Wizards

Standard Code Library

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1. 数论

1.1 $O(m^2 \log n)$ 线性递推

```
Given a_0, a_1, \ldots, a_{m-1}

a_n = c_0 \times a_{n-m} + \cdots + c_{m-1} \times a_{n-1}

Solve for a_n = v_0 \times a_0 + v_1 \times a_1 + \cdots + v_{m-1} \times a_{m-1}
```

```
1
   void linear_recurrence(long long n, int m, int a[], int
      \hookrightarrow c[], int p) {
2
     long long v[M] = \{1 \% p\}, u[M << 1], msk = !!n;
     for(long long i(n); i > 1; i >>= 1) {
3
       msk <<= 1;
5
     }
6
     for(long long x(0); msk; msk >>= 1, x <<= 1) {
7
       fill_n(u, m \ll 1, 0);
       int b(!!(n & msk));
8
9
       x \mid = b:
10
       if(x < m) {
11
         u[x] = 1 \% p;
12
       }else {
13
         for(int i(0); i < m; i++) {
            for(int j(0), t(i + b); j < m; j++, t++) {
14
              u[t] = (u[t] + v[i] * v[j]) % p;
16
         }
17
          for(int i((m << 1) - 1); i >= m; i--) {
18
19
            for(int j(0), t(i - m); j < m; j++, t++) {
              u[t] = (u[t] + c[j] * u[i]) % p;
20
21
         }
22
23
24
       copy(u, u + m, v);
25
     //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] *
26

→ a[m - 1].

27
     for(int i(m); i < 2 * m; i++) {</pre>
28
       a[i] = 0:
29
       for(int j(0); j < m; j++) {</pre>
30
          a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31
     }
32
33
     for(int j(0); j < m; j++) {
34
       b[j] = 0;
35
       for(int i(0); i < m; i++) {</pre>
36
         b[j] = (b[j] + v[i] * a[i + j]) % p;
37
     }
38
39
     for(int j(0); j < m; j++) {</pre>
40
       a[j] = b[j];
     }
41
42 }
```

1.2 求逆元

```
void ex_gcd(long long a, long long b, long long &x, long
      \hookrightarrow long &y) {
2
     if (b == 0) {
       x = 1;
3
       y = 0;
5
       return;
6
     }
7
     long long xx, yy;
8
     ex_gcd(b, a % b, xx, yy);
9
     y = xx - a / b * yy;
10
     x = yy;
11
12
   long long inv(long long x, long long MODN) {
13
     long long inv_x, y;
     ex_gcd(x, MODN, inv_x, y);
14
     return (inv_x % MODN + MODN) % MODN;
15
16 }
```

1.3 中国剩余定理

```
//返回 (ans, M), 其中 ans 是模 M 意义下的解
  std::pair<long long, long long> CRT(const std::vector<long</pre>
     \hookrightarrow long>& m, const std::vector<long long>& a) {
     long long M = 1, ans = 0;
     int n = m.size();
5
     for (int i = 0; i < n; i++) M *= m[i];</pre>
6
     for (int i = 0; i < n; i++) {
7
       ans = (ans + (M / m[i]) * a[i] % M * inv(M / m[i],
          →m[i])) % M; // 可能需要大整数相乘取模
8
    }
9
     return std::make_pair(ans, M);
10
```

1.4 素性测试

```
int strong_pseudo_primetest(long long n,int base) {
 2
       long long n2=n-1,res;
 3
       int s=0:
       while(n2\%2==0) n2>>=1,s++;
 4
5
       res=powmod(base,n2,n);
6
       if((res==1)||(res==n-1)) return 1;
 8
       while(s \ge 0) {
9
           res=mulmod(res,res,n);
10
            if(res==n-1) return 1;
11
12
       }
13
       return 0; // n is not a strong pseudo prime
14
   }
15
   int isprime(long long n) {
16
     static LL testNum[]={2,3,5,7,11,13,17,19,23,29,31,37};
17
     static LL lim[]={4,0,1373653LL,25326001LL,25000000000LL,
18
     2152302898747LL.

→ 3474749660383LL,341550071728321LL,0,0,0,0);

19
     if(n<2||n==3215031751LL) return 0;
20
     for(int i=0;i<12;++i){</pre>
21
       if(n<lim[i]) return 1;</pre>
       if(strong_pseudo_primetest(n,testNum[i])==0) return 0;
22
23
     }
24
     return 1;
25
   }
```

1.5 质因数分解

```
int ansn; LL ans[1000];
   LL func(LL x,LL n){ return(mod_mul(x,x,n)+1)%n; }
   LL Pollard(LL n){
     LL i,x,y,p;
     if(Rabin_Miller(n)) return n;
6
     if(!(n&1)) return 2;
 7
     for(i=1;i<20;i++){
8
       x=i; y=func(x,n); p=gcd(y-x,n);
 9
       while(p==1) {x=func(x,n); y=func(func(y,n),n);
          \hookrightarrow p=\gcd((y-x+n)%n,n)%n;
10
       if(p==0||p==n) continue;
11
       return p;
     }
12
13
   }
14
   void factor(LL n){
15
     LL x;
16
     x=Pollard(n);
     if(x==n){ ans[ansn++]=x; return; }
17
     factor(x), factor(n/x);
18
19
```

1.6 佩尔方程

```
import java.math.BigInteger;
import java.util.Scanner;
```

```
3/a[n]=(g[n]+a[0])/h[n]
4 //g[n] = a[n-1] *h[n-1] -g[n-1]
5 / h[n] = (N-g[n]*g[n])/h[n-1]
6 //p[n]=a[n-1]*p[n-1]+p[n-2]
7 / q[n] = a[n-1] * q[n-1] + q[n-2]
8 //so:
   //p[n]*q[n-1]-p[n-1]*q[n]=(-1)^(n+1);
9
   //p[n]^2-N*q[n]^2=(-1)^(n+1)*h[n+1];
11
   public class Main {
12
       public static BigInteger p, q;
13
       public static void solve(int n) {
14
            BigInteger N, p1, p2, q1, q2, a0, a1, a2, g1, g2,
               \hookrightarrow h1, h2;
            g1 = q2 = p1 = BigInteger.ZERO;
15
16
            h1 = q1 = p2 = BigInteger.ONE;
17
            a0 = a1 =

→ BigInteger.valueOf((long)Math.sqrt(1.0*n));
            N = BigInteger.valueOf(n);
18
19
            while (true) {
                g2 = a1.multiply(h1).subtract(g1);
20
21
                h2 = N.subtract(g2.pow(2)).divide(h1);
                a2 = g2.add(a0).divide(h2);
22
                p = a1.multiply(p2).add(p1);
23
24
                q = a1.multiply(q2).add(q1);
25
                if
                   \hookrightarrow (p.pow(2).subtract(N.multiply(q.pow(2))).compareTo(BigInteger.ONE)
                   \hookrightarrow == 0) return;
26
          g1 = g2;h1 = h2;a1 = a2;
27
                p1 = p2; p2 = p;
28
                q1 = q2; q2 = q;
29
            }
30
31
32
        public static void main(String[] args) {
            Scanner cin = new Scanner(System.in);
33
            int t=cin.nextInt();
34
            while (t--!=0) {
35
36
                solve(cin.nextInt());
                System.out.println(p + " " + q);
37
38
39
       }
40 }
```

1.7 二次剩余

```
1 // x^2 = a (mod p), 0 <= a < p, 返回 true or false 代表
     →是否存在解
  // p 必须是质数, 若是多个单次质数的乘积, 可以分别
     →求解再用 CRT 合并
  // 复杂度为 D(log n)
3
   void multiply(ll &c, ll &d, ll a, ll b, ll w) {
       int cc = (a * c + b * d % MOD * w) % MOD;
      int dd = (a * d + b * c) % MOD;
6
7
      c = cc, d = dd;
8 }
  bool solve(int n, int &x) {
10
      if (MOD == 2) return x = 1, true;
11
      if (power(n, MOD / 2, MOD) == MOD - 1) return false;
12
      11 c = 1, d = 0, b = 1, a, w;
13
      // finding a such that a^2 - n is not a square
14
15
      do { a = rand() % MOD;
          w = (a * a - n + MOD) \% MOD;
16
          if (w == 0) return x = a, true;
17
      } while (power(w, MOD / 2, MOD) != MOD - 1);
18
      for (int times = (MOD + 1) / 2; times; times >>= 1) {
19
20
          if (times & 1) multiply(c, d, a, b, w);
```

```
21
           multiply(a, b, a, b, w);
22
23
       // x = (a + sqrt(w)) ^ ((p + 1) / 2)
24
       return x = c, true;
25
```

1.8 一元三次方程

```
double a(p[3]), b(p[2]), c(p[1]), d(p[0]);
   double k(b / a), m(c / a), n(d / a);
   double p(-k * k / 3. + m);
   double q(2. * k * k * k / 27 - k * m / 3. + n);
   Complex omega[3] = \{Complex(1, 0), Complex(-0.5, 0.5 *
      \hookrightarrow sqrt(3)), Complex(-0.5, -0.5 * sqrt(3))};
   Complex r1, r2;
   double delta(q * q / 4 + p * p * p / 27);
 8
   if (delta > 0) {
9
       r1 = cubrt(-q / 2. + sqrt(delta));
10
       r2 = cubrt(-q / 2. - sqrt(delta));
11
   } else {
12
       r1 = pow(-q / 2. + pow(Complex(delta), 0.5), 1. / 3);
13
       r2 = pow(-q / 2. - pow(Complex(delta), 0.5), 1. / 3);
   }
14
15
   for(int _(0); _ < 3; _++) {
       Complex x = -k / 3. + r1 * omega[_ * 1] + r2 * omega[_
16
          \hookrightarrow * 2 % 3];
17
```

线下整点 1.9

```
// \sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor, n, m, a, b > 0
2
  LL solve(LL n,LL a,LL b,LL m){
     if(b==0) return n*(a/m);
     if(a>=m) return n*(a/m)+solve(n,a%m,b,m);
5
     if(b>=m) return (n-1)*n/2*(b/m)+solve(n,a,b%m,m);
6
     return solve((a+b*n)/m,(a+b*n)%m,m,b);
```

1.10 线性同余不等式

```
// Find the minimal non-negtive solutions for
      rightarrow l \le d \cdot x \mod m \le r
   // 0 \le d, l, r < m; l \le r, O(\log n)
   11 cal(11 m, 11 d, 11 1, 11 r) {
       if (1 == 0) return 0;
       if (d == 0) return MXL; // 无解
5
       if (d * 2 > m) return cal(m, m - d, m - r, m - 1);
6
7
       if ((1 - 1) / d < r / d) return (1 - 1) / d + 1;
8
       ll k = cal(d, (-m % d + d) % d, l % d, r % d);
9
       return k == MXL ? MXL : (k * m + l - 1) / d + 1; // 无
          →解 2
10 }
```

组合数取模 1.11

```
LL prod=1,P;
   pair<LL,LL> comput(LL n,LL p,LL k){
 3
        if(n<=1)return make_pair(0,1);</pre>
       LL ans=1,cnt=0;
        ans=pow(prod,n/P,P);
 6
 7
        pair<LL,LL>res=comput(n/p,p,k);
8
        cnt+=res.first:
9
        ans=ans*res.second%P:
10
        for(int i=n-n%P+1;i<=n;i++)if(i%p){</pre>
11
12
            ans=ans*i%P;
       }
13
14
        return make_pair(cnt,ans);
15 }
```

```
pair<LL,LL> calc(LL n,LL p,LL k){
       prod=1;P=pow(p,k,1e18);
17
18
       for(int i=1;i<P;i++)if(i%p)prod=prod*i%P;</pre>
19
       pair<LL,LL> res=comput(n,p,k);
20
       return res;
21 }
  LL calc(LL n,LL m,LL p,LL k){
       pair<LL,LL>A,B,C;
23
24
       LL P=pow(p,k,1e18);
25
       A=calc(n,p,k);
26
       B=calc(m,p,k);
       C=calc(n-m,p,k);
27
28
       LL ans=1:
       ans=pow(p,A.first-B.first-C.first,P);
29
30
           \hookrightarrow ans=ans*A.second%P*inv(B.second,P)%P*inv(C.second,P)%P;60
31
       return ans:
32 }
```

1.12 Schreier-Sims

```
struct Perm{
     vector<int> P; Perm() {} Perm(int n) { P.resize(n); }
2
3
     Perm inv()const{
       Perm ret(P.size());
5
       for(int i = 0; i < int(P.size()); ++i) ret.P[P[i]] =</pre>
6
       return ret;
7
     }
8
     int &operator [](const int &dn){ return P[dn]; }
9
     void resize(const size_t &sz){ P.resize(sz); }
10
     size_t size()const{ return P.size(); }
     const int &operator [](const int &dn)const{ return
11
        \hookrightarrow P[dn]; }
12 };
13 | Perm operator *(const Perm &a, const Perm &b){
14
     Perm ret(a.size()):
15
     for(int i = 0; i < (int)a.size(); ++i) ret[i] = b[a[i]];</pre>
16
17 }
18 typedef vector<Perm> Bucket;
19
   typedef vector<int> Table;
20 typedef pair<int,int> PII;
21
   int n, m;
   vector<Bucket> buckets, bucketsInv; vector<Table>

→ lookupTable;

23
   int fastFilter(const Perm &g, bool addToGroup = true) {
     int n = buckets.size();
24
25
     Perm p(g);
    for(int i = 0; i < n; ++i){
26
27
       int res = lookupTable[i][p[i]];
28
       if(res == -1){
29
         if (addToGroup) {
30
           buckets[i].push_back(p);
              → bucketsInv[i].push_back(p.inv());
           lookupTable[i][p[i]] = (int)buckets[i].size() - 1;
31
32
33
         return i;
34
       p = p * bucketsInv[i][res];
35
36
37
     return -1;
38 }
   long long calcTotalSize(){
39
40
     long long ret = 1;
    for(int i = 0; i < n; ++i) ret *= buckets[i].size();</pre>
41
42
     return ret;
43
   }
   bool inGroup(const Perm &g){ return fastFilter(g, false)
45 | void solve(const Bucket &gen,int _n){// m perm[0..n - 1]s
    n = _n, m = gen.size();
46
    {//clear all
47
```

```
vector<Bucket> _buckets(n); swap(buckets, _buckets);
49
       vector<Bucket> _bucketsInv(n); swap(bucketsInv,

→ _bucketsInv);
50
       vector<Table> _lookupTable(n); swap(lookupTable,
          51
     for(int i = 0; i < n; ++i){
52
53
       lookupTable[i].resize(n);
54
       fill(lookupTable[i].begin(), lookupTable[i].end(),
          \hookrightarrow -1):
55
     }
56
     Perm id(n);
57
     for(int i = 0; i < n; ++i) id[i] = i;
     for(int i = 0; i < n; ++i){
58
59
       buckets[i].push_back(id); bucketsInv[i].push_back(id);
       lookupTable[i][i] = 0;
61
62
     for(int i = 0; i < m; ++i) fastFilter(gen[i]);</pre>
     queue<pair<PII,PII> > toUpdate;
63
     for(int i = 0; i < n; ++i)
65
       for(int j = i; j < n; ++j)
         for(int k = 0; k < (int)buckets[i].size(); ++k)</pre>
66
           for(int 1 = 0; 1 < (int)buckets[j].size(); ++1)</pre>
67
68
              toUpdate.push(make_pair(PII(i,k), PII(j,l)));
69
     while(!toUpdate.empty()){
70
       PII a = toUpdate.front().first, b =
          71
       toUpdate.pop();
       int res = fastFilter(buckets[a.first][a.second] *
72

    buckets[b.first][b.second]);

73
       if(res==-1) continue;
       PII newPair(res, (int)buckets[res].size() - 1);
74
75
       for(int i = 0; i < n; ++i)
         for(int j = 0; j < (int)buckets[i].size(); ++j){
76
77
           if(i <= res) toUpdate.push(make_pair(PII(i, j),</pre>
              → newPair));
78
           if(res <= i) toUpdate.push(make_pair(newPair,</pre>
              \hookrightarrow PII(i, j)));
79
         }
80
     }
81
   }
```

2. 代数

2.1 快速傅里叶变换

```
void fft(Complex a[], int n, int f) {
 2
     for (int i = 0; i < n; ++i)
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
3
 4
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
 5
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
 6
       Complex w = Complex(1, 0);
 7
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
8
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
9
         for (int k = 0; k < i; ++k) {
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
10
11
           a[j + k] = x + y; a[j + k + i] = x - y;
12
13
       }
14
     }
15
   }
```

2.2 分治卷积

```
1 // n 必须是 2 的次幂
2 void fft(Complex a[], int n, int f) {
3 for (int i = 0; i < n; ++i)
4 if (R[i] < i) swap(a[i], a[R[i]]);
5 for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6 Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7 Complex w = Complex(1, 0);
```

```
for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
for (int p = i << 1, j = 0; j < n; j += p) {
  for (int k = 0; k < i; ++k) {
    Complex x = a[j + k], y = a[j + k + i] * tmp[k];
    a[j + k] = x + y; a[j + k + i] = x - y;
}
}
}
}
}
}
}
</pre>
```

2.3 快速数论变换

```
// n 必须是 2 的次幂
   void fft(Complex a[], int n, int f) {
2
     for (int i = 0; i < n; ++i)
3
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
5
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
       Complex w = Complex(1, 0);
7
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
10
         for (int k = 0; k < i; ++k) {
11
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
       }
     }
15
16 }
```

2.4 快速沃尔什变换

```
void FWT(LL a[],int n,int ty){ //the length is 2^n
2
     for(int d=1;d<n;d<<=1){</pre>
3
       for(int m=(d<<1),i=0;i<n;i+=m){</pre>
4
          if(ty==1){
5
            for(int j=0; j<d; j++){
6
              LL x=a[i+j], y=a[i+j+d];
7
              a[i+j]=x+y;
8
              a[i+j+d]=x-y;
9
                       //and:a[i+j]=x+y; or:a[i+j+d]=x+y;
10
            }
11
         }else{
            for(int j=0;j<d;j++){
12
13
              LL x=a[i+j], y=a[i+j+d];
14
              a[i+j]=(x+y)/2;
15
              a[i+j+d]=(x-y)/2;
16
                       //and:a[i+j]=x-y; or:a[i+j+d]=y-x;
17
18
19
20
```

2.5 自适应辛普森积分

```
namespace adaptive_simpson {
2
    template<typename function>
    inline double area(function f, const double &left, const
3
       double mid = (left + right) / 2;
5
      return (right - left) * (f(left) + 4 * f(mid) +
         \hookrightarrow f(right)) / 6;
6
    template<typename function>
8
    inline double simpson(function f, const double &left,
       \hookrightarrow double &area_sum) {
9
      double mid = (left + right) / 2;
10
      double area_left = area(f, left, mid);
      double area_right = area(f, mid, right);
11
      double area_total = area_left + area_right;
12
```

```
13
      if (fabs(area_total - area_sum) <= 15 * eps) {</pre>
14
        return area_total + (area_total - area_sum) / 15;
15
16
       return simpson(f, left, right, eps / 2, area_left) +
         17
18
     template<typename function>
19
     inline double simpson(function f, const double &left,
       \hookrightarrow const double &right, const double &eps) {
20
       return simpson(f, left, right, eps, area(f, left,

    right));
21
22
```

2.6 单纯形

```
const double eps = 1e-8;
   // max{c * x | Ax <= b, x >= 0} 的解, 无解返回空的
      → vector, 否则就是解.
   vector<double> simplex(vector<vector<double> > &A,
      \hookrightarrow vector<double> b, vector<double> c) {
     int n = A.size(), m = A[0].size() + 1, r = n, s = m - 1;
     vector < vector < double > D(n + 2, vector < double > (m + 1));
 5
     vector<int> ix(n + m);
6
     for(int i = 0; i < n + m; i++) {</pre>
 7
 8
       ix[i] = i;
9
10
     for(int i = 0; i < n; i++) {
11
       for(int j = 0; j < m - 1; j++) {
12
         D[i][j] = -A[i][j];
13
       D[i][m - 1] = 1;
14
15
       D[i][m] = b[i];
       if (D[r][m] > D[i][m]) {
16
17
         r = i;
18
       }
19
     }
20
     for(int j = 0; j < m - 1; j++) {
21
       D[n][j] = c[j];
22
     D[n + 1][m - 1] = -1;
23
     for(double d; ;) {
24
25
       if (r < n) {</pre>
26
         swap(ix[s], ix[r + m]);
27
         D[r][s] = 1. / D[r][s];
28
         for(int j = 0; j \le m; j++) {
           if (j != s) {
29
              D[r][j] *= -D[r][s];
30
31
32
         }
33
         for(int i = 0; i <= n + 1; i++) {
34
            if (i != r) {
35
             for(int j = 0; j \le m; j++) {
                if (j != s) {
36
37
                  D[i][j] += D[r][j] * D[i][s];
38
39
             D[i][s] *= D[r][s];
40
41
            }
         }
42
       }
43
       r = -1, s = -1;
44
45
       for(int j = 0; j < m; j++) {
46
         if (s < 0 || ix[s] > ix[j]) {
            if (D[n + 1][j] > eps || D[n + 1][j] > -eps &&
47
               \hookrightarrow D[n][j] > eps) {
              s = j;
48
            }
49
50
         }
       }
51
       if (s < 0) break:
52
       for(int i = 0; i < n; i++) {
53
```

```
if (D[i][s] < -eps) {</pre>
            if (r < 0 \mid | (d = D[r][m] / D[r][s] - D[i][m] /
55
               \hookrightarrow D[i][s]) < -eps
56
               || d < eps && ix[r + m] > ix[i + m]) {
57
              r = i;
58
          }
59
60
        }
61
        if (r < 0) return vector<double> ();
62
      }
63
      if (D[n + 1][m] < -eps) return vector<double> ();
      vector<double> x(m - 1);
64
65
      for(int i = m; i < n + m; i++) {</pre>
        if (ix[i] < m - 1) {</pre>
66
          x[ix[i]] = D[i - m][m];
67
68
     }
69
70
     return x:
71 }
```

3. 计算几何

3.1 二维

3.1.1 点类

```
1 int sign(DB x) {
2
    return (x > eps) - (x < -eps);
3
  ۱,
4 DB msqrt(DB x) {
5
    return sign(x) > 0 ? sqrt(x) : 0;
6 }
7
   struct Point {
8
    DB x, y;
    Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
Q
       return Point(cos(ang) * x - sin(ang) * y, cos(ang) * y
10
          \hookrightarrow + sin(ang) * x);
11
    }
    Point turn90() const { // 逆时针旋转 90 度
12
13
       return Point(-y, x);
14
15
    Point unit() const {
16
       return *this / len();
17
18
  1:
19
   DB dot(const Point& a, const Point& b) {
20
    return a.x * b.x + a.y * b.y;
21 }
22
  DB det(const Point& a, const Point& b) {
23
    return a.x * b.y - a.y * b.x;
24 }
   #define cross(p1,p2,p3)
     \leftrightarrow ((p2.x-p1.x)*(p3.y-p1.y)-(p3.x-p1.x)*(p2.y-p1.y))
   #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
  bool isLL(const Line& 11, const Line& 12, Point& p) { //
     → 直线与直线交点
28
     DB s1 = det(12.b - 12.a, 11.a - 12.a),
29
        s2 = -det(12.b - 12.a, 11.b - 12.a);
    if (!sign(s1 + s2)) return false;
30
     p = (11.a * s2 + 11.b * s1) / (s1 + s2);
31
32
     return true;
33
   bool onSeg(const Line& 1, const Point& p) { // 点在线段
34
     return sign(det(p - 1.a, 1.b - 1.a)) == 0 && sign(dot(p
35
       \hookrightarrow - 1.a, p - 1.b)) <= 0;
36
   Point projection(const Line & 1, const Point& p) {
37
    return 1.a + (1.b - 1.a) * (dot(p - 1.a, 1.b - 1.a) /
        \hookrightarrow (1.b - 1.a).len2());
39 }
40 DB disToLine(const Line& 1, const Point& p) { // 点到 *
     → 直线 * 距离
```

```
return fabs(det(p - 1.a, 1.b - 1.a) / (1.b -
        \hookrightarrow 1.a).len());
42
   }
43
   DB disToSeg(const Line& 1, const Point& p) { // 点到线段
     return sign(dot(p - 1.a, 1.b - 1.a)) * sign(dot(p - 1.b,
        \rightarrow l.a - l.b)) == 1 ? disToLine(l, p) : std::min((p -
        \hookrightarrow 1.a).len(), (p - 1.b).len());
45
   }
   // 圆与直线交点
46
47
   bool isCL(Circle a, Line 1, Point& p1, Point& p2) {
     DB x = dot(1.a - a.o, 1.b - 1.a),
48
49
        y = (1.b - 1.a).len2(),
        d = x * x - y * ((1.a - a.o).len2() - a.r * a.r);
51
     if (sign(d) < 0) return false;</pre>
52
     Point p = 1.a - ((1.b - 1.a) * (x / y)), delta = (1.b - 1.a) * (x / y))
        \hookrightarrow 1.a) * (msqrt(d) / y);
     p1 = p + delta; p2 = p - delta;
53
54
     return true;
55
   //圆与圆的交面积
56
   DB areaCC(const Circle& c1, const Circle& c2) {
57
58
     DB d = (c1.o - c2.o).len();
     if (sign(d - (c1.r + c2.r)) >= 0) return 0;
59
     if (sign(d - std::abs(c1.r - c2.r)) \le 0) {
60
61
       DB r = std::min(c1.r, c2.r);
62
       return r * r * PI;
63
64
     DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
65
       t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
     return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r *
66
        \hookrightarrow \sin(t1);
67
68
   // 圆与圆交点
   bool isCC(Circle a, Circle b, P& p1, P& p2) {
     DB s1 = (a.o - b.o).len();
     if (sign(s1 - a.r - b.r) > 0 \mid \mid sign(s1 - std::abs(a.r -
        \hookrightarrow b.r)) < 0) return false;
72
     DB s2 = (a.r * a.r - b.r * b.r) / s1;
     DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
73
     P \circ = (b.o - a.o) * (aa / (aa + bb)) + a.o;
75
     P delta = (b.o - a.o).unit().turn90() * msqrt(a.r * a.r
        \hookrightarrow - aa * aa);
76
     p1 = o + delta, p2 = o - delta;
77
     return true;
78
   // 求点到圆的切点,按关于点的顺时针方向返回两个点
79
   bool tanCP(const Circle &c, const Point &p0, Point &p1,
      → Point &p2) {
     double x = (p0 - c.o).len2(), d = x - c.r * c.r;
     if (d < eps) return false; // 点在圆上认为没有切点
83
     Point p = (p0 - c.o) * (c.r * c.r / x);
     Point delta = ((p0 - c.o) * (-c.r * sqrt(d) /
        \hookrightarrow x)).turn90();
85
     p1 = c.o + p + delta;
     p2 = c.o + p - delta;
     return true:
88
   }
   // 求圆到圆的外共切线, 按关于 c1.o 的顺时针方向返
89
   vector<Line> extanCC(const Circle &c1, const Circle &c2) {
     vector<Line> ret:
     if (sign(c1.r - c2.r) == 0) {
93
       Point dir = c2.o - c1.o;
       dir = (dir * (c1.r / dir.len())).turn90();
94
95
       ret.push_back(Line(c1.o + dir, c2.o + dir));
96
       ret.push_back(Line(c1.o - dir, c2.o - dir));
97
     } else {
       Point p = (c1.0 * -c2.r + c2.o * c1.r) / (c1.r - c2.r + c2.o * c1.r) / (c1.r - c2.r + c2.o * c1.r)
98
          \hookrightarrow c2.r):
99
       Point p1, p2, q1, q2;
```

```
if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) {
100
          if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);
101
102
          ret.push_back(Line(p1, q1));
103
          ret.push_back(Line(p2, q2));
104
     }
105
      return ret;
   }
107
    // 求圆到圆的内共切线, 按关于 c1.o 的顺时针方向返
109 std::vector<Line> intanCC(const Circle &c1, const Circle
      std::vector<Line> ret;
111
     Point p = (c1.0 * c2.r + c2.o * c1.r) / (c1.r + c2.r);
112
      Point p1, p2, q1, q2;
      if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { //
113
        →两圆相切认为没有切线
114
        ret.push_back(Line(p1, q1));
        ret.push_back(Line(p2, q2));
115
116
     }
117
     return ret;
118
   |bool contain(vector<Point> polygon, Point p) { // 判断点
119
      →p 是否被多边形包含,包括落在边界上
      int ret = 0, n = polygon.size();
120
      for(int i = 0; i < n; ++ i) {
121
122
        Point u = polygon[i], v = polygon[(i + 1) \% n];
123
        if (onSeg(Line(u, v), p)) return true; // Here I
124
        if (sign(u.y - v.y) \le 0) swap(u, v);
        if (sign(p.y - u.y) > 0 \mid \mid sign(p.y - v.y) \le 0)
125
           ret += sign(det(p, v, u)) > 0;
126
127
128
      return ret & 1;
129 }
   |// 用半平面 (q1,q2) 的逆时针方向去切凸多边形
130
131 | std::vector<Point> convexCut(const std::vector<Point>&ps,
      \hookrightarrow Point q1, Point q2) {
      std::vector<Point> qs; int n = ps.size();
132
      for (int i = 0; i < n; ++i) {</pre>
133
134
        Point p1 = ps[i], p2 = ps[(i + 1) % n];
135
        int d1 = crossOp(q1,q2,p1), d2 = crossOp(q1,q2,p2);
136
        if (d1 \ge 0) qs.push_back(p1);
        if (d1 * d2 < 0) qs.push_back(isSS(p1, p2, q1, q2));
137
138
139
     return qs;
140 }
   // 求凸包
141
    std::vector<Point> convexHull(std::vector<Point> ps) {
143
      int n = ps.size(); if (n <= 1) return ps;</pre>
144
      std::sort(ps.begin(), ps.end());
      std::vector<Point> qs;
145
      for (int i = 0; i < n; qs.push_back(ps[i ++]))</pre>
146
147
        while (qs.size() > 1 && sign(det(qs[qs.size() - 2],
           \hookrightarrow qs.back(), ps[i])) <= 0)
148
          qs.pop_back();
      for (int i = n - 2, t = qs.size(); i >= 0;
149
        \hookrightarrow \texttt{qs.push\_back(ps[i --]))}
        while ((int)qs.size() > t && sign(det(qs[qs.size() -
150
           \hookrightarrow 2], qs.back(), ps[i])) <= 0)
151
          qs.pop_back();
152
      return qs;
153 }
```

3.1.2 凸包

```
1 // 凸包中的点接逆时针方向
2 struct Convex {
3 int n;
4 std::vector<Point> a, upper, lower;
5 void make_shell(const std::vector<Point>& p,
```

```
std::vector<Point>& shell) { // p needs to be
6
 7
       clear(shell); int n = p.size();
       for (int i = 0, j = 0; i < n; i++, j++) {
8
         for (; j \ge 2 \&\& sign(det(shell[j-1] - shell[j-2]),
9
                 p[i] - shell[j-2])) \le 0; --j)
10
                    \hookrightarrow shell.pop_back();
11
         shell.push_back(p[i]);
12
       }
13
     }
14
     void make_convex() {
15
       std::sort(a.begin(), a.end());
16
       make_shell(a, lower);
17
       std::reverse(a.begin(), a.end());
18
       make_shell(a, upper);
19
       a = lower; a.pop_back();
20
       a.insert(a.end(), upper.begin(), upper.end());
21
       if ((int)a.size() >= 2) a.pop_back();
22
       n = a.size();
23
     }
24
     void init(const std::vector<Point>& _a) {
25
       clear(a); a = _a; n = a.size();
26
       make convex():
27
28
     void read(int _n) { // Won't make convex.
29
       clear(a); n = _n; a.resize(n);
       for (int i = 0; i < n; i++)
30
         a[i].read();
31
32
     std::pair<DB, int> get_tangent(
33
         const std::vector<Point>& convex, const Point& vec)
35
       int l = 0, r = (int)convex.size() - 2;
36
       assert(r >= 0);
37
       for (; l + 1 < r; ) {
38
         int mid = (1 + r) / 2;
         if (sign(det(convex[mid + 1] - convex[mid], vec)) >
            \rightarrow 0)
40
           r = mid;
         else 1 = mid;
41
42
43
       return std::max(std::make_pair(det(vec, convex[r]),
44
           std::make_pair(det(vec, convex[0]), 0));
45
     int binary_search(Point u, Point v, int 1, int r) {
46
47
       int s1 = sign(det(v - u, a[1 % n] - u));
48
       for (; 1 + 1 < r; ) {
49
         int mid = (1 + r) / 2;
50
         int smid = sign(det(v - u, a[mid % n] - u));
51
         if (smid == s1) l = mid;
52
         else r = mid;
53
54
       return 1 % n;
55
     // 求凸包上和向量 vec 叉积最大的点,返回编号,共
56
        →线的多个切点返回任意一个
     int get_tangent(Point vec) {
57
       std::pair<DB, int> ret = get_tangent(upper, vec);
58
59
       ret.second = (ret.second + (int)lower.size() - 1) % n;
60
       ret = std::max(ret, get_tangent(lower, vec));
61
       return ret.second;
62
     // 求凸包和直线 u, v 的交点, 如果不相交返回 false,
        →如果有则是和 (i, next(i)) 的交点,交在点上不
        →确定返回前后两条边其中之-
     bool get_intersection(Point u, Point v, int &i0, int
        \hookrightarrow \texttt{\&i1)} \ \{
65
       int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
       if (sign(det(v - u, a[p0] - u)) * sign(det(v - u, a[p0] - u))
66
          \hookrightarrow \texttt{a[p1] - u)) <= 0) \ \{
         if (p0 > p1) std::swap(p0, p1);
```

3.1.3 凸包最近点对

```
1 //判断点是否在多边形内
2
   int isPointInPolygon(point p, point *a, int n) {
3
       int cnt = 0;
4
       for(int i=0; i<n; ++i) {</pre>
5
           if(OnSegment(p, a[i], a[(i+1)%n])) return -1;
6
            double k = cross(a[(i+1)%n]-a[i], p-a[i]);
           double d1 = a[i].y - p.y;
7
       double d2 = a[(i+1)].y - p.y;
8
9
           if(k>0 &&d1<=0 &&d2>0) cnt++;
10
           if(k<0 &&d2<=0 &&d1>0) cnt++;
           //k==0, 点和线段共线的情况不考虑
11
12
13
       if(cnt&1)return 1;
14
       return 0;
15 }
  1/1判断凸包是否相离
16
   bool two_getaway_ConvexHull(point *cha, int n1, point
17
      \hookrightarrow *chb, int m1) {
18
       if(n1==1 && m1==1) {
19
            if (cha[0] == chb[0])
20
                return false;
21
       } else if(n1==1 && m1==2) {
22
           if(OnSegment(cha[0], chb[0], chb[1]))
23
                return false:
24
       } else if(n1==2 && m1==1) {
25
           if(OnSegment(chb[0], cha[0], cha[1]))
26
                return false;
27
       } else if(n1==2 && m1==2) {
           if(SegmentIntersection(cha[0], cha[1], chb[0],
28
              \hookrightarrow chb[1])
29
                return false;
30
       } else if(n1==2) {
31
           for(int i=0; i<n1; ++i)</pre>
32
                if(isPointInPolygon(cha[i], chb, m1))
33
                    return false:
       } else if(m1==2) {
34
35
           for(int i=0; i<m1; ++i)</pre>
36
                if(isPointInPolygon(chb[i], cha, n1))
37
                    return false;
38
       } else {
39
           for(int i=0; i<n1; ++i) {</pre>
                for(int j=0; j<m1; ++j) {</pre>
40
41
                    if(SegmentIntersection(cha[i],
                       \hookrightarrow cha[(i+1)%n1], chb[j],
                       \hookrightarrow chb[(j+1)\%m1]))
42
                        return false;
                }
43
           }
44
           for(int i=0; i<n1; ++i)</pre>
45
                if(isPointInPolygon(cha[i], chb, m1))
46
                    return false;
47
48
            for(int i=0; i<m1; ++i)</pre>
49
                if(isPointInPolygon(chb[i], cha, n1))
50
                    return false;
51
       }
52
       return true;
53 }
   //旋转卡壳求两个凸包最近距离
   double solve(point *P, point *Q, int n, int m) {
55
56
       if(n==1 \&\& m==1) {
           return length(P[0] - Q[0]);
57
       } else if(n==1 && m==2) {
58
```

```
return DistanceToSegment(P[0], Q[0], Q[1]);
59
60
       } else if(n==2 && m==1) {
61
            return DistanceToSegment(Q[0], P[0], P[1]);
62
       } else if(n==2 \&\& m==2) {
63
            return SegmentToSegment(P[0], P[1], Q[0], Q[1]);
64
65
66
       int yminP = 0, ymaxQ = 0;
       for(int i=0; i<n; ++i) if(P[i].y < P[yminP].y) yminP =</pre>
68
       for(int i=0; i < m; ++i) if(Q[i].y > Q[ymaxQ].y) ymaxQ =
          \hookrightarrow i:
69
       P[n] = P[0];
70
       Q[n] = Q[0];
       double INF2 = 1e100;
71
72
       double arg, ans = INF2;
73
       for(int i=0; i<n; ++i) {</pre>
            //当叉积负正转正时,说明点 ymaxQ 就是对踵点
74
75
            while((arg=cross(P[yminP] - P[yminP+1],Q[ymaxQ+1]
               \hookrightarrow - Q[ymaxQ])) < -eps)
                ymaxQ = (ymaxQ+1)%m;
76
77
            double ret;
            if(arg > eps) { //卡住第二个凸包上的点。
78
79
                ret = DistanceToSegment(Q[ymaxQ], P[yminP],
                   \hookrightarrow P[yminP+1]);
                ans = min(ans,ret);
80
            } else { //arg==0, 卡住第二个凸包的边
81
82

→ SegmentToSegment(P[yminP],P[yminP+1],Q[ymaxQ],Q
83
                ans = min(ans,ret);
84
85
            yminP = (yminP+1)%n;
       }
86
87
       return ans:
88
   }
89
   double mindis_twotubao(point *P, point *Q, int n, int m){
       //return min(solve(P, Q, n, m),solve(Q,P,m,n));
90
91
       if(two_getaway_ConvexHull(P,n,Q,m)==true) return
           \hookrightarrow \min(\text{solve}(P, Q, n, m), \text{solve}(Q, P, m, n));
92
       else return 0.0;
93
```

3.1.4 三角形的心

```
1 Point inCenter(const Point &A, const Point &B, const Point
      →&C) { // 内心
     double a = (B - C).len(), b = (C - A).len(), c = (A -
 2
        \hookrightarrow B).len(),
3
       s = fabs(det(B - A, C - A)),
       r = s / p;
4
     return (A * a + B * b + C * c) / (a + b + c);
6
  }
 7
   Point circumCenter(const Point &a, const Point &b, const
     → Point &c) { // 外心
     Point bb = b - a, cc = c - a;
 9
     double db = bb.len2(), dc = cc.len2(), d = 2 * det(bb,
     return a - Point(bb.y * dc - cc.y * db, cc.x * db - bb.x
        \rightarrow * dc) / d:
11
   Point othroCenter(const Point &a, const Point &b, const
     → Point &c) { // 垂心
     Point ba = b - a, ca = c - a, bc = b - c;
13
14
     double Y = ba.y * ca.y * bc.y,
15
          A = ca.x * ba.y - ba.x * ca.y,
16
          x0 = (Y + ca.x * ba.y * b.x - ba.x * ca.y * c.x) /
17
          y0 = -ba.x * (x0 - c.x) / ba.y + ca.y;
18
     return Point(x0, y0);
  }
```

3.1.5 半平面交

```
struct Point {
1
2
     int quad() const { return sign(y) == 1 || (sign(y) == 0
        \hookrightarrow && sign(x) >= 0);}
3 };
4
   struct Line {
5
     bool include(const Point &p) const { return sign(det(b -
        \hookrightarrow a, p - a)) > 0; }
     Line push() const{ // 将半平面向外推 eps
6
7
       const double eps = 1e-6;
8
       Point delta = (b - a).turn90().norm() * eps;
       return Line(a - delta, b - delta);
9
10
     }
11
  };
   bool sameDir(const Line &10, const Line &11) { return
12
      \hookrightarrow parallel(10, 11) && sign(dot(10.b - 10.a, 11.b -
      \hookrightarrow 11.a)) == 1; }
   bool operator < (const Point &a, const Point &b) {</pre>
13
    if (a.quad() != b.quad()) {
14
15
       return a.quad() < b.quad();</pre>
16
     } else {
       return sign(det(a, b)) > 0;
17
18
     }
19
20
   bool operator < (const Line &10, const Line &11) {
     if (sameDir(10, 11)) {
21
       return 11.include(10.a);
22
     } else {
23
       return (10.b - 10.a) < (11.b - 11.a);</pre>
24
25
26 }
27
   bool check(const Line &u, const Line &v, const Line &w) {
      vector<Point> intersection(vector<Line> &1) {
28
29
     sort(1.begin(), 1.end());
     deque<Line> q;
30
     for (int i = 0; i < (int)l.size(); ++i) {</pre>
31
32
       if (i && sameDir(l[i], l[i - 1])) {
33
         continue;
34
       while (q.size() > 1 && !check(q[q.size() - 2],
35
          \hookrightarrow q[q.size() - 1], l[i])) q.pop_back();
       while (q.size() > 1 && !check(q[1], q[0], 1[i]))
36
           \hookrightarrow q.pop\_front();
37
       q.push_back(l[i]);
     }
38
39
     while (q.size() > 2 && !check(q[q.size() - 2],
        \rightarrow q[q.size() - 1], q[0])) q.pop_back();
40
     while (q.size() > 2 \&\& !check(q[1], q[0], q[q.size() -
        \hookrightarrow 1])) q.pop_front();
41
     vector<Point> ret;
     for (int i = 0; i < (int)q.size(); ++i)</pre>
42
        \hookrightarrow ret.push_back(intersect(q[i], q[(i + 1) %
        \hookrightarrow q.size()]));
43
     return ret;
44 | }
```

3.1.6 最大空凸包

```
inline double eq(double x, double y) {
1
2
       return fabs(x-y)<eps;</pre>
3 }
   double xmult(point a, point b, point o) {
5
       return (a.x-o.x)*(o.y-b.y)-(a.y-o.y)*(o.x-b.x);
6
   1
7
   double dist(point a, point b) {
       return (a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y);
8
9
   point o;
11
   bool cmp_angle(point a,point b) {
       if(eg(xmult(a,b,o),0.0)) {
12
13
           return dist(a,o)<dist(b,o);</pre>
```

```
14
15
        return xmult(a,o,b)>0;
16
   }
17
   double empty_convex(point *p, int pn) {
        double ans=0:
18
        for(int i=0; i<pn; i++) {</pre>
19
20
            for(int j=0; j<pn; j++) {</pre>
21
                 dp[i][j]=0;
22
            }
        }
23
24
        for(int i=0; i<pn; i++) {</pre>
25
            int i = i-1:
26
            while(j>=0 && eq(xmult(p[i], p[j],
               \hookrightarrow o),0.0))j--;//coline
27
            bool flag= j==i-1;
28
            while(j \ge 0) {
                 int k = j-1;
29
30
                 while(k \geq 0 && xmult(p[i],p[k],p[j])>0)k--;
31
                 double area = fabs(xmult(p[i],p[j],o))/2;
                 if(k >= 0)area+=dp[j][k];
32
33
                 if(flag) dp[i][j]=area;
34
                 ans=max(ans, area);
                 j=k;
35
36
            }
37
            if(flag) {
38
                 for(int j=1; j<i; j++) {</pre>
39
                     dp[i][j] = max(dp[i][j],dp[i][j-1]);
40
            }
41
42
        }
43
        return ans;
44
45
   double largest_empty_convex(point *p, int pn) {
46
        point data[maxn];
47
        double ans=0;
48
        for(int i=0; i<pn; i++) {</pre>
49
            o=p[i];
50
            int dn=0;
51
            for(int j=0; j<pn; j++) {
52
                 if(p[j].y>o.y||(p[j].y==o.y&&p[j].x>=o.x)) {
53
                     data[dn++]=p[j];
54
55
56
            sort(data, data+dn, cmp_angle);
57
            ans=max(ans, empty_convex(data, dn));
        }
58
59
        return ans;
60
```

3.1.7 平面最近点对

```
double Dis(Point a. Point b) {
     return sqrt((a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y));
 3
   }
 4
   double Closest_Pair(int left, int right) {
5
     double d = INF;
6
     if(left == right) return d;
7
     if(left +1 == right)
8
       return Dis(p[left],p[right]);
9
     int mid = (left+right)>>1;
10
     double d1 = Closest_Pair(left,mid);
11
     double d2 = Closest_Pair(mid,right);
12
     d = min(d1,d2);
13
     int k = 0:
     for(int i = left; i <= right; i++) {</pre>
14
15
       if(fabs(p[mid].x - p[i].x) \le d)
16
         temp[k++] = p[i];
17
18
     sort(temp,temp+k,cmpy);
     for(int i = 0; i < k; i++) {</pre>
19
20
       for(int j = i+1; j < k && temp[j].y - temp[i].y < d;
          → j++) {
```

3.1.8 最小覆盖圆

```
#include<cmath>
 2
   #include<cstdio>
3
   #include<algorithm>
4 using namespace std;
5 const double eps=1e-6:
6 struct couple {
7
     double x, y;
 8
     couple(){}
9
      couple(const double &xx, const double &yy) {
10
       x = xx; y = yy;
     }
11
   } a[100001];
12
13
   int n;
   //dis means distance, dis2 means square of it
15 struct circle {
16
    double r; couple c;
17 } cir:
18 inline bool inside(const couple & x) {
19
    return di2(x, cir.c) < cir.r*cir.r+eps;</pre>
20 }
21 inline void p2c(int x, int y) {
22
    cir.c.x = (a[x].x+a[y].x)/2;
23
     cir.c.y = (a[x].y+a[y].y)/2;
24
     cir.r = dis(cir.c, a[x]);
25
26
   inline void p3c(int i, int j, int k) {
27
      couple x = a[i], y = a[j], z = a[k];
28
      cir.r =
         \rightarrow \mathsf{sqrt}(\mathsf{di2}(\mathtt{x},\mathtt{y}) * \mathsf{di2}(\mathtt{y},\mathtt{z}) * \mathsf{di2}(\mathtt{z},\mathtt{x})) / \mathsf{fabs}(\mathtt{x} * \mathtt{y} + \mathtt{y} * \mathtt{z} + \mathtt{z} * \mathtt{x}) / 2;
29
      couple t1((x-y).x, (y-z).x), t2((x-y).y, (y-z).y),
         \hookrightarrow t3((len(x)-len(y))/2, (len(y)-len(z))/2);
30
      cir.c = couple(t3*t2, t1*t3)/(t1*t2);
31 }
32 inline circle mi() {
33
     sort(a + 1, a + 1 + n);
      n = unique(a + 1, a + 1 + n) - a - 1;
34
35
     if(n == 1) {
36
        cir.c = a[1];
37
        cir.r = 0;
38
        return cir;
     }
39
     random_shuffle(a + 1, a + 1 + n);
40
41
      p2c(1, 2);
      for(int i = 3; i <= n; i++)
42
43
       if(!inside(a[i])) {
44
          p2c(1, i);
          for(int j = 2; j < i; j++)
45
46
             if(!inside(a[j])) {
47
               p2c(i, j);
               for(int k = 1; k < j; k++)
48
49
                  if(!inside(a[k]))
50
                    p3c(i,j, k);
51
        }
52
53
      return cir;
```

3.1.9 多边形内部可视

```
4 }
 5
   int Onleft(const Point & a, const Point &b, const Point &
6
     return dcmp(Cross(b - c, a - c)) > 0;
7
  }
   int visible(int x, int y) {
8
     int P = (x + n - 1) \% n, Q = (x + 1) \% n;
10
     Point u = p[y] - p[x], v = p[x] - p[P], w = p[x] - p[Q];
11
     if (Onleft(p[Q], p[x], p[P])) {
       return dcmp(Cross(v, u)) > 0 && dcmp(Cross(w, u)) < 0;
12
13
     } else {
14
       return !(dcmp(Cross(v, u)) < 0 && dcmp(Cross(w, u)) >
15
16
   }
17
   int solve(int x, int y) {
     if (vis[x][y] == dfn) return g[x][y];
18
19
     vis[x][y] = dfn;
     if (x == y | | y == x + 1) return g[x][y] = 1;
     for (int i = x; i + 1 <= y; i++) {
21
22
       if (C(p[x], p[y], p[i], p[i + 1])) return g[x][y] = 0;
23
     for (int i = x + 1; i < y; i++) {
24
25
       if (OnLine(p[x], p[i], p[y])) {
26
         return g[x][y] = solve(x, i) && solve(i, y);
27
28
     if (!visible(x, y) || !visible(y, x)) return g[x][y] =
29
        → 0:
30
     return g[x][y] = 1;
31
```

3.2 三维

3.2.1 三维点类

```
1 // 三维绕轴旋转,大拇指指向 axis 向量方向,四指弯曲
     →方向转 w 弧度
  Point rotate(const Point& s, const Point& axis, DB w) {
     DB x = axis.x, y = axis.y, z = axis.z;
 3
     DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
4
        cosw = cos(w), sinw = sin(w);
6
    DB a[4][4];
7
     memset(a, 0, sizeof a);
8
     a[3][3] = 1;
9
     a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
10
     a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
11
     a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
12
     a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
13
     a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
14
     a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
     a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
15
     a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
16
17
     a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
     DB ans [4] = \{0, 0, 0, 0\}, c[4] = \{s.x, s.y, s.z, 1\};
18
19
     for (int i = 0; i < 4; ++ i)
20
       for (int j = 0; j < 4; ++ j)
21
         ans[i] += a[j][i] * c[j];
22
    return Point(ans[0], ans[1], ans[2]);
23
```

3.2.2 凸包

```
10 }
   __inline DB volume(const P& a, const P& b, const P& c,
11
      12
     return mix(b - a, c - a, d - a);
13 }
   struct Face {
14
15
    int a, b, c;
     __inline Face() {}
16
17
     __inline Face(int _a, int _b, int _c):
18
       a(_a), b(_b), c(_c) {}
     __inline DB area() const {
19
20
       return 0.5 * cross(p[b] - p[a], p[c] - p[a]).len();
21
     __inline P normal() const {
22
23
       return cross(p[b] - p[a], p[c] - p[a]).unit();
24
     __inline DB dis(const P& p0) const {
25
26
       return dot(normal(), p0 - p[a]);
27
     }
28 }:
29
   std::vector<Face> face, tmp; // Should be O(n).
30 int mark[N][N], Time, n;
   __inline void add(int v) {
31
32
     ++ Time:
33
     clear(tmp);
34
     for (int i = 0; i < (int)face.size(); ++ i) {</pre>
35
       int a = face[i].a, b = face[i].b, c = face[i].c;
       if (sign(volume(p[v], p[a], p[b], p[c])) > 0) {
36
         mark[a][b] = mark[b][a] = mark[a][c] =
37
           mark[c][a] = mark[b][c] = mark[c][b] = Time;
38
39
         tmp.push_back(face[i]);
40
41
       }
42
     }
43
     clear(face); face = tmp;
     for (int i = 0; i < (int)tmp.size(); ++ i) {</pre>
44
45
       int a = face[i].a, b = face[i].b, c = face[i].c;
       if (mark[a][b] == Time) face.emplace_back(v, b, a);
46
47
       if (mark[b][c] == Time) face.emplace_back(v, c, b);
       if (mark[c][a] == Time) face.emplace_back(v, a, c);
48
49
       assert(face.size() < 500u);</pre>
50
     }
51 }
   void reorder() {
     for (int i = 2; i < n; ++ i) {
53
54
       P \text{ tmp} = cross(p[i] - p[0], p[i] - p[1]);
55
       if (sign(tmp.len())) {
56
         std::swap(p[i], p[2]);
57
         for (int j = 3; j < n; ++ j)
58
           if (sign(volume(p[0], p[1], p[2], p[j]))) {
59
              std::swap(p[j], p[3]);
60
              return;
61
           }
       }
62
63
     }
64
  |}
   void build_convex() {
65
66
     reorder();
67
     clear(face);
     face.emplace_back(0, 1, 2);
68
     face.emplace_back(0, 2, 1);
69
70
     for (int i = 3; i < n; ++ i)
71
       add(i);
72 | }
```

3.2.3 最小覆盖球

```
const int eps = 1e-8;
truct Tpoint {
   double x, y, z;
};
int npoint, nouter;
Tpoint pt[200000], outer[4],res;
```

```
7 double radius,tmp;
   inline double dist(Tpoint p1, Tpoint p2) {
 8
     double dx=p1.x-p2.x, dy=p1.y-p2.y, dz=p1.z-p2.z;
10
     return ( dx*dx + dy*dy + dz*dz );
11
   7
12
   inline double dot(Tpoint p1, Tpoint p2) {
13
     return p1.x*p2.x + p1.y*p2.y + p1.z*p2.z;
14
   }
15
   void ball() {
16
     Tpoint q[3]; double m[3][3], sol[3], L[3], det;
17
     int i,j;
18
     res.x = res.v = res.z = radius = 0:
19
     switch ( nouter ) {
20
       case 1: res=outer[0]; break;
21
       case 2:
           res.x=(outer[0].x+outer[1].x)/2;
22
           res.y=(outer[0].y+outer[1].y)/2;
23
24
           res.z=(outer[0].z+outer[1].z)/2;
25
            radius=dist(res, outer[0]);
26
27
       case 3:
           for (i=0; i<2; ++i ) {
28
29
              q[i].x=outer[i+1].x-outer[0].x;
30
              q[i].y=outer[i+1].y-outer[0].y;
31
              q[i].z=outer[i+1].z-outer[0].z;
32
33
            for (i=0; i<2; ++i) for(j=0; j<2; ++j)
34
              m[i][j]=dot(q[i], q[j])*2;
            for (i=0; i<2; ++i ) sol[i]=dot(q[i], q[i]);</pre>
35
            if (fabs(det=m[0][0]*m[1][1]-m[0][1]*m[1][0])<eps)</pre>
36
37
38
           L[0]=(sol[0]*m[1][1]-sol[1]*m[0][1])/det;
39
           L[1]=(sol[1]*m[0][0]-sol[0]*m[1][0])/det;
40
           res.x=outer[0].x+q[0].x*L[0]+q[1].x*L[1];
41
           res.y=outer[0].y+q[0].y*L[0]+q[1].y*L[1];
42
            res.z=outer[0].z+q[0].z*L[0]+q[1].z*L[1];
43
            radius=dist(res, outer[0]);
           break;
44
45
           for (i=0; i<3; ++i) {
46
47
              q[i].x=outer[i+1].x-outer[0].x;
              q[i].y=outer[i+1].y-outer[0].y;
48
49
              q[i].z=outer[i+1].z-outer[0].z;
50
              sol[i]=dot(q[i], q[i]);
51
           }
52
           for (i=0;i<3;++i)</pre>
53
              for(j=0;j<3;++j) m[i][j]=dot(q[i],q[j])*2;</pre>
54
            det= m[0][0]*m[1][1]*m[2][2]
55
              + m[0][1]*m[1][2]*m[2][0]
56
              + m[0][2]*m[2][1]*m[1][0]
57
              - m[0][2]*m[1][1]*m[2][0]
58
              - m[0][1]*m[1][0]*m[2][2]
59
               - m[0][0]*m[1][2]*m[2][1];
            if ( fabs(det) < eps ) return;</pre>
60
61
            for (j=0; j<3; ++j) {
62
              for (i=0; i<3; ++i) m[i][j]=sol[i];
              L[j]=(m[0][0]*m[1][1]*m[2][2]
63
64
                  + m[0][1]*m[1][2]*m[2][0]
65
                  + m[0][2]*m[2][1]*m[1][0]
                  - m[0][2]*m[1][1]*m[2][0]
66
67
                  - m[0][1]*m[1][0]*m[2][2]
68
                  - m[0][0]*m[1][2]*m[2][1]
69
                 ) / det;
70
              for (i=0; i<3; ++i)</pre>
71
                m[i][j]=dot(q[i], q[j])*2;
72
73
           res=outer[0];
            for (i=0; i<3; ++i ) {
74
75
             res.x += q[i].x * L[i];
76
              res.y += q[i].y * L[i];
77
              res.z += q[i].z * L[i];
78
79
            radius=dist(res, outer[0]);
```

```
80
81
   }
82
    void minball(int n) {
83
      ball():
      if ( nouter<4 )</pre>
84
        for (int i=0; i<n; ++i)
85
           if (dist(res, pt[i])-radius>eps) {
             outer[nouter]=pt[i];
87
88
             ++nouter:
89
             minball(i);
90
             --nouter:
             if (i>0) {
91
92
               Tpoint Tt = pt[i];
93
               memmove(&pt[1], &pt[0], sizeof(Tpoint)*i);
94
               pt[0]=Tt;
             }
95
          }
96
97 }
98
    void solve() {
      for (int i=0;i<npoint;i++)</pre>
         \rightarrow scanf("%lf%lf",&pt[i].x,&pt[i].y,&pt[i].z);
100
      random_shuffle(pt, pt + npoint);
101
      radius=-1:
102
      for (int i=0;i<npoint;i++){</pre>
103
        if (dist(res,pt[i])-radius>eps){
104
          nouter=1;
105
           outer[0]=pt[i];
          minball(i);
106
107
      }
108
109
      printf("%.5f\n",sqrt(radius));
110 }
111
    int main(){
112
      for( ; cin >> npoint && npoint; )
113
        solve():
114
      return 0;
115 }
```

4. 字符串

4.1 AC 自动机

```
1 int newnode() {
2
     ++tot;
3
     memset(ch[tot], 0, sizeof(ch[tot]));
     fail[tot] = 0;
5
     dep[tot] = 0;
6
     par[tot] = 0;
     return tot;
7
8 }
   void insert(char *s,int x) {
10
    if(*s == '\0') return;
11
     else {
12
       int &y = ch[x][*s - 'a'];
13
       if(y == 0) {
         y = newnode();
14
15
         par[y] = x;
16
         dep[y] = dep[x] + 1;
17
18
       insert(s + 1, y);
     }
19
20 | }
   void build() {
    int line[maxn];
22
23
     int f = 0, r = 0;
     fail[root] = root;
24
     for(int i = 0; i < alpha; i++) {</pre>
25
       if(ch[root][i]) {
26
         fail[ch[root][i]] = root;
         line[r++] = ch[root][i];
28
29
       } else {
         ch[root][i] = root;
30
31
```

```
32
33
     while(f != r) {
34
       int x = line[f++];
        for(int i = 0; i < alpha; i++) {</pre>
35
36
          if(ch[x][i]) {
            fail[ch[x][i]] = ch[fail[x]][i];
37
            line[r++] = ch[x][i];
39
          } else {
40
            ch[x][i] = ch[fail[x]][i];
41
42
43
     }
44
```

4.2 后缀数组

```
const int MAXN = MAXL * 2 + 1;
   int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN],
      \hookrightarrow rank[MAXN], height[MAXN];
 3
   void calc_sa(int n) {
     int m = alphabet, k = 1;
4
5
     memset(c, 0, sizeof(*c) * (m + 1));
6
     for (int i = 1; i <= n; ++i) c[x[i] = a[i]]++;
7
     for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
8
     for (int i = n; i; --i) sa[c[x[i]]--] = i;
9
     for (; k <= n; k <<= 1) {
10
       int tot = k;
11
       for (int i = n - k + 1; i \le n; ++i) y[i - n + k] = i;
12
       for (int i = 1; i <= n; ++i)
         if (sa[i] > k) y[++tot] = sa[i] - k;
13
       memset(c, 0, sizeof(*c) * (m + 1));
       for (int i = 1; i \le n; ++i) c[x[i]]++;
16
       for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
17
       for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];
18
       for (int i = 1; i \le n; ++i) y[i] = x[i];
19
       tot = 1; x[sa[1]] = 1;
20
       for (int i = 2; i <= n; ++i) {
         if (max(sa[i], sa[i - 1]) + k > n || y[sa[i]] !=
             \hookrightarrow y[sa[i - 1]] || y[sa[i] + k] != y[sa[i - 1] +
             \hookrightarrow kl) ++tot:
22
         x[sa[i]] = tot;
23
24
       if (tot == n) break; else m = tot;
25
     }
26
27
   void calc_height(int n) {
     for (int i = 1; i <= n; ++i) rank[sa[i]] = i;</pre>
28
     for (int i = 1; i <= n; ++i) {
30
       height[rank[i]] = max(0, height[rank[i - 1]] - 1);
       if (rank[i] == 1) continue;
31
       int j = sa[rank[i] - 1];
32
       while (max(i, j) + height[rank[i]] <= n && a[i +</pre>
33

    height[rank[i]]] == a[j + height[rank[i]]])

    ++height[rank[i]];

34
     }
35
```

4.3 后缀自动机

```
memset(par, 0, sizeof(*par) * (1 * 2 + 1));
8
     memset(mxl, 0, sizeof(*mxl) * (1 * 2 + 1));
9
10
     memset(size, 0, sizeof(*size) * (1 * 2 + 1));
11 | }
  inline void extend(int pos, int c) {
12
     int p = last, np = last = ++cnt;
13
     mxl[np] = mxl[p] + 1; size[np] = 1;
     for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
15
16
     if (!p) par[np] = 1;
17
     else {
18
       int q = trans[p][c];
19
       if (mxl[p] + 1 == mxl[q]) par[np] = q;
20
       else {
21
         int ng = ++cnt;
         mxl[nq] = mxl[p] + 1;
22
23
         memcpy(trans[nq], trans[q], sizeof(trans[nq]));
         par[nq] = par[q];
24
25
         par[np] = par[q] = nq;
26
         for (; trans[p][c] == q; p = par[p]) trans[p][c] =
27
       }
     }
28
29
  | }
   inline void buildsam() {
30
31
     for (int i = 1; i <= 1; ++i) extend(i, str[i] - 'a');</pre>
32
     memset(sum, 0, sizeof(*sum) * (1 * 2 + 1));
33
     for (int i = 1; i <= cnt; ++i) sum[mxl[i]]++;</pre>
     for (int i = 1; i <= 1; ++i) sum[i] += sum[i - 1];
34
     for (int i = cnt; i; --i) seq[sum[mxl[i]]--] = i;
35
     for (int i = cnt; i; --i) size[par[seq[i]]] +=
36
        \hookrightarrow size[seq[i]];
37 }
```

4.4 广义后缀自动机

```
inline void add_node(int x, int &last) {
2
     int lastnode = last;
3
     if (c[lastnode][x]) {
       int nownode = c[lastnode][x];
5
       if (l[nownode] == l[lastnode] + 1) last = nownode;
6
7
         int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
         for (int i = 0; i < alphabet; ++i) c[auxnode][i] =</pre>
8
            \hookrightarrow c[nownode][i];
9
         par[auxnode] = par[nownode]; par[nownode] = auxnode;
10
         for (; lastnode && c[lastnode][x] == nownode;
             \hookrightarrow lastnode = par[lastnode]) {
            c[lastnode][x] = auxnode;
11
         }
12
13
         last = auxnode;
14
       }
15
     } else {
16
       int newnode = ++cnt; l[newnode] = l[lastnode] + 1;
17
       for (; lastnode && !c[lastnode][x]; lastnode =

    par[lastnode]) c[lastnode][x] = newnode:
18
       if (!lastnode) par[newnode] = 1;
19
         int nownode = c[lastnode][x];
20
21
         if (l[lastnode] + 1 == l[nownode]) par[newnode] =

→ nownode:

22
         else {
23
            int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
            for (int i = 0; i < alphabet; ++i) c[auxnode][i] =</pre>

    c [nownode] [i];

25
            par[auxnode] = par[nownode]; par[nownode] =
              → par[newnode] = auxnode;
            for (; lastnode && c[lastnode][x] == nownode;
26
               → lastnode = par[lastnode]) {
              c[lastnode][x] = auxnode;
27
28
            }
29
         }
30
31
       last = newnode:
```

```
32 }
33 }
```

4.5 manacher

```
void Manacher(std::string s,int p[]) {
       string t = "$#";
3
       for (int i = 0; i < s.size(); i++) {
4
            t += s[i];
            t += "#";
5
6
7
       std::vector<int> p(t.size(), 0);
8
       int mx = 0, id = 0;
9
       for (int i = 1; i < t.size(); i++) {</pre>
10
           p[i] = mx > i ? min(p[2 * id - i], mx - i) : 1;
            while (t[i + p[i]] == t[i - p[i]]) ++p[i];
11
12
            if (mx < i + p[i]) {</pre>
13
                mx = i + p[i];
14
                id = i;
15
           }
16
17
```

4.6 回文自动机

```
int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN],
      \hookrightarrow 1[MAXN], s[MAXN];
   int allocate(int len) {
     l[nT] = len:
     r[nT] = 0;
     fail[nT] = 0;
     memset(c[nT], 0, sizeof(c[nT]));
 6
 7
     return nT++:
8
9
   void init() {
     nT = nStr = 0;
11
     int newE = allocate(0);
12
     int new0 = allocate(-1);
     last = newE:
13
     fail[newE] = newO;
14
15
     fail[new0] = newE;
16
     s[0] = -1;
17
   }
18
   void add(int x) {
     s[++nStr] = x;
19
     int now = last;
20
21
     while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now];
22
23
       int newnode = allocate(l[now] + 2), &newfail =
          newfail = fail[now]:
24
       while (s[nStr - l[newfail] - 1] != s[nStr]) newfail =
25

    fail[newfail];
       newfail = c[newfail][x];
26
27
       c[now][x] = newnode;
28
29
     last = c[now][x]:
30
     r[last]++;
31
   void count() {
32
33
     for (int i = nT - 1; i \ge 0; i--) {
       r[fail[i]] += r[i];
34
35
36 }
```

4.7 循环串的最小表示

```
1 // n 必须是 2 的次幂
2 void fft(Complex a[], int n, int f) {
  for (int i = 0; i < n; ++i)
```

```
if (R[i] < i) swap(a[i], a[R[i]]);</pre>
4
5
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
       Complex w = Complex(1, 0);
7
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
8
9
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
         for (int k = 0; k < i; ++k) {
10
11
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
       }
15
     }
16
```

5. 数据结构

5.1 可并堆

```
1 int merge(int x,int y) {
  //p[i] 结点 i 的权值,这里是维护大根堆
  //d[i] 在 i 的子树中, i 到右叶子结点的最远距离.
      if(!x) return y;
      if(!y) return x;
5
      if(p[x] < p[y]) std::swap(x, y);
6
7
      r[x] = merge(r[x], y);
8
      if(r[x]) fa[r[x]] = x;
      if(d[l[x]] < d[r[x]]) std::swap(l[x], r[x]);//调整树
9
        → 的结构, 使其满足左偏性质
      d[x] = d[r[x]] + 1;
10
11
      return x;
12 }
```

5.2 KD-Tree

```
long long norm(const long long &x) {
1
2
       return std::abs(x);
3
       return x * x;
4 }
5
   struct Point {
6
       int x, y, id;
7
       const int& operator [] (int index) const {
8
           if (index == 0) {
9
                return x;
            } else {
11
                return y;
12
       }
13
       friend long long dist(const Point &a, const Point &b)
14
15
            long long result = 0;
16
            for (int i = 0; i < 2; ++i) {
17
                result += norm(a[i] - b[i]);
            }
18
19
            return result:
20
   } point[N];
21
   struct Rectangle {
23
       int min[2], max[2];
24
       Rectangle() {
25
            min[0] = min[1] = INT_MAX; // sometimes int is
               \hookrightarrow \mathtt{not} enough
            max[0] = max[1] = INT_MIN;
26
       }
27
       void add(const Point &p) {
28
           for (int i = 0; i < 2; ++i) {
29
                min[i] = std::min(min[i], p[i]);
30
31
                max[i] = std::max(max[i], p[i]);
32
33
34
       long long dist(const Point &p) {
            long long result = 0;
35
            for (int i = 0; i < 2; ++i) {
36
```

```
result += norm(std::min(std::max(p[i],
37
                    38
                 result += std::max(norm(max[i] - p[i]),
                   \hookrightarrow \texttt{norm}(\texttt{min[i] - p[i]));}
            }
39
40
            return result:
41
42
   };
43
    struct Node {
44
        Point seperator;
45
        Rectangle rectangle;
46
        int child[2];
47
        void reset(const Point &p) {
48
            seperator = p;
49
            rectangle = Rectangle();
50
            rectangle.add(p);
51
            child[0] = child[1] = 0;
        }
52
   } tree[N << 1];
    int size, pivot;
    bool compare(const Point &a, const Point &b) {
55
        if (a[pivot] != b[pivot]) {
56
57
            return a[pivot] < b[pivot];</pre>
58
59
        return a.id < b.id;
60
    // 左閉右開: build(1, n + 1)
    int build(int 1, int r, int type = 1) {
        pivot = type;
63
        if (1 >= r) {
64
65
            return 0:
66
        int x = ++size;
67
68
        int mid = 1 + r >> 1;
69
        std::nth_element(point + 1, point + mid, point + r,

→ compare);
70
        tree[x].reset(point[mid]);
71
        for (int i = 1; i < r; ++i) {
72
             tree[x].rectangle.add(point[i]);
73
        tree[x].child[0] = build(1, mid, type ^ 1);
74
75
        tree[x].child[1] = build(mid + 1, r, type ^ 1);
76
        return x;
77
   }
78
    int insert(int x, const Point &p, int type = 1) {
79
        pivot = type;
80
        if (x == 0) {
81
            tree[++size].reset(p);
82
            return size:
        tree[x].rectangle.add(p);
        if (compare(p, tree[x].seperator)) {
85
86
            tree[x].child[0] = insert(tree[x].child[0], p,
               \hookrightarrow type ^ 1);
87
        } else {
88
            tree[x].child[1] = insert(tree[x].child[1], p,
               \hookrightarrow type ^ 1);
89
        }
90
        return x;
91
    // For minimum distance
92
    // For maximum: 下面递归 query 时 0, 1 换顺序;< and
       \leftrightarrow >:min and max
    void query(int x, const Point &p, std::pair<long long,</pre>
       \hookrightarrow int> &answer, int type = 1) {
95
        pivot = type;
        if (x == 0 || tree[x].rectangle.dist(p) >
96
           → answer.first) {
97
            return:
98
99
        answer = std::min(answer,
100
                  std::make_pair(dist(tree[x].seperator, p),
```

```
101
       if (compare(p, tree[x].seperator)) {
           query(tree[x].child[0], p, answer, type ^ 1);
102
103
           query(tree[x].child[1], p, answer, type ^ 1);
104
       } else {
105
           query(tree[x].child[1], p, answer, type ^ 1);
           query(tree[x].child[0], p, answer, type ^ 1);
106
107
   }
108
109
   std::priority_queue<std::pair<long long, int> > answer;
110
   void query(int x, const Point &p, int k, int type = 1) {
111
       pivot = type;
       if (x == 0 || (int)answer.size() == k &&
112
          113
114
115
       answer.push(std::make_pair(dist(tree[x].seperator, p),
          116
       if ((int)answer.size() > k) {
117
           answer.pop();
118
119
       if (compare(p, tree[x].seperator)) {
           query(tree[x].child[0], p, k, type ^{1});
120
           query(tree[x].child[1], p, k, type ^ 1);
121
122
       } else {
123
           query(tree[x].child[1], p, k, type ^ 1);
124
           query(tree[x].child[0], p, k, type ^ 1);
125
126 | }
```

5.3 Treap

```
1
   struct Node{
 2
      int mn, key, size, tag;
 3
 4
      Node* ch[2];
      Node(int mn, int key, int size): mn(mn), key(key),
 5
        \hookrightarrow \text{size}(\text{size}), \text{ rev}(0), \text{ tag}(0)\{\}
 6
      void downtag();
      Node* update(){
 8
        mn = min(ch[0] \rightarrow mn, min(key, ch[1] \rightarrow mn));
 9
        size = ch[0] \rightarrow size + 1 + ch[1] \rightarrow size;
10
        return this;
     }
11
12
   };
13
   typedef pair<Node*, Node*> Pair;
   Node *null, *root;
15
   void Node::downtag(){
16
     if(rev){
        for(int i = 0; i < 2; i++)
17
          if(ch[i] != null){
18
            ch[i] -> rev ^= 1;
19
20
            swap(ch[i] -> ch[0], ch[i] -> ch[1]);
21
          }
22
        rev = 0;
     }
23
24
      if(tag){
25
       for(int i = 0; i < 2; i++)
26
          if(ch[i] != null){
27
            ch[i] -> key += tag;
            ch[i] -> mn += tag;
28
29
            ch[i] -> tag += tag;
          }
30
        tag = 0;
31
32
     }
33 }
34 | int r(){
     static int s = 3023192386;
35
36
     return (s += (s << 3) + 1) & (^{\circ}0u >> 1);
37
   bool random(int x, int y){
38
39
     return r() % (x + y) < x;
40 }
41 | Node* merge(Node *p, Node *q){
```

```
42
      if(p == null) return q;
43
      if(q == null) return p;
44
      p -> downtag();
45
      q -> downtag();
      if(random(p \rightarrow size, q \rightarrow size)){}
46
47
        p -> ch[1] = merge(p -> ch[1], q);
        return p -> update();
48
49
50
        q \rightarrow ch[0] = merge(p, q \rightarrow ch[0]);
51
        return q -> update();
52
53
54
   Pair split(Node *x, int n){
      if(x == null) return make_pair(null, null);
55
56
      x -> downtag();
      if(n \le x \rightarrow ch[0] \rightarrow size){
57
        Pair ret = split(x -> ch[0], n);
58
59
        x \rightarrow ch[0] = ret.second;
60
        return make_pair(ret.first, x -> update());
61
62
      Pair ret = split(x \rightarrow ch[1], n - x \rightarrow ch[0] \rightarrow size -
         \hookrightarrow 1);
      x \rightarrow ch[1] = ret.first;
63
      return make_pair(x -> update(), ret.second);
64
65
66
   pair<Node*, Pair> get_segment(int 1, int r){
67
      Pair ret = split(root, l - 1);
68
      return make_pair(ret.first, split(ret.second, r - 1 +
         \hookrightarrow 1)):
69
   }
    int main(){
      null = new Node(INF, INF, 0);
72
      null \rightarrow ch[0] = null \rightarrow ch[1] = null;
73
      root = null;
74
```

5.4 Splay

```
template<class T>void checkmin(T &x,T y) {
2
     if(y < x) x = y;
3
   }
4
   struct Node {
5
     Node *c[2], *fa;
6
     int size, rev;
     LL val, add, min;
8
     Node *init(LL v) {
       val = min = v;
9
       add = rev = 0;
10
       c[0] = c[1] = fa = NULL;
11
12
       size = 1:
13
       return this;
14
     }
15
     void rvs() {
       std::swap(c[0], c[1]);
16
17
       rev ^= 1;
18
19
     void inc(LL x) {
       val += x;
20
21
       add += x;
22
       min += x;
23
24
     void pushdown() {
       if(rev) {
25
26
         if(c[0]) c[0]->rvs();
         if(c[1]) c[1]->rvs();
27
28
         rev = 0:
29
30
       if(add) {
         if(c[0]) c[0]->inc(add);
31
32
         if(c[1]) c[1]->inc(add);
         add = 0;
33
34
```

```
}
35
36
      void update() {
37
        min = val;
38
        if(c[0]) checkmin(min, c[0]->min);
        if(c[1]) checkmin(min, c[1]->min);
39
40
        size = 1:
        if(c[0]) size += c[0]->size;
        if(c[1]) size += c[1]->size;
42
43
     }
44 | } *root;
45 Node* newnode(LL x) {
46
      static Node pool[maxs], *p = pool;
47
      return (++p)->init(x);
48 }
49
    void setc(Node *x,int t,Node *y) {
50
      x->c[t] = y;
      if(y) y->fa = x;
51
52 }
53 Node *find(int k) {
     Node *now = root:
55
     while(true) {
56
       now->pushdown();
        int t = (now->c[0] ? now->c[0]->size : 0) + 1;
57
        if(t == k) break;
58
59
        if(t > k) now = now->c[0];
60
        else now = now->c[1], k \rightarrow t;
61
     }
62
      return now;
63 }
64
   void rotate(Node *x,Node* &k) {
65
     Node *y = x-fa, *z = y-fa;
      if(y != k) z->c[z->c[1] == y] = x;
      else k = x;
67
68
      x->fa = z;
      int i = (y->c[1] == x);
69
70
      setc(y, i, x->c[i ^ 1]);
71
      setc(x, i ^ 1, y);
72
      y->update(), x->update();
73 }
74
    void spaly(Node *x,Node* &k) {
75
     static Node *st[maxs]:
76
     int top = 0;
77
      Node *y, *z;
78
      y = x;
79
      while(y != k) st[++top] = y, y = y->fa;
80
      st[++top] = y;
81
      while(top) st[top]->pushdown(), top--;
82
      while(x != k) {
83
        y = x-fa, z = y-fa;
        if(y != k) {
85
          if((y == z-c[1]) ^ (x == y-c[1])) rotate(x, k);
86
          else rotate(y, k);
87
88
        rotate(x, k);
89
90 }
   Node *subtree(int 1,int r) {
91
92
     assert((++1) <= (++r));
93
      spaly(find(l - 1), root);
94
      spaly(find(r + 1), root->c[1]);
95
      return root->c[1]->c[0];
96
97
    void ins(int pos,int v) {
98
      spaly(find(pos), root);
99
      spaly(find(pos + 1), root->c[1]);
100
      setc(root->c[1], 0, newnode(v));
101
      root->c[1]->update();
102
103
      root->update();
104 }
105 | void del(int pos) {
106
     pos++:
     spaly(find(pos - 1), root);
107
```

```
108
      spaly(find(pos + 1), root->c[1]);
      root->c[1]->c[0] = NULL;
109
110
      root->c[1]->update();
111
      root->update();
112
   }
    void init() {
113
      root = newnode(0);
115
      setc(root, 1, newnode(0));
116
      root->update();
117
```

5.5 Link cut Tree

```
inline void reverse(int x) {
2
    tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
3
  }
  inline void rotate(int x, int k) {
4
5
    int y = tr[x].fa, z = tr[y].fa;
       tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;
6
       tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^
 7
8
       tr[x].c[k ^ 1] = y; tr[y].fa = x;
9
10
   inline void splay(int x, int w) {
11
     int z = x; pushdown(x);
     while (tr[x].fa != w) {
12
       int y = tr[x].fa; z = tr[y].fa;
13
14
       if (z == w) {
         pushdown(z = y); pushdown(x);
15
16
         rotate(x, tr[y].c[1] == x);
17
         update(y); update(x);
18
       } else {
19
         pushdown(z); pushdown(y); pushdown(x);
20
         int t1 = tr[y].c[1] == x, t2 = tr[z].c[1] == y;
         if (t1 == t2) rotate(y, t2), rotate(x, t1);
21
22
         else rotate(x, t1), rotate(x, t2);
23
         update(z); update(y); update(x);
24
       }
25
26
     update(x);
27
     if (x != z) par[x] = par[z], par[z] = 0;
28 }
29
   inline void access(int x) {
30
    for (int y = 0; x; y = x, x = par[x]) {
31
       splay(x, 0);
32
       if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa
33
       tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
34
  }
35
   inline void makeroot(int x) {
36
37
    access(x); splay(x, 0); reverse(x);
38
39
   inline void link(int x, int y) {
40
    makeroot(x); par[x] = y;
41
  }
   inline void cut(int x, int y) {
42
43
     access(x); splay(y, 0);
44
     if (par[y] != x) swap(x, y), access(x), splay(y, 0);
45
     par[y] = 0;
46
   47
     \hookrightarrow of the tree
    makeroot(y); access(x); splay(x, 0);
```

5.6 树上莫队

```
void dfs(int u) {
  dep[u] = dep[fa[u][0]] + 1;
  for(int i = 1; i < logn; i++)
  fa[u][i] = fa[fa[u][i - 1]][i - 1];</pre>
```

```
5
     stk.push(u);
6
     for(int i = 0; i < vec[u].size(); i++) {</pre>
7
       int v = vec[u][i];
       if(v == fa[u][0]) continue;
8
9
       fa[v][0] = u, dfs(v);
       size[u] += size[v];
10
       if(size[u] >= bufsize) {
11
12
13
         while(stk.top() != u) {
14
           block[stk.top()] = bcnt;
15
           stk.pop();
16
17
         size[u] = 0;
18
19
     }
20
     size[u]++;
21 }
22
   void prework() {
23
     dfs(1):
24
     ++bcnt:
25
     while(!stk.empty()) {
26
       block[stk.top()] = bcnt;
       stk.pop();
27
28
29
30
   void rev(int u) {
31
     now -= (cnt[val[u]] > 0);
     if(used[u]) {
32
       cnt[val[u]]--:
33
       used[u] = false;
34
35
     } else {
       cnt[val[u]]++;
36
37
       used[u] = true;
38
     now += (cnt[val[u]] > 0);
39
40
  }
41
   void move(int &x,int y,int z) {
     int fwd = y;
42
43
     rev(getlca(x, z));
44
     rev(getlca(y, z));
     while(x != y) {
45
       if(dep[x] < dep[y]) std::swap(x, y);</pre>
46
47
       rev(x), x = fa[x][0];
     }
48
49
     x = fwd;
50 }
51
   void solve() {
52
     std::sort(query + 1, query + m + 1);
53
     int L = 1, R = 1;
     rev(1);
55
     for(int i = 1; i <= m; i++) {
56
       int 1 = query[i].u;
57
       int r = query[i].v;
       move(L, 1, R);
58
59
       move(R, r, L);
       ans[query[i].t] = now;
60
61
62 }
```

5.7 CDQ 分治

```
1 struct Node {
2
     int x, y, z, idx;
     friend bool operator == (const Node &a,const Node &b) {
       return a.x == b.x && a.y == b.y && a.z == b.z;
     }
5
6
     friend bool operator < (const Node &a,const Node &b) {
7
       return a.y < b.y;</pre>
8
   } triple[maxn];
9
10
  | bool cmpx(const Node &a,const Node &b) {
     if(a.x != b.x) return a.x < b.x;</pre>
11
12
    if(a.y != b.y) return a.y < b.y;</pre>
```

```
13
     return a.z < b.z:
14
15
   void solve(int l,int r) {
     if(1 == r) return;
16
     int mid = (1 + r) >> 1;
17
     solve(1, mid);
18
     static std::pair<Node,int> Lt[maxn], Rt[maxn];
     int Ls = 0, Rs = 0;
20
21
     for(int i = 1; i <= mid; i++)</pre>
22
       Lt[++Ls] = std::make_pair(triple[i], i);
23
     for(int i = mid + 1; i <= r; i++)</pre>
24
       Rt[++Rs] = std::make_pair(triple[i], i);
25
     int pos = 1;
26
     std::sort(Lt + 1, Lt + Ls + 1);
27
     std::sort(Rt + 1, Rt + Rs + 1);
28
     backup.clear();
     for(int i = 1; i <= Rs; i++) {
29
30
       while(pos <= Ls && !(Rt[i].first < Lt[pos].first)) {</pre>
31
         insert(Lt[pos].first.z, 1);
32
33
       f[Rt[i].second] += query(Rt[i].first.z);
34
35
36
     for(int i = 0; i < backup.size(); i++) pre[backup[i]] =</pre>
37
     solve(mid + 1, r);
38
```

5.8 整体二分

```
void solve(int l,int r,std::vector<int> q) {
     if(1 == r || q.empty()) {
3
        for(int i = 0; i < q.size(); i++) {</pre>
4
          ans[q[i]] = 1;
5
6
     } else {
7
       int mid = (1 + r) >> 1;
8
        backup.clear();
        for(int i = 1; i <= mid; i++) {</pre>
9
10
          Event e = event[i];
11
          if(e.1 <= e.r) {
            add(e.1, e.v);
12
13
           add(e.r + 1, -e.v);
14
          } else {
15
            add(1, e.v);
            add(e.r + 1, -e.v);
16
17
            add(e.1, e.v);
          }
18
19
20
        std::vector<int> qL, qR;
21
        for(int i = 0; i < q.size(); i++) {</pre>
22
          LL val = 0;
23
          for(int j = 0; j < vec[q[i]].size(); j++) {</pre>
24
            val += count(vec[q[i]][j]);
25
            if(val >= p[q[i]]) break;
26
27
          if(cnt[q[i]] + val >= p[q[i]]) {
28
            qL.push_back(q[i]);
29
          } else {
30
            cnt[q[i]] += val;
31
            qR.push_back(q[i]);
32
33
        for(int i = 0; i < backup.size(); i++) sum[backup[i]]</pre>
35
        solve(1, mid, qL);
36
        solve(mid + 1, r, qR);
37
38
   }
```

6. 图论

6.1 2-SAT tarjan

```
template<class TAT>void checkmin(TAT &x,TAT y) {
2
     if(y < x) x = y;
3 | }
4 void tarjan(int u) {
    dfn[u] = low[u] = ++dt;
     flag[u] = true;
7
     stk.push(u);
8
     for(int i = 0; i < vec[u].size(); i++) {</pre>
Q
       int v = vec[u][i];
10
       if(!dfn[v]) {
11
         tarian(v):
12
         checkmin(low[u], low[v]);
13
14
       else if(flag[v]) {
15
         checkmin(low[u], dfn[v]);
16
17
     }
18
     if(low[u] == dfn[u]) {
19
20
       while(stk.top() != u) {
21
         block[stk.top()] = bcnt;
         flag[stk.top()] = false;
22
23
         stk.pop();
24
25
       block[u] = bcnt;
26
       flag[u] = false;
       stk.pop();
28
29 }
30
   bool solve() {
       for(int i = 1; i <= 2 * n; i++)
31
32
         if(!dfn[i]) tarjan(i);
33
       bool ans = true;
34
       for(int i = 1; i <= n; i++)
         if(block[2 * i] == block[2 * i - 1]) {
35
           ans = false:
36
37
           break;
38
         }
39
       return ans;
40 }
```

6.2 KM

```
struct KM {
1
2
    // Truly O(n^3)
    // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时
3
       →边权取负,但不能连的还是 -INF,使用时先对 1
       → -> n 调用 hungary() , 再 get_ans() 求值
4
    int w[N][N]:
    int lx[N], ly[N], match[N], way[N], slack[N];
5
6
    bool used[N];
7
    void init() {
8
      for (int i = 1; i <= n; i++) {
9
        match[i] = 0;
10
        lx[i] = 0:
        ly[i] = 0;
11
12
        way[i] = 0;
13
    }
14
15
    void hungary(int x) {
      match[0] = x:
16
      int j0 = 0;
17
18
      for (int j = 0; j \le n; j++) {
        slack[j] = INF;
19
20
        used[j] = false;
21
      do {
22
23
        used[j0] = true;
```

```
int i0 = match[j0], delta = INF, j1 = 0;
25
          for (int j = 1; j \le n; j++) {
26
            if (used[j] == false) {
              int cur = -w[i0][j] - lx[i0] - ly[j];
27
              if (cur < slack[j]) {</pre>
28
29
                slack[j] = cur;
30
                way[j] = j0;
31
              }
32
              if (slack[j] < delta) {</pre>
33
                delta = slack[j];
34
                j1 = j;
35
36
           }
37
         }
38
         for (int j = 0; j \le n; j++) {
30
            if (used[j]) {
              lx[match[j]] += delta;
40
              ly[j] -= delta;
41
            else slack[j] -= delta;
43
         }
44
45
         j0 = j1;
       } while (match[j0] != 0);
46
47
       do {
48
         int j1 = way[j0];
49
         match[j0] = match[j1];
50
          j0 = j1;
51
       } while (j0);
52
53
     int get_ans() {
       int sum = 0;
       for(int i = 1; i <= n; i++) {
55
          if (w[match[i]][i] == -INF); // 无解
57
          if (match[i] > 0) sum += w[match[i]][i];
58
       }
59
       return sum;
60
     }
   } km;
61
```

6.3 点双连通分量

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
 2
   struct BCC { // N = NO + MO. Remember to call
      \hookrightarrow init(&raw_graph).
     Graph *g, forest; // g is raw graph ptr.
 4
     int dfn[N], DFN, low[N];
5
     int stack[N], top;
                            // Where edge i is expanded to in
6
     int expand_to[N];

    ⇔ expaned graph.

     // Vertex i expaned to i.
8
     int compress_to[N]; \ //\ Where vertex i is compressed to.
Q
     bool vertex_type[N], cut[N], compress_cut[N], branch[M];
     //std::vector<int> BCC_component[N]; // Cut vertex
10
        \hookrightarrow belongs to none.
11
     __inline void init(Graph *raw_graph) {
       g = raw_graph;
12
13
14
     void DFS(int u, int pe) {
       dfn[u] = low[u] = ++DFN; cut[u] = false;
15
16
       if (!~g->adj[u]) {
17
         cut[u] = 1;
         compress_to[u] = forest.new_node();
18
19
         compress_cut[compress_to[u]] = 1;
20
       for (int e = g->adj[u]; ~e; e = g->nxt[e]) {
21
22
         int v = g->v[e];
23
         if ((e ^ pe) > 1 && dfn[v] > 0 && dfn[v] < dfn[u]) {
24
            stack[top++] = e;
25
           low[u] = std::min(low[u], dfn[v]);
         }
26
         else if (!dfn[v]) {
27
28
            stack[top++] = e; branch[e] = 1;
```

```
29
           DFS(v, e);
30
           low[u] = std::min(low[v], low[u]);
31
            if (low[v] >= dfn[u]) {
32
             if (!cut[u]) {
                cut[u] = 1:
33
                compress_to[u] = forest.new_node();
34
35
                compress_cut[compress_to[u]] = 1;
             }
36
37
             int cc = forest.new_node();
38
             forest.bi_ins(compress_to[u], cc);
              compress_cut[cc] = 0;
39
              //BCC_component[cc].clear();
40
41
                int cur_e = stack[--top];
42
43
                compress_to[expand_to[cur_e]] = cc;
44
                compress_to[expand_to[cur_e^1]] = cc;
                if (branch[cur e]) {
45
46
                  int v = g->v[cur_e];
47
                  if (cut[v])
48
                    forest.bi_ins(cc, compress_to[v]);
49
50
                    //BCC_component[cc].push_back(v);
                    compress_to[v] = cc;
51
52
53
             } while (stack[top] != e);
55
56
       }
57
     }
58
59
     void solve() {
       forest.init(g->base);
60
61
       int n = g->n;
62
       for (int i = 0; i < g > e; i + +) {
         expand_to[i] = g->new_node();
63
64
65
       memset(branch, 0, sizeof(*branch) * g->e);
       memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
66
67
       for (int i = 0; i < n; i++)
         if (!dfn[i + g->base]) {
68
           top = 0:
69
70
           DFS(i + g - base, -1);
71
72
     }
73
  } bcc;
74
75
  bcc.init(&raw_graph);
76 bcc.solve();
   // Do something with bcc.forest ...
```

6.4 边双连通分量

```
1
   struct BCC {
2
     Graph *g, forest;
     int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N],
3
        // tot[] is the size of each BCC, belong[] is the BCC
        \hookrightarrow that each node belongs to
5
     pair<int, int > ori[M]; // bridge in raw_graph(raw node)
6
     bool is_bridge[M];
7
     __inline void init(Graph *raw_graph) {
8
       g = raw_graph;
9
       memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
       memset(vis + g \rightarrow base, 0, sizeof(*vis) * g \rightarrow n);
10
11
     }
     void tarjan(int u, int from) {
12
       dfn[u] = low[u] = ++dfs\_clock; vis[u] = 1;
13

    stack[++top] = u;

       for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
14
         if ((p ^ 1) == from) continue;
15
16
         int v = g \rightarrow v[p];
         if (vis[v]) {
17
           if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
18
```

```
19
          } else {
20
             tarjan(v, p);
21
             low[u] = min(low[u], low[v]);
             if (low[v] > dfn[u]) is_bridge[p / 2] = true;
22
23
          }
        }
24
        if (dfn[u] != low[u]) return;
25
        tot[forest.new_node()] = 0;
26
27
        do {
28
          belong[stack[top]] = forest.n;
29
          vis[stack[top]] = 2;
30
          tot[forest.n]++;
31
           --top;
32
        } while (stack[top + 1] != u);
      }
33
34
      void solve() {
        forest.init(g -> base);
35
36
        int n = g \rightarrow n;
37
        for (int i = 0; i < n; ++i)
          if (!vis[i + g -> base]) {
38
39
             top = dfs_clock = 0;
40
             tarjan(i + g \rightarrow base, -1);
41
42
        for (int i = 0; i < g -> e / 2; ++i)
43
          if (is_bridge[i]) {
             int e = forest.e;
45
             forest.bi_ins(belong[g -> v[i * 2]], belong[g ->
                \hookrightarrow v[i * 2 + 1]], g \rightarrow w[i * 2]);
             ori[e] = make_pair(g \rightarrow v[i * 2 + 1], g \rightarrow v[i *
46
                \hookrightarrow 2]);
             ori[e + 1] = make_pair(g -> v[i * 2], g -> v[i * 2
                \hookrightarrow + 1]);
48
          }
49
     }
   } bcc;
50
```

6.5 最小树形图

```
const int MAXN,INF;// INF >= sum( W_ij )
   int from [MAXN + 10] [MAXN * 2 + 10], n, m, edge [MAXN +
      \hookrightarrow 10] [MAXN * 2 + 10];
   int sel[MAXN * 2 + 10], fa[MAXN * 2 + 10], vis[MAXN * 2 +
 3

→ 10];

   int getfa(int x){if(x == fa[x]) return x; return fa[x] =
       \hookrightarrow getfa(fa[x]);
   void liuzhu(){ // 1-base: root is 1, answer = (sel[i], i)
      \hookrightarrow for i in [2..n]
     fa[1] = 1;
 6
 7
     for(int i = 2; i <= n; ++i){
        sel[i] = 1; fa[i] = i;
9
        for(int j = 1; j \le n; ++j) if(fa[j] != i)
          if(from[j][i] = i, edge[sel[i]][i] > edge[j][i])
             \hookrightarrow sel[i] = j;
11
     }
12
     int limit = n;
13
        int prelimit = limit; memset(vis, 0, sizeof(vis));
           \hookrightarrow vis[1] = 1;
        for(int i = 2; i <= prelimit; ++i) if(fa[i] == i &&</pre>
15
          16
          int j = i; while(!vis[j]) vis[j] = i, j =
             \hookrightarrow getfa(sel[j]);
          if(j == 1 || vis[j] != i) continue; vector<int> C;
17
             \hookrightarrow int k = j;
18
          do C.push_back(k), k = getfa(sel[k]); while(k != j);
19
          ++limit:
20
          for(int i = 1; i \le n; ++i){
            edge[i][limit] = INF, from[i][limit] = limit;
21
22
23
          fa[limit] = vis[limit] = limit;
          for(int i = 0; i < int(C.size()); ++i){</pre>
24
            int x = C[i], fa[x] = limit;
25
```

```
for(int j = 1; j \le n; ++j)
26
27
              if(edge[j][x] != INF && edge[j][limit] >
                 \hookrightarrow edge[j][x] - edge[sel[x]][x]){
                edge[j][limit] = edge[j][x] - edge[sel[x]][x];
28
29
                from[j][limit] = x;
30
         }
31
          for(int j=1;j<=n;++j) if(getfa(j)==limit)</pre>
32
             33
          sel[limit] = 1;
          for(int j = 1; j \le n; ++j)
34
            if(edge[sel[limit]][limit] > edge[j][limit])
35
               \hookrightarrow sel[limit] = j;
36
37
        if(prelimit == limit) break;
     7
38
     for(int i = limit; i > 1; --i) sel[from[sel[i]][i]] =
39
        \hookrightarrow sel[i];
40 }
```

6.6 带花树

```
1 vector<int> link[maxn];
  int n,match[maxn],Queue[maxn],head,tail;
3 int pred[maxn],base[maxn],start,finish,newbase;
4 bool InQueue [maxn], InBlossom [maxn];
5 | void push(int u){ Queue[tail++]=u;InQueue[u]=true; }
  int pop(){ return Queue[head++]; }
7
   int FindCommonAncestor(int u,int v){
8
     bool InPath[maxn];
     for(int i=0:i<n:i++) InPath[i]=0:</pre>
9
10
     while(true){ u=base[u];InPath[u]=true;if(u==start)
        while(true){ v=base[v];if(InPath[v])
11
       12
     return v;
13 }
14
   void ResetTrace(int u){
15
16
     while(base[u]!=newbase){
17
       v=match[u];
18
       InBlossom[base[u]] = InBlossom[base[v]] = true;
19
       u=pred[v]:
20
       if(base[u]!=newbase) pred[u]=v;
21
22
  ۱,
23
   void BlossomContract(int u,int v){
    newbase=FindCommonAncestor(u.v):
24
     for (int i=0;i<n;i++)</pre>
25
     InBlossom[i]=0;
26
27
     ResetTrace(u);ResetTrace(v);
     if(base[u]!=newbase) pred[u]=v;
28
29
     if(base[v]!=newbase) pred[v]=u;
30
     for(int i=0:i<n:++i)
31
     if(InBlossom[base[i]]){
32
       base[i]=newbase;
33
       if(!InQueue[i]) push(i);
34
     }
  }
35
   bool FindAugmentingPath(int u){
36
37
     bool found=false:
38
     for(int i=0;i<n;++i) pred[i]=-1,base[i]=i;</pre>
     for (int i=0;i<n;i++) InQueue[i]=0;</pre>
39
     start=u;finish=-1; head=tail=0; push(start);
40
41
     while(head<tail){
       int u=pop();
42
       for(int i=link[u].size()-1:i>=0:i--){
43
44
         int v=link[u][i];
         if (base[u]!=base[v]&&match[u]!=v)
45
           if(v==start||(match[v]>=0&&pred[match[v]]>=0))
46
47
             BlossomContract(u,v);
           else if(pred[v]==-1){
48
49
             pred[v]=u;
```

```
if(match[v]>=0) push(match[v]);
50
51
             else{ finish=v; return true; }
52
       }
53
    }
54
55
     return found:
56
   }
57
   void AugmentPath(){
58
     int u=finish, v, w;
59
     while (u \ge 0) {

    v=pred[u]; w=match[v]; match[v]=u; match[u]=v; u=w; }

60
61
   void FindMaxMatching(){
62
     for(int i=0;i<n;++i) match[i]=-1;</pre>
63
     for(int i=0;i<n;++i) if(match[i]==-1)</pre>
        64
```

6.7 支配树

```
vector<int> prec[N], succ[N];
   vector<int> ord;
 2
   int stamp, vis[N];
 3
   int num[N];
   int fa[N];
   void dfs(int u) {
 7
     vis[u] = stamp;
8
     num[u] = ord.size();
9
     ord.push_back(u);
10
     for (int i = 0; i < (int)succ[u].size(); ++i) {</pre>
11
       int v = succ[u][i];
12
        if (vis[v] != stamp) {
13
          fa[v] = u;
14
          dfs(v);
15
16
     }
17
   }
   int fs[N], mins[N], dom[N], sem[N];
   int find(int u) {
20
     if (u != fs[u]) {
       int v = fs[u];
21
       fs[u] = find(fs[u]);
22
        if (mins[v] != -1 && num[sem[mins[v]]] <</pre>
23
           \hookrightarrow num[sem[mins[u]]]) {
24
          mins[u] = mins[v];
25
       }
     }
26
27
     return fs[u];
28
   void merge(int u, int v) { fs[u] = v; }
30
   vector<int> buf[N];
31
   int buf2[N];
32
   void mark(int source) {
33
     ord.clear():
34
     ++stamp;
35
     dfs(source):
     for (int i = 0; i < (int)ord.size(); ++i) {</pre>
36
37
       int u = ord[i];
38
       fs[u] = u, mins[u] = -1, buf2[u] = -1;
39
40
     for (int i = (int) ord.size() - 1; i > 0; --i) {
41
       int u = ord[i], p = fa[u];
        sem[u] = p;
42
43
        for (int j = 0; j < (int)prec[u].size(); ++j) {</pre>
         int v = prec[u][j];
44
45
          if (use[v] != stamp) continue;
46
          if (num[v] > num[u]) {
            find(v); v = sem[mins[v]];
47
48
          if (num[v] < num[sem[u]]) {</pre>
49
            sem[u] = v;
50
51
```

```
52
53
       buf[sem[u]].push_back(u);
54
       mins[u] = u:
55
       merge(u, p);
       while (buf[p].size()) {
56
         int v = buf[p].back();
57
58
         buf[p].pop_back();
59
         find(v);
         if (sem[v] == sem[mins[v]]) {
60
61
            dom[v] = sem[v];
62
          } else {
63
            buf2[v] = mins[v];
64
65
       }
66
     }
     dom[ord[0]] = ord[0];
67
     for (int i = 0; i < (int)ord.size(); ++i) {</pre>
68
69
       int u = ord[i];
70
       if (~buf2[u]) {
          dom[u] = dom[buf2[u]];
71
72
       }
73
     }
74 }
```

6

无向图最小割 6.8

```
int cost[maxn] [maxn], seq[maxn], len[maxn], n, m, pop, ans;
   bool used[maxn]:
3
  void Init(){
4
    int i,j,a,b,c;
5
     for(i=0;i<n;i++) for(j=0;j<n;j++) cost[i][j]=0;</pre>
6
     for(i=0;i<m;i++){</pre>
7
       scanf("%d %d %d",&a,&b,&c); cost[a][b]+=c;
           \hookrightarrow cost[b][a]+=c;
8
     pop=n; for(i=0;i<n;i++) seq[i]=i;</pre>
9
  ۱,
10
   void Work(){
11
     ans=inf; int i,j,k,l,mm,sum,pk;
12
13
     while(pop > 1){
       for(i=1;i<pop;i++) used[seq[i]]=0; used[seq[0]]=1;</pre>
15
       for(i=1;i<pop;i++) len[seq[i]]=cost[seq[0]][seq[i]];</pre>
16
       pk=0; mm=-inf; k=-1;
       for(i=1;i < pop;i++) \ if(len[seq[i]] > mm){} \{
17
          18
       for(i=1;i<pop;i++){</pre>
19
         used[seq[l=k]]=1;
20
          if(i==pop-2) pk=k;
21
         if(i==pop-1) break;
22
         mm=-inf:
         for(j=1;j<pop;j++) if(!used[seq[j]])</pre>
23
24
            if((len[seq[j]]+=cost[seq[1]][seq[j]]) > mm)
25
              mm=len[seq[j]], k=j;
26
       }
27
       sum=0;
       for(i=0;i<pop;i++) if(i != k)</pre>
28

    sum+=cost[seq[k]][seq[i]];

29
       ans=min(ans,sum);
30
        for(i=0:i<pop:i++)
31
          cost[seq[k]][seq[i]]=cost[seq[i]][seq[k]]+=cost[seq[pk]][seq[i]];
32
        seq[pk]=seq[--pop];
33
     printf("%d\n",ans);
34
35 }
```

最大团搜索 6.9

```
const int N = 1000 + 7;
  vector<vector<bool> > adj;
3
  class MaxClique {
4
      const vector<vector<bool> > adj;
5
      const int n:
```

```
vector<int> result, cur_res;
 7
       vector<vector<int> > color_set;
8
       const double t_limit; // MAGIC
9
     int para, level;
10
     vector<pair<int, int> > steps;
   public:
11
       class Vertex {
12
13
       public:
14
            int i, d;
15
            Vertex(int i, int d = 0) : i(i), d(d) \{ \}
16
       ን:
17
       void reorder(vector<Vertex> &p) {
18
            for (auto &u : p) {
19
                u.d = 0;
20
                for (auto v : p) u.d += adj[v.i][u.i];
21
            sort(p.begin(), p.end(), [&](const Vertex &a,
22
              }
23
     // reuse p[i].d to denote the maximum possible clique
24
        \hookrightarrow \texttt{for first i vertices.}
25
       void init_color(vector<Vertex> &p) {
26
            int maxd = p[0].d:
27
            for (int i = 0; i < p.size(); i++) p[i].d = min(i,</pre>
              \rightarrow maxd) + 1;
28
29
       bool bridge(const vector<int> &s, int x) {
30
            for (auto v : s) if (adj[v][x]) return true;
31
            return false;
32
       }
33
     // approximate estimate the p[i].d
     // Do not care about first mink color class (For better
34
        \hookrightarrow \texttt{result}, we must get some vertex in some color class
        \hookrightarrow larger than mink )
       void color_sort(vector<Vertex> &cur) {
35
            int totc = 0, ptr = 0, mink =
36
              37
            for (int i = 0; i < cur.size(); i++) {</pre>
                int x = cur[i].i, k = 0;
38
39
                while (k < totc && bridge(color_set[k], x))</pre>
                   if (k == totc) color_set[totc++].clear();
41
                color_set[k].push_back(x);
42
                if (k < mink) cur[ptr++].i = x;</pre>
           }
43
44
           if (ptr) cur[ptr - 1].d = 0;
45
            for (int i = mink; i < totc; i ++) {</pre>
46
                for (auto v : color_set[i]) {
47
                    cur[ptr++] = Vertex(v, i + 1);
48
           }
49
       }
50
51
       void expand(vector<Vertex> &cur) {
       steps[level].second = steps[level].second -

    steps[level].first + steps[level - 1].first;

53
       steps[level].first = steps[level - 1].second;
54
            while (cur.size()) {
                if (cur_res.size() + cur.back().d <=</pre>

    result.size()) return :
                int x = cur.back().i;
                cur_res.push_back(x); cur.pop_back();
58
                vector<Vertex> remain;
59
                for (auto v : cur) {
60
                    if (adj[v.i][x]) remain.push_back(v.i);
61
62
                if (remain.size() == 0) {
                    if (cur_res.size() > result.size()) result
                       \hookrightarrow = cur_res;
                } else {
64
            // Magic ballance.
65
66
            if (1. * steps[level].second / ++para < t_limit)</pre>
              \hookrightarrow reorder(remain);
```

```
67
                     color_sort(remain);
68
            steps[level++].second++;
69
                     expand(remain);
70
            level--:
                }
71
72
                cur_res.pop_back();
            }
73
74
       }
75
   public:
76
       MaxClique(const vector<vector<bool> > &_adj, int n,
          \hookrightarrow double tt = 0.025) : adj(_adj), n(n), t_limit(tt)
77
            result.clear();
78
            cur_res.clear();
79
            color_set.resize(n);
80
       steps.resize(n + 1);
       fill(steps.begin(), steps.end(), make_pair(0, 0));
81
82
       level = 1;
83
       para = 0;
84
       }
85
       vector<int> solve() {
86
            vector<Vertex> p;
            for (int i = 0; i < n; i++)
87

    p.push_back(Vertex(i));
88
            reorder(p);
89
            init_color(p);
90
            expand(p);
91
            return result;
92
93 };
```

6.10 斯坦纳树

```
void SPFA(int *dist) {
2
        static int line[maxn + 5];
 3
        static bool hash[maxn + 5];
 4
        int f = 0, r = 0;
 5
        for(int i = 1; i <= N; i++)</pre>
            if(dist[i] < inf) {</pre>
 7
                 line[r] = i;
 8
                 hash[i] = true;
9
                 r = (r + 1) \% (N + 1);
            }
10
        while(f != r) {
11
12
            int t = line[f];
13
            hash[t] = false;
14
            f = (f + 1) \% (N + 1);
            for(int i = head[t]; i ; i = edge[i].next) {
15
                 int v = edge[i].v, dt = dist[t] + edge[i].w;
16
17
                 if(dt < dist[v]) {</pre>
18
                     dist[v] = dt;
19
                     if(!hash[v]) {
20
                          if(dist[v] < dist[line[f]]) {</pre>
                              f = (f + N) \% (N + 1);
21
                              line[f] = v;
22
                          }
23
24
25
                              line[r] = v;
26
                               r = (r + 1) \% (N + 1);
28
                          hash[v] = true;
                     }
29
                 }
30
31
            }
32
        }
   | }
33
   void solve() {
34
35
        for(int i = 1; i <= S; i++) {</pre>
            for(int j = 1; j <= N; j++)</pre>
36
37
                 for(int k = (i - 1) \& i; k; k = (k - 1) \& i)
38
                     G[i][j] = std::min(G[i][j], G[k][j] + G[k]
                        \hookrightarrow ^ i][j]);
            SPFA(G[i]);
39
```

```
40 }
41 }
```

6.11 虚树

```
bool cmp(const int lhs,const int rhs) {
 2
     return dfn[lhs] < dfn[rhs];</pre>
3
   }
4
   void build() {
     std::sort(h + 1, h + 1 + m, cmp);
 6
     int top = 0;
 7
     for (int i = 1; i <= m; i++) {
8
       if (!top) father[st[++top] = h[i]] = 0;
9
       else {
            int p = h[i], lca = LCA(h[i],st[top]);
10
11
            while(d[st[top]] > d[lca]) {
12
                if (d[st[top - 1]] <= d[lca])</pre>
13
                    father[st[top]] = lca;
14
                top--;
15
            }
16
            if (st[top] != lca) {
17
                t[++tot] = 1ca;
18
                father[lca] = st[top];
19
                st[++top] = lca;
20
21
            father[p] = lca;
            st[++top] = p;
22
23
24
     }
25
   }
```

6.12 点分治

```
template<class TAT>void checkmax(TAT &x,TAT y) {
2
     if(x < y) x = y;
3
   }
   template<class TAT>void checkmin(TAT &x,TAT y) {
5
     if(y < x) x = y;
6
   }
7
   void getsize(int u,int fa) {
8
     size[u] = 1:
9
     smax[u] = 0;
     for(int i = 0; i < G[u].size(); i++) {</pre>
11
       int v = G[u][i];
12
       if(v == fa || ban[v]) continue;
13
       getsize(v, u);
       size[u] += size[v];
14
15
       checkmax(smax[u], size[v]);
16
17
   }
18
   int getroot(int u,int ts,int fa) {
19
     checkmax(smax[u], ts - size[u]);
20
     int res = u:
21
     for(int i = 0; i < G[u].size(); i++) {</pre>
22
       int v = G[u][i];
23
       if(v == fa || ban[v]) continue;
24
       int w = getroot(v, ts, u);
25
       if(smax[w] < smax[res]) res = w;</pre>
     }
26
27
     return res;
28
   }
29
   void solve() {
30
     static int line[maxn];
     static std::vector<int> vec;
31
     int f = 0, r = 0;
32
33
     line[r++] = 1;
     while(f != r) {
35
       int u = line[f++];
36
       getsize(u, 0);
37
       u = getroot(u, size[u], 0);
       ban[u] = true;
38
```

```
39     vec.clear();
40     for(int i = 0; i < G[u].size(); i++)
41         if(!ban[G[u][i]])     vec.push_back(G[u][i]);
42     for(int i = 0; i < vec.size(); i++)
43         line[r++] = vec[i];
44     }
45 }</pre>
```

6.13 最小割最大流

```
1
   bool BFS() {
2
       for(int i = 1; i <= ind; i++) dep[i] = 0;</pre>
3
       dep[S] = 1, line.push(S);
4
       while(!line.empty()) {
5
           int now = line.front();
           line.pop();
7
            for(int i = head[now], p; i ; i = edge[i].next)
8
                if(edge[i].cap && !dep[p = edge[i].v])
9
                    dep[p] = dep[now] + 1, line.push(p);
10
       }
11
       if(dep[T]) {
            for(int i = 1; i <= ind; i++)</pre>
12
13
                cur[i] = head[i];
14
            return true;
       } else return false;
15
16 }
17
   int DFS(int a,int flow) {
       if(a == T) return flow;
18
19
       int ret = 0;
20
       for(int &i = cur[a], p; i ; i = edge[i].next)
           if(dep[p = edge[i].v] == dep[a] + 1 &&
21
              \hookrightarrow edge[i].cap) {
                int ff = DFS(p, std::min(flow, edge[i].cap));
22
23
                flow -= ff, edge[i].cap -= ff;
24
                ret += ff, edge[i ^ 1].cap += ff;
25
                if(!flow) break;
            }
26
27
            return ret;
28 }
29
   int solve() {
       int totflow = 0;
30
31
       while(BFS())
32
            totflow += DFS(S, INF);
33
34
       return totflow:
35
```

6.14 最小费用流

```
bool SPFA() {
1
2
       static int line[maxv];
3
       static bool hash[maxv];
4
       register int f = 0, r = 0;
5
     for(int i = 1; i <= ind; i++) {</pre>
         dist[i] = inf;
6
7
         from[i] = 0;
8
9
       dist[S] = 0, line[r] = S, r = (r + 1) % maxv;
       hash[S] = true;
10
       while(f != r) {
11
           int x = line[f];
12
13
           line[f] = 0, f = (f + 1) \% maxv;
14
           hash[x] = false;
           for(int i = head[x]; i; i = edge[i].next)
15
                if(edge[i].cap) {
17
                    int v = edge[i].v;
                    int w = dist[x] + edge[i].cost;
18
19
                    if(w < dist[v]) {</pre>
                        dist[v] = w;
20
                        from[v] = i:
21
                        if(!hash[v]) {
22
```

```
if(f != r && dist[v] <=</pre>
                                   \hookrightarrow dist[line[f]])
24
                                    f = (f - 1 + maxv) \% maxv,
                                        \hookrightarrow line[f] = v;
25
                                else line[r] = v, r = (r + 1) %
                                   \hookrightarrow \mathtt{maxv}:
                                hash[v] = true;
26
27
                           }
28
                      }
29
                 }
30
31
        return from[T];
32
33
    int back(int x,int flow) {
34
     if(from[x]) {
        flow = back(edge[from[x] ^ 1].v, std::min(flow,
35

    edge[from[x]].cap)):
36
        edge[from[x]].cap -= flow;
        edge[from[x] ^ 1].cap += flow;
37
38
39
     return flow;
40
   }
41
   int solve() {
42
        int mincost = 0, maxflow = 0;
43
        while(SPFA()) {
44
             int flow = back(T, inf);
45
             mincost += dist[T] * flow;
             maxflow += flow;
46
47
        }
48
        return mincost;
```

6.15 zkw 费用流

```
int S, T, totFlow, totCost;
   int dis[N], slack[N], visit[N];
 3
   int modlable () {
       int delta = INF;
        for (int i = 1; i <= T; i++) {
 6
            if (!visit[i] && slack[i] < delta) delta =</pre>
               \hookrightarrow \mathtt{slack[i]};
 7
            slack[i] = INF;
       }
8
9
        if (delta == INF) return 1;
        for (int i = 1; i <= T; i++)
10
11
            if (visit[i]) dis[i] += delta;
12
        return 0;
   }
13
   int dfs (int x, int flow) {
14
15
       if (x == T) {
16
            totFlow += flow:
17
            totCost += flow * (dis[S] - dis[T]);
18
            return flow;
19
       }
20
       visit[x] = 1:
21
        int left = flow;
        for (int i = e.last[x]; ~i; i = e.succ[i])
22
23
            if (e.cap[i] > 0 && !visit[e.other[i]]) {
24
                int y = e.other[i];
                if (dis[y] + e.cost[i] == dis[x]) {
25
26
                     int delta = dfs (y, min (left, e.cap[i]));
27
                     e.cap[i] -= delta;
                     e.cap[i ^ 1] += delta;
28
29
                     left -= delta;
30
                     if (!left) { visit[x] = 0; return flow; }
31
                } else {
32
                     slack[y] = min (slack[y], dis[y] +
                        \hookrightarrow e.cost[i] - dis[x]);
33
34
            }
35
        return flow - left;
36
37 pair <int, int> minCost () {
```

```
totFlow = 0; totCost = 0;
38
39
       fill (dis + 1, dis + T + 1, 0);
40
       do {
41
               fill (visit + 1, visit + T + 1, 0);
42
           } while (dfs (S, INF));
43
       } while (!modlable ());
45
       return make_pair (totFlow, totCost);
46
  }
```

6.16 最小割树

```
\hookrightarrow cnt,n,m,dis[N],last[N],a[N],tmp[N],ans[N][N],s,t,mark[N];
   struct edge{int to,c,next;}e[N*200];
   queue <int> q;
   void addedge(int u,int v,int c) {
       e[++cnt].to=v;e[cnt].c=c;
5
     e[cnt].next=last[u];last[u]=cnt;
6
7
       e[++cnt].to=u;e[cnt].c=c;
8
     e[cnt].next=last[v];last[v]=cnt;
   }
9
10
   bool bfs() {
       memset(dis,0,sizeof(dis));
11
12
       dis[s]=2;
13
       while (!q.empty()) q.pop();
14
       q.push(s);
15
       while (!q.empty()) {
16
           int u=q.front();
17
           q.pop();
           for (int i=last[u]:i:i=e[i].next)
18
19
                if (e[i].c&&!dis[e[i].to]) {
20
                    dis[e[i].to]=dis[u]+1;
21
                    if (e[i].to==t) return 1;
22
                    q.push(e[i].to);
23
24
       }
25
       return 0;
26 }
27
   int dfs(int x,int maxf) {
28
       if (x==t||!maxf) return maxf;
       int ret=0:
29
       for (int i=last[x];i;i=e[i].next)
30
            if (e[i].c&&dis[e[i].to]==dis[x]+1) {
31
32
                int f=dfs(e[i].to,min(e[i].c,maxf-ret));
33
                e[i].c-=f;
34
                e[i^1].c+=f;
                ret+=f:
35
                if (ret==maxf) break;
36
37
       if (!ret) dis[x]=0;
38
39
       return ret;
40
  |}
41
   void dfs(int x) {
42
       mark[x]=1:
43
       for (int i=last[x];i;i=e[i].next)
            if (e[i].c&&!mark[e[i].to]) dfs(e[i].to);
44
45
  }
46
   void solve(int l,int r) {
       if (l==r) return;
47
48
       s=a[1]:t=a[r]:
49
       for (int i=2;i<=cnt;i+=2)</pre>
            e[i].c=e[i^1].c=(e[i].c+e[i^1].c)/2;
51
       int flow=0;
       while (bfs()) flow+=dfs(s,inf);
52
       memset(mark,0,sizeof(mark));
53
       dfs(s):
54
55
       for (int i=1;i<=n;i++)</pre>
            if (mark[i])
56
57
                for (int j=1;j<=n;j++)</pre>
58
                    if (!mark[j])
                ans[i][j]=ans[j][i]=min(ans[i][j],flow);
59
60
       int i=1,j=r;
```

```
for (int k=1;k<=r;k++)
    if (mark[a[k]]) tmp[i++]=a[k];
    else tmp[j--]=a[k];

for (int k=1;k<=r;k++)
        a[k]=tmp[k];

solve(1,i-1);
    solve(j+1,r);

8 }</pre>
```

6.17 上下界网络流建图

B(u,v) 表示边 (u,v) 流量的下界,C(u,v) 表示边 (u,v) 流量的上界,F(u,v) 表示边 (u,v) 的流量。设 G(u,v)=F(u,v)-B(u,v),显然有

$$0 \le G(u, v) \le C(u, v) - B(u, v)$$

6.17.1 无源汇的上下界可行流

建立超级源点 S^* 和超级汇点 T^* ,对于原图每条边 (u,v) 在新网络中连如下三条边: $S^* \to v$,容量为 B(u,v); $u \to T^*$,容量为 B(u,v); $u \to v$,容量为 C(u,v) - B(u,v)。 最后求新网络的最大流,判断从超级源点 S^* 出发的边是否都满流即可,边 (u,v) 的最终解中的实际流量为 G(u,v) + B(u,v)。

6.17.2 有源汇的上下界可行流

从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的边。 按照**无源汇的上下界可行流**一样做即可,流量即为 $T \to S$ 边上的流量。

6.17.3 有源汇的上下界最大流

- 1. 在**有源汇的上下界可行流**中,从汇点 T 到源点 S 的 边改为连一条上界为 ∞ ,下届为 x 的边。x 满足二分性质,找到最大的 x 使得新网络存在**无源汇的上下界可行流**即为原图的最大流。
- 2. 从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的 边,变成无源汇的网络。按照**无源汇的上下界可行流** 的方法,建立超级源点 S^* 和超级汇点 T^* ,求一遍 $S^* \to T^*$ 的最大流,再将从汇点 T 到源点 S 的这条边拆掉,求一次 $S \to T$ 的最大流即可。

6.17.4 有源汇的上下界最小流

- 1. 在**有源汇的上下界可行流**中,从汇点 T 到源点 S 的 边改为连一条上界为 x,下界为 0 的边。x 满足二分性质,找到最小的 x 使得新网络存在**无源汇的上下界可行流**即为原图的最小流。
- 2. 按照无源汇的上下界可行流的方法,建立超级源点 S^* 与超级汇点 T^* ,求一遍 $S^* \to T^*$ 的最大流,但 是注意这一次不加上汇点 T 到源点 S 的这条边,即 不使之改为无源汇的网络去求解。求完后,再加上那 条汇点 T 到源点 S 上界 ∞ 的边。因为这条边下界为 0,所以 S^* , T^* 无影响,再直接求一次 $S^* \to T^*$ 的最大流。若超级源点 S^* 出发的边全部满流,则 $T \to S$ 边上的流量即为原图的最小流,否则无解。

7. 其他

7.1 Dancing Links

7.1.1 精确覆盖

```
#pragma comment(linker, "/STACK:1024000000,1024000000")
#define maxn 1000005

using namespace std;
int head,sz;
int U[maxn],D[maxn],L[maxn],R[maxn];
int H[maxn],ROW[maxn],C[maxn],O[maxn];

void remove(int c) {
```

```
8
       L[R[c]]=L[c];
9
       R[L[c]]=R[c];
10
       for(int i=D[c]; i!=c; i=D[i])
            for(int j=R[i]; j!=i; j=R[j]) {
11
                U[D[j]]=U[j];
12
                D[U[i]]=D[i];
13
                --S[C[j]];
15
            }
16 }
   void resume(int c) {
17
       for(int i=U[c]; i!=c; i=U[i]) {
18
19
            for(int j=L[i]; j!=i; j=L[j]) {
20
                ++S[C[j]];
21
                U[D[j]]=j;
22
                D[U[j]]=j;
23
24
25
       L[R[c]]=c;
       R[L[c]]=c;
26
27
  }
   void init(int m) {
28
       head=0;//头指针为 0
29
       for(int i=0; i<=m; i++) {</pre>
30
           U[i]=i:
31
           D[i]=i;//建立双向十字链表
32
           L[i]=i-1;
33
            R[i]=i+1;
34
35
            S[i]=0;
36
37
       R[m] = 0;
       L[0]=m;
38
       S[0] = INF + 1:
39
40
       sz=m+1;
41
       memset(H,0,sizeof(H));
42
  | }
43
   void insert(int i, int j) {
44
       if(H[i]) {
           L[sz] = L[H[i]];
45
            R[sz] = H[i];
46
            L[R[sz]] = sz;
47
            R[L[sz]] = sz;
48
49
       }
       else {
50
            L[sz] = sz;
51
            R[sz] = sz;
52
            H[i] = sz;
53
55
       U[sz] = U[j];
       D[sz] = j;
56
       U[D[sz]] = sz;
57
       D[U[sz]] = sz;
58
59
       C[sz] = j;
       ROW[sz] = i;
60
61
        ++S[j];
62
       ++sz;
63 }
   bool dfs(int k,int len) {
64
65
       if(R[head] == head) return true;
66
        int s=INF,c;
67
       for (int t=R[head]; t!=head; t=R[t])
68
            if (S[t] < s) s = S[t], c = t;</pre>
69
       remove(c);
       for(int i=D[c]; i!=c; i=D[i]) {
70
71
            0[k]=ROW[i];
72
            for(int j=R[i]; j!=i; j=R[j])
73
                remove(C[j]);
74
            if(dfs(k+1,len))
75
                return true:
            for(int j=L[i]; j!=i; j=L[j])
76
77
                resume(C[j]);
       }
78
79
       resume(c);
80
       return false;
```

```
81 | }
```

7.1.2 重复覆盖

```
int h()
 2
   {
 3
        int i,j,k,count=0;
        bool visit[N]:
 4
 5
        memset(visit,0,sizeof(visit));
 6
        for(i=R[0];i;i=R[i])
 7
 8
            if(visit[i]) continue;
 9
            count++;
10
            visit[i]=1;
            for(j=D[i];j!=i;j=D[j])
11
12
            {
13
                 for(k=R[j];k!=j;k=R[k])
14
                     visit[C[k]]=1;
15
            }
        }
16
17
        return count;
18
   }
   void Dance(int k)
19
20
21
        int i,j,c,Min,ans;
22
        ans=h();
23
        if(k+ans>K || k+ans>=ak) return;
24
        if(!R[0])
25
        {
26
            if(k<ak) ak=k;</pre>
27
            return;
28
        for(Min=N,i=R[0];i;i=R[i])
29
30
            if(S[i]<Min) Min=S[i],c=i;</pre>
31
        for(i=D[c];i!=c;i=D[i])
32
33
            remove(i);
34
            for(j=R[i];j!=i;j=R[j])
35
                remove(j);
36
            Dance(k+1):
37
            for(j=L[i];j!=i;j=L[j])
38
                resume(j);
39
            resume(i);
        7
40
41
        return:
42
```

7.2 蔡勒公式

```
int zeller(int y,int m,int d) {
  if (m<=2) y--,m+=12; int c=y/100; y%=100;
  int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
  if (w<0) w+=7; return(w);
}</pre>
```

7.3 五边形数定理

$$p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k-1} p(n - \frac{k(3k-1)}{2})$$

```
1 LL dp[N],fi[N];
2
  LL five(LL x) { return (3*x*x-x)/2; }
   void wbxs(){
3
4
       dp[0]=1:
       int t=1000; //其实可以等于 sqrt(N)
5
6
       for(int i=-t;i<=t;++i)</pre>
7
           fi[i+t]=five(i);
       for(int i=1;i<=100000;++i){
8
9
           int flag=1;
10
           for(int j=1;;++j){
               LL a=fi[j+t],b=fi[-j+t];
11
12
               if(a>i && b>i) break;
```

7.4 凸包闵可夫斯基和

```
1 // cv[0..1] 为两个顺时针凸包, 其中起点等于终点, 求
     →出的闵可夫斯基和不一定是严格凸包
  int i[2] = \{0, 0\}, len[2] = \{(int)cv[0].size() - 1,
     \hookrightarrow (int)cv[1].size() - 1};
  vector<P> mnk;
  mnk.push_back(cv[0][0] + cv[1][0]);
5
 J op
    int d((cv[0][i[0] + 1] - cv[0][i[0]]) * (cv[1][i[1] + 1]
6
       \hookrightarrow - cv[1][i[1]]) >= 0);
    mnk.push_back(cv[d][i[d] + 1] - cv[d][i[d]] +
      \hookrightarrow mnk.back());
8
    i[d] = (i[d] + 1) % len[d];
  } while(i[0] || i[1]);
```

8. 技巧

8.1 STL 归还空间

```
template <typename T>
2 __inline void clear(T& container) {
    container.clear(); // 或者删除了一堆元素
    T(container).swap(container);
}
```

8.2 大整数取模

8.3 读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
  |// 用 isdigit() 会小幅变慢
  // 返回 false 表示读到文件尾
   namespace Reader {
5
       const int L = (1 << 15) + 5;
6
       char buffer[L], *S, *T;
7
       __inline bool getchar(char &ch) {
8
           if (S == T) {
               T = (S = buffer) + fread(buffer, 1, L, stdin);
9
               if (S == T) {
11
           ch = EOF;
12
           return false;
         }
13
           }
14
       ch = *S++;
15
16
       return true;
17
       __inline bool getint(int &x) {
18
19
       char ch; bool neg = 0;
20
       for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^=
          \hookrightarrow ch == '-';
21
       if (ch == EOF) return false;
       x = ch - '0';
22
       for (; getchar(ch), ch >= '0' && ch <= '9'; )
23
         x = x * 10 + ch - '0';
24
```

8.4 二次随机法

```
#include <random>

int main() {
    std::mt19937 g(seed); // std::mt19937_64
    std::cout << g() << std::endl;
}</pre>
```

8.5 vimrc

```
set ruler
   set number
   set smartindent
   set autoindent
   set tabstop=4
   set softtabstop=4
   set shiftwidth=4
   set hlsearch
   set autoread
   set backspace=2
   set mouse=a
13
   syntax on
16
   nmap <C-A> ggVG
   vmap <C-C> "+y
17
18
19
   filetype plugin indent on
   autocmd FileType cpp set cindent
   autocmd FileType cpp map <F9> :!g++ % -o %< -g -std=c++11
      \hookrightarrow -Wall -Wextra -Wconversion && size %< <CR>
   autocmd FileType cpp map <C-F9> :!g++ \% -o \%< -std=c++11
      \hookrightarrow -02 && size %< <CR>
   autocmd FileType cpp map <F8> :!time ./%< < %<.in <CR>
   autocmd FileType cpp map <F5> :!time ./%< <CR>
26
   map <F3> :vnew %<.in <CR>
   map <F4> :!gedit % <CR>
```

8.6 控制 cout 输出实数精度

```
std::cout << std::fixed << std::setprecision(5);
```

8.7 汇编技巧

```
1 03优化
   #define __ _attribute__ ((optimize("-03")))
   #define _ __ _inline __attribute__ ((__gnu_inline__,
      \hookrightarrow __always_inline__, __artificial__))
   汇编开栈
5
   #pragma comment(linker, "/STACK:256000000")
 6
8
   int __size = 256 << 20;</pre>
   char* __p_ = (char *) malloc(__size__) + __size__;
11
   int main() {
12
     __asm__("movl %0, %%esp\n" :: "r"(__p__));
13
     return 0;
14
```

9. 提示

9.1 线性规划转对偶

 $\begin{array}{l} \text{maximize } \mathbf{c}^T \mathbf{x} \\ \text{subject to } \mathbf{A} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0 \\ \end{array} \Longrightarrow \begin{array}{l} \text{minimize } \mathbf{y}^T \mathbf{b} \\ \text{subject to } \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T, \mathbf{y} \geq 0 \end{array}$

9.2 NTT 素数及其原根

Prime	Primitive root
1053818881	7
1051721729	6
1045430273	3
1012924417	5
1007681537	3

9.3 积分表

1.
$$\int \frac{\mathrm{d}x}{ax^2 + bx + c} = \begin{cases} \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}} + C & (b^2 < 4ac) \\ \frac{1}{\sqrt{b^2 - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^2 - 4ac}}{2ax + b + \sqrt{b^2 - 4ac}} \right| + C & (b^2 > 4ac) \end{cases}$$

2.
$$\int \frac{x}{ax^2+bx+c} dx = \frac{1}{2a} \ln|ax^2+bx+c| - \frac{b}{2a} \int \frac{dx}{ax^2+bx+c}$$

1.
$$\int \frac{dx}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

2.
$$\int \sqrt{ax^2 + bx + c} dx = \frac{\frac{2ax + b}{4a} \sqrt{ax^2 + bx + c}}{\frac{4ac - b^2}{8\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C} + C$$

3.
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a} \sqrt{ax^2 + bx + c}| + C$$

4.
$$\int \frac{\mathrm{d}x}{\sqrt{c+bx-ax^2}} = -\frac{1}{\sqrt{a}}\arcsin\frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

5.
$$\int \sqrt{c + bx - ax^2} dx = \frac{2ax - b}{4a} \sqrt{c + bx - ax^2} + \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

1.
$$\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$$

2.

$$\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C, (a < b) \quad (1)$$

- 1. $\int \tan x dx = -\ln|\cos x| + C$
- 2. $\int \cot x dx = \ln|\sin x| + C$
- 3. $\int \sec x dx = \ln \left| \tan \left(\frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$
- 4. $\int \csc x dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x \cot x \right| + C$
- 5. $\int \sec^2 x dx = \tan x + C$
- 6. $\int \csc^2 x dx = -\cot x + C$
- 7. $\int \sec x \tan x dx = \sec x + C$
- 8. $\int \csc x \cot x dx = -\csc x + C$
- 9. $\int \sin^2 x dx = \frac{x}{2} \frac{1}{4} \sin 2x + C$
- 10. $\int \cos^2 x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$
- 11. $\int \sin^n x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$
- 12. $\int \cos^n x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$
- 13. $\int \frac{dx}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\sin^{n-2} x}$

14.
$$\int \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$$

15

$$\int \cos^m x \sin^n x dx$$

$$= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^n x dx$$

$$= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^m x \sin^{n-2} x dx$$

16.
$$\int \frac{\mathrm{d}x}{a+b\sin x} = \begin{cases} \frac{2}{\sqrt{a^2 - b^2}} \arctan \frac{a\tan \frac{x}{2} + b}{\sqrt{a^2 - b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2 - a^2}} \ln \left| \frac{a\tan \frac{x}{2} + b - \sqrt{b^2 - a^2}}{a\tan \frac{x}{2} + b + \sqrt{b^2 - a^2}} \right| + C & (a^2 < b^2) \end{cases}$$

17.
$$\int \frac{dx}{a+b\cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan\left(\sqrt{\frac{a-b}{a+b}} \tan\frac{x}{2}\right) + C & (a^2 > b^2) \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln\left|\frac{\tan\frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan\frac{x}{2} - \sqrt{\frac{a+b}{b-a}}}\right| + C & (a^2 < b^2) \end{cases}$$

18.
$$\int \frac{\mathrm{d}x}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan\left(\frac{b}{a} \tan x\right) + C$$

19.
$$\int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x - a} \right| + C$$

20.
$$\int x \sin ax dx = \frac{1}{a^2} \sin ax - \frac{1}{a} x \cos ax + C$$

21.
$$\int x^2 \sin ax dx = -\frac{1}{a}x^2 \cos ax + \frac{2}{a^2}x \sin ax + \frac{2}{a^3}\cos ax + C$$

22.
$$\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$$

23.
$$\int x^2 \cos ax dx = \frac{1}{a}x^2 \sin ax + \frac{2}{a^2}x \cos ax - \frac{2}{a^3} \sin ax + C$$

1.
$$\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 - x^2} + C$$

2.
$$\int x \arcsin \frac{x}{a} dx = (\frac{x^2}{2} - \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 - x^2} + C$$

3.
$$\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{3} \arcsin \frac{x}{a} + \frac{1}{9}(x^2 + 2a^2)\sqrt{a^2 - x^2} + C$$

4.
$$\int \arccos \frac{x}{a} dx = x \arccos \frac{x}{a} - \sqrt{a^2 - x^2} + C$$

5.
$$\int x \arccos \frac{x}{a} dx = (\frac{x^2}{2} - \frac{a^2}{4}) \arccos \frac{x}{a} - \frac{x}{4} \sqrt{a^2 - x^2} + C$$

6.
$$\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} - \frac{1}{9}(x^2 + 2a^2)\sqrt{a^2 - x^2} + C$$

7.
$$\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} - \frac{a}{2} \ln(a^2 + x^2) + C$$

8.
$$\int x \arctan \frac{x}{a} dx = \frac{1}{2}(a^2 + x^2) \arctan \frac{x}{a} - \frac{a}{2}x + C$$

9.
$$\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} - \frac{a}{6}x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$$

1.
$$\int a^x dx = \frac{1}{\ln a} a^x + C$$

2.
$$\int e^{ax} dx = \frac{1}{a} a^{ax} + C$$

3.
$$\int xe^{ax} dx = \frac{1}{a^2}(ax-1)a^{ax} + C$$

4.
$$\int x^n e^{ax} dx = \frac{1}{a} x^n e^{ax} - \frac{n}{a} \int x^{n-1} e^{ax} dx$$

5.
$$\int xa^x dx = \frac{x}{\ln a}a^x - \frac{1}{(\ln a)^2}a^x + C$$

6.
$$\int x^n a^x dx = \frac{1}{\ln a} x^n a^x - \frac{n}{\ln a} \int x^{n-1} a^x dx$$

7.
$$\int e^{ax} \sin bx dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx - b \cos bx) + C$$

8.
$$\int e^{ax} \cos bx dx = \frac{1}{a^2+b^2} e^{ax} (b \sin bx + a \cos bx) + C$$

9.
$$\int e^{ax} \sin^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \sin^{n-1} bx (a \sin bx - nb \cos bx) + \frac{n(n-1)b^2}{a^2 + b^2 n^2} \int e^{ax} \sin^{n-2} bx dx$$

10.
$$\int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) + \frac{n(n-1)b^2}{a^2 + b^2 n^2} \int e^{ax} \cos^{n-2} bx dx$$

1.
$$\int \ln x dx = x \ln x - x + C$$

$$2. \int \frac{\mathrm{d}x}{x \ln x} = \ln \left| \ln x \right| + C$$

3.
$$\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x - \frac{1}{n+1}) + C$$

4.
$$\int (\ln x)^n dx = x(\ln x)^n - n \int (\ln x)^{n-1} dx$$

5.
$$\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n - \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$$