# 算法与数据结构体系课程

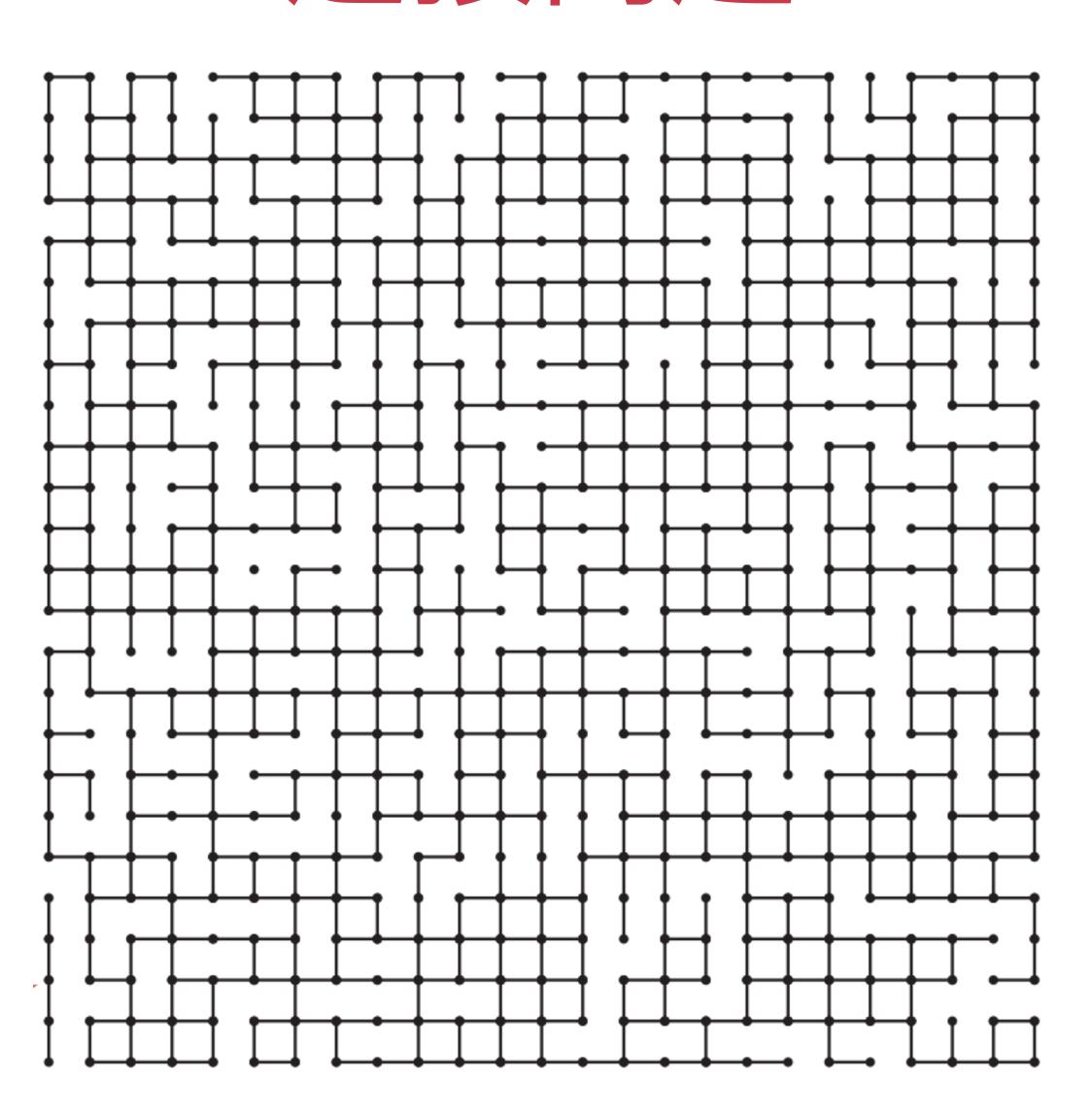
liuyubobobo

# 并查集 Union Find

# 一种很不一样的树形结构

# 连接问题 Connectivity Problem

# 连接问题



# 连接问题

网络中节点间的连接状态

• 网络是个抽象的概念: 用户之间形成的网络

数学中的集合类实现

# 连接问题和路径问题

比路径问题要回答的问题少

• 和堆作比较

## 并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)
- isConnected(p,q)

# 实践:并查集的接口设计

# Quick Find

## 并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)
- isConnected(p,q)

## 并查集的基本数据表示

0 1 2 3 4 5 6 7 8 9

000011111

## 并查集的基本数据表示

id 0 1 2 3 4 5 6 7 8 9 0 1 0 1 0 1 0 1

## 并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)
- isConnected(p,q)

# 并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)
- isConnected(p,q) == find(q)

#### Quick Find

id 0 1 2 3 4 5 6 7 8 9 0 1 0 1 0 1 0 1

Quick Find 时间复杂度 O(1)

# Quick Find 下的 Union

union(1, 4)

id 0 1 2 3 4 5 6 7 8 9

0 1 0 1 0 1 0 1

# Quick Find 下的 Union

union(1, 4)

id 0 1 2 3 4 5 6 7 8 9

1 1 1 1 1 1 1 1 1

# Quick Find 下的 Union

id 0 1 2 3 4 5 6 7 8 9 1 1 1 1 1 1 1 1 1

Quick Find 下的 Union 时间复杂度 O(n)

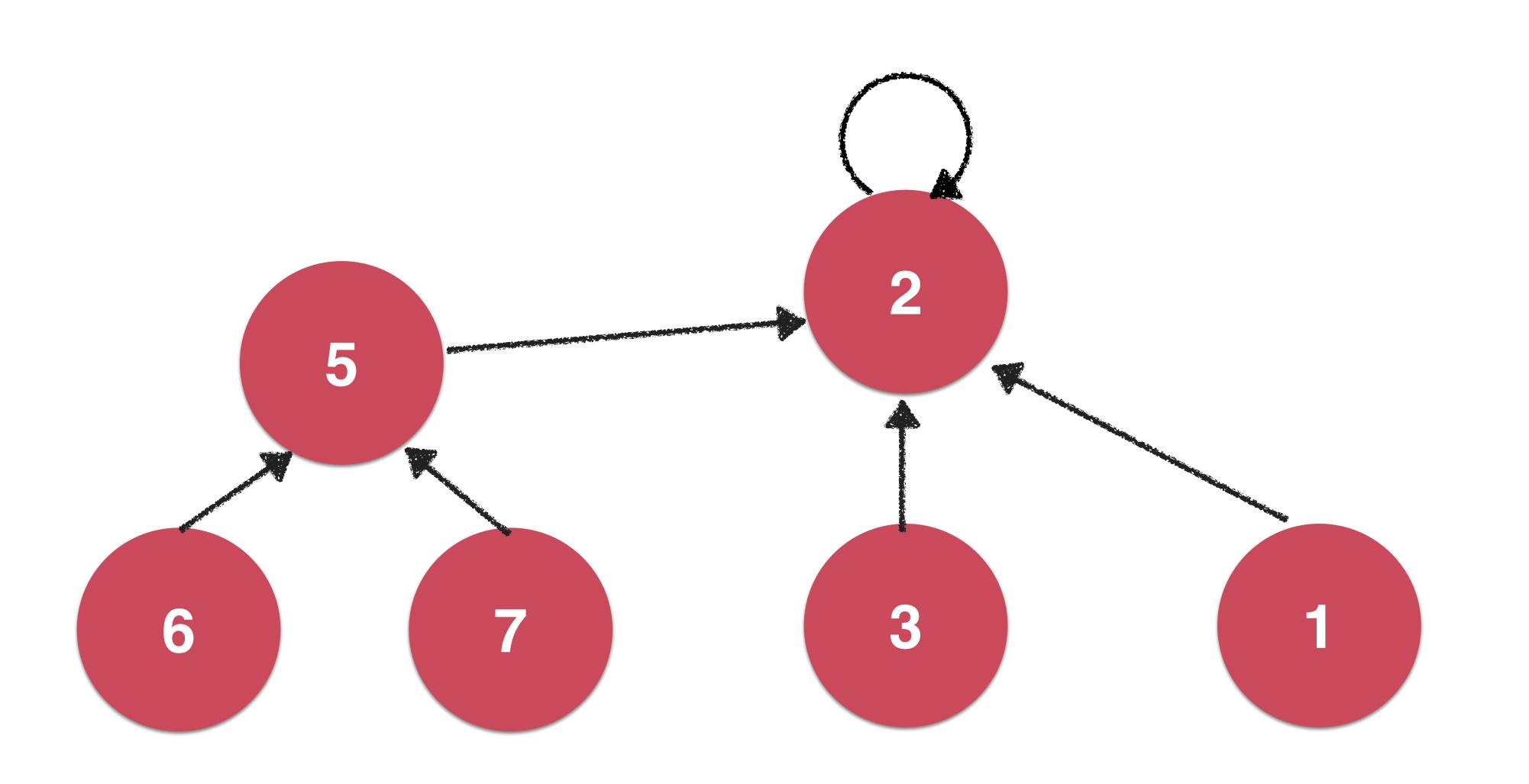
# 实践: Quick Find

#### Quick Find

- unionElements(p,q)O(n)
- isConnected(p,q) O(1

# Quick Union

# 将每一个元素,看做是一个节点



# Quick Union 下的数据表示

parent 0 1 2 3 4 5 6 7 8 9

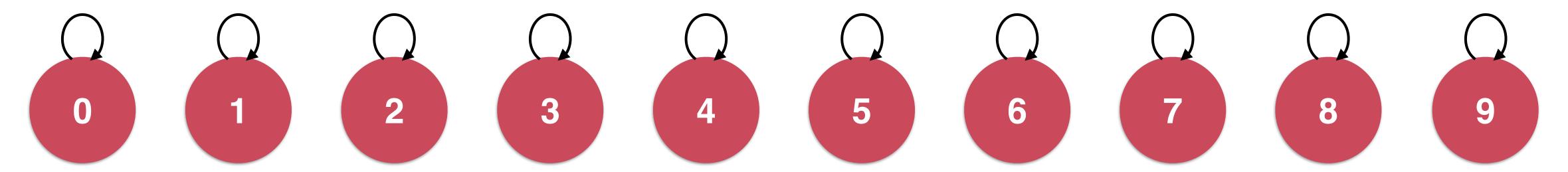
0 1 2 3 4 5 6 7 8 9

#### Quick Union



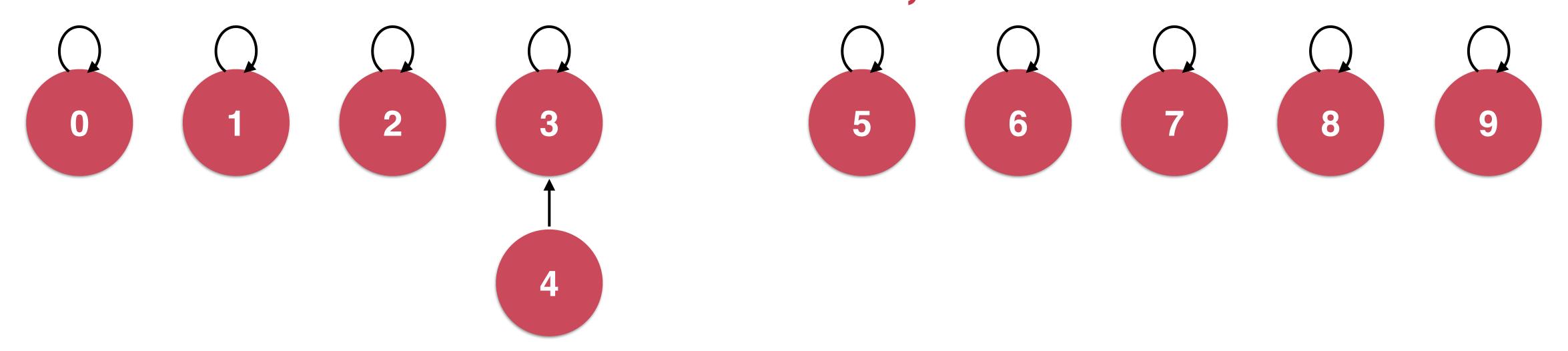
parent 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9

### union 4, 3



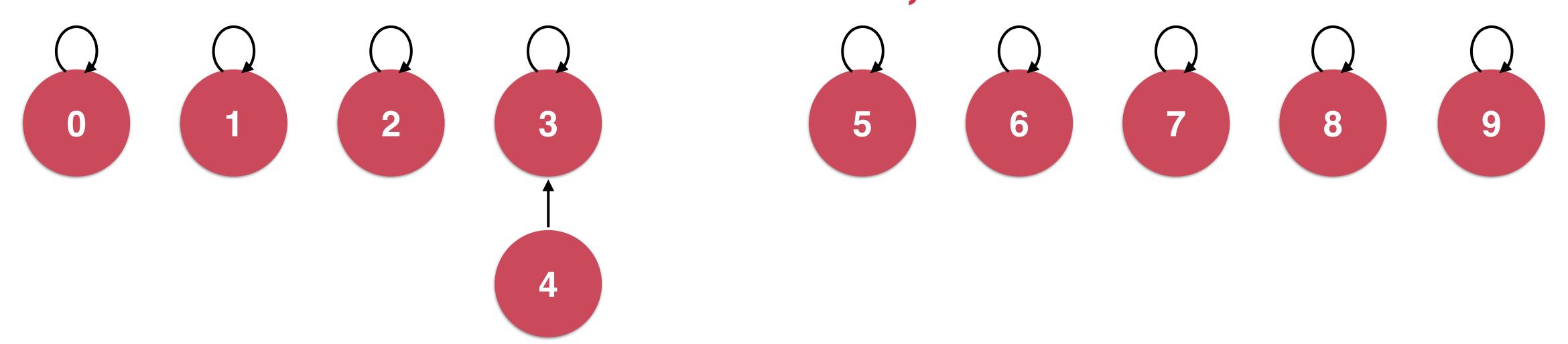
parent 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9

### union 4, 3



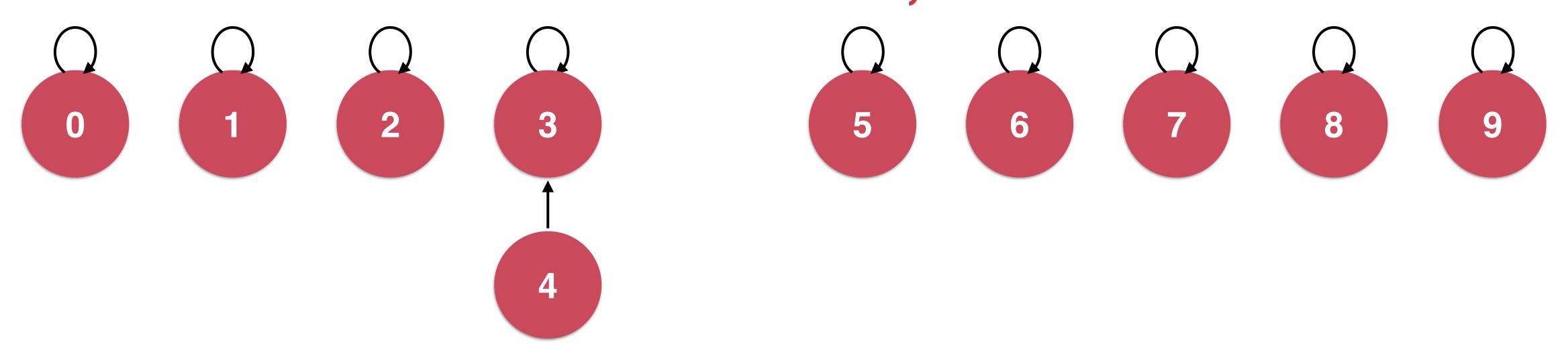


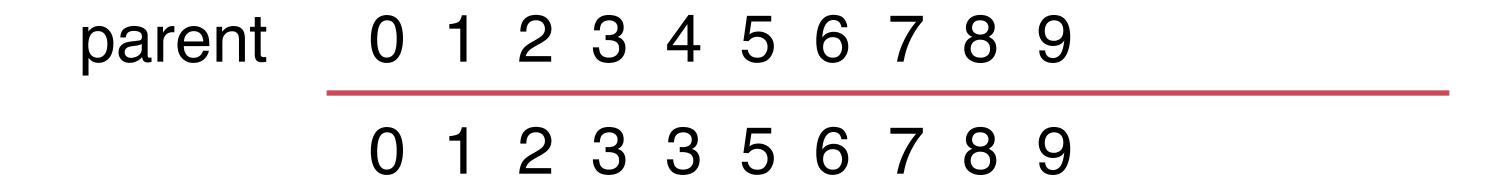
#### union 4,3



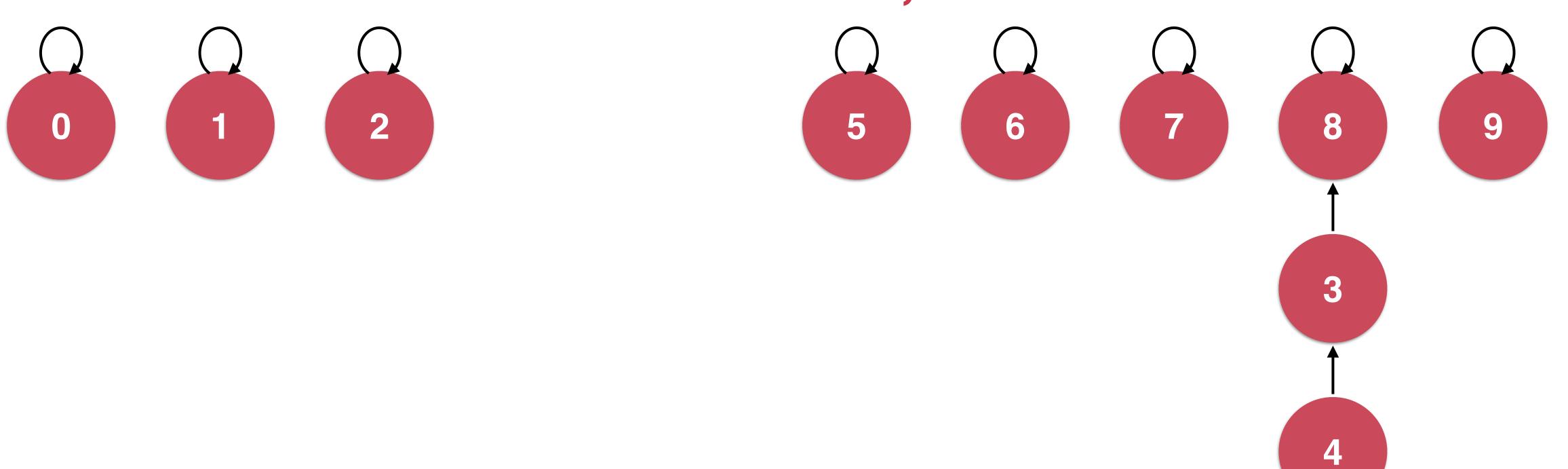
parent 0 1 2 3 4 5 6 7 8 9 0 1 2 3 3 5 6 7 8 9

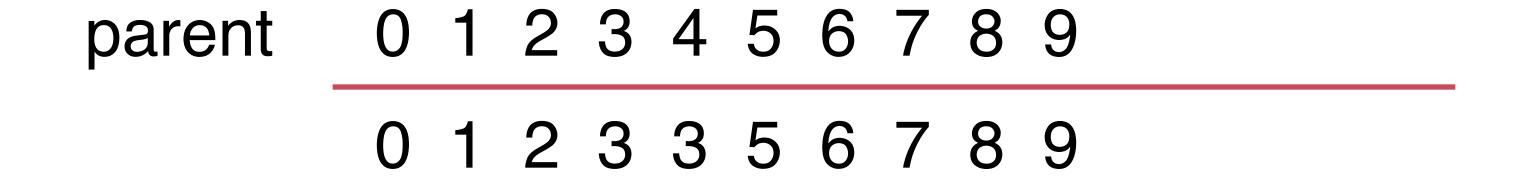
#### union 3,8



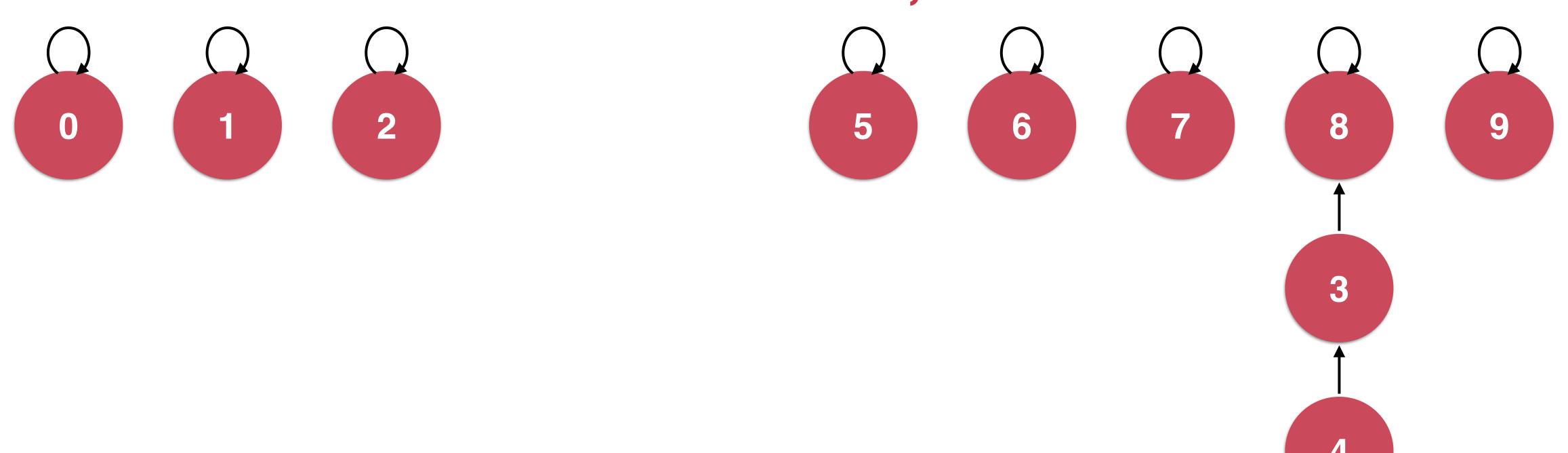


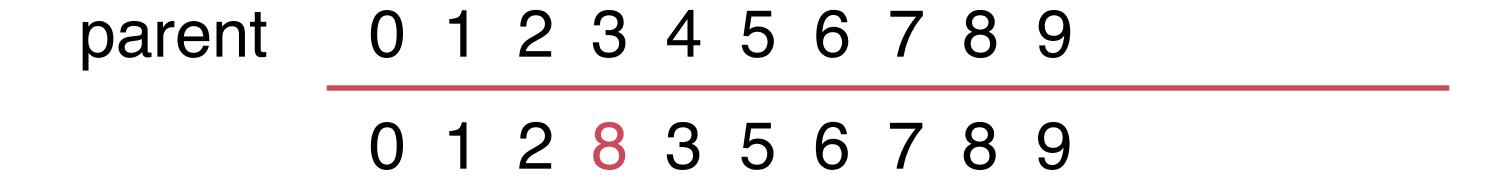
## union 3,8



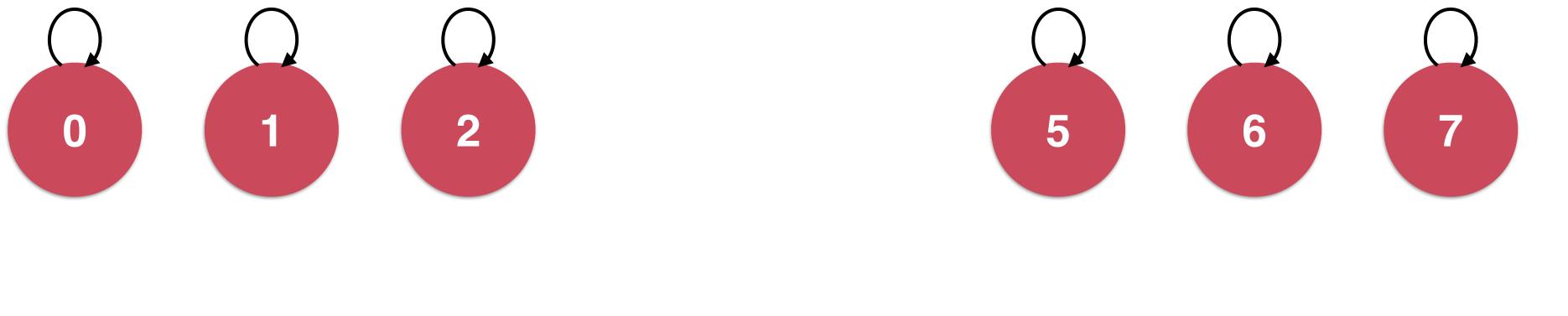


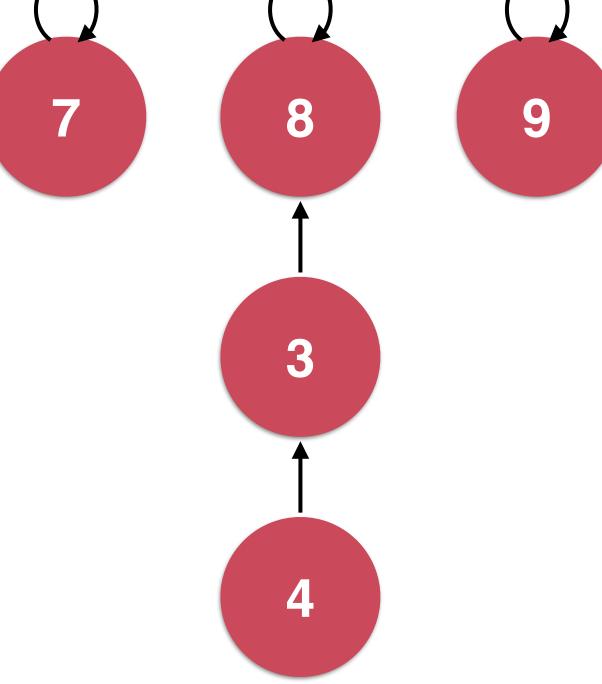
## union 3,8





### union 6,5

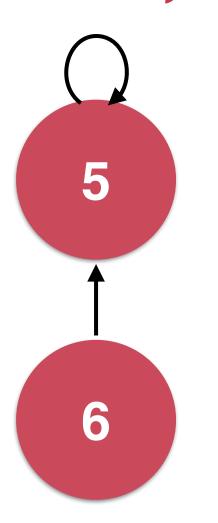


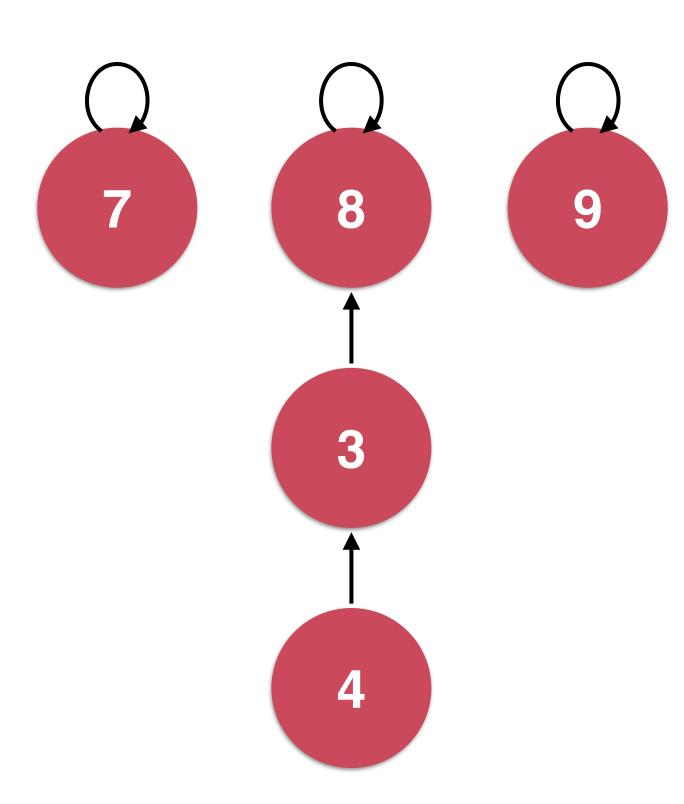


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 6 7 8 9

## union 6,5



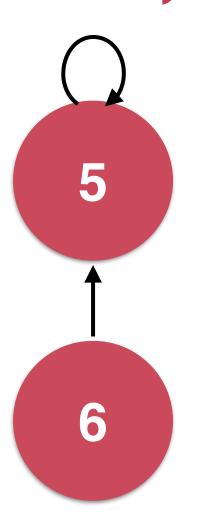


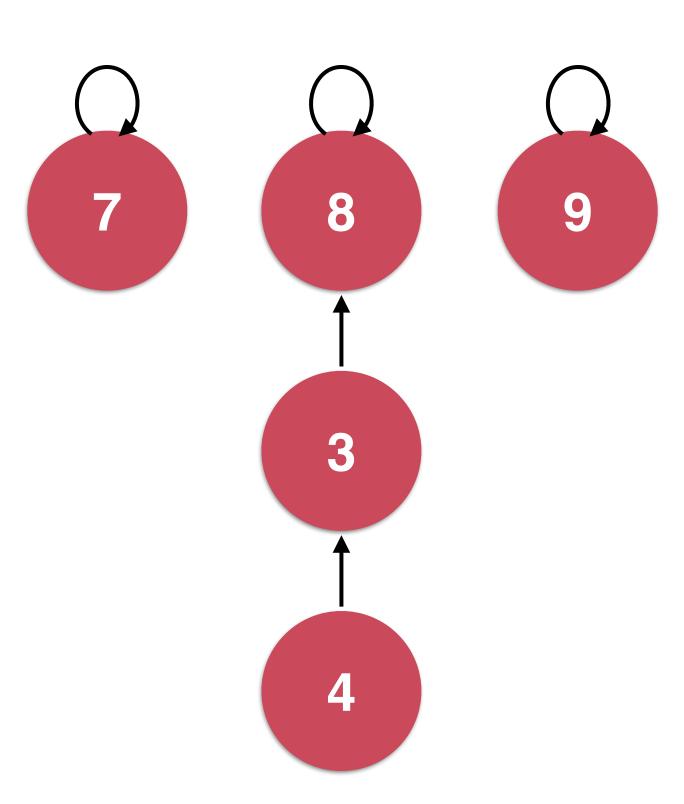


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 6 7 8 9

## union 6,5



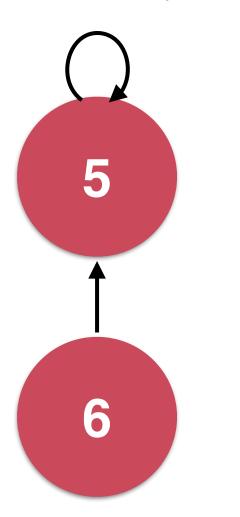


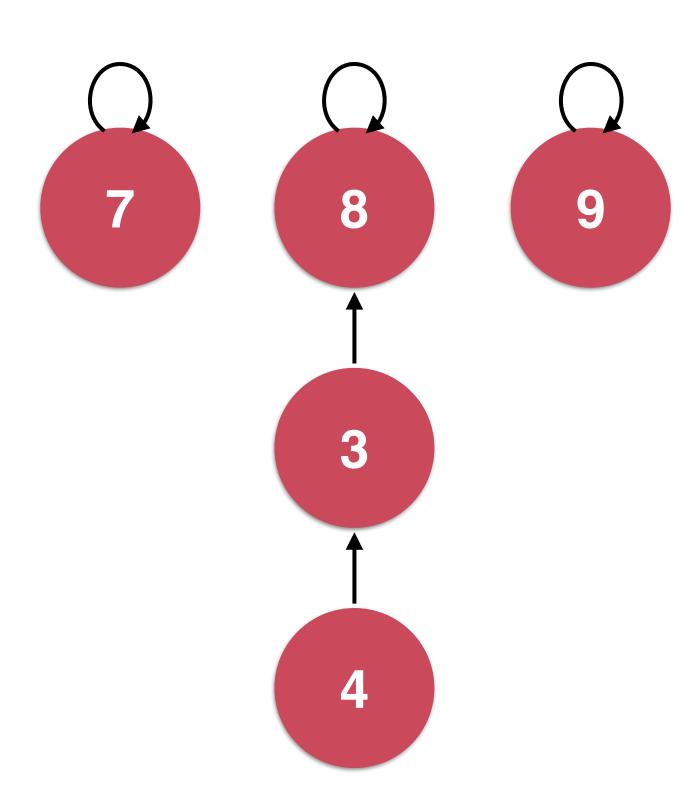


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 5 7 8 9

## union 9,4



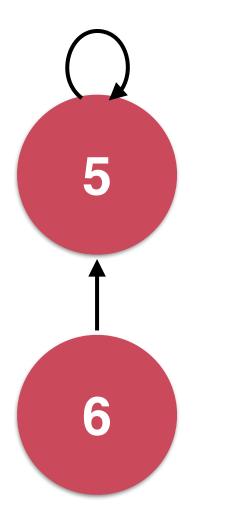


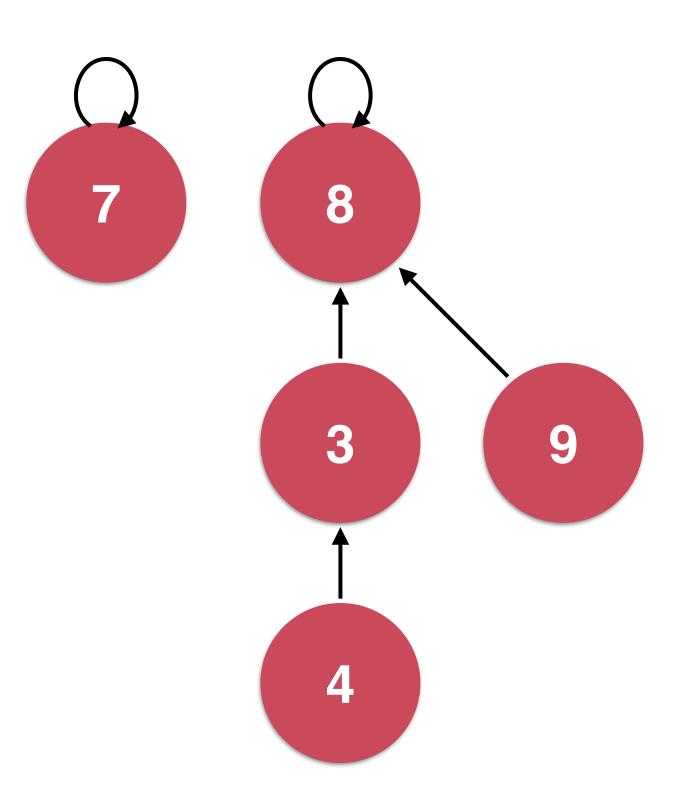


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 5 7 8 9

# union 9,4



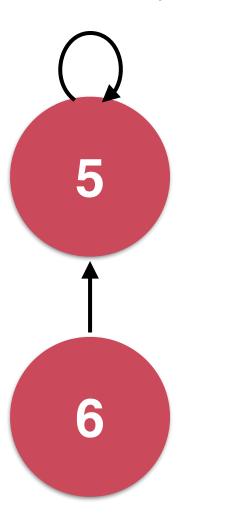


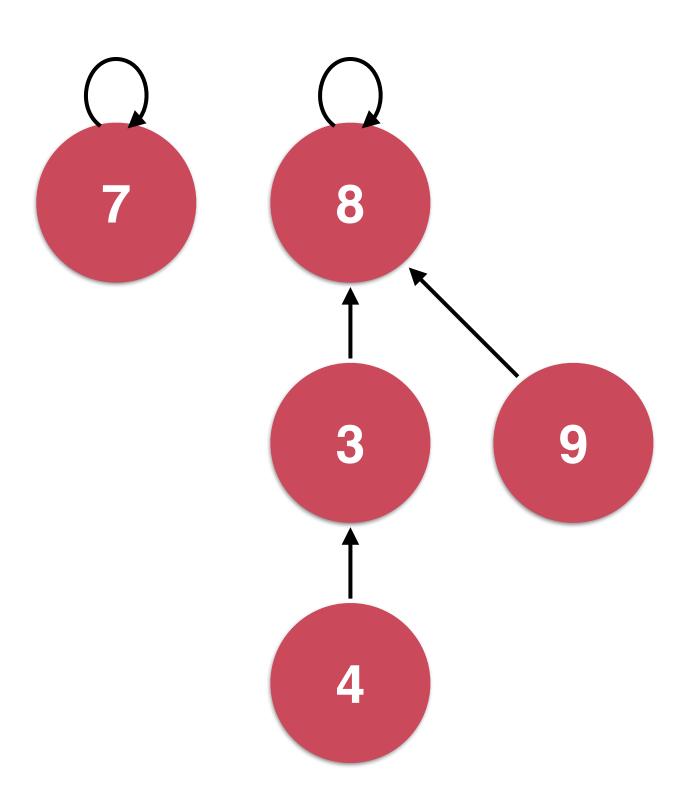


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 5 7 8 9

## union 9,4



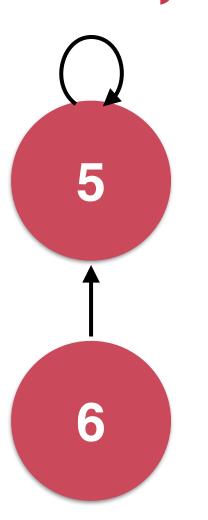


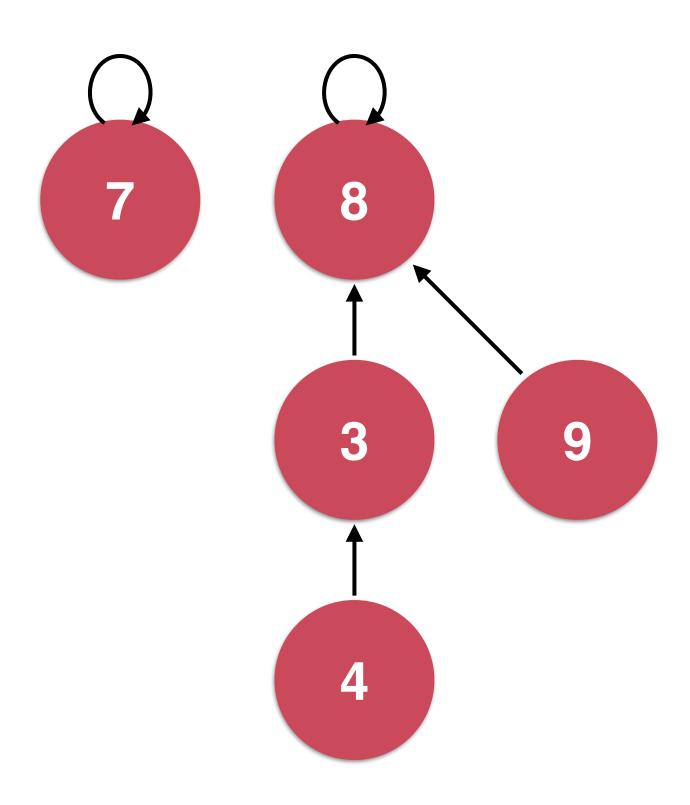


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 5 7 8 8

## union 2, 1

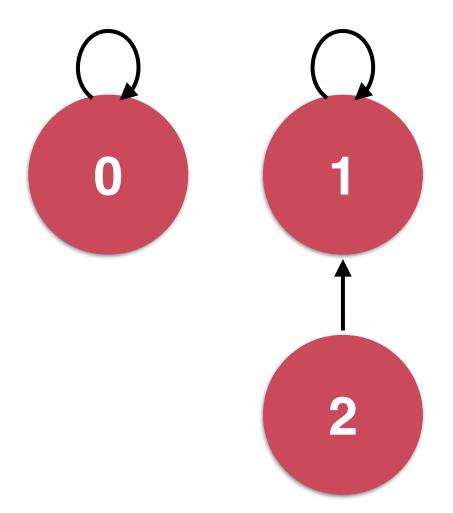


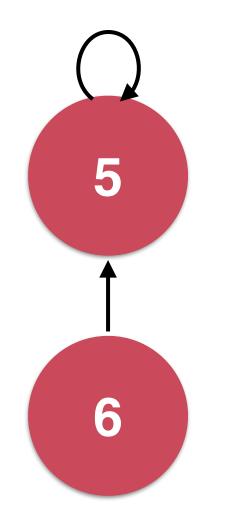


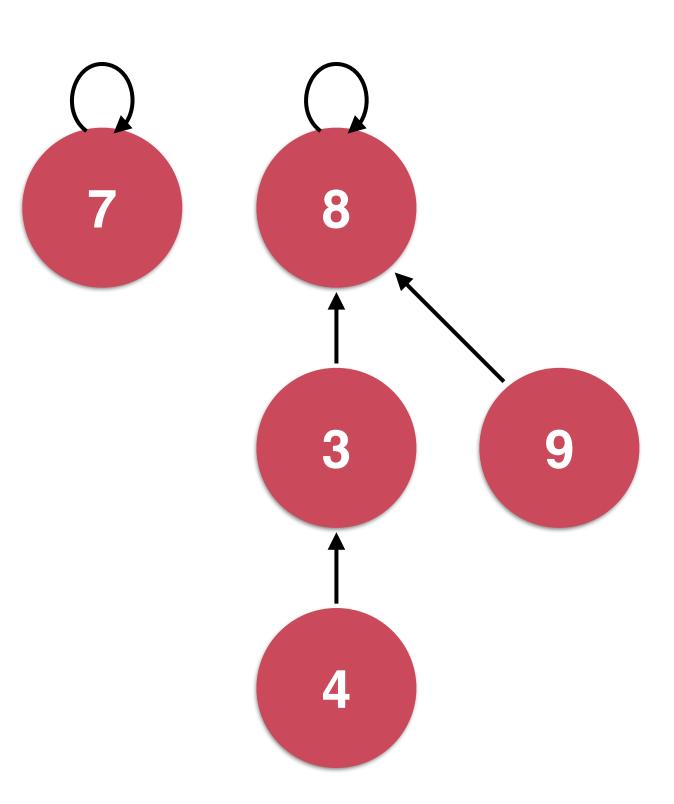


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 5 7 8 8

## union 2, 1

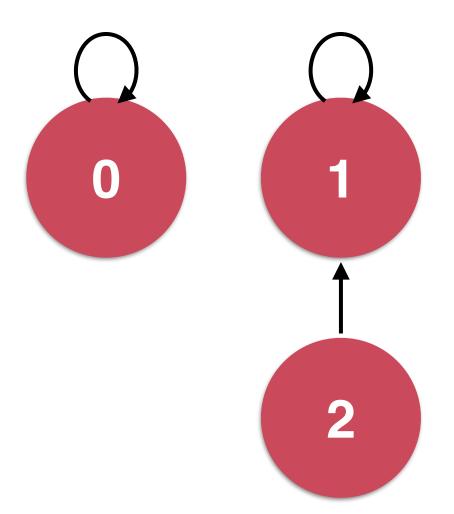


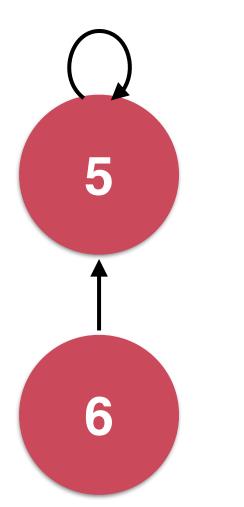


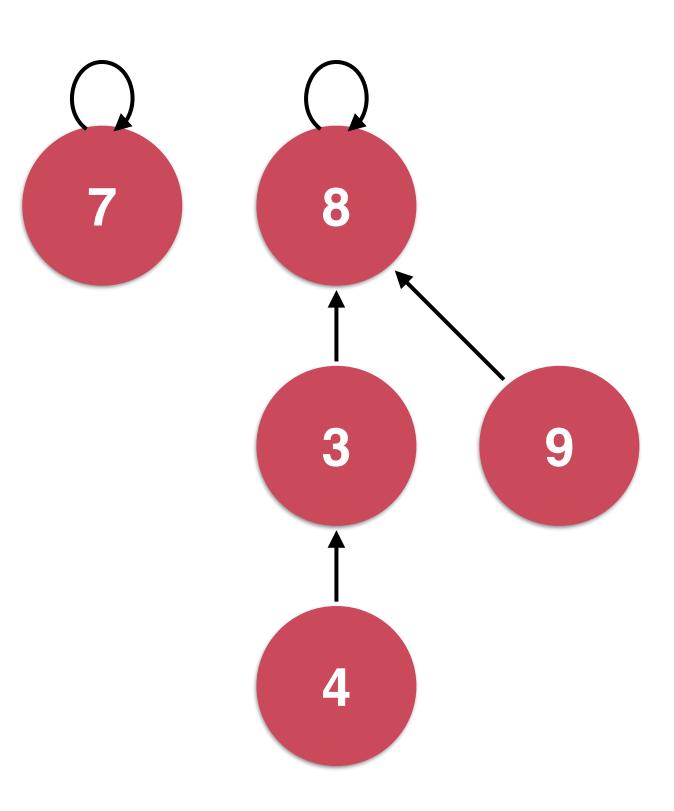


parent 0 1 2 3 4 5 6 7 8 9 0 1 2 8 3 5 5 7 8 8

## union 2, 1

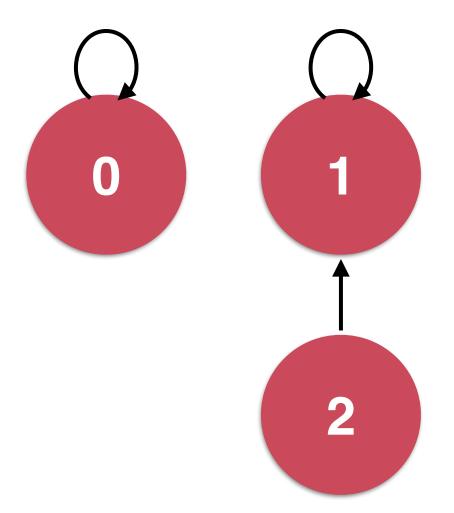


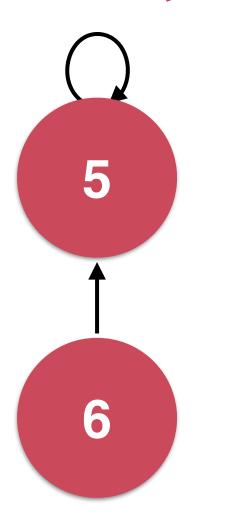


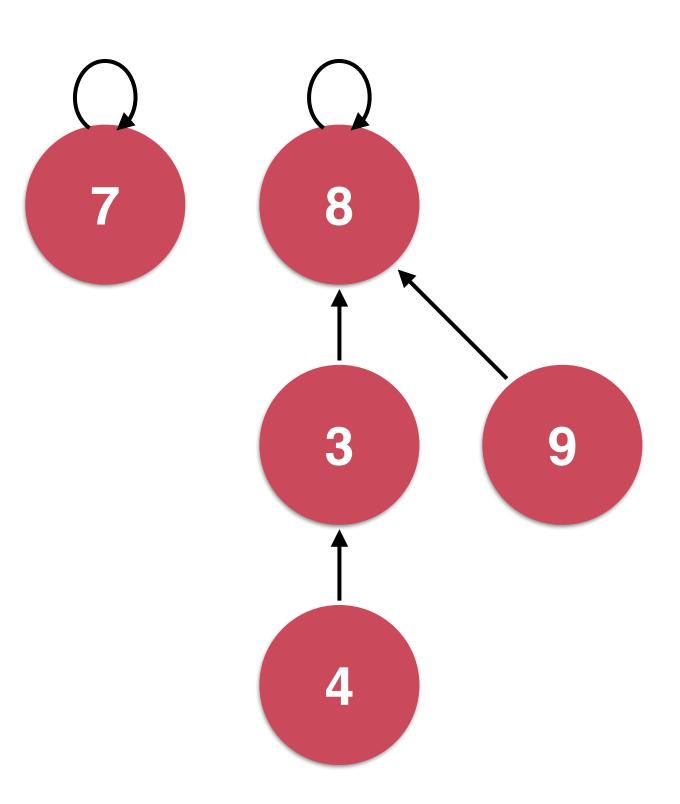


parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 5 5 7 8 8

## union 5, 0

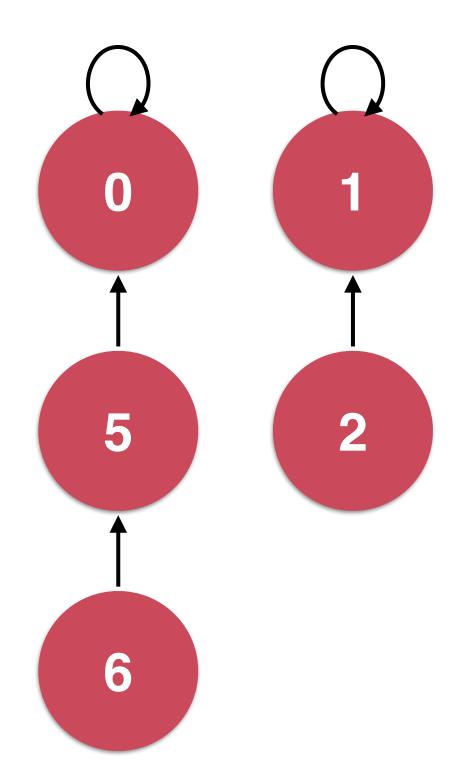


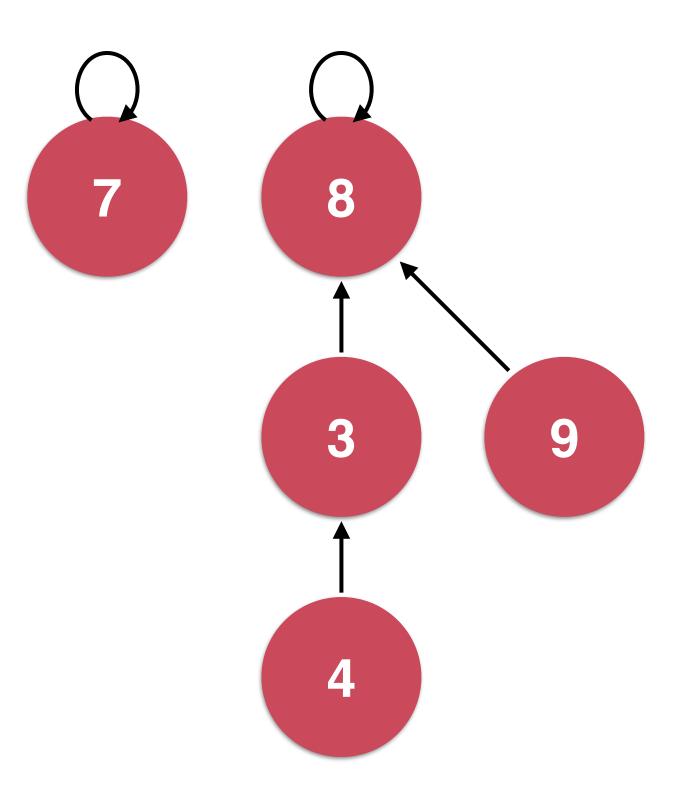




parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 5 5 7 8 8

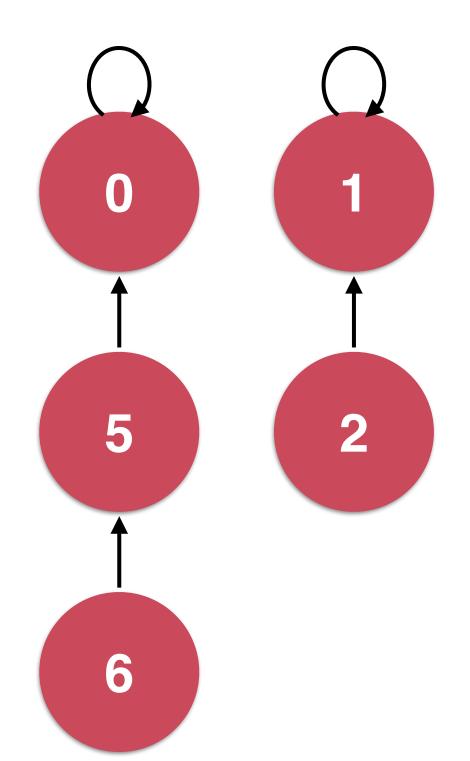
## union 5, 0

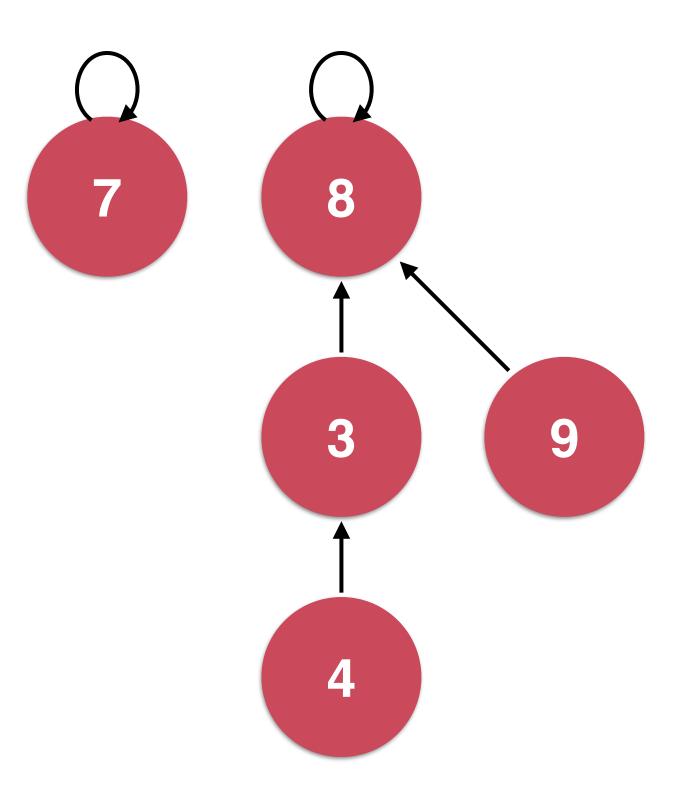




parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 5 5 7 8 8

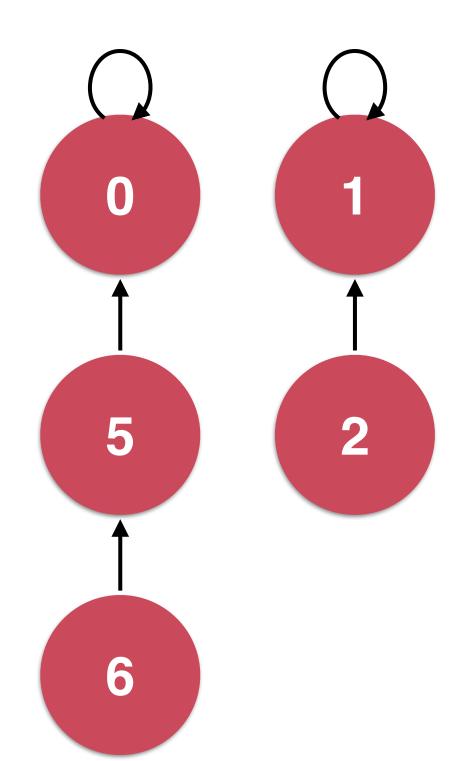
## union 5, 0

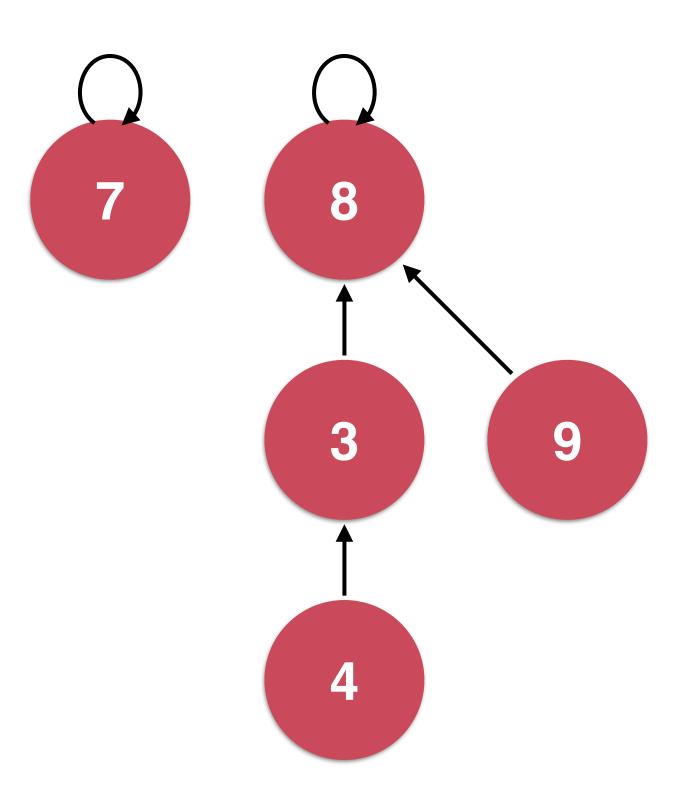




parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 0 5 7 8 8

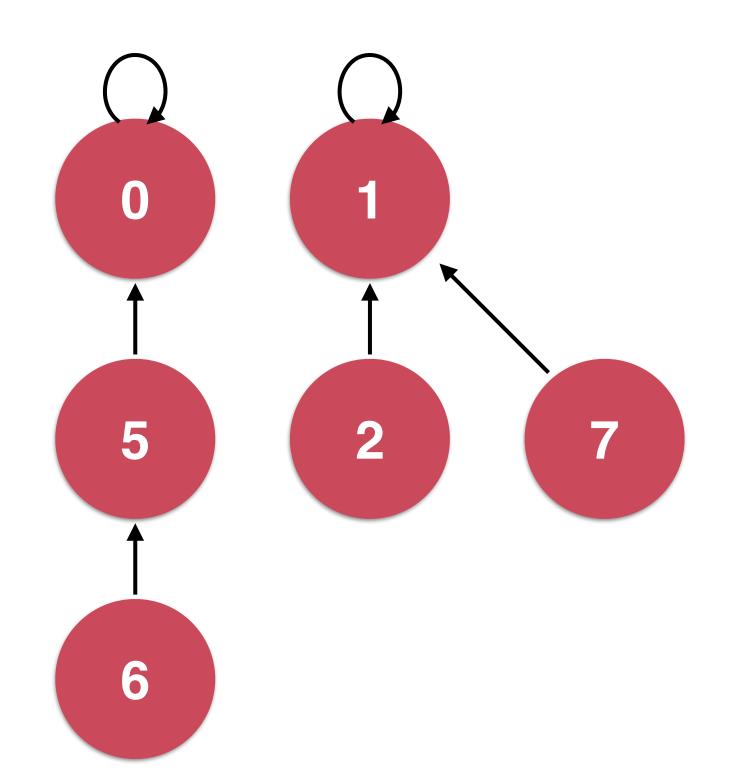
### union 7, 2

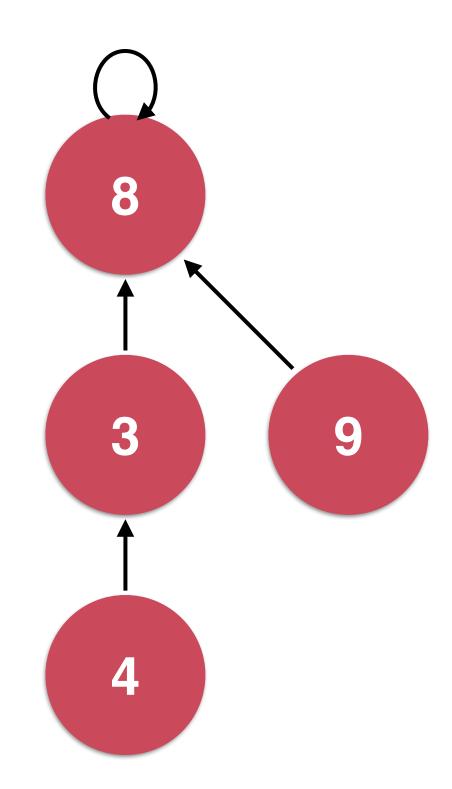




parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 0 5 7 8 8

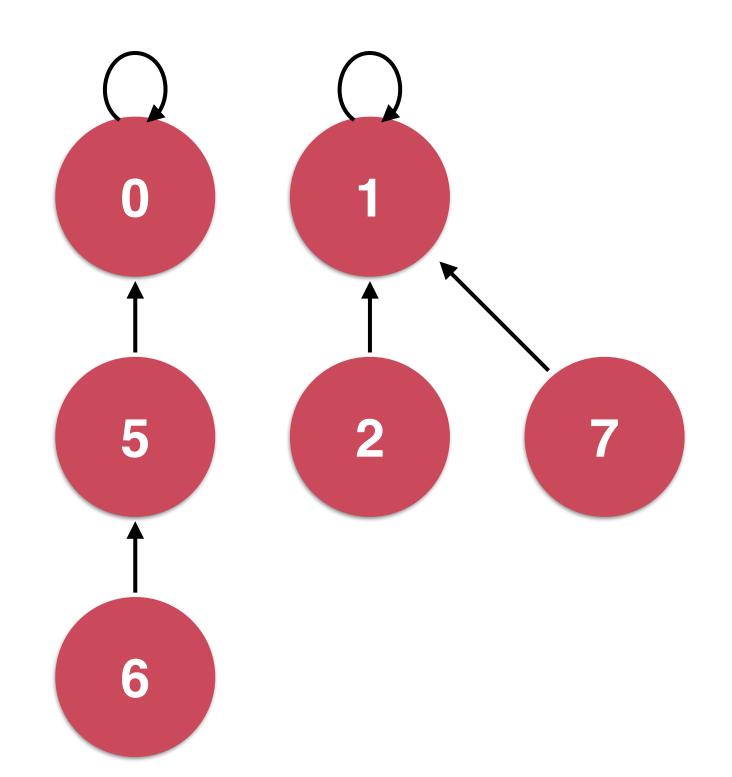
### union 7, 2

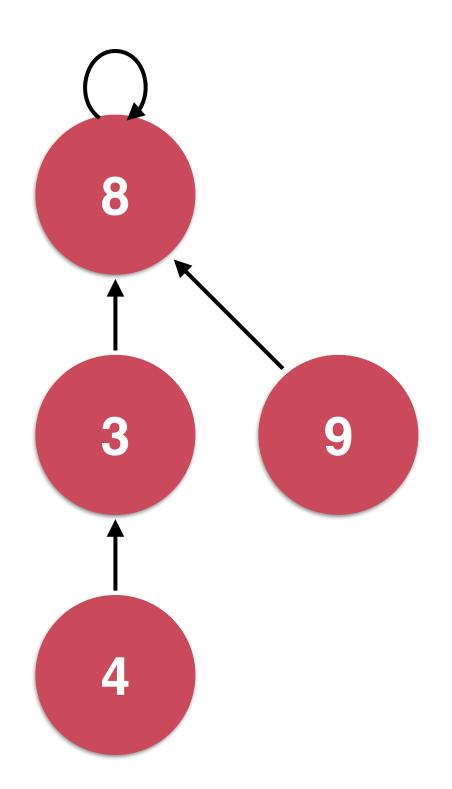




parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 0 5 7 8 8

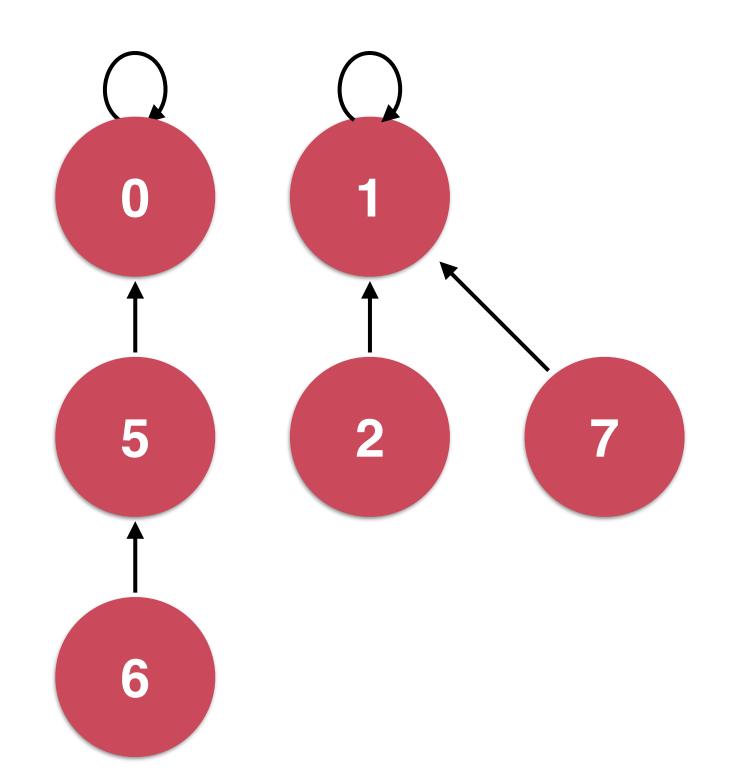
### union 7, 2

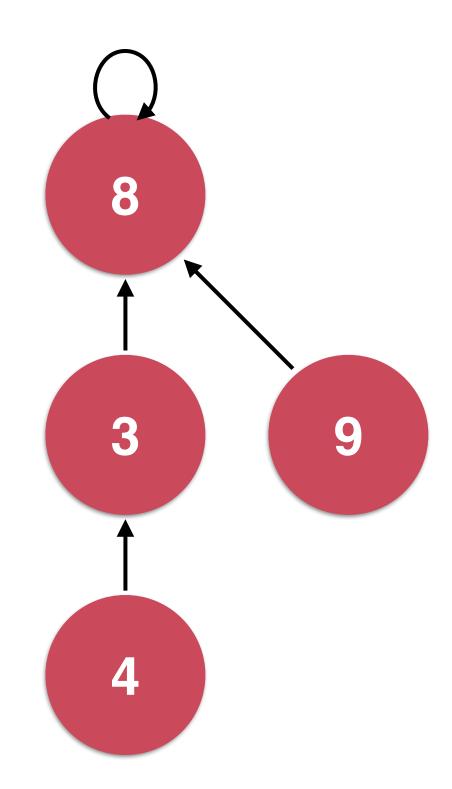




parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 0 5 1 8 8

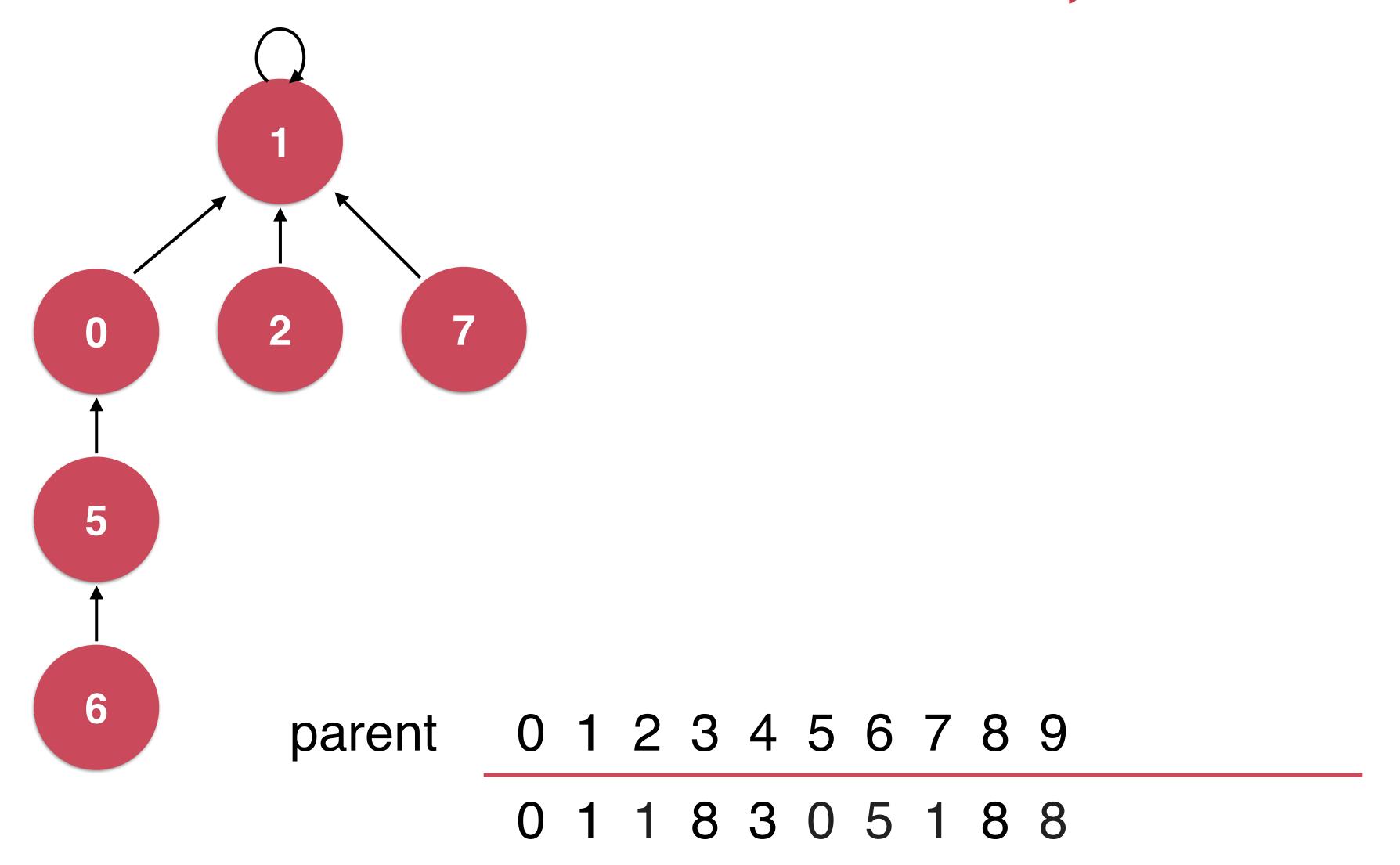
## union 6, 2

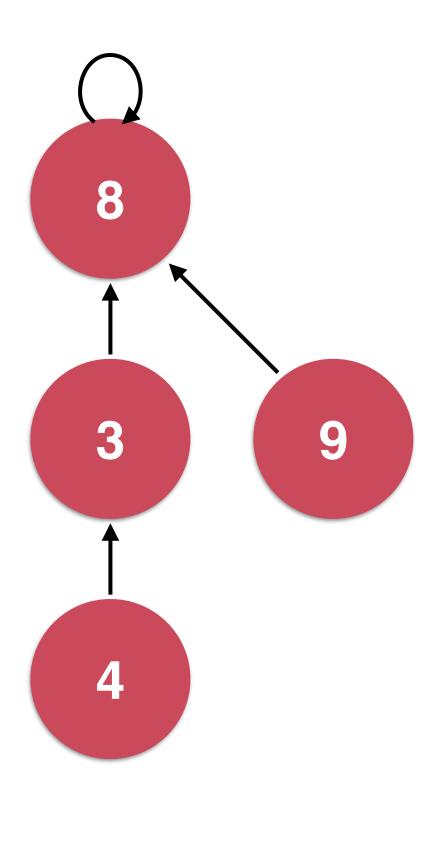




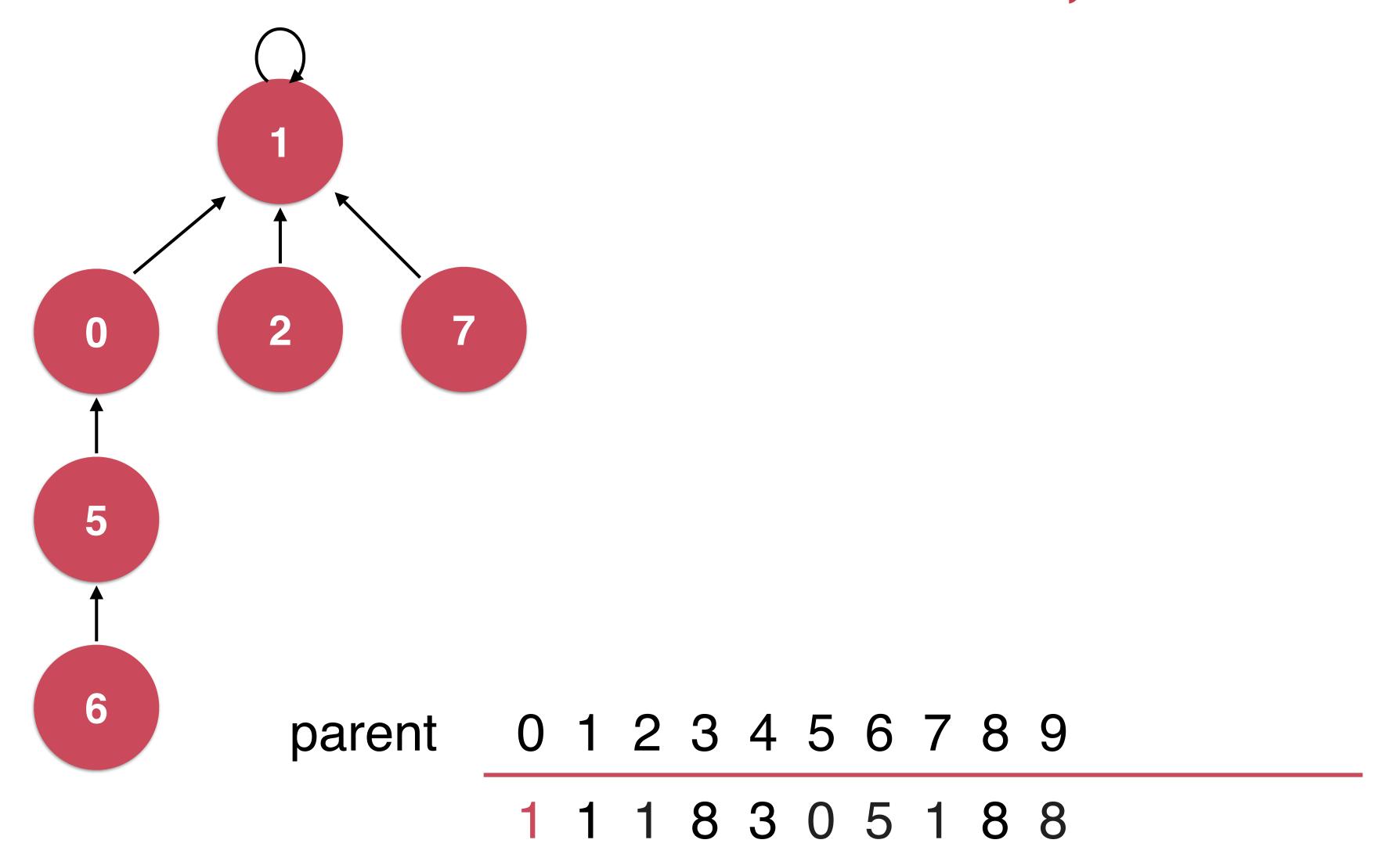
parent 0 1 2 3 4 5 6 7 8 9 0 1 1 8 3 0 5 1 8 8

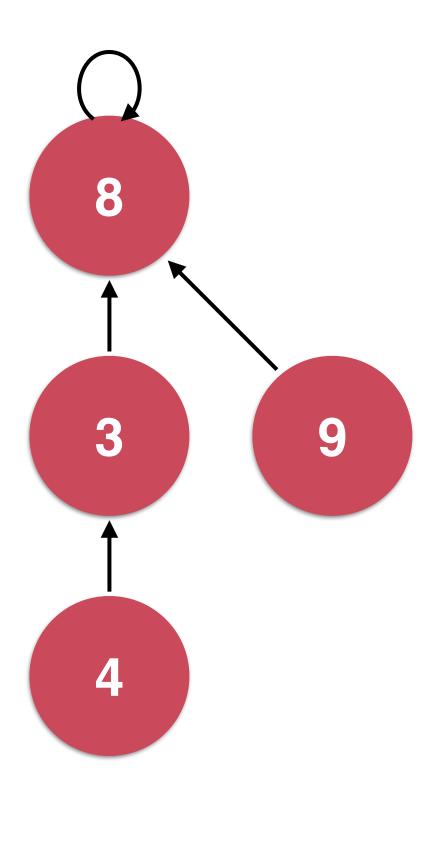
## union 6, 2



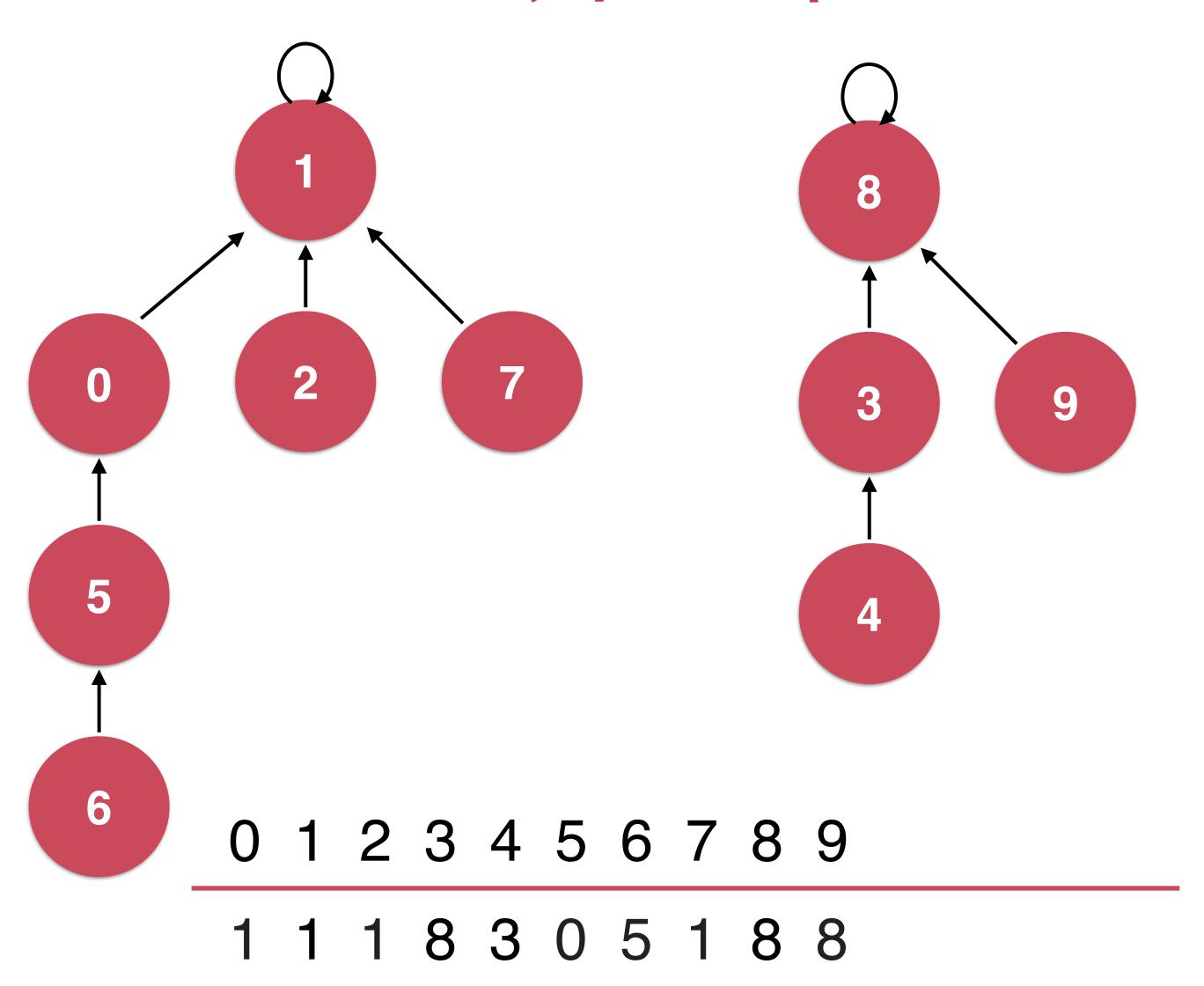


## union 6, 2





## 并查集



## 实践: Quick Union 的并查集实现

# 并查集的优化

实践: Quick Union 和 Quick Find 的时间效率比较

## 之前Quick Union的问题

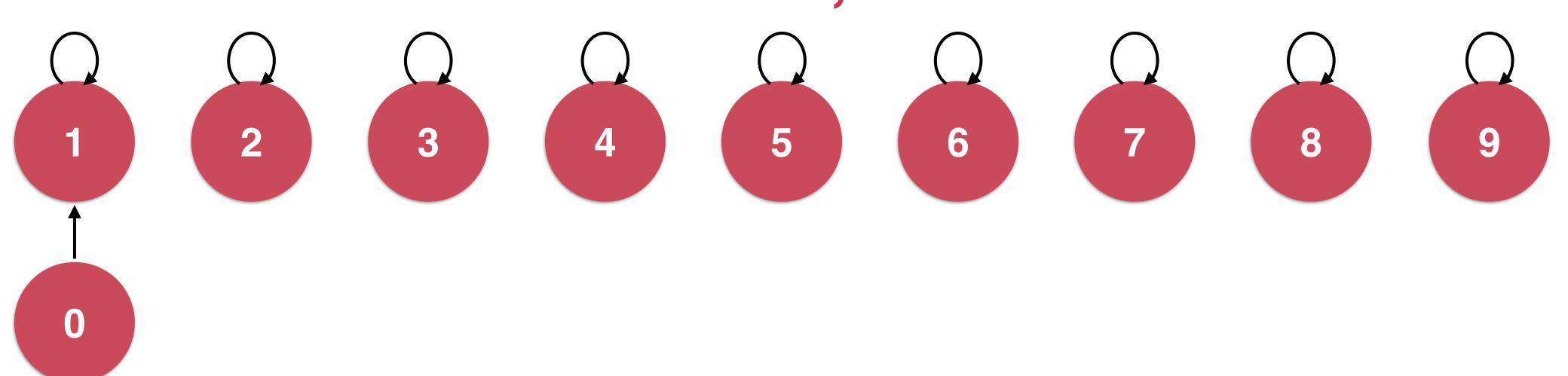
#### Quick Union



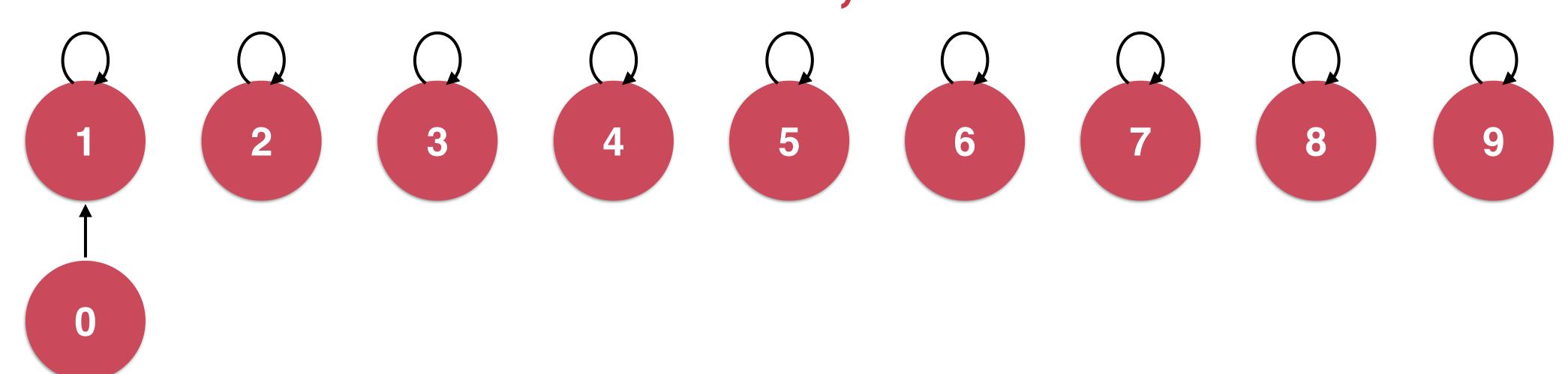
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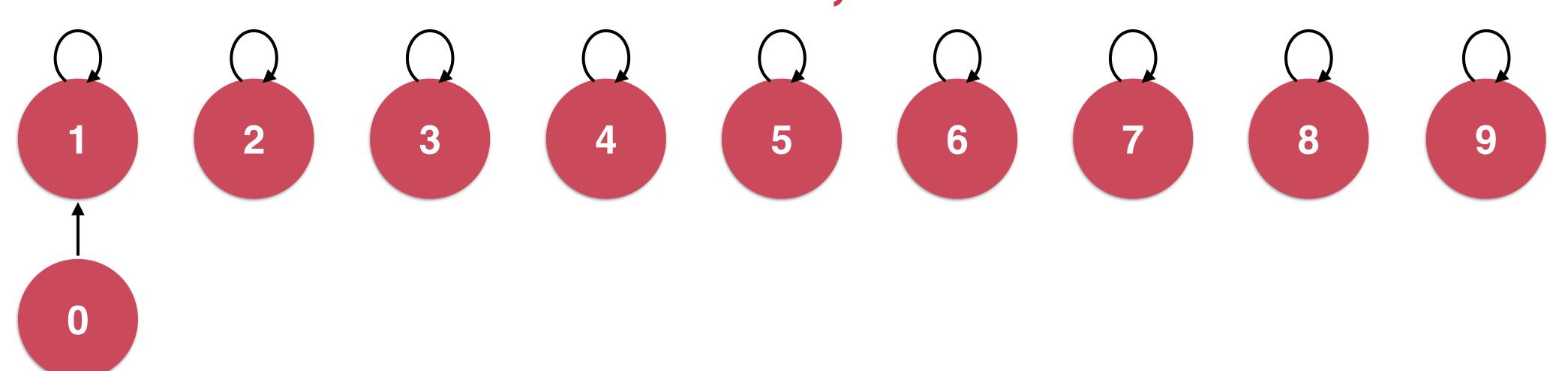
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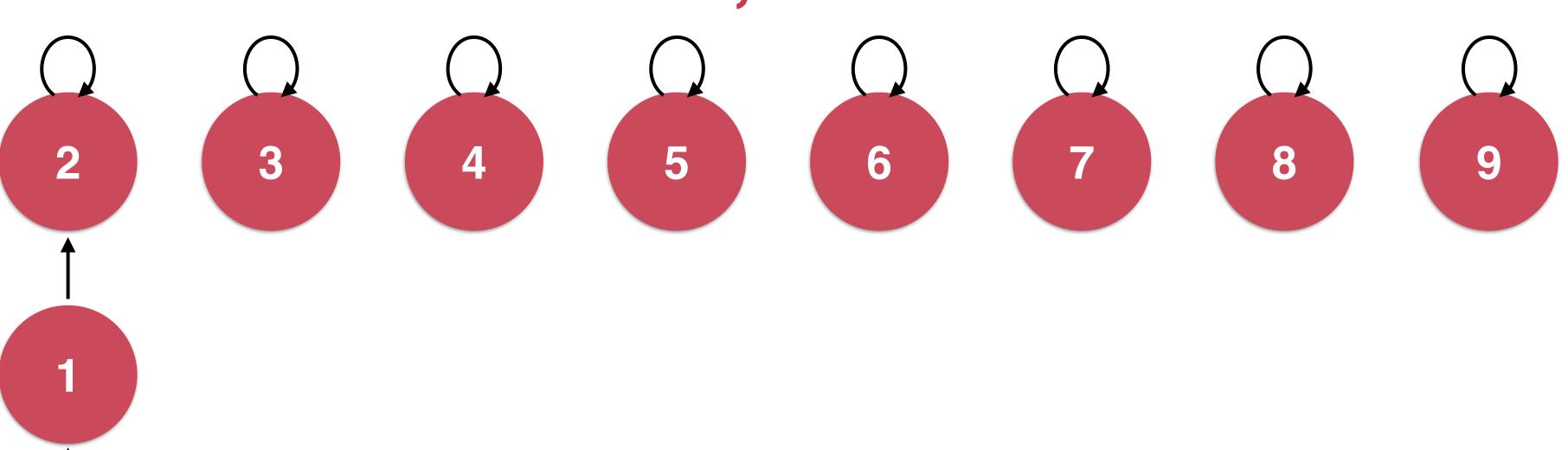
parent 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9



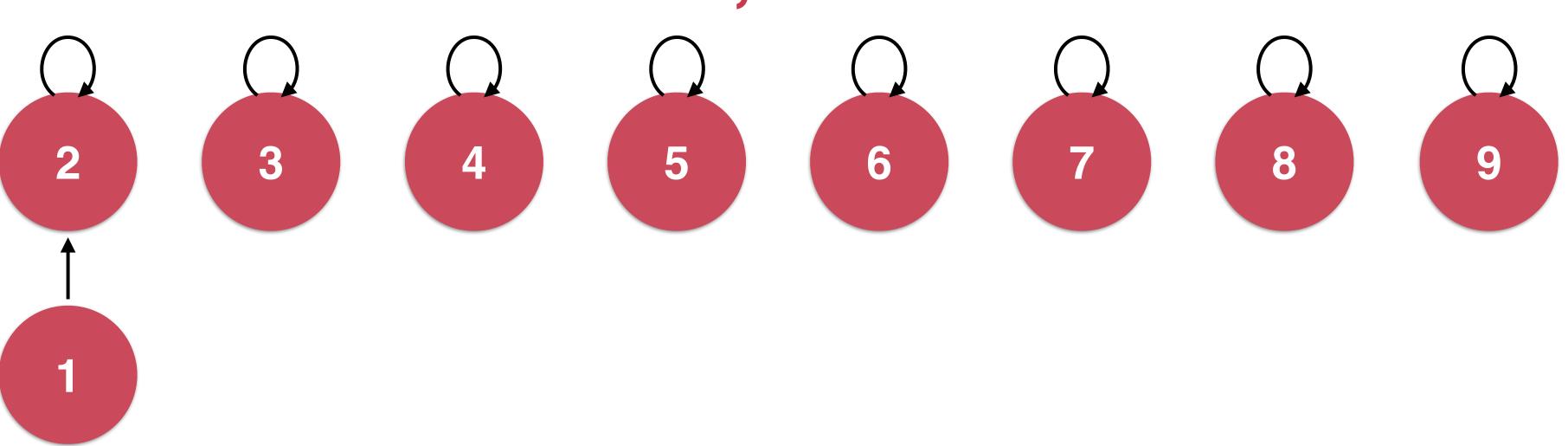
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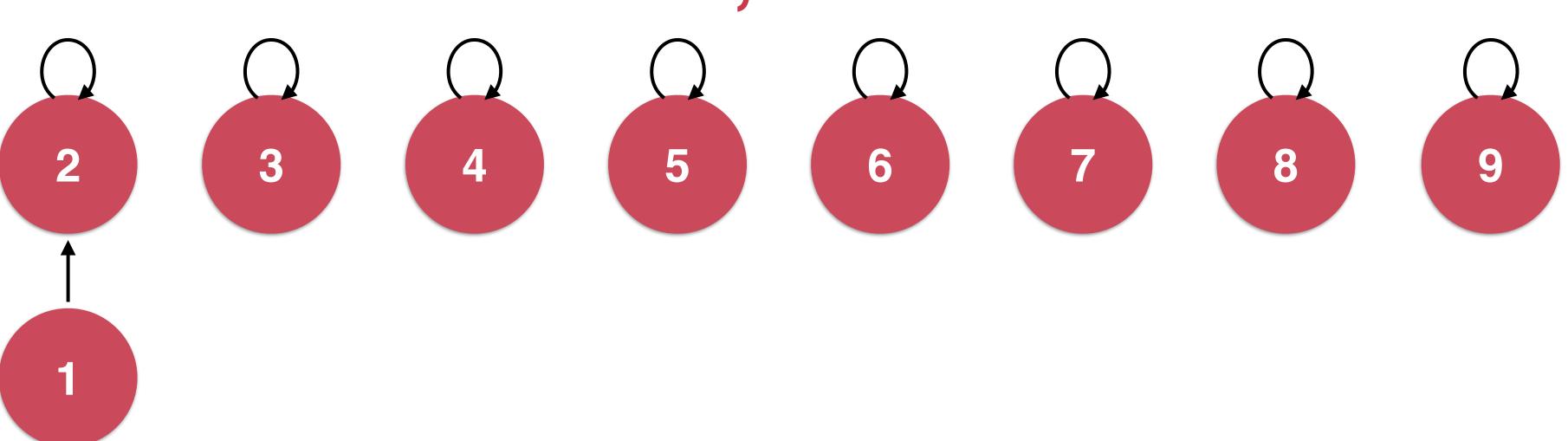
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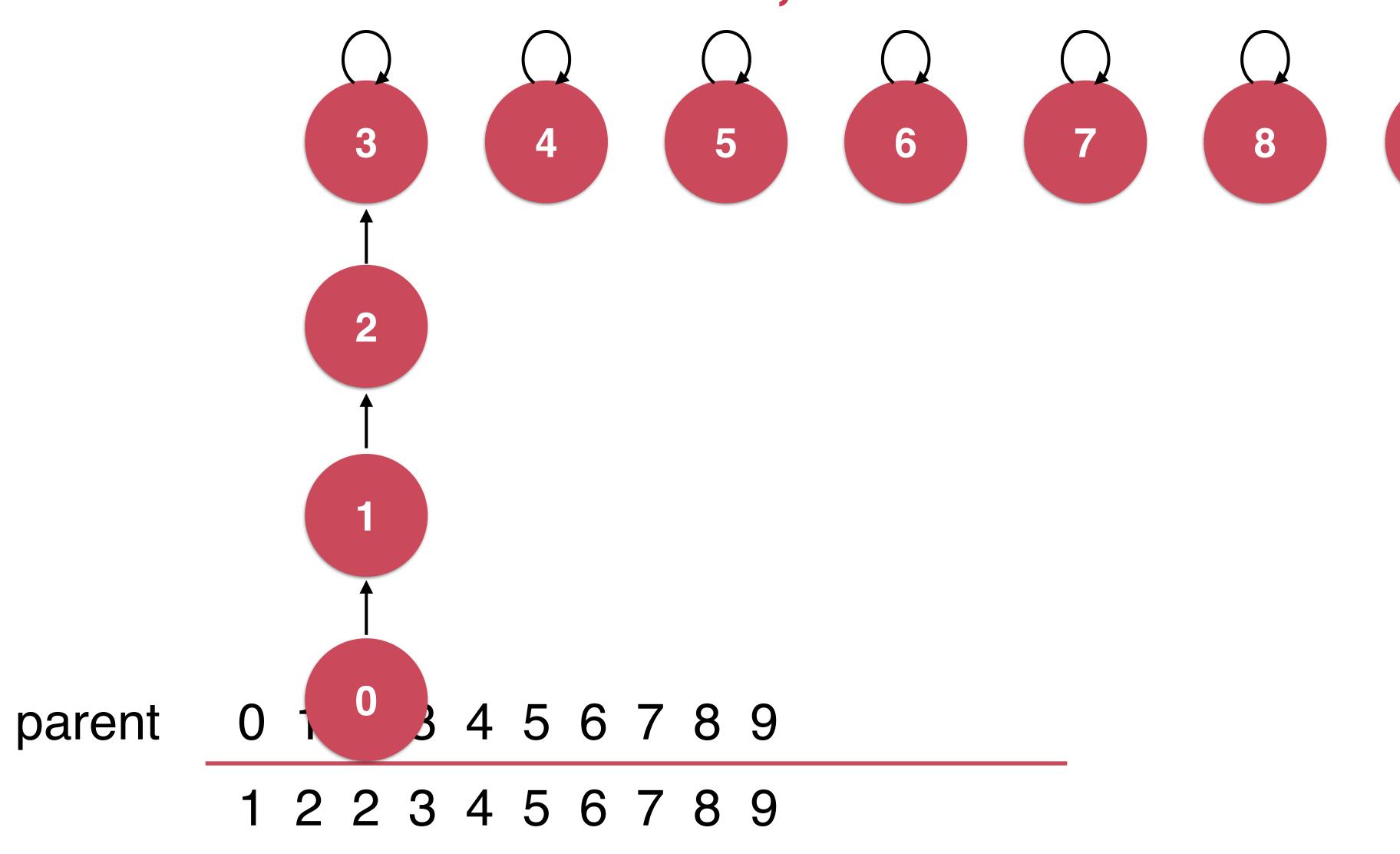


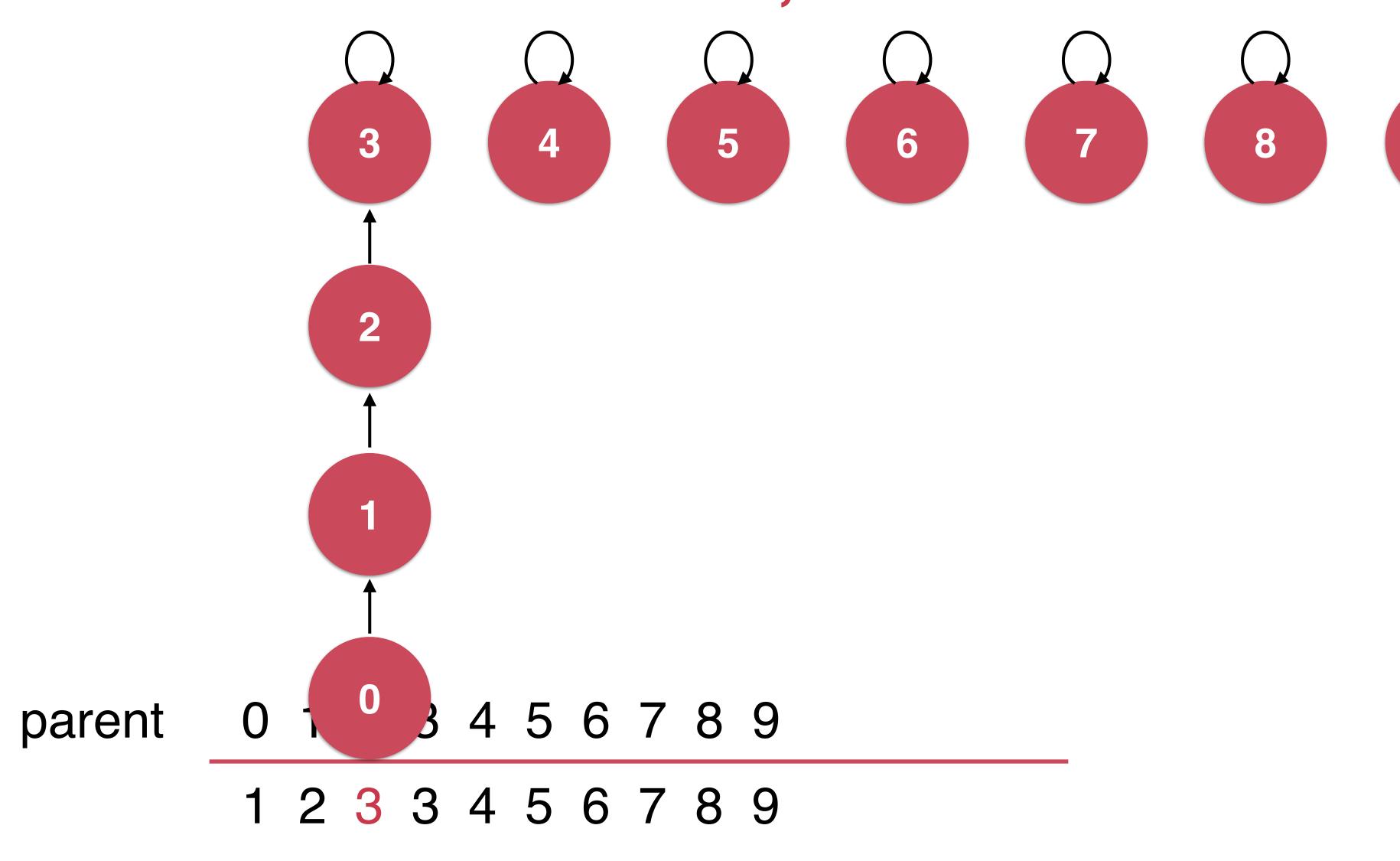




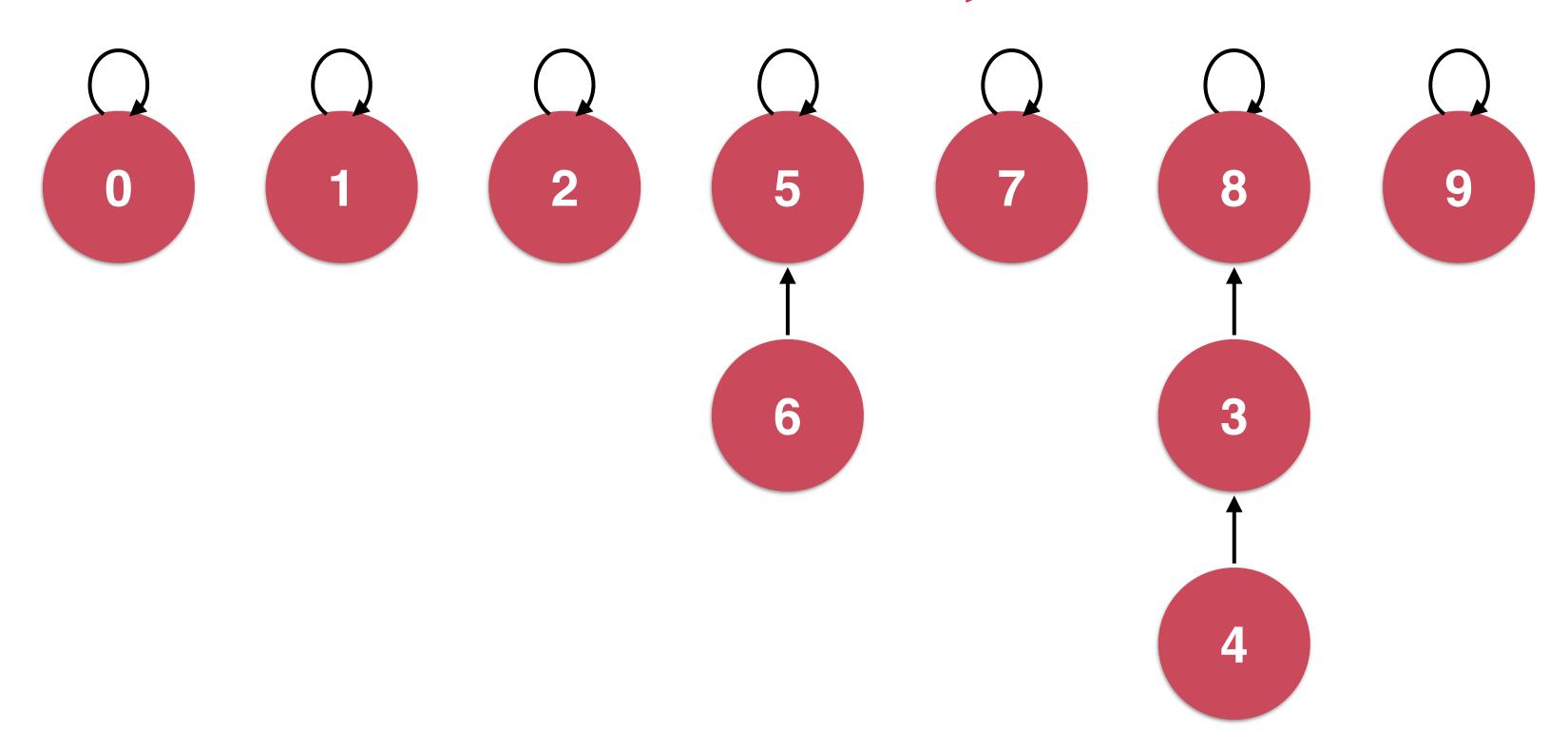




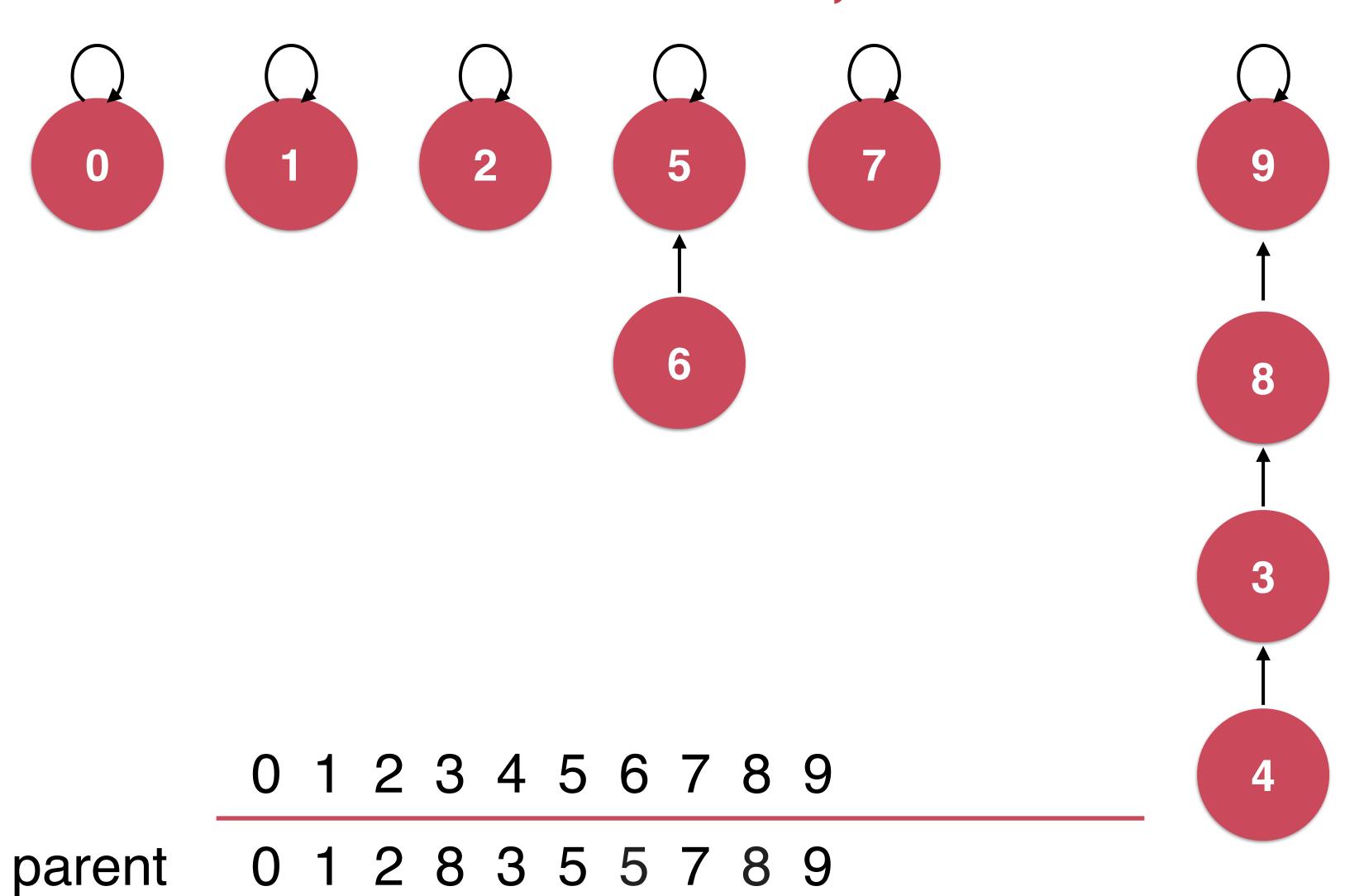


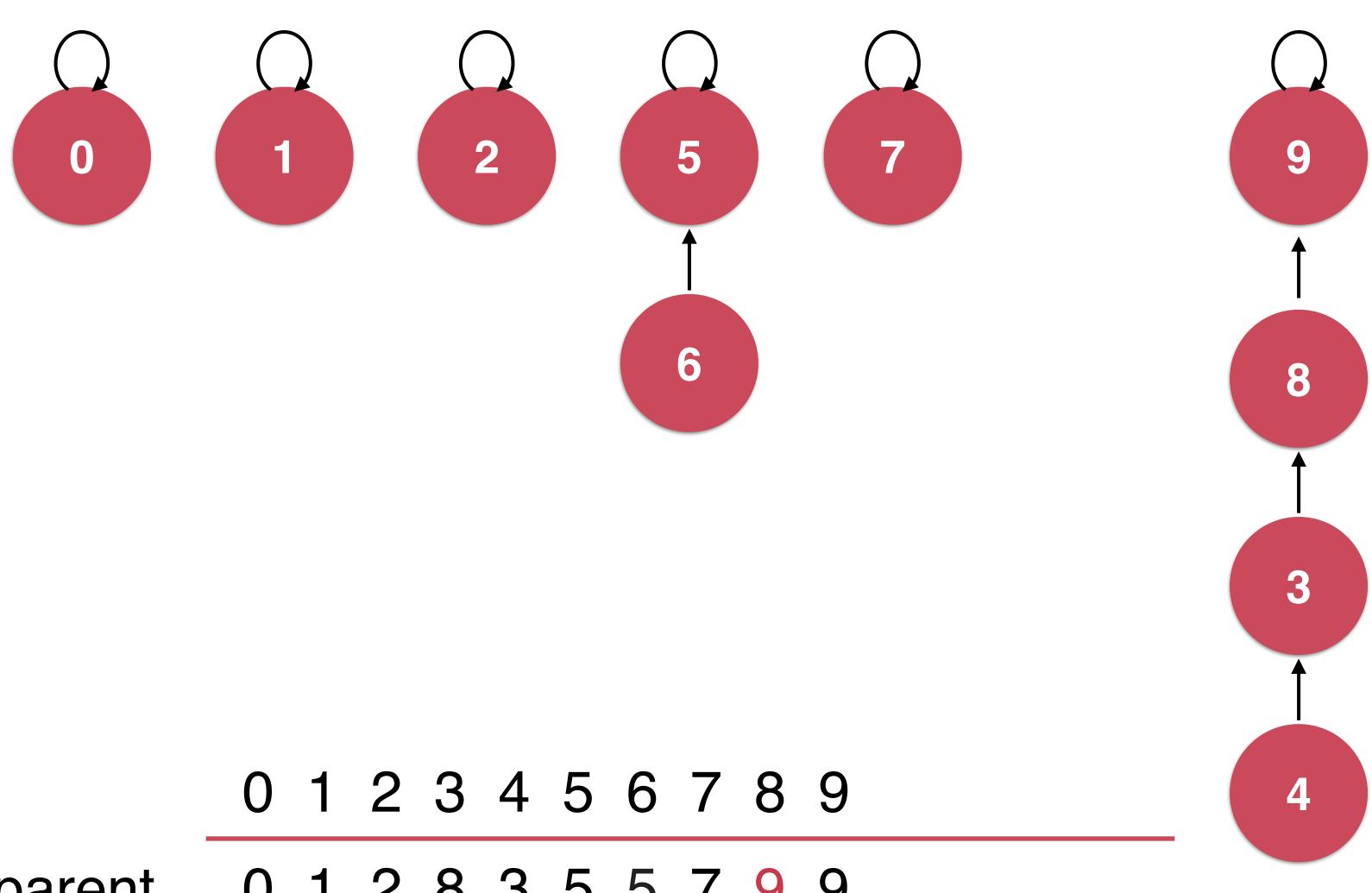


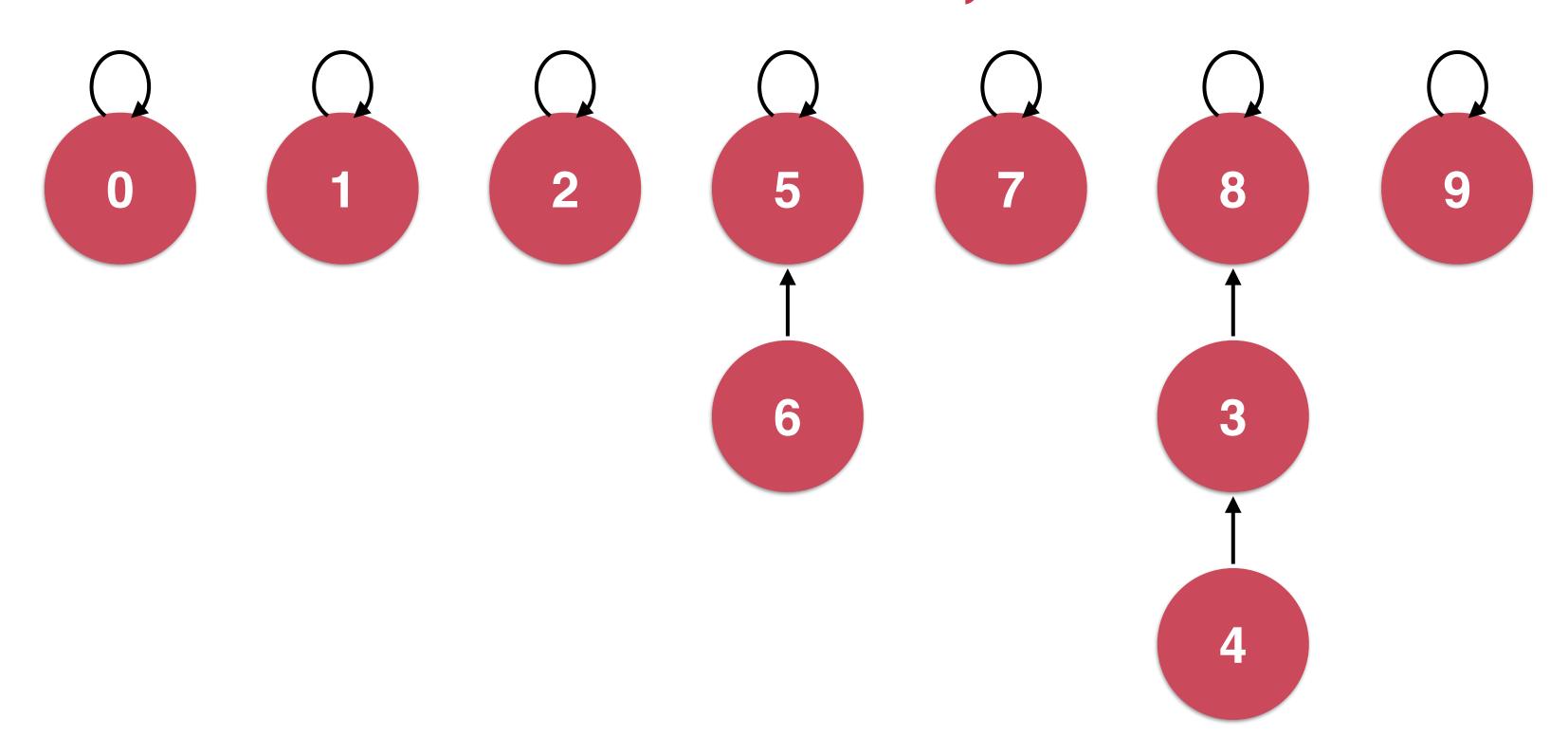
## 解决方案:考虑size



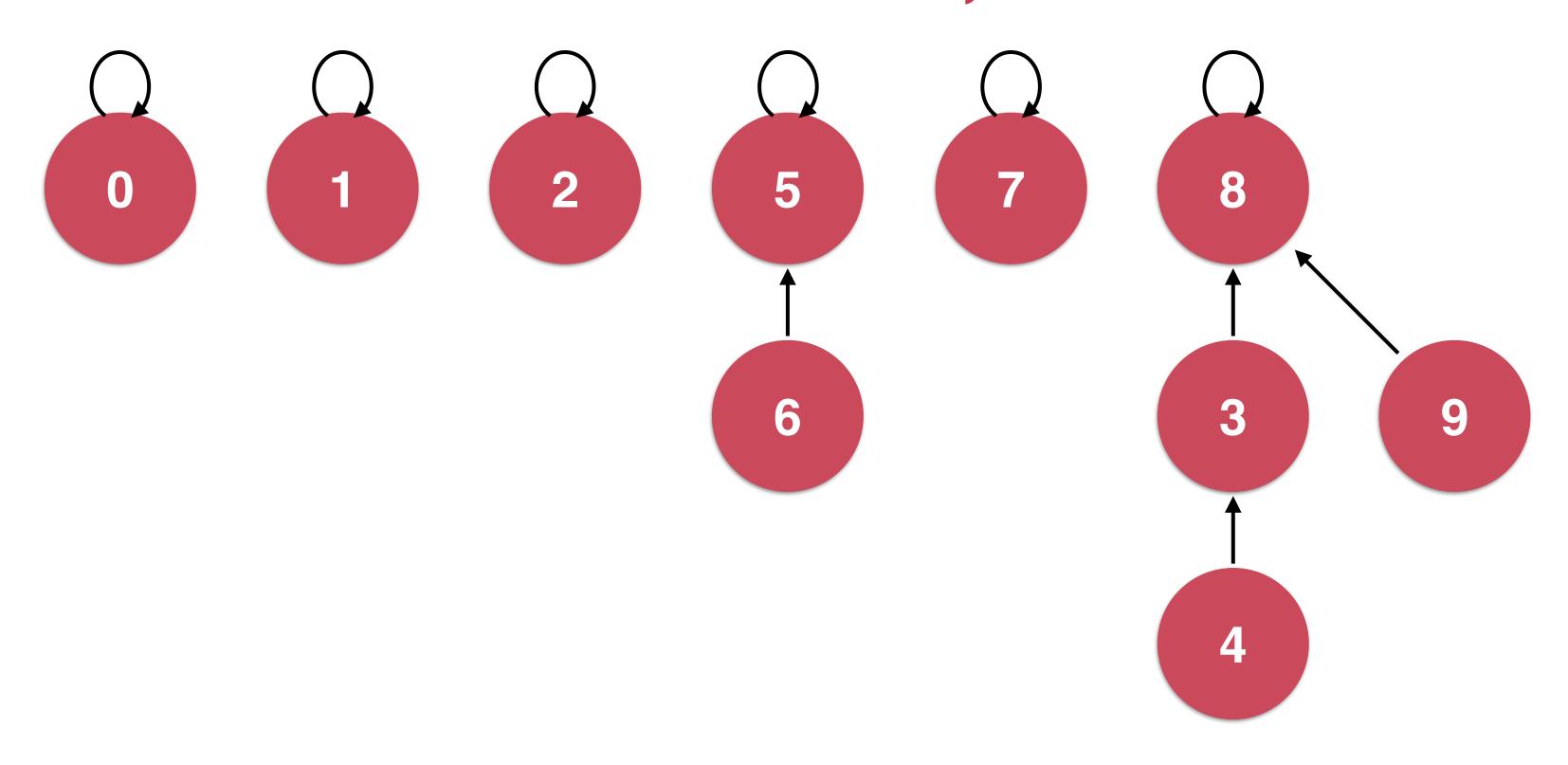
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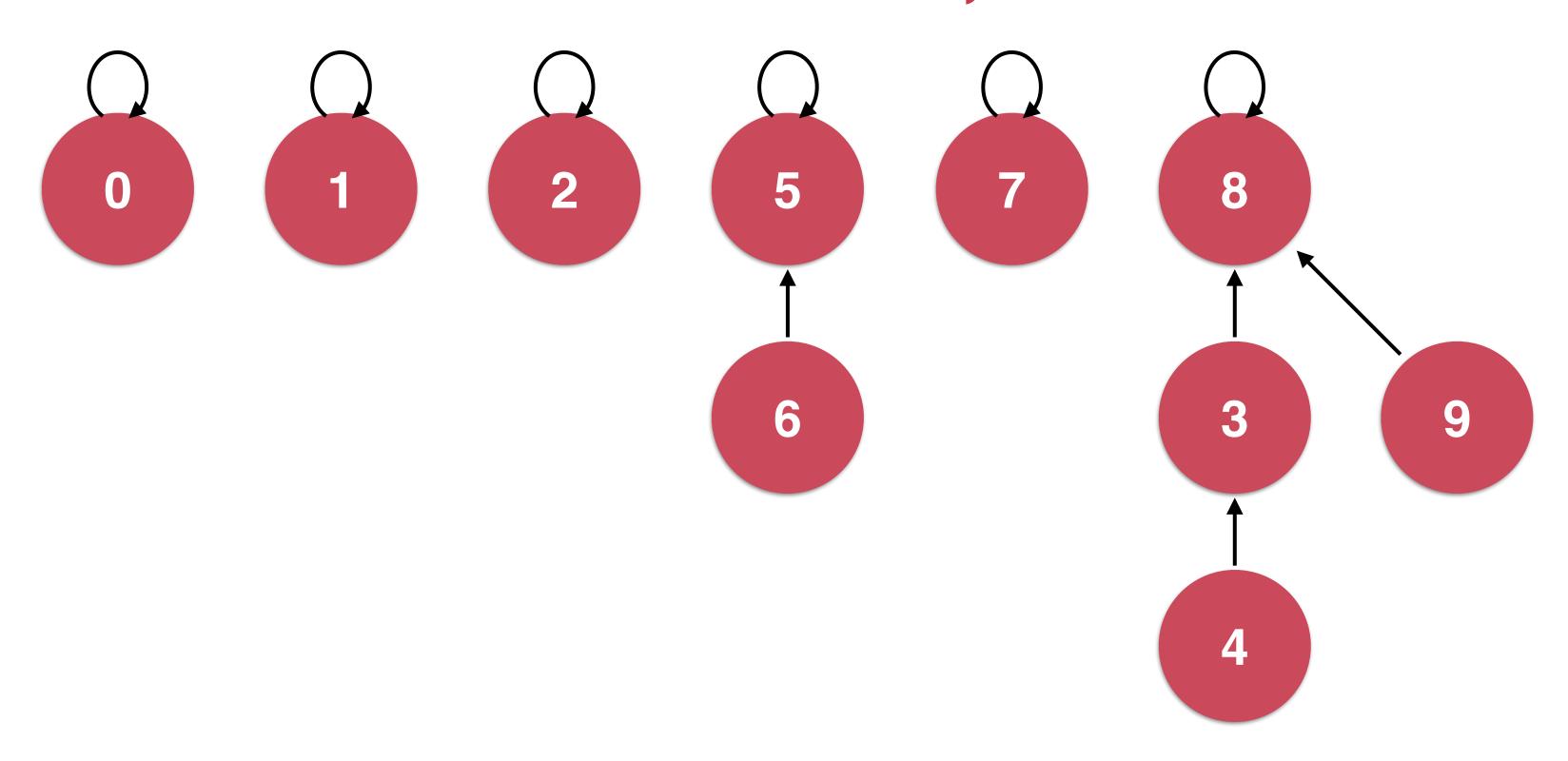




0 1 2 3 4 5 6 7 8 9



0 1 2 3 4 5 6 7 8 9

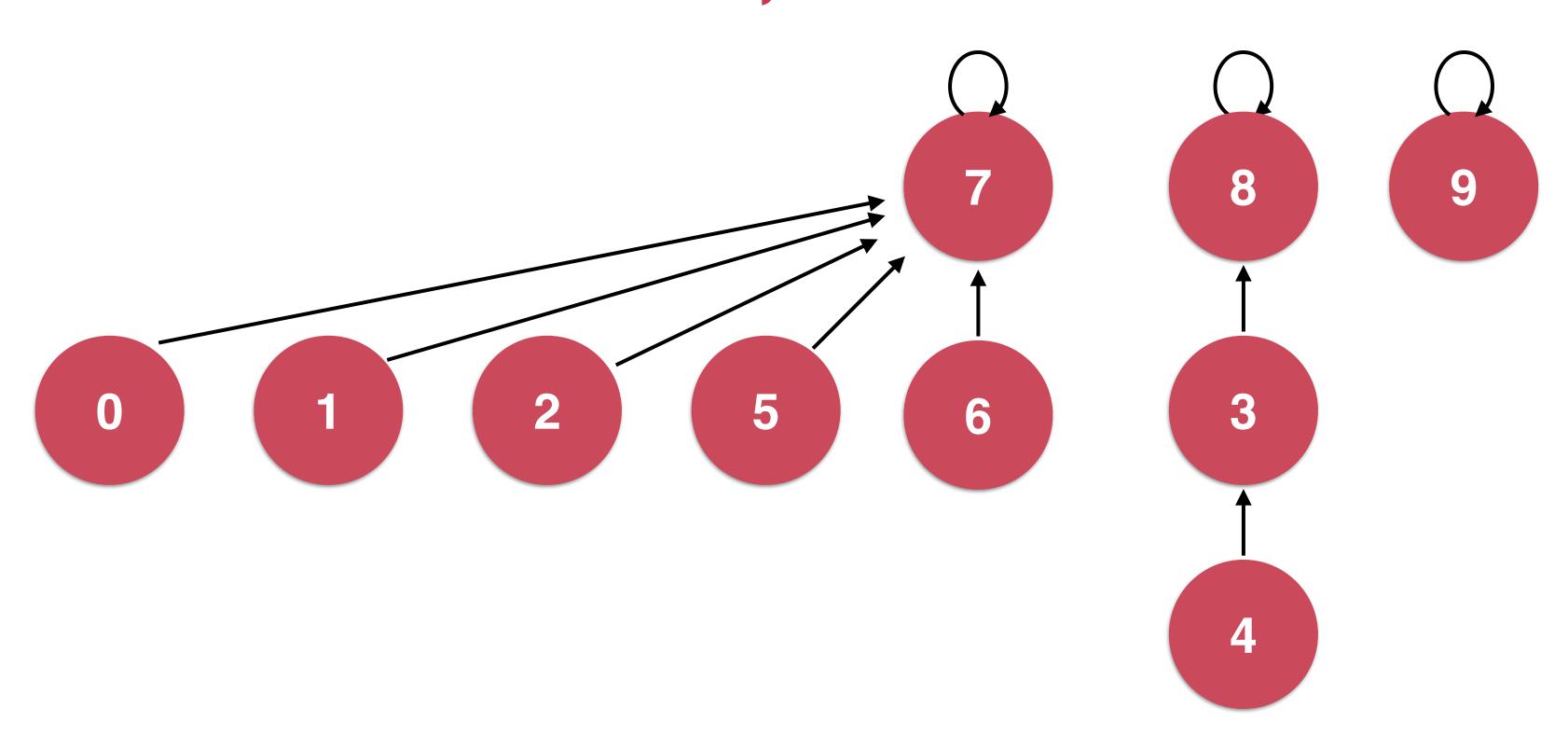


0 1 2 3 4 5 6 7 8 9

## 实践:基于size的优化

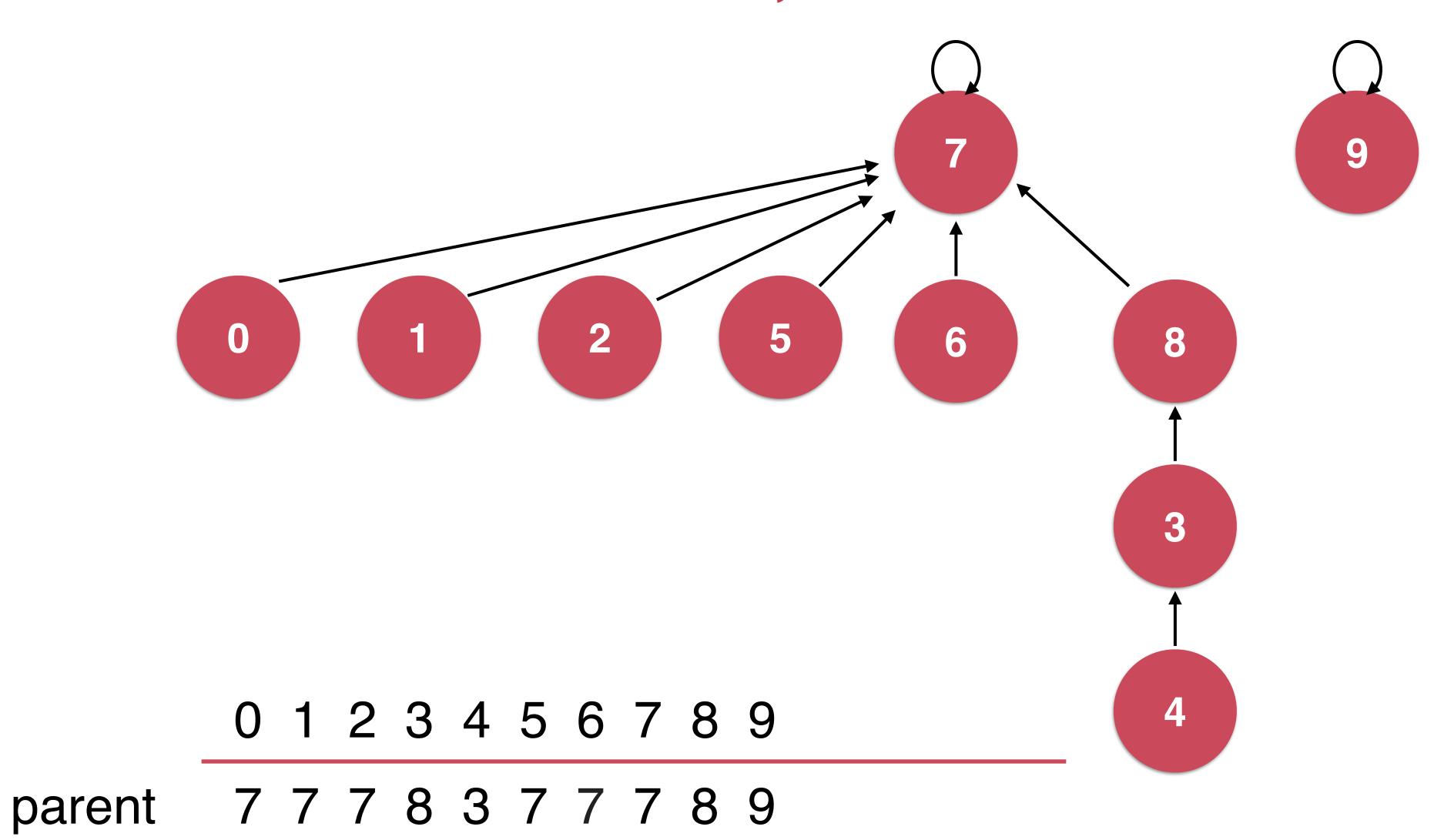
#### 实践: Quick Union 和优化后的时间效率比较

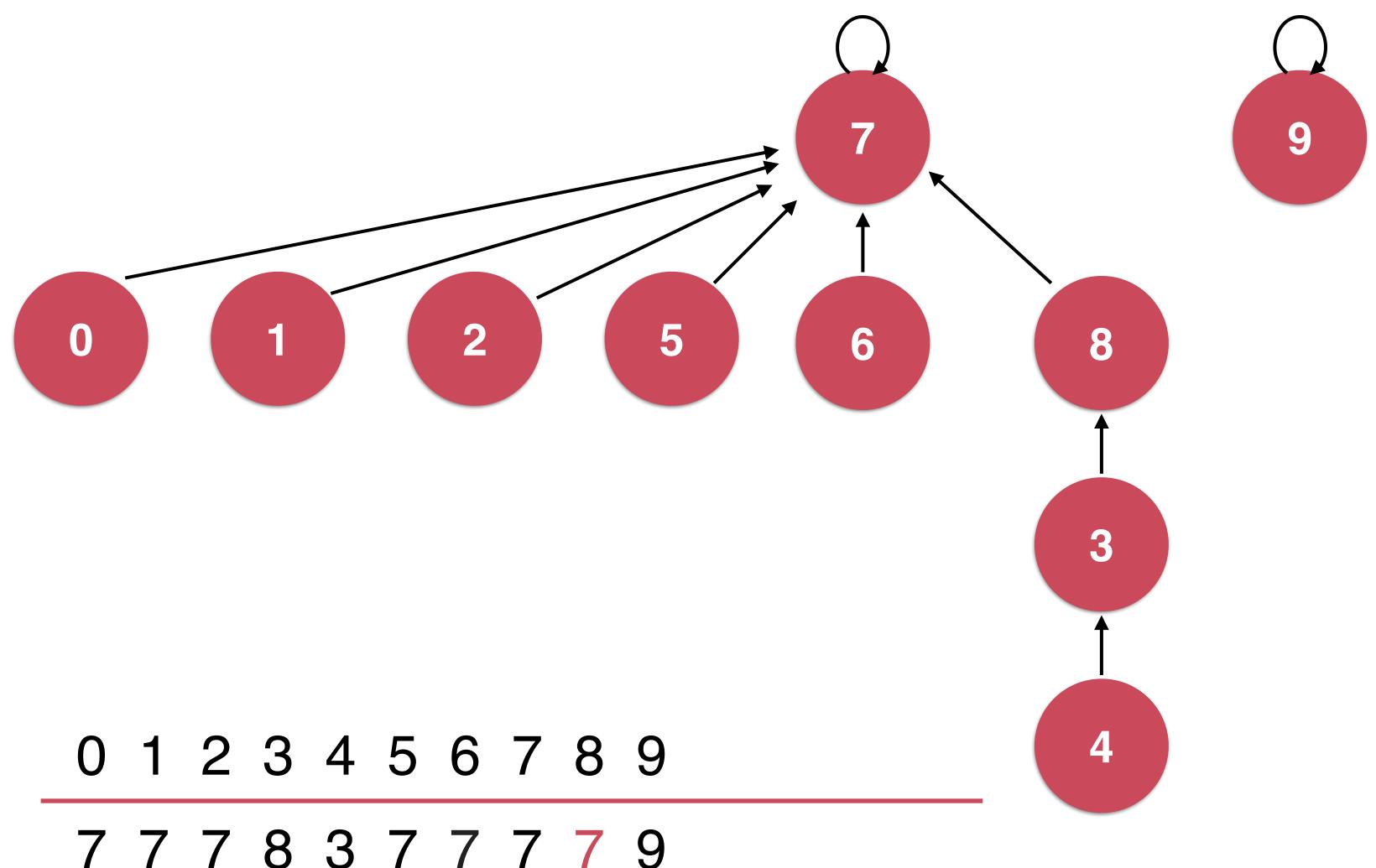
# 基于rank的优化



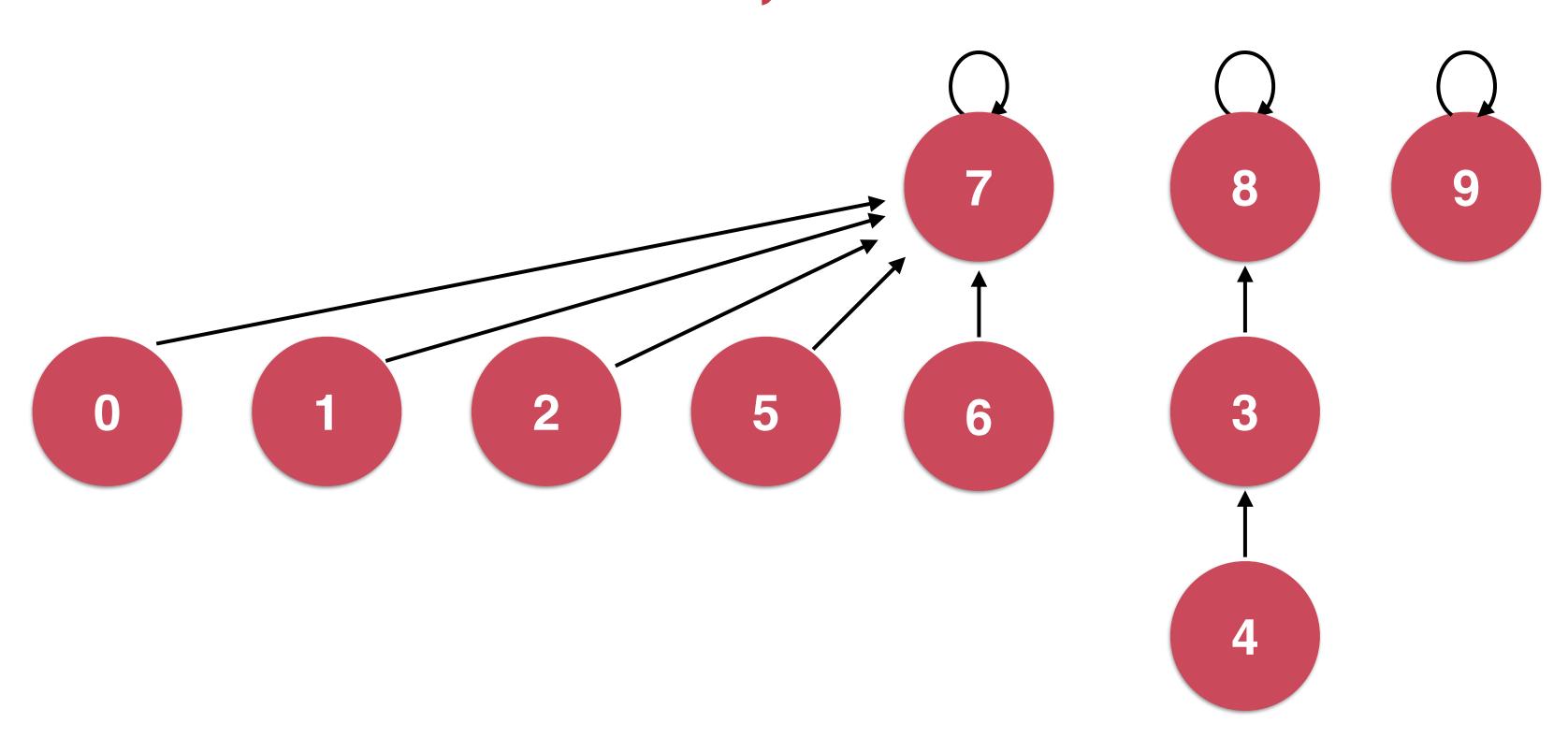
0 1 2 3 4 5 6 7 8 9

parent 7 7 7 8 3 7 7 8 9



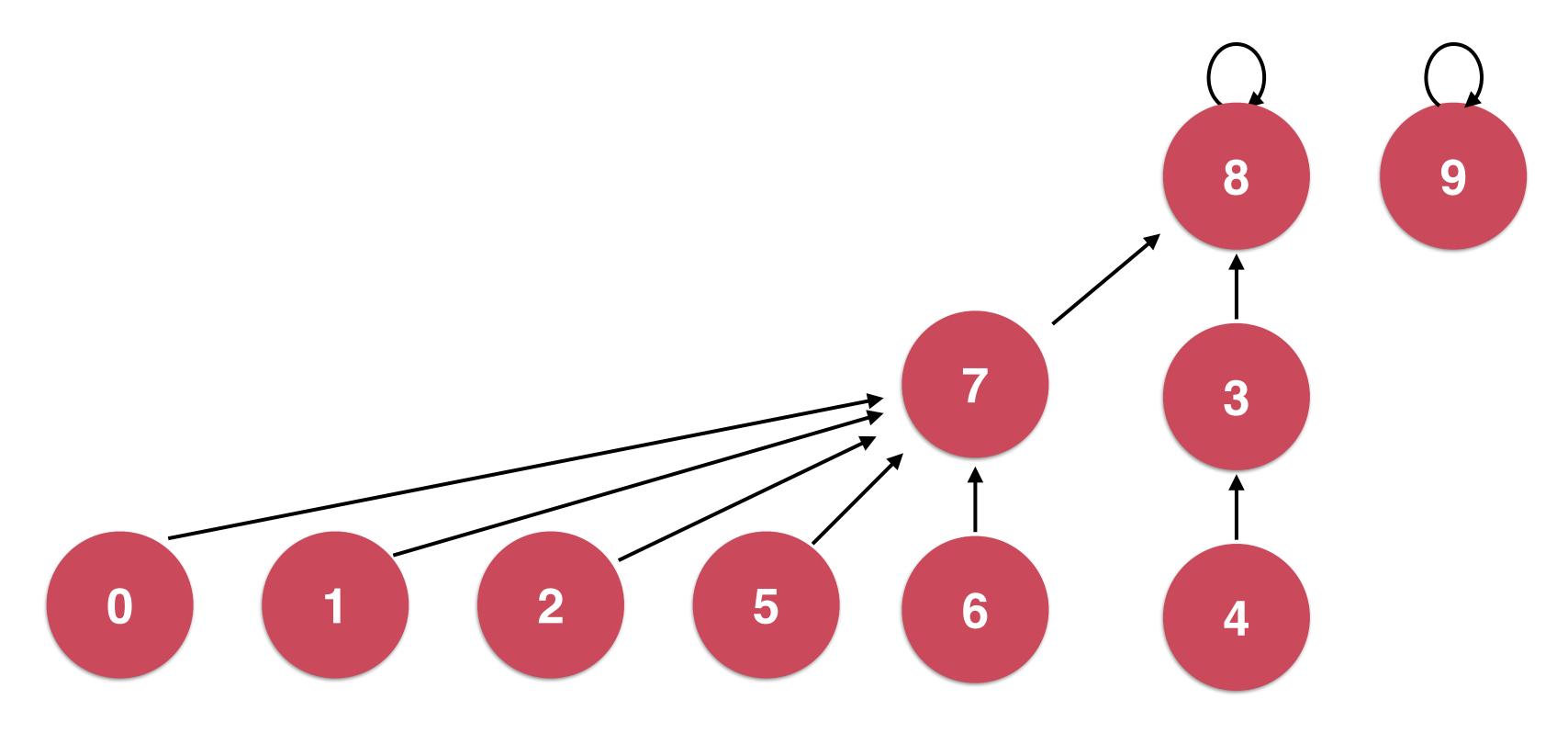


parent 7 7 7 8 3 7 7 7 9



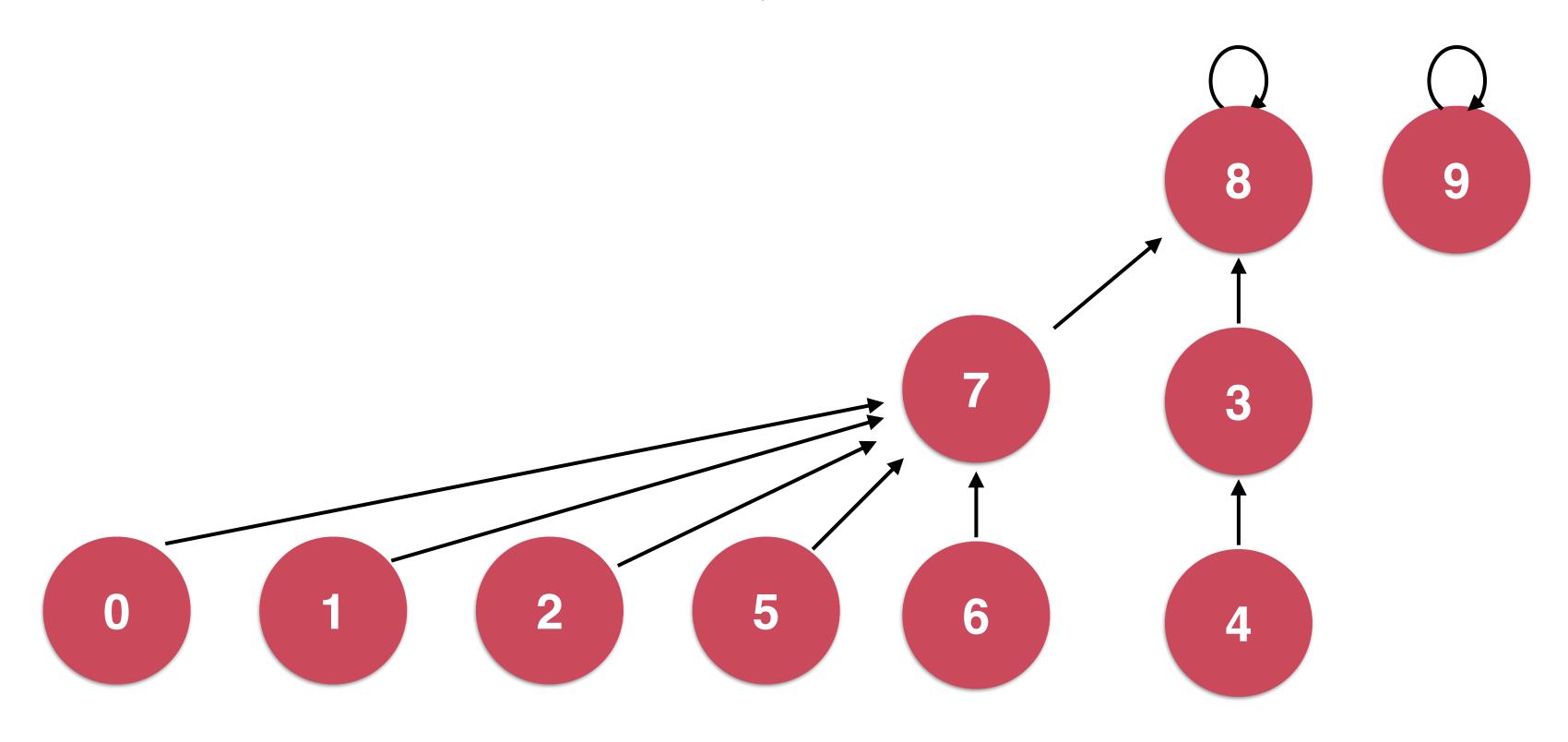
0 1 2 3 4 5 6 7 8 9

parent 7 7 7 8 3 7 7 8 9



0 1 2 3 4 5 6 7 8 9

parent 7 7 7 8 3 7 7 8 9



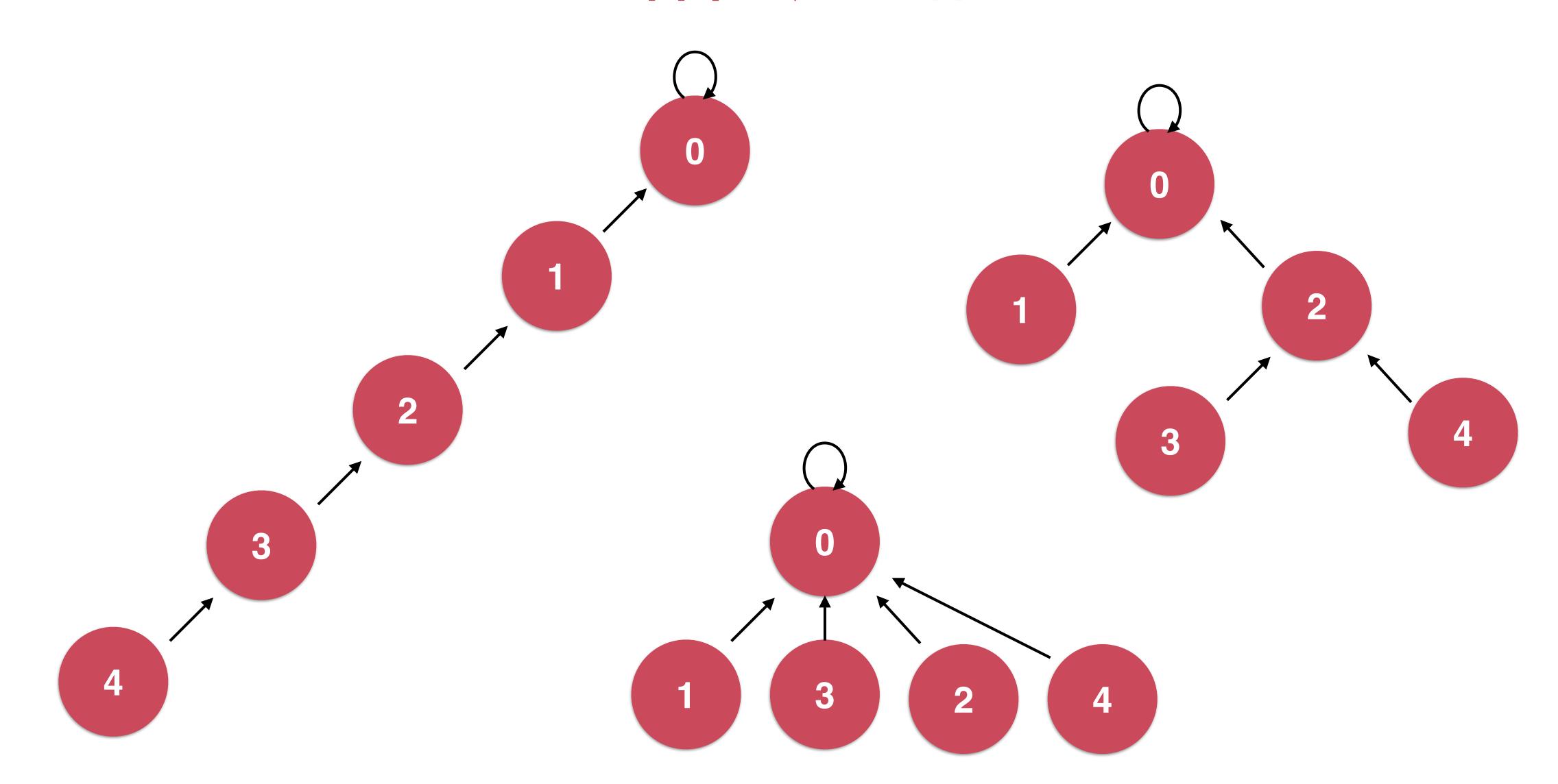
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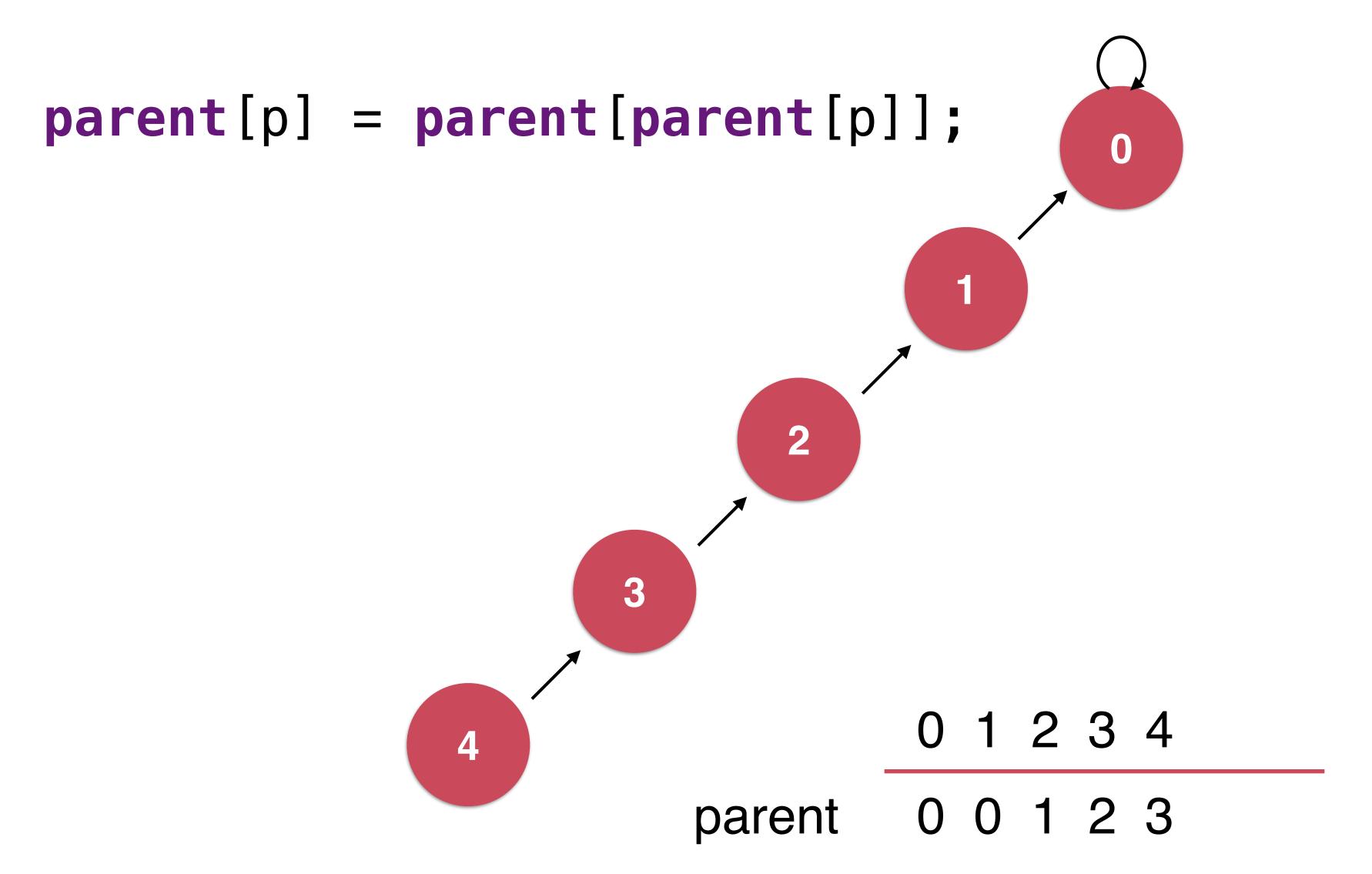
parent 7 7 7 8 3 7 7 8 8 9

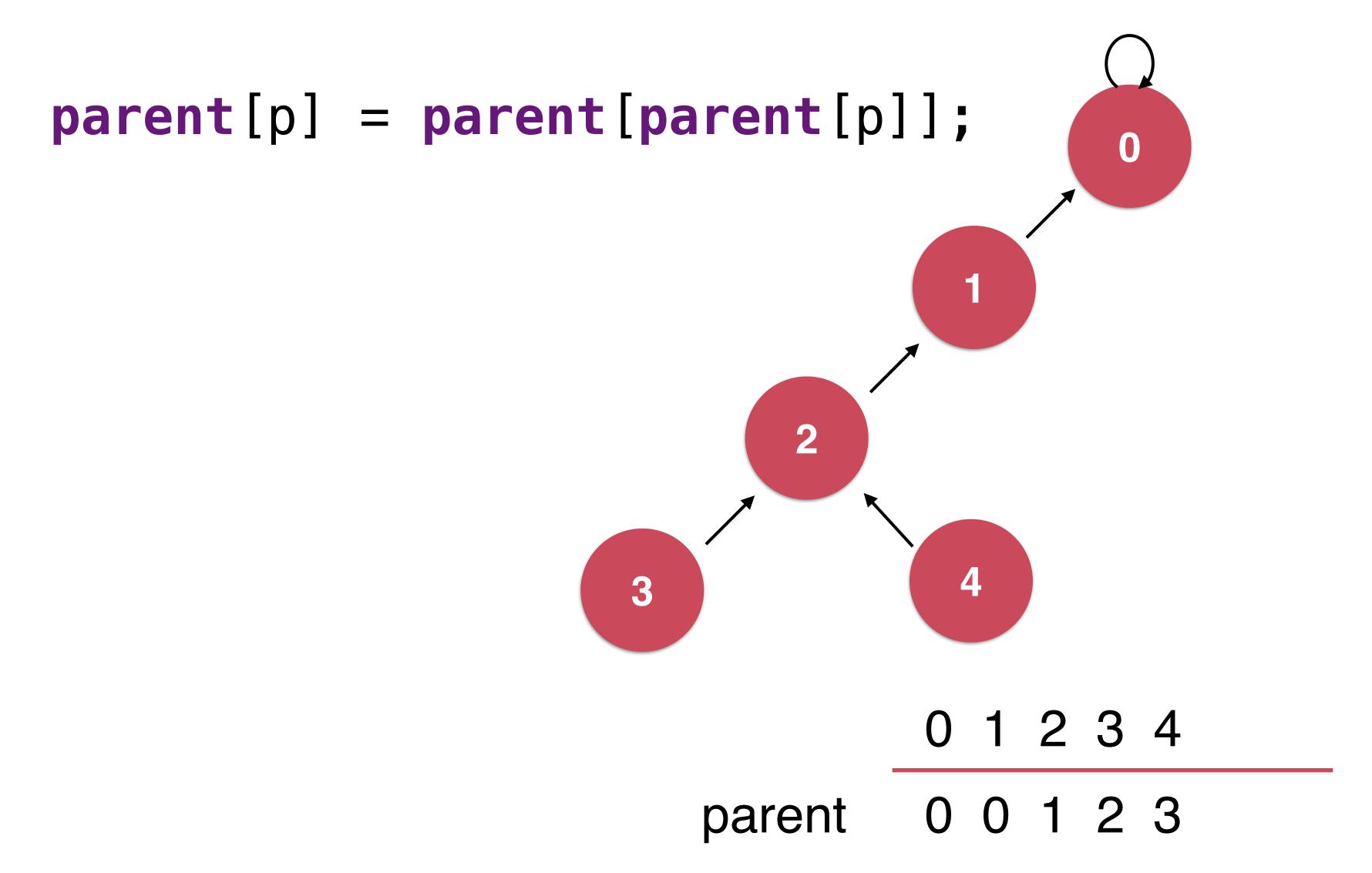
# 基于rank的优化 rank[i] 表示根节点为i的树的高度

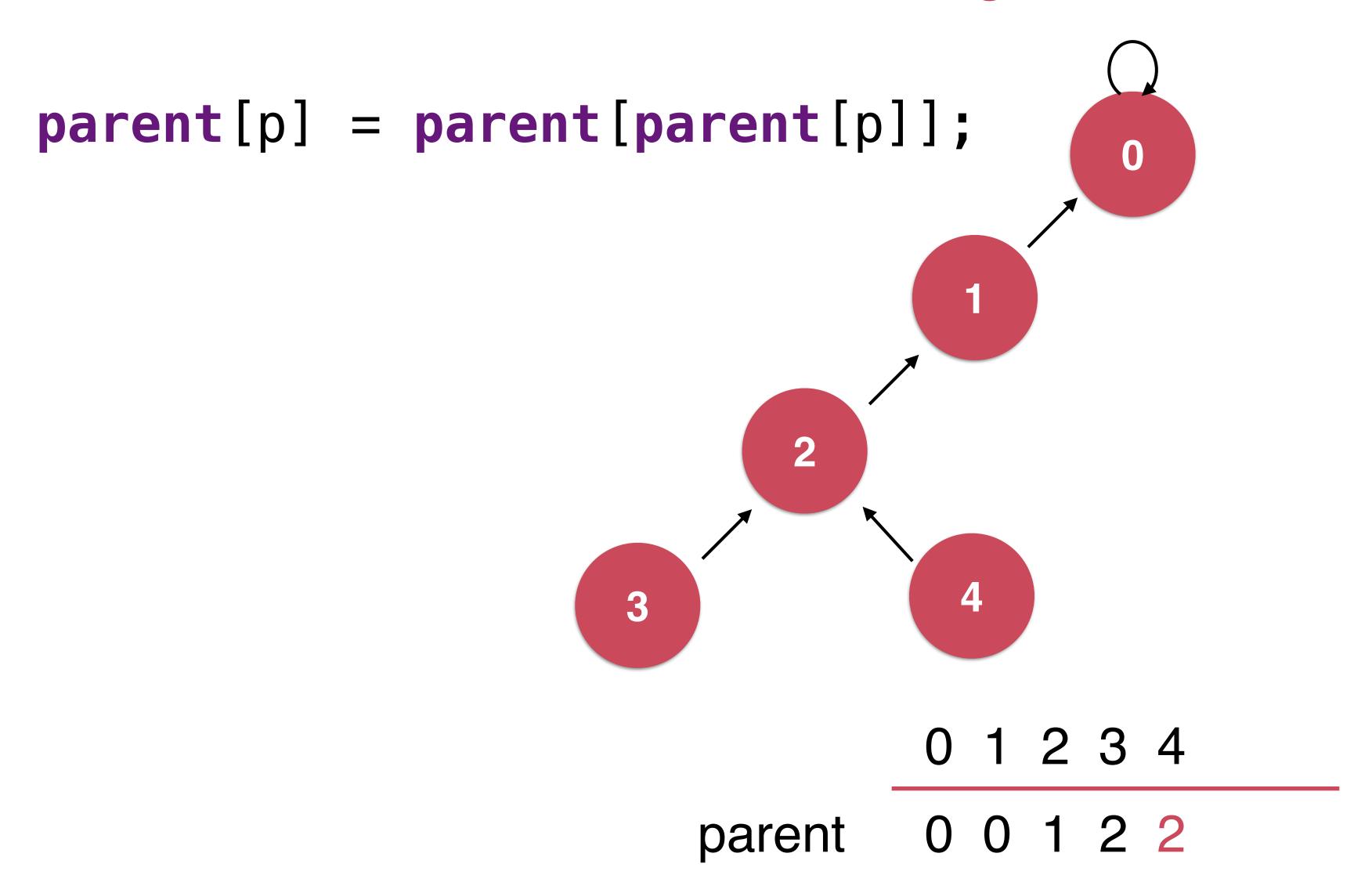
# 实践:基于rank的优化

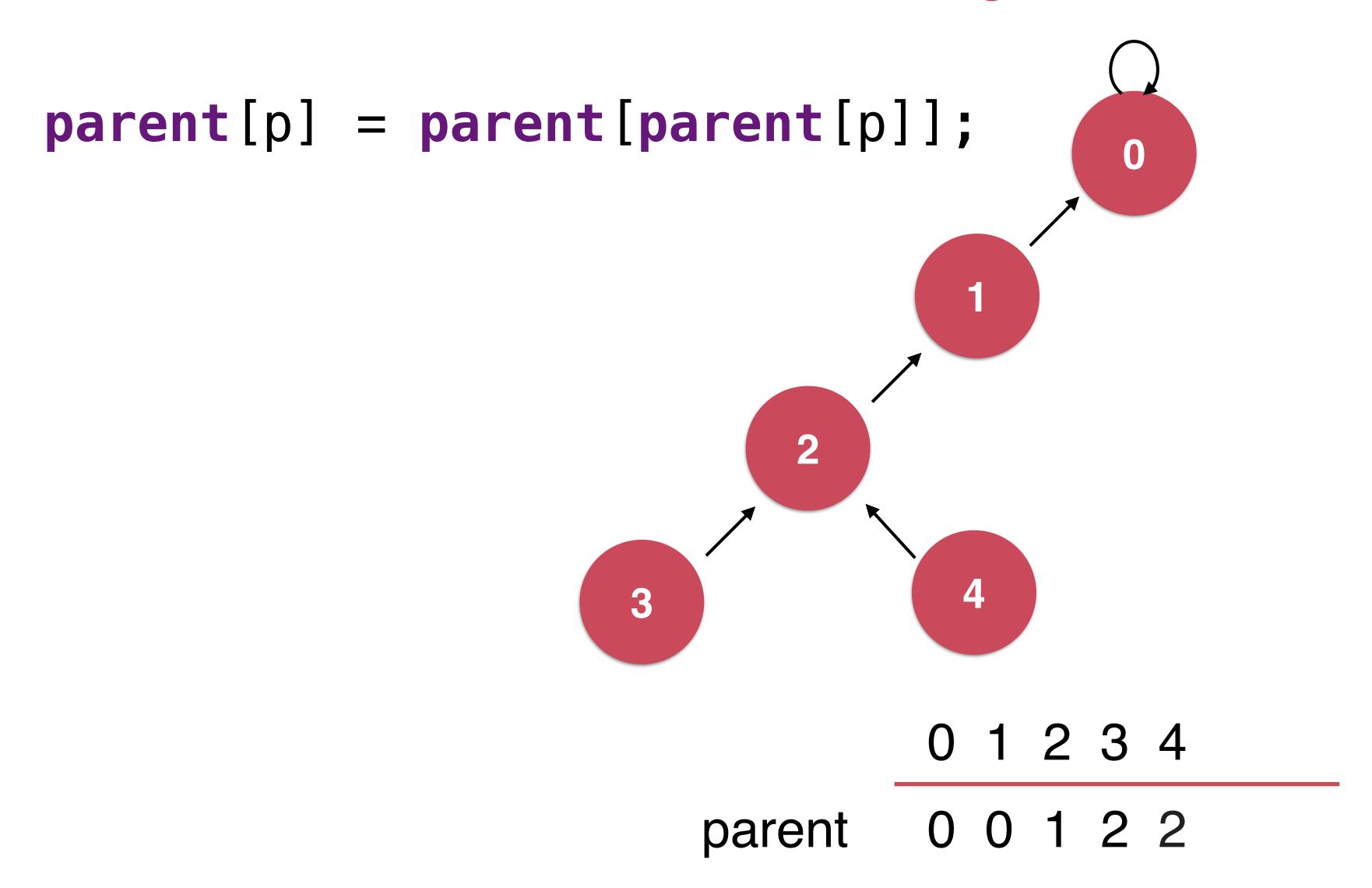
### 路径压缩











parent[p] = parent[parent[p]];

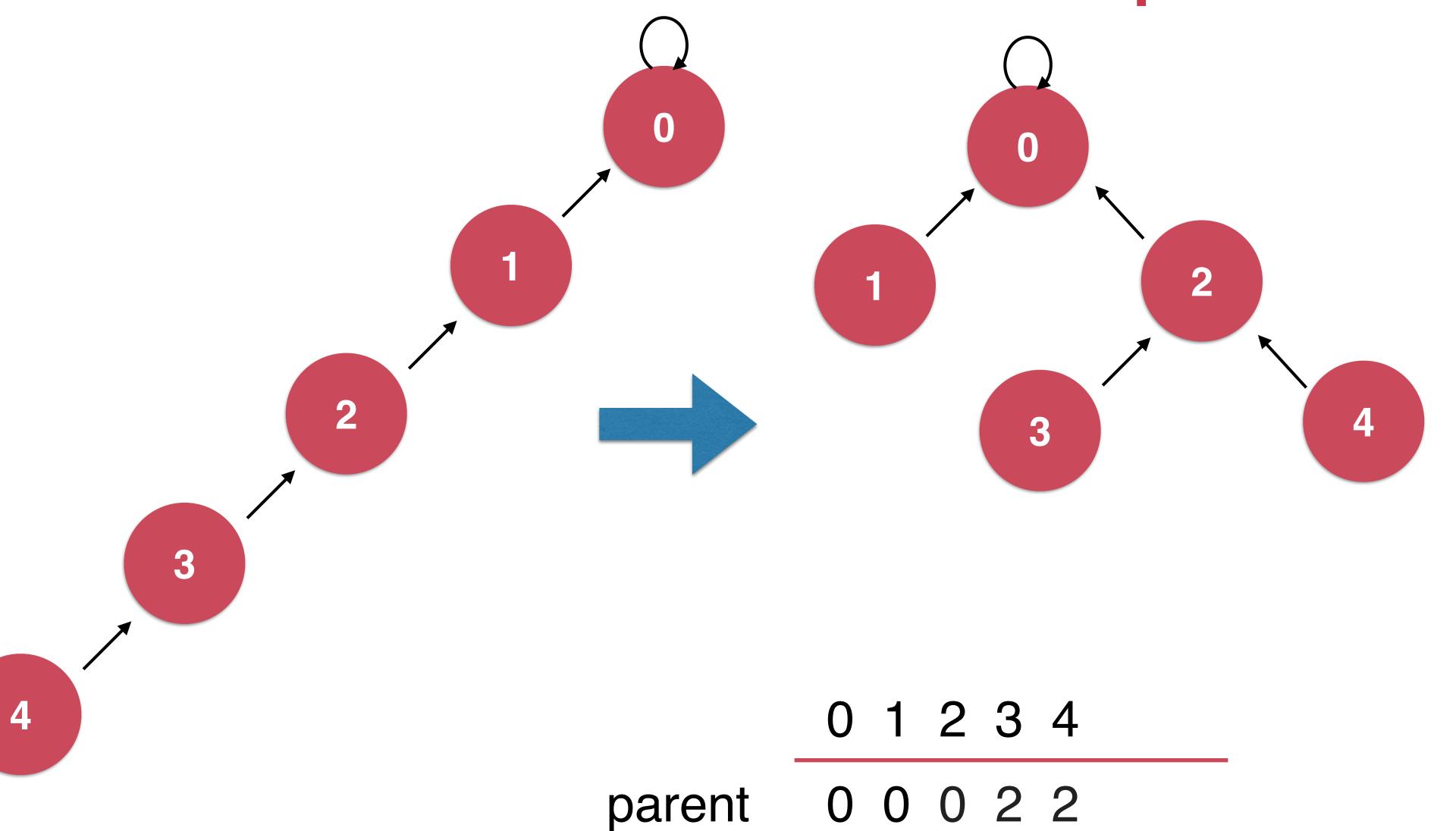
1
2
4

parent[p] = parent[parent[p]];

1
2
4

parent[p] = parent[parent[p]];

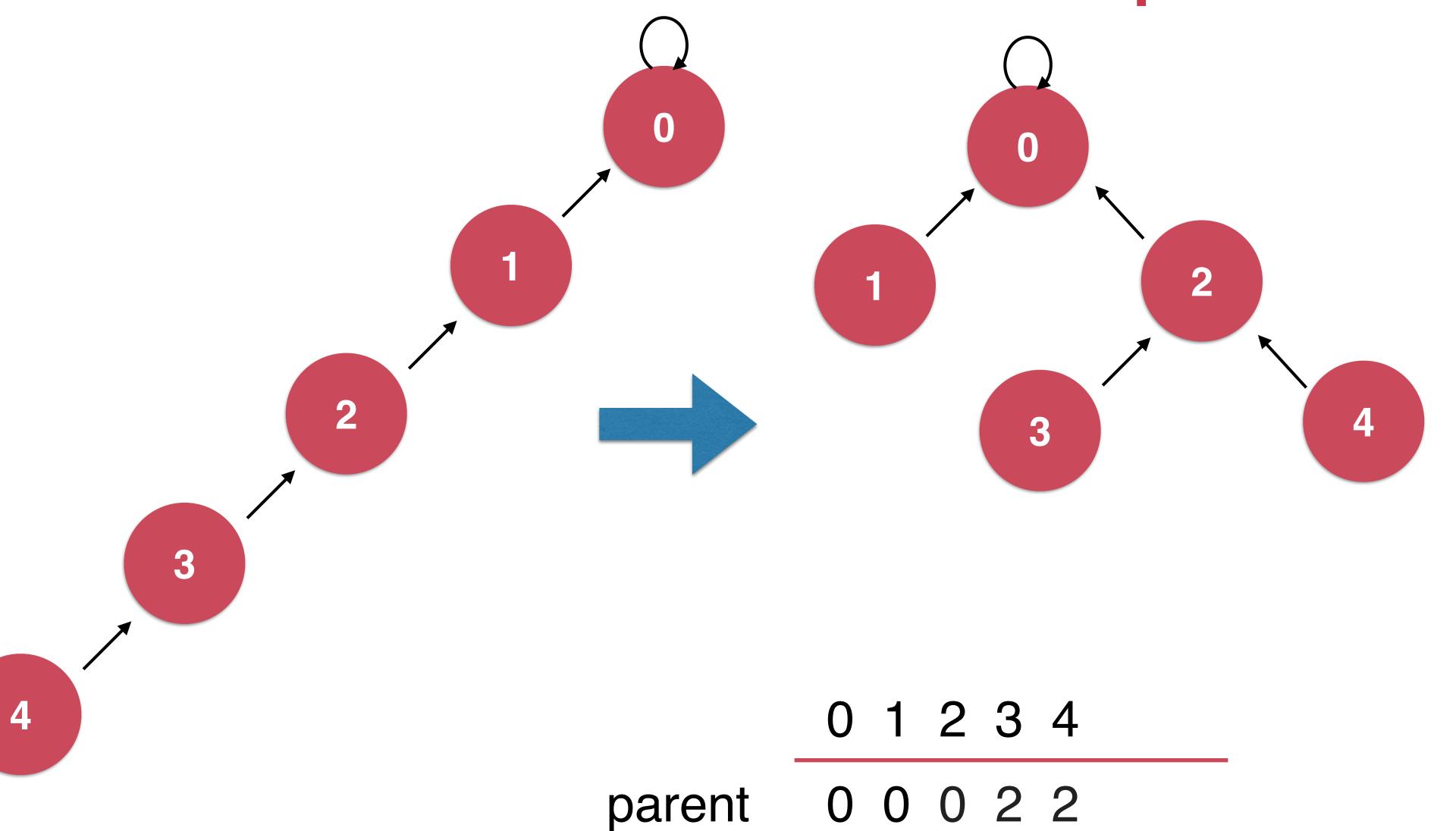
1
2
4

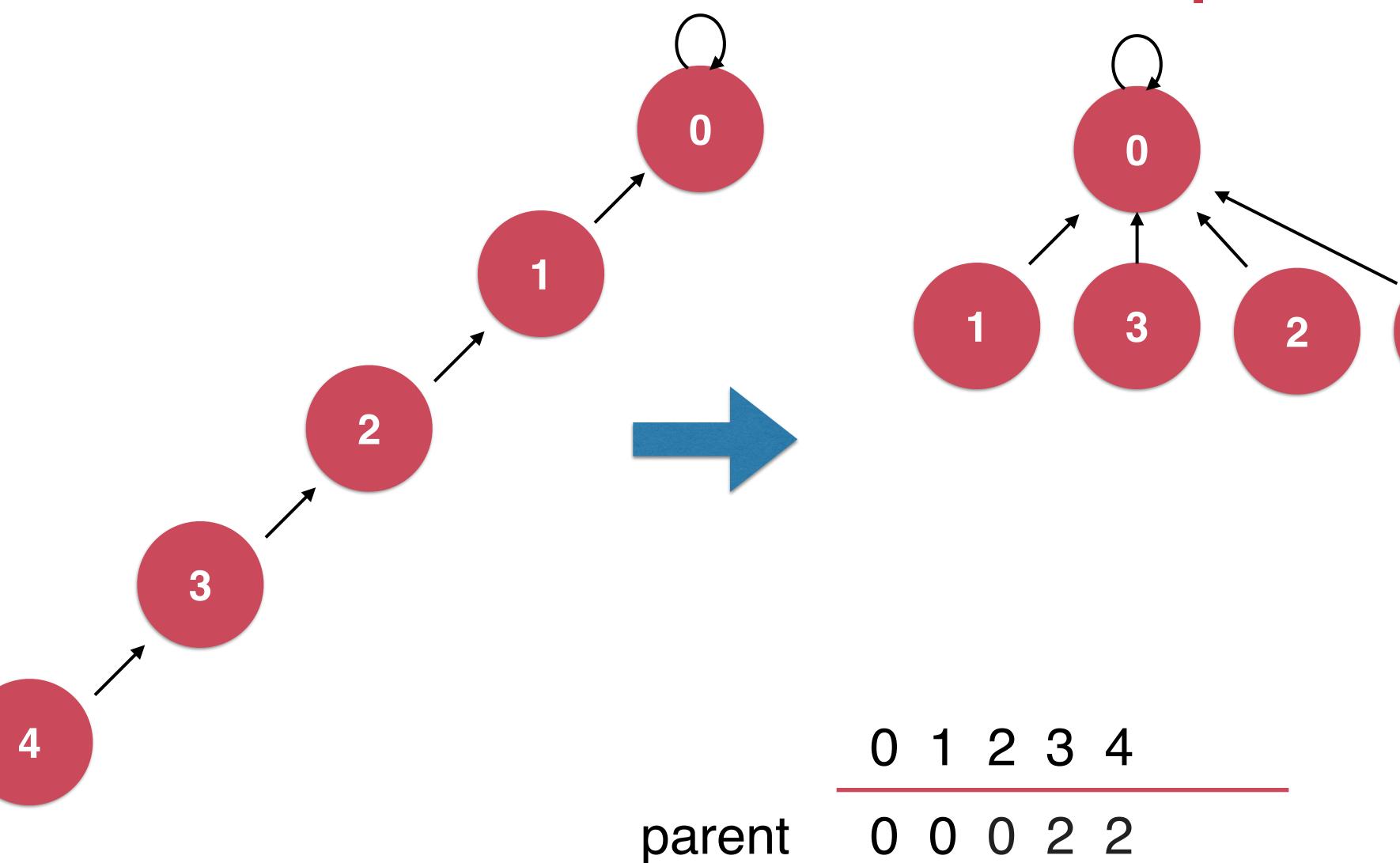


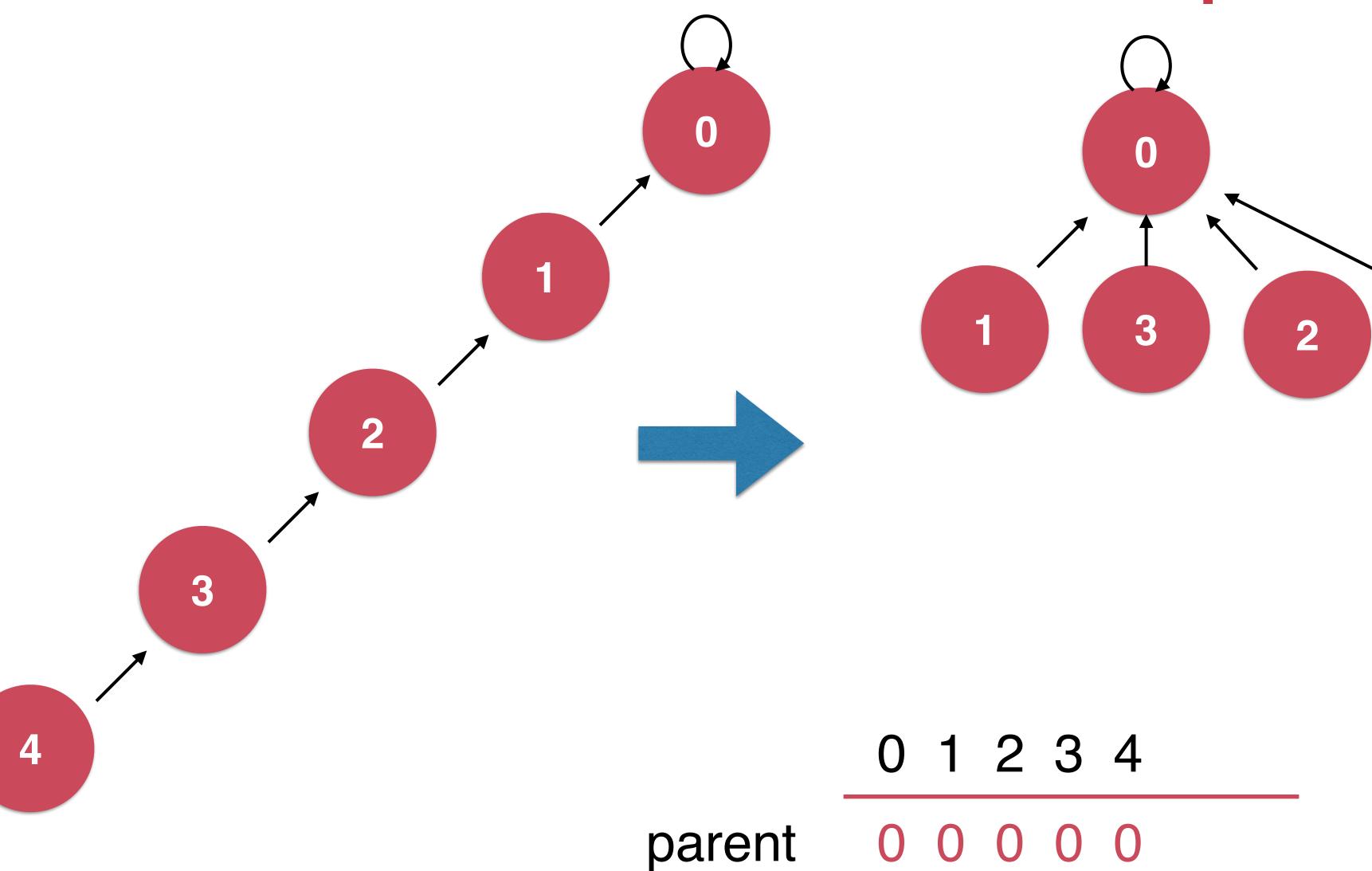
# 实践:路径压缩

### 实践:比较路径压缩和非路径压缩的效率

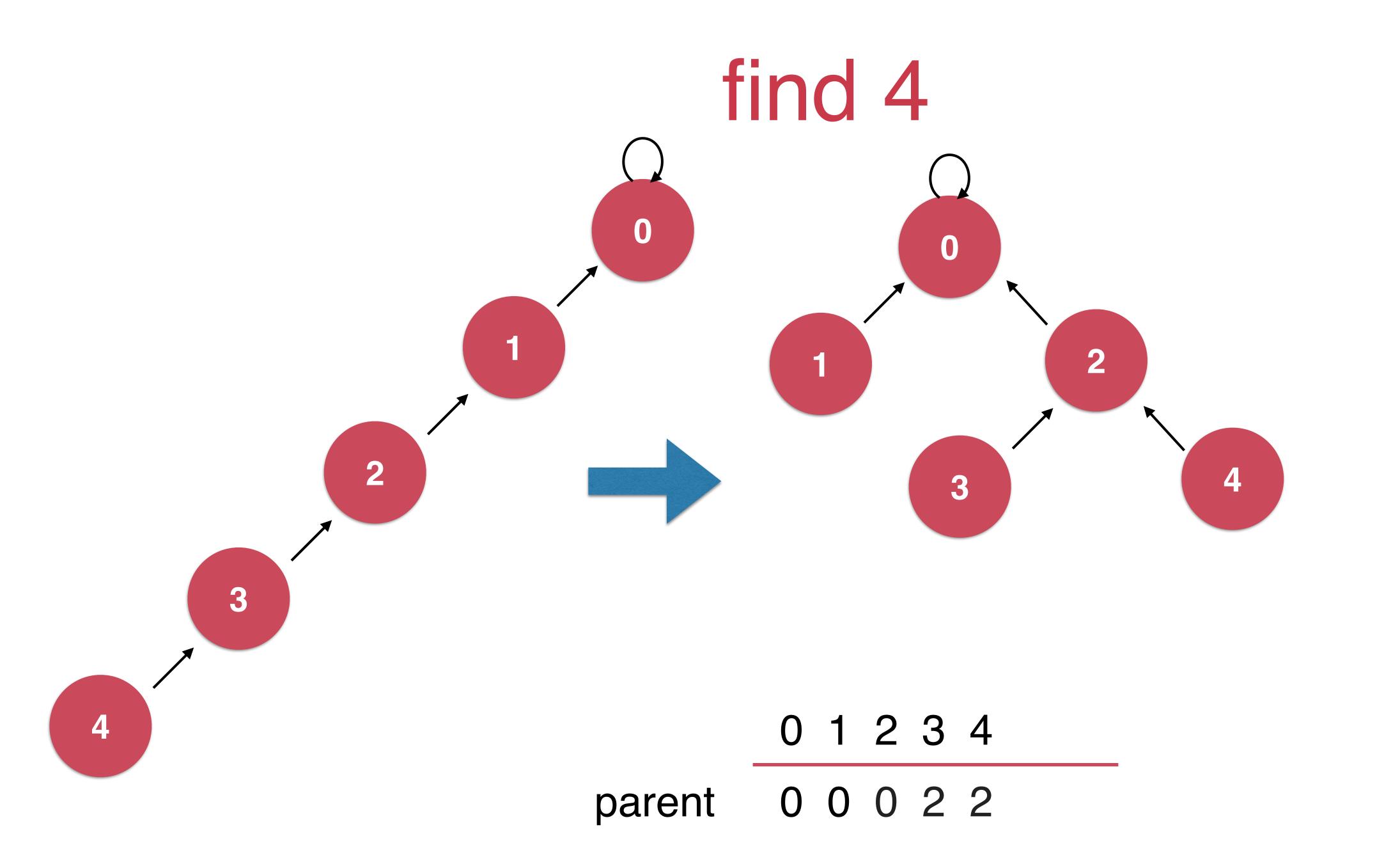
# 更多和并查集相关的话题

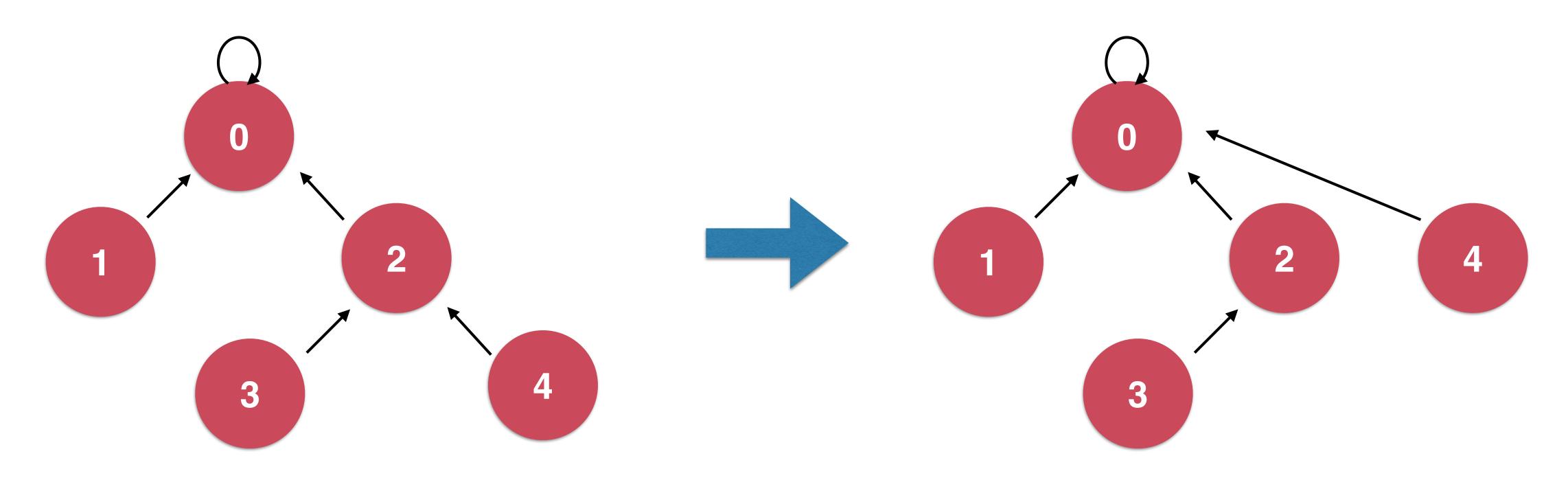


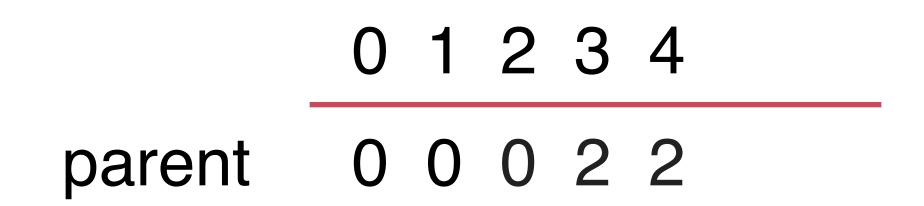


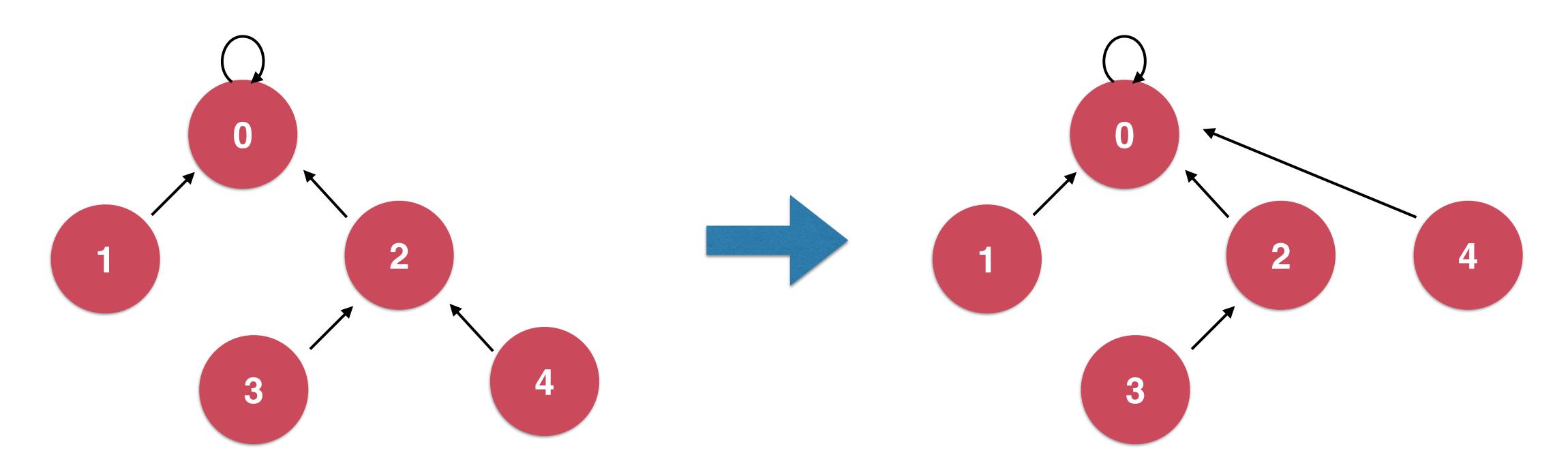


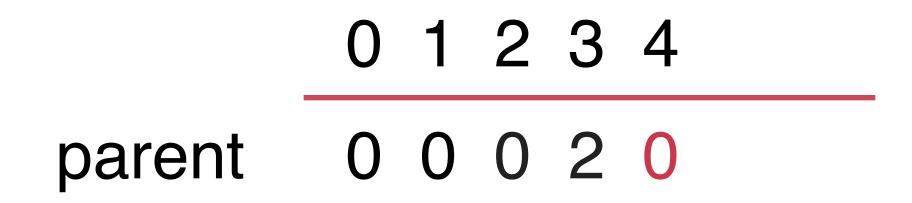
# 实践:递归的路径压缩

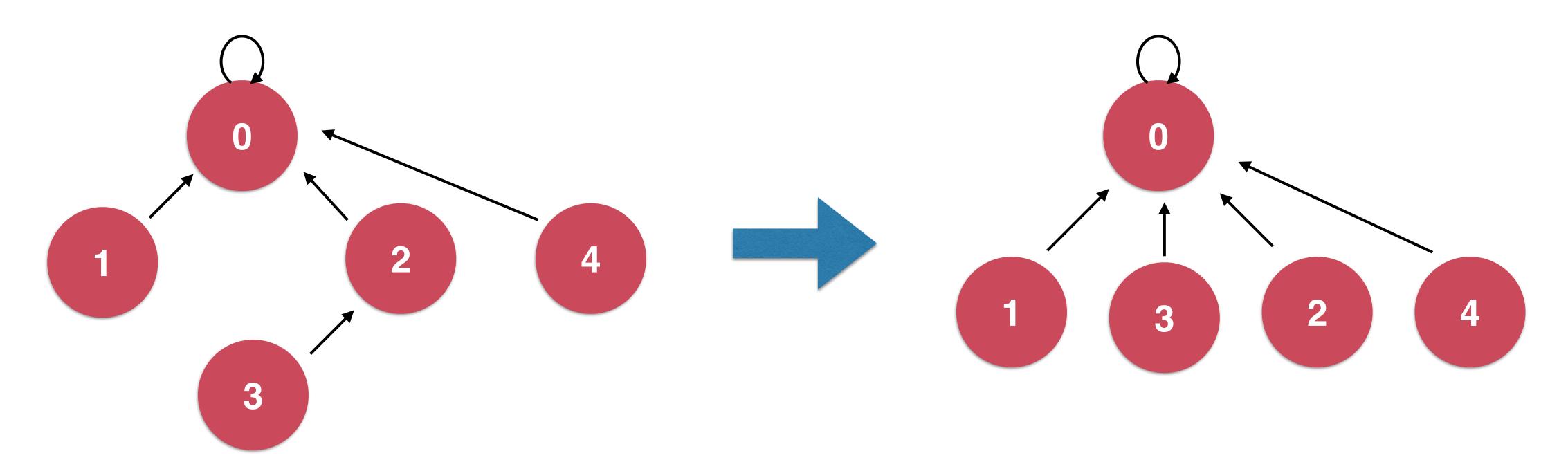


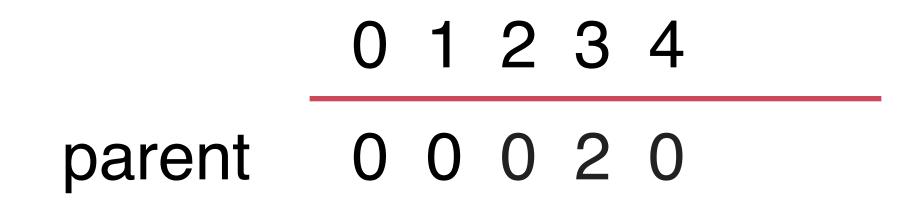


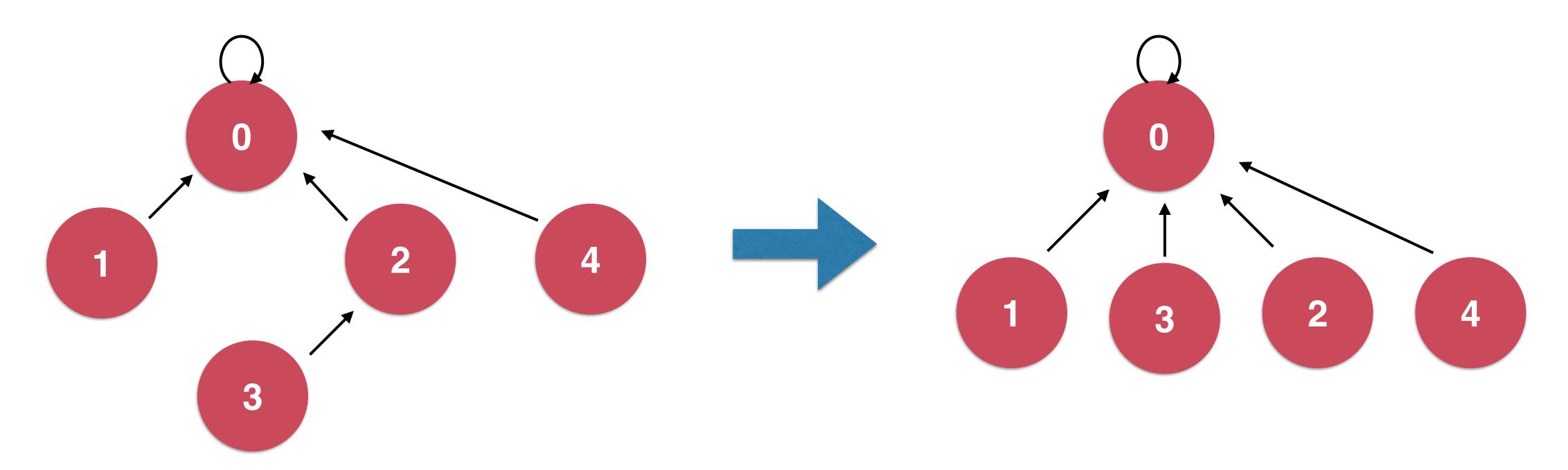


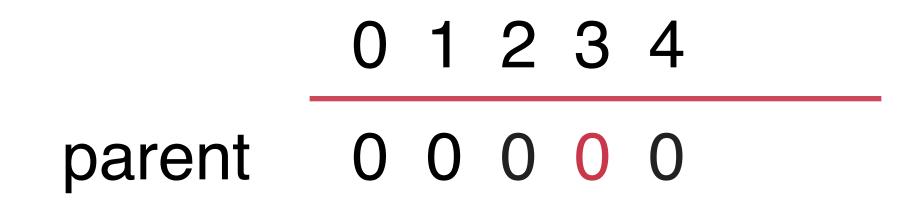












### 并查集的时间复杂度分析

**O**(h)

严格意义上: O(log\*n) iterated logarithm

### 并查集的时间复杂度分析

严格意义上: O(log\*n) iterated logarithm

$$\log^* n = \begin{cases} 0 & if (n \le 1) \\ 1 + \log^* (\log n) & if (n > 1) \end{cases}$$

近乎是O(1)级别的

### 更多Leetcode上Union Find相关的问题

### 并查集 Union Find

### 其他

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# 算法与数据结构体系课程

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