

Stacks

linear data structure hai jo LIFO principle follow karta hai.

* (Last in → First out)

→ last wala element phle bhar.

basic operations :-

① push(x)	→ add	(TC) ↴ O(1)
② pop(x)	→ remove top()	O(1)
③ top(x)	→ dekhna top()	O(1)
④ isEmpty(x)	→ bohot T/F	O(1)

Sx → plate stack

Undo/Redo

browser back button

function calls (Call stack)

Stack implementation

stack<int>s;
s.push(10);
s.pop();
int x = s.top();



Use Cases :

- ① Expression problems
 → Valid parenthesis
 → infix → postfix
 → Evaluate postfix
- ② Next greater/Smaller
 → while s. empty if
 s.top() <= cur

③ Monotonic Stack

- increasing {nearest greater}
- decreasing {smallest}

④ Largest Area/Range

- Maximal Rectangle
 → Largest Rectangle in histogram

⇒ Widely used in function calls, recursion, expression evaluation, Undo-redo operations and monotonic stack problem like greater/smaller element;

Can be implemented using arrays, linked list or STL containers & they play a crucial role when a problem requires reverse, backtracking or previous state.

Although linear DS, But proves (Monotonicity)
reduces TC from $O(n^2)$ to $\underline{O(n)}$

LC \rightarrow 20

Valid parenthesis :

- open bracket must be closed by same type of closing bracket
- must be in same order
- closing bracket has same opening bracket as that of closing.

Ex: ()[]{}
 ↑ if opening bracket \rightarrow push.

 else check karo. same hai ki ni hai

{ Stack se?

→ Implement back:

→ Ex: I/P: a#b, b

given 2 strings $s = a \# b$, and $t = b$.

→ As we get '#' pop the recent element
 → make the string

and check if both are equal or not

→ Make a string Great (LC \rightarrow 1544)

$s[i]$ & $s[i+1]$ where $0 < i < s.\text{length}() - 2$

where $s[i] == \text{Upper}(s[i])$
 $(as s[i+1])$

simple get the element push in stack, check top of

remove the element

"leetcode"

= Stack

e
d
o
c
t
e
e
e

→ "leetode"

"abABCC"

a
b
B

l
A
B

" "
empty string

difference b/w a uppercase & lowercase char = 32.

Stack<char> st;

for (int i = 0; i < s.size(); i++)

 char ch = s[i];

 if (!st.empty() && abs(st.top() - ch) == 32)

 // bad pair;

 st.pop();

 else

 st.push(ch);

Since adjacent
dependency check
LIFO is best fit

} making / filling
stack

// string creation;

while (!st.empty())

 char c = st.top()

 st.pop();

 ans = ans + c;

 reverse(ans.begin(), ans.end());

return ans;

};

} string creation

LC → 1021

Remove outer parenthesis

I/P → "((())())"

→ decomposition → "(()())" + "(()")"

after removal becomes

"()() + ()"

→ So, answer → ()()(),

Intuition :

→ Check for the s.top()

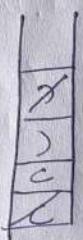
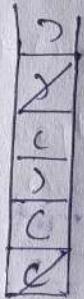
if same as (ch) pop the element;
and push the "ch"

∴ "((())())()

if (st.top() == ch)

st.pop();

st.push(ch);



→ ()()() Hence
O/P

→ Since adjacency is to checked

LIFO is the best used/ comfortable to use

“() (X) (X X) (X) (X) X X ”

→ ()()()()()

```

stack<char> st;
string ans = "";
for (char ch : s)
    if (ch == '(')
        if (!st.empty()) // not outer // parens pair
            ans = ans + ch;
        st.push(ch);
    else {
        // ab closing agya
        st.pop();
        if (!st.empty()) // outer
            ans = ans + ch;
    }
    return ans;
}

```

ans = C)C

LC → 682 → Baseball Game

$[5 \ 2 \ C \ D \ +]$

→ ⑤ ②

20
2
5

$$5 + 2 \Rightarrow 7$$

C → Ø D

st.push(2 * st.top());

$[5 \ -2 \ 4 \ C \ D \ 9 \ +]$

$\boxed{5 \ | -2} \times$

5	-2	4	9
			5

$\boxed{5 \ | 3}$

$$\text{sum} = 4 + 9$$

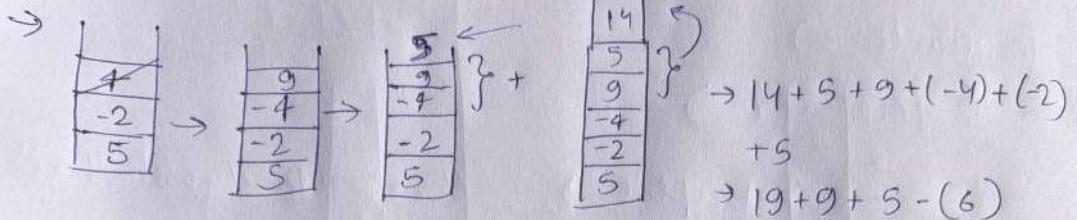
→ 13

3

8

8 -

$[S, -2, 4, C, D, 9, +, +]$



$$\rightarrow 14 + 5 + 9 + (-4) + (-2)$$

+ 5

$$\rightarrow 19 + 9 + 5 - (6)$$

→ 27

at sum = 0;

Stack < int > st;

for (auto c : operations)

if ($C == '+' \& C == 'D' \& C == 'C'$)

// Matlab integer hai;

int x = stoi(c);

st.push(x);

else if ($c == '+'$)

a = st.top();

st.pop();

b = st.top();

st.pop();

st.push(b); // phale (b) kyonki wo nikla phale

st.push(a);

st.push(a+b);

else if ($c == 'D'$)

int x = st.top();

st.push(x * 2);

else st.pop();

// Stack is ready, Sum up Karo

while (!st.empty())

sum = sum + st.top();

st.pop();

return sum;