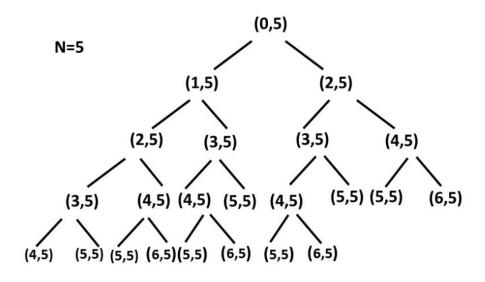


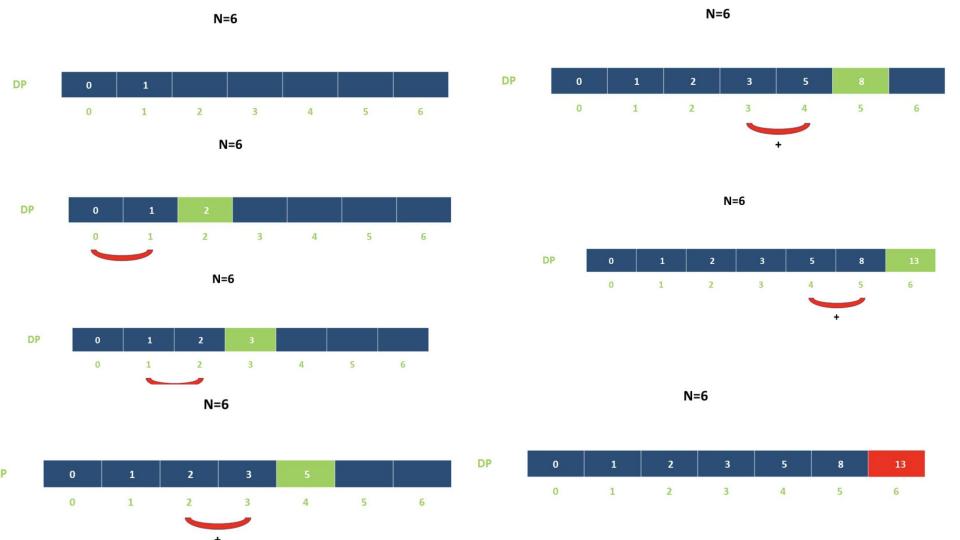
Approach 1: Brute Force

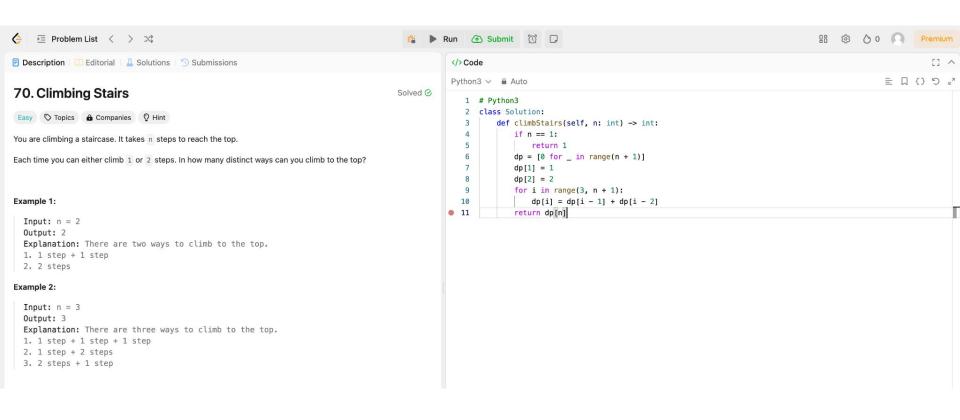
```
# Python3
class Solution:
    def climbStairs(self, n: int) -> int:
        return self.climb_Stairs(0, n)
    def climb_Stairs(self, i: int, n: int) -> int:
        if i > n:
            return 0
        if i == n:
            return 1
        return self.climb_Stairs(i + 1, n) + self.climb_Stairs(i + 2, n)
```

Time complexity : $O(2^n)$.



Number of Nodes = $O(2^n)$





```
E3 <
                                                                                                         </>Code
300. Longest Increasing Subsequence
 Medium ♥ Topics ② Companies
Given an integer array nums, return the length of the longest strictly increasing subsequence.
                                                                                                           5
                                                                                                           6 };
Example 1:
  Input: nums = [10,9,2,5,3,7,101,18]
  Output: 4
  Explanation: The longest increasing subsequence is [2,3,7,101], therefore the length is 4.
Example 2:
  Input: nums = [0,1,0,3,2,3]
  Output: 4
Example 3:
  Input: nums = [7,7,7,7,7,7,7]
  Output: 1
Constraints:
• 1 <= nums.length <= 2500
• -10^4 <= nums[i] <= 10^4
Follow up: Can you come up with an algorithm that runs in O(n \log(n)) time complexity?
```

```
1 □ ( ) □ □
C++ v · Auto
 1 class Solution {
 2 public:
        int lengthOfLIS(vector<int>& nums) {
```

nums

| | 10 | 9 | 2 | 5 | 3 | 7 | |
|---|----|---|---|---|---|----|--|
| 1 | | | _ | | | ٠. | |

dp

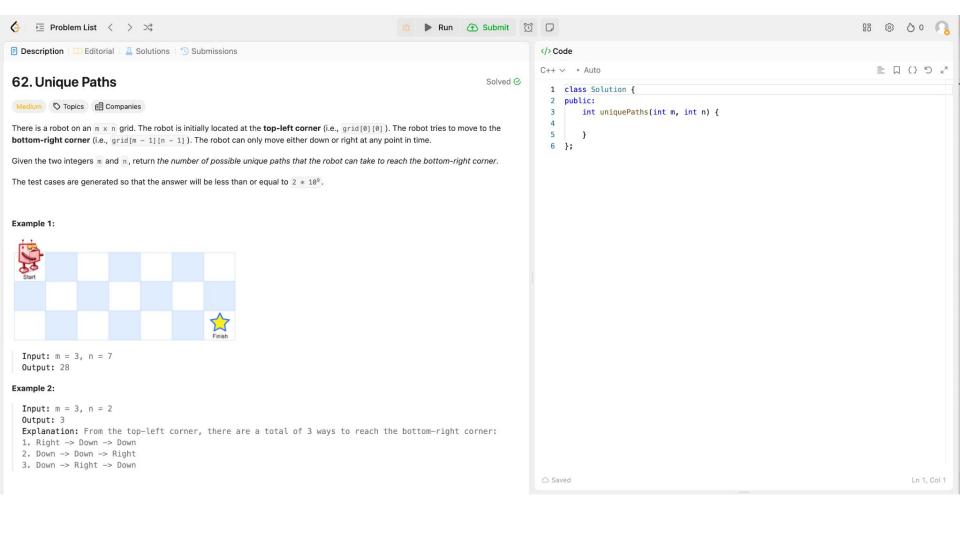
| 1 | 1 | 1 | 1 | 1 | 1 |
|-----|---|---|---|---|---|
| - 5 | 0 | | | - | |

dp[i] represents the length of the longest increasing subsequence that ends at index i

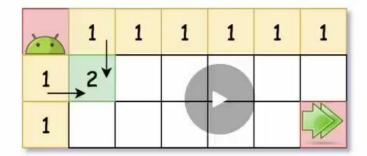
A Framework to Solve Dynamic Programming Problems

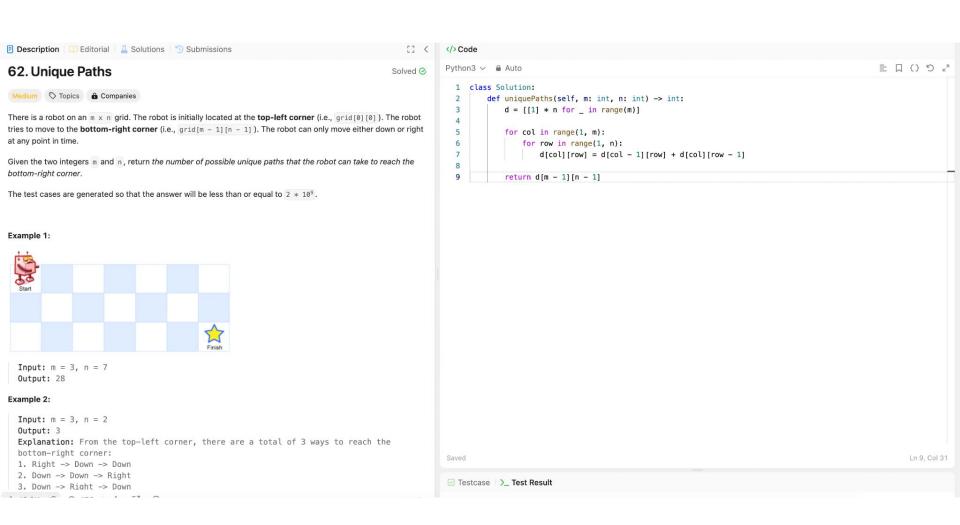
- 1. First, we need some function or array that represents the answer to the problem from a given state
- 2. Second, we need a way to transition between states, such as dp[5] and dp[7]. This is called a recurrence relation
- 3. The third component is: we need a base case. For this problem, we can initialize every element of dp to 1, since every element on its own is technically an increasing subsequence.

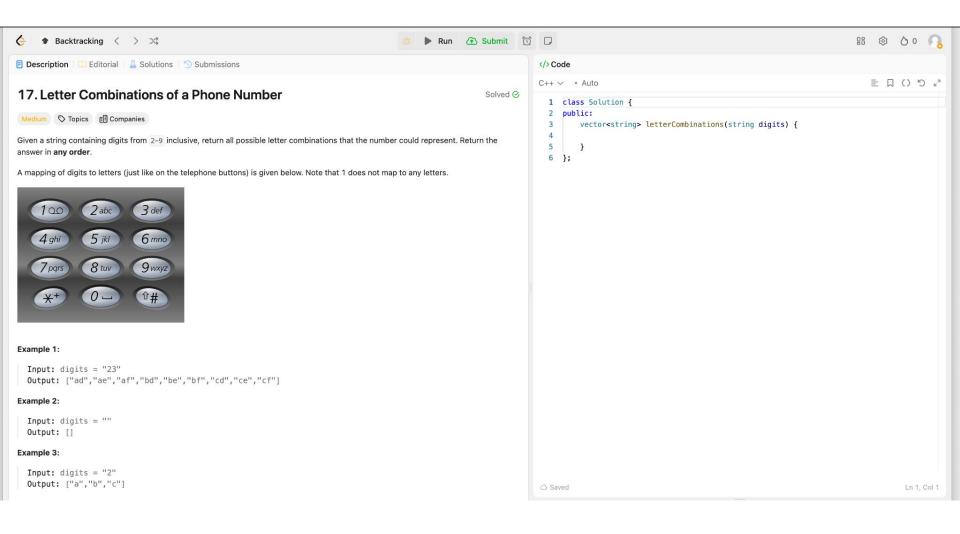


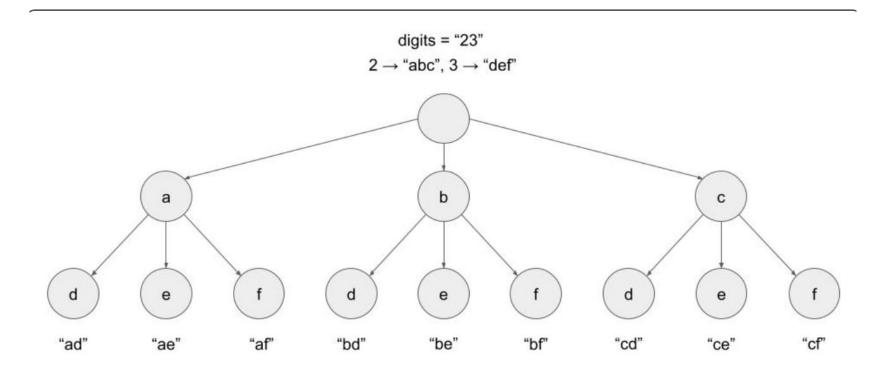


$$d[1][1] = d[0][1] + d[1][0] = 2$$

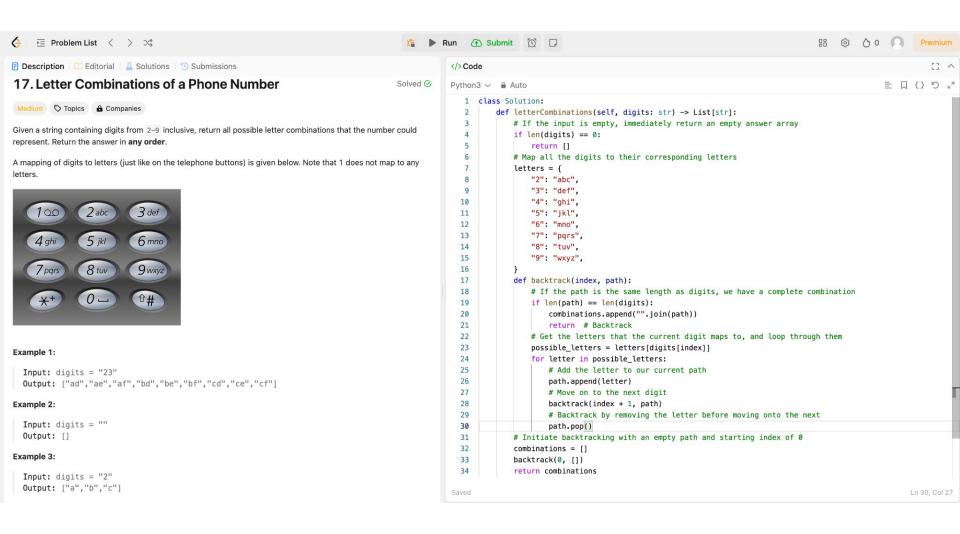


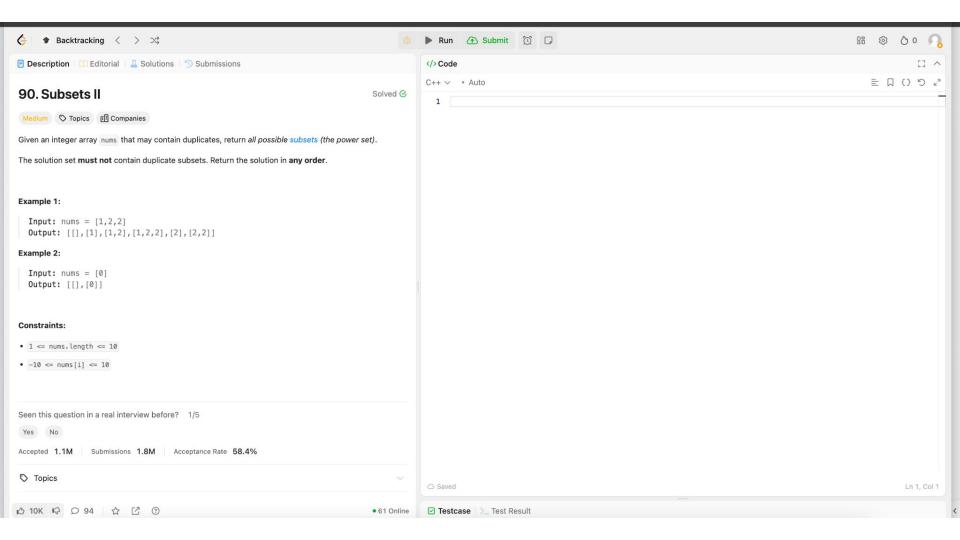


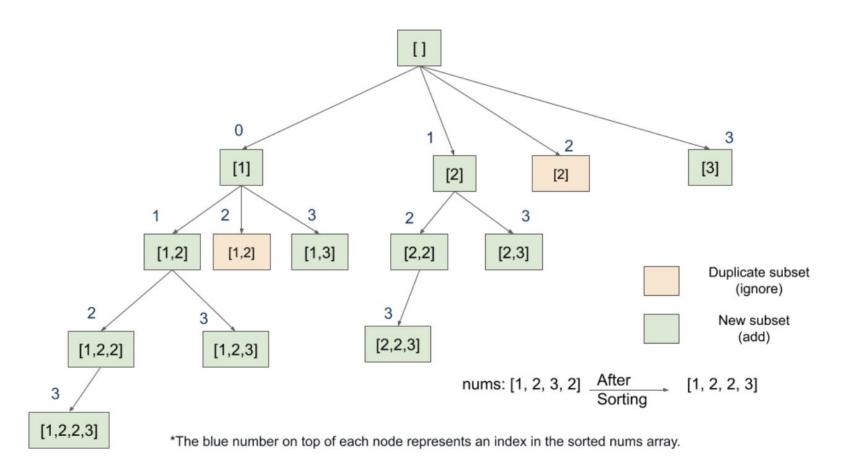


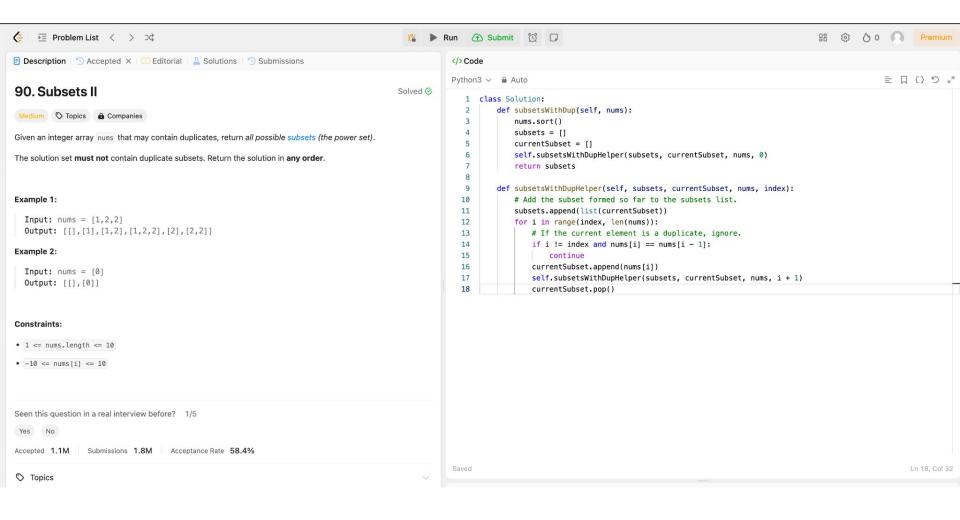


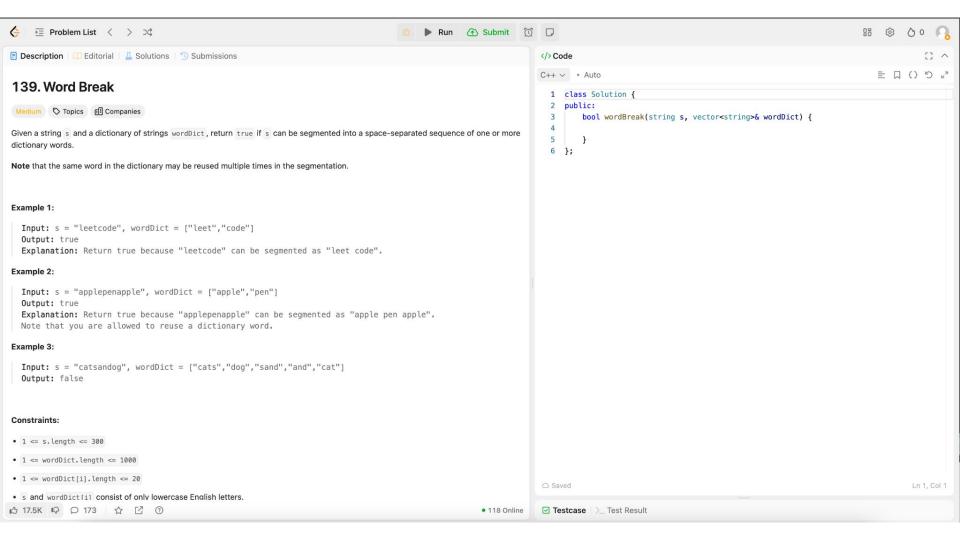
| "ad" | "ae" | "af" | "bd" | "be" | "bf" | "cd" | "ce" | "cf" | |
|------|------|------|------|------|------|------|------|------|---|
| | | | | | | | | | - |



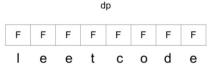








1.

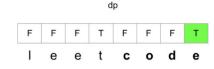


s = "leetcode" wordDict = ["leet", "code"]

The table holds boolean values T (true) and F (false). Under each element is the corresponding letter from the input string.

dp(i) indicates if it is possible to use words from wordDict to build the input string up to index i.

3.



s = "leetcode" wordDict = ["leet", "code"]

.

dp

The next time that the first criteria is satisfied is at index 7. The word "code" can end here.

The second criteria is also satisfied - the word "code" starts at index 4. The index before that, 3, is also "true". Therefore, dp(7) = true.

Since this is also the last index, the answer to the problem is "true".

s = "leetcode" wordDict = ["leet", "code"]

The first criteria is that a word from wordDict can end at s[i]. The first occurrence of this is at index 3. The word "leet" can end here.

The second criteria is that the input string is formable up to the point before where the current string starts.

Since this is the first word being used, the criteria is also met.

Both criteria are satisfied, dp(3) = true.

