

# Stacks

linear data structure hai jo LIFO principle follow Krta hai,

★ (Last in  $\rightarrow$  First out)

$\rightarrow$  last wala element phle bahar,

basic operations:

- |              |                            |        |
|--------------|----------------------------|--------|
| ① push(x)    | $\rightarrow$ add          | $O(1)$ |
| ② pop(x)     | $\rightarrow$ remove top() | $O(1)$ |
| ③ top(x)     | $\rightarrow$ dekhna top() | $O(1)$ |
| ④ isempty(x) | $\rightarrow$ bool T/F     | $O(1)$ |

Ex  $\rightarrow$  plate stack  
Undo/Redo  
browser back button  
function calls (call stack)

## Stack implementation

stack < int > s;  
s.push(10);  
s.pop();  
int x = s.top();

## Use Cases

- ① Expression problems  
 $\rightarrow$  Valid parenthesis  
 $\rightarrow$  infix  $\rightarrow$  postfix  
 $\rightarrow$  Evaluate postfix

- ② Next greater/  
Smaller

$\rightarrow$  while s.empty() {  
s.top() <= cur

- ③ Monotonic Stack

$\rightarrow$  increasing {nearest greater}  
 $\rightarrow$  decreasing {smallest}

- ④ Largest Area/range

$\rightarrow$  Maximal Rectangle  
 $\rightarrow$  largest Rectangle in histogram

$\Rightarrow$  Widely used in function calls, recursion, expression evaluation, undo-redo operations and monotonic stack problem like greater/smaller element;

Can be implemented using arrays, linked list or STL containers & they play a crucial role when a problem requires reverse, backtracking or previous state.



Although linear DS, But power (Monotonicity)  
reduces TC from  $O(n^2)$  to  $O(n)$

LC  $\rightarrow 20$

Valid parenthesis :

- $\rightarrow$  open bracket must be closed by same type of closing bracket
- $\rightarrow$  must be in same order
- $\rightarrow$  closing bracket has same opening bracket as that of closing.

eg  $\rightarrow$   $()[]\{\}$

$\rightarrow$  if opening bracket  $\rightarrow$  push.

$\leftarrow$  Stack set

else check karo, same hai ki ni hai

$\rightarrow$  Implement back:

$\rightarrow$  eg  $\rightarrow$  I/P  $\rightarrow a\#b, b$

given 2 strings  $s = a\#b$ , and  $t = b$ .

$\rightarrow$  As we get '#' pop the recent element  
 $\rightarrow$  make the string

and check if both are equal or not

$\rightarrow$  Make a string great (LC  $\rightarrow 1544$ )

$s[i]$  &  $s[i+1]$  where  $0 < i < s.length() - 2$

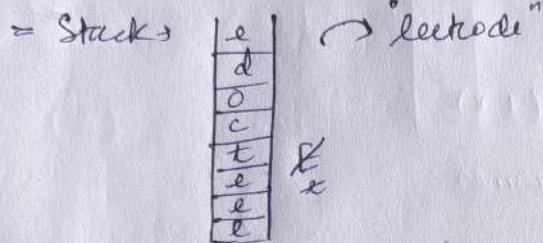
where  $s[i] == \text{upper}(s[i])$   
(as  $s[i+1]$ )

simple get the element push in stack, check top 04

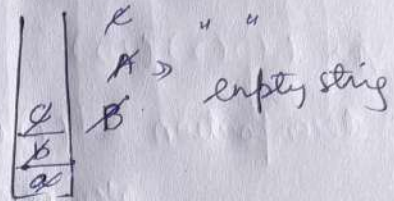


remove the element

"leFetcode"



"abABcC"



# difference b/w a uppercase & lowercase char = (32)

stack <char> st;

for (int i = 0; i < s.size(); i++)

↳ char ch = s[i];

if (!st.empty() && abs(st.top() - ch) == 32)

↳ // bad pair;  
st.pop();

else ↳  
st.push(ch);

# Since adjacent dependancy check  
LIFO is best fit

} making / filling stack

// string creation;

while (!st.empty())

↳ char c = st.top();

st.pop();

ans = ans + c;

↳

reverse(ans.begin(), ans.end());

return ans;

↳

} string creation



LC → 1021

Remove outer parenthesis

Ex: "(())(())"

→ decomposition ⇒ "(())" + "(())"

after removal becomes

"()() + ()"

→ So, answer ⇒ "()()()",

Intuition:

→ Check for the s.top()

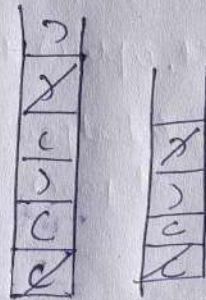
if same as (ch) pop the element;  
and push the "ch"

Ex: "(())(())"

if (st.top() == ch)

st.pop();

st.push(ch);



→ "()()()" Hence  
O/P

→ Since adjacency is to be checked  
LIFO is the best used / comfortable to use

Ex: "(())(())(())(())(())(())"

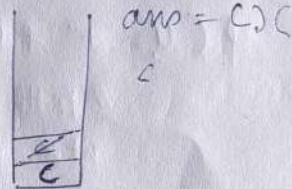
→ "()()()()()()"



```

stack<char> st;
string ans = "";
for (char ch : s)
{
    if (ch == '(')
    {
        if (!st.empty()) // not enter // pare & hai
        {
            ans = ans + ch;
            st.push(ch);
        }
        else // ab closing agya → (XL)
        {
            st.pop();
            if (!st.empty()) // enter
            {
                ans = ans + ch;
            }
        }
    }
    return ans;
}

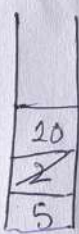
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LC → 682 → Baseball Game

["5", "2", "C", "D", "+"]

→ ⑤ ②

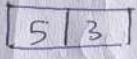
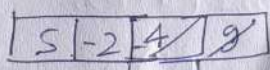
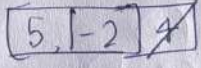


→ 5 + 2 ⇒ ⑦

C → ~~②~~ ⑦

st.push(2<sup>nd</sup> st. top());

["5", "-2", "4", "C", "D", "9", "+", "+"]



sum = 4 + 9

⇒ 13

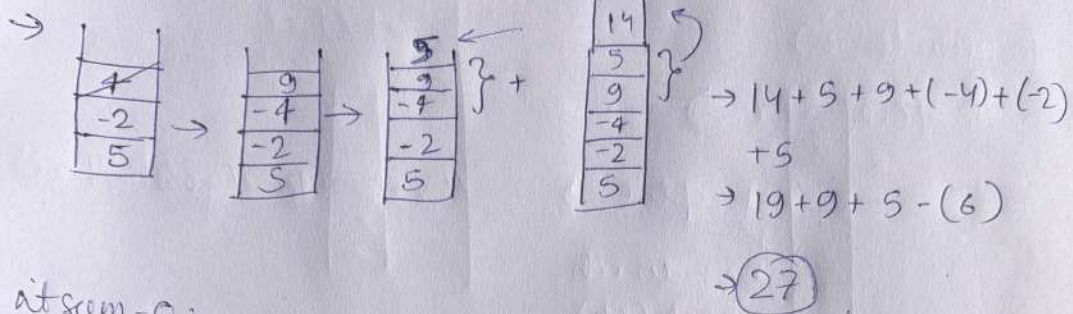
3

8

8



[5, -2, 4, C, D, 9, +, +]



int sum = 0;

Stack<int> st;

for(auto c: operations)

if (c == "+" || c == "D" || c == "C")

// mtlb integer hai;

int x = stoi(c);

st.push(x);

else if (c == "+")

a = st.top();

st.pop();

b = st.top();

st.pop();

st.push(b);

st.push(a);

st.push(a+b);

else if (c == "D")

int x = st.top();

st.push(x\*2);

else st.pop();

// phere (b) kizuki me nikla phere

// stack is ready, sumup kardo

while (!st.empty())

sum = sum + st.top();

st.pop();

return sum;