算法与数据结构体系课程

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快速排序法

快速排序 Quick Sort

 4
 6
 2
 3
 1
 5
 7
 8

 2
 3
 1
 4
 6
 5
 7
 8

4

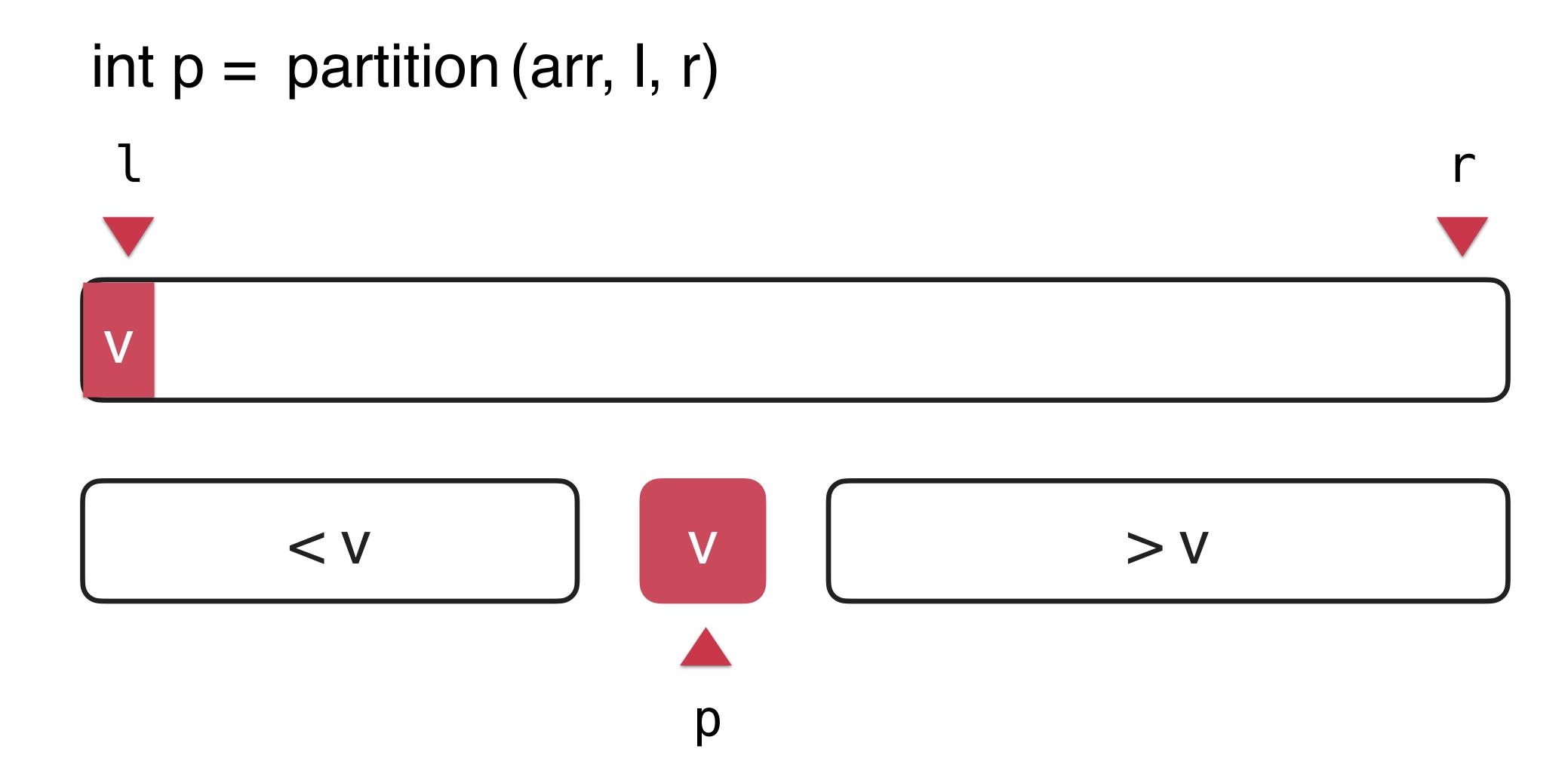
> 4

快速排序 Quick Sort



partition

快速排序 Quick Sort



快速排序法

int p = partition (arr, I, r)

注意递归函数的"宏观"语意

QuickSort(arr, l, r)

对 arr 的 [I, r] 部分排序

```
QuickSort(arr, l, r){
    if(l >= r) return;
    int p = partition(arr, l, r);
   // 对 arr[l, p - 1] 进行排序
   QuickSort(arr, l, p - 1);
   // 对 arr[p + 1, r] 进行排序
    QuickSort(arr, p + 1, r);
```

快速排序法

```
QuickSort(arr, l, r){
   if(l >= r) return; 水解最基本问题
   int p = partition(arr, l, r); 如何partiton?
   // 对 arr[l, p - 1] 进行排序
                                     把原问题转化成
   QuickSort(arr, l, p - 1);
                                     更小的问题
   // 对 arr[p + 1, r] 进行排序
   QuickSort(arr, p + 1, r);
```

V

最基础的 partition

4 6 2 3 1 5 7 8

left:

right:

 4

 2

 3

 1

 5

 7

 8

left:

right: 6

3] [1] [5] [7] [8]

4

left: 2 3 1

right: 6

5

7

8

left: [2] [3] [1]

left: [2] [3] [1]

4

left: [2] [3] [1]

right: [6] [5] [7] [8]

4

left: [2] [3] [1]

right: [6] [5] [7] [8]

 $\begin{bmatrix} 2 \end{bmatrix} \begin{bmatrix} 3 \end{bmatrix} \begin{bmatrix} 1 \end{bmatrix} \begin{bmatrix} 4 \end{bmatrix} \begin{bmatrix} 6 \end{bmatrix} \begin{bmatrix} 5 \end{bmatrix} \begin{bmatrix} 7 \end{bmatrix} \begin{bmatrix} 8 \end{bmatrix}$

使用了额外空间,非原地排序

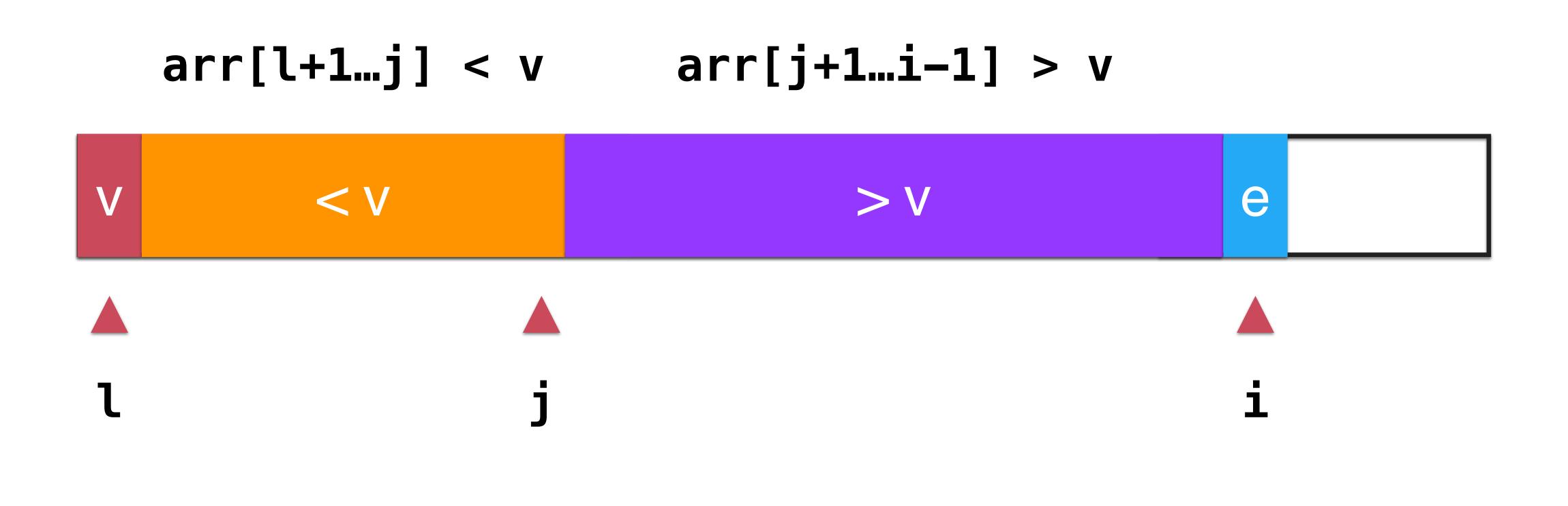
如何原地进行 partition?

arr[l+1...j] < v arr[j+1...i-1] > v

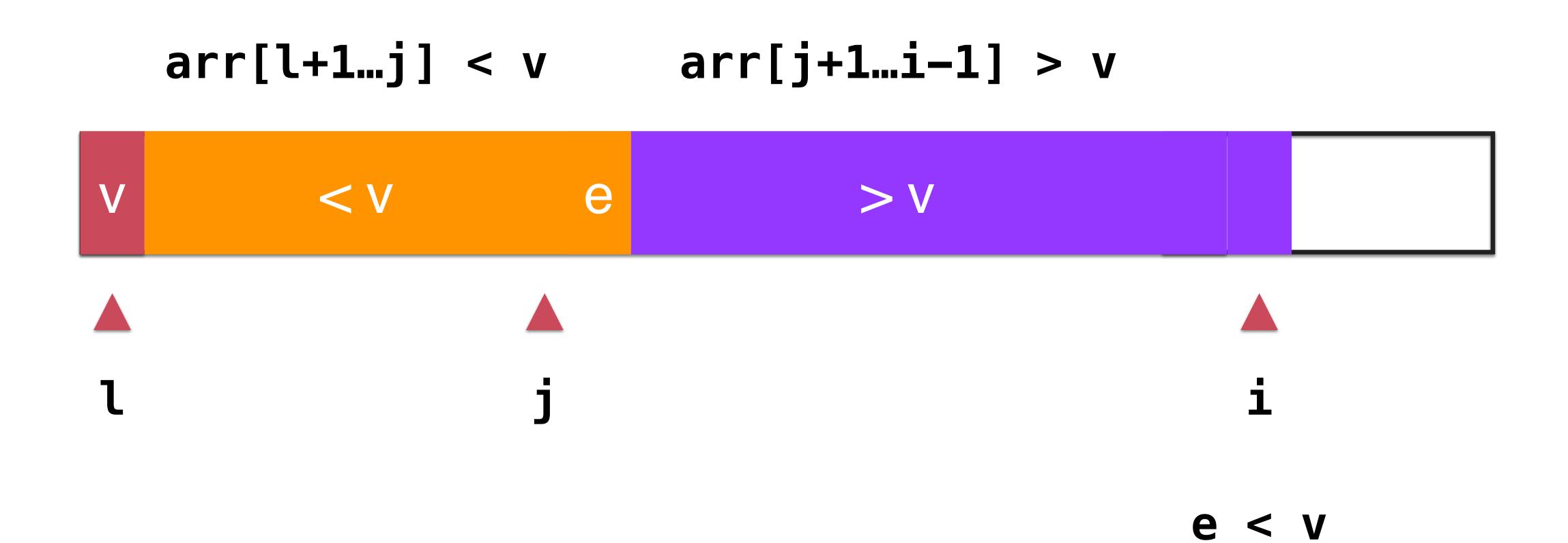
v < v > v e

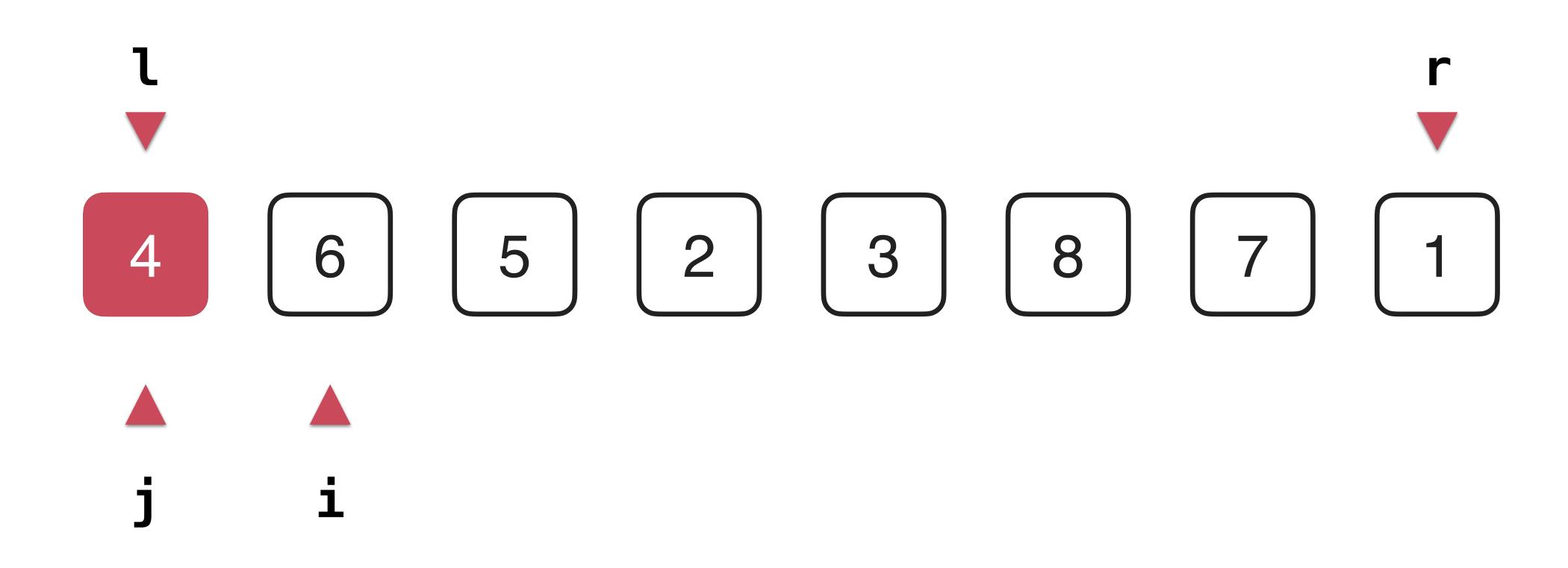
l j i

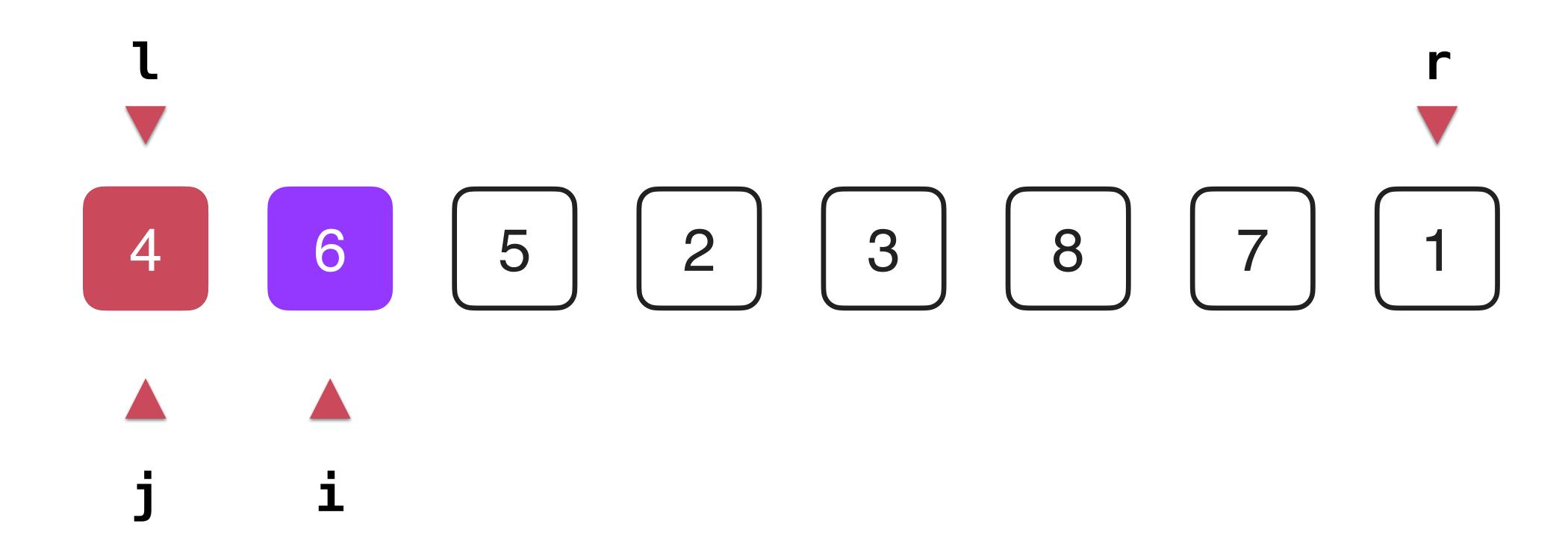
e > v

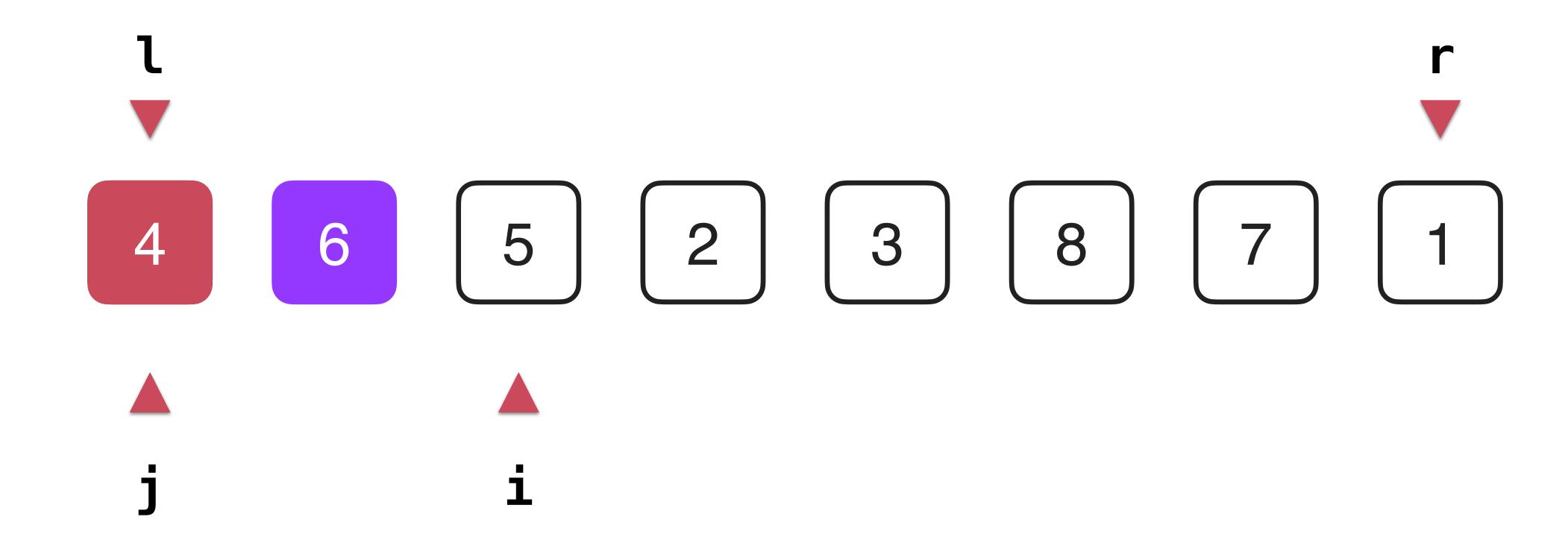


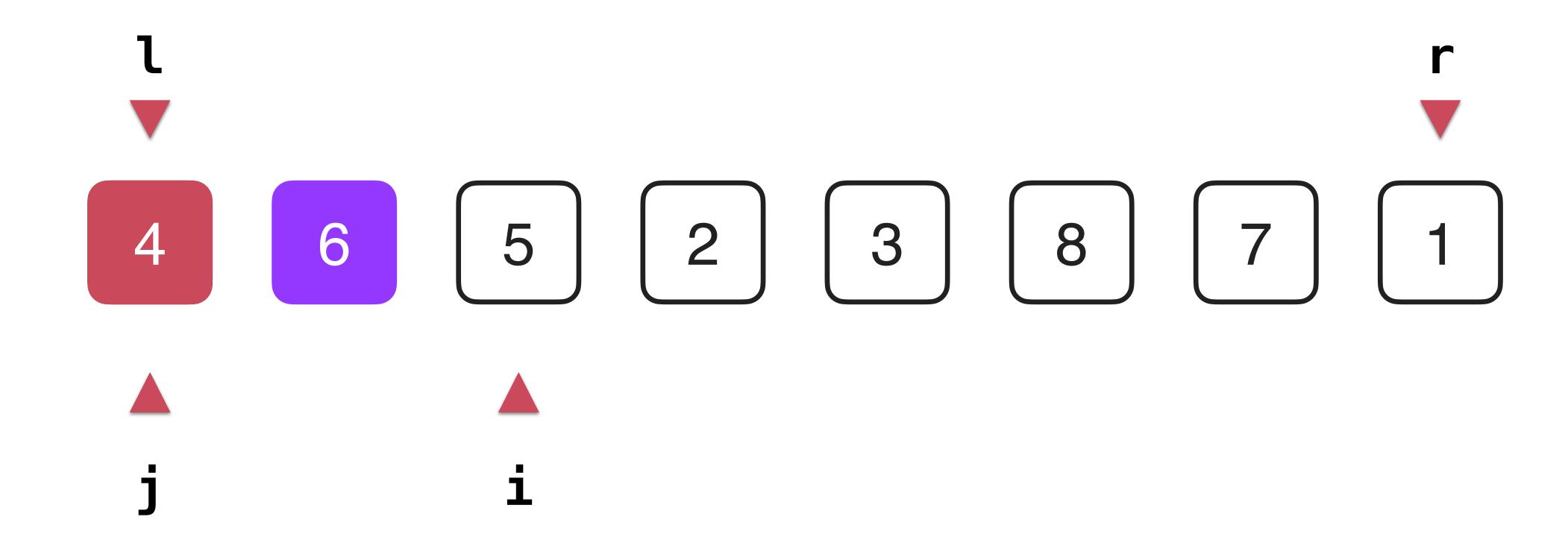
e < v

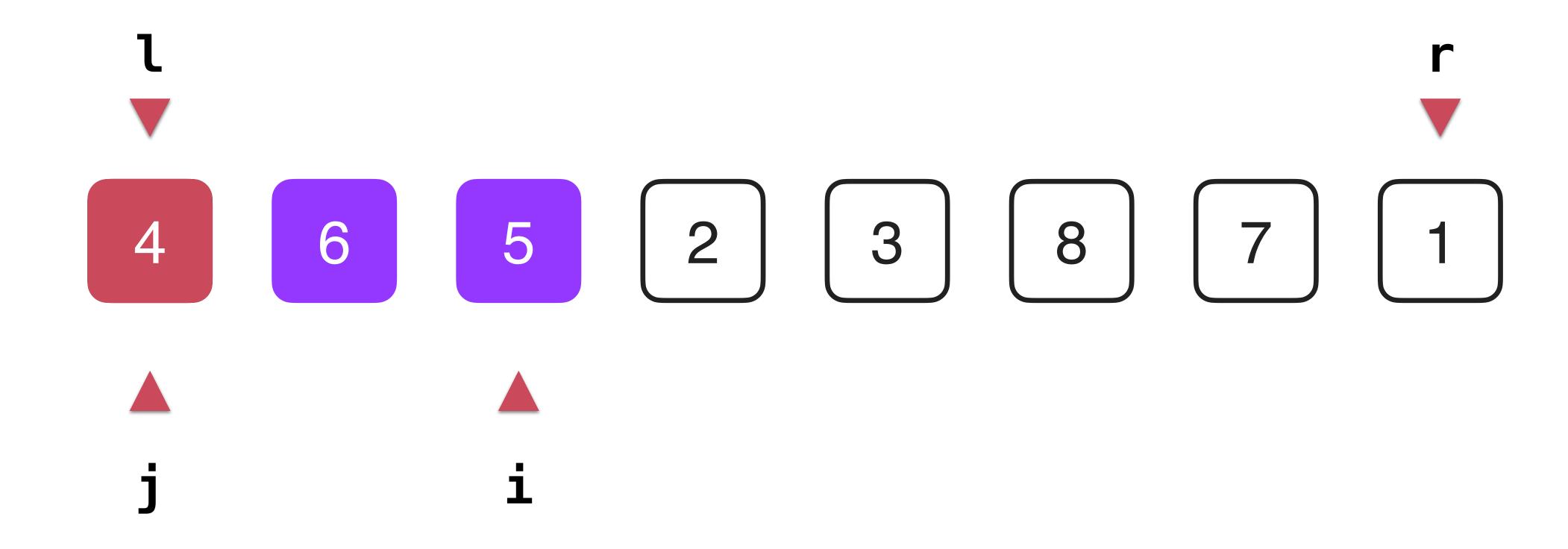


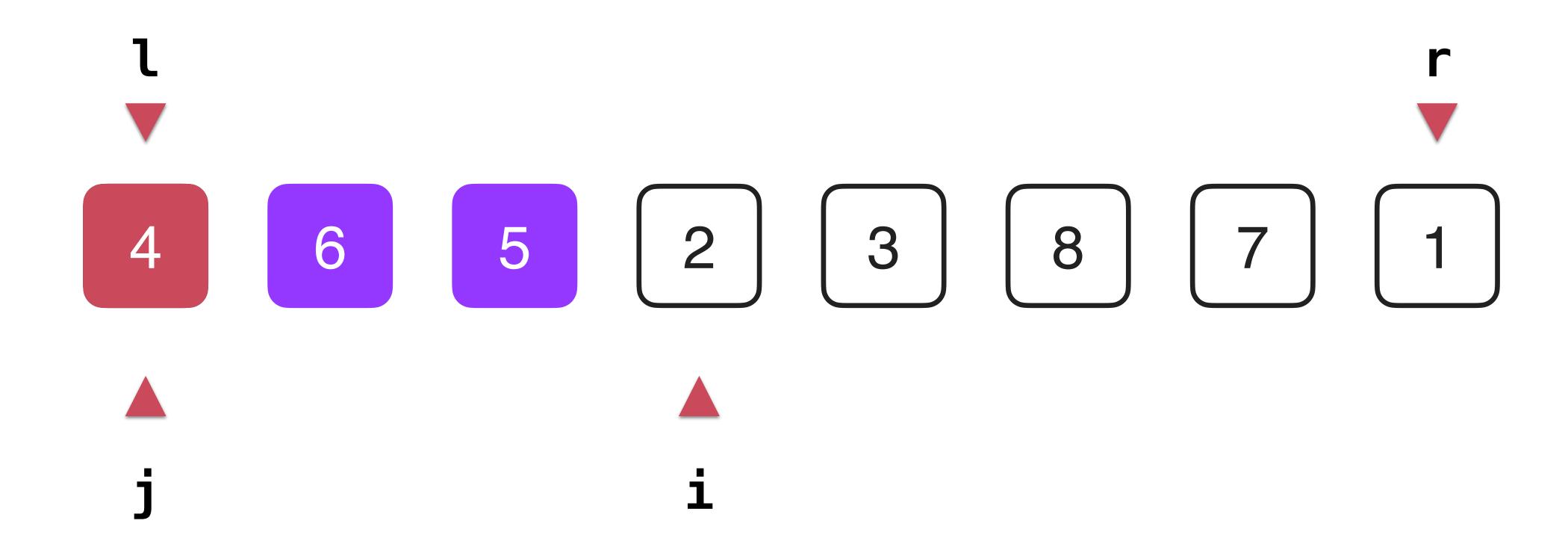


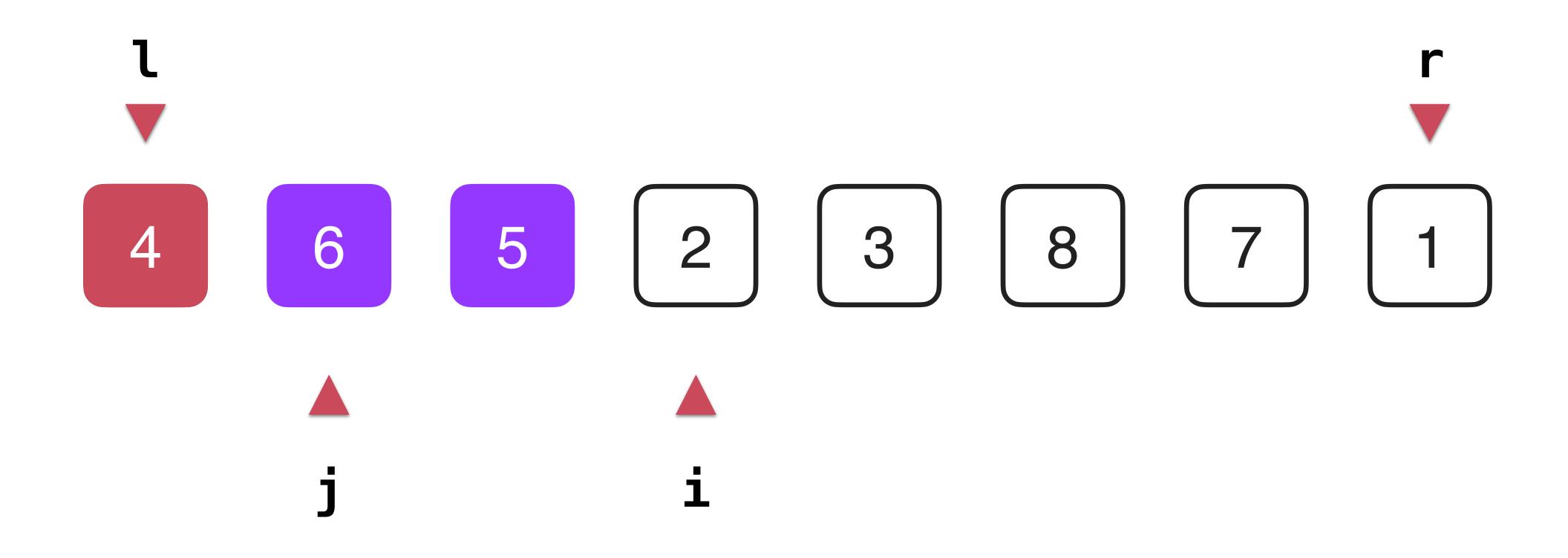


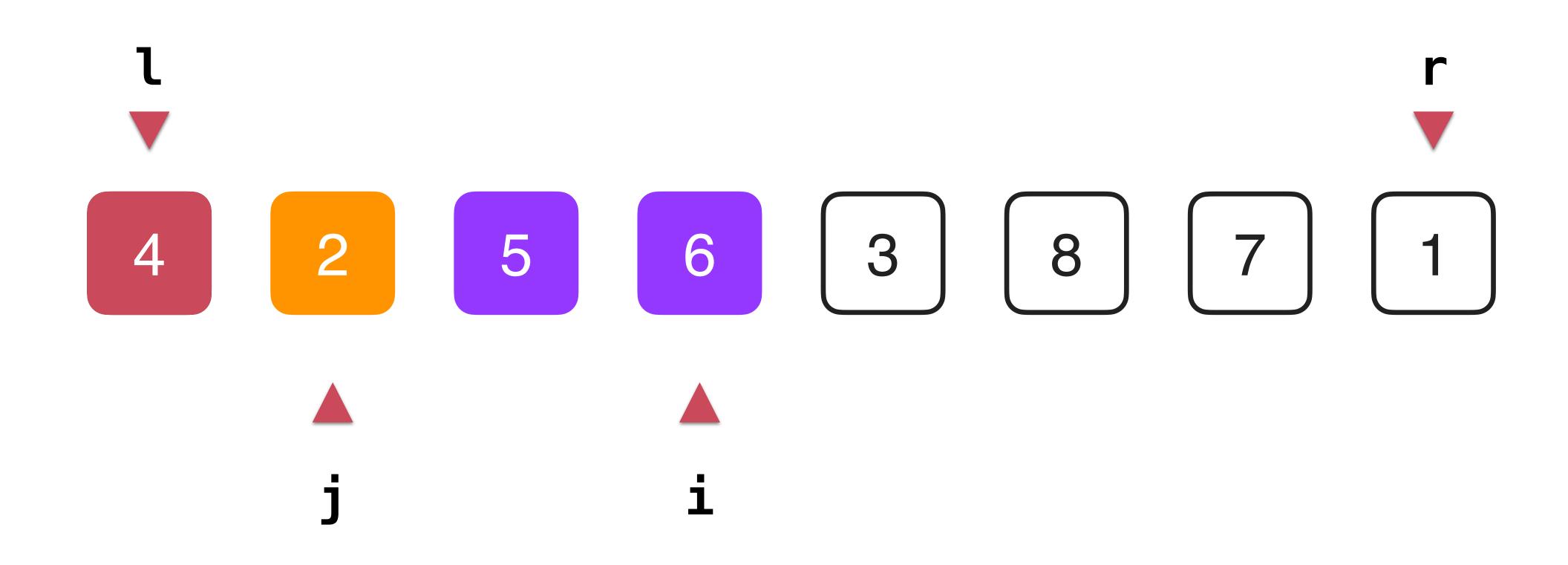


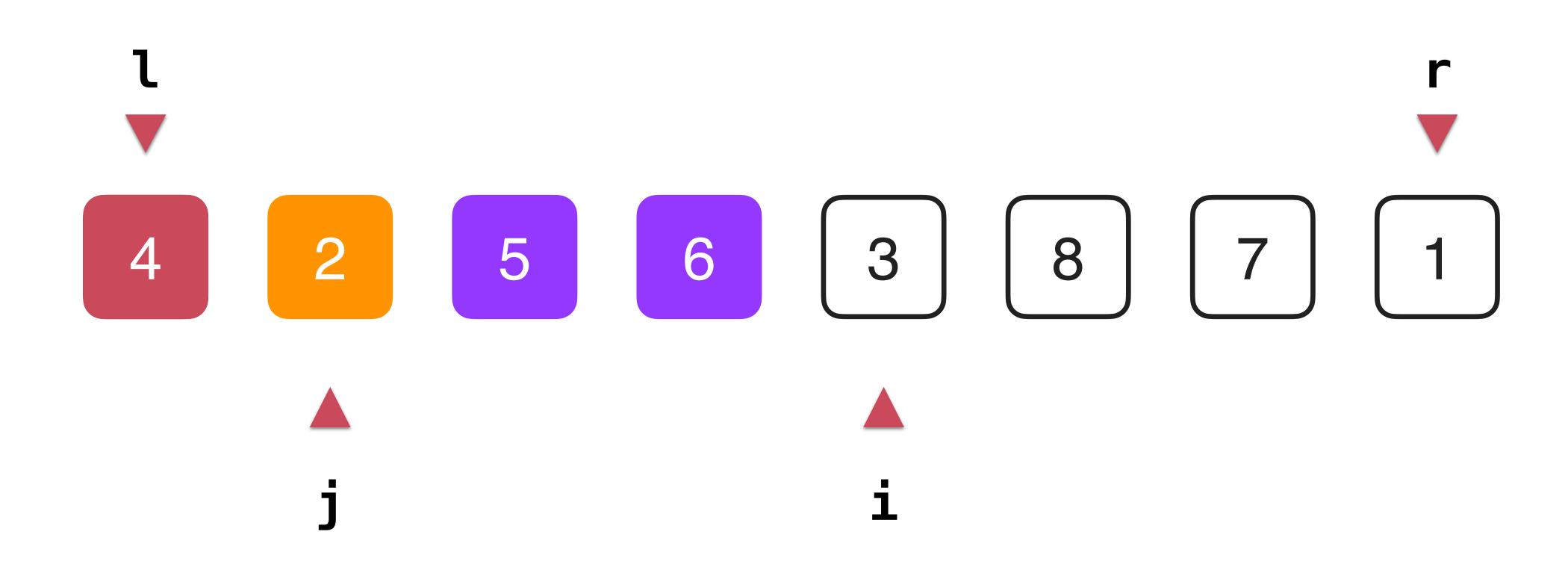


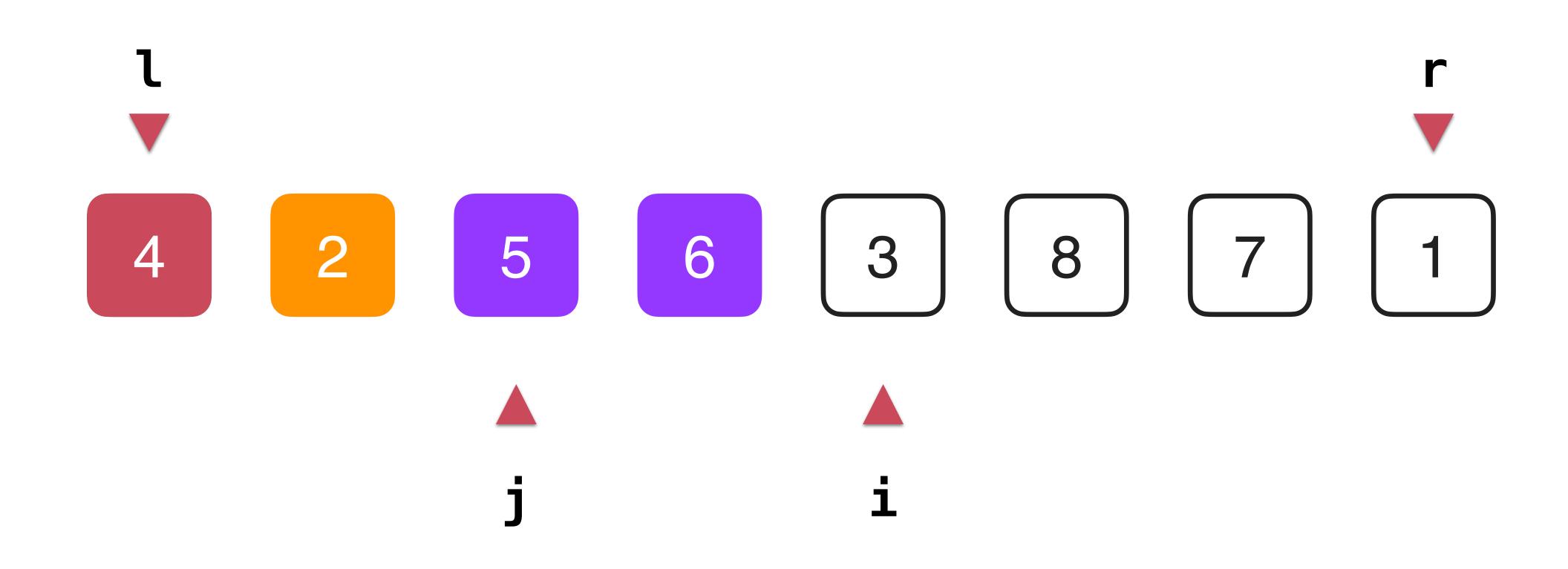


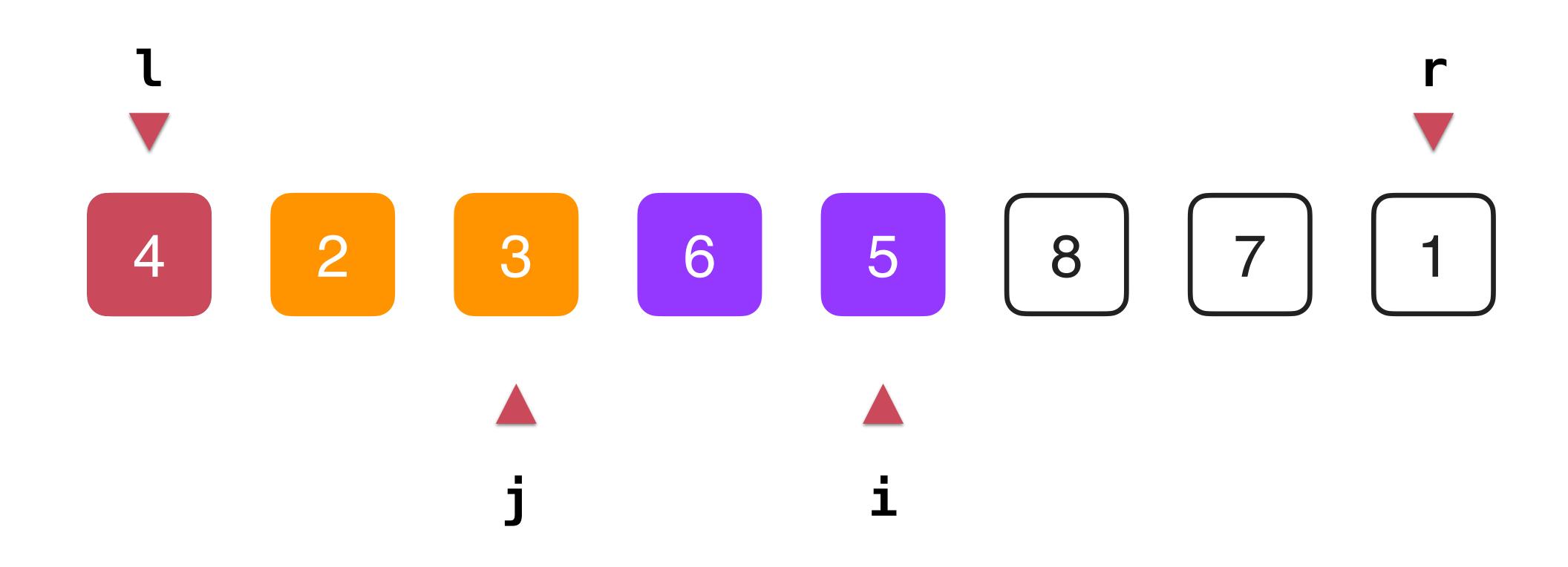


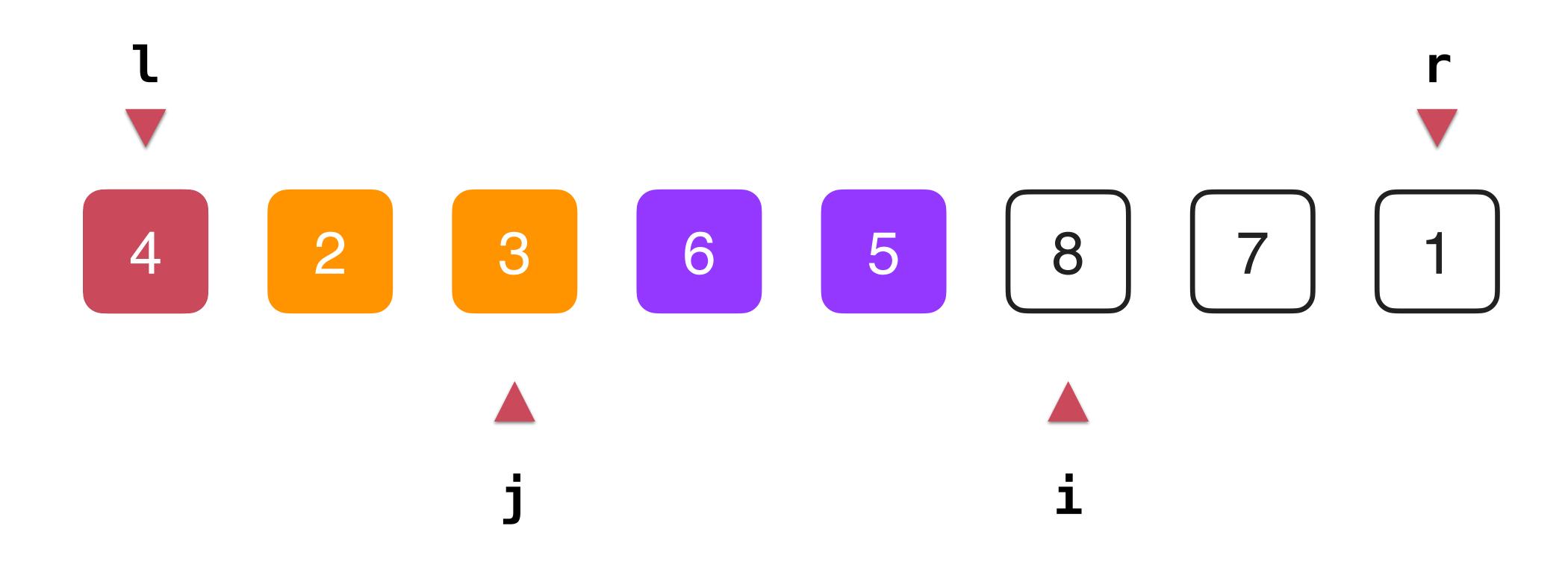


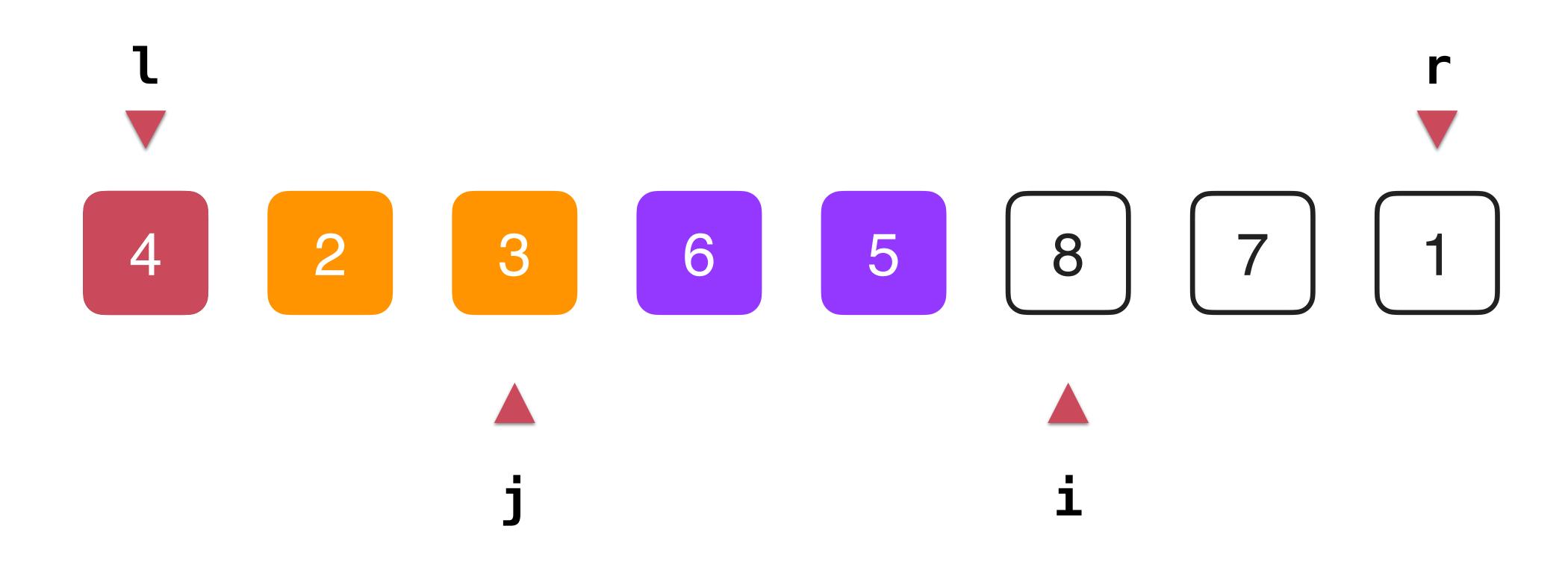


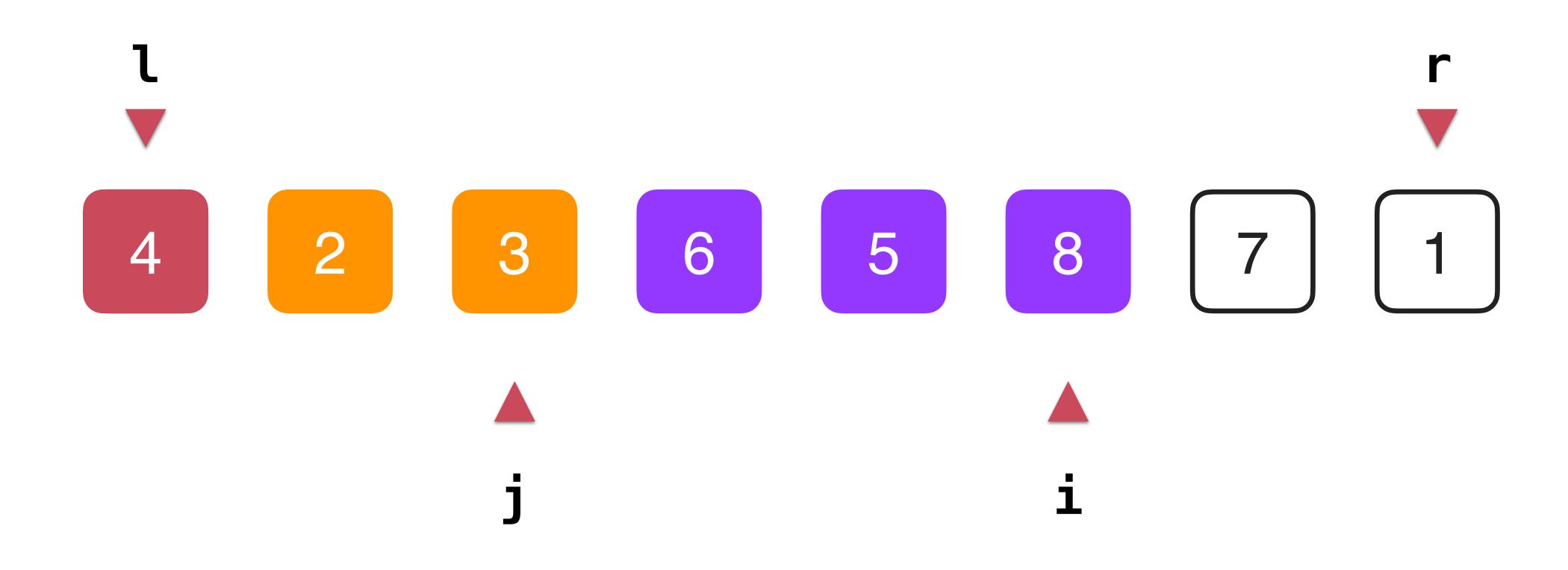








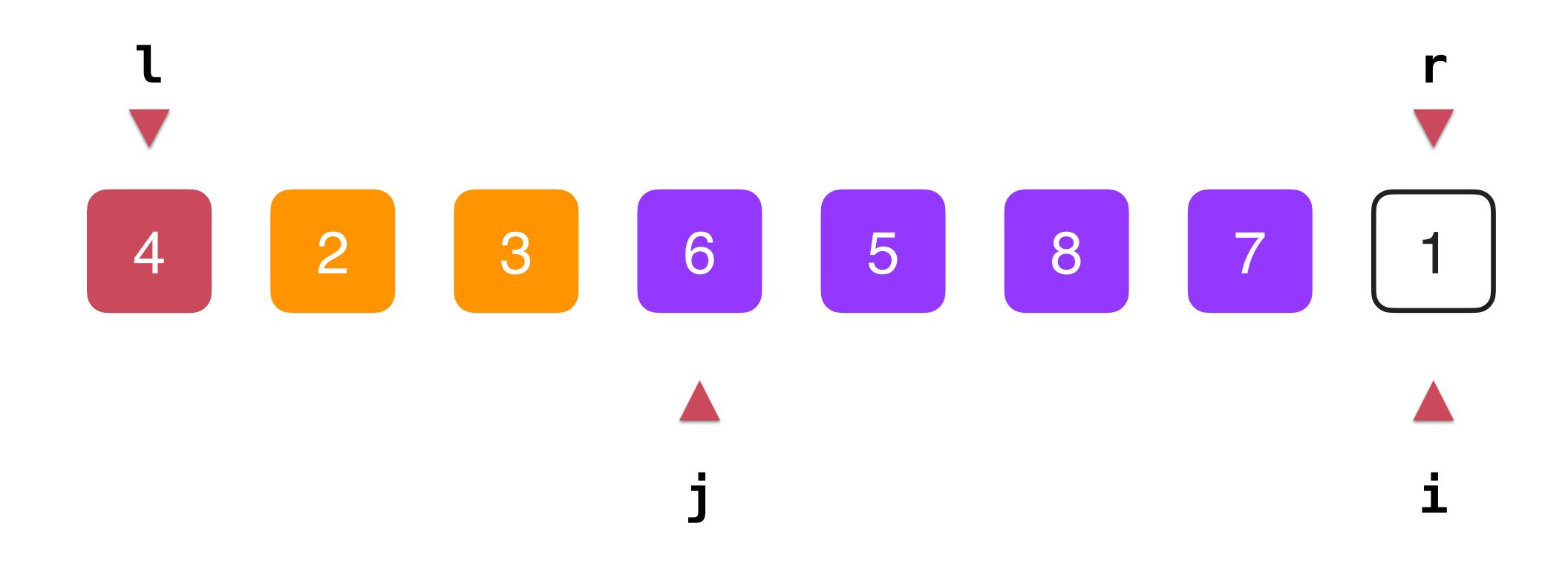


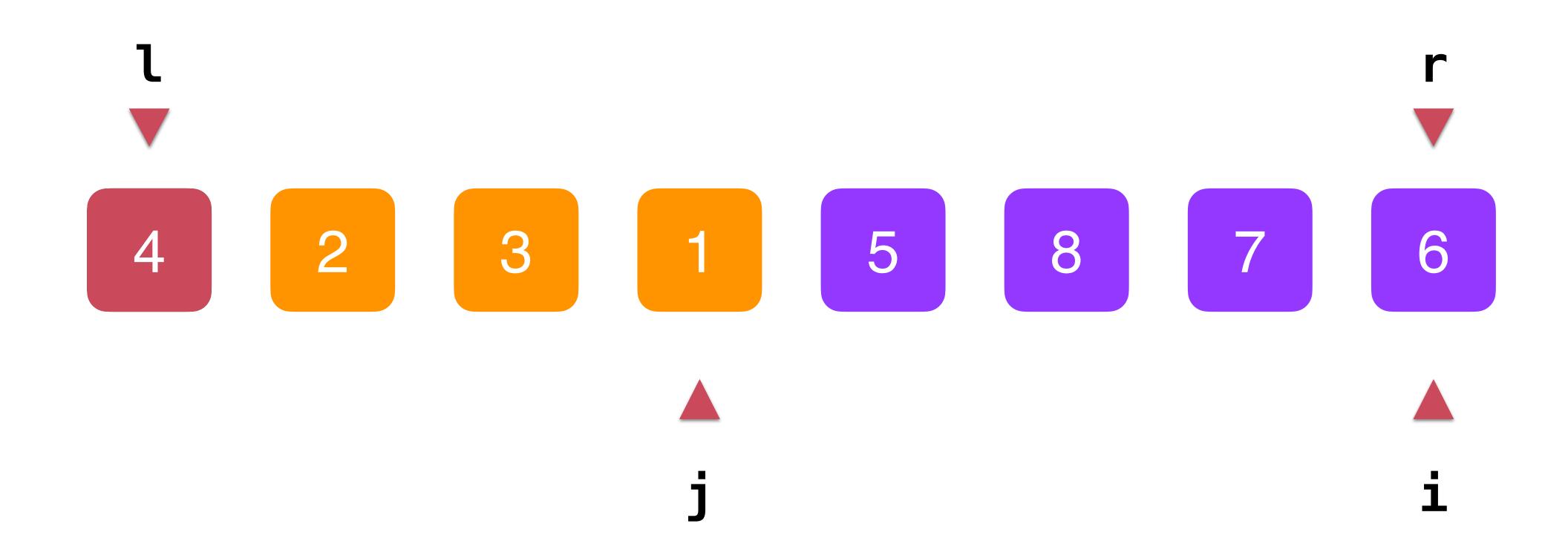


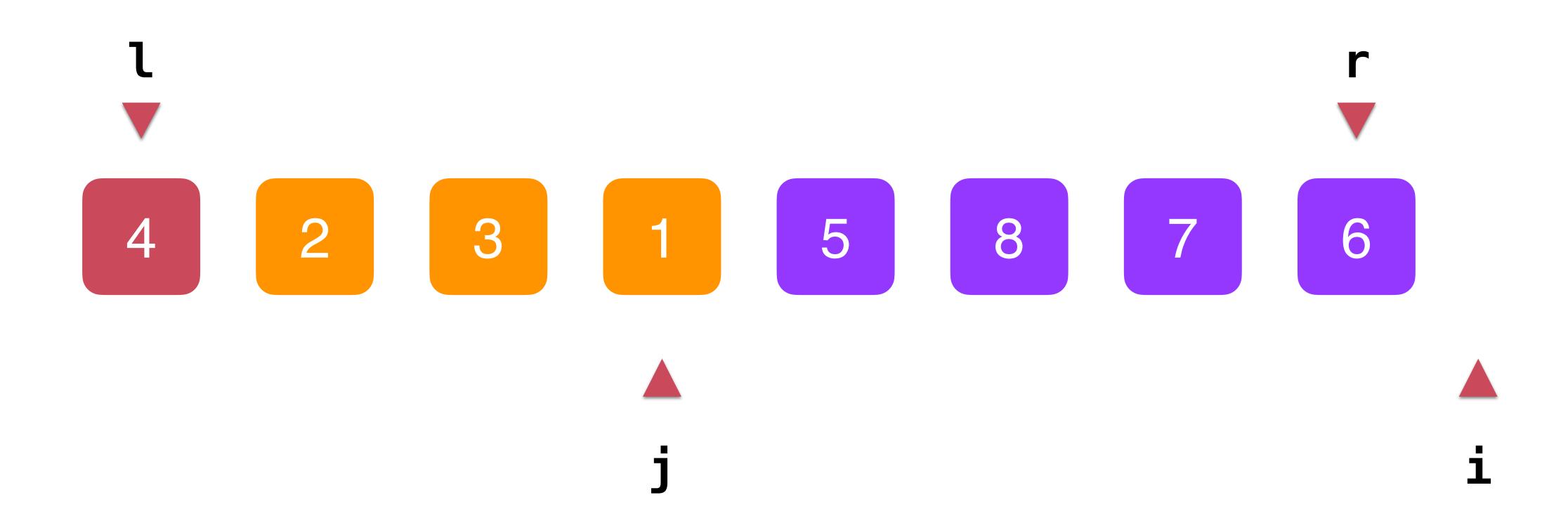


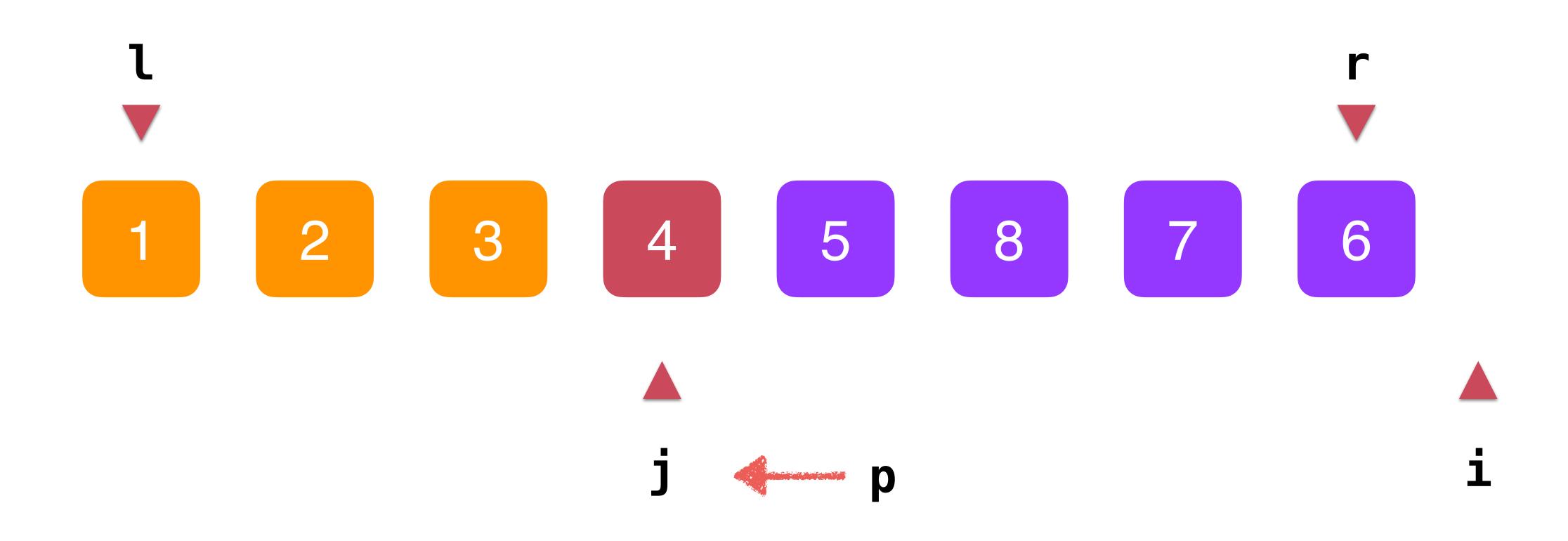












实现 partition

实践: 实现 partition

QuickSort 相关的作业

QuickSort 相关的作业

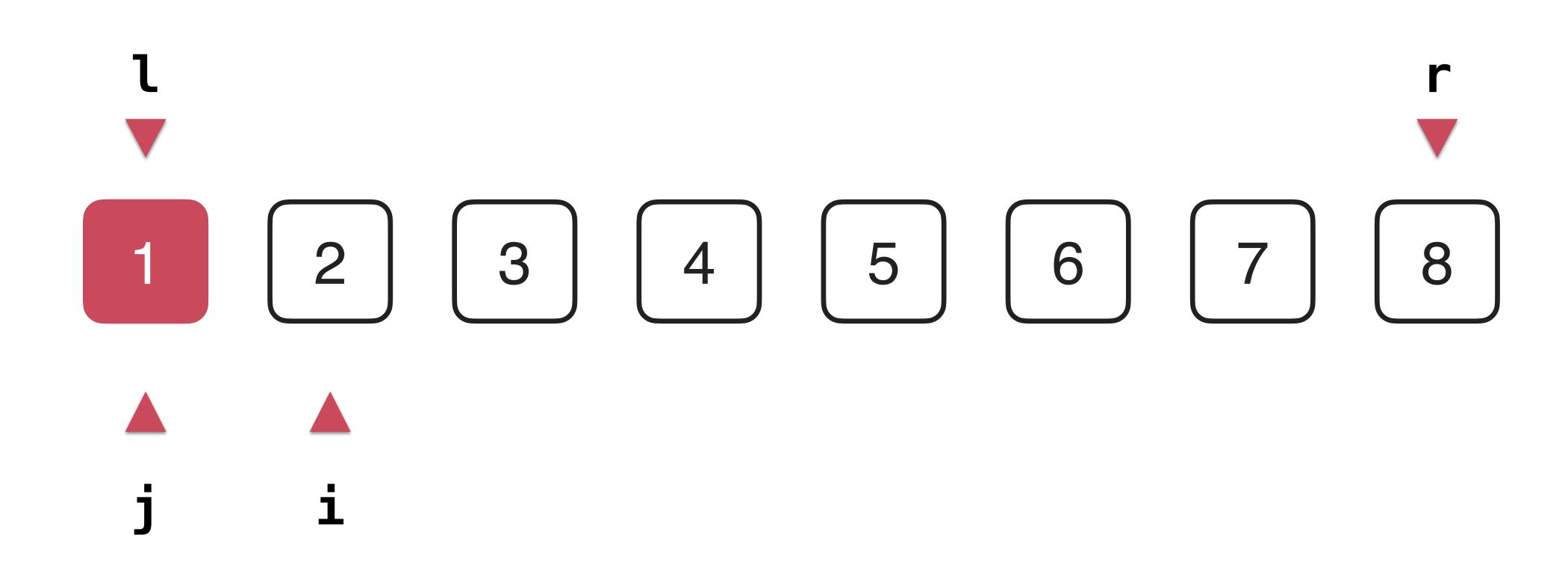
QuickSort 递归的微观解读

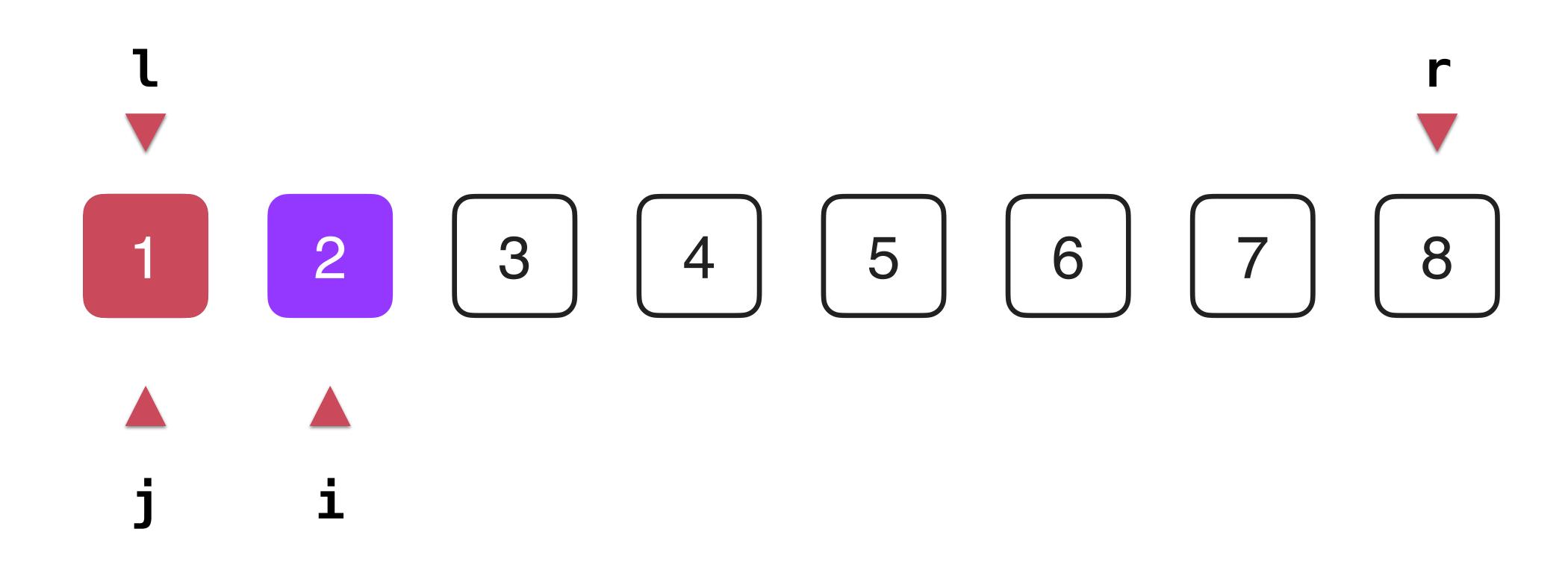
在 QuickSort 的递归函数中添加打印输出,

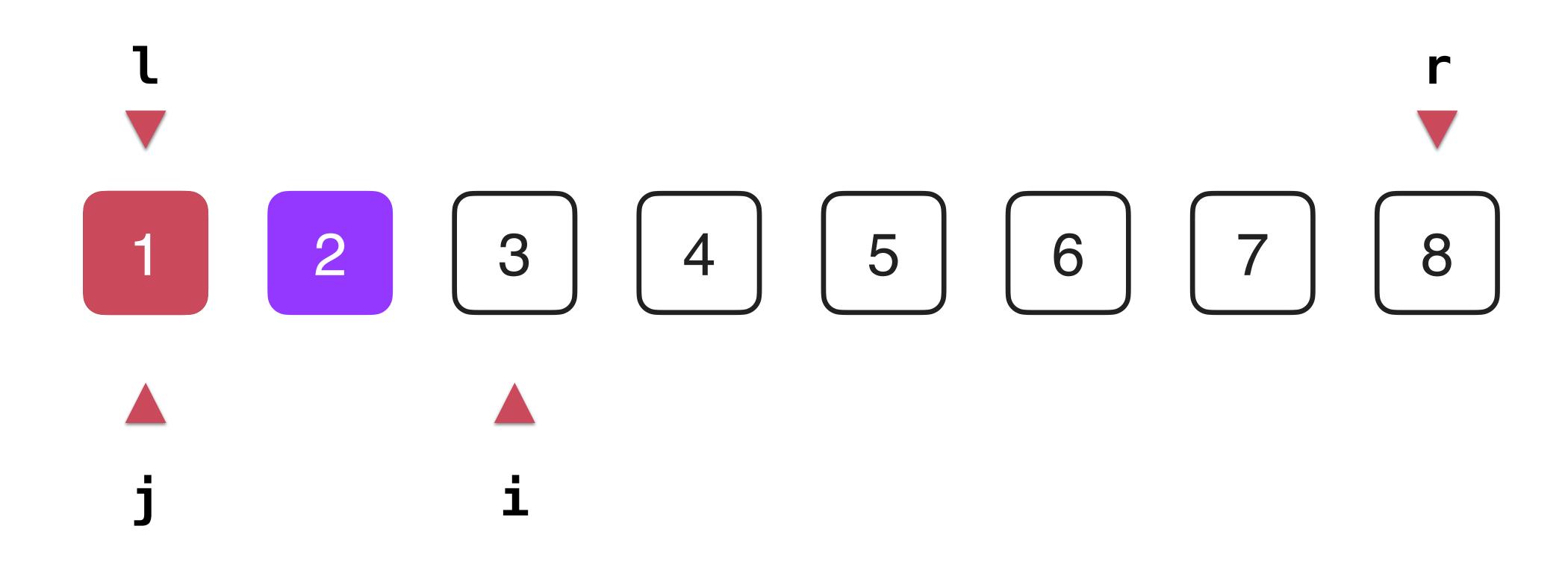
仔细理解 QuickSort 的递归执行过程

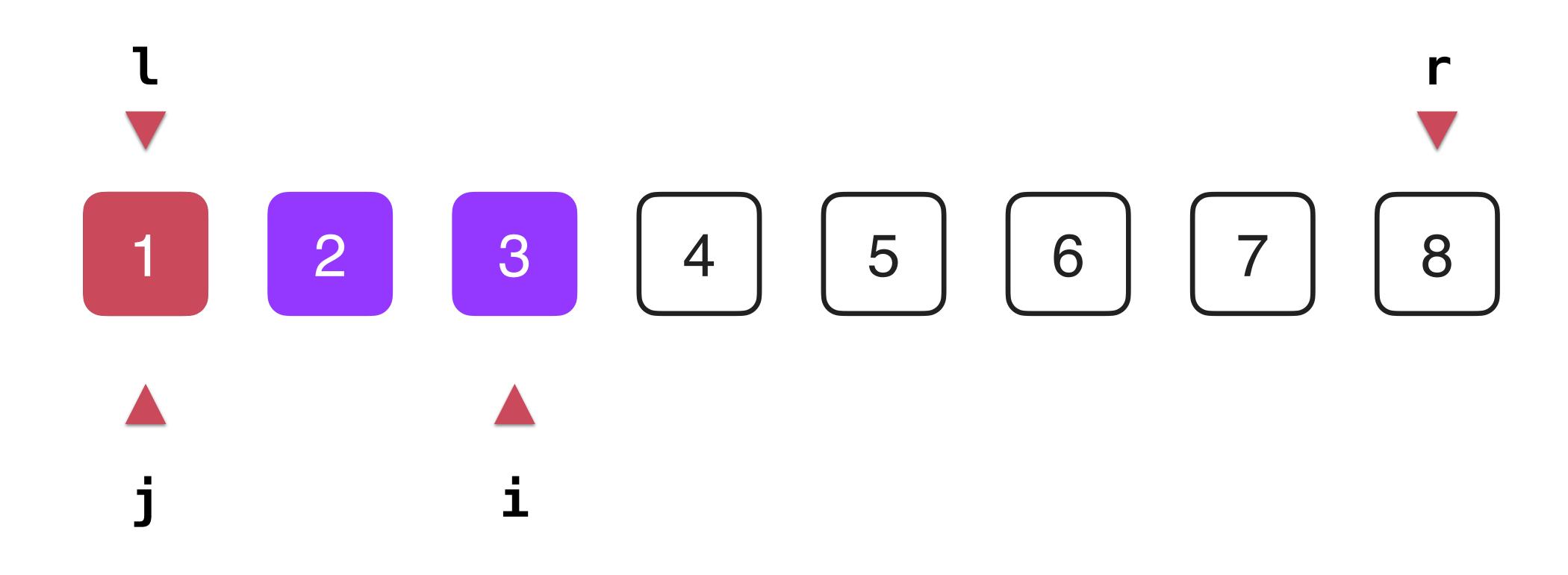
使用 InsertionSort 来优化 QuickSort

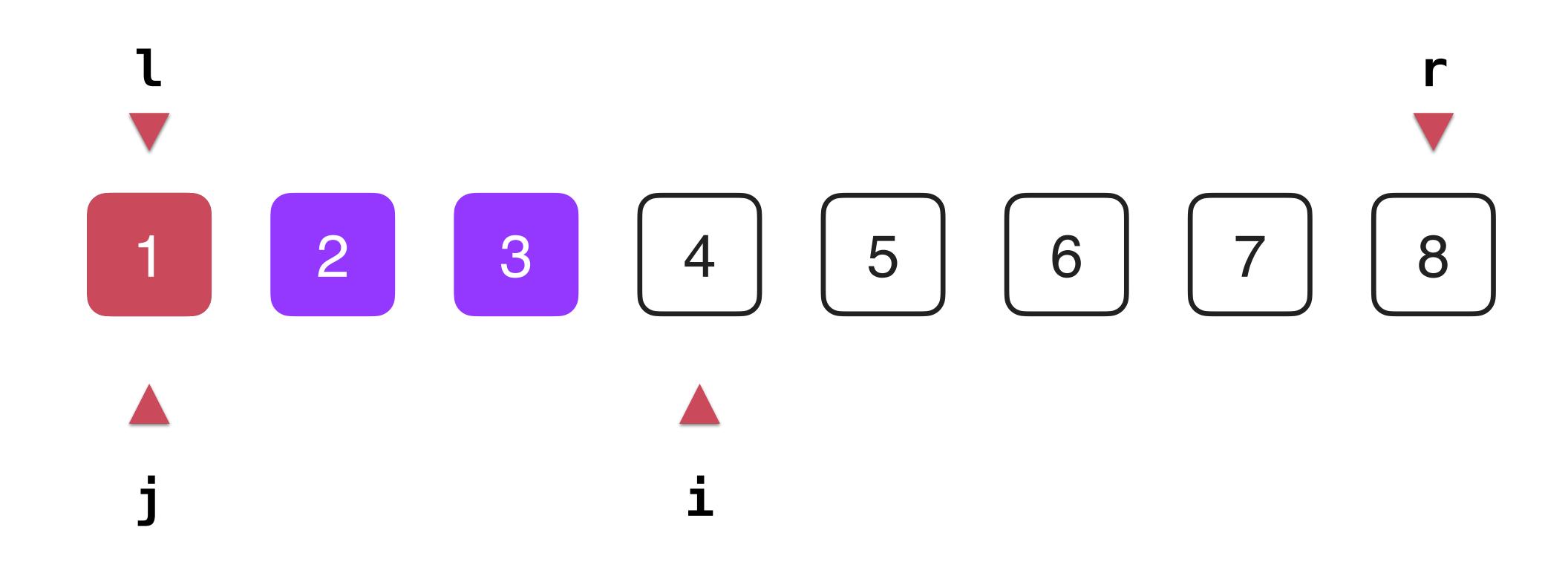
实践:有序数组的问题



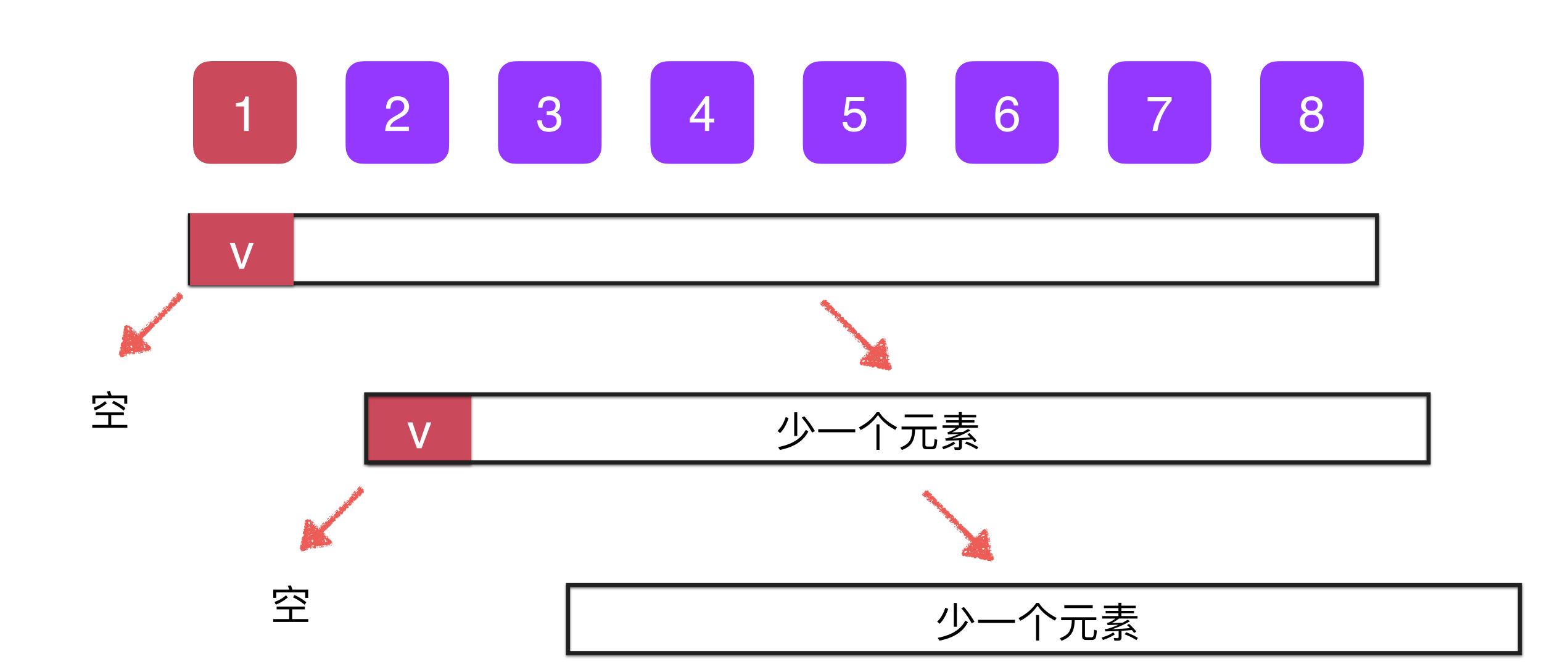


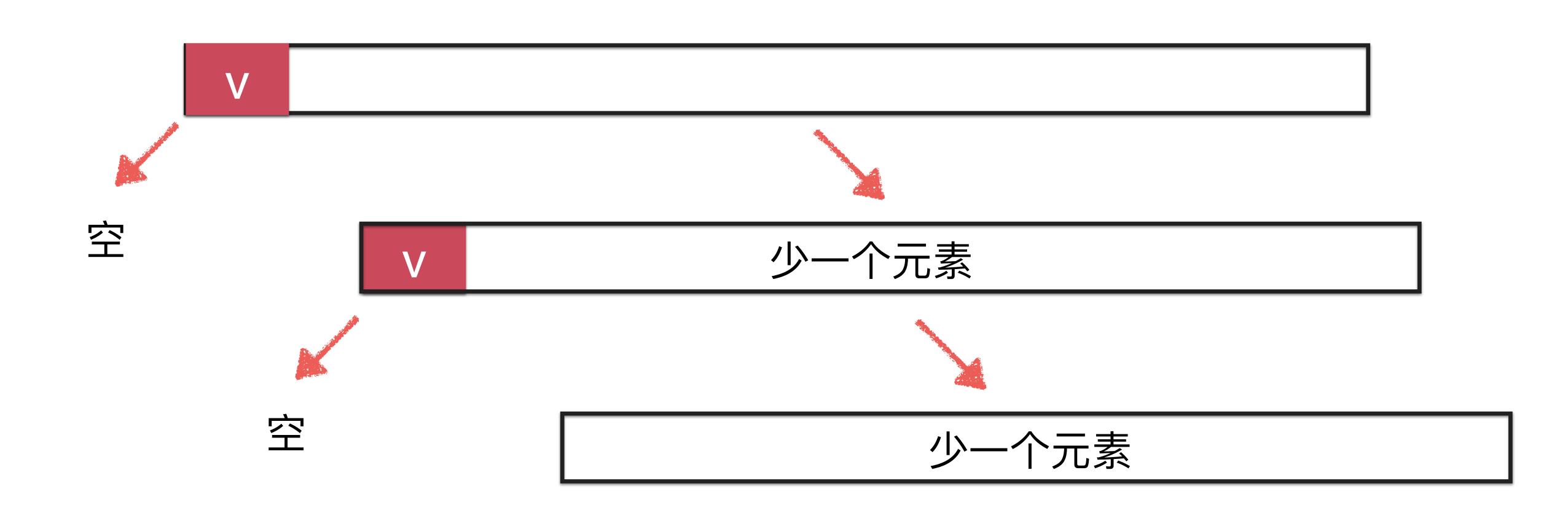












时间复杂度: O(n^2)

递归深度:O(n)

如何解决? 时间复杂度: O(n^2) 递归深度: O(n)

1 2 3 4 5 6 7 8

如何解决? 时间复杂度: O(n^2) 递归深度: O(n)

1 2 3 4 5 6 7 8

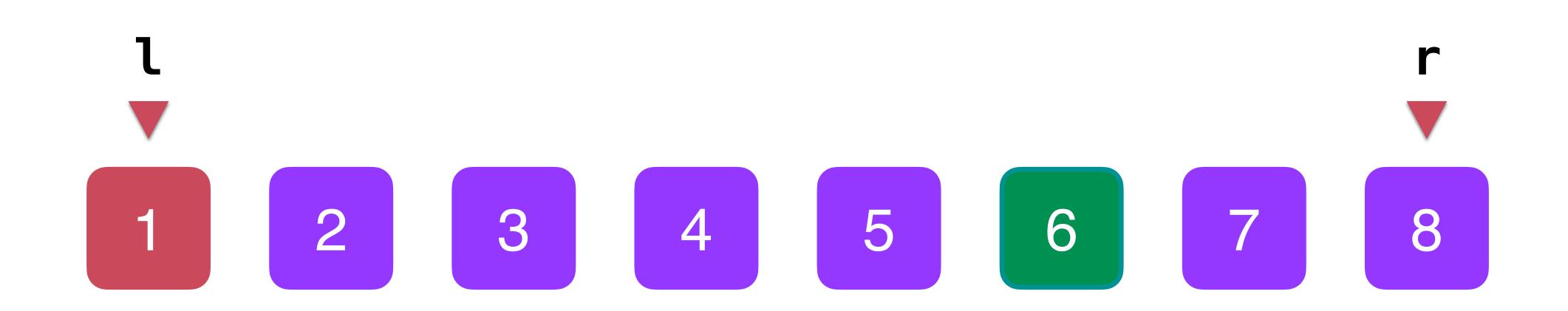
如何解决? 时间复杂度: O(n^2) 递归深度: O(n)

6 2 3 4 5 1 7 8

更具体的理论分析,后续揭晓

为快速排序添加随机化

为快速排序添加随机化



目标: 生成一个 [I, r] 区间的随机值

生成一个 [0, r - I] 区间的随机值

I + [0, r - I] 区间的随机值 -> [I, r] 区间的随机值

实践:为快速排序添加随机化

为快速排序添加随机化

1 2 3 4 5 6 7 8

每次取中间的值作为标定点?

万一随机每次都取最小值?

$$\frac{1}{n} \times \frac{1}{n-1} \times \frac{1}{n-2} \times \dots \times \frac{1}{1} = \frac{1}{n!}$$

两个作业:深入玩转快速排序

作业一

private static <E extends Comparable<E>> int partition(E[] arr, int I, int r){

```
// 生成 [/, 引之间的随机索引
int p = I + (new Random()).nextInt(r - I + 1);
swap(arr, t, p),
// arr[l+1...i] < v ; arr[j+1...i] >= v
int j = l;
for(int i = I + 1; i \le r; i + +)
   if(arr[i].compareTo(arr[l]) < 0){</pre>
      j ++;
      swap(arr, i, j);
swap(arr, I, j);
return j;
```

作业二

private static <E extends Comparable<E>> int partition(E[] arr, int I, int r){

```
swap(arr, I, (I + r) / 2);
// arr[l+1...i] < v ; arr[j+1...i] >= v
int j = l;
for(int i = I + 1; i \le r; i + +)
   if(arr[i].compareTo(arr[l]) < 0){</pre>
      j ++;
      swap(arr, i, j);
swap(arr, I, j);
return j;
```

ArrayGenerator.generateSpecialArray(n)

其他

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