3 4 5 6

