

Wizards

Standard Code Library

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# 1. 数论

## 1.1 $O(m^2 \log n)$ 线性递推

Given  $a_0, a_1, \dots, a_{m-1}$   
 $a_n = c_0 \times a_{n-m} + \dots + c_{m-1} \times a_{n-1}$   
 Solve for  $a_n = v_0 \times a_0 + v_1 \times a_1 + \dots + v_{m-1} \times a_{m-1}$

```
1 void linear_recurrence(long long n, int m, int a[], int
  ↳ c[], int p) {
2   long long v[M] = {1 % p}; u[M << 1], msk = !n;
3   for(long long i(n); i > 1; i >= 1) {
4     msk <= 1;
5   }
6   for(long long x(0); msk; msk >>= 1, x <= 1) {
7     fill_n(u, m << 1, 0);
8     int b(!(n & msk));
9     x |= b;
10    if(x < m) {
11      u[x] = 1 % p;
12    } else {
13      for(int i(0); i < m; i++) {
14        for(int j(0), t(i + b); j < m; j++, t++) {
15          u[t] = (u[t] + v[i] * v[j]) % p;
16        }
17      }
18      for(int i((m << 1) - 1); i >= m; i--) {
19        for(int j(0), t(i - m); j < m; j++, t++) {
20          u[t] = (u[t] + c[j] * u[i]) % p;
21        }
22      }
23    }
24    copy(u, u + m, v);
25  }
26  //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] *
  ↳ a[m - 1].
27  for(int i(m); i < 2 * m; i++) {
28    a[i] = 0;
29    for(int j(0); j < m; j++) {
30      a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31    }
32  }
33  for(int j(0); j < m; j++) {
34    b[j] = 0;
35    for(int i(0); i < m; i++) {
36      b[j] = (b[j] + v[i] * a[i + j]) % p;
37    }
38  }
39  for(int j(0); j < m; j++) {
40    a[j] = b[j];
41  }
42 }
```

## 1.2 求逆元

```
1 void ex_gcd(long long a, long long b, long long &x, long
  ↳ long &y) {
2   if (b == 0) {
3     x = 1;
4     y = 0;
5     return;
6   }
7   long long xx, yy;
8   ex_gcd(b, a % b, xx, yy);
9   y = xx - a / b * yy;
10  x = yy;
11 }
12 long long inv(long long x, long long MODN) {
13   long long inv_x, y;
14   ex_gcd(x, MODN, inv_x, y);
15   return (inv_x % MODN + MODN) % MODN;
16 }
```

## 1.3 中国剩余定理

```
1 //返回 (ans, M), 其中 ans 是模 M 意义下的解
2 std::pair<long long, long long> CRT(const std::vector<long
  ↳ long>& m, const std::vector<long long>& a) {
3   long long M = 1, ans = 0;
4   int n = m.size();
5   for (int i = 0; i < n; i++) M *= m[i];
6   for (int i = 0; i < n; i++) {
7     ans = (ans + (M / m[i]) * a[i] % M * inv(M / m[i],
  ↳ m[i])) % M; // 可能需要大整数相乘取模
8   }
9   return std::make_pair(ans, M);
10 }
```

## 1.4 素性测试

```
1 int strong_pseudo_primetest(long long n, int base) {
2   long long n2=n-1, res;
3   int s=0;
4   while(n%2==0) n2>>=1, s++;
5   res=powmod(base, n2, n);
6   if((res==1) || (res==n-1)) return 1;
7   s--;
8   while(s>0) {
9     res=mulmod(res, res, n);
10    if(res==n-1) return 1;
11    s--;
12  }
13  return 0; // n is not a strong pseudo prime
14 }
15 int isprime(long long n) {
16   static LL testNum[]={2,3,5,7,11,13,17,19,23,29,31,37};
17   static LL lim[]={4,0,1373653LL,25326001LL,25000000000LL,
18     ↳ 3474749660383LL,341550071728321LL,0,0,0,0};
19   if(n<2 || n==3215031751LL) return 0;
20   for(int i=0; i<12; ++i){
21     if(n<lim[i]) return 1;
22     if(strong_pseudo_primetest(n, testNum[i])==0) return 0;
23   }
24   return 1;
25 }
```

## 1.5 质因数分解

```
1 int ansn; LL ans[1000];
2 LL func(LL x, LL n){ return(mod_mul(x,x,n)+1)%n; }
3 LL Pollard(LL n){
4   LL i,x,y,p;
5   if(Rabin_Miller(n)) return n;
6   if(!(n&1)) return 2;
7   for(i=1; i<20; i++){
8     x=i; y=func(x,n); p=gcd(y-x,n);
9     while(p==1) {x=func(x,n); y=func(y,n,n);
10      ↳ p=gcd((y-x+n)%n,n)%n;}
11     if(p==0 || p==n) continue;
12     return p;
13   }
14 void factor(LL n){
15   LL x;
16   x=Pollard(n);
17   if(x==n){ ans[ansn++]=x; return; }
18   factor(x), factor(n/x);
19 }
```

## 1.6 佩尔方程

```
1 import java.math.BigInteger;
2 import java.util.Scanner;
```

```

3 public class Main {
4     public static BigInteger p, q;
5     public static void solve(int n) {
6         BigInteger N, p1, p2, q1, q2, a0, a1, a2, g1, g2,
            ↪ h1, h2;
7         g1 = q2 = p1 = BigInteger.ZERO;
8         h1 = q1 = p2 = BigInteger.ONE;
9         a0 = a1 =
            ↪ BigInteger.valueOf((long)Math.sqrt(1.0*n));
10        N = BigInteger.valueOf(n);
11        while (true) {
12            g2 = a1.multiply(h1).subtract(g1);
13            h2 = N.subtract(g2.pow(2)).divide(h1);
14            a2 = g2.add(a0).divide(h2);
15            p = a1.multiply(p2).add(p1);
16            q = a1.multiply(q2).add(q1);
17            if (p.pow(2).subtract(N.multiply(q.pow(2)))
18                .compareTo(BigInteger.ONE) == 0) return;
19            g1 = g2; h1 = h2; a1 = a2;
20            p1 = p2; p2 = p;
21            q1 = q2; q2 = q;
22        }
23    }
24    public static void main(String[] args) {
25        Scanner cin = new Scanner(System.in);
26        int t=cin.nextInt();
27        while (t--!=0) {
28            solve(cin.nextInt());
29            System.out.println(p + " " + q);
30        }
31    }
32 }

```

## 1.7 二次剩余

```

1 // x^2 = a (mod p), 0 <= a < p, 返回 true or false 代表
    ↪ 是否存在解
2 // p 必须是质数, 若是多个单次质数的乘积, 可以分别
    ↪ 求解再用 CRT 合并
3 void multiply(ll &c, ll &d, ll a, ll b, ll w) {
4     int cc = (a * c + b * d % MOD * w) % MOD;
5     int dd = (a * d + b * c) % MOD;
6     c = cc, d = dd;
7 }
8 bool solve(int n, int &x) {
9     if (MOD == 2) return x = 1, true;
10    if (power(n, MOD / 2, MOD) == MOD - 1) return false;
11    ll c = 1, d = 0, b = 1, a, w;
12    do { a = rand() % MOD;
13        w = (a * a - n + MOD) % MOD;
14        if (w == 0) return x = a, true;
15    } while (power(w, MOD / 2, MOD) != MOD - 1);
16    for (int times = (MOD + 1) / 2; times; times >>= 1) {
17        if (times & 1) multiply(c, d, a, b, w);
18        multiply(a, b, a, b, w);
19    }
20    // x = (a + sqrt(w)) ^ ((p + 1) / 2)
21    return x = c, true;
22 }

```

## 1.8 一元三次方程

```

1 double a[p[3]], b[p[2]], c[p[1]], d[p[0]];
2 double k(b / a), m(c / a), n(d / a);
3 double p(-k * k / 3. + m);
4 double q(2. * k * k * k / 27 - k * m / 3. + n);
5 Complex omega[3] = {Complex(1, 0), Complex(-0.5, 0.5 *
    ↪ sqrt(3)), Complex(-0.5, -0.5 * sqrt(3))};

```

```

6 Complex r1, r2;
7 double delta(q * q / 4 + p * p * p / 27);
8 if (delta > 0) {
9     r1 = cubrt(-q / 2. + sqrt(delta));
10    r2 = cubrt(-q / 2. - sqrt(delta));
11 } else {
12     r1 = pow(-q / 2. + pow(Complex(delta), 0.5), 1. / 3);
13     r2 = pow(-q / 2. - pow(Complex(delta), 0.5), 1. / 3);
14 }
15 for(int _ (0); _ < 3; _++) {
16     Complex x = -k / 3. + r1 * omega[_ * 1] + r2 * omega[_
    ↪ * 2 % 3];
17 }

```

## 1.9 线下整点

```

1 //  $\sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor$ ,  $n, m, a, b > 0$ 
2 LL solve(LL n, LL a, LL b, LL m) {
3     if (b==0) return n*(a/m);
4     if (a>=m) return n*(a/m)+solve(n, a%m, b, m);
5     if (b>=m) return (n-1)*n/2*(b/m)+solve(n, a, b%m, m);
6     return solve((a+b*n)/m, (a+b*n)%m, m, b);
7 }

```

## 1.10 线性同余不等式

```

1 // Find the minimal non-negative solutions for
    ↪  $l \leq d \cdot x \bmod m \leq r$ 
2 //  $0 \leq d, l, r < m; l \leq r, O(\log n)$ 
3 ll cal(ll m, ll d, ll l, ll r) {
4     if (l == 0) return 0;
5     if (d == 0) return MXL; // 无解
6     if (d * 2 > m) return cal(m, m - d, m - r, m - l);
7     if ((l - 1) / d < r / d) return (l - 1) / d + 1;
8     ll k = cal(d, (-m % d + d) % d, l % d, r % d);
9     return k == MXL ? MXL : (k * m + l - 1) / d + 1;
10 }

```

## 1.11 组合数取模

```

1 LL prod=1,P;
2 pair<LL,LL> comput(LL n,LL p,LL k){
3     if(n<=1)return make_pair(0,1);
4     LL ans=1,cnt=0;
5     ans=pow(prod,n/P,P);
6     cnt=n/p;
7     pair<LL,LL>res=comput(n/p,p,k);
8     cnt+=res.first;
9     ans=ans*res.second%P;
10    for(int i=n-n%P+1;i<=n;i++)if(i%p){
11
12        ans=ans*i%P;
13    }
14    return make_pair(cnt,ans);
15 }
16 pair<LL,LL> calc(LL n,LL p,LL k){
17     prod=1;P=pow(p,k,1e18);
18     for(int i=1;i<P;i++)if(i%p)prod=prod*i%P;
19     pair<LL,LL> res=comput(n,p,k);
20     return res;
21 }
22 LL calc(LL n,LL m,LL p,LL k){
23     pair<LL,LL>A,B,C;
24     LL P=pow(p,k,1e18);
25     A=calc(n,p,k);
26     B=calc(m,p,k);
27     C=calc(n-m,p,k);
28     LL ans=1;
29     ans=pow(p,A.first-B.first-C.first,P);

```

```

30     ↪ ans=ans*A.second%P*inv(B.second,P)%P*inv(C.second,P)%P;
31     return ans;
32 }

```

## 2. 代数

### 2.1 快速傅里叶变换

```

1 void fft(Complex a[], int n, int f) {
2     for (int i = 0; i < n; ++i)
3         if (R[i] < i) swap(a[i], a[R[i]]);
4     for (int i = 1, h = 0; i < n; i <= 1, h++) {
5         Complex wn = Complex(cos(pi / i), f * sin(pi / i));
6         Complex w = Complex(1, 0);
7         for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
8         for (int p = i < 1, j = 0; j < n; j += p) {
9             for (int k = 0; k < i; ++k) {
10                Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11                a[j + k] = x + y; a[j + k + i] = x - y;
12            }
13        }
14    }
15 }

```

### 2.2 快速数论变换

```

1 // n 必须是 2 的次幂
2 void fft(Complex a[], int n, int f) {
3     for (int i = 0; i < n; ++i)
4         if (R[i] < i) swap(a[i], a[R[i]]);
5     for (int i = 1, h = 0; i < n; i <= 1, h++) {
6         Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7         Complex w = Complex(1, 0);
8         for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9         for (int p = i < 1, j = 0; j < n; j += p) {
10            for (int k = 0; k < i; ++k) {
11                Complex x = a[j + k], y = a[j + k + i] * tmp[k];
12                a[j + k] = x + y; a[j + k + i] = x - y;
13            }
14        }
15    }
16 }

```

### 2.3 快速沃尔什变换

```

1 void FWT(LL a[], int n, int ty) { //the length is 2^n
2     for (int d = 1; d < n; d <= 1) {
3         for (int m = (d < 1), i = 0; i < n; i += m) {
4             if (ty == 1) {
5                 for (int j = 0; j < d; j++) {
6                     LL x = a[i + j], y = a[i + j + d];
7                     a[i + j] = x + y;
8                     a[i + j + d] = x - y;
9                     //and: a[i + j] = x + y; or: a[i + j + d] = x + y;
10                }
11            } else {
12                for (int j = 0; j < d; j++) {
13                    LL x = a[i + j], y = a[i + j + d];
14                    a[i + j] = (x + y) / 2;
15                    a[i + j + d] = (x - y) / 2;
16                    //and: a[i + j] = x - y; or: a[i + j + d] = y - x;
17                }
18            }
19        }
20    }
21 }

```

## 2.4 自适应辛普森积分

```

1 namespace adaptive_simpson {
2     template<typename function>
3     inline double area(function f, const double &left, const
4         ↪ double &right) {
5         double mid = (left + right) / 2;
6         return (right - left) * (f(left) + 4 * f(mid) +
7             ↪ f(right)) / 6;
8     }
9     template<typename function>
10    inline double simpson(function f, const double &left,
11        ↪ const double &right, const double &eps, const
12        ↪ double &area_sum) {
13        double mid = (left + right) / 2;
14        double area_left = area(f, left, mid);
15        double area_right = area(f, mid, right);
16        double area_total = area_left + area_right;
17        if (fabs(area_total - area_sum) <= 15 * eps) {
18            return area_total + (area_total - area_sum) / 15;
19        }
20        return simpson(f, left, right, eps / 2, area_left) +
21            ↪ simpson(f, mid, right, eps / 2, area_right);
22    }
23     template<typename function>
24     inline double simpson(function f, const double &left,
25        ↪ const double &right, const double &eps) {
26        return simpson(f, left, right, eps, area(f, left,
27            ↪ right));
28    }
29 }

```

## 2.5 单纯形

```

1 const double eps = 1e-8;
2 // max{c * x | Ax <= b, x >= 0} 的解, 无解返回空的
3 ↪ vector, 否则就是解.
4 vector<double> simplex(vector<vector<double>> &A,
5     ↪ vector<double> b, vector<double> c) {
6     int n = A.size(), m = A[0].size() + 1, r = n, s = m - 1;
7     vector<vector<double>> D(n + 2, vector<double>(m + 1));
8     vector<int> ix(n + m);
9     for (int i = 0; i < n + m; i++) {
10         ix[i] = i;
11     }
12     for (int i = 0; i < n; i++) {
13         for (int j = 0; j < m - 1; j++) {
14             D[i][j] = -A[i][j];
15         }
16         D[i][m - 1] = 1;
17         D[i][m] = b[i];
18         if (D[r][m] > D[i][m]) {
19             r = i;
20         }
21     }
22     for (int j = 0; j < m - 1; j++) {
23         D[n][j] = c[j];
24     }
25     D[n + 1][m - 1] = -1;
26     for (double d; ;) {
27         if (r < n) {
28             swap(ix[s], ix[r + m]);
29             D[r][s] = 1. / D[r][s];
30             for (int j = 0; j <= m; j++) {
31                 if (j != s) {
32                     D[r][j] *= -D[r][s];
33                 }
34             }
35             for (int i = 0; i <= n + 1; i++) {
36                 if (i != r) {
37                     for (int j = 0; j <= m; j++) {
38                         if (j != s) {

```

```

37         D[i][j] += D[r][j] * D[i][s];
38     }
39 }
40 D[i][s] *= D[r][s];
41 }
42 }
43 }
44 r = -1, s = -1;
45 for(int j = 0; j < m; j++) {
46     if (s < 0 || ix[s] > ix[j]) {
47         if (D[n + 1][j] > eps || D[n + 1][j] > -eps &&
48             ↪ D[n][j] > eps) {
49             s = j;
50         }
51     }
52     if (s < 0) break;
53     for(int i = 0; i < n; i++) {
54         if (D[i][s] < -eps) {
55             if (r < 0 || (d = D[r][m] / D[r][s] - D[i][m] /
56                 ↪ D[i][s]) < -eps
57                 || d < eps && ix[r + m] > ix[i + m]) {
58                 r = i;
59             }
60         }
61         if (r < 0) return vector<double> ();
62     }
63     if (D[n + 1][m] < -eps) return vector<double> ();
64     vector<double> x(m - 1);
65     for(int i = m; i < n + m; i++) {
66         if (ix[i] < m - 1) {
67             x[ix[i]] = D[i - m][m];
68         }
69     }
70     return x;
71 }

```

## 3. 计算几何

### 3.1 二维

#### 3.1.1 点类

```

1 int sign(DB x) {
2     return (x > eps) - (x < -eps);
3 }
4 DB msqrt(DB x) {
5     return sign(x) > 0 ? sqrt(x) : 0;
6 }
7 struct Point {
8     DB x, y;
9     Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
10         return Point(cos(ang) * x - sin(ang) * y, cos(ang) * y
11             ↪ + sin(ang) * x);
12     }
13     Point turn90() const { // 逆时针旋转 90 度
14         return Point(-y, x);
15     }
16     Point unit() const {
17         return *this / len();
18     }
19 };
20 DB dot(const Point& a, const Point& b) {
21     return a.x * b.x + a.y * b.y;
22 }
23 DB det(const Point& a, const Point& b) {
24     return a.x * b.y - a.y * b.x;
25 }
26 #define cross(p1,p2,p3)
27     ↪ ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
28 #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
29 bool isLL(const Line& l1, const Line& l2, Point& p) {

```

```

30     DB s1 = det(l2.b - l2.a, l1.a - l2.a),
31     s2 = -det(l2.b - l2.a, l1.b - l2.a);
32     if (!sign(s1 + s2)) return false;
33     p = (l1.a * s2 + l1.b * s1) / (s1 + s2);
34     return true;
35 }
36 bool onSeg(const Line& l, const Point& p) {
37     return sign(det(p - l.a, l.b - l.a)) == 0 && sign(dot(p
38         ↪ - l.a, p - l.b)) <= 0;
39 }
40 Point projection(const Line & l, const Point& p) {
41     return l.a + (l.b - l.a) * (dot(p - l.a, l.b - l.a) /
42         ↪ (l.b - l.a).len2());
43 }
44 DB disToLine(const Line& l, const Point& p) {
45     return fabs(det(p - l.a, l.b - l.a) / (l.b -
46         ↪ l.a).len());
47 }
48 DB disToSeg(const Line& l, const Point& p) {
49     return sign(dot(p - l.a, l.b - l.a)) * sign(dot(p - l.b,
50         ↪ l.a - l.b)) == 1 ? disToLine(l, p) : std::min((p -
51         ↪ l.a).len(), (p - l.b).len());
52 }
53 // 圆与直线交点
54 bool isCL(Circle a, Line l, Point& p1, Point& p2) {
55     DB x = dot(l.a - a.o, l.b - l.a),
56     y = (l.b - l.a).len2(),
57     d = x * x - y * ((l.a - a.o).len2() - a.r * a.r);
58     if (sign(d) < 0) return false;
59     Point p = l.a - ((l.b - l.a) * (x / y)), delta = (l.b -
60         ↪ l.a) * (msqrt(d) / y);
61     p1 = p + delta; p2 = p - delta;
62     return true;
63 }
64 // 圆与圆的交面积
65 DB areaCC(const Circle& c1, const Circle& c2) {
66     DB d = (c1.o - c2.o).len();
67     if (sign(d - (c1.r + c2.r)) >= 0) return 0;
68     if (sign(d - std::abs(c1.r - c2.r)) <= 0) {
69         DB r = std::min(c1.r, c2.r);
70         return r * r * PI;
71     }
72     DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
73     t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
74     return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r *
75         ↪ sin(t1);
76 }
77 // 圆与圆交点
78 bool isCC(Circle a, Circle b, P& p1, P& p2) {
79     DB s1 = (a.o - b.o).len();
80     if (sign(s1 - a.r - b.r) > 0 || sign(s1 - std::abs(a.r -
81         ↪ b.r)) < 0) return false;
82     DB s2 = (a.r * a.r - b.r * b.r) / s1;
83     DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
84     P o = (b.o - a.o) * (aa / (aa + bb)) + a.o;
85     P delta = (b.o - a.o).unit().turn90() * msqrt(a.r * a.r
86         ↪ - aa * aa);
87     p1 = o + delta, p2 = o - delta;
88     return true;
89 }
90 // 求点到圆的切点, 按关于点的顺时针方向
91 bool tanCP(const Circle &c, const Point &p0, Point &p1,
92     ↪ Point &p2) {
93     double x = (p0 - c.o).len2(), d = x - c.r * c.r;
94     if (d < eps) return false; // 点在圆上认为没有切点
95     Point p = (p0 - c.o) * (c.r * c.r / x);
96     Point delta = ((p0 - c.o) * (-c.r * sqrt(d) /
97         ↪ x)).turn90();
98     p1 = c.o + p + delta;
99     p2 = c.o + p - delta;
100     return true;
101 }

```

```

89 // 求圆到圆的外共切线, 按关于 c1.o 的顺时针方向
90 vector<Line> extanCC(const Circle &c1, const Circle &c2) {
91     vector<Line> ret;
92     if (sign(c1.r - c2.r) == 0) {
93         Point dir = c2.o - c1.o;
94         dir = (dir * (c1.r / dir.len())).turn90();
95         ret.push_back(Line(c1.o + dir, c2.o + dir));
96         ret.push_back(Line(c1.o - dir, c2.o - dir));
97     } else {
98         Point p = (c1.o * -c2.r + c2.o * c1.r) / (c1.r -
99             ↪ c2.r);
100         Point p1, p2, q1, q2;
101         if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) {
102             if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);
103             ret.push_back(Line(p1, q1));
104             ret.push_back(Line(p2, q2));
105         }
106     }
107     return ret;
108 }
109 // 求圆到圆的内共切线, 按关于 c1.o 的顺时针方向
110 std::vector<Line> intanCC(const Circle &c1, const Circle
111     ↪ &c2) {
112     std::vector<Line> ret;
113     Point p = (c1.o * c2.r + c2.o * c1.r) / (c1.r + c2.r);
114     Point p1, p2, q1, q2;
115     if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { //
116         ↪ 两圆相切认为没有切线
117         ret.push_back(Line(p1, q1));
118         ret.push_back(Line(p2, q2));
119     }
120     return ret;
121 }
122 // 判断点
123 bool contain(vector<Point> polygon, Point p) { // 判断点
124     ↪ p 是否被多边形包含, 包括落在边界上
125     int ret = 0, n = polygon.size();
126     for (int i = 0; i < n; ++i) {
127         Point u = polygon[i], v = polygon[(i + 1) % n];
128         if (onSeg(Line(u, v), p)) return true; // Here I
129         ↪ guess.
130         if (sign(u.y - v.y) <= 0) swap(u, v);
131         if (sign(p.y - u.y) > 0 || sign(p.y - v.y) <= 0)
132             ↪ continue;
133         ret += sign(det(p, v, u)) > 0;
134     }
135     return ret & 1;
136 }
137 // 用半平面 (q1,q2) 的逆时针方向去切凸多边形
138 std::vector<Point> convexCut(const std::vector<Point>&ps,
139     ↪ Point q1, Point q2) {
140     std::vector<Point> qs; int n = ps.size();
141     for (int i = 0; i < n; ++i) {
142         Point p1 = ps[i], p2 = ps[(i + 1) % n];
143         int d1 = crossOp(q1, q2, p1), d2 = crossOp(q1, q2, p2);
144         if (d1 >= 0) qs.push_back(p1);
145         if (d1 * d2 < 0) qs.push_back(isSS(p1, p2, q1, q2));
146     }
147     return qs;
148 }
149 // 求凸包
150 std::vector<Point> convexHull(std::vector<Point> ps) {
151     int n = ps.size(); if (n <= 1) return ps;
152     std::sort(ps.begin(), ps.end());
153     std::vector<Point> qs;
154     for (int i = 0; i < n; qs.push_back(ps[i ++]))
155         while (qs.size() > 1 && sign(det(qs[qs.size() - 2],
156             ↪ qs.back(), ps[i])) <= 0)
157             qs.pop_back();
158     for (int i = n - 2, t = qs.size(); i >= 0;
159         ↪ qs.push_back(ps[i --]))
160         while ((int)qs.size() > t && sign(det(qs[qs.size() -
161             ↪ 2], qs.back(), ps[i])) <= 0)
162             qs.pop_back();

```

```

152     return qs;
153 }

```

### 3.1.2 凸包

```

1 // 凸包中的点按逆时针方向
2 struct Convex {
3     int n;
4     std::vector<Point> a, upper, lower;
5     void make_shell(const std::vector<Point>& p,
6         std::vector<Point>& shell) { // p needs to be
7         ↪ sorted.
8         clear(shell); int n = p.size();
9         for (int i = 0, j = 0; i < n; i++, j++) {
10             for (; j >= 2 && sign(det(shell[j-1] - shell[j-2],
11                 ↪ p[i] - shell[j-2])) <= 0; --j)
12                 shell.pop_back();
13             shell.push_back(p[i]);
14         }
15     }
16     void make_convex() {
17         std::sort(a.begin(), a.end());
18         make_shell(a, lower);
19         std::reverse(a.begin(), a.end());
20         make_shell(a, upper);
21         a = lower; a.pop_back();
22         a.insert(a.end(), upper.begin(), upper.end());
23         if ((int)a.size() >= 2) a.pop_back();
24         n = a.size();
25     }
26     void init(const std::vector<Point>& _a) {
27         clear(a); a = _a; n = a.size();
28         make_convex();
29     }
30     void read(int _n) { // Won't make convex.
31         clear(a); n = _n; a.resize(n);
32         for (int i = 0; i < n; i++)
33             a[i].read();
34     }
35     std::pair<DB, int> get_tangent(
36         const std::vector<Point>& convex, const Point& vec)
37         ↪ {
38         int l = 0, r = (int)convex.size() - 2;
39         assert(r >= 0);
40         for (; l + 1 < r; ) {
41             int mid = (l + r) / 2;
42             if (sign(det(convex[mid + 1] - convex[mid], vec)) >
43                 ↪ 0)
44                 r = mid;
45             else l = mid;
46         }
47         return std::max(std::make_pair(det(vec, convex[r]),
48             ↪ r),
49             std::make_pair(det(vec, convex[0]), 0));
50     }
51     void binary_search(Point u, Point v, int l, int r) {
52         int s1 = sign(det(v - u, a[l % n] - u));
53         for (; l + 1 < r; ) {
54             int mid = (l + r) / 2;
55             int smid = sign(det(v - u, a[mid % n] - u));
56             if (smid == s1) l = mid;
57             else r = mid;
58         }
59         return l % n;
60     }
61     // 求凸包上和向量 vec 叉积最大的点, 返回编号, 共
62     ↪ 线的多个切点返回任意一个
63     int get_tangent(Point vec) {
64         std::pair<DB, int> ret = get_tangent(upper, vec);
65         ret.second = (ret.second + (int)lower.size() - 1) % n;
66         ret = std::max(ret, get_tangent(lower, vec));
67         return ret.second;

```



```

62 }
63 // 求凸包和直线 u, v 的交点, 如果不相交返回 false,
    ↳ 如果有则是和 (i, next(i)) 的交点, 交在点上不
    ↳ 确定返回前后两条边其中之一
64 bool get_intersection(Point u, Point v, int &i0, int
    ↳ &i1) {
65     int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
66     if (sign(det(v - u, a[p0] - u)) * sign(det(v - u,
        ↳ a[p1] - u)) <= 0) {
67         if (p0 > p1) std::swap(p0, p1);
68         i0 = binary_search(u, v, p0, p1);
69         i1 = binary_search(u, v, p1, p0 + n);
70         return true;
71     }
72     else return false;
73 }
74 };

```

### 3.1.3 凸包最近点对

```

1 //判断点是否在多边形内
2 int isPointInPolygon(point p, point *a, int n) {
3     int cnt = 0;
4     for(int i=0; i<n; ++i) {
5         if(OnSegment(p, a[i], a[(i+1)%n])) return -1;
6         double k = cross(a[(i+1)%n]-a[i], p-a[i]);
7         double d1 = a[i].y - p.y;
8         double d2 = a[(i+1)].y - p.y;
9         if(k>0 &&d1<=0 &&d2>0) cnt++;
10        if(k<0 &&d2<=0 &&d1>0) cnt++;
11        //k=0, 点和线段共线的情况不考虑
12    }
13    if(cnt&1)return 1;
14    return 0;
15 }
16 //判断凸包是否相离
17 bool two_getaway_ConvexHull(point *cha, int n1, point
    ↳ *chb, int m1) {
18     if(n1==1 && m1==1) {
19         if(cha[0]==chb[0])
20             return false;
21     } else if(n1==1 && m1==2) {
22         if(OnSegment(cha[0], chb[0], chb[1]))
23             return false;
24     } else if(n1==2 && m1==1) {
25         if(OnSegment(chb[0], cha[0], cha[1]))
26             return false;
27     } else if(n1==2 && m1==2) {
28         if(SegmentIntersection(cha[0], cha[1], chb[0],
        ↳ chb[1]))
29             return false;
30     } else if(n1==2) {
31         for(int i=0; i<n1; ++i)
32             if(isPointInPolygon(cha[i], chb, m1))
33                 return false;
34     } else if(m1==2) {
35         for(int i=0; i<m1; ++i)
36             if(isPointInPolygon(chb[i], cha, n1))
37                 return false;
38     } else {
39         for(int i=0; i<n1; ++i) {
40             for(int j=0; j<m1; ++j) {
41                 if(SegmentIntersection(cha[i],
        ↳ cha[(i+1)%n1], chb[j],
        ↳ chb[(j+1)%m1]))
42                     return false;
43             }
44         }
45         for(int i=0; i<n1; ++i)
46             if(isPointInPolygon(cha[i], chb, m1))
47                 return false;
48         for(int i=0; i<m1; ++i)

```

```

49         if(isPointInPolygon(chb[i], cha, n1))
50             return false;
51     }
52     return true;
53 }
54 //旋转卡壳求两个凸包最近距离
55 double solve(point *P, point *Q, int n, int m) {
56     if(n==1 && m==1) {
57         return length(P[0] - Q[0]);
58     } else if(n==1 && m==2) {
59         return DistanceToSegment(P[0], Q[0], Q[1]);
60     } else if(n==2 && m==1) {
61         return DistanceToSegment(Q[0], P[0], P[1]);
62     } else if(n==2 && m==2) {
63         return SegmentToSegment(P[0], P[1], Q[0], Q[1]);
64     }
65
66     int yminP = 0, ymaxQ = 0;
67     for(int i=0; i<n; ++i) if(P[i].y < P[yminP].y) yminP =
        ↳ i;
68     for(int i=0; i<m; ++i) if(Q[i].y > Q[ymaxQ].y) ymaxQ =
        ↳ i;
69     P[n] = P[0];
70     Q[m] = Q[0];
71     double INF2 = 1e100;
72     double arg, ans = INF2;
73     for(int i=0; i<n; ++i) {
74         //当叉积负正转正时, 说明点 ymaxQ 就是对踵点
75         while((arg=cross(P[yminP] - P[yminP+1], Q[ymaxQ+1]
        ↳ - Q[ymaxQ])) < -eps)
76             ymaxQ = (ymaxQ+1)%m;
77         double ret;
78         if(arg > eps) { //卡住第二个凸包上的点。
79             ret = DistanceToSegment(Q[ymaxQ], P[yminP],
        ↳ P[yminP+1]);
80             ans = min(ans, ret);
81         } else { //arg==0, 卡住第二个凸包的边
82             ret = SegmentToSegment(P[yminP], P[yminP+1],
        ↳ Q[ymaxQ], Q[ymaxQ+1]);
83             ans = min(ans, ret);
84         }
85         yminP = (yminP+1)%n;
86     }
87     return ans;
88 }
89 }
90 double mindis_twotubao(point *P, point *Q, int n, int m){
91     //return min(solve(P, Q, n, m), solve(Q, P, m, n));
92     if(two_getaway_ConvexHull(P, n, Q, m)==true) return
        ↳ min(solve(P, Q, n, m), solve(Q, P, m, n));
93     else return 0.0;
94 }

```

### 3.1.4 三角形的心

```

1 Point inCenter(const Point &A, const Point &B, const Point
    ↳ &C) { // 内心
2     double a = (B - C).len(), b = (C - A).len(), c = (A -
        ↳ B).len(),
3     s = fabs(det(B - A, C - A)),
4     r = s / p;
5     return (A * a + B * b + C * c) / (a + b + c);
6 }
7 Point circumCenter(const Point &a, const Point &b, const
    ↳ Point &c) { // 外心
8     Point bb = b - a, cc = c - a;
9     double db = bb.len2(), dc = cc.len2(), d = 2 * det(bb,
        ↳ cc);
10    return a - Point(bb.y * dc - cc.y * db, cc.x * db - bb.x
        ↳ * dc) / d;
11 }
12 Point orthoCenter(const Point &a, const Point &b, const
    ↳ Point &c) { // 垂心

```



```

13 Point ba = b - a, ca = c - a, bc = b - c;
14 double Y = ba.y * ca.y * bc.y,
15     A = ca.x * ba.y - ba.x * ca.y,
16     x0 = (Y + ca.x * ba.y * b.x - ba.x * ca.y * c.x) /
17         ↪ A,
18     y0 = -ba.x * (x0 - c.x) / ba.y + ca.y;
19 return Point(x0, y0);
20 }

```

### 3.1.5 半平面交

```

1 struct Point {
2     int quad() const { return sign(y) == 1 || (sign(y) == 0
3         ↪ && sign(x) >= 0); }
4 };
5 struct Line {
6     bool include(const Point &p) const { return sign(det(b -
7         ↪ a, p - a)) > 0; }
8     Line push() const { // 将半平面向外推 eps
9         const double eps = 1e-6;
10        Point delta = (b - a).turn90().norm() * eps;
11        return Line(a - delta, b - delta);
12    }
13 };
14 bool sameDir(const Line &l0, const Line &l1) { return
15     ↪ parallel(l0, l1) && sign(dot(l0.b - l0.a, l1.b -
16     ↪ l1.a)) == 1; }
17 bool operator < (const Point &a, const Point &b) {
18     if (a.quad() != b.quad()) {
19         return a.quad() < b.quad();
20     } else {
21         return sign(det(a, b)) > 0;
22     }
23 }
24 bool operator < (const Line &l0, const Line &l1) {
25     if (sameDir(l0, l1)) {
26         return l1.include(l0.a);
27     } else {
28         return (l0.b - l0.a) < (l1.b - l1.a);
29     }
30 }
31 bool check(const Line &u, const Line &v, const Line &w) {
32     ↪ return w.include(intersect(u, v)); }
33 vector<Point> intersection(vector<Line> &l) {
34     sort(l.begin(), l.end());
35     deque<Line> q;
36     for (int i = 0; i < (int)l.size(); ++i) {
37         if (i && sameDir(l[i], l[i - 1])) {
38             continue;
39         }
40         while (q.size() > 1 && !check(q[q.size() - 2],
41             ↪ q[q.size() - 1], l[i])) q.pop_back();
42         while (q.size() > 1 && !check(q[1], q[0], l[i]))
43             ↪ q.pop_front();
44         q.push_back(l[i]);
45     }
46     while (q.size() > 2 && !check(q[q.size() - 2],
47         ↪ q[q.size() - 1], q[0])) q.pop_back();
48     while (q.size() > 2 && !check(q[1], q[0], q[q.size() -
49         ↪ 1])) q.pop_front();
50     vector<Point> ret;
51     for (int i = 0; i < (int)q.size(); ++i)
52         ↪ ret.push_back(intersect(q[i], q[(i + 1) %
53         ↪ q.size()]));
54     return ret;
55 }

```

### 3.1.6 最大空凸包

```

1 inline double eq(double x, double y) {
2     return fabs(x-y)<eps;
3 }

```

```

4 double xmult(point a, point b, point o) {
5     return (a.x-o.x)*(o.y-b.y)-(a.y-o.y)*(o.x-b.x);
6 }
7 double dist(point a, point b) {
8     return (a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y);
9 }
10 point o;
11 bool cmp_angle(point a, point b) {
12     if(eq(xmult(a,b,o),0.0)) {
13         return dist(a,o)<dist(b,o);
14     }
15     return xmult(a,o,b)>0;
16 }
17 double empty_convex(point *p, int pn) {
18     double ans=0;
19     for(int i=0; i<pn; i++) {
20         for(int j=0; j<pn; j++) {
21             dp[i][j]=0;
22         }
23     }
24     for(int i=0; i<pn; i++) {
25         int j = i-1;
26         while(j>=0 && eq(xmult(p[i], p[j],
27             ↪ o),0.0))j--;//coline
28         bool flag= j==i-1;
29         while(j>=0) {
30             int k = j-1;
31             while(k >= 0 && xmult(p[i],p[k],p[j])>0)k--;
32             double area = fabs(xmult(p[i],p[j],o))/2;
33             if(k >= 0)area+=dp[j][k];
34             if(flag) dp[i][j]=area;
35             ans=max(ans,area);
36             j=k;
37         }
38         if(flag) {
39             for(int j=1; j<i; j++) {
40                 dp[i][j] = max(dp[i][j],dp[i][j-1]);
41             }
42         }
43     }
44     return ans;
45 }
46 double largest_empty_convex(point *p, int pn) {
47     point data[maxn];
48     double ans=0;
49     for(int i=0; i<pn; i++) {
50         o=p[i];
51         int dn=0;
52         for(int j=0; j<pn; j++) {
53             if(p[j].y>o.y|| (p[j].y==o.y&&p[j].x>o.x)) {
54                 data[dn++]=p[j];
55             }
56         }
57         sort(data, data+dn, cmp_angle);
58         ans=max(ans, empty_convex(data, dn));
59     }
60     return ans;
61 }

```

### 3.1.7 平面最近点对

```

1 double Dis(Point a, Point b) {
2     return sqrt((a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y));
3 }
4 double Closest_Pair(int left, int right) {
5     double d = INF;
6     if(left == right) return d;
7     if(left + 1 == right)
8         return Dis(p[left],p[right]);
9     int mid = (left+right)>>1;
10    double d1 = Closest_Pair(left,mid);
11    double d2 = Closest_Pair(mid,right);

```

```

12 d = min(d1,d2);
13 int k = 0;
14 for(int i = left; i <= right; i++) {
15     if(fabs(p[mid].x - p[i].x) <= d)
16         temp[k++] = p[i];
17 }
18 sort(temp,temp+k,cmp);
19 for(int i = 0; i < k; i++) {
20     for(int j = i+1; j < k && temp[j].y - temp[i].y < d;
21         ↪ j++) {
22         double d3 = Dis(temp[i],temp[j]);
23         d = min(d,d3);
24     }
25 }
26 return d;
27 }

```

### 3.1.8 最小覆盖圆

```

1 #include<cmath>
2 #include<cstdio>
3 #include<algorithm>
4 using namespace std;
5 const double eps=1e-6;
6 struct couple {
7     double x, y;
8     couple(){}
9     couple(const double &xx, const double &yy) {
10         x = xx; y = yy;
11     }
12 } a[100001];
13 int n;
14 //dis means distance, dis2 means square of it
15 struct circle {
16     double r; couple c;
17 } cir;
18 inline bool inside(const couple &x) {
19     return di2(x, cir.c) < cir.r*cir.r+eps;
20 }
21 inline void p2c(int x, int y) {
22     cir.c.x = (a[x].x+a[y].x)/2;
23     cir.c.y = (a[x].y+a[y].y)/2;
24     cir.r = dis(cir.c, a[x]);
25 }
26 inline void p3c(int i, int j, int k) {
27     couple x = a[i], y = a[j], z = a[k];
28     cir.r =
29         ↪ sqrt(di2(x,y)*di2(y,z)*di2(z,x))/fabs(x*y+y*z+z*x)/2;
30     couple t1((x-y).x, (y-z).x), t2((x-y).y, (y-z).y),
31         ↪ t3((len(x)-len(y))/2, (len(y)-len(z))/2);
32     cir.c = couple(t3*t2, t1*t3)/(t1*t2);
33 }
34 inline circle mi() {
35     sort(a + 1, a + 1 + n);
36     n = unique(a + 1, a + 1 + n) - a - 1;
37     if(n == 1) {
38         cir.c = a[1];
39         cir.r = 0;
40         return cir;
41     }
42     random_shuffle(a + 1, a + 1 + n);
43     p2c(1, 2);
44     for(int i = 3; i <= n; i++)
45         if(!inside(a[i])) {
46             p2c(1, i);
47             for(int j = 2; j < i; j++)
48                 if(!inside(a[j])) {
49                     p2c(i, j);
50                     for(int k = 1; k < j; k++)
51                         if(!inside(a[k]))
52                             p3c(i,j, k);
53                 }
54         }
55 }

```

```

53     return cir;
54 }

```

### 3.1.9 多边形内部可视

```

1 int C(const Point & P, const Point & A, const Point & Q,
2     ↪ const Point & B) {
3     Point C = GetIntersection(P, A - P, Q, Q - B);
4     return OnLine(Q, C, B);
5 }
6 int Onleft(const Point & a, const Point &b, const Point &
7     ↪ c) {
8     return dcmp(Cross(b - c, a - c)) > 0;
9 }
10 int visible(int x, int y) {
11     int P = (x + n - 1) % n, Q = (x + 1) % n;
12     Point u = p[y] - p[x], v = p[x] - p[P], w = p[x] - p[Q];
13     if (Onleft(p[Q], p[x], p[P])) {
14         return dcmp(Cross(v, u)) > 0 && dcmp(Cross(w, u)) < 0;
15     } else {
16         return !(dcmp(Cross(v, u)) < 0 && dcmp(Cross(w, u)) >
17             ↪ 0);
18     }
19 }
20 int solve(int x, int y) {
21     if (vis[x][y] == dfn) return g[x][y];
22     vis[x][y] = dfn;
23     if (x == y || y == x + 1) return g[x][y] = 1;
24     for (int i = x; i + 1 <= y; i++) {
25         if (C(p[x], p[y], p[i], p[i + 1])) return g[x][y] = 0;
26     }
27     for (int i = x + 1; i < y; i++) {
28         if (OnLine(p[x], p[i], p[y])) {
29             return g[x][y] = solve(x, i) && solve(i, y);
30         }
31     }
32     if (!visible(x, y) || !visible(y, x)) return g[x][y] =
33         ↪ 0;
34     return g[x][y] = 1;
35 }

```

## 3.2 三维

### 3.2.1 三维点类

```

1 // 三维绕轴旋转, 大拇指指向 axis 向量方向, 四指弯曲
2     ↪ 方向转 w 弧度
3 Point rotate(const Point& s, const Point& axis, DB w) {
4     DB x = axis.x, y = axis.y, z = axis.z;
5     DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
6         cosw = cos(w), sinw = sin(w);
7     DB a[4][4];
8     memset(a, 0, sizeof a);
9     a[3][3] = 1;
10     a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
11     a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
12     a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
13     a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
14     a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
15     a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
16     a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
17     a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
18     a[2][2] = ((x * x + y * y) * cosw + z * z) / s1;
19     DB ans[4] = {0, 0, 0, 0}, c[4] = {s.x, s.y, s.z, 1};
20     for (int i = 0; i < 4; ++ i)
21         for (int j = 0; j < 4; ++ j)
22             ans[i] += a[j][i] * c[j];
23     return Point(ans[0], ans[1], ans[2]);
24 }

```

## 3.2.2 凸包

```

1  __inline P cross(const P& a, const P& b) {
2      return P(
3          a.y * b.z - a.z * b.y,
4          a.z * b.x - a.x * b.z,
5          a.x * b.y - a.y * b.x
6      );
7  }
8  __inline DB mix(const P& a, const P& b, const P& c) {
9      return dot(cross(a, b), c);
10 }
11 __inline DB volume(const P& a, const P& b, const P& c,
12     ↪ const P& d) {
13     return mix(b - a, c - a, d - a);
14 }
15 struct Face {
16     int a, b, c;
17     __inline Face() {}
18     __inline Face(int _a, int _b, int _c):
19         a(_a), b(_b), c(_c) {}
20     __inline DB area() const {
21         return 0.5 * cross(p[b] - p[a], p[c] - p[a]).len();
22     }
23     __inline P normal() const {
24         return cross(p[b] - p[a], p[c] - p[a]).unit();
25     }
26     __inline DB dis(const P& p0) const {
27         return dot(normal(), p0 - p[a]);
28     }
29 };
30 std::vector<Face> face, tmp; // Should be O(n).
31 int mark[N][N], Time, n;
32 __inline void add(int v) {
33     ++ Time;
34     clear(tmp);
35     for (int i = 0; i < (int)face.size(); ++ i) {
36         int a = face[i].a, b = face[i].b, c = face[i].c;
37         if (sign(volume(p[v], p[a], p[b], p[c])) > 0) {
38             mark[a][b] = mark[b][a] = mark[a][c] =
39             mark[c][a] = mark[b][c] = mark[c][b] = Time;
40         } else {
41             tmp.push_back(face[i]);
42         }
43     }
44     clear(face); face = tmp;
45     for (int i = 0; i < (int)tmp.size(); ++ i) {
46         int a = face[i].a, b = face[i].b, c = face[i].c;
47         if (mark[a][b] == Time) face.emplace_back(v, b, a);
48         if (mark[b][c] == Time) face.emplace_back(v, c, b);
49         if (mark[c][a] == Time) face.emplace_back(v, a, c);
50         assert(face.size() < 500u);
51     }
52 }
53 void reorder() {
54     for (int i = 2; i < n; ++ i) {
55         P tmp = cross(p[i] - p[0], p[i] - p[1]);
56         if (sign(tmp.len())) {
57             std::swap(p[i], p[2]);
58             for (int j = 3; j < n; ++ j)
59                 if (sign(volume(p[0], p[1], p[2], p[j]))) {
60                     std::swap(p[j], p[3]);
61                     return;
62                 }
63     }
64 }
65 void build_convex() {
66     reorder();
67     clear(face);
68     face.emplace_back(0, 1, 2);
69     face.emplace_back(0, 2, 1);
70     for (int i = 3; i < n; ++ i)

```

```

71     add(i);
72 }

```

## 3.2.3 最小覆盖球

```

1  const int eps = 1e-8;
2  struct Tpoint {
3      double x, y, z;
4  };
5  int npoint, nouter;
6  Tpoint pt[200000], outer[4], res;
7  double radius, tmp;
8  inline double dist(Tpoint p1, Tpoint p2) {
9      double dx=p1.x-p2.x, dy=p1.y-p2.y, dz=p1.z-p2.z;
10     return ( dx*dx + dy*dy + dz*dz );
11 }
12 inline double dot(Tpoint p1, Tpoint p2) {
13     return p1.x*p2.x + p1.y*p2.y + p1.z*p2.z;
14 }
15 void ball() {
16     Tpoint q[3]; double m[3][3], sol[3], L[3], det;
17     int i, j;
18     res.x = res.y = res.z = radius = 0;
19     switch ( nouter ) {
20         case 1: res=outer[0]; break;
21         case 2:
22             res.x=(outer[0].x+outer[1].x)/2;
23             res.y=(outer[0].y+outer[1].y)/2;
24             res.z=(outer[0].z+outer[1].z)/2;
25             radius=dist(res, outer[0]);
26             break;
27         case 3:
28             for (i=0; i<2; ++i) {
29                 q[i].x=outer[i+1].x-outer[0].x;
30                 q[i].y=outer[i+1].y-outer[0].y;
31                 q[i].z=outer[i+1].z-outer[0].z;
32             }
33             for (i=0; i<2; ++i) for(j=0; j<2; ++j)
34                 m[i][j]=dot(q[i], q[j])*2;
35             for (i=0; i<2; ++i) sol[i]=dot(q[i], q[i]);
36             if (fabs(det=m[0][0]*m[1][1]-m[0][1]*m[1][0])<eps)
37                 return;
38             L[0]=(sol[0]*m[1][1]-sol[1]*m[0][1])/det;
39             L[1]=(sol[1]*m[0][0]-sol[0]*m[1][0])/det;
40             res.x=outer[0].x+q[0].x*L[0]+q[1].x*L[1];
41             res.y=outer[0].y+q[0].y*L[0]+q[1].y*L[1];
42             res.z=outer[0].z+q[0].z*L[0]+q[1].z*L[1];
43             radius=dist(res, outer[0]);
44             break;
45         case 4:
46             for (i=0; i<3; ++i) {
47                 q[i].x=outer[i+1].x-outer[0].x;
48                 q[i].y=outer[i+1].y-outer[0].y;
49                 q[i].z=outer[i+1].z-outer[0].z;
50                 sol[i]=dot(q[i], q[i]);
51             }
52             for (i=0; i<3; ++i)
53                 for(j=0; j<3; ++j) m[i][j]=dot(q[i], q[j])*2;
54             det= m[0][0]*m[1][1]*m[2][2]
55                 + m[0][1]*m[1][2]*m[2][0]
56                 + m[0][2]*m[1][0]*m[2][1]
57                 - m[0][2]*m[1][1]*m[2][0]
58                 - m[0][1]*m[1][0]*m[2][2]
59                 - m[0][0]*m[1][2]*m[2][1];
60             if ( fabs(det)<eps ) return;
61             for (j=0; j<3; ++j) {
62                 for (i=0; i<3; ++i) m[i][j]=sol[i];
63                 L[j]=( m[0][0]*m[1][1]*m[2][2]
64                     + m[0][1]*m[1][2]*m[2][0]
65                     + m[0][2]*m[1][0]*m[2][1]
66                     - m[0][2]*m[1][1]*m[2][0]
67                     - m[0][1]*m[1][0]*m[2][2]
68                     - m[0][0]*m[1][2]*m[2][1]

```

```

69         ) / det;
70         for (i=0; i<3; ++i)
71             m[i][j]=dot(q[i], q[j])*2;
72     }
73     res=outer[0];
74     for (i=0; i<3; ++i) {
75         res.x += q[i].x * L[i];
76         res.y += q[i].y * L[i];
77         res.z += q[i].z * L[i];
78     }
79     radius=dist(res, outer[0]);
80 }
81 }
82 void minball(int n) {
83     ball();
84     if ( nouter<4 )
85         for (int i=0; i<n; ++i)
86             if (dist(res, pt[i])-radius>eps) {
87                 outer[nouter]=pt[i];
88                 ++nouter;
89                 minball(i);
90                 --nouter;
91                 if (i>0) {
92                     Tpoint Tt = pt[i];
93                     memmove(&pt[1], &pt[0], sizeof(Tpoint)*i);
94                     pt[0]=Tt;
95                 }
96             }
97 }
98 void solve() {
99     for (int i=0; i<npoint; i++)
100         scanf("%lf%lf%lf", &pt[i].x, &pt[i].y, &pt[i].z);
101     random_shuffle(pt, pt + npoint);
102     radius=-1;
103     for (int i=0; i<npoint; i++){
104         if (dist(res, pt[i])-radius>eps){
105             nouter=1;
106             outer[0]=pt[i];
107             minball(i);
108         }
109     }
110     printf("%.5f\n", sqrt(radius));
111 }
112 int main(){
113     for( ; cin >> npoint && npoint; )
114         solve();
115     return 0;
116 }

```

## 4. 字符串

### 4.1 AC 自动机

```

1 int newnode() {
2     ++tot;
3     memset(ch[tot], 0, sizeof(ch[tot]));
4     fail[tot] = dep[tot] = par[tot] = 0;
5     return tot;
6 }
7 void insert(char *s, int x) {
8     if(*s == '\0') return;
9     else {
10         int &y = ch[x][*s - 'a'];
11         if(y == 0) y = newnode(), par[y] = x, dep[y] = dep[x]
12             ↳ + 1;
13         insert(s + 1, y);
14     }
15 }
16 void build() {
17     int line[maxn];
18     int f = 0, r = 0;
19     fail[root] = root;
20     for(int i = 0; i < alpha; i++) {

```

```

21         if(ch[root][i]) fail[ch[root][i]] = root, line[r++] =
22             ↳ ch[root][i];
23         else ch[root][i] = root;
24     }
25     while(f != r) {
26         int x = line[f++];
27         for(int i = 0; i < alpha; i++) {
28             if(ch[x][i]) fail[ch[x][i]] = ch[fail[x]][i],
29                 ↳ line[r++] = ch[x][i];
30             else ch[x][i] = ch[fail[x]][i];
31         }
32     }
33 }

```

### 4.2 后缀数组

```

1 const int MAXN = MAXL * 2 + 1;
2 int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN],
3     ↳ rank[MAXN], height[MAXN];
4 void calc_sa(int n) {
5     int m = alphabet, k = 1;
6     memset(c, 0, sizeof(*c) * (m + 1));
7     for (int i = 1; i <= n; ++i) c[x[i] = a[i]]++;
8     for (int i = 1; i <= m; ++i) c[i] += c[i - 1];
9     for (int i = n; i; --i) sa[c[x[i]]--] = i;
10    for (; k <= n; k <= 1) {
11        int tot = k;
12        for (int i = n - k + 1; i <= n; ++i) y[i - n + k] = i;
13        for (int i = 1; i <= n; ++i)
14            if (sa[i] > k) y[++tot] = sa[i] - k;
15        memset(c, 0, sizeof(*c) * (m + 1));
16        for (int i = 1; i <= n; ++i) c[x[i]]++;
17        for (int i = 1; i <= m; ++i) c[i] += c[i - 1];
18        for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];
19        for (int i = 1; i <= n; ++i) y[i] = x[i];
20        tot = 1; x[sa[1]] = 1;
21        for (int i = 2; i <= n; ++i) {
22            if (max(sa[i], sa[i - 1]) + k > n || y[sa[i]] !=
23                ↳ y[sa[i - 1]] || y[sa[i] + k] != y[sa[i - 1] +
24                    ↳ k]) ++tot;
25            x[sa[i]] = tot;
26        }
27        if (tot == n) break; else m = tot;
28    }
29 }
30 void calc_height(int n) {
31     for (int i = 1; i <= n; ++i) rank[sa[i]] = i;
32     for (int i = 1; i <= n; ++i) {
33         height[rank[i]] = max(0, height[rank[i - 1]] - 1);
34         int j = sa[rank[i] - 1];
35         while (max(i, j) + height[rank[i]] <= n && a[i +
36             ↳ height[rank[i]]] == a[j + height[rank[i]]])
37             ++height[rank[i]];
38     }
39 }

```

### 4.3 后缀自动机

```

1 static const int MAXL = MAXN * 2; // MAXN is original
2 static const int alphabet = 26; // sometimes need
3 static const int cnt, trans[MAXN][alphabet], par[MAXN],
4     ↳ sum[MAXN], seq[MAXN], mxl[MAXN], size[MAXN]; // mxl
5 char str[MAXN];
6 inline void init() {
7     l = strlen(str + 1); cnt = last = 1;
8     for (int i = 0; i <= l * 2; ++i) memset(trans[i], 0,
9         ↳ sizeof(trans[i]));

```

```

8  memset(par, 0, sizeof(*par) * (1 * 2 + 1));
9  memset(mx1, 0, sizeof(*mx1) * (1 * 2 + 1));
10 memset(size, 0, sizeof(*size) * (1 * 2 + 1));
11 }
12 inline void extend(int pos, int c) {
13     int p = last, np = last = ++cnt;
14     mx1[np] = mx1[p] + 1; size[np] = 1;
15     for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
16     if (!p) par[np] = 1;
17     else {
18         int q = trans[p][c];
19         if (mx1[p] + 1 == mx1[q]) par[np] = q;
20         else {
21             int nq = ++cnt;
22             mx1[nq] = mx1[p] + 1;
23             memcpy(trans[nq], trans[q], sizeof(trans[nq]));
24             par[nq] = par[q];
25             par[np] = par[q] = nq;
26             for (; trans[p][c] == q; p = par[p]) trans[p][c] =
                ↳ nq;
27         }
28     }
29 }
30 inline void buildsam() {
31     for (int i = 1; i <= l; ++i) extend(i, str[i] - 'a');
32     memset(sum, 0, sizeof(*sum) * (1 * 2 + 1));
33     for (int i = 1; i <= cnt; ++i) sum[mx1[i]]++;
34     for (int i = 1; i <= l; ++i) sum[i] += sum[i - 1];
35     for (int i = cnt; i; --i) seq[sum[mx1[i]]--] = i;
36     for (int i = cnt; i; --i) size[par[seq[i]]] +=
        ↳ size[seq[i]];
37 }

```

#### 4.4 广义后缀自动机

```

1  inline void add_node(int x, int &last) {
2      int lastnode = last;
3      if (c[lastnode][x]) {
4          int nownode = c[lastnode][x];
5          if (l[nownode] == l[lastnode] + 1) last = nownode;
6          else {
7              int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
8              for (int i = 0; i < alphabet; ++i) c[auxnode][i] =
                ↳ c[nownode][i];
9              par[auxnode] = par[nownode]; par[nownode] = auxnode;
10             for (; lastnode && c[lastnode][x] == nownode;
                ↳ lastnode = par[lastnode]) {
11                 c[lastnode][x] = auxnode;
12             }
13             last = auxnode;
14         }
15     } else {
16         int newnode = ++cnt; l[newnode] = l[lastnode] + 1;
17         for (; lastnode && !c[lastnode][x]; lastnode =
            ↳ par[lastnode]) c[lastnode][x] = newnode;
18         if (!lastnode) par[newnode] = 1;
19         else {
20             int nownode = c[lastnode][x];
21             if (l[lastnode] + 1 == l[nownode]) par[newnode] =
                ↳ nownode;
22             else {
23                 int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
24                 for (int i = 0; i < alphabet; ++i) c[auxnode][i] =
                    ↳ c[nownode][i];
25                 par[auxnode] = par[nownode]; par[nownode] =
                    ↳ par[newnode] = auxnode;
26                 for (; lastnode && c[lastnode][x] == nownode;
                    ↳ lastnode = par[lastnode]) {
27                     c[lastnode][x] = auxnode;
28                 }
29             }
30         }
31         last = newnode;

```

```

32 }
33 }

```

#### 4.5 manacher

```

1  void Manacher(std::string s, int p[]) {
2      std::string t = "$#";
3      for (int i = 0; i < s.size(); i++) {
4          t += s[i];
5          t += "#";
6      }
7      int mx = 0, id = 0;
8      for (int i = 1; i < t.size(); i++) {
9          p[i] = mx > i ? min(p[2 * id - i], mx - i) : 1;
10         while (t[i + p[i]] == t[i - p[i]]) ++p[i];
11         if (mx < i + p[i]) {
12             mx = i + p[i];
13             id = i;
14         }
15     }
16 }

```

#### 4.6 回文自动机

```

1  int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN],
    ↳ l[MAXN], s[MAXN];
2  int allocate(int len) {
3      l[nT] = len;
4      r[nT] = 0;
5      fail[nT] = 0;
6      memset(c[nT], 0, sizeof(c[nT]));
7      return nT++;
8  }
9  void init() {
10     nT = nStr = 0;
11     int newE = allocate(0);
12     int newO = allocate(-1);
13     last = newE;
14     fail[newE] = newO;
15     fail[newO] = newE;
16     s[0] = -1;
17 }
18 void add(int x) {
19     s[++nStr] = x;
20     int now = last;
21     while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now];
22     if (!c[now][x]) {
23         int newnode = allocate(l[now] + 2), &newfail =
            ↳ fail[newnode];
24         newfail = fail[now];
25         while (s[nStr - l[newfail] - 1] != s[nStr]) newfail =
            ↳ fail[newfail];
26         newfail = c[newfail][x];
27         c[now][x] = newnode;
28     }
29     last = c[now][x];
30     r[last]++;
31 }
32 void count() {
33     for (int i = nT - 1; i >= 0; i--) {
34         r[fail[i]] += r[i];
35     }
36 }

```

#### 4.7 循环串的最小表示

```

1  std::string find(std::string s) {
2      int i, j, k, l, n = s.length(); s += s;
3      for (i = 0, j = 1; j < n; j++) {
4          for (k = 0; k < n && s[i + k] == s[j + k]; k++);

```

```

5     if (k >= n) break;
6     if (s[i + k] < s[j + k]) j += k + 1; // 如果求最大表
    ↪ 示, 换成 '>'
7     else l = i + k, i = j, j = max(l, j) + 1;
8 }
9 return s.substr(i, n);
10 }

```

## 5. 数据结构

### 5.1 可并堆

```

1 int merge(int x, int y) {
2 //p[i] 结点 i 的权值, 这里是维护大根堆
3 //a[i] 在 i 的子树中, i 到右叶子结点的最远距离.
4     if(!x) return y;
5     if(!y) return x;
6     if(p[x] < p[y]) std::swap(x, y);
7     r[x] = merge(r[x], y);
8     if(r[x]) fa[r[x]] = x;
9     if(d[l[x]] < d[r[x]]) std::swap(l[x], r[x]); //调整树
    ↪ 的结构, 使其满足左偏性质
10    d[x] = d[r[x]] + 1;
11    return x;
12 }

```

### 5.2 KD-Tree

```

1 long long norm(const long long &x) {
2     return std::abs(x);
3     return x * x;
4 }
5 struct Point {
6     int x, y, id;
7     const int& operator [] (int index) const {
8         if (index == 0) {
9             return x;
10        } else {
11            return y;
12        }
13    }
14    friend long long dist(const Point &a, const Point &b)
    ↪ {
15        long long result = 0;
16        for (int i = 0; i < 2; ++i) {
17            result += norm(a[i] - b[i]);
18        }
19        return result;
20    }
21 } point[N];
22 struct Rectangle {
23     int min[2], max[2];
24     Rectangle() {
25         min[0] = min[1] = INT_MAX; // sometimes int is
    ↪ not enough
26         max[0] = max[1] = INT_MIN;
27     }
28     void add(const Point &p) {
29         for (int i = 0; i < 2; ++i) {
30             min[i] = std::min(min[i], p[i]);
31             max[i] = std::max(max[i], p[i]);
32         }
33     }
34     long long dist(const Point &p) {
35         long long result = 0;
36         for (int i = 0; i < 2; ++i) {
37             result += norm(std::min(std::max(p[i],
    ↪ min[i]), max[i]) - p[i]);
38             result += std::max(norm(max[i] - p[i]),
    ↪ norm(min[i] - p[i]));
39     }

```

```

40     return result;
41 }
42 };
43 struct Node {
44     Point separator;
45     Rectangle rectangle;
46     int child[2];
47     void reset(const Point &p) {
48         separator = p;
49         rectangle = Rectangle();
50         rectangle.add(p);
51         child[0] = child[1] = 0;
52     }
53 } tree[N << 1];
54 int size, pivot;
55 bool compare(const Point &a, const Point &b) {
56     if (a[pivot] != b[pivot]) {
57         return a[pivot] < b[pivot];
58     }
59     return a.id < b.id;
60 }
61 // 左闭右开: build(1, n + 1)
62 int build(int l, int r, int type = 1) {
63     pivot = type;
64     if (l >= r) {
65         return 0;
66     }
67     int x = ++size;
68     int mid = l + r >> 1;
69     std::nth_element(point + l, point + mid, point + r,
    ↪ compare);
70     tree[x].reset(point[mid]);
71     for (int i = l; i < r; ++i) {
72         tree[x].rectangle.add(point[i]);
73     }
74     tree[x].child[0] = build(l, mid, type ^ 1);
75     tree[x].child[1] = build(mid + 1, r, type ^ 1);
76     return x;
77 }
78 int insert(int x, const Point &p, int type = 1) {
79     pivot = type;
80     if (x == 0) {
81         tree[++size].reset(p);
82         return size;
83     }
84     tree[x].rectangle.add(p);
85     if (compare(p, tree[x].separator)) {
86         tree[x].child[0] = insert(tree[x].child[0], p,
    ↪ type ^ 1);
87     } else {
88         tree[x].child[1] = insert(tree[x].child[1], p,
    ↪ type ^ 1);
89     }
90     return x;
91 }
92 // For minimum distance
93 // For maximum: 下面递归 query 时 0, 1 换顺序; < and
    ↪ >; min and max
94 void query(int x, const Point &p, std::pair<long long,
    ↪ int> &answer, int type = 1) {
95     pivot = type;
96     if (x == 0 || tree[x].rectangle.dist(p) >
    ↪ answer.first) {
97         return;
98     }
99     answer = std::min(answer,
100         std::make_pair(dist(tree[x].separator, p),
    ↪ tree[x].separator.id));
101     if (compare(p, tree[x].separator)) {
102         query(tree[x].child[0], p, answer, type ^ 1);
103         query(tree[x].child[1], p, answer, type ^ 1);
104     } else {
105         query(tree[x].child[1], p, answer, type ^ 1);

```



```

106     query(tree[x].child[0], p, answer, type ^ 1);
107 }
108 }
109 std::priority_queue<std::pair<long long, int> > answer;
110 void query(int x, const Point &p, int k, int type = 1) {
111     pivot = type;
112     if (x == 0 || (int)answer.size() == k &&
113         ↪ tree[x].rectangle.dist(p) > answer.top().first) {
114         return;
115     }
116     answer.push(std::make_pair(dist(tree[x].seperator, p),
117         ↪ tree[x].seperator.id));
118     if ((int)answer.size() > k) {
119         answer.pop();
120     }
121     if (compare(p, tree[x].seperator)) {
122         query(tree[x].child[0], p, k, type ^ 1);
123         query(tree[x].child[1], p, k, type ^ 1);
124     } else {
125         query(tree[x].child[1], p, k, type ^ 1);
126         query(tree[x].child[0], p, k, type ^ 1);
127     }
128 }

```

### 5.3 Treap

```

1 struct Node{
2     int mn, key, size, tag;
3     bool rev;
4     Node* ch[2];
5     Node(int mn, int key, int size): mn(mn), key(key),
6         ↪ size(size), rev(0), tag(0){}
7     void downtag();
8     Node* update(){
9         mn = min(ch[0] -> mn, min(key, ch[1] -> mn));
10        size = ch[0] -> size + 1 + ch[1] -> size;
11        return this;
12    }
13 };
14 typedef pair<Node*, Node*> Pair;
15 Node *null, *root;
16 void Node::downtag(){
17     if(rev){
18         for(int i = 0; i < 2; i++){
19             if(ch[i] != null){
20                 ch[i] -> rev ^= 1;
21                 swap(ch[i] -> ch[0], ch[i] -> ch[1]);
22             }
23             rev = 0;
24         }
25     }
26     if(tag){
27         for(int i = 0; i < 2; i++){
28             if(ch[i] != null){
29                 ch[i] -> key += tag;
30                 ch[i] -> mn += tag;
31                 ch[i] -> tag += tag;
32             }
33             tag = 0;
34         }
35     }
36 }
37 int r(){
38     static int s = 3023192386;
39     return (s += (s << 3) + 1) & (~0u >> 1);
40 }
41 bool random(int x, int y){
42     return r() % (x + y) < x;
43 }
44 Node* merge(Node *p, Node *q){
45     if(p == null) return q;
46     if(q == null) return p;
47     p -> downtag();
48     q -> downtag();
49     if(random(p -> size, p -> size)){

```

```

47     p -> ch[1] = merge(p -> ch[1], q);
48     return p -> update();
49 }else{
50     q -> ch[0] = merge(p, q -> ch[0]);
51     return q -> update();
52 }
53 }
54 Pair split(Node *x, int n){
55     if(x == null) return make_pair(null, null);
56     x -> downtag();
57     if(n <= x -> ch[0] -> size){
58         Pair ret = split(x -> ch[0], n);
59         x -> ch[0] = ret.second;
60         return make_pair(ret.first, x -> update());
61     }
62     Pair ret = split(x -> ch[1], n - x -> ch[0] -> size -
63         ↪ 1);
64     x -> ch[1] = ret.first;
65     return make_pair(x -> update(), ret.second);
66 }
67 pair<Node*, Pair> get_segment(int l, int r){
68     Pair ret = split(root, l - 1);
69     return make_pair(ret.first, split(ret.second, r - l +
70         ↪ 1));
71 }
72 int main(){
73     null = new Node(INF, INF, 0);
74     null -> ch[0] = null -> ch[1] = null;
75     root = null;
76 }

```

### 5.4 Splay

```

1 template<class T>void checkmin(T &x, T y) { if(y < x) x =
2     ↪ y; }
3 struct Node {
4     Node *c[2], *fa;
5     int size, rev; LL val, add, min;
6
7     Node *init(LL v) {
8         val = min = v, add = rev = 0;
9         c[0] = c[1] = fa = NULL, size = 1;
10        return this;
11    }
12    void rvs() { std::swap(c[0], c[1]), rev ^= 1; }
13    void inc(LL x) { val += x, add += x, min += x; }
14    void pushdown() {
15        if(rev) {
16            if(c[0]) c[0]->rvs();
17            if(c[1]) c[1]->rvs();
18            rev = 0;
19        }
20        if(add) {
21            if(c[0]) c[0]->inc(add);
22            if(c[1]) c[1]->inc(add);
23            add = 0;
24        }
25    }
26    void update() {
27        min = val, size = 1;
28        if(c[0]) checkmin(min, c[0]->min), size += c[0]->size;
29        if(c[1]) checkmin(min, c[1]->min), size += c[1]->size;
30    }
31 } *root;
32 Node* newnode(LL x) {
33     static Node pool[100000], *p = pool;
34     return (++p)->init(x);
35 }
36 void setc(Node *x, int t, Node *y) {
37     x->c[t] = y;
38     if(y) y->fa = x;

```



```

39 }
40 Node *find(int k) {
41     Node *now = root;
42     while(true) {
43         now->pushdown();
44         int t = (now->c[0] ? now->c[0]->size : 0) + 1;
45         if(t == k) break;
46         if(t > k) now = now->c[0];
47         else now = now->c[1], k -= t;
48     }
49     return now;
50 }
51 void rotate(Node *x, Node* &k) {
52     Node *y = x->fa, *z = y->fa;
53     if(y != k) z->c[z->c[1] == y] = x;
54     else k = x;
55     x->fa = z;
56     int i = (y->c[1] == x);
57     setc(y, i, x->c[i ^ 1]);
58     setc(x, i ^ 1, y);
59     y->update(), x->update();
60 }
61 void spaly(Node *x, Node* &k) {
62     static Node *st[maxs], *y, *z;
63     int top = 0;
64     y = x;
65     while(y != k) st[++top] = y, y = y->fa;
66     st[++top] = y;
67     while(top) st[top]->pushdown(), top--;
68     while(x != k) {
69         y = x->fa, z = y->fa;
70         if(y != k) {
71             if((y == z->c[1]) ^ (x == y->c[1])) rotate(x, k);
72             else rotate(y, k);
73         }
74         rotate(x, k);
75     }
76 }
77 Node *subtree(int l, int r) {
78     assert((++l) <= (++r));
79     spaly(find(l - 1), root), spaly(find(r + 1),
80         ↪ root->c[1]);
81     return root->c[1]->c[0];
82 }
83 void ins(int pos, int v) {
84     pos++;
85     spaly(find(pos), root), spaly(find(pos + 1),
86         ↪ root->c[1]);
87     setc(root->c[1], 0, newnode(v));
88     root->c[1]->update(), root->update();
89 }
90 void del(int pos) {
91     pos++;
92     spaly(find(pos - 1), root), spaly(find(pos + 1),
93         ↪ root->c[1]);
94     root->c[1]->c[0] = NULL, root->c[1]->update(),
95     ↪ root->update();
96 }
97 void init() {
98     root = newnode(0);
99     setc(root, 1, newnode(0));
100     root->update();
101 }

```

## 5.5 Link cut Tree

```

1 inline void reverse(int x) {
2     tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
3 }
4 inline void rotate(int x, int k) {
5     int y = tr[x].fa, z = tr[y].fa;
6     tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;

```

```

7     tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^
8     ↪ 1];
9     tr[x].c[k ^ 1] = y; tr[y].fa = x;
10 }
11 inline void splay(int x, int w) {
12     int z = x; pushdown(x);
13     while (tr[x].fa != w) {
14         int y = tr[x].fa; z = tr[y].fa;
15         if (z == w) {
16             pushdown(z = y); pushdown(x);
17             rotate(x, tr[y].c[1] == x);
18             update(y); update(x);
19         } else {
20             pushdown(z); pushdown(y); pushdown(x);
21             int t1 = tr[y].c[1] == x, t2 = tr[z].c[1] == y;
22             if (t1 == t2) rotate(y, t2), rotate(x, t1);
23             else rotate(x, t1), rotate(x, t2);
24             update(z); update(y); update(x);
25         }
26     }
27     update(x);
28     if (x != z) par[x] = par[z], par[z] = 0;
29 }
30 inline void access(int x) {
31     for (int y = 0; x; y = x, x = par[x]) {
32         splay(x, 0);
33         if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa
34         ↪ = 0;
35         tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
36     }
37 }
38 inline void makeroot(int x) {
39     access(x); splay(x, 0); reverse(x);
40 }
41 inline void link(int x, int y) {
42     makeroot(x); par[x] = y;
43 }
44 inline void cut(int x, int y) {
45     access(x); splay(y, 0);
46     if (par[y] != x) swap(x, y), access(x), splay(y, 0);
47     par[y] = 0;
48 }
49 inline void split(int x, int y) { // x will be the root
50     ↪ of the tree
51     makeroot(y); access(x); splay(x, 0);
52 }

```

## 5.6 树上莫队

```

1 void dfs(int u) {
2     dep[u] = dep[fa[u][0]] + 1, stk.push(u);
3     for(int i = 1; i < logn; i++) fa[u][i] = fa[fa[u][i -
4     ↪ 1]][i - 1];
5     for(int i = 0; i < vec[u].size(); i++) {
6         int v = vec[u][i];
7         if(v == fa[u][0]) continue;
8         fa[v][0] = u, dfs(v), size[u] += size[v];
9         if(size[u] >= bufsize) {
10             ++bcnt;
11             while(stk.top() != u) {
12                 block[stk.top()] = bcnt;
13                 stk.pop();
14             }
15             size[u] = 0;
16         }
17     }
18     size[u]++;
19 }
20 void prework() {
21     dfs(1), ++bcnt;
22     while(!stk.empty()) block[stk.top()] = bcnt, stk.pop();
23 }
24 void rev(int u) {

```

```

24 now -= (cnt[val[u]] > 0);
25 if(used[u]) cnt[val[u]]--, used[u] = false;
26 else cnt[val[u]]++, used[u] = true;
27 now += (cnt[val[u]] > 0);
28 }
29 void move(int &x,int y,int z) {
30     int fwd = y;
31     rev(getlca(x, z)), rev(getlca(y, z));
32     while(x != y) {
33         if(dep[x] < dep[y]) std::swap(x, y);
34         rev(x), x = fa[x][0];
35     }
36     x = fwd;
37 }
38 void solve() {
39     int L = 1, R = 1;
40     std::sort(query + 1, query + m + 1);
41     rev(1);
42     for(int i = 1; i <= m; i++) {
43         int l = query[i].u, r = query[i].v;
44         move(L, l, R), move(R, r, L);
45         ans[query[i].t] = now;
46     }
47 }

```

```

7     std::vector<int> qL, qR;
8     backup.clear();
9     for(int i = 1; i <= mid; i++) {
10         Event e = event[i];
11         if(e.l <= e.r) add(e.l, e.v), add(e.r + 1, -e.v);
12         else add(1, e.v), add(e.r + 1, -e.v), add(e.l, e.v);
13     }
14     for(int i = 0; i < q.size(); i++) {
15         LL val = 0;
16         for(int j = 0; j < vec[q[i]].size(); j++) {
17             val += count(vec[q[i]][j]);
18             if(val >= p[q[i]]) break;
19         }
20         if(cnt[q[i]] + val >= p[q[i]]) qL.push_back(q[i]);
21         else cnt[q[i]] += val, qR.push_back(q[i]);
22     }
23     for(int i = 0; i < backup.size(); i++) sum[backup[i]]
        += 0;
24     solve(l, mid, qL), solve(mid + 1, r, qR);
25 }
26 }

```

## 5.7 CDQ 分治

```

1 struct Node {
2     int x, y, z, idx;
3     friend bool operator == (const Node &a,const Node &b) {
4         return a.x == b.x && a.y == b.y && a.z == b.z; }
5     friend bool operator < (const Node &a,const Node &b) {
6         return a.y < b.y; }
7 } triple[maxn];
8 bool cmpx(const Node &a,const Node &b) {
9     if(a.x != b.x) return a.x < b.x;
10    if(a.y != b.y) return a.y < b.y;
11    return a.z < b.z;
12 }
13 void solve(int l,int r) {
14     if(l == r) return;
15     int mid = (l + r) >> 1;
16     solve(l, mid);
17     static std::pair<Node,int> Lt[maxn], Rt[maxn];
18     int Ls = 0, Rs = 0, pos = 1;
19     for(int i = 1; i <= mid; i++) Lt[++Ls] =
20         triple[i];
21     for(int i = mid + 1; i <= r; i++) Rt[++Rs] =
22         triple[i];
23     std::sort(Lt + 1, Lt + Ls + 1);
24     std::sort(Rt + 1, Rt + Rs + 1);
25     backup.clear();
26     for(int i = 1; i <= Rs; i++) {
27         while(pos <= Ls && (Rt[i].first < Lt[pos].first)) {
28             insert(Lt[pos].first.z, 1);
29             pos++;
30         }
31         f[Rt[i].second] += query(Rt[i].first.z);
32     }
33     for(int i = 0; i < backup.size(); i++) pre[backup[i]] =
34         0;
35     solve(mid + 1, r);
36 }

```

## 5.8 整体二分

```

1 void solve(int l,int r,std::vector<int> q) {
2     if(l == r || q.empty()) {
3         for(int i = 0; i < q.size(); i++) ans[q[i]] = l;
4     }
5     else {
6         int mid = (l + r) >> 1;

```

# 6. 图论

## 6.1 2-SAT tarjan

```

1 template<class TAT>void checkmin(TAT &x,TAT y) { if(y < x)
2     x = y; }
3 void tarjan(int u) {
4     dfn[u] = low[u] = ++dt, flag[u] = true, stk.push(u);
5     for(int i = 0; i < vec[u].size(); i++) {
6         int v = vec[u][i];
7         if(!dfn[v]) tarjan(v), checkmin(low[u], low[v]);
8         else if(flag[v]) checkmin(low[u], dfn[v]);
9     }
10    if(low[u] == dfn[u]) {
11        ++bcnt;
12        while(stk.top() != u) block[stk.top()] = bcnt,
13            flag[stk.top()] = false, stk.pop();
14        block[u] = bcnt, flag[u] = false, stk.pop();
15    }
16 }
17 bool solve() {
18     bool ans = true;
19     for(int i = 1; i <= 2 * n; i++) if(!dfn[i]) tarjan(i);
20     for(int i = 1; i <= n; i++) if(block[2 * i] == block[2 *
21         i - 1]) { ans = false; break; }
22     return ans;
23 }

```

## 6.2 KM

```

1 struct KM {
2     // Truly O(n^3)
3     // 邻接矩阵，不能连的边设为 -INF，求最小权匹配时
4     // 边权取负，但不能连的还是 -INF，使用时先对 1
5     // -> n 调用 hungary()，再 get_ans() 求值
6     int w[N][N];
7     int lx[N], ly[N], match[N], way[N], slack[N];
8     bool used[N];
9     void init() {
10         for (int i = 1; i <= n; i++) {
11             match[i] = 0;
12             lx[i] = 0;
13             ly[i] = 0;
14             way[i] = 0;
15         }
16     }
17     void hungary(int x) {
18         match[0] = x;

```

```

17 int j0 = 0;
18 for (int j = 0; j <= n; j++) {
19     slack[j] = INF;
20     used[j] = false;
21 }
22 do {
23     used[j0] = true;
24     int i0 = match[j0], delta = INF, j1 = 0;
25     for (int j = 1; j <= n; j++) {
26         if (used[j] == false) {
27             int cur = -w[i0][j] - lx[i0] - ly[j];
28             if (cur < slack[j]) {
29                 slack[j] = cur;
30                 way[j] = j0;
31             }
32             if (slack[j] < delta) {
33                 delta = slack[j];
34                 j1 = j;
35             }
36         }
37     }
38     for (int j = 0; j <= n; j++) {
39         if (used[j]) {
40             lx[match[j]] += delta;
41             ly[j] -= delta;
42         }
43         else slack[j] -= delta;
44     }
45     j0 = j1;
46 } while (match[j0] != 0);
47 do {
48     int j1 = way[j0];
49     match[j0] = match[j1];
50     j0 = j1;
51 } while (j0);
52 }
53 int get_ans() {
54     int sum = 0;
55     for(int i = 1; i <= n; i++) {
56         if (w[match[i]][i] == -INF); // 无解
57         if (match[i] > 0) sum += w[match[i]][i];
58     }
59     return sum;
60 }
61 } km;

```

### 6.3 点双连通分量

```

1 void tarjan(int u,int f) {
2     low[u] = dfn[u] = ++df;
3     for(int i = head[u]; i ; i = edge[i].next) {
4         int v = edge[i].v;
5         if(v == f) continue;
6         if(!dfn[v]) {
7             stack[++top] = i, tarjan(v, u);
8             low[u] = std::min(low[u], low[v]);
9             if(low[v] >= dfn[u]) {
10                 ++tot;
11                 while(stack[top] != i) pbc[stack[top--]] = tot;
12                 pbc[stack[top--]] = tot;
13             }
14         }
15         else if(dfn[v] < dfn[u]) {
16             stack[++top] = i;
17             low[u] = std::min(low[u], dfn[v]);
18         }
19     }
20 }
21 void work() {
22     for(int i = 1; i <= n; i++) if(!dfn[i]) tarjan(i, 0);
23     for(int i = 1; i <= m; i++) pbc[i * 2 - 1] = pbc[i * 2]
        ⇨ = std::max(pbc[i * 2 - 1], pbc[i * 2]);

```

24 }

### 6.4 边双连通分量

```

1 struct BCC {
2     Graph *g, forest;
3     int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N],
        ⇨ top, dfs_clock;
4     // tot[] is the size of each BCC, belong[] is the BCC
        ⇨ that each node belongs to
5     pair<int, int> ori[M]; // bridge in raw_graph(raw node)
6     bool is_bridge[M];
7     __inline void init(Graph *raw_graph) {
8         g = raw_graph;
9         memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
10        memset(vis + g -> base, 0, sizeof(*vis) * g -> n);
11    }
12    void tarjan(int u, int from) {
13        dfn[u] = low[u] = ++dfs_clock; vis[u] = 1;
14        ⇨ stack[++top] = u;
15        for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
16            if ((p ^ 1) == from) continue;
17            int v = g -> v[p];
18            if (vis[v]) {
19                if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
20            } else {
21                tarjan(v, p);
22                low[u] = min(low[u], low[v]);
23                if (low[v] > dfn[u]) is_bridge[p / 2] = true;
24            }
25        }
26        if (dfn[u] != low[u]) return;
27        tot[forest.new_node()] = 0;
28        do {
29            belong[stack[top]] = forest.n;
30            vis[stack[top]] = 2;
31            tot[forest.n]++;
32            --top;
33        } while (stack[top + 1] != u);
34    }
35    void solve() {
36        forest.init(g -> base);
37        int n = g -> n;
38        for (int i = 0; i < n; ++i)
39            if (!vis[i + g -> base]) {
40                top = dfs_clock = 0;
41                tarjan(i + g -> base, -1);
42            }
43        for (int i = 0; i < g -> e / 2; ++i)
44            if (is_bridge[i]) {
45                int e = forest.e;
46                forest.bi_ins(belong[g -> v[i * 2]], belong[g ->
                    ⇨ v[i * 2 + 1]], g -> w[i * 2]);
47                ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i *
                    ⇨ 2]);
48                ori[e + 1] = make_pair(g -> v[i * 2], g -> v[i * 2
                    ⇨ + 1]);
49            }
50    } bcc;

```

### 6.5 树上路径求交

```

1 bool cmp(int x,int y) {
2     return dep[x] > dep[y];
3 }
4 std::pair<int,int> path_intersection(int a,int b,int c,int
    ⇨ d) {
5     std::vector<int> poi;
6     poi.push_back(getlca(a, c));
7     poi.push_back(getlca(a, d));
8     poi.push_back(getlca(b, c));

```

```

9   poi.push_back(getlca(b, d));
10  std::sort(poi.begin(), poi.end(), cmp);
11  return std::make_pair(poi[0], poi[1]);
12 }

```

## 6.6 最小树形图

```

1  const int MAXN, INF; // INF >= sum( W_ij )
2  int from[MAXN + 10][MAXN * 2 + 10], n, m, edge[MAXN +
   ↪ 10][MAXN * 2 + 10];
3  int sel[MAXN * 2 + 10], fa[MAXN * 2 + 10], vis[MAXN * 2 +
   ↪ 10];
4  int getfa(int x){if(x == fa[x]) return x; return fa[x] =
   ↪ getfa(fa[x]);}
5  void liuzhu(){ // 1-base: root is 1, answer = (sel[i], i)
   ↪ for i in [2..n]
6      fa[1] = 1;
7      for(int i = 2; i <= n; ++i){
8          sel[i] = 1; fa[i] = i;
9          for(int j = 1; j <= n; ++j) if(fa[j] != i)
10             if(from[j][i] = i, edge[sel[i]][i] > edge[j][i])
               ↪ sel[i] = j;
11     }
12     int limit = n;
13     while(1){
14         int prelimit = limit; memset(vis, 0, sizeof(vis));
               ↪ vis[1] = 1;
15         for(int i = 2; i <= prelimit; ++i) if(fa[i] == i &&
               ↪ !vis[i]){
16             int j = i; while(!vis[j]) vis[j] = i, j =
               ↪ getfa(sel[j]);
17             if(j == 1 || vis[j] != i) continue; vector<int> C;
               ↪ int k = j;
18             do C.push_back(k), k = getfa(sel[k]); while(k != j);
               ↪ ++limit;
19             for(int i = 1; i <= n; ++i){
20                 edge[i][limit] = INF, from[i][limit] = limit;
21             }
22             fa[limit] = vis[limit] = limit;
23             for(int i = 0; i < int(C.size()); ++i){
24                 int x = C[i], fa[x] = limit;
25                 for(int j = 1; j <= n; ++j)
26                     if(edge[j][x] != INF && edge[j][limit] >
               ↪ edge[j][x] - edge[sel[x]][x]){
27                         edge[j][limit] = edge[j][x] - edge[sel[x]][x];
28                         from[j][limit] = x;
29                     }
30             }
31             for(int j=1; j<=n; ++j) if(getfa(j)==limit)
               ↪ edge[j][limit] = INF;
32             sel[limit] = 1;
33             for(int j = 1; j <= n; ++j)
34                 if(edge[sel[limit]][limit] > edge[j][limit])
               ↪ sel[limit] = j;
35             if(prelimit == limit) break;
36         }
37     }
38     for(int i = limit; i > 1; --i) sel[from[sel[i]][i]] =
   ↪ sel[i];
39 }
40 }

```

## 6.7 带花树

```

1  vector<int> link[maxn];
2  int n, match[maxn], Queue[maxn], head, tail;
3  int pred[maxn], base[maxn], start, finish, newbase;
4  bool InQueue[maxn], InBlossom[maxn];
5  void push(int u){ Queue[tail++] = u; InQueue[u] = true; }
6  int pop(){ return Queue[head++]; }
7  int FindCommonAncestor(int u, int v){
8      bool InPath[maxn];
9      for(int i=0; i<n; ++i) InPath[i] = 0;

```

```

10  while(true){ u=base[u]; InPath[u]=true; if(u==start)
   ↪ break; u=pred[match[u]]; }
11  while(true){ v=base[v]; if(InPath[v])
   ↪ break; v=pred[match[v]]; }
12  return v;
13 }
14 void ResetTrace(int u){
15     int v;
16     while(base[u] != newbase){
17         v=match[u];
18         InBlossom[base[u]] = InBlossom[base[v]] = true;
19         u=pred[v];
20         if(base[u] != newbase) pred[u] = v;
21     }
22 }
23 void BlossomContract(int u, int v){
24     newbase = FindCommonAncestor(u, v);
25     for (int i=0; i<n; ++i)
26         InBlossom[i] = 0;
27     ResetTrace(u); ResetTrace(v);
28     if(base[u] != newbase) pred[u] = v;
29     if(base[v] != newbase) pred[v] = u;
30     for(int i=0; i<n; ++i)
31         if(InBlossom[base[i]]){
32             base[i] = newbase;
33             if(!InQueue[i]) push(i);
34         }
35 }
36 bool FindAugmentingPath(int u){
37     bool found = false;
38     for(int i=0; i<n; ++i) pred[i] = -1, base[i] = i;
39     for (int i=0; i<n; ++i) InQueue[i] = 0;
40     start = u; finish = -1; head = tail = 0; push(start);
41     while(head < tail){
42         int u = pop();
43         for(int i=link[u].size()-1; i>=0; i--){
44             int v = link[u][i];
45             if(base[u] != base[v] && match[u] != v)
46                 if(v == start || (match[v] >= 0 && pred[match[v]] >= 0))
47                     BlossomContract(u, v);
48             else if(pred[v] == -1){
49                 pred[v] = u;
50                 if(match[v] >= 0) push(match[v]);
51                 else{ finish = v; return true; }
52             }
53         }
54     }
55     return found;
56 }
57 void AugmentPath(){
58     int u = finish, v, w;
59     while(u >= 0){
60         ↪ v = pred[u]; w = match[v]; match[v] = u; match[u] = v; u = w; }
61 }
62 void FindMaxMatching(){
63     for(int i=0; i<n; ++i) match[i] = -1;
64     for(int i=0; i<n; ++i) if(match[i] == -1)
   ↪ if(FindAugmentingPath(i)) AugmentPath();

```

## 6.8 支配树

```

1  vector<int> prec[N], succ[N];
2  vector<int> ord;
3  int stamp, vis[N];
4  int num[N];
5  int fa[N];
6  void dfs(int u) {
7      vis[u] = stamp;
8      num[u] = ord.size();
9      ord.push_back(u);
10     for (int i = 0; i < (int)succ[u].size(); ++i) {

```

```

11     int v = succ[u][i];
12     if (vis[v] != stamp) {
13         fa[v] = u;
14         dfs(v);
15     }
16 }
17 }
18 int fs[N], mins[N], dom[N], sem[N];
19 int find(int u) {
20     if (u != fs[u]) {
21         int v = fs[u];
22         fs[u] = find(fs[u]);
23         if (mins[v] != -1 && num[sem[mins[v]]] <
            ↪ num[sem[mins[u]]]) {
24             mins[u] = mins[v];
25         }
26     }
27     return fs[u];
28 }
29 void merge(int u, int v) { fs[u] = v; }
30 vector<int> buf[N];
31 int buf2[N];
32 void mark(int source) {
33     ord.clear();
34     ++stamp;
35     dfs(source);
36     for (int i = 0; i < (int)ord.size(); ++i) {
37         int u = ord[i];
38         fs[u] = u, mins[u] = -1, buf2[u] = -1;
39     }
40     for (int i = (int)ord.size() - 1; i > 0; --i) {
41         int u = ord[i], p = fa[u];
42         sem[u] = p;
43         for (int j = 0; j < (int)prec[u].size(); ++j) {
44             int v = prec[u][j];
45             if (use[v] != stamp) continue;
46             if (num[v] > num[u]) {
47                 find(v); v = sem[mins[v]];
48             }
49             if (num[v] < num[sem[u]]) {
50                 sem[u] = v;
51             }
52         }
53         buf[sem[u]].push_back(u);
54         mins[u] = u;
55         merge(u, p);
56         while (buf[p].size()) {
57             int v = buf[p].back();
58             buf[p].pop_back();
59             find(v);
60             if (sem[v] == sem[mins[v]]) {
61                 dom[v] = sem[v];
62             } else {
63                 buf2[v] = mins[v];
64             }
65         }
66     }
67     dom[ord[0]] = ord[0];
68     for (int i = 0; i < (int)ord.size(); ++i) {
69         int u = ord[i];
70         if (~buf2[u]) {
71             dom[u] = dom[buf2[u]];
72         }
73     }
74 }

```

## 6.9 随机最大权匹配

```

1 bool dfs(int i) {
2     path[len++] = i;
3     if (v[i]) return true;
4     v[i] = true;
5     for (int j = 0; j < k; ++j) {

```

```

6         if (i != j && match[i] != j && !v[j]) {
7             int kok = match[j];
8             if (d[kok] < d[i] + w[i][j] - w[j][kok]) {
9                 d[kok] = d[i] + w[i][j] - w[j][kok];
10                if (dfs(kok)) return true;
11            }
12        }
13    }
14    --len; v[i] = false;
15    return false;
16 }
17 void solve() {
18     for (int i = 0; i < k; ++i) p[i] = i, match[i] = i ^
        ↪ 1;
19     int cnt = 0;
20     for (;;) {
21         len = 0;
22         bool flag = false;
23         memset(d, 0, sizeof(d));
24         memset(v, 0, sizeof(v));
25         for (int i = 0; i < k; ++i) {
26             if (dfs(p[i])) {
27                 flag = true;
28                 int t = match[path[len - 1]], j = len - 2;
29                 while (path[j] != path[len - 1]) {
30                     match[t] = path[j];
31                     swap(t, match[path[j]]);
32                     --j;
33                 }
34                 match[t] = path[j];
35                 match[path[j]] = t;
36                 break;
37             }
38         }
39         if (!flag) {
40             if (++cnt >= 3) break;
41             random_shuffle(p, p+k);
42         }
43     }
44 }

```

## 6.10 无向图最小割

```

1 int cost[maxn][maxn], seq[maxn], len[maxn], n, m, pop, ans;
2 bool used[maxn];
3 void Init() {
4     int i, j, a, b, c;
5     for (i = 0; i < n; ++i) for (j = 0; j < n; ++j) cost[i][j] = 0;
6     for (i = 0; i < m; ++i) {
7         scanf("%d %d %d", &a, &b, &c); cost[a][b] += c;
            ↪ cost[b][a] += c;
8     }
9     pop = n; for (i = 0; i < n; ++i) seq[i] = i;
10 }
11 void Work() {
12     ans = inf; int i, j, k, l, mm, sum, pk;
13     while (pop > 1) {
14         for (i = 1; i < pop; ++i) used[seq[i]] = 0; used[seq[0]] = 1;
15         for (i = 1; i < pop; ++i) len[seq[i]] = cost[seq[0]][seq[i]];
16         pk = 0; mm = -inf; k = -1;
17         for (i = 1; i < pop; ++i) if (len[seq[i]] > mm) {
            ↪ mm = len[seq[i]]; k = i; }
18         for (i = 1; i < pop; ++i) {
19             used[seq[i]] = 1;
20             if (i == pop - 2) pk = k;
21             if (i == pop - 1) break;
22             mm = -inf;
23             for (j = 1; j < pop; ++j) if (!used[seq[j]])
24                 if ((len[seq[j]] + cost[seq[i]][seq[j]]) > mm)
25                     mm = len[seq[j]], k = j;
26         }
27         sum = 0;

```

```

28     for(i=0;i<pop;i++) if(i != k)
        ↳ sum+=cost[seq[k]][seq[i]];
29     ans=min(ans,sum);
30     for(i=0;i<pop;i++)
31         cost[seq[k]][seq[i]]=cost[seq[i]][seq[k]]+=cost[seq[pk]][seq[i]];
32     seq[pk]=seq[--pop];
33 }
34 printf("%d\n",ans);
35 }

```

## 6.11 最大团搜索

```

1  const int N = 1000 + 7;
2  vector<vector<bool>> adj;
3  class MaxClique {
4      const vector<vector<bool>> adj;
5      const int n;
6      vector<int> result, cur_res;
7      vector<vector<int>> color_set;
8      const double t_limit; // MAGIC
9      int para, level;
10     vector<pair<int, int>> steps;
11 public:
12     class Vertex {
13     public:
14         int i, d;
15         Vertex(int i, int d = 0) : i(i), d(d) {}
16     };
17     void reorder(vector<Vertex> &p) {
18         for (auto &u : p) {
19             u.d = 0;
20             for (auto v : p) u.d += adj[v.i][u.i];
21         }
22         sort(p.begin(), p.end(), [&](const Vertex &a,
        ↳ const Vertex &b) { return a.d > b.d; } );
23     }
24     void init_color(vector<Vertex> &p) {
25         int maxd = p[0].d;
26         for (int i = 0; i < p.size(); i++) p[i].d = min(i,
        ↳ maxd) + 1;
27     }
28     bool bridge(const vector<int> &s, int x) {
29         for (auto v : s) if (adj[v][x]) return true;
30         return false;
31     }
32     void color_sort(vector<Vertex> &cur) {
33         int totc = 0, ptr = 0, mink =
        ↳ max((int)result.size() - (int)cur_res.size(),
        ↳ 0);
34         for (int i = 0; i < cur.size(); i++) {
35             int x = cur[i].i, k = 0;
36             while (k < totc && bridge(color_set[k], x))
        ↳ k++;
37             if (k == totc) color_set[totc++].clear();
38             color_set[k].push_back(x);
39             if (k < mink) cur[ptr++].i = x;
40         }
41         if (ptr) cur[ptr - 1].d = 0;
42         for (int i = mink; i < totc; i++) {
43             for (auto v : color_set[i]) {
44                 cur[ptr++] = Vertex(v, i + 1);
45             }
46         }
47     }
48     void expand(vector<Vertex> &cur) {
49         steps[level].second = steps[level].second -
        ↳ steps[level].first + steps[level - 1].first;
50         steps[level].first = steps[level - 1].second;
51         while (cur.size()) {
52             if (cur_res.size() + cur.back().d <=
        ↳ result.size()) return ;
53             int x = cur.back().i;
54             cur_res.push_back(x); cur.pop_back();

```

```

55     vector<Vertex> remain;
56     for (auto v : cur) {
57         if (adj[v.i][x]) remain.push_back(v.i);
58     }
59     if (remain.size() == 0) {
60         if (cur_res.size() > result.size()) result
        ↳ = cur_res;
61     } else {
62         if (1. * steps[level].second / ++para < t_limit)
        ↳ reorder(remain);
63         color_sort(remain);
64         steps[level++].second++;
65         expand(remain);
66         level--;
67     }
68     cur_res.pop_back();
69 }
70 }
71 public:
72     MaxClique(const vector<vector<bool>> &_adj, int n,
        ↳ double tt = 0.025) : adj(_adj), n(n), t_limit(tt)
        ↳ {
73         result.clear();
74         cur_res.clear();
75         color_set.resize(n);
76         steps.resize(n + 1);
77         fill(steps.begin(), steps.end(), make_pair(0, 0));
78         level = 1;
79         para = 0;
80     }
81     vector<int> solve() {
82         vector<Vertex> p;
83         for (int i = 0; i < n; i++)
        ↳ p.push_back(Vertex(i));
84         reorder(p);
85         init_color(p);
86         expand(p);
87         return result;
88     }
89 };

```

## 6.12 弦图判定

```

1  // n 必须是 2 的次幂
2  void fft(Complex a[], int n, int f) {
3      for (int i = 0; i < n; ++i)
4          if (R[i] < i) swap(a[i], a[R[i]]);
5      for (int i = 1, h = 0; i < n; i <= 1, h++) {
6          Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7          Complex w = Complex(1, 0);
8          for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9          for (int p = i < 1, j = 0; j < n; j += p) {
10             for (int k = 0; k < i; ++k) {
11                 Complex x = a[j + k], y = a[j + k + i] * tmp[k];
12                 a[j + k] = x + y; a[j + k + i] = x - y;
13             }
14         }
15     }
16 }

```

## 6.13 斯坦纳树

```

1  void SPFA(int *dist) {
2      static int line[maxn + 5];
3      static bool hash[maxn + 5];
4      int f = 0, r = 0;
5      for(int i = 1; i <= N; i++) if(dist[i] < inf) line[r]
        ↳ = i, hash[i] = true, r = (r + 1) % (N + 1);
6      while(f != r) {
7          int t = line[f]; hash[t] = false, f = (f + 1) % (N
        ↳ + 1);

```



```

8     for(int i = head[t]; i ; i = edge[i].next) {
9         int v = edge[i].v, dt = dist[t] + edge[i].w;
10        if(dt < dist[v]) {
11            dist[v] = dt;
12            if(!hash[v]) {
13                if(dist[v] < dist[line[f]]) f = (f +
14                    ↪ N) % (N + 1), line[f] = v;
15                else line[r] = v, r = (r + 1) % (N +
16                    ↪ 1);
17                hash[v] = true;
18            }
19        }
20    }
21    void solve()
22    {
23        for(int i = 1; i <= S; i++) {
24            for(int j = 1; j <= N; j++)
25                for(int k = (i - 1) & i; k ; k = (k - 1) & i)
26                    ↪ G[i][j] = std::min(G[i][j], G[k][j] + G[k
27                    ↪ i][j]);
28            SPFA(G[i]);
29        }
30    }

```

## 6.14 虚树

```

1 bool cmp(const int lhs,const int rhs) { return dfn[lhs] <
    ↪ dfn[rhs]; }
2 void build()
3 {
4     int top = 0;
5     std::sort(h + 1, h + 1 + m, cmp);
6     for (int i = 1; i <= m; i++) {
7         if (!top) father[st[++top] = h[i]] = 0;
8         else {
9             int p = h[i], lca = LCA(h[i],st[top]);
10            while(d[st[top]] > d[lca]) {
11                if (d[st[top - 1]] <= d[lca]) father[st[top]]
12                    ↪ = lca;
13                top--;
14            }
15            if (st[top] != lca) t[st[top]] = lca, father[lca] =
16                ↪ st[top], st[top] = lca;
17            father[p] = lca, st[top] = p;
18        }
19    }
20 }

```

## 6.15 点分治

```

1 template<class TAT>void checkmax(TAT &x,TAT y) { if(x < y)
    ↪ x = y; }
2 template<class TAT>void checkmin(TAT &x,TAT y) { if(y < x)
    ↪ x = y; }
3 void getsize(int u,int fa) {
4     size[u] = 1, smax[u] = 0;
5     for(int i = 0; i < G[u].size(); i++) {
6         int v = G[u][i];
7         if(v == fa || ban[v]) continue;
8         getsize(v, u);
9         size[u] += size[v];
10        checkmax(smax[u], size[v]);
11    }
12 }
13 int getroot(int u,int ts,int fa) {
14     int res = u;
15     checkmax(smax[u], ts - size[u]);
16     for(int i = 0; i < G[u].size(); i++) {
17         int v = G[u][i];
18         if(v == fa || ban[v]) continue;

```

```

19     int w = getroot(v, ts, u);
20     if(smax[w] < smax[res]) res = w;
21 }
22 return res;
23 }
24 void solve() {
25     static int line[maxn];
26     static std::vector<int> vec;
27     int f = 0, r = 0; line[r++] = 1;
28     while(f != r) {
29         int u = line[f++];
30         getsize(u, 0), u = getroot(u, size[u], 0);
31         ban[u] = true, vec.clear();
32         for(int i = 0; i < G[u].size(); i++)
33             if(!ban[G[u][i]]) vec.push_back(G[u][i]);
34         /*do something you like...*/
35         for(int i = 0; i < vec.size(); i++) line[r++] =
36             ↪ vec[i];
37     }

```

## 6.16 最小割最大流

```

1 bool BFS() {
2     for(int i = 1; i <= ind; i++) dep[i] = 0;
3     dep[S] = 1, line.push(S);
4     while(!line.empty()) {
5         int now = line.front(); line.pop();
6         for(int i = head[now], p; i ; i = edge[i].next)
7             if(edge[i].cap && !dep[p = edge[i].v])
8                 dep[p] = dep[now] + 1, line.push(p);
9     }
10    if(dep[T]) {
11        for(int i = 1; i <= ind; i++) cur[i] = head[i];
12        return true;
13    }
14    else return false;
15 }
16 int DFS(int a,int flow) {
17     if(a == T) return flow;
18     int ret = 0;
19     for(int &i = cur[a], p; i ; i = edge[i].next)
20         if(dep[p = edge[i].v] == dep[a] + 1 &&
21             ↪ edge[i].cap) {
22             int ff = DFS(p, std::min(flow, edge[i].cap));
23             flow -= ff, edge[i].cap -= ff;
24             ret += ff, edge[i ^ 1].cap += ff;
25             if(!flow) break;
26         }
27     return ret;
28 }
29 int solve() {
30     int totflow = 0;
31     while(BFS()) {
32         totflow += DFS(S, INF);
33     }
34     return totflow;
35 }

```

## 6.17 最小费用流

```

1 bool SPFA() {
2     static int line[maxv];
3     static bool hash[maxv];
4     register int f = 0, r = 0;
5     for(int i = 1; i <= ind; i++) dist[i] = inf, from[i] =
6         ↪ 0;
7     dist[S] = 0, line[r] = S, r = (r + 1) % maxv, hash[S]
8         ↪ = true;
9     while(f != r) {
10        int x = line[f];

```



```

9      line[f] = 0, f = (f + 1) % maxv, hash[x] = false;
10      for(int i = head[x]; i; i = edge[i].next)
11          if(edge[i].cap) {
12              int v = edge[i].v, w = dist[x] +
                  ↳ edge[i].cost;
13              if(w < dist[v]) {
14                  dist[v] = w, from[v] = i;
15                  if(!hash[v]) {
16                      if(f != r && dist[v] <=
                          ↳ dist[line[f]]) f = (f - 1 +
                          ↳ maxv) % maxv, line[f] = v;
17                      else line[r] = v, r = (r + 1) %
                          ↳ maxv;
18                      hash[v] = true;
19                  }
20              }
21          }
22      }
23      return from[T];
24  }
25  int back(int x,int flow) {
26      if(from[x]) {
27          flow = back(edge[from[x] ^ 1].v, std::min(flow,
                  ↳ edge[from[x]].cap));
28          edge[from[x]].cap -= flow, edge[from[x] ^ 1].cap +=
                  ↳ flow;
29      }
30      return flow;
31  }
32  int solve() {
33      int mincost = 0, maxflow = 0;
34      while(SPFA()) {
35          int flow = back(T, inf);
36          mincost += dist[T] * flow, maxflow += flow;
37      }
38      return mincost;
39  }

```

## 6.18 zkw 费用流

```

1  int S, T, totFlow, totCost;
2  int dis[N], slack[N], visit[N];
3  int modlable () {
4      int delta = INF;
5      for (int i = 1; i <= T; i++) {
6          if (!visit[i] && slack[i] < delta) delta =
                  ↳ slack[i];
7          slack[i] = INF;
8      }
9      if (delta == INF) return 1;
10     for (int i = 1; i <= T; i++)
11         if (visit[i]) dis[i] += delta;
12     return 0;
13 }
14 int dfs (int x, int flow) {
15     if (x == T) {
16         totFlow += flow;
17         totCost += flow * (dis[S] - dis[T]);
18         return flow;
19     }
20     visit[x] = 1;
21     int left = flow;
22     for (int i = e.last[x]; ~i; i = e.succ[i])
23         if (e.cap[i] > 0 && !visit[e.other[i]]) {
24             int y = e.other[i];
25             if (dis[y] + e.cost[i] == dis[x]) {
26                 int delta = dfs (y, min (left, e.cap[i]));
27                 e.cap[i] -= delta;
28                 e.cap[i ^ 1] += delta;
29                 left -= delta;
30                 if (!left) { visit[x] = 0; return flow; }
31             } else {

```

```

32         slack[y] = min (slack[y], dis[x] +
                  ↳ e.cost[i] - dis[y]);
33     }
34 }
35 return flow - left;
36 }
37 pair <int, int> minCost () {
38     totFlow = 0; totCost = 0;
39     fill (dis + 1, dis + T + 1, 0);
40     do {
41         do {
42             fill (visit + 1, visit + T + 1, 0);
43         } while (dfs (S, INF));
44     } while (!modlable ());
45     return make_pair (totFlow, totCost);
46 }

```

## 6.19 最小割树

```

1  int
    ↳ cnt,n,m,dis[N],last[N],a[N],tmp[N],ans[N][N],s,t,mark[N];
2  struct edge{int to,c,next;}e[N*200];
3  queue <int> q;
4  void addedge(int u,int v,int c) {
5      e[++cnt].to=v;e[cnt].c=c;
6      e[cnt].next=last[u];last[u]=cnt;
7      e[++cnt].to=u;e[cnt].c=c;
8      e[cnt].next=last[v];last[v]=cnt;
9  }
10 bool bfs() {
11     memset(dis,0,sizeof(dis));
12     dis[s]=2;
13     while (!q.empty()) q.pop();
14     q.push(s);
15     while (!q.empty()) {
16         int u=q.front();
17         q.pop();
18         for (int i=last[u];i;i=e[i].next)
19             if (e[i].c&&!dis[e[i].to]) {
20                 dis[e[i].to]=dis[u]+1;
21                 if (e[i].to==t) return 1;
22                 q.push(e[i].to);
23             }
24     }
25     return 0;
26 }
27 int dfs(int x,int maxf) {
28     if (x==t||!maxf) return maxf;
29     int ret=0;
30     for (int i=last[x];i;i=e[i].next)
31         if (e[i].c&&dis[e[i].to]==dis[x]+1) {
32             int f=dfs(e[i].to,min(e[i].c,maxf-ret));
33             e[i].c-=f;
34             e[i^1].c+=f;
35             ret+=f;
36             if (ret==maxf) break;
37         }
38     if (!ret) dis[x]=0;
39     return ret;
40 }
41 void dfs(int x) {
42     mark[x]=1;
43     for (int i=last[x];i;i=e[i].next)
44         if (e[i].c&&!mark[e[i].to]) dfs(e[i].to);
45 }
46 void solve(int l,int r) {
47     if (l==r) return;
48     s=a[l];t=a[r];
49     for (int i=2;i<=cnt;i+=2)
50         e[i].c=e[i^1].c=(e[i].c+e[i^1].c)/2;
51     int flow=0;
52     while (bfs()) flow+=dfs(s,inf);
53     memset(mark,0,sizeof(mark));

```

```

54     dfs(s);
55     for (int i=1;i<=n;i++)
56         if (mark[i])
57             for (int j=1;j<=n;j++)
58                 if (!mark[j])
59                     ans[i][j]=ans[j][i]=min(ans[i][j],flow);
60     int i=1,j=r;
61     for (int k=1;k<=r;k++)
62         if (mark[a[k]]) tmp[i++]=a[k];
63         else tmp[j--]=a[k];
64     for (int k=1;k<=r;k++)
65         a[k]=tmp[k];
66     solve(l,i-1);
67     solve(j+1,r);
68 }

```

## 6.20 上下界网络流建图

$B(u, v)$  表示边  $(u, v)$  流量的下界,  $C(u, v)$  表示边  $(u, v)$  流量的上界,  $F(u, v)$  表示边  $(u, v)$  的流量。设  $G(u, v) = F(u, v) - B(u, v)$ , 显然有

$$0 \leq G(u, v) \leq C(u, v) - B(u, v)$$

### 6.20.1 无源汇的上下界可行流

建立超级源点  $S^*$  和超级汇点  $T^*$ , 对于原图每条边  $(u, v)$  在新网络中连如下三条边:  $S^* \rightarrow v$ , 容量为  $B(u, v)$ ;  $u \rightarrow T^*$ , 容量为  $B(u, v)$ ;  $u \rightarrow v$ , 容量为  $C(u, v) - B(u, v)$ 。最后求新网络的最大流, 判断从超级源点  $S^*$  出发的边是否都满流即可, 边  $(u, v)$  的最终解中的实际流量为  $G(u, v) + B(u, v)$ 。

### 6.20.2 有源汇的上下界可行流

从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边。按照无源汇的上下界可行流一样做即可, 流量即为  $T \rightarrow S$  边上的流量。

### 6.20.3 有源汇的上下界最大流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $\infty$ , 下届为  $x$  的边。  $x$  满足二分性质, 找到最大的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最大流。
2. 从汇点  $T$  到源点  $S$  连一条上界为  $\infty$ , 下界为 0 的边, 变成无源汇的网络。按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  和超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 再将汇点  $T$  到源点  $S$  的这条边拆掉, 求一次  $S \rightarrow T$  的最大流即可。

### 6.20.4 有源汇的上下界最小流

1. 在有源汇的上下界可行流中, 从汇点  $T$  到源点  $S$  的边改为连一条上界为  $x$ , 下界为 0 的边。  $x$  满足二分性质, 找到最小的  $x$  使得新网络存在无源汇的上下界可行流即为原图的最小流。
2. 按照无源汇的上下界可行流的方法, 建立超级源点  $S^*$  与超级汇点  $T^*$ , 求一遍  $S^* \rightarrow T^*$  的最大流, 但是注意这一次不加上汇点  $T$  到源点  $S$  的这条边, 即不使之改为无源汇的网络去求解。求完后, 再加上那条汇点  $T$  到源点  $S$  上界  $\infty$  的边。因为这条边下界为 0, 所以  $S^*, T^*$  无影响, 再直接求一次  $S^* \rightarrow T^*$  的最大流。若超级源点  $S^*$  出发的边全部满流, 则  $T \rightarrow S$  边上的流量即为原图的最小流, 否则无解。

## 7. 其他

### 7.1 Dancing Links

#### 7.1.1 精确覆盖

```

1 #pragma comment(linker, "/STACK:1024000000,1024000000")
2 int head,sz;

```

```

3 int U[maxn],D[maxn],L[maxn],R[maxn];
4 int H[maxn],ROW[maxn],C[maxn],S[maxn],O[maxn];
5 void remove(int c) {
6     L[R[c]]=L[c];
7     R[L[c]]=R[c];
8     for(int i=D[c]; i!=c; i=D[i])
9         for(int j=R[i]; j!=i; j=R[j]) {
10             U[D[j]]=U[j];
11             D[U[j]]=D[j];
12             --S[C[j]];
13         }
14 }
15 void resume(int c) {
16     for(int i=U[c]; i!=c; i=U[i]) {
17         for(int j=L[i]; j!=i; j=L[j]) {
18             ++S[C[j]];
19             U[D[j]]=j;
20             D[U[j]]=j;
21         }
22     }
23     L[R[c]]=c;
24     R[L[c]]=c;
25 }
26 void init(int m) {
27     head=0; //头指针为 0
28     for(int i=0; i<=m; i++) {
29         U[i]=i; D[i]=i; L[i]=i-1; R[i]=i+1; S[i]=0;
30     }
31     R[m]=0; L[0]=m; S[0]=INF+1; sz=m+1;
32     memset(H,0,sizeof(H));
33 }
34 void insert(int i, int j) {
35     if(H[i]) {
36         L[sz] = L[H[i]]; R[sz] = H[i];
37         L[R[sz]] = sz; R[L[sz]] = sz;
38     } else {
39         L[sz] = sz; R[sz] = sz;
40         H[i] = sz;
41     }
42     U[sz] = U[j]; D[sz] = j;
43     U[D[sz]] = sz; D[U[sz]] = sz;
44     C[sz] = j; ROW[sz] = i;
45     ++S[j]; ++sz;
46 }
47 bool dfs(int k, int len) {
48     if(R[head]==head) return true;
49     int s=INF, c;
50     for (int t=R[head]; t!=head; t=R[t])
51         if (S[t]<s) s=S[t], c=t;
52     remove(c);
53     for(int i=D[c]; i!=c; i=D[i]) {
54         O[k]=ROW[i];
55         for(int j=R[i]; j!=i; j=R[j])
56             remove(C[j]);
57         if(dfs(k+1,len)) return true;
58         for(int j=L[i]; j!=i; j=L[j])
59             resume(C[j]);
60     }
61     resume(c);
62     return false;
63 }

```

#### 7.1.2 重复覆盖

```

1 int h() {
2     int i,j,k,count=0;
3     bool visit[N];
4     memset(visit,0,sizeof(visit));
5     for(i=R[0]; i!=R[i]) {
6         if(visit[i]) continue;
7         count++;

```

```

8     visit[i]=1;
9     for(j=D[i];j!=i;j=D[j]) {
10         for(k=R[j];k!=j;k=R[k])
11             visit[C[k]]=1;
12     }
13 }
14 return count;
15 }
16 void Dance(int k) {
17     int i,j,c,Min,ans;
18     ans=h();
19     if(k+ans>K || k+ans>=ak) return;
20     if(!R[0]) {
21         if(k<ak) ak=k;
22         return;
23     }
24     for(Min=N,i=R[0];i=R[i])
25         if(S[i]<Min) Min=S[i],c=i;
26     for(i=D[c];i!=c;i=D[i]) {
27         remove(i);
28         for(j=R[i];j!=i;j=R[j])
29             remove(j);
30         Dance(k+1);
31         for(j=L[i];j!=i;j=L[j])
32             resume(j);
33         resume(i);
34     }
35     return;
36 }

```

## 7.2 蔡勒公式

```

1 int zeller(int y,int m,int d) {
2     if (m<=2) y--,m+=12; int c=y/100; y%=100;
3     int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
4     if (w<0) w+=7; return(w);
5 }

```

## 7.3 五边形数定理

$$p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k-1} p(n - \frac{k(3k-1)}{2})$$

```

1 LL dp[N],fi[N];
2 LL five(LL x){ return (3*x*x-x)/2; }
3 void wbxs(){
4     dp[0]=1;
5     int t=1000; //其实可以等于 sqrt(N)
6     for(int i=-t;i<=t;++i)
7         fi[i+t]=five(i); //Q
8     for(int i=1;i<=100000;++i){
9         int flag=1;
10        for(int j=1;++j){
11            LL a=fi[j+t],b=fi[-j+t];
12            if(a>i && b>i) break;
13            if(a<=i) dp[i]=(dp[i]+dp[i-a]*flag+MOD)%MOD;
14                //p
15            if(b<=i) dp[i]=(dp[i]+dp[i-b]*flag+MOD)%MOD;
16            flag*=-1;
17        }
18    }

```

## 7.4 凸包闵可夫斯基和

```

1 // cv[0..1] 为两个顺时针凸包，其中起点等于终点，求
2 // 出的闵可夫斯基和不一定严格凸包
3 int i[2] = {0, 0}, len[2] = {(int)cv[0].size() - 1,
4 // (int)cv[1].size() - 1};
5 vector<P> mnk;
6 mnk.push_back(cv[0][0] + cv[1][0]);
7 do {

```

```

6     int d((cv[0][i[0] + 1] - cv[0][i[0]]) * (cv[1][i[1] + 1]
7 // - cv[1][i[1]]) >= 0);
8     mnk.push_back(cv[d][i[d] + 1] - cv[d][i[d]] +
9 // - mnk.back());
10    i[d] = (i[d] + 1) % len[d];
11 } while(i[0] || i[1]);

```

# 8. 技巧

## 8.1 STL 归还空间

```

1 template <typename T>
2 __inline void clear(T& container) {
3     container.clear(); // 或者删除了一堆元素
4     T(container).swap(container);
5 }

```

## 8.2 大整数取模

```

1 // 需要保证 x 和 y 非负
2 long long mult(long long x, long long y, long long MODN) {
3     long long t = (x * y - (long long)((long double)x / MODN
4 // * y + 1e-3) * MODN) % MODN;
5     return t < 0 ? t + MODN : t;
6 }

```

## 8.3 读入优化

```

1 // getchar() 读入优化 << 关同步 cin << 此优化
2 // 用 isdigit() 会小幅变慢
3 // 返回 false 表示读到文件尾
4 namespace Reader {
5     const int L = (1 << 15) + 5;
6     char buffer[L], *S, *T;
7     __inline bool getchar(char &ch) {
8         if (S == T) {
9             T = (S = buffer) + fread(buffer, 1, L, stdin);
10            if (S == T) {
11                ch = EOF;
12                return false;
13            }
14        }
15        ch = *S++;
16        return true;
17    }
18    __inline bool getint(int &x) {
19        char ch; bool neg = 0;
20        for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^=
21 // ch == '-';
22        if (ch == EOF) return false;
23        x = ch - '0';
24        for (; getchar(ch), ch >= '0' && ch <= '9'; )
25            x = x * 10 + ch - '0';
26        if (neg) x = -x;
27        return true;
28    }

```

## 8.4 汇编技巧

```

1 03优化
2 #define __ __attribute__ ((optimize("-O3")))
3 #define __ __inline __attribute__ ((__gnu_inline__,
4 // __always_inline__, __artificial__))
5 汇编开栈
6 #pragma comment(linker, "/STACK:256000000")
7
8 int __size = 256 << 20;

```

```

9 char* __p__ = (char *) malloc(__size__) + __size__;
10
11 int main() {
12     __asm__("movl %0, %%esp\n" :: "r"(__p__));
13     return 0;
14 }

```

## 9. 提示

### 9.1 线性规划转对偶

$\text{maximize } \mathbf{c}^T \mathbf{x}$   
 $\text{subject to } \mathbf{Ax} \leq \mathbf{b}, \mathbf{x} \geq 0 \iff \text{minimize } \mathbf{y}^T \mathbf{b}$   
 $\text{subject to } \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T, \mathbf{y} \geq 0$

### 9.2 NTT 素数及其原根

Prime	Primitive root
1053818881	7
1051721729	6
1045430273	3
1012924417	5
1007681537	3

### 9.3 积分表

$$1. \int \frac{dx}{ax^2+bx+c} = \begin{cases} \frac{2}{\sqrt{4ac-b^2}} \arctan \frac{2ax+b}{\sqrt{4ac-b^2}} + C & (b^2 < 4ac) \\ \frac{1}{\sqrt{b^2-4ac}} \ln \left| \frac{2ax+b-\sqrt{b^2-4ac}}{2ax+b+\sqrt{b^2-4ac}} \right| + C & (b^2 > 4ac) \end{cases}$$

$$2. \int \frac{x}{ax^2+bx+c} dx = \frac{1}{2a} \ln |ax^2+bx+c| - \frac{b}{2a} \int \frac{dx}{ax^2+bx+c}$$

$$1. \int \frac{dx}{\sqrt{ax^2+bx+c}} = \frac{1}{\sqrt{a}} \ln |2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}| + C$$

$$2. \int \sqrt{ax^2+bx+c} dx = \frac{2ax+b}{4a} \sqrt{ax^2+bx+c} + \frac{4ac-b^2}{8\sqrt{a^3}} \ln |2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}| + C$$

$$3. \int \frac{x}{\sqrt{ax^2+bx+c}} dx = \frac{1}{a} \sqrt{ax^2+bx+c} - \frac{b}{2\sqrt{a^3}} \ln |2ax+b+2\sqrt{a}\sqrt{ax^2+bx+c}| + C$$

$$4. \int \frac{dx}{\sqrt{c+bx-ax^2}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

$$5. \int \sqrt{c+bx-ax^2} dx = \frac{2ax-b}{4a} \sqrt{c+bx-ax^2} + \frac{b^2+4ac}{8\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

$$6. \int \frac{x}{\sqrt{c+bx-ax^2}} dx = -\frac{1}{a} \sqrt{c+bx-ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

$$1. \int \frac{dx}{\sqrt{(x-a)(b-x)}} = 2 \arcsin \sqrt{\frac{x-a}{b-x}} + C \quad (a < b)$$

2.

$$\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C, (a < b) \quad (1)$$

$$1. \int \tan x dx = -\ln |\cos x| + C$$

$$2. \int \cot x dx = \ln |\sin x| + C$$

$$3. \int \sec x dx = \ln \left| \tan \left( \frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln |\sec x + \tan x| + C$$

$$4. \int \csc x dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln |\csc x - \cot x| + C$$

$$5. \int \sec^2 x dx = \tan x + C$$

$$6. \int \csc^2 x dx = -\cot x + C$$

$$7. \int \sec x \tan x dx = \sec x + C$$

$$8. \int \csc x \cot x dx = -\csc x + C$$

$$9. \int \sin^2 x dx = \frac{x}{2} - \frac{1}{4} \sin 2x + C$$

$$10. \int \cos^2 x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$$

$$11. \int \sin^n x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$$

$$12. \int \cos^n x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$$

$$13. \int \frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$$

$$14. \int \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$$

15.

$$\begin{aligned} & \int \cos^m x \sin^n x dx \\ &= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx \\ &= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx \end{aligned}$$

$$16. \int \frac{dx}{a+b \sin x} = \begin{cases} \frac{2}{\sqrt{a^2-b^2}} \arctan \frac{a \tan \frac{x}{2} + b}{\sqrt{a^2-b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2-a^2}} \ln \left| \frac{a \tan \frac{x}{2} + b - \sqrt{b^2-a^2}}{a \tan \frac{x}{2} + b + \sqrt{b^2-a^2}} \right| + C & (a^2 < b^2) \end{cases}$$

$$17. \int \frac{dx}{a+b \cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan \left( \sqrt{\frac{a-b}{a+b}} \tan \frac{x}{2} \right) + C & (a^2 > b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln \left| \frac{\tan \frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan \frac{x}{2} - \sqrt{\frac{a+b}{b-a}}} \right| + C & (a^2 < b^2) \end{cases}$$

$$18. \int \frac{dx}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan \left( \frac{b}{a} \tan x \right) + C$$

$$19. \int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x - a} \right| + C$$

$$20. \int x \sin ax dx = \frac{1}{a^2} \sin ax - \frac{1}{a} x \cos ax + C$$

$$21. \int x^2 \sin ax dx = -\frac{1}{a} x^2 \cos ax + \frac{2}{a^2} x \sin ax + \frac{2}{a^3} \cos ax + C$$

$$22. \int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$$

$$23. \int x^2 \cos ax dx = \frac{1}{a} x^2 \sin ax + \frac{2}{a^2} x \cos ax - \frac{2}{a^3} \sin ax + C$$

$$1. \int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 - x^2} + C$$

$$2. \int x \arcsin \frac{x}{a} dx = \left( \frac{x^2}{2} - \frac{a^2}{4} \right) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 - a^2} + C$$

$$3. \int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 - x^2} + C$$

$$4. \int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} - \sqrt{a^2 - x^2} + C$$

$$5. \int x \arccos \frac{x}{a} dx = \left( \frac{x^2}{2} - \frac{a^2}{4} \right) \arccos \frac{x}{a} - \frac{x}{4} \sqrt{a^2 - x^2} + C$$

$$6. \int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} - \frac{1}{9} (x^2 + 2a^2) \sqrt{a^2 - x^2} + C$$

$$7. \int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2 + x^2) + C$$

$$8. \int x \arctan \frac{x}{a} dx = \frac{1}{2} (a^2 + x^2) \arctan \frac{x}{a} - \frac{a}{2} x + C$$

$$9. \int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} - \frac{a}{6} x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$$

$$1. \int a^x dx = \frac{1}{\ln a} a^x + C$$

$$2. \int e^{ax} dx = \frac{1}{a} e^{ax} + C$$

$$3. \int x e^{ax} dx = \frac{1}{a^2} (ax - 1) e^{ax} + C$$

$$4. \int x^n e^{ax} dx = \frac{1}{a} x^n e^{ax} - \frac{n}{a} \int x^{n-1} e^{ax} dx$$

$$5. \int x a^x dx = \frac{x}{\ln a} a^x - \frac{1}{(\ln a)^2} a^x + C$$

$$6. \int x^n a^x dx = \frac{1}{\ln a} x^n a^x - \frac{n}{\ln a} \int x^{n-1} a^x dx$$

$$7. \int e^{ax} \sin bx dx = \frac{1}{a^2+b^2} e^{ax} (a \sin bx - b \cos bx) + C$$

$$8. \int e^{ax} \cos bx dx = \frac{1}{a^2+b^2} e^{ax} (b \sin bx + a \cos bx) + C$$

$$9. \int e^{ax} \sin^n bxdx = \frac{1}{a^2+b^2n^2} e^{ax} \sin^{n-1} bx (a \sin bx - nb \cos bx) + \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \sin^{n-2} bxdx$$

$$10. \int e^{ax} \cos^n bxdx = \frac{1}{a^2+b^2n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) + \frac{n(n-1)b^2}{a^2+b^2n^2} \int e^{ax} \cos^{n-2} bxdx$$

$$1. \int \ln x dx = x \ln x - x + C$$

$$2. \int \frac{dx}{x \ln x} = \ln |\ln x| + C$$

$$3. \int x^n \ln x dx = \frac{1}{n+1} x^{n+1} \left( \ln x - \frac{1}{n+1} \right) + C$$

$$4. \int (\ln x)^n dx = x (\ln x)^n - n \int (\ln x)^{n-1} dx$$

$$5. \int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n - \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$$