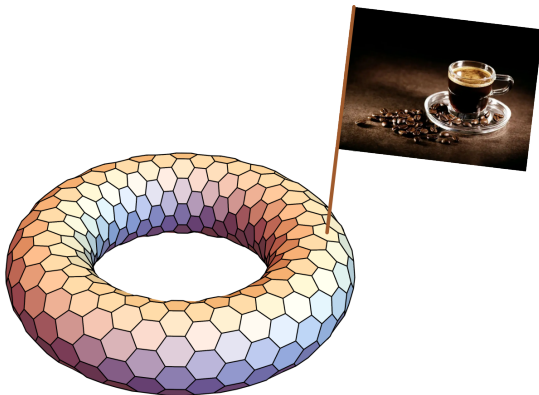


A Tutte-type canonical decomposition of 3- and 4-connected graphs

Jan Kurkofka (TU Freiberg)

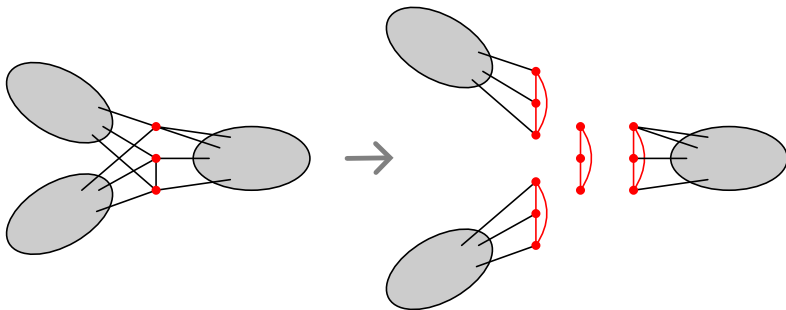


Joint work with Tim Planken
BWAG '25

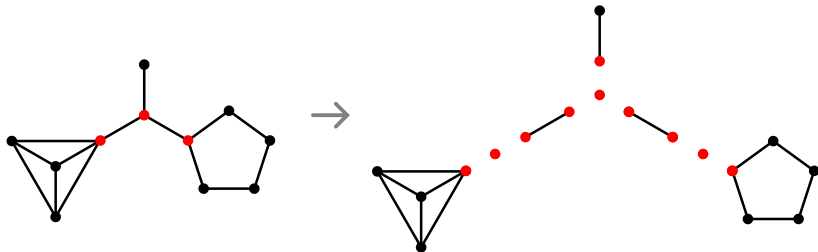
Problem: Canonically decompose k -con'd G along k -separators
into parts that are $(k + 1)$ -con'd or 'basic'.

canonical : \iff $\text{Aut}(G)$ -invariant

Decomposing G along a k -separator:

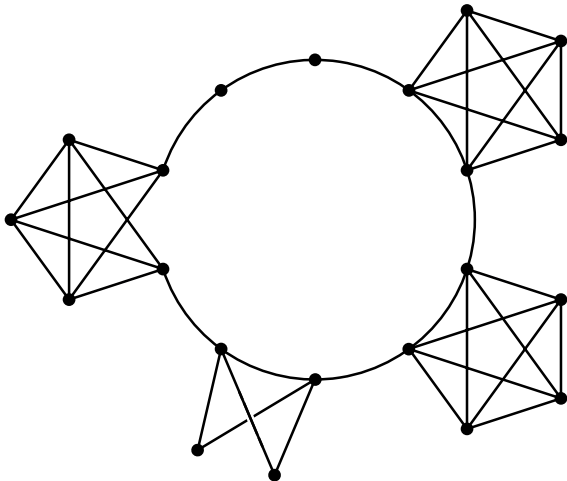


Problem: Canonically decompose k -con'd G along k -separators into parts that are $(k + 1)$ -con'd or 'basic'.

 $k = 1:$ 

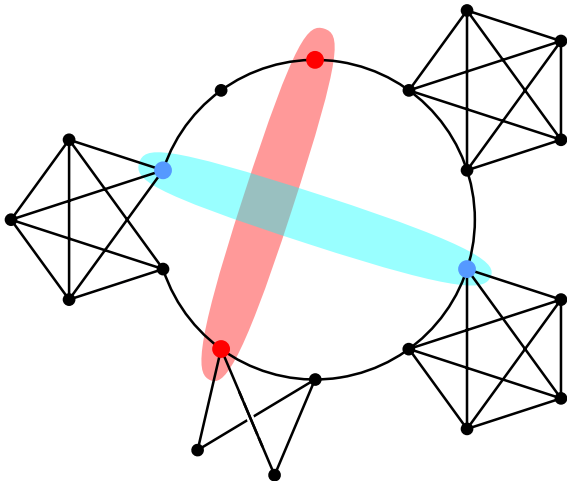
Problem: Canonically decompose k -con'd G along k -separators into parts that are $(k + 1)$ -con'd or 'basic'.

$k = 2$:

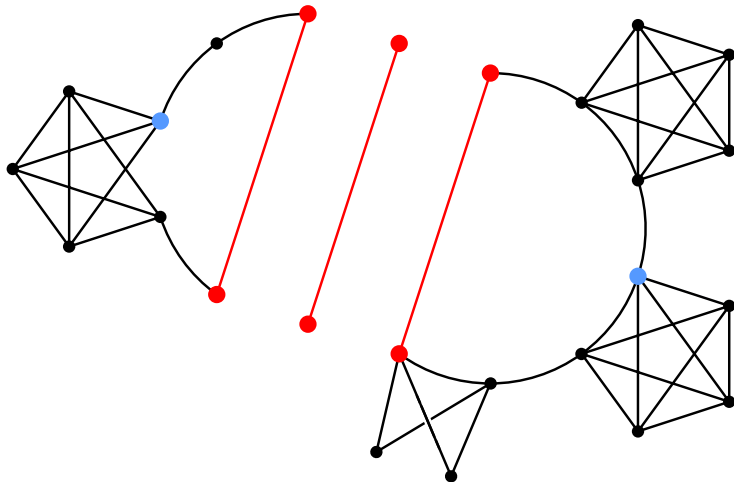


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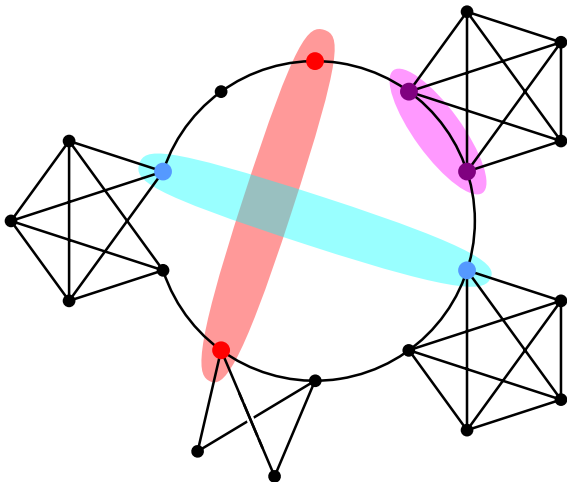
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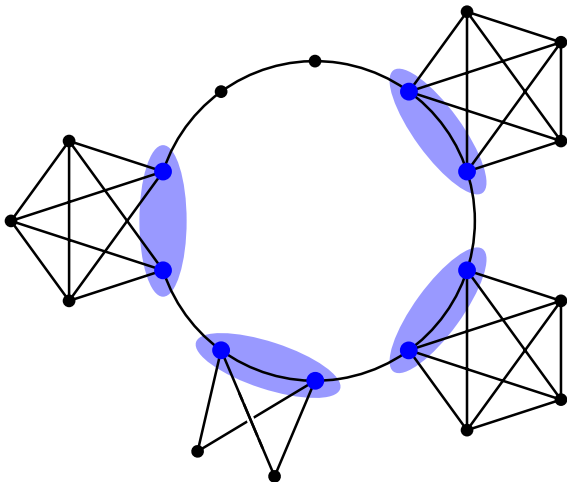
 $k = 2:$ 

Two k -separators *cross* if they separate each other;
otherwise they are *nested*.



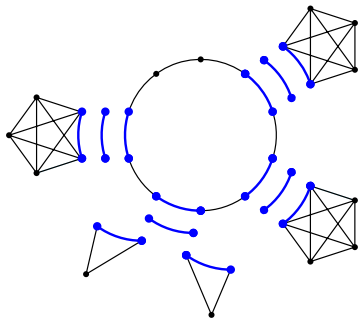
Two k -separators *cross* if they separate each other;
otherwise they are *nested*.

A k -separator is *totally-nested* if it is nested with every k -separator.



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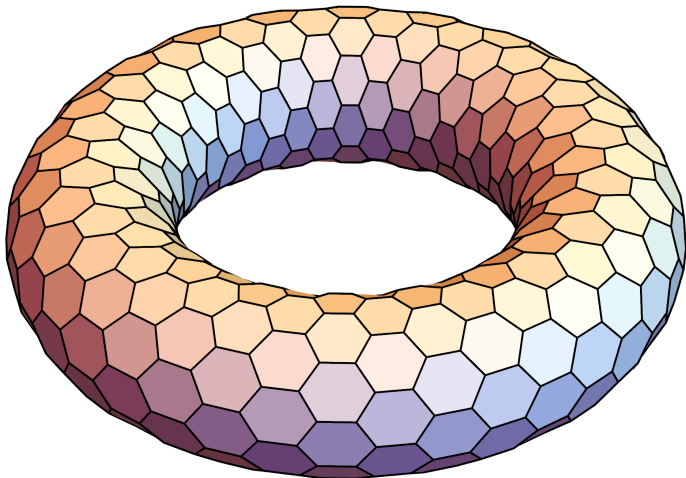
Theorem (Tutte 66)

Every 2-con'd G decomposes along its totally-nested 2-separators
into 3-con'd graphs, cycles and K_2 's.

Guess: Every k -con'd G decomposes along its totally-nested k -separators into $(k + 1)$ -con'd graphs and 'basic' graphs.

Guess: Every k -con'd G decomposes along its totally-nested k -separators into $(k + 1)$ -con'd graphs and 'basic' graphs.

Challenge 1 (Figure: $k = 3$)



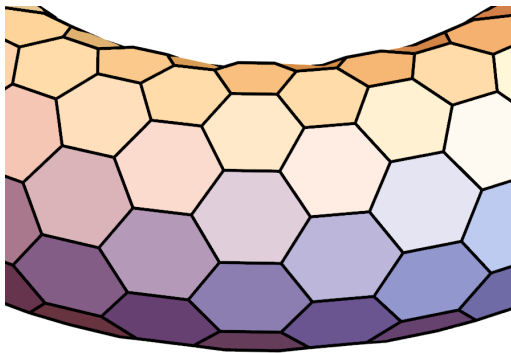
Guess: Every k -con'd G decomposes along its totally-nested k -separators into $(k + 1)$ -con'd graphs and 'basic' graphs.

Challenge 1 (Figure: $k = 3$)



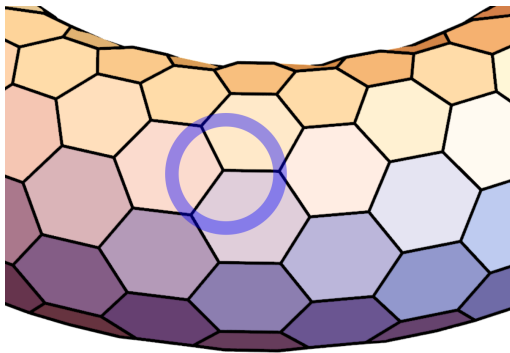
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Challenge 1 (Figure: $k = 3$)



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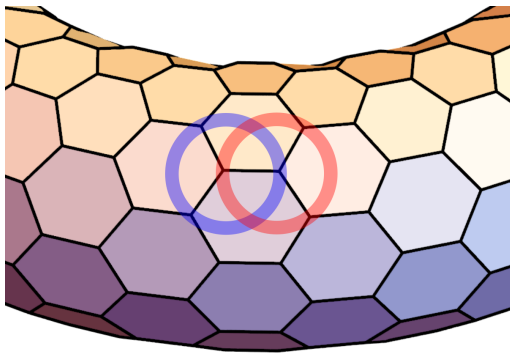
Challenge 1 (Figure: $k = 3$)



$$(\text{set of all } k\text{-separators}) = \{ N(v) : v \in V(G) \}$$

Guess: Every k -con'd G decomposes along its totally-nested k -separators into $(k + 1)$ -con'd graphs and 'basic' graphs.

Challenge 1 (Figure: $k = 3$)

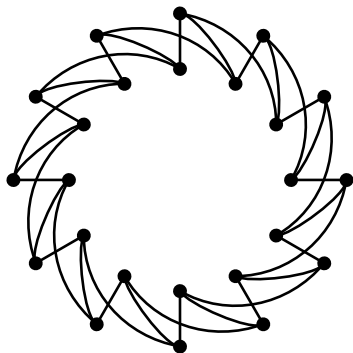


(set of all k -separators) = $\{ N(v) : v \in V(G) \}$

\implies for every edge uv : $N(u)$ crosses $N(v)$.

Guess: Every k -con'd G decomposes along its totally-nested k -separators into $(k + 1)$ -con'd graphs and 'basic' graphs.

Challenge 1 (Figure: $k = 3$)



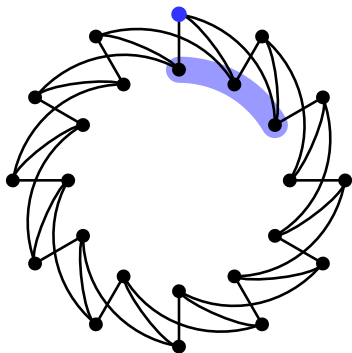
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Generalises to $k \geq 3$!

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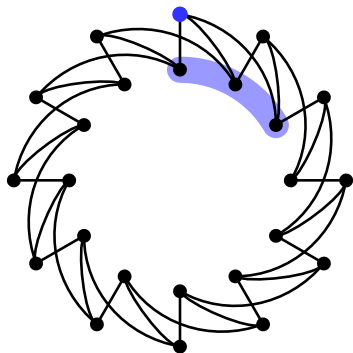
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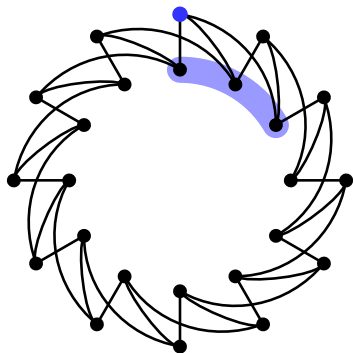
\implies for every edge uv : $N(u)$ crosses $N(v)$.

Generalises to $k \geq 3$!

Guess: Every k -con'd G decomposes along its totally-nested k -separators into **quasi- $(k + 1)$ -con'd** graphs and 'basic' graphs.

\iff k -con'd and every k -sep'r cuts off only one vertex

Challenge 1 (Figure: $k = 3$)



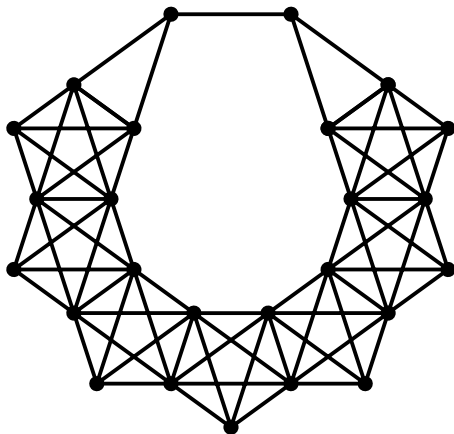
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Generalises to $k \geq 3$!

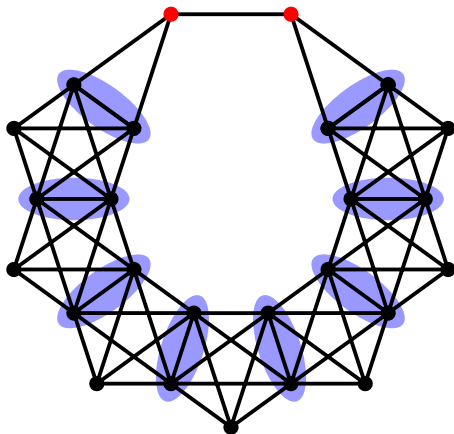
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Challenge 2 (Figure: $k = 3$)



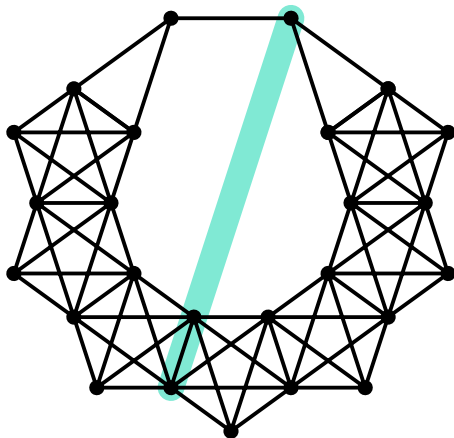
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Challenge 2 (Figure: $k = 3$)



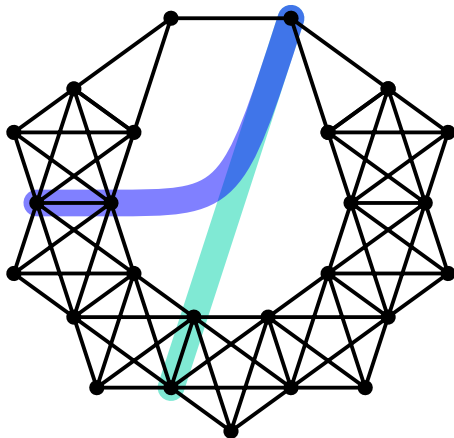
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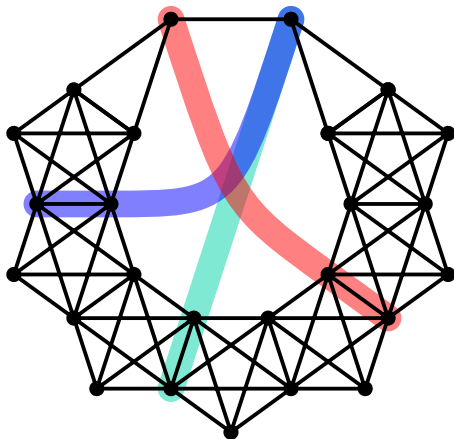
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Challenge 2 (Figure: $k = 3$)



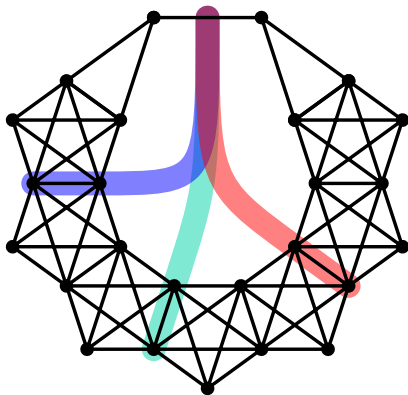
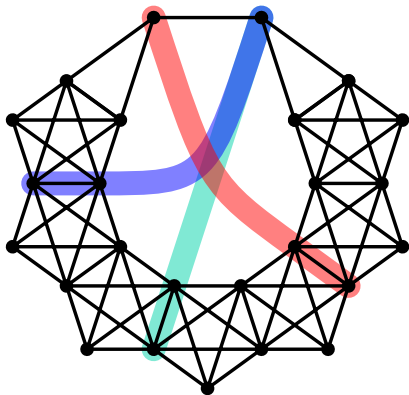
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Challenge 2 (Figure: $k = 3$)



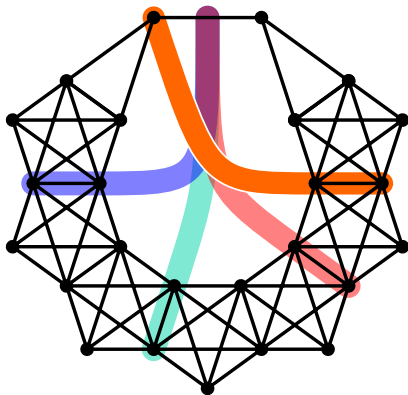
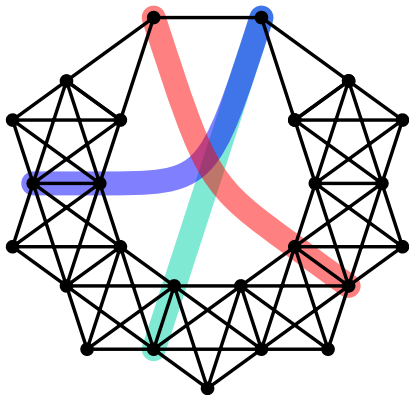
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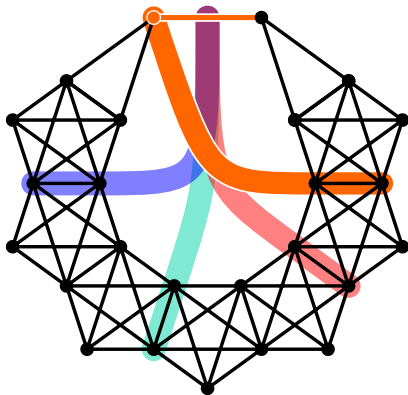
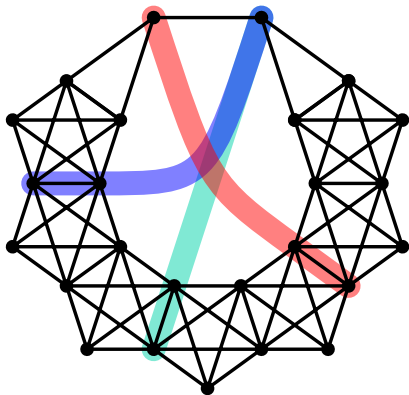
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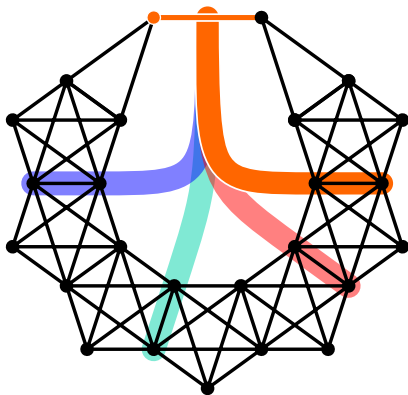
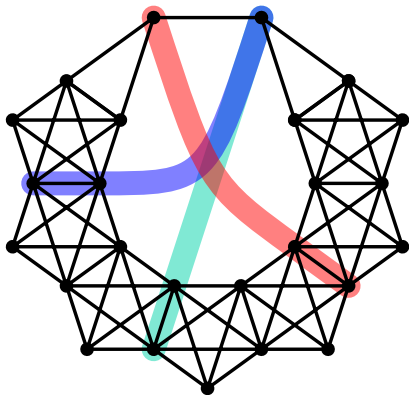
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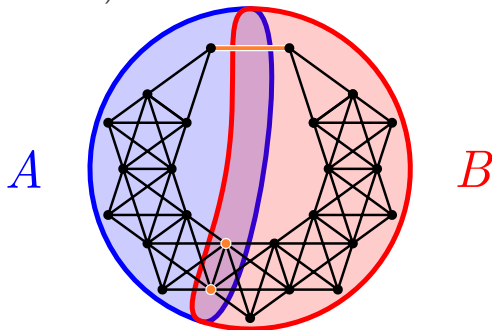
Guess: Every k -con'd G decomposes along its totally-nested k -separators into quasi- $(k + 1)$ -con'd graphs and 'basic' graphs.

Challenge 2 (Figure: $k = 3$)



Guess: Every k -con'd G decomposes along its totally-nested k -separators into quasi- $(k + 1)$ -con'd graphs and 'basic' graphs.

Challenge 2 (Figure: $k = 3$)



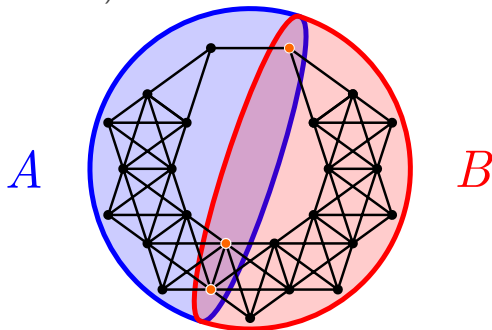
mixed-separation of G : (A, B) with $A \cup B = V(G)$ and $A, B \neq V(G)$

separator of (A, B) : $(A \cap B) \cup E(A \setminus B, B \setminus A)$

A **tri-separation** of G is a mixed-sep'n (A, B) with $|\text{sep'r}| = 3$ s.t.
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Challenge 2 (Figure: $k = 3$)



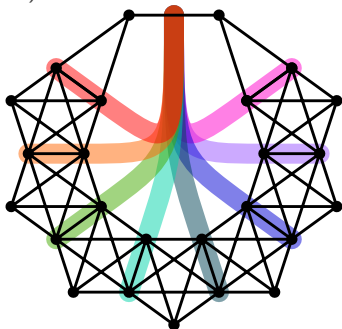
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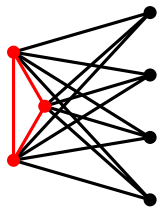
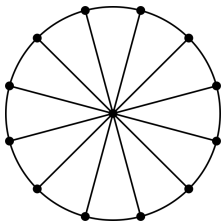
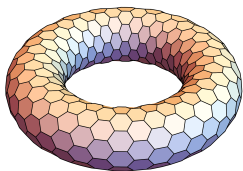
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Theorem (Carmesin & K. 23)

Every 3-con'd G decomposes along its totally-nested nontrivial tri-separations into parts that are

- quasi-4-con'd
- wheels
- thickened $K_{3,m}$



3

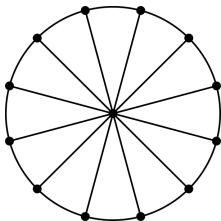
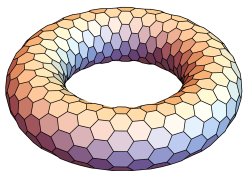
m

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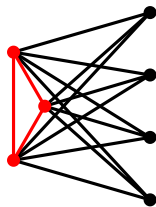
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4-connectivity?

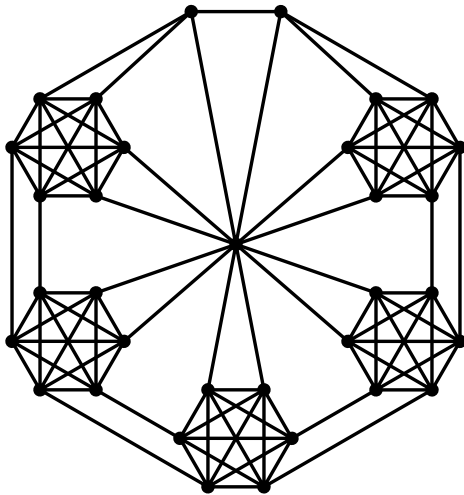


3

m

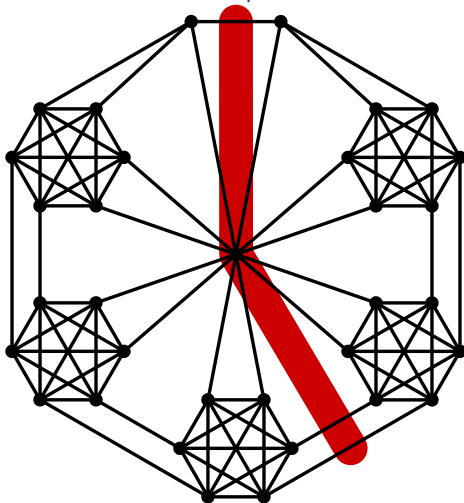
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Challenge 2: Verbatim extension of tri-separations to $k = 4$



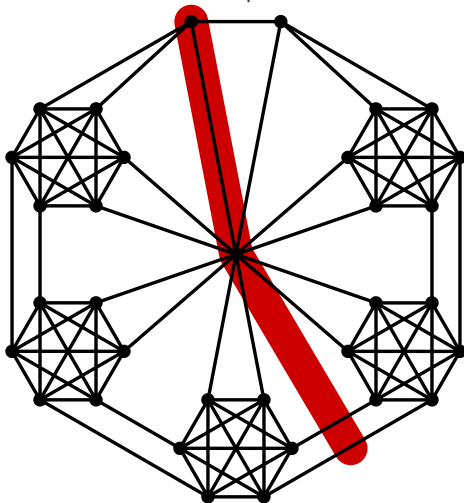
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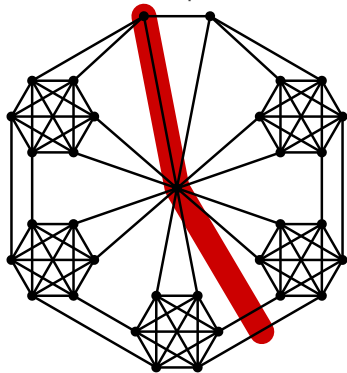
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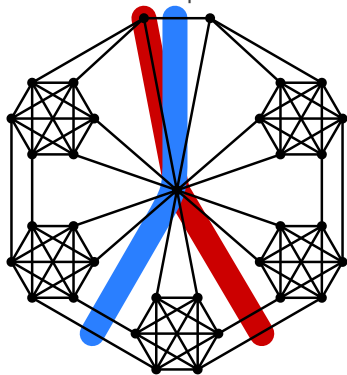
Challenge 2: Verbatim extension of tri-separations to $k = 4$



A tri-separation of G is a mixed-sep'n (A, B) with $|\text{sep}'r| = 3$ s.t.
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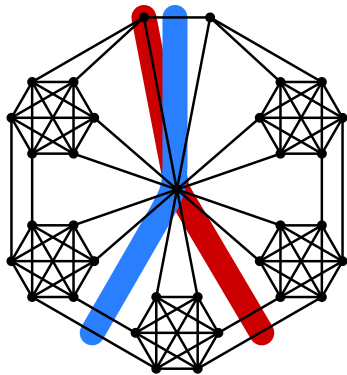
Challenge 2: Verbatim extension of tri-separations to $k = 4$ **fails!**



A tri-separation of G is a mixed-sep'n (A, B) with $|\text{sep}'r| = 3$ s.t.
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Challenge 2 for $k = 4$:

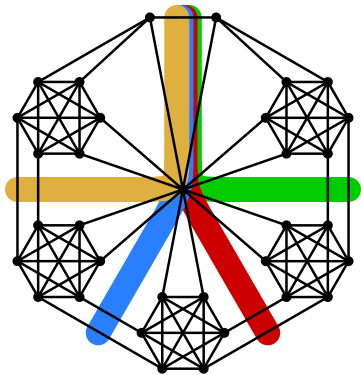


A **tetra-separation** of G is a mixed-sep'n (A, B) with $|\text{sep}'r| = 4$ s.t.:

- every v_x in $A \cap B$ has ≥ 2 neighb's in $A \setminus B$ and $B \setminus A$
- the edges in the sep'r form a matching

Guess: Every k -con'd G decomposes along its totally-nested k -separators into quasi- $(k + 1)$ -con'd graphs and 'basic' graphs.

Challenge 2 for $k = 4$:



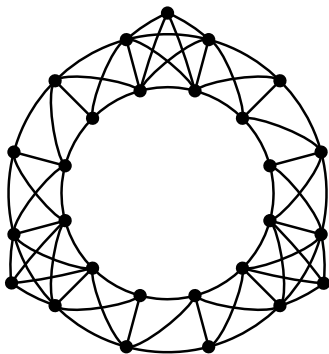
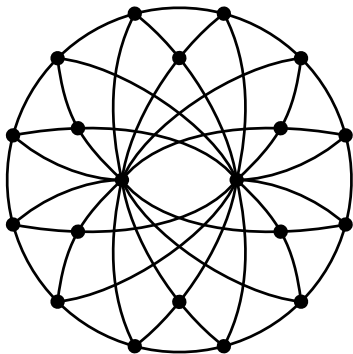
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Main result (K. & Planken 25)

Every 4-con'd G decomposes along its totally-nested tetra-separations into parts that are

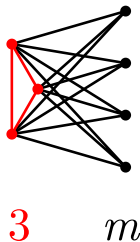
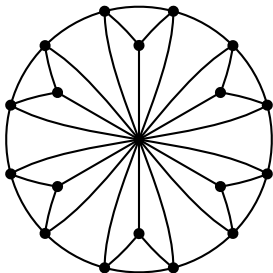
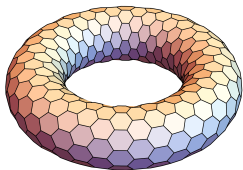
- quasi-5-con'd
- generalised double-wheels
- thickened $K_{4,m}$
- cycles of triangles and 3-con'd graphs on ≤ 5 vxs.



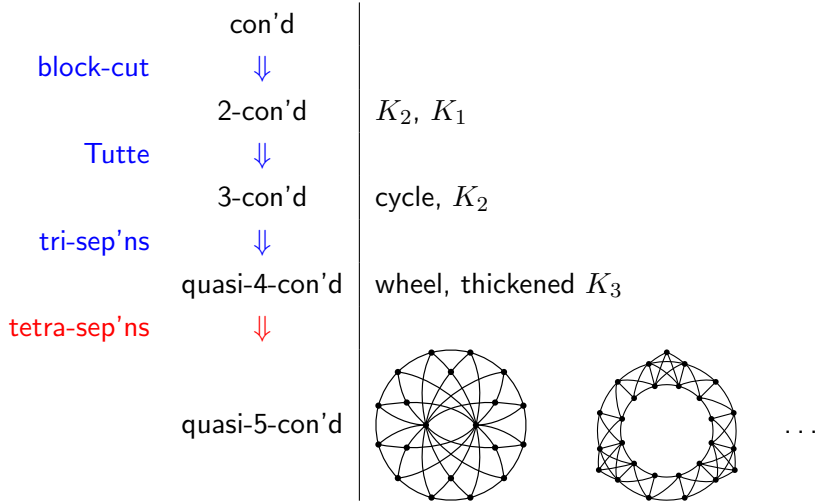
Corollary (K. & Planken 25)

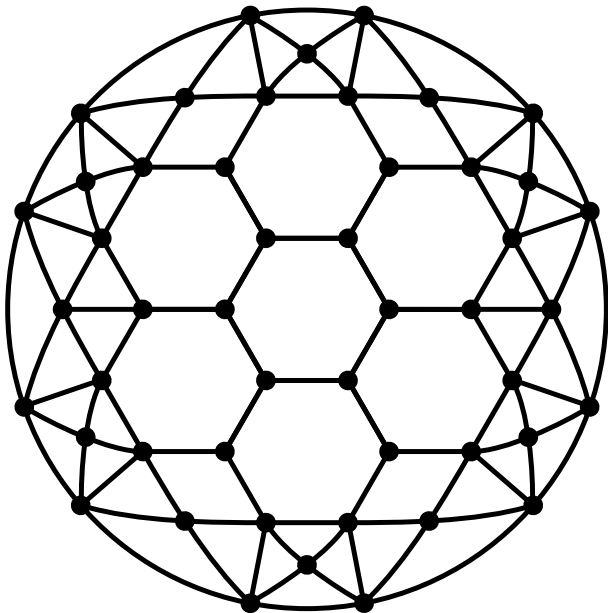
Every 3-con'd G decomposes along its totally-nested
strict tri-separations into parts that are

- quasi-4-con'd
- **generalised wheels**
- thickened $K_{3,m}$



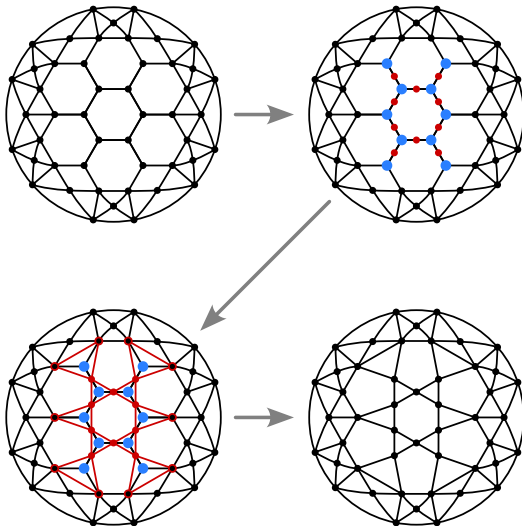
	con'd	
block-cut	↓	
	2-con'd	K_2, K_1
Tutte	↓	
	3-con'd	cycle, K_2
tri-sep'ns	↓	
	quasi-4-con'd	wheel, thickened K_3



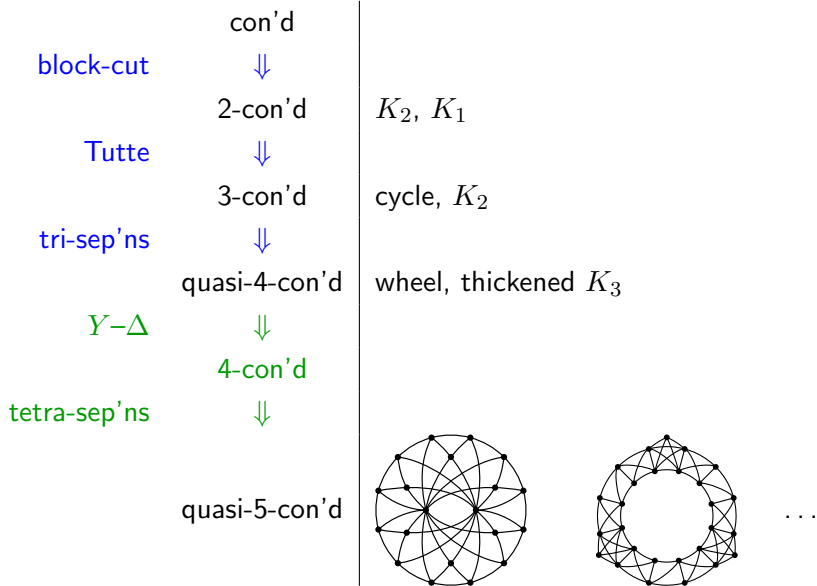


canonical $Y-\Delta$ transformation

quasi-4-con'd



4-con'd



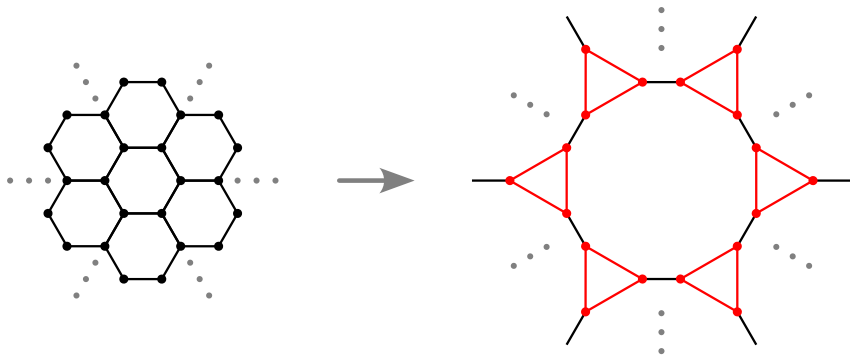
Problem: Classify all **vertex-transitive** finite con'd G

Theorem (Carmesin & K. 23)

Every vertex-transitive finite con'd G is either

- a **cycle**, K_2 , K_1 ,
- **quasi-4-con'd** or

K_3 -expansion of a 3-regular **quasi-4-con'd** arc-transitive graph.



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- **quasi-4-con'd** or

K_3 -expansion of a 3-regular **quasi-4-con'd** arc-transitive graph.

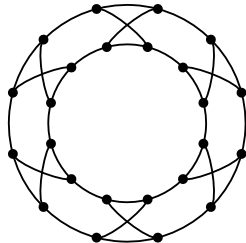
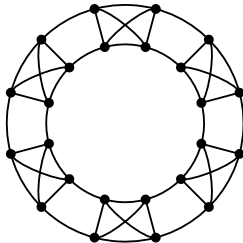
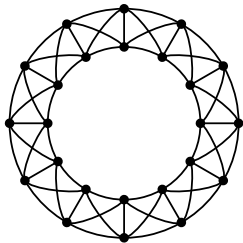
Theorem (K. & Planken 25)

Every **quasi-4-con'd** vertex-transitive finite G is either

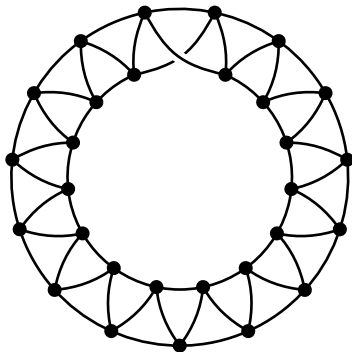
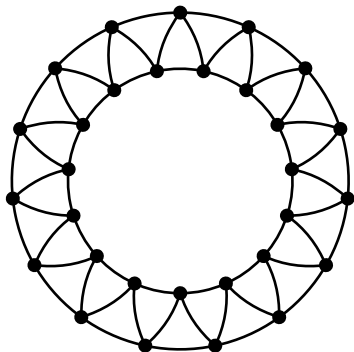
- **bagel-like or cube-like**,
- quasi-5-con'd or

K_4/C_4 -expansion of a 4-regular quasi-5-con'd arc-transitive graph.

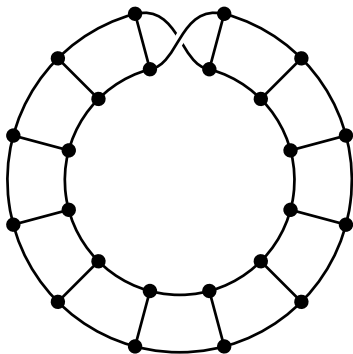
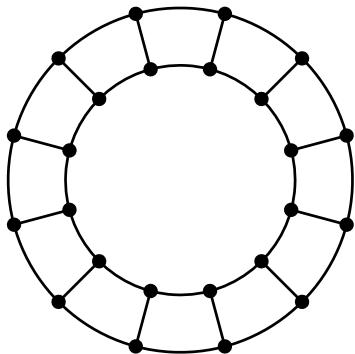
bagel-like



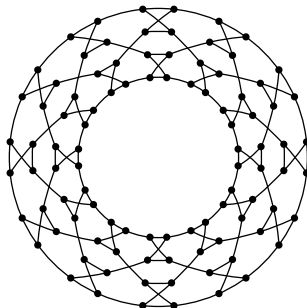
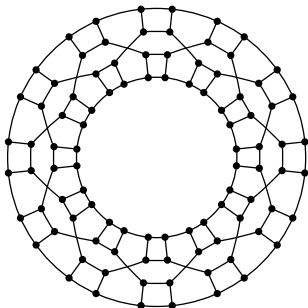
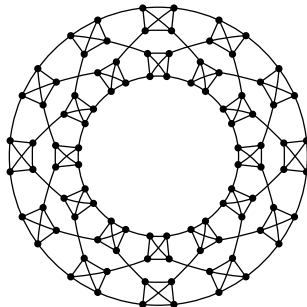
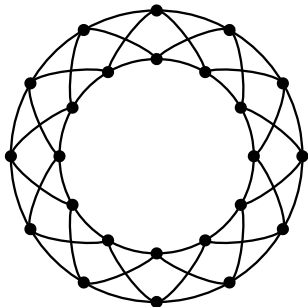
bagel-like



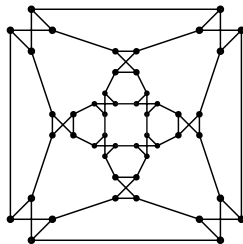
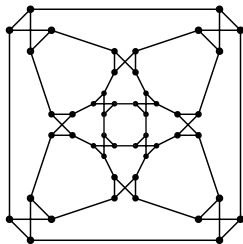
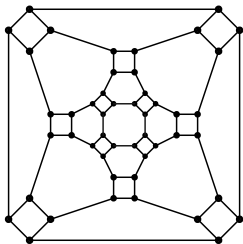
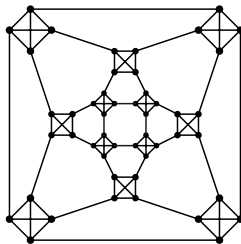
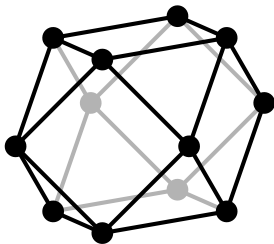
bagel-like



bagel-like



cube-like



Open: Extend Tutte's decomposition to *all* k .

Open: Directed graphs?

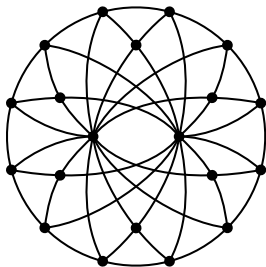
$k = 1$: Bowler, Gut, Hatzel, Kawarabayashi, Muzi, Reich 23

$k \geq 2$: ???

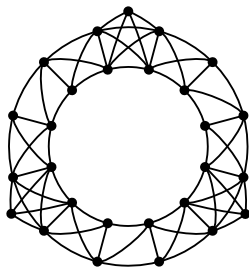
Tetra-separation: mixed-sep'n (A, B) with $|\text{sep}'r| = 4$ such that every vx in $A \cap B$ has ≥ 2 neighb's in $A \setminus B$ and in $B \setminus A$, and cross-edges form matching.

Main result (K. & Planken 25)

Every 4-con'd G decomposes along its totally-nested tetra-separations into parts that are quasi-5-con'd, thickened $K_{4,m}$'s,

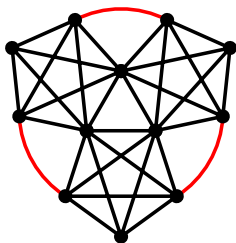


or



Open: Graphs for $k \geq 5$. Digraphs for $k \geq 2$.

Application: Connectivity Augmentation from 0 to 4



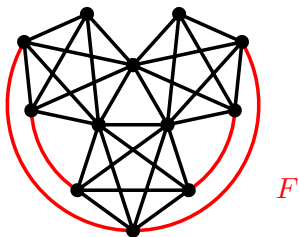
Theorem (Carmesin & Sridharan 25+)

\exists FPT-algorithm with runtime $C(\ell) \cdot \text{Poly}(|V(G)|)$ and

Input: Graph G , $\ell \in \mathbb{N}$ and $F \subseteq E(\overline{G})$

Output: No, or $\leq \ell$ -sized $X \subseteq F$ such that $G + X$ is 4-con'd

Application: Connectivity Augmentation from 0 to 4



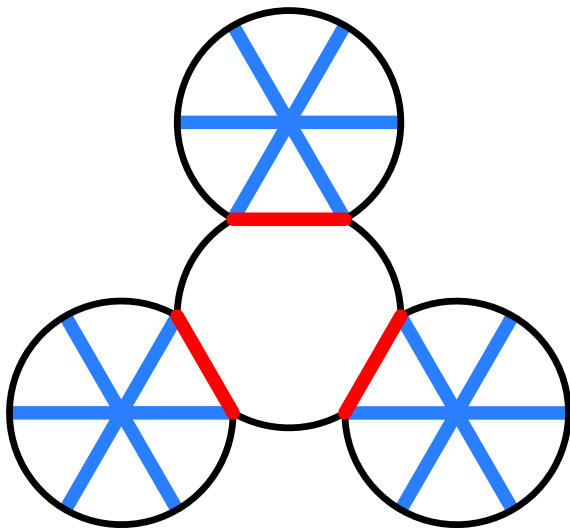
Theorem (Carmesin & Sridharan 25+)

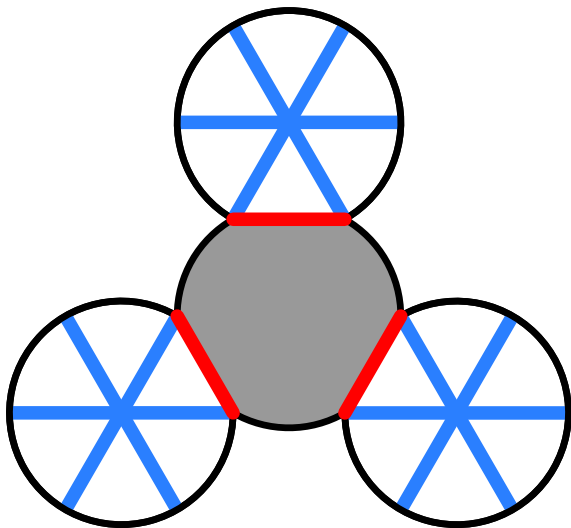
\exists FPT-algorithm with runtime $C(\ell) \cdot \text{Poly}(|V(G)|)$ and

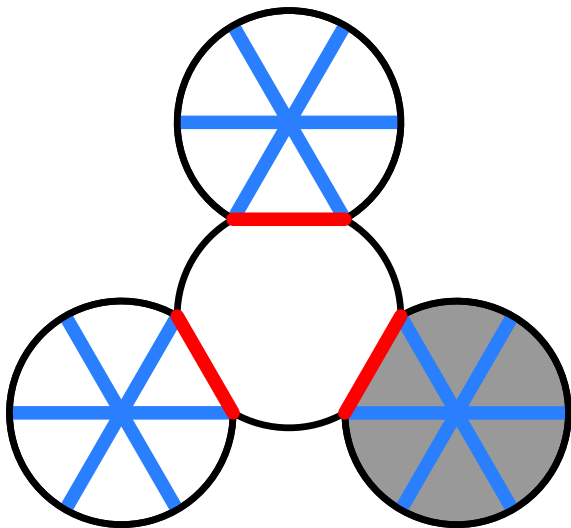
Input: Graph G , $\ell \in \mathbb{N}$ and $F \subseteq E(\overline{G})$

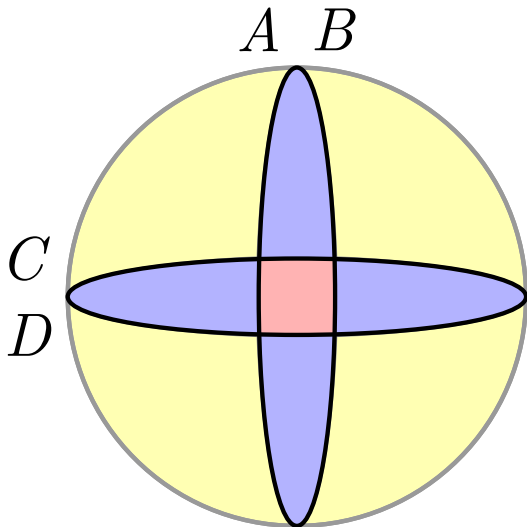
Output: No, or $\leq \ell$ -sized $X \subseteq F$ such that $G + X$ is 4-con'd

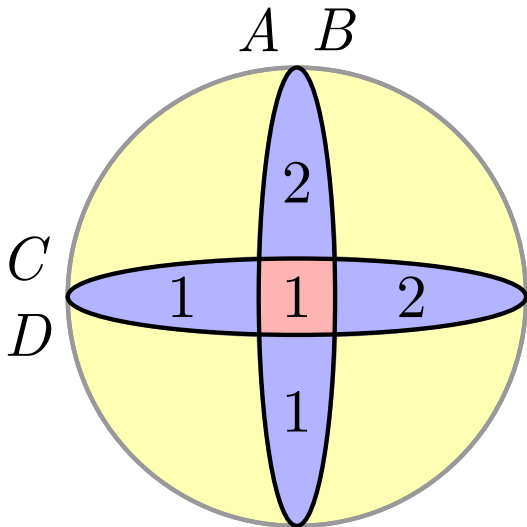
Proof of tetra-decomposition



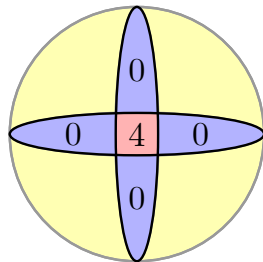
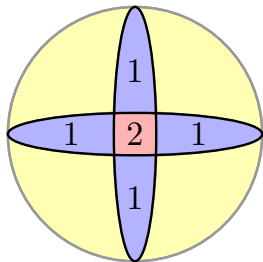
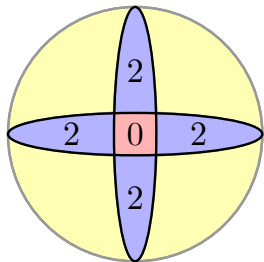




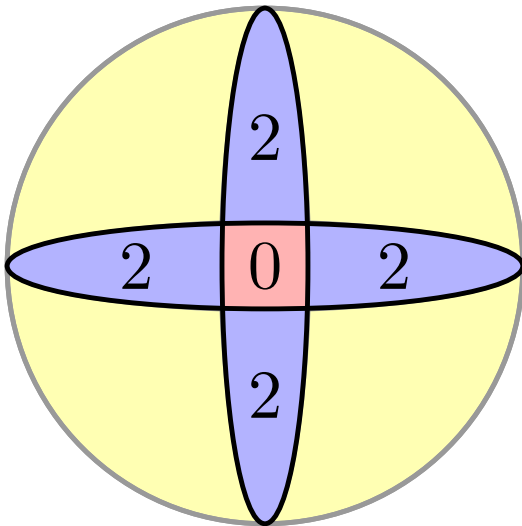


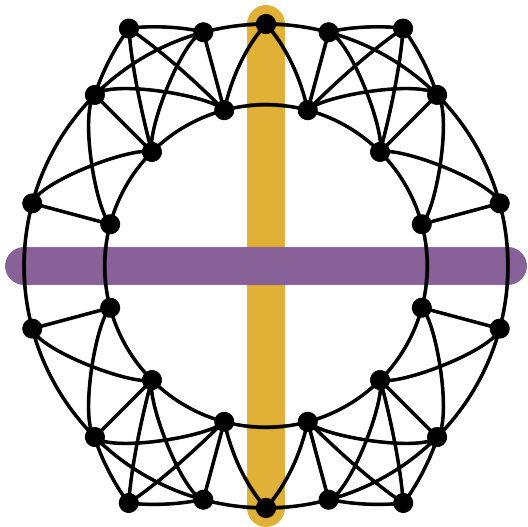


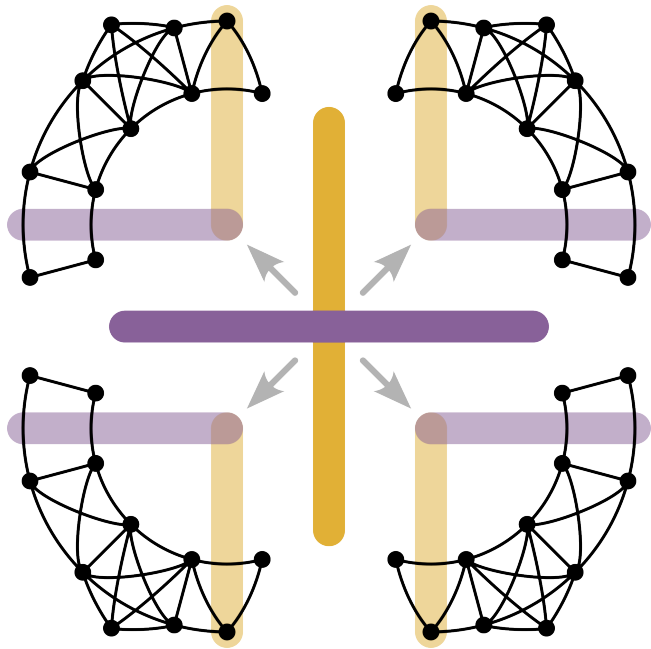
Crossing Lemma. Tetra-sep'ns only cross symmetrically:

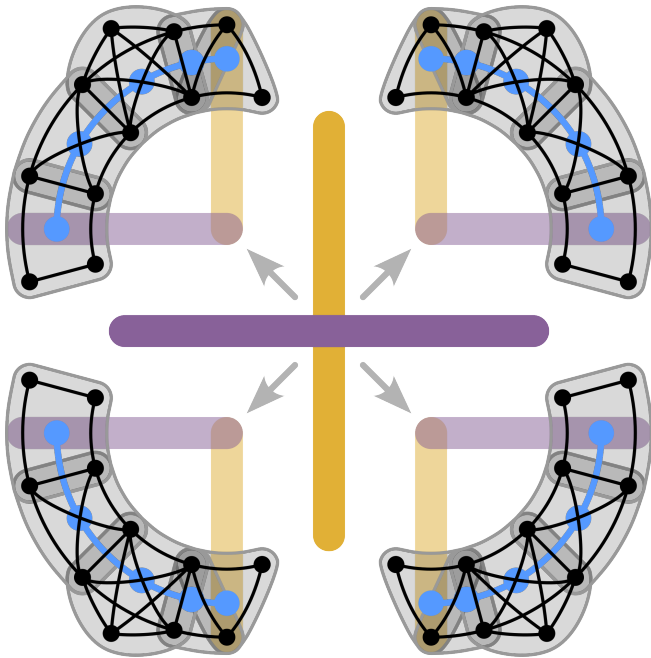


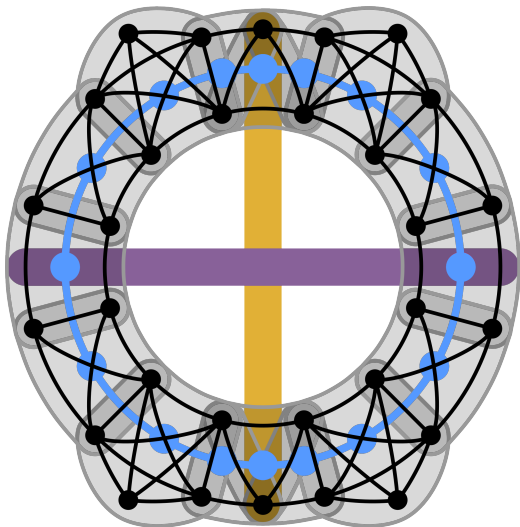
focus

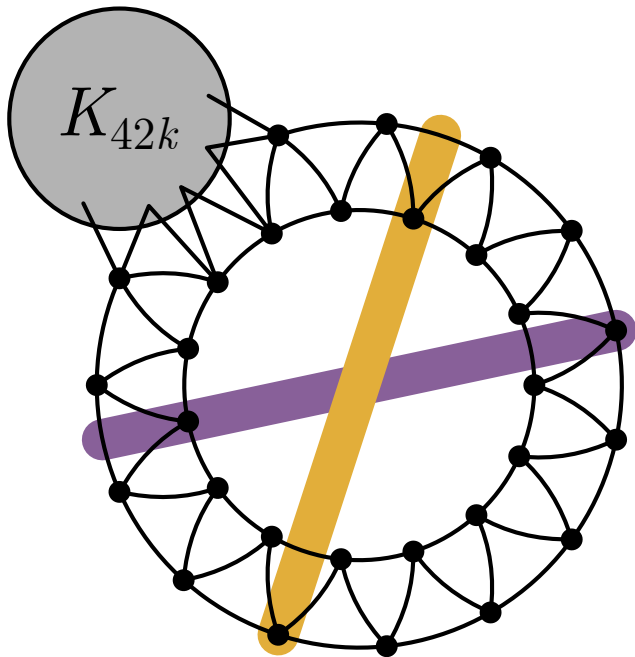


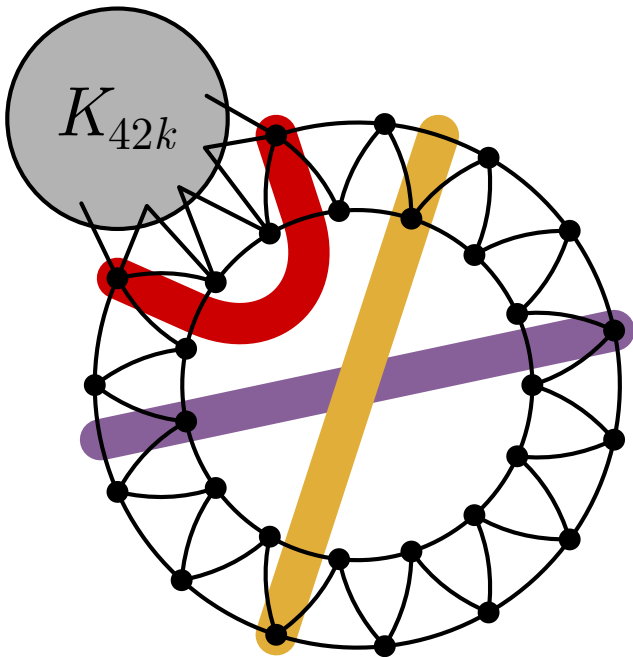












(A, B) totally-nested

\iff the sep'r of (A, B) is highly con'd:

