

# Algorithmics Problem Sheet 3

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This problem sheet is due on Monday 29th April. The tutorial will be 14:30-16:00 in LAM-2090. Please have a go at as many problems as you can. As well as going through the problems, the session on Monday can be used to answer any other questions you have about the course.

## Problem 1

For the set  $\{1, 4, 5, 10, 16, 17, 21\}$  of keys, draw BSTs of heights 2, 3, 4, 5, and 6.

## Problem 2

Consider the following algorithm which finds the node with minimum key in a BST.

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**Require:** BST  $T$

```
1: initialise  $x \leftarrow \text{root}(T)$ 
2: while node  $\text{left}(x)$  exists do
3:    $x \leftarrow \text{left}(x)$ 
4: return  $x$ 
```

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Write a recursive version of the above algorithm. Recall that a recursive function is defined in terms of itself. That is, somewhere in the pseudocode, the algorithm should run itself.

## Problem 3

Consider using the following algorithm to find the number 363 in a BST.

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**Require:** BST  $T$ , key  $k$

```
1: initialise  $x \leftarrow \text{root}(T)$ 
2: while True do
3:   if  $T[x] = k$  then
4:     return  $x$ 
5:   if  $T[x] > k$  then
6:      $x \leftarrow \text{left}(x)$ 
7:   if  $T[x] > k$  then
8:      $x \leftarrow \text{right}(x)$ 
```

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Which of the following sequences cannot be the sequence of nodes examined?

(a) 2, 252, 401, 398, 330, 344, 397, 363.

- (b) 924, 220, 911, 244, 898, 258, 362, 363.
- (c) 925, 202, 911, 240, 912, 245, 363.
- (d) 2, 399, 387, 219, 266, 382, 381, 278, 363.
- (e) 935, 278, 347, 621, 299, 392, 358, 363.

#### Problem 4

The following is the algorithm for inserting a new node into a BST:

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**Require:** BST  $T$ , new key value  $k$

```

1: initialise  $x \leftarrow 0$ ,  $p \leftarrow 0$ 
2: while node  $x$  exists do
3:    $p \leftarrow x$ 
4:   if  $k < T[x]$  then
5:      $x \leftarrow \text{left}(x)$ 
6:   if  $k > T[x]$  then
7:      $x \leftarrow \text{right}(x)$ 
8: if  $k < T[p]$  then
9:   create left child of  $p$  with key  $k$ 
10: if  $k > T[p]$  then
11:   create right child of  $p$  with key  $k$ 

```

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Give a recursive version of the above algorithm.

#### Problem 5

Give a non-recursive algorithm that ‘traverses’ a BST in order. That is, it takes a BST as an input and outputs each key in order.

#### Problem 6

Given a BST  $T$  and a node  $x$  in  $T$ , we write  $\text{succ}(x)$  to denote the node with the next-highest key value in  $T$ . Show that if a node in a BST has two children, then  $\text{succ}(x)$  has no left child.

#### Problem 7

It is possible to sort a set of numbers by first building a BST containing them and then traversing the BST in order. To build a BST, we can repeatedly run our node insertion algorithm on the unsorted set. What are the worst-case and best-case running times for this sorting algorithm?

#### Problem 8

Consider the following algorithm. It takes in an array  $A[1 : 2n]$  of even length  $2n$  and swaps the key  $A[i]$  with the key  $A[2n + 1 - i]$  for each  $i = 1, \dots, n$ .

State a loop invariant for this procedure. Use it to show that the algorithm reverses the order of the keys in the array  $A$ .

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**Require:** array  $A[1 : 2n]$

```
1: for  $i = 1$  upto  $n$  do  
2:   swap the keys  $A[i]$  and  $A[2n + 1 - i]$   
3: return  $A[1 : 2n]$ 
```

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HINT: You need to come up with a true/false statement that holds at each iteration of the loop and is dependent on  $i$ . Then you need to show that this is true at the first step. Then you need to show that, if it is true at the  $i$ th iteration, it's also true at the  $(i + 1)$ th iteration. Finally, you must show that when the algorithm ends, that this loop invariant tells you that the algorithm has swapped all of the numbers in the array correctly.