Wizards

Standard Code Library

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1. 数论

1.1 $O(m^2 \log n)$ 线性递推

```
Given a_0, a_1, \ldots, a_{m-1}

a_n = c_0 \times a_{n-m} + \cdots + c_{m-1} \times a_{n-1}

Solve for a_n = v_0 \times a_0 + v_1 \times a_1 + \cdots + v_{m-1} \times a_{m-1}
```

```
1
   void linear_recurrence(long long n, int m, int a[], int
      \hookrightarrow c[], int p) {
 2
     long long v[M] = \{1 \% p\}, u[M << 1], msk = !!n;
     for(long long i(n); i > 1; i >>= 1) {
 3
 4
       msk <<= 1:
 5
 6
     for(long long x(0); msk; msk >>= 1, x <<= 1) {
 7
       fill_n(u, m << 1, 0);
       int b(!!(n & msk));
 8
9
       x \mid = b;
10
        if(x < m) {
11
          u[x] = 1 \% p;
12
        }else {
13
          for(int i(0); i < m; i++) {</pre>
14
            for(int j(0), t(i + b); j < m; j++, t++) {
15
              u[t] = (u[t] + v[i] * v[j]) % p;
            }
16
17
          for(int i((m << 1) - 1); i >= m; i--) {
18
19
            for(int j(0), t(i - m); j < m; j++, t++) {
              u[t] = (u[t] + c[j] * u[i]) % p;
20
21
22
          }
23
        }
24
        copy(u, u + m, v);
     }
25
     //a[n] = v[0] * a[0] + v[1] * a[1] + ... + v[m - 1] *
26
        \hookrightarrow a[m - 1].
27
     for(int i(m); i < 2 * m; i++) {</pre>
28
        a[i] = 0:
29
        for(int j(0); j < m; j++) {</pre>
30
          a[i] = (a[i] + (long long)c[j] * a[i + j - m]) % p;
31
     }
32
33
     for(int j(0); j < m; j++) {
       b[j] = 0;
        for(int i(0); i < m; i++) {</pre>
35
36
          b[j] = (b[j] + v[i] * a[i + j]) % p;
37
38
     }
39
     for(int j(0); j < m; j++) {
40
        a[j] = b[j];
41
  }
42
```

1.2 求逆元

```
void ex_gcd(long long a, long long b, long long &x, long
      \hookrightarrow long &y) {
     if (b == 0) {
       x = 1;
3
       y = 0;
4
5
       return;
7
     long long xx, yy;
8
     ex_gcd(b, a % b, xx, yy);
9
     y = xx - a / b * yy;
10
     x = yy;
11
   }
12
   long long inv(long long x, long long MODN) {
13
14
     long long inv_x, y;
     ex_gcd(x, MODN, inv_x, y);
15
     return (inv_x % MODN + MODN) % MODN;
16
```

17 }

1.3 中国剩余定理

```
//返回 (ans, M), 其中 ans 是模 M 意义下的解
   std::pair<long long, long long> CRT(const std::vector<long
     \hookrightarrow long>& m, const std::vector<long long>& a) {
3
    long long M = 1, ans = 0;
    int n = m.size();
4
5
    for (int i = 0; i < n; i++) M *= m[i];</pre>
6
    for (int i = 0; i < n; i++) {
       ans = (ans + (M / m[i]) * a[i] % M * inv(M / m[i],
         → m[i])) % M; // 可能需要大整数相乘取模
8
9
    return std::make_pair(ans, M);
10 }
```

1.4 素性测试

```
int strong_pseudo_primetest(long long n,int base) {
2
       long long n2=n-1,res;
3
        int s=0:
 4
        while(n2\%2==0) n2>>=1,s++;
 5
        res=powmod(base,n2,n);
 6
        if((res==1)||(res==n-1)) return 1;
 7
        s--;
 8
        while(s>=0) {
 g
            res=mulmod(res,res,n);
10
            if(res==n-1) return 1;
11
12
       }
13
        return 0; // n is not a strong pseudo prime
14
15
   int isprime(long long n) {
16
     static LL testNum[]={2,3,5,7,11,13,17,19,23,29,31,37};
17
        \rightarrow \lim[]=\{4,0,1373653LL,25326001LL,25000000000LL,21523028987\}
         \rightarrow 3474749660383LL,341550071728321LL,0,0,0,0);
     if(n<2||n==3215031751LL) return 0;
18
     for(int i=0;i<12;++i){</pre>
19
20
        if(n<lim[i]) return 1;</pre>
        if(strong_pseudo_primetest(n,testNum[i])==0) return 0;
21
22
     }
23
     return 1;
24
```

1.5 质因数分解

```
int ansn; LL ans[1000];
   LL func(LL x,LL n){ return(mod_mul(x,x,n)+1)%n; }
   LL Pollard(LL n){
     LL i,x,y,p;
     if(Rabin_Miller(n)) return n;
     if(!(n&1)) return 2;
6
 7
     for(i=1;i<20;i++){</pre>
8
       x=i; y=func(x,n); p=gcd(y-x,n);
9
        while(p==1) {x=func(x,n); y=func(func(y,n),n);
           \hookrightarrow p=\gcd((y-x+n)%n,n)%n;
10
        if(p==0||p==n) continue;
11
        return p;
12
13
14
   void factor(LL n){
     LL x;
16
     x=Pollard(n);
17
     if(x==n){ ans[ansn++]=x; return; }
     factor(x), factor(n/x);
18
19
   }
```

1.6 佩尔方程

```
import java.math.BigInteger;
1
  import java.util.Scanner;
2
3/a[n]=(g[n]+a[0])/h[n]
4 //g[n] = a[n-1] *h[n-1] -g[n-1]
5 / h[n] = (N-g[n]*g[n])/h[n-1]
6 //p[n] = a[n-1] *p[n-1] + p[n-2]
  //q[n]=a[n-1]*q[n-1]+q[n-2]
  //so:
8
9
  //p[n]*q[n-1]-p[n-1]*q[n]=(-1)^(n+1);
10 //p[n]^2-N*q[n]^2=(-1)^(n+1)*h[n+1];
11
   public class Main {
12
       public static BigInteger p, q;
       public static void solve(int n) {
13
           BigInteger N, p1, p2, q1, q2, a0, a1, a2, g1, g2,
14
              \hookrightarrow h1, h2;
15
           g1 = q2 = p1 = BigInteger.ZERO;
           h1 = q1 = p2 = BigInteger.ONE;
16
           a0 = a1 =
17
              → BigInteger.valueOf((long)Math.sqrt(1.0*n));
18
           N = BigInteger.valueOf(n);
           while (true) {
19
20
               g2 = a1.multiply(h1).subtract(g1);
               h2 = N.subtract(g2.pow(2)).divide(h1);
21
22
               a2 = g2.add(a0).divide(h2);
23
               p = a1.multiply(p2).add(p1);
24
               q = a1.multiply(q2).add(q1);
25
               if
                  \hookrightarrow == 0) return;
         g1 = g2;h1 = h2;a1 = a2;
26
27
               p1 = p2; p2 = p;
               q1 = q2; q2 = q;
28
           }
29
30
       }
31
       public static void main(String[] args) {
32
           Scanner cin = new Scanner(System.in);
33
           int t=cin.nextInt();
34
35
           while (t--!=0) {
               solve(cin.nextInt());
36
               System.out.println(p + " " + q);
37
38
39
40 }
```

1.7 二次剩余

```
1 // x^2 = a (mod p), 0 <= a < p, 返回 true or false 代表
     →是否存在解
  // p 必须是质数, 若是多个单次质数的乘积, 可以分别
     →求解再用 CRT 合并
3 // 复杂度为 D(log n)
  void multiply(11 &c, 11 &d, 11 a, 11 b, 11 w) {
      int cc = (a * c + b * d % MOD * w) % MOD;
      int dd = (a * d + b * c) % MOD;
7
      c = cc, d = dd;
8
  1
9
10
  bool solve(int n, int &x) {
11
      if (MOD == 2) return x = 1, true;
      if (power(n, MOD / 2, MOD) == MOD - 1) return false;
12
      11 c = 1, d = 0, b = 1, a, w;
13
14
      // finding a such that a^2 - n is not a square
      do { a = rand() % MOD;
15
          w = (a * a - n + MOD) \% MOD;
16
```

```
17
           if (w == 0) return x = a, true;
18
       } while (power(w, MOD / 2, MOD) != MOD - 1);
19
       for (int times = (MOD + 1) / 2; times; times >>= 1) {
20
           if (times & 1) multiply(c, d, a, b, w);
21
           multiply(a, b, a, b, w);
22
23
       // x = (a + sqrt(w)) ^ ((p + 1) / 2)
24
       return x = c, true;
25
```

1.8 一元三次方程

```
double a(p[3]), b(p[2]), c(p[1]), d(p[0]);
   double k(b / a), m(c / a), n(d / a);
   double p(-k * k / 3. + m);
 3
   double q(2. * k * k * k / 27 - k * m / 3. + n);
   Complex omega[3] = \{Complex(1, 0), Complex(-0.5, 0.5 *
      \rightarrow sqrt(3)), Complex(-0.5, -0.5 * sqrt(3))};
   Complex r1, r2;
 6
   double delta(q * q / 4 + p * p * p / 27);
   if (delta > 0) {
       r1 = cubrt(-q / 2. + sqrt(delta));
10
       r2 = cubrt(-q / 2. - sqrt(delta));
11
   } else {
12
       r1 = pow(-q / 2. + pow(Complex(delta), 0.5), 1. / 3);
13
       r2 = pow(-q / 2. - pow(Complex(delta), 0.5), 1. / 3);
14
15
   for(int _(0); _ < 3; _++) {
       Complex x = -k / 3. + r1 * omega[_ * 1] + r2 * omega[_
          \hookrightarrow * 2 % 3];
17 }
```

1.9 线下整点

```
//\sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor, n, m, a, b > 0
LL solve(LL n,LL a,LL b,LL m){
  if (b==0) return n*(a/m);
  if(a>=m) return n*(a/m)+solve(n,a%m,b,m);
  if (b>=m) return (n-1)*n/2*(b/m)+solve(n,a,b,m,m);
  return solve((a+b*n)/m,(a+b*n)%m,m,b);
```

1.10 线性同余不等式

```
\ensuremath{//} Find the minimal non-negtive solutions for
      // 0 \le d, l, r \le m; l \le r, O(\log n)
   11 cal(11 m, 11 d, 11 1, 11 r) {
 3
       if (1 == 0) return 0;
       if (d == 0) return MXL; // 无解
 6
       if (d * 2 > m) return cal(m, m - d, m - r, m - 1);
       if ((1 - 1) / d < r / d) return (1 - 1) / d + 1;
 7
8
       ll k = cal(d, (-m % d + d) % d, 1 % d, r % d);
 9
       return k == MXL ? MXL : (k * m + 1 - 1) / d + 1; // \pi
          →解 2
10
11
12
   // return all x satisfying l1<=x<=r1 and
      \hookrightarrow 12<=(x*mul+add)%LIM<=r2
  // here LIM = 2^32 so we use UI instead of "%".
   // O(\log p + \#solutions)
15
   struct Jump {
       UI val, step;
16
17
       Jump(UI val, UI step) : val(val), step(step) { }
18
       Jump operator + (const Jump & b) const {
           return Jump(val + b.val, step + b.step); }
19
20
       Jump operator - (const Jump & b) const {
           return Jump(val - b.val, step + b.step);
21
22
23 inline Jump operator * (UI x, const Jump & a) {
```

```
return Jump(x * a.val, x * a.step);
24
25
  }
26
   vector<UI> solve(UI 11, UI r1, UI 12, UI r2, pair<UI, UI>

    muladd) {
       UI mul = muladd.first, add = muladd.second, w = r2 -
          Jump up(mul, 1), dn(-mul, 1);
28
       UI s(11 * mul + add);
29
30
       Jump lo(r2 - s, 0), hi(s - 12, 0);
31
       function<void(Jump &, Jump &)> sub = [&](Jump & a,
          \hookrightarrow Jump & b) {
            if (a.val > w) {
32
33
                UI t(((long long)a.val - max(0ll, w + 1ll -
                   \hookrightarrow b.val)) / b.val);
34
                a = a - t * b:
            7
35
       }:
36
37
       sub(lo, up), sub(hi, dn);
38
       while (up.val > w || dn.val > w) {
            sub(up, dn); sub(lo, up);
39
40
            sub(dn, up); sub(hi, dn); }
       assert(up.val + dn.val > w);
41
       vector<UI> res;
42
43
        Jump bg(s + mul * min(lo.step, hi.step), min(lo.step,

    hi.step));
       while (bg.step <= r1 - 11) {</pre>
44
45
            if (12 <= bg.val && bg.val <= r2)</pre>
46
                res.push_back(bg.step + 11);
            if (12 <= bg.val - dn.val && bg.val - dn.val <=
47
               \hookrightarrow r2) {
48
                bg = bg - dn;
            } else bg = bg + up;
49
50
       } return res;
51
  | }
```

1.11 组合数取模

```
1 LL prod=1,P;
2
   pair<LL,LL> comput(LL n,LL p,LL k){
       if(n<=1)return make_pair(0,1);</pre>
3
4
       LL ans=1,cnt=0;
5
       ans=pow(prod,n/P,P);
6
       cnt=n/p;
7
       pair<LL,LL>res=comput(n/p,p,k);
8
       cnt+=res.first;
9
       ans=ans*res.second%P:
10
       for(int i=n-n%P+1;i<=n;i++)if(i%p){</pre>
11
            ans=ans*i%P;
12
13
       }
14
       return make_pair(cnt,ans);
15
  }
  pair<LL,LL> calc(LL n,LL p,LL k){
16
17
       prod=1;P=pow(p,k,1e18);
       for(int i=1;i<P;i++)if(i%p)prod=prod*i%P;</pre>
18
19
       pair<LL,LL> res=comput(n,p,k);
20
   // res.second=res.second*pow(p,res.first%k,P)%P;
  // res.first-=res.first%k;
21
22
       return res;
23 }
24
  LL calc(LL n,LL m,LL p,LL k){
25
       pair<LL,LL>A,B,C;
26
       LL P=pow(p,k,1e18);
       A=calc(n,p,k);
28
       B=calc(m,p,k);
29
       C=calc(n-m,p,k);
30
31
       ans=pow(p,A.first-B.first-C.first,P);
32

→ ans=ans*A.second%P*inv(B.second,P)%P*inv(C.second,P)%P;

59
33
       return ans;
34 }
```

1.12 Schreier-Sims

```
struct Perm{
 2
     vector<int> P; Perm() {} Perm(int n) { P.resize(n); }
 3
     Perm inv()const{
 4
       Perm ret(P.size()):
       for(int i = 0; i < int(P.size()); ++i) ret.P[P[i]] =</pre>
 6
       return ret;
 7
8
     int &operator [](const int &dn){ return P[dn]; }
9
     void resize(const size_t &sz){ P.resize(sz); }
10
     size_t size()const{ return P.size(); }
     const int &operator [](const int &dn)const{ return
11
        \hookrightarrow P[dn]; }
12
   };
13
   Perm operator *(const Perm &a, const Perm &b){
     Perm ret(a.size());
14
15
     for(int i = 0; i < (int)a.size(); ++i) ret[i] = b[a[i]];</pre>
16
     return ret;
17
18
   typedef vector<Perm> Bucket;
   typedef vector<int> Table;
19
   typedef pair<int,int> PII;
20
   int n, m;
   vector<Bucket> buckets, bucketsInv; vector<Table>
      \hookrightarrow \texttt{lookupTable;}
23
   int fastFilter(const Perm &g, bool addToGroup = true) {
     int n = buckets.size();
24
25
     Perm p(g);
26
     for(int i = 0; i < n; ++i){
27
       int res = lookupTable[i][p[i]];
28
       if(res == -1){
29
         if (addToGroup) {
30
           buckets[i].push_back(p);
               → bucketsInv[i].push_back(p.inv());
           lookupTable[i][p[i]] = (int)buckets[i].size() - 1;
31
32
         }
33
         return i;
34
       }
       p = p * bucketsInv[i][res];
35
36
37
     return -1;
38
39
   long long calcTotalSize(){
40
     long long ret = 1;
     for(int i = 0; i < n; ++i) ret *= buckets[i].size();</pre>
41
42
     return ret:
43
   }
   bool inGroup(const Perm &g){ return fastFilter(g, false)
      45
   void solve(const Bucket &gen,int _n){// m perm[0..n - 1]s
46
     n = _n, m = gen.size();
47
     {//clear all
48
       vector<Bucket> _buckets(n); swap(buckets, _buckets);
       vector<Bucket> _bucketsInv(n); swap(bucketsInv,

→ _bucketsInv);
50
       vector<Table> _lookupTable(n); swap(lookupTable,
          → _lookupTable);
51
52
     for(int i = 0; i < n; ++i){
53
       lookupTable[i].resize(n):
       fill(lookupTable[i].begin(), lookupTable[i].end(),
54
          \hookrightarrow -1):
     }
55
56
     Perm id(n):
57
     for(int i = 0; i < n; ++i) id[i] = i;</pre>
     for(int i = 0; i < n; ++i){
       buckets[i].push_back(id); bucketsInv[i].push_back(id);
       lookupTable[i][i] = 0;
61
62
     for(int i = 0; i < m; ++i) fastFilter(gen[i]);</pre>
```

```
queue<pair<PII,PII> > toUpdate;
63
     for(int i = 0; i < n; ++i)
64
65
       for(int j = i; j < n; ++j)
          for(int k = 0; k < (int)buckets[i].size(); ++k)</pre>
66
            for(int 1 = 0; 1 < (int)buckets[j].size(); ++1)</pre>
67
              toUpdate.push(make_pair(PII(i,k), PII(j,1)));
68
69
     while(!toUpdate.empty()){
70
        PII a = toUpdate.front().first, b =
           \hookrightarrow toUpdate.front().second;
71
        toUpdate.pop();
72
        int res = fastFilter(buckets[a.first][a.second] *

    buckets[b.first][b.second]);
73
        if(res==-1) continue:
74
        PII newPair(res, (int)buckets[res].size() - 1);
75
        for(int i = 0; i < n; ++i)
          for(int j = 0; j < (int)buckets[i].size(); ++j){</pre>
76
            if(i <= res) toUpdate.push(make_pair(PII(i, j),</pre>
77
               → newPair));
78
            if(res <= i) toUpdate.push(make_pair(newPair,</pre>
               \hookrightarrow PII(i, j)));
79
     }
80
81 }
```

2. 代数

2.1 快速傅里叶变换

```
1 // n 必须是 2 的次幂
  void fft(Complex a[], int n, int f) {
3
    for (int i = 0; i < n; ++i)
4
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
5
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
6
7
       Complex w = Complex(1, 0);
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
       for (int p = i \ll 1, j = 0; j < n; j += p) {
         for (int k = 0; k < i; ++k) {
10
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11
           a[j + k] = x + y; a[j + k + i] = x - y;
12
13
14
       }
15
    }
16
  }
```

2.2 分治卷积

```
1 // n 必须是 2 的次幂
  void fft(Complex a[], int n, int f) {
3
     for (int i = 0; i < n; ++i)
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
       Complex w = Complex(1, 0);
7
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
9
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
10
         for (int k = 0; k < i; ++k) {
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
       }
     }
15
16 }
```

2.3 快速数论变换

```
1 // n 必须是 2 的次幂
2 void fft(Complex a[], int n, int f) {
3 for (int i = 0; i < n; ++i)
4 if (R[i] < i) swap(a[i], a[R[i]]);
5 for (int i = 1, h = 0; i < n; i <<= 1, h++) {
```

```
Complex wn = Complex(cos(pi / i), f * sin(pi / i));
6
 7
       Complex w = Complex(1, 0);
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
       for (int p = i << 1, j = 0; j < n; j += p) {
9
         for (int k = 0; k < i; ++k) {
10
11
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
12
           a[j + k] = x + y; a[j + k + i] = x - y;
13
14
       }
15
     }
16
```

2.4 快速沃尔什变换

```
void FWT(LL a[],int n,int ty){ //the length is 2^n
     for(int d=1;d<n;d<<=1){</pre>
 3
        for(int m=(d<<1), i=0; i< n; i+=m){
          if(tv==1){
 5
            for(int j=0; j<d; j++){</pre>
               LL x=a[i+j], y=a[i+j+d];
 7
               a[i+j]=x+y;
 8
               a[i+j+d]=x-y;
 q
                        //and:a[i+j]=x+y; or:a[i+j+d]=x+y;
10
            }
11
          }else{
12
            for(int j=0; j<d; j++){</pre>
13
              LL x=a[i+j], y=a[i+j+d];
14
               a[i+j]=(x+y)/2;
15
               a[i+j+d]=(x-y)/2;
                       //and:a[i+j]=x-y; or:a[i+j+d]=y-x;
16
17
18
          }
19
        }
20
     }
21
```

2.5 自适应辛普森积分

```
namespace adaptive_simpson {
 2
     template<typename function>
     inline double area(function f, const double &left, const
 3
        double mid = (left + right) / 2;
 4
       return (right - left) * (f(left) + 4 * f(mid) +
          \hookrightarrow f(right)) / 6;
 6
 7
8
     template<typename function>
9
     inline double simpson(function f, const double &left,
        → const double &right, const double &eps, const

    double &area_sum) {
10
       double mid = (left + right) / 2;
11
       double area_left = area(f, left, mid);
       double area_right = area(f, mid, right);
12
       double area_total = area_left + area_right;
13
14
       if (fabs(area_total - area_sum) <= 15 * eps) {</pre>
15
         return area_total + (area_total - area_sum) / 15;
16
17
       return simpson(f, left, right, eps / 2, area_left) +
          \hookrightarrow simpson(f, mid, right, eps / 2, area_right);
18
19
20
     template<typename function>
     inline double simpson(function f, const double &left,
        return simpson(f, left, right, eps, area(f, left,
22
          \hookrightarrow right));
23
24
  }
```

2.6 单纯形

```
const double eps = 1e-8;
   // max{c * x | Ax <= b, x >= 0} 的解, 无解返回空的
      → vector, 否则就是解.
   vector<double> simplex(vector<vector<double> > &A,
3

    vector<double> b, vector<double> c) {
     int n = A.size(), m = A[0].size() + 1, r = n, s = m - 1;
     vector < vector < double > D(n + 2, vector < double > (m + 1));
     vector<int> ix(n + m);
7
     for(int i = 0; i < n + m; i++) {</pre>
       ix[i] = i;
8
9
     }
10
     for(int i = 0; i < n; i++) {</pre>
       for(int j = 0; j < m - 1; j++) {
11
12
         D[i][j] = -A[i][j];
13
       D[i][m - 1] = 1;
14
       D[i][m] = b[i];
15
       if (D[r][m] > D[i][m]) {
16
17
         r = i;
18
     }
19
20
     for(int j = 0; j < m - 1; j++) {
21
22
       D[n][j] = c[j];
23
24
     D[n + 1][m - 1] = -1;
25
     for(double d; ;) {
26
       if (r < n) {
         swap(ix[s], ix[r + m]);
27
28
         D[r][s] = 1. / D[r][s];
29
         for(int j = 0; j \le m; j++) {
30
            if (j != s) {
              D[r][j] *= -D[r][s];
31
32
33
34
         for(int i = 0; i <= n + 1; i++) {
35
            if (i != r) {
36
              for(int j = 0; j \le m; j++) {
37
                if (j != s) {
                  D[i][j] += D[r][j] * D[i][s];
38
39
40
41
              D[i][s] *= D[r][s];
42
         }
43
       }
44
       r = -1, s = -1;
45
       for(int j = 0; j < m; j++) {
46
47
          if (s < 0 || ix[s] > ix[j]) {
48
            if (D[n + 1][j] > eps || D[n + 1][j] > -eps &&
               \hookrightarrow D[n][j] > eps) {
49
              s = j;
            }
50
51
52
53
        if (s < 0) {
         break;
54
55
56
       for(int i = 0; i < n; i++) {</pre>
57
         if (D[i][s] < -eps) {</pre>
            if (r < 0 \mid | (d = D[r][m] / D[r][s] - D[i][m] /
58
               \hookrightarrow D[i][s]) < -eps
50
              || d < eps && ix[r + m] > ix[i + m]) {
60
61
              r = i;
62
            }
         }
63
64
65
       if (r < 0) {
66
         return vector<double> ();
67
```

```
69
70
     if (D[n + 1][m] < -eps) {
71
       return vector<double> ();
72
73
     vector<double> x(m - 1);
74
75
     for(int i = m; i < n + m; i++) {
76
       if (ix[i] < m - 1) {</pre>
77
          x[ix[i]] = D[i - m][m];
78
79
     }
80
     return x;
81
```

3. 计算几何

3.1 二维

3.1.1 点类

```
int sign(DB x) {
2
    return (x > eps) - (x < -eps);
3
   DB msqrt(DB x) {
    return sign(x) > 0 ? sqrt(x) : 0;
6
7
   struct Point {
     DB x, y;
8
     Point rotate(DB ang) const { // 逆时针旋转 ang 弧度
9
       return Point(cos(ang) * x - sin(ang) * y, cos(ang) * y
10
          \hookrightarrow + sin(ang) * x);
11
     Point turn90() const { // 逆时针旋转 90 度
12
13
       return Point(-y, x);
14
15
     Point unit() const {
16
       return *this / len();
17
18
   };
   DB dot(const Point& a, const Point& b) {
19
20
    return a.x * b.x + a.y * b.y;
21
   DB det(const Point& a, const Point& b) {
23
     return a.x * b.y - a.y * b.x;
24
   #define cross(p1,p2,p3)
      \hookrightarrow ((\texttt{p2.x-p1.x})*(\texttt{p3.y-p1.y}) - (\texttt{p3.x-p1.x})*(\texttt{p2.y-p1.y}))
   #define crossOp(p1,p2,p3) sign(cross(p1,p2,p3))
   bool isLL(const Line& 11, const Line& 12, Point& p) { //
      →直线与直线交点
28
     DB s1 = det(12.b - 12.a, 11.a - 12.a),
        s2 = -det(12.b - 12.a, 11.b - 12.a);
29
30
     if (!sign(s1 + s2)) return false;
31
     p = (11.a * s2 + 11.b * s1) / (s1 + s2);
32
     return true;
33
   bool onSeg(const Line& 1, const Point& p) { // 点在线段
35
     return sign(det(p - 1.a, 1.b - 1.a)) == 0 && sign(dot(p
        \hookrightarrow - 1.a, p - 1.b)) <= 0;
36
37
   Point projection(const Line & 1, const Point& p) {
     return 1.a + (1.b - 1.a) * (dot(p - 1.a, 1.b - 1.a) /
38
        \hookrightarrow (1.b - 1.a).len2());
39
   | DB disToLine(const Line& 1, const Point& p) { // 点到 *
40
      →直线 * 距离
     return fabs(det(p - 1.a, 1.b - 1.a) / (1.b -
41
        \hookrightarrow 1.a).len());
42 }
```

```
43 DB disToSeg(const Line& 1, const Point& p) { // 点到线段
       →距离
      return sign(dot(p - 1.a, 1.b - 1.a)) * sign(dot(p - 1.b,
44
         \hookrightarrow l.a - l.b)) == 1 ? disToLine(l, p) : std::min((p -
         \hookrightarrow 1.a).len(), (p - 1.b).len());
45 }
    // 圆与直线交点
46
    bool isCL(Circle a, Line 1, Point& p1, Point& p2) {
47
      DB x = dot(1.a - a.o, 1.b - 1.a),
48
         y = (1.b - 1.a).len2(),
49
         d = x * x - y * ((1.a - a.o).len2() - a.r * a.r);
50
      if (sign(d) < 0) return false;</pre>
51
      Point p = 1.a - ((1.b - 1.a) * (x / y)), delta = (1.b - 1.a) * (x / y))
52
         \hookrightarrow 1.a) * (msqrt(d) / y);
53
      p1 = p + delta; p2 = p - delta;
54
      return true;
55
   ۱,
    |//圆与圆的交面积
56
    DB areaCC(const Circle& c1, const Circle& c2) {
57
58
      DB d = (c1.o - c2.o).len();
      if (sign(d - (c1.r + c2.r)) >= 0) return 0;
59
60
      if (sign(d - std::abs(c1.r - c2.r)) \le 0) {
        DB r = std::min(c1.r, c2.r);
61
        return r * r * PI;
62
63
      DB x = (d * d + c1.r * c1.r - c2.r * c2.r) / (2 * d),
       t1 = acos(x / c1.r), t2 = acos((d - x) / c2.r);
65
      return c1.r * c1.r * t1 + c2.r * c2.r * t2 - d * c1.r *
66
         \hookrightarrow \sin(t.1):
   }
67
   // 圆与圆交点
68
    bool isCC(Circle a, Circle b, P& p1, P& p2) {
69
70
      DB s1 = (a.o - b.o).len();
71
      if (sign(s1 - a.r - b.r) > 0 \mid \mid sign(s1 - std::abs(a.r - a.r - b.r))
         \hookrightarrow b.r)) < 0) return false;
      DB s2 = (a.r * a.r - b.r * b.r) / s1;
72
      DB aa = (s1 + s2) * 0.5, bb = (s1 - s2) * 0.5;
73
      P \circ = (b.o - a.o) * (aa / (aa + bb)) + a.o;
74
75
      P delta = (b.o - a.o).unit().turn90() * msqrt(a.r * a.r
         \rightarrow - aa * aa):
76
      p1 = o + delta, p2 = o - delta;
77
      return true;
78 }
    // 求点到圆的切点,按关于点的顺时针方向返回两个点
79
    bool tanCP(const Circle &c, const Point &p0, Point &p1,
80
      → Point &p2) {
81
      double x = (p0 - c.o).len2(), d = x - c.r * c.r;
      if (d < eps) return false; // 点在圆上认为没有切点
82
      Point p = (p0 - c.o) * (c.r * c.r / x);
83
      Point delta = ((p0 - c.o) * (-c.r * sqrt(d) /
         \rightarrow x)).turn90();
85
      p1 = c.o + p + delta;
      p2 = c.o + p - delta;
86
87
      return true:
88 }
    // 求圆到圆的外共切线, 按关于 c1.o 的顺时针方向返
    vector<Line> extanCC(const Circle &c1, const Circle &c2) {
90
      vector<Line> ret;
91
92
      if (sign(c1.r - c2.r) == 0) {
        Point dir = c2.o - c1.o;
        dir = (dir * (c1.r / dir.len())).turn90();
94
95
        ret.push_back(Line(c1.o + dir, c2.o + dir));
        ret.push_back(Line(c1.o - dir, c2.o - dir));
96
      } else {
97
98
        Point p = (c1.0 * -c2.r + c2.o * c1.r) / (c1.r - c2.r + c2.o * c1.r) / (c1.r - c2.r + c2.o * c1.r)
           \hookrightarrow c2.r):
99
        Point p1, p2, q1, q2;
        if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) {
100
          if (c1.r < c2.r) swap(p1, p2), swap(q1, q2);
101
102
          ret.push_back(Line(p1, q1));
103
          ret.push_back(Line(p2, q2));
104
```

```
106
      return ret;
107
    // 求圆到圆的内共切线, 按关于 c1.o 的顺时针方向返
108
    std::vector<Line> intanCC(const Circle &c1, const Circle
109
       \hookrightarrow &c2) {
110
      std::vector<Line> ret;
111
      Point p = (c1.0 * c2.r + c2.o * c1.r) / (c1.r + c2.r);
112
      Point p1, p2, q1, q2;
      if (tanCP(c1, p, p1, p2) && tanCP(c2, p, q1, q2)) { //
113
         →两圆相切认为没有切线
        ret.push_back(Line(p1, q1));
114
115
        ret.push_back(Line(p2, q2));
      }
116
117
      return ret;
118 }
    bool contain(vector<Point> polygon, Point p) { // 判断点
119
       →p 是否被多边形包含,包括落在边界上
      int ret = 0, n = polygon.size();
120
      for(int i = 0; i < n; ++ i) {
121
122
        Point u = polygon[i], v = polygon[(i + 1) % n];
        if (onSeg(Line(u, v), p)) return true; // Here I
123
124
        if (sign(u.y - v.y) \le 0) swap(u, v);
125
        if (sign(p.y - u.y) > 0 \mid \mid sign(p.y - v.y) \le 0)
           \hookrightarrow continue;
126
        ret += sign(det(p, v, u)) > 0;
127
      }
128
      return ret & 1;
129
    // 用半平面 (q1,q2) 的逆时针方向去切凸多边形
130
    std::vector<Point> convexCut(const std::vector<Point>&ps,
       \hookrightarrow Point q1, Point q2) {
      std::vector<Point> qs; int n = ps.size();
132
      for (int i = 0; i < n; ++i) {
133
        Point p1 = ps[i], p2 = ps[(i + 1) % n];
134
135
        int d1 = crossOp(q1,q2,p1), d2 = crossOp(q1,q2,p2);
136
        if (d1 \ge 0) qs.push_back(p1);
        if (d1 * d2 < 0) qs.push_back(isSS(p1, p2, q1, q2));
137
138
139
      return qs;
140
    // 求凸包
141
142
    std::vector<Point> convexHull(std::vector<Point> ps) {
143
      int n = ps.size(); if (n <= 1) return ps;</pre>
      std::sort(ps.begin(), ps.end());
144
      std::vector<Point> qs;
145
      for (int i = 0; i < n; qs.push_back(ps[i ++]))</pre>
146
        while (qs.size() > 1 && sign(det(qs[qs.size() - 2],
           \hookrightarrow qs.back(), ps[i])) <= 0)
148
          qs.pop_back();
149
      for (int i = n - 2, t = qs.size(); i \ge 0;

    qs.push_back(ps[i --]))
150
        while ((int)qs.size() > t && sign(det(qs[qs.size() -
           \hookrightarrow 2], qs.back(), ps[i])) <= 0)
151
          qs.pop_back();
152
      return qs;
153
```

3.1.2 凸包

```
for (; j \ge 2 \&\& sign(det(shell[j-1] - shell[j-2],
9
10
                 p[i] - shell[j-2])) \le 0; --j)

    shell.pop_back();
11
         shell.push_back(p[i]);
12
    }
13
14
     void make_convex() {
       std::sort(a.begin(), a.end());
15
16
       make_shell(a, lower);
17
       std::reverse(a.begin(), a.end());
18
       make_shell(a, upper);
19
       a = lower; a.pop_back();
20
       a.insert(a.end(), upper.begin(), upper.end());
21
       if ((int)a.size() >= 2) a.pop_back();
22
       n = a.size();
     7
23
     void init(const std::vector<Point>& _a) {
24
25
       clear(a); a = _a; n = a.size();
26
       make_convex();
27
     void read(int _n) { // Won't make convex.
28
29
       clear(a); n = _n; a.resize(n);
       for (int i = 0; i < n; i++)
30
31
         a[i].read();
32
33
     std::pair<DB, int> get_tangent(
34
         const std::vector<Point>& convex, const Point& vec)
       int l = 0, r = (int)convex.size() - 2;
35
       assert(r >= 0);
36
37
       for (; 1 + 1 < r; ) {
         int mid = (1 + r) / 2;
38
39
         if (sign(det(convex[mid + 1] - convex[mid], vec)) >
           \rightarrow 0)
           r = mid:
40
         else 1 = mid;
41
42
       return std::max(std::make_pair(det(vec, convex[r]),
43
44
           std::make_pair(det(vec, convex[0]), 0));
45
     int binary_search(Point u, Point v, int 1, int r) {
46
47
       int s1 = sign(det(v - u, a[1 % n] - u));
       for (; 1 + 1 < r; ) {
48
49
         int mid = (1 + r) / 2;
50
         int smid = sign(det(v - u, a[mid % n] - u));
         if (smid == s1) l = mid;
51
52
         else r = mid;
53
       return 1 % n;
     }
55
     // 求凸包上和向量 vec 叉积最大的点,返回编号,共
56
       →线的多个切点返回任意一个
57
     int get_tangent(Point vec) {
58
       std::pair<DB, int> ret = get_tangent(upper, vec);
       ret.second = (ret.second + (int)lower.size() - 1) % n;
59
60
       ret = std::max(ret, get_tangent(lower, vec));
61
       return ret.second;
    }
62
     // 求凸包和直线 u, v 的交点, 如果不相交返回 false,
63
       →如果有则是和 (i, next(i)) 的交点,交在点上不
        →确定返回前后两条边其中之一
     bool get_intersection(Point u, Point v, int &i0, int
64
       int p0 = get_tangent(u - v), p1 = get_tangent(v - u);
65
       if (sign(det(v - u, a[p0] - u)) * sign(det(v - u,
66
         \hookrightarrow \texttt{a[p1] - u)) <= 0) \ \{
67
         if (p0 > p1) std::swap(p0, p1);
68
         i0 = binary_search(u, v, p0, p1);
         i1 = binary_search(u, v, p1, p0 + n);
69
70
         return true;
71
72
       else return false;
```

```
73 }
74 };
```

3.1.3 凸包最近点对

```
//判断点是否在多边形内
 2
   int isPointInPolygon(point p, point *a, int n) {
 3
       int cnt = 0:
       for(int i=0; i<n; ++i) {</pre>
 5
            if(OnSegment(p, a[i], a[(i+1)%n])) return -1;
6
            double k = cross(a[(i+1)%n]-a[i], p-a[i]);
           double d1 = a[i].y - p.y;
 7
8
       double d2 = a[(i+1)].y - p.y;
9
            if(k>0 &&d1<=0 &&d2>0) cnt++;
10
            if(k<0 &&d2<=0 &&d1>0) cnt++;
            //k==0, 点和线段共线的情况不考虑
11
12
13
       if(cnt&1)return 1;
14
       return 0;
15
   }
   //判断凸包是否相离
16
   bool two_getaway_ConvexHull(point *cha, int n1, point
17
      \hookrightarrow *chb, int m1) {
       if(n1==1 && m1==1) {
18
19
            if(cha[0]==chb[0])
20
                return false;
21
       } else if(n1==1 && m1==2) {
22
            if(OnSegment(cha[0], chb[0], chb[1]))
23
                return false;
24
       } else if(n1==2 && m1==1) {
25
            if(OnSegment(chb[0], cha[0], cha[1]))
26
                return false;
27
       } else if(n1==2 && m1==2) {
28
            if(SegmentIntersection(cha[0], cha[1], chb[0],
              \hookrightarrow chb[1])
29
               return false;
30
       } else if(n1==2) {
31
            for(int i=0; i<n1; ++i)</pre>
32
                if(isPointInPolygon(cha[i], chb, m1))
33
                    return false:
34
       } else if(m1==2) {
35
            for(int i=0; i<m1; ++i)</pre>
36
                if(isPointInPolygon(chb[i], cha, n1))
37
                    return false;
       } else {
38
           for(int i=0; i<n1; ++i) {</pre>
39
40
                for(int j=0; j<m1; ++j) {</pre>
                    if (SegmentIntersection(cha[i],
                       \hookrightarrow cha[(i+1)%n1], chb[j],
                       \hookrightarrow chb[(j+1)\%m1]))
42
                        return false;
43
44
            for(int i=0; i<n1; ++i)</pre>
45
                if(isPointInPolygon(cha[i], chb, m1))
46
47
                    return false;
48
            for(int i=0; i<m1; ++i)</pre>
49
                if(isPointInPolygon(chb[i], cha, n1))
50
                    return false;
52
       return true;
53
   }
   //旋转卡壳求两个凸包最近距离
   double solve(point *P, point *Q, int n, int m) {
       if (n==1 && m==1) {
56
           return length(P[0] - Q[0]);
57
58
       } else if(n==1 && m==2) {
           return DistanceToSegment(P[0], Q[0], Q[1]);
59
60
       } else if(n==2 && m==1) {
61
           return DistanceToSegment(Q[0], P[0], P[1]);
       } else if(n==2 && m==2) {
62
            return SegmentToSegment(P[0], P[1], Q[0], Q[1]);
63
```

```
65
66
       int yminP = 0, ymaxQ = 0;
67
       for(int i=0; i<n; ++i) if(P[i].y < P[yminP].y) yminP =</pre>
       for(int i=0; i < m; ++i) if(Q[i].y > Q[ymaxQ].y) ymaxQ =
68
       P[n] = P[0];
69
70
       Q[n] = Q[0];
71
       double INF2 = 1e100;
72
       double arg, ans = INF2;
73
74
       for(int i=0; i<n; ++i) {</pre>
           //当叉积负正转正时,说明点 ymaxQ 就是对踵点
75
           while((arg=cross(P[yminP] - P[yminP+1],Q[ymaxQ+1]
76
              \hookrightarrow - Q[ymaxQ])) < -eps)
               ymaxQ = (ymaxQ+1)%m;
77
           double ret:
78
79
           if(arg > eps) { //卡住第二个凸包上的点。
80
               ret = DistanceToSegment(Q[ymaxQ], P[yminP],
81
                  \hookrightarrow P[yminP+1]);
                ans = min(ans,ret);
82
           } else { //arg==0, 卡住第二个凸包的边
83
85
                ans = min(ans,ret);
           }
86
87
           yminP = (yminP+1)%n;
88
89
       return ans:
90
   }
91
   double mindis_twotubao(point *P, point *Q, int n, int m){
       //return min(solve(P, Q, n, m),solve(Q,P,m,n));
92
       if(two_getaway_ConvexHull(P,n,Q,m)==true) return
93
          \hookrightarrow min(solve(P, Q, n, m), solve(Q,P,m,n));
       else return 0.0;
94
95
  }
```

3.1.4 三角形的心

```
1 Point inCenter(const Point &A, const Point &B, const Point
      →&C) { // 内心
2
     double a = (B - C).len(), b = (C - A).len(), c = (A - C).len()
        \hookrightarrow B).len(),
3
       s = fabs(det(B - A, C - A)),
       r = s / p;
4
     return (A * a + B * b + C * c) / (a + b + c);
5
6
  }
   Point circumCenter(const Point &a, const Point &b, const
7
      → Point &c) { // 外心
8
     Point bb = b - a, cc = c - a;
     double db = bb.len2(), dc = cc.len2(), d = 2 * det(bb,
9
     return a - Point(bb.y * dc - cc.y * db, cc.x * db - bb.x
10
        \rightarrow * dc) / d:
11
   }
12
   Point othroCenter(const Point &a, const Point &b, const
      → Point &c) { // 垂心
     Point ba = b - a, ca = c - a, bc = b - c;
13
14
     double Y = ba.y * ca.y * bc.y,
15
          A = ca.x * ba.y - ba.x * ca.y,
16
          x0 = (Y + ca.x * ba.y * b.x - ba.x * ca.y * c.x) /
17
          y0 = -ba.x * (x0 - c.x) / ba.y + ca.y;
     return Point(x0, y0);
18
19 }
```

3.1.5 半平面交

```
struct Point {
                                                                                                                                 int quad() const { return sign(y) == 1 || (sign(y) == 0
                                                                                                                                        \hookrightarrow && sign(x) >= 0);}
                                                                                                                       3
                                                                                                                           };
                                                                                                                       4
                                                                                                                            struct Line {
                                                                                                                       5
                                                                                                                                bool include(const Point &p) const { return sign(det(b -
                                                                                                                                        \hookrightarrow a, p - a)) > 0; }
                                                                                                                                 Line push() const{ // 将半平面向外推 eps
                                                                                                                       6
                                                                                                                                      const double eps = 1e-6;
                                                                                                                       7
                                                                                                                      8
                                                                                                                                      Point delta = (b - a).turn90().norm() * eps;
                                                                                                                       9
                                                                                                                                      return Line(a - delta, b - delta);
                                                                                                                     10
                                                                                                                     11
                                                                                                                            };
                                                                                                                            bool sameDir(const Line &10, const Line &11) { return
                                                                                                                     12
                                                                                                                                  \hookrightarrow parallel(10, 11) && sign(dot(10.b - 10.a, 11.b -
                                                                                                                                  \hookrightarrow 11.a)) == 1; }
                                                                                                                            bool operator < (const Point &a, const Point &b) {</pre>
                                                                                                                                 if (a.quad() != b.quad()) {
                                                                                                                    15
                                                                                                                                     return a.quad() < b.quad();</pre>
                                                                                                                    16
                                                                                                                                } else {
                                                                                                                    17
                                                                                                                                     return sign(det(a, b)) > 0;
                                                                                                                     18
                                                                                                                     19
 \hspace{0.1in} \hookrightarrow \texttt{SegmentToSegment(P[yminP],P[yminP+1],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],P[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ymaxQ],Q[ym
                                                                                                                                 if (sameDir(10, 11)) {
                                                                                                                                     return 11.include(10.a);
                                                                                                                                } else {
                                                                                                                    23
                                                                                                                    24
                                                                                                                                     return (10.b - 10.a) < (11.b - 11.a);
                                                                                                                     25
                                                                                                                     26
                                                                                                                           }
                                                                                                                     27
                                                                                                                            bool check(const Line &u, const Line &v, const Line &w) {

    return w.include(intersect(u, v)); }

                                                                                                                             vector<Point> intersection(vector<Line> &1) {
                                                                                                                    28
                                                                                                                    29
                                                                                                                                 sort(l.begin(), l.end());
                                                                                                                                 deque<Line> q;
                                                                                                                     30
                                                                                                                                 for (int i = 0; i < (int)1.size(); ++i) {</pre>
                                                                                                                     31
                                                                                                                     32
                                                                                                                                     if (i && sameDir(l[i], l[i - 1])) {
                                                                                                                    33
                                                                                                                                          continue:
                                                                                                                     34
                                                                                                                                      while (q.size() > 1 && !check(q[q.size() - 2],
                                                                                                                     35
                                                                                                                                            \hookrightarrow q[q.size() - 1], l[i])) q.pop_back();
                                                                                                                                      while (q.size() > 1 && !check(q[1], q[0], l[i]))
                                                                                                                     36
                                                                                                                                            \hookrightarrow q.pop\_front();
                                                                                                                    37
                                                                                                                                      q.push_back(l[i]);
                                                                                                                     38
                                                                                                                     39
                                                                                                                                 while (q.size() > 2 && !check(q[q.size() - 2],
                                                                                                                                        \rightarrow q[q.size() - 1], q[0])) q.pop_back();
                                                                                                                     40
                                                                                                                                 while (q.size() > 2 \&\& !check(q[1], q[0], q[q.size() -
                                                                                                                                        \leftrightarrow 1])) q.pop_front();
                                                                                                                    41
                                                                                                                                 vector<Point> ret;
                                                                                                                                 for (int i = 0; i < (int)q.size(); ++i)</pre>
                                                                                                                    42
                                                                                                                                        \hookrightarrow ret.push_back(intersect(q[i], q[(i + 1) %
                                                                                                                                        \hookrightarrow q.size()]));
                                                                                                                     43
                                                                                                                                 return ret;
                                                                                                                     44
                                                                                                                           }
```

3.1.6 最大空凸包

```
inline double eq(double x, double y) {
2
       return fabs(x-y)<eps;</pre>
3
  }
   double xmult(point a, point b, point o) {
5
       return (a.x-o.x)*(o.y-b.y)-(a.y-o.y)*(o.x-b.x);
6
   }
7
   double dist(point a, point b) {
       return (a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y);
8
9
10
   point o;
11
   bool cmp_angle(point a,point b) {
       if(eq(xmult(a,b,o),0.0)) {
12
13
           return dist(a.o) < dist(b.o):
```

```
14
15
        return xmult(a,o,b)>0;
16
   }
17
   double empty_convex(point *p, int pn) {
18
        double ans=0:
        for(int i=0; i<pn; i++) {</pre>
19
20
            for(int j=0; j<pn; j++) {</pre>
                 dp[i][j]=0;
21
22
            }
23
       }
24
        for(int i=0; i<pn; i++) {</pre>
25
            int i = i-1:
26
            while(j>=0 && eq(xmult(p[i], p[j],
               \hookrightarrow o),0.0))j--;//coline
27
            bool flag= j==i-1;
28
            while(j \ge 0) {
                 int k = i-1:
29
30
                 while(k >= 0 && xmult(p[i],p[k],p[j])>0)k--;
31
                 double area = fabs(xmult(p[i],p[j],o))/2;
                 if(k \ge 0)area+=dp[j][k];
32
33
                 if(flag) dp[i][j]=area;
34
                 ans=max(ans, area);
35
                 j=k;
            }
36
37
            if(flag) {
38
                 for(int j=1; j<i; j++) {</pre>
39
                     dp[i][j] = max(dp[i][j],dp[i][j-1]);
40
            }
41
        }
42
43
        return ans;
44
   }
45
   double largest_empty_convex(point *p, int pn) {
46
        point data[maxn];
47
        double ans=0;
48
        for(int i=0; i<pn; i++) {</pre>
49
            o=p[i];
            int dn=0;
50
51
            for(int j=0; j<pn; j++) {
52
                 if(p[j].y>o.y||(p[j].y=o.y&&p[j].x>=o.x)) {
                     data[dn++]=p[j];
53
                 }
54
55
56
            sort(data, data+dn, cmp_angle);
57
            ans=max(ans, empty_convex(data, dn));
58
        }
59
        return ans;
60
   }
```

3.1.7 平面最近点对

```
1
   double Dis(Point a, Point b) {
2
     return sqrt((a.x-b.x)*(a.x-b.x) + (a.y-b.y)*(a.y-b.y));
3
  | }
   double Closest_Pair(int left, int right) {
4
5
     double d = INF;
6
     if(left == right) return d;
     if(left +1 == right)
8
       return Dis(p[left],p[right]);
     int mid = (left+right)>>1;
9
10
     double d1 = Closest Pair(left.mid):
     double d2 = Closest_Pair(mid,right);
11
12
     d = min(d1.d2):
     int k = 0;
13
14
     for(int i = left; i <= right; i++) {</pre>
       if(fabs(p[mid].x - p[i].x) \le d)
15
         temp[k++] = p[i];
17
     }
18
     sort(temp,temp+k,cmpy);
19
     for(int i = 0; i < k; i++) {
20
       for(int j = i+1; j < k && temp[j].y - temp[i].y < d;
          \hookrightarrow j++) {
21
         double d3 = Dis(temp[i],temp[j]);
```

3.1.8 最小覆盖圆

```
#include<cmath>
   #include<cstdio>
   #include<algorithm>
   using namespace std;
   const double eps=1e-6;
6
   struct couple
7
   {
8
     double x, y;
9
     couple(){}
10
     couple(const double &xx, const double &yy)
11
12
       x = xx; y = yy;
13
14
   } a[100001];
15
   //dis means distance, dis2 means square of it
17
   struct circle {
     double r; couple c;
18
   } cir;
19
20
   inline bool inside(const couple & x) {
     return di2(x, cir.c) < cir.r*cir.r+eps;</pre>
22
   }
23
   inline void p2c(int x, int y) {
24
     cir.c.x = (a[x].x+a[y].x)/2;
     cir.c.y = (a[x].y+a[y].y)/2;
25
26
     cir.r = dis(cir.c, a[x]);
27
28
   inline void p3c(int i, int j, int k) {
29
     couple x = a[i], y = a[j], z = a[k];
30
        \hookrightarrow sqrt(di2(x,y)*di2(y,z)*di2(z,x))/fabs(x*y+y*z+z*x)/2;
     couple t1((x-y).x, (y-z).x), t2((x-y).y, (y-z).y),
        \hookrightarrow t3((len(x)-len(y))/2, (len(y)-len(z))/2);
32
     cir.c = couple(t3*t2, t1*t3)/(t1*t2);
33
   }
34
   inline circle mi() {
35
     sort(a + 1, a + 1 + n);
36
     n = unique(a + 1, a + 1 + n) - a - 1;
37
     if(n == 1) {
38
       cir.c = a[1];
       cir.r = 0;
39
40
       return cir;
41
42
     random_shuffle(a + 1, a + 1 + n);
43
     p2c(1, 2);
44
     for(int i = 3; i <= n; i++)
45
       if(!inside(a[i])) {
46
         p2c(1, i);
         for(int j = 2; j < i; j++)
47
48
            if(!inside(a[j])) {
49
              p2c(i, j);
50
              for(int k = 1; k < j; k++)
51
                if(!inside(a[k]))
52
                  p3c(i,j, k);
53
54
       }
55
     return cir:
56
```

3.1.9 多边形内部可视

```
return OnLine(Q, C, B);
3
4
  }
5
   int Onleft(const Point & a, const Point &b, const Point &
     return dcmp(Cross(b - c, a - c)) > 0;
6
7 }
8 int visible(int x, int y) {
     int P = (x + n - 1) \% n, Q = (x + 1) \% n;
9
10
     Point u = p[y] - p[x], v = p[x] - p[P], w = p[x] - p[Q];
11
     if (Onleft(p[Q], p[x], p[P])) {
      return dcmp(Cross(v, u)) > 0 && dcmp(Cross(w, u)) < 0;
12
13
     } else {
14
       return !(dcmp(Cross(v, u)) < 0 && dcmp(Cross(w, u)) >
15
     }
16 }
17
   int solve(int x, int y) {
     if (vis[x][y] == dfn) return g[x][y];
18
19
     vis[x][y] = dfn;
     if (x == y || y == x + 1) return g[x][y] = 1;
20
21
     for (int i = x; i + 1 \le y; i++) {
22
      if (C(p[x], p[y], p[i], p[i + 1])) return g[x][y] = 0;
     }
23
     for (int i = x + 1; i < y; i++) {
24
25
       if (OnLine(p[x], p[i], p[y])) {
26
         return g[x][y] = solve(x, i) && solve(i, y);
27
     }
28
     if (!visible(x, y) || !visible(y, x)) return g[x][y] =
29
30
     return g[x][y] = 1;
31 }
```

3.1.10 V 图

```
const int AIX = 5;
2
   const int MAXM = AIX * MAXN:
3
 4 struct point {
 5
     double x, y;
 6
     int index;
 7
     struct Edge *in;
8
     point(double _x = 0, double _y = 0) : x(_x), y(_y) {}
9
   };
10
   inline bool operator< (const point &a, const point &b) {</pre>
11
     return a.x < b.x \mid \mid (sgn(a.x - b.x) == 0 && a.y < b.y);
12 }
13 inline double cross(const point &a, const point &b, const
      14
       (b - a, c - a); }
15
   struct Edge {
     point *Org, *Dest;
16
17
     Edge *Onext, *Oprev, *Dnext, *Dprev;
18 | };
   inline point* Other(const Edge *e, const point *p) {
19

    return e->Org == p ?
20
         e->Dest : e->Org; }
21
   inline Edge* Next(const Edge *e, const point *p) { return
      \hookrightarrow e->0rg == p ? e->0next
22
        : e->Dnext; }
   inline Edge* Prev(const Edge *e, const point *p) { return
      \hookrightarrow e->0rg == p ? e->0prev
        : e->Dprev; }
   struct gEdge {
25
26
     int u, v;
27
     double w:
     gEdge() {}
     \label{eq:gedge} \mbox{gEdge(int $\_$u, int $\_$v, double $\_$w)} \; : \; \mbox{u($\_$u)}, \; \mbox{v($\_$v)}, \; \mbox{w($\_$w)}
29
30 };
  inline bool operator< (const gEdge &a, const gEdge &b) {
31
      → return a.w < b.w: }</pre>
32 point p[MAXN], *Q[MAXN];
```

```
33 Edge mem[AIX * MAXN], *elist[AIX * MAXN];
    static int nfree;
    //Alloc memory
35
36
    inline void Alloc_Memory(const int &n) {
      nfree = AIX * n;
      Edge *e = mem:
38
      for (int i = 0; i < nfree; ++i)</pre>
40
        elist[i] = e++;
41
   }
42
    //Add an edge to a ring of edges
43
    inline void Splice(Edge *a, Edge *b, point *v) {
44
      Edge *next:
45
      if (a->0rg == v)
        next = a->Onext, a->Onext = b;
46
47
48
        next = a->Dnext, a->Dnext = b;
      if (next->0rg == v)
49
50
       next->Oprev = b;
51
        next->Dprev = b;
52
53
      if (b->0rg == v)
54
        b->Onext = next, b->Oprev = a;
55
56
        b->Dnext = next, b->Dprev = a;
57
58
    //Initialise a new edge
    inline Edge *MakeEdge(point *u, point *v) {
59
60
      Edge *e = elist[--nfree];
      e->Onext = e->Oprev = e->Dnext = e->Dprev = e;
61
      e->0rg = u, e->Dest = v;
62
63
      if (!u->in)
64
        u->in = e;
65
      if (!v->in)
66
        v->in = e;
67
      return e;
68
69
    //Creates a new edge and adds it to two rings of edges.
    inline Edge *Join(Edge *a, point *u, Edge *b, point *v,
       \hookrightarrow int side) {
71
      Edge *e = MakeEdge(u, v);
72
      if (side == 1) {
73
        if (a->0rg == u)
74
          Splice(a->Oprev, e, u);
75
76
          Splice(a->Dprev, e, u);
77
        Splice(b, e, v);
78
79
      else {
80
        Splice(a, e, u);
        if (b->0rg == v)
82
          Splice(b->Oprev, e, v);
83
        else
84
          Splice(b->Dprev, e, v);
      }
85
86
      return e;
87
    }
88
    //Remove an edge
89
    inline void Remove(Edge *e) {
90
      point *u = e \rightarrow Org, *v = e \rightarrow Dest;
91
      if (u->in == e)
92
        u->in = e->0next;
93
      if (v->in == e)
94
        v->in = e->Dnext;
95
      if (e->0next->0rg == u)
96
        e->Onext->Oprev = e->Oprev;
97
98
        e->Onext->Dprev = e->Oprev;
99
      if (e->Oprev->Org == u)
100
        e->Oprev->Onext = e->Onext;
        e->Oprev->Dnext = e->Onext;
103
      if (e->Dnext->Org == v)
        e->Dnext->Oprev = e->Dprev;
104
```

```
105
                                                                       175
                                                                                    vv = *D - *Other(prev, D);
        e->Dnext->Dprev = e->Dprev;
                                                                       176
106
                                                                                    P1 = det(uu, vv);
107
      if (e->Dprev->Org == v)
                                                                       177
                                                                                    if (!(P1 > EPS))
        e->Dprev->Onext = e->Dnext;
108
                                                                       178
                                                                                      break;
                                                                       179
                                                                                    cot_P = dot(uu, vv) / P1;
109
      else
        e->Dprev->Dnext = e->Dnext;
                                                                                    if (cot_P > cot_R)
110
                                                                       180
111
      elist[nfree++] = e;
                                                                                      break:
112 }
                                                                                    Remove(Er);
                                                                       182
113 //Determines the lower tangent of two triangulations
                                                                       183
                                                                                    Er = prev;
114 inline void Low_tangent(Edge *e_1, point *o_1, Edge *e_r,
                                                                       184
                                                                                    cot_R = cot_P;
                                                                                  }
       → point *o_r, Edge
                                                                       185
         **1_low, point **OL, Edge **r_low, point **OR) {
                                                                       186
                                                                                  while (true);
115
116
      point *d_1 = Other(e_1, o_1), *d_r = Other(e_r, o_r);
                                                                       187
117
                                                                                1 = Other(E1, 0); r = Other(Er, D);
      while (true) {
                                                                       188
118
        if (cross(*o_1, *o_r, *d_1) < -EPS) {</pre>
                                                                       189
                                                                                if (!BL || (BL && BR && cot_R < cot_L)) {</pre>
                                                                                  B = Join(B, 0, Er, r, 0);
119
          e_l = Prev(e_l, d_l);
                                                                       190
          o_1 = d_1;
                                                                                  D = r;
                                                                       191
120
121
          d_1 = Other(e_1, o_1);
                                                                       192
                                                                                }
122
                                                                       193
                                                                                else {
        else if (cross(*o_l, *o_r, *d_r) < -EPS) {</pre>
                                                                                  B = Join(E1, 1, B, D, 0);
123
                                                                       194
124
          e_r = Next(e_r, d_r);
                                                                       195
                                                                                  0 = 1:
                                                                                }
125
          o_r = d_r;
                                                                       196
          d_r = Other(e_r, o_r);
                                                                       197
                                                                              }
126
127
                                                                       198
                                                                              while (true);
128
        else
                                                                       199
129
          break;
                                                                       200
                                                                           inline void Divide(int s, int t, Edge **L, Edge **R) {
130
                                                                       201
                                                                              Edge *a, *b, *c, *ll, *lr, *rl, *rr, *tangent;
                                                                              int n = t - s + 1;
      *OL = o_1, *OR = o_r;
131
                                                                       202
                                                                              if (n == 2)
      *l_low = e_l, *r_low = e_r;
                                                                       203
132
133 }
                                                                                *L = *R = MakeEdge(Q[s], Q[t]);
                                                                       204
                                                                              else if (n == 3) {
134
   inline void Merge(Edge *lr, point *s, Edge *rl, point *u,
                                                                                a = MakeEdge(Q[s], Q[s + 1]);
       135
      double cot_L, cot_R, N1, cot_N, P1, cot_P;
                                                                       207
                                                                                b = MakeEdge(Q[s + 1], Q[t]);
136
      point 11, 12, r1, r2, uu, vv;
                                                                       208
                                                                                Splice(a, b, Q[s + 1]);
      point *0, *D, *OR, *OL;
                                                                                double v = cross(*Q[s], *Q[s + 1], *Q[t]);
137
                                                                       209
138
      Edge *B, *L, *R;
                                                                       210
                                                                                if (v > EPS) {
139
      Low_tangent(lr, s, rl, u, &L, &OL, &R, &OR);
                                                                       211
                                                                                  c = Join(a, Q[s], b, Q[t], 0);
      *tangent = B = Join(L, OL, R, OR, O);
                                                                       212
                                                                                  *L = a, *R = b;
140
141
      0 = OL, D = OR;
                                                                       213
                                                                                else if (v < -EPS) {</pre>
142
      do {
                                                                       214
                                                                                  c = Join(a, Q[s], b, Q[t], 1);
143
        Edge *El = Next(B, O), *Er = Prev(B, D), *next, *prev;
                                                                       215
        point *l = Other(El, O), *r = Other(Er, D);
                                                                                  *L = c, *R = c;
144
                                                                       216
        11 = *0 - *1, 12 = *D - *1, r1 = *0 - *r, r2 = *D -
145
                                                                       217
                                                                       218
146
        double cl = det(11, 12), cr = det(r1, r2);
                                                                       219
                                                                                  *L = a, *R = b;
        bool BL = cl > EPS, BR = cr > EPS;
147
                                                                       220
        if (!BL && !BR)
                                                                              else if (n > 3) {
148
                                                                       221
149
          break;
                                                                       222
                                                                                int split = (s + t) / 2;
150
        if (BL) {
                                                                       223
                                                                                Divide(s, split, &ll, &lr);
151
          double dl = dot(11, 12);
                                                                       224
                                                                                Divide(split + 1, t, &rl, &rr);
152
          cot_L = dl / cl;
                                                                                Merge(lr, Q[split], rl, Q[split + 1], &tangent);
153
          do {
                                                                       226
                                                                                if (tangent->Org == Q[s])
                                                                                 11 = tangent;
154
            next = Next(E1, 0);
                                                                       227
            uu = *0 - *0ther(next, 0);
                                                                                if (tangent->Dest == Q[t])
                                                                       228
155
            vv = *D - *Other(next, 0);
156
                                                                       229
                                                                                  rr = tangent;
            N1 = det(uu, vv);
157
                                                                       230
                                                                                *L = 11; *R = rr;
            if (!(N1 > EPS))
                                                                       231
                                                                              }
158
159
              break;
                                                                       232
                                                                           int task, n, m, k, root[MAXN];
160
            cot_N = dot(uu, vv) / N1;
                                                                       233
161
            if (cot_N > cot_L)
                                                                           gEdge E[MAXM], MST[MAXN];
                                                                       234
162
                                                                       235
                                                                           inline int Make_Graph() {
              break;
163
            Remove(E1);
                                                                       236
                                                                             Edge *start, *e;
            El = next;
                                                                       237
                                                                             int M = 0;
164
165
            cot_L = cot_N;
                                                                       238
                                                                              point *u, *v;
                                                                       239
                                                                              for(int i = 0; i < n; ++i) {
166
167
          while (true);
                                                                       240
                                                                                u = p + i:
168
                                                                       241
                                                                                start = e = u->in;
        if (BR) {
169
                                                                       242
                                                                                do {
170
          double dr = dot(r1, r2);
                                                                                 v = Other(e, u);
                                                                       243
171
          cot_R = dr / cr;
                                                                       244
                                                                                 if (u < v)
                                                                                    E[M++] = gEdge(u - p + 1, v - p + 1, dis(*u, *v));
172
          do ₹
                                                                       245
            prev = Prev(Er, D);
173
                                                                       246
                                                                                  e = Next(e, u):
174
            uu = *0 - *0ther(prev, D);
                                                                       247
```

```
248
        while(e != start);
249
250
      return M;
251 }
    int find_root(const int &x) { return root[x] ? root[x] =
252
       \hookrightarrow find_root(root[x]) : x;
253
    inline bool merge(const int &x, const int &y) {
254
255
      int p = find_root(x), q = find_root(y);
256
      if (p != q) {
        root[p] = q;
257
        return true;
258
259
260
      else
261
        return false;
262 }
    inline void kruskal(gEdge *E, int m, int n, gEdge* MST) {
263
264
     for (int i = 1; i <= n; ++i)
265
       root[i] = 0;
      sort(E, E + m);
266
267
      int tot = 0;
      for (int i = 0; i < m; ++i)
268
        if (merge(E[i].u, E[i].v))
269
270
          MST[tot++] = E[i];
271
272
    inline void MinimumEuclideanSpaningTree(point* p, int n,
       273
      Alloc_Memory(n);
      sort(p, p + n);
274
      for (int i = 0; i < n; ++i)
275
276
        Q[i] = p + i;
      Edge *L, *R;
277
278
      Divide(0, n - 1, &L, &R);
279
      m = Make_Graph();
      kruskal(E, m, n, MST);
280
281
   | }
282
    int main() {
      for (scanf("%d", &task); task--; ) {
283
284
        scanf("%d", &k);
        for (n = 0; scanf("%lf", &p[n].x) == 1 && p[n].x !=
285
           \hookrightarrow -1: ++n) {
          scanf("%lf", &p[n].y);
286
287
          p[n].in = NULL;
288
          p[n].index = n;
289
        }
290
        if (n == 1) {
          printf("0\n");
291
292
          continue;
293
294
        MinimumEuclideanSpaningTree(p, n, MST);
295
        printf("%d\n", int(ceil(k > n ? 0 : MST[n - k - 1].w))
           \hookrightarrow + EPS));
296
297 }
```

3.2 三维

3.2.1 三维点类

```
1 // 三维绕轴旋转,大拇指指向 axis 向量方向,四指弯曲
     →方向转 w 弧度
2 | Point rotate(const Point& s, const Point& axis, DB w) {
    DB x = axis.x, y = axis.y, z = axis.z;
    DB s1 = x * x + y * y + z * z, ss1 = msqrt(s1),
5
       cosw = cos(w), sinw = sin(w);
    DB a[4][4]:
6
7
    memset(a, 0, sizeof a);
8
    a[3][3] = 1:
    a[0][0] = ((y * y + z * z) * cosw + x * x) / s1;
    a[0][1] = x * y * (1 - cosw) / s1 + z * sinw / ss1;
11
    a[0][2] = x * z * (1 - cosw) / s1 - y * sinw / ss1;
    a[1][0] = x * y * (1 - cosw) / s1 - z * sinw / ss1;
12
    a[1][1] = ((x * x + z * z) * cosw + y * y) / s1;
13
```

```
14
     a[1][2] = y * z * (1 - cosw) / s1 + x * sinw / ss1;
     a[2][0] = x * z * (1 - cosw) / s1 + y * sinw / ss1;
16
     a[2][1] = y * z * (1 - cosw) / s1 - x * sinw / ss1;
     a[2][2] = ((x * x + y * y) * cos(w) + z * z) / s1;
17
     DB ans[4] = \{0, 0, 0, 0\}, c[4] = \{s.x, s.y, s.z, 1\};
18
     for (int i = 0; i < 4; ++ i)
19
20
       for (int j = 0; j < 4; ++ j)
21
         ans[i] += a[j][i] * c[j];
22
     return Point(ans[0], ans[1], ans[2]);
23
```

3.2.2 凸包

```
__inline P cross(const P& a, const P& b) {
 2
     return P(
 3
         a.y * b.z - a.z * b.y,
         a.z * b.x - a.x * b.z,
4
5
         a.x * b.y - a.y * b.x
6
           );
7
8
    __inline DB mix(const P& a, const P& b, const P& c) {
9
10
    return dot(cross(a, b), c);
11
   __inline DB volume(const P& a, const P& b, const P& c,
      \hookrightarrow \texttt{const} P& d) {
14
     return mix(b - a, c - a, d - a);
15
16
17
   struct Face {
18
     int a, b, c;
     __inline Face() {}
19
20
     __inline Face(int _a, int _b, int _c):
21
       a(_a), b(_b), c(_c) {}
22
     inline DB area() const {
23
       return 0.5 * cross(p[b] - p[a], p[c] - p[a]).len();
25
     __inline P normal() const {
26
      return cross(p[b] - p[a], p[c] - p[a]).unit();
27
     __inline DB dis(const P& p0) const {
28
29
       return dot(normal(), p0 - p[a]);
30
     }
31
32
   std::vector<Face> face, tmp; // Should be O(n).
33
   int mark[N][N], Time, n;
34
35
   __inline void add(int v) {
37
     ++ Time:
38
     clear(tmp);
39
     for (int i = 0; i < (int)face.size(); ++ i) {</pre>
40
       int a = face[i].a, b = face[i].b, c = face[i].c;
41
       if (sign(volume(p[v], p[a], p[b], p[c])) > 0) {
         mark[a][b] = mark[b][a] = mark[a][c] =
42
43
            mark[c][a] = mark[b][c] = mark[c][b] = Time;
44
45
       else {
         tmp.push_back(face[i]);
46
47
       }
48
     clear(face); face = tmp;
49
50
     for (int i = 0; i < (int)tmp.size(); ++ i) {</pre>
       int a = face[i].a, b = face[i].b, c = face[i].c;
51
       if (mark[a][b] == Time) face.emplace_back(v, b, a);
52
53
       if (mark[b][c] == Time) face.emplace_back(v, c, b);
       if (mark[c][a] == Time) face.emplace_back(v, a, c);
55
       assert(face.size() < 500u);
56
57
   }
58
```

```
void reorder() {
     for (int i = 2; i < n; ++ i) {
60
61
       P \text{ tmp} = cross(p[i] - p[0], p[i] - p[1]);
62
       if (sign(tmp.len())) {
63
         std::swap(p[i], p[2]);
         for (int j = 3; j < n; ++ j)
64
65
            if (sign(volume(p[0], p[1], p[2], p[j]))) {
66
              std::swap(p[j], p[3]);
67
              return;
68
            }
69
       }
70
     }
71
   }
72
73
   void build_convex() {
     reorder():
74
     clear(face):
75
76
     face.emplace_back(0, 1, 2);
77
     face.emplace_back(0, 2, 1);
     for (int i = 3; i < n; ++ i)
78
79
       add(i):
80 }
```

3.2.3 最小覆盖球

```
1 #include <iostream>
2
  #include<cstring>
3
  #include<algorithm>
   #include<cstdio>
5
   #include<cmath>
6
7
   using namespace std;
8
9
   const int eps = 1e-8;
10
11
   struct Tpoint
12
  | {
13
     double x, y, z;
14
  };
15
16
   int npoint, nouter;
17
  Tpoint pt[200000], outer[4],res;
18
19
   double radius, tmp;
20
   inline double dist(Tpoint p1, Tpoint p2) {
21
     double dx=p1.x-p2.x, dy=p1.y-p2.y, dz=p1.z-p2.z;
22
     return ( dx*dx + dy*dy + dz*dz );
23 }
   inline double dot(Tpoint p1, Tpoint p2) {
24
    return p1.x*p2.x + p1.y*p2.y + p1.z*p2.z;
25
26 }
27
   void ball() {
28
     Tpoint q[3]; double m[3][3], sol[3], L[3], det;
29
     int i, j;
30
     res.x = res.v = res.z = radius = 0:
31
     switch ( nouter ) {
32
       case 1: res=outer[0]; break;
33
       case 2:
           res.x=(outer[0].x+outer[1].x)/2;
34
           res.y=(outer[0].y+outer[1].y)/2;
35
           res.z=(outer[0].z+outer[1].z)/2;
36
37
           radius=dist(res, outer[0]);
38
           break:
39
       case 3:
           for (i=0; i<2; ++i ) {
40
              q[i].x=outer[i+1].x-outer[0].x;
41
              q[i].y=outer[i+1].y-outer[0].y;
42
43
              q[i].z=outer[i+1].z-outer[0].z;
44
45
            for (i=0; i<2; ++i) for(j=0; j<2; ++j)
46
             m[i][j]=dot(q[i], q[j])*2;
            for (i=0; i<2; ++i ) sol[i]=dot(q[i], q[i]);</pre>
47
            if (fabs(det=m[0][0]*m[1][1]-m[0][1]*m[1][0])<eps)</pre>
48
```

```
49
               return:
 50
             L[0]=(sol[0]*m[1][1]-sol[1]*m[0][1])/det;
 51
             L[1]=(sol[1]*m[0][0]-sol[0]*m[1][0])/det;
 52
             res.x=outer[0].x+q[0].x*L[0]+q[1].x*L[1];
53
             res.y=outer[0].y+q[0].y*L[0]+q[1].y*L[1];
             res.z=outer[0].z+q[0].z*L[0]+q[1].z*L[1];
54
             radius=dist(res, outer[0]);
 55
 56
             break;
 57
         case 4:
 58
             for (i=0; i<3; ++i) {
59
               q[i].x=outer[i+1].x-outer[0].x;
60
               q[i].y=outer[i+1].y-outer[0].y;
61
               q[i].z=outer[i+1].z-outer[0].z;
 62
               sol[i]=dot(q[i], q[i]);
63
64
             for (i=0;i<3;++i)
               for(j=0;j<3;++j) m[i][j]=dot(q[i],q[j])*2;</pre>
65
66
             det= m[0][0]*m[1][1]*m[2][2]
 67
               + m[0][1]*m[1][2]*m[2][0]
               + m[0][2]*m[2][1]*m[1][0]
 68
 69
               - m[0][2]*m[1][1]*m[2][0]
               - m[0][1]*m[1][0]*m[2][2]
 70
 71
               - m[0][0]*m[1][2]*m[2][1];
 72
             if ( fabs(det) < eps ) return;</pre>
 73
             for (j=0; j<3; ++j) {
 74
               for (i=0; i<3; ++i) m[i][j]=sol[i];</pre>
 75
               L[j]=(m[0][0]*m[1][1]*m[2][2]
 76
                   + m[0][1]*m[1][2]*m[2][0]
 77
                   + m[0][2]*m[2][1]*m[1][0]
78
                   - m[0][2]*m[1][1]*m[2][0]
 79
                   - m[0][1]*m[1][0]*m[2][2]
80
                   - m[0][0]*m[1][2]*m[2][1]
81
                  ) / det;
82
               for (i=0; i<3; ++i)</pre>
83
                 m[i][j]=dot(q[i], q[j])*2;
84
85
             res=outer[0];
             for (i=0; i<3; ++i ) {
 86
 87
               res.x += q[i].x * L[i];
88
               res.y += q[i].y * L[i];
89
               res.z += q[i].z * L[i];
90
 91
             radius=dist(res, outer[0]);
 92
93
    }
    void minball(int n) {
94
95
      ball():
96
      //printf("(%.31f,%.31f,%.31f) %.31f\n",

    res.x,res.y,res.z,radius);
97
      if ( nouter<4 )
98
         for (int i=0; i<n; ++i)
99
          if (dist(res, pt[i])-radius>eps) {
             outer[nouter]=pt[i];
100
101
             ++nouter:
102
             minball(i):
103
             --nouter;
104
             if (i>0) {
               Tpoint Tt = pt[i];
106
               memmove(&pt[1], &pt[0], sizeof(Tpoint)*i);
               pt[0]=Tt;
107
108
             }
          }
109
110
111
    void solve()
112
113
      for (int i=0:i<npoint:i++)</pre>
         \hookrightarrow scanf("%lf%lf",&pt[i].x,&pt[i].y,&pt[i].z);
114
      random_shuffle(pt, pt + npoint);
115
      radius=-1:
116
      for (int i=0;i<npoint;i++){</pre>
        if (dist(res,pt[i])-radius>eps){
117
118
          nouter=1:
119
          outer[0]=pt[i];
```

```
120     minball(i);
121     }
122     }
123     printf("%.5f\n",sqrt(radius));
124     }
125     int main(){
126     for(; cin >> npoint && npoint; )
127          solve();
128     return 0;
129     }
```

4. 字符串

4.1 AC 自动机

```
1 int newnode()
2 {
3
     ++tot:
     memset(ch[tot], 0, sizeof(ch[tot]));
5
     fail[tot] = 0;
     dep[tot] = 0;
7
     par[tot] = 0;
8
9
     return tot;
10 }
   void insert(char *s,int x)
11
12
     if(*s == '\0') return;
13
14
     else
15
       int &y = ch[x][*s - 'a'];
16
17
       if(y == 0)
18
19
         y = newnode();
20
         par[y] = x;
21
22
         dep[y] = dep[x] + 1;
23
25
       insert(s + 1, y);
26
     }
27
  | }
   void build()
29
30
     int line[maxn];
31
     int f = 0, r = 0;
32
     fail[root] = root;
33
34
35
     for(int i = 0; i < alpha; i++)</pre>
36
37
       if(ch[root][i])
38
39
         fail[ch[root][i]] = root;
         line[r++] = ch[root][i];
40
41
42
43
          ch[root][i] = root;
44
45
     }
46
47
     while(f != r)
48
49
       int x = line[f++];
50
51
       for(int i = 0; i < alpha; i++)
52
53
         if(ch[x][i])
55
            fail[ch[x][i]] = ch[fail[x]][i];
56
           line[r++] = ch[x][i];
57
58
```

```
59 else
60 {
61 ch[x][i] = ch[fail[x]][i];
62 }
63 }
64 }
```

4.2 后缀数组

```
const int MAXN = MAXL * 2 + 1;
   int a[MAXN], x[MAXN], y[MAXN], c[MAXN], sa[MAXN],
      \hookrightarrow \texttt{rank} \texttt{[MAXN]} , <code>height[MAXN]</code> ;
   void calc_sa(int n) {
     int m = alphabet, k = 1;
     memset(c, 0, sizeof(*c) * (m + 1));
     for (int i = 1; i <= n; ++i) c[x[i] = a[i]]++;</pre>
     for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
8
     for (int i = n; i; --i) sa[c[x[i]]--] = i;
9
     for (; k <= n; k <<= 1) {
10
       int tot = k;
11
       for (int i = n - k + 1; i \le n; ++i) y[i - n + k] = i;
12
       for (int i = 1; i \le n; ++i)
13
         if (sa[i] > k) y[++tot] = sa[i] - k;
       memset(c, 0, sizeof(*c) * (m + 1));
14
       for (int i = 1; i <= n; ++i) c[x[i]]++;</pre>
15
16
       for (int i = 1; i \le m; ++i) c[i] += c[i - 1];
       for (int i = n; i; --i) sa[c[x[y[i]]]--] = y[i];
       for (int i = 1; i \le n; ++i) y[i] = x[i];
18
19
       tot = 1; x[sa[1]] = 1;
20
       for (int i = 2; i \le n; ++i) {
         if (max(sa[i], sa[i - 1]) + k > n || y[sa[i]] !=
21
             \hookrightarrow y[sa[i - 1]] || y[sa[i] + k] != y[sa[i - 1] +
             \hookrightarrow k]) ++tot;
22
          x[sa[i]] = tot;
23
       if (tot == n) break; else m = tot;
24
25
26
   }
27
   void calc_height(int n) {
    for (int i = 1; i <= n; ++i) rank[sa[i]] = i;
     for (int i = 1; i <= n; ++i) {
30
       height[rank[i]] = max(0, height[rank[i - 1]] - 1);
       if (rank[i] == 1) continue;
31
32
       int j = sa[rank[i] - 1];
        while (max(i, j) + height[rank[i]] <= n && a[i +</pre>

    height[rank[i]]] == a[j + height[rank[i]]])

    ++height[rank[i]];

34
35
```

4.3 后缀自动机

```
static const int MAXL = MAXN * 2; // MAXN is original
  static const int alphabet = 26; // sometimes need
  int 1, last, cnt, trans[MAXL][alphabet], par[MAXL],
     \hookrightarrow is maxlength, size is the size of right
  char str[MAXL]:
  inline void init() {
    l = strlen(str + 1); cnt = last = 1;
    for (int i = 0; i <= 1 * 2; ++i) memset(trans[i], 0,

    sizeof(trans[i])):
8
    memset(par, 0, sizeof(*par) * (1 * 2 + 1));
    memset(mxl, 0, sizeof(*mxl) * (1 * 2 + 1));
10
    memset(size, 0, sizeof(*size) * (1 * 2 + 1));
11
12
  inline void extend(int pos, int c) {
    int p = last, np = last = ++cnt;
```

```
mxl[np] = mxl[p] + 1; size[np] = 1;
     for (; p && !trans[p][c]; p = par[p]) trans[p][c] = np;
15
16
     if (!p) par[np] = 1;
17
     else {
       int q = trans[p][c];
18
       if (mxl[p] + 1 == mxl[q]) par[np] = q;
19
20
         int nq = ++cnt;
21
22
         mxl[nq] = mxl[p] + 1;
23
         memcpy(trans[nq], trans[q], sizeof(trans[nq]));
24
         par[nq] = par[q];
25
         par[np] = par[q] = nq;
26
         for (; trans[p][c] == q; p = par[p]) trans[p][c] =
27
     }
28
29 }
30
   inline void buildsam() {
31
    for (int i = 1; i <= 1; ++i) extend(i, str[i] - 'a');</pre>
     memset(sum, 0, sizeof(*sum) * (1 * 2 + 1));
33
     for (int i = 1; i <= cnt; ++i) sum[mxl[i]]++;</pre>
     for (int i = 1; i <= 1; ++i) sum[i] += sum[i - 1];
34
     for (int i = cnt; i; --i) seq[sum[mxl[i]]--] = i;
35
     for (int i = cnt; i; --i) size[par[seq[i]]] +=
36
        \hookrightarrow size[seq[i]];
37
  }
```

4.5 manacher

```
void Manacher(std::string s,int p[])
 2
   {
 3
        string t = "$#";
 4
 5
        for (int i = 0; i < s.size(); i++)</pre>
 6
 7
            t += s[i]:
 8
            t += "#";
 9
10
11
        std::vector<int> p(t.size(), 0);
12
        int mx = 0, id = 0;
13
14
15
        for (int i = 1; i < t.size(); i++)</pre>
16
17
            p[i] = mx > i ? min(p[2 * id - i], mx - i) : 1;
            while (t[i + p[i]] == t[i - p[i]]) ++p[i];
18
19
20
            if (mx < i + p[i])
21
            {
22
                 mx = i + p[i];
23
                 id = i;
            }
24
25
        }
26
```

4.4 广义后缀自动机

```
1 inline void add_node(int x, int &last) {
     int lastnode = last;
2
3
     if (c[lastnode][x]) {
       int nownode = c[lastnode][x];
5
       if (l[nownode] == l[lastnode] + 1) last = nownode;
6
7
         int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
         for (int i = 0; i < alphabet; ++i) c[auxnode][i] =</pre>
8
            \hookrightarrow c[nownode][i];
9
         par[auxnode] = par[nownode]; par[nownode] = auxnode;
         for (; lastnode && c[lastnode][x] == nownode;
10
            \hookrightarrow lastnode = par[lastnode]) {
            c[lastnode][x] = auxnode;
11
         }
12
13
         last = auxnode;
       }
14
     } else {
15
       int newnode = ++cnt; l[newnode] = l[lastnode] + 1;
16
17
       for (; lastnode && !c[lastnode][x]; lastnode =
          18
       if (!lastnode) par[newnode] = 1;
19
         int nownode = c[lastnode][x];
20
21
         if (l[lastnode] + 1 == l[nownode]) par[newnode] =
            \hookrightarrow nownode;
22
         else {
            int auxnode = ++cnt; l[auxnode] = l[lastnode] + 1;
23
24
            for (int i = 0; i < alphabet; ++i) c[auxnode][i] =</pre>
              \hookrightarrow \texttt{c[nownode][i];}
25
            par[auxnode] = par[nownode]; par[nownode] =
              \hookrightarrow par[newnode] = auxnode;
            for (; lastnode && c[lastnode][x] == nownode;
26
              → lastnode = par[lastnode]) {
27
              c[lastnode][x] = auxnode;
28
29
         }
30
31
       last = newnode;
32
33 }
```

4.6 回文自动机

```
int nT, nStr, last, c[MAXT][26], fail[MAXT], r[MAXN],
      \hookrightarrow 1[MAXN], s[MAXN];
   int allocate(int len) {
 3
     l[nT] = len;
 4
     r[nT] = 0:
5
     fail[nT] = 0;
6
     memset(c[nT], 0, sizeof(c[nT]));
 7
     return nT++;
8
9
   void init() {
10
     nT = nStr = 0;
     int newE = allocate(0);
11
     int new0 = allocate(-1);
12
     last = newE;
13
     fail[newE] = newO;
     fail[new0] = newE;
15
16
     s[0] = -1;
17
18
   void add(int x) {
19
     s[++nStr] = x;
20
     int now = last;
     while (s[nStr - l[now] - 1] != s[nStr]) now = fail[now];
21
22
     if (!c[now][x]) {
       int newnode = allocate(l[now] + 2), &newfail =
23
          \hookrightarrow fail[newnode]:
       newfail = fail[now];
24
        while (s[nStr - l[newfail] - 1] != s[nStr]) newfail =
           \hookrightarrow fail[newfail];
26
       newfail = c[newfail][x];
27
       c[now][x] = newnode;
28
     }
29
     last = c[now][x];
30
     r[last]++;
31
32
   void count() {
     for (int i = nT - 1; i \ge 0; i--) {
33
34
       r[fail[i]] += r[i];
35
     }
36
   }
```

4.7 循环串的最小表示

```
// n 必须是 2 的次幂
   void fft(Complex a[], int n, int f) {
2
3
     for (int i = 0; i < n; ++i)
       if (R[i] < i) swap(a[i], a[R[i]]);</pre>
5
     for (int i = 1, h = 0; i < n; i <<= 1, h++) {
6
       Complex wn = Complex(cos(pi / i), f * sin(pi / i));
7
       Complex w = Complex(1, 0);
8
       for (int k = 0; k < i; ++k, w = w * wn) tmp[k] = w;
g
       for (int p = i \ll 1, j = 0; j \ll n; j += p) {
10
         for (int k = 0; k < i; ++k) {
           Complex x = a[j + k], y = a[j + k + i] * tmp[k];
11
           a[j + k] = x + y; a[j + k + i] = x - y;
12
13
       }
14
15
     }
  |}
16
```

5. 数据结构

5.1 可并堆

```
1 int merge(int x,int y)
2 {
  |//p[i] 结点 i 的权值,这里是维护大根堆
3
  //d[i] 在 i 的子树中, i 到右叶子结点的最远距离.
5
6
      if(!x) return y;
7
      if(!y) return x;
8
9
      if(p[x] < p[y]) std::swap(x, y);
10
11
      r[x] = merge(r[x], y);
      if(r[x]) fa[r[x]] = x;
12
13
      if(d[l[x]] < d[r[x]]) std::swap(l[x], r[x]);//调整树
14
        →的结构,使其满足左偏性质
15
      d[x] = d[r[x]] + 1;
16
17
      return x:
18 }
```

5.2 KD-Tree

```
1
   long long norm(const long long &x) {
       //
2
            For manhattan distance
3
       return std::abs(x):
            For euclid distance
       return x * x;
6
  }
7
8
   struct Point {
9
       int x, y, id;
10
       const int& operator [] (int index) const {
11
12
            if (index == 0) {
13
                return x;
           } else {
14
15
                return y;
16
           }
17
18
19
       friend long long dist(const Point &a, const Point &b)
          - √ {
20
           long long result = 0;
21
           for (int i = 0; i < 2; ++i) {
                result += norm(a[i] - b[i]);
22
23
24
           return result;
25
26 } point[N];
```

```
28
   struct Rectangle {
29
       int min[2], max[2];
30
31
       Rectangle() {
           min[0] = min[1] = INT_MAX; // sometimes int is
32
              \hookrightarrow not enough
33
            max[0] = max[1] = INT_MIN;
34
35
36
       void add(const Point &p) {
37
           for (int i = 0; i < 2; ++i) {
                min[i] = std::min(min[i], p[i]);
38
39
                max[i] = std::max(max[i], p[i]);
40
           }
       7
41
42
43
       long long dist(const Point &p) {
           long long result = 0;
           for (int i = 0; i < 2; ++i) {
45
                // For minimum distance
46
47
                result += norm(std::min(std::max(p[i],
                  \hookrightarrow \min[i]), \max[i]) - p[i]);
48
                // For maximum distance
49
                result += std::max(norm(max[i] - p[i]),
                   50
           }
51
            return result;
52
       }
   };
53
54
   struct Node {
55
56
       Point seperator;
57
       Rectangle rectangle;
       int child[2];
58
59
60
       void reset(const Point &p) {
            seperator = p;
62
           rectangle = Rectangle();
63
           rectangle.add(p);
           child[0] = child[1] = 0;
64
65
66
   } tree[N << 1];
68
   int size, pivot;
69
70
   bool compare(const Point &a, const Point &b) {
71
       if (a[pivot] != b[pivot]) {
72
           return a[pivot] < b[pivot];</pre>
73
74
       return a.id < b.id;</pre>
75
76
   // 左閉右開: build(1, n + 1)
77
   int build(int 1, int r, int type = 1) {
78
79
       pivot = type;
       if (1 >= r) {
81
           return 0;
82
       }
83
       int x = ++size;
84
       int mid = 1 + r \gg 1;
85
       std::nth_element(point + 1, point + mid, point + r,
          \hookrightarrow compare);
       tree[x].reset(point[mid]);
86
87
       for (int i = 1; i < r; ++i) {
            tree[x].rectangle.add(point[i]);
88
89
90
       tree[x].child[0] = build(1, mid, type ^ 1);
       tree[x].child[1] = build(mid + 1, r, type ^ 1);
92
       return x;
93
   }
94
```

```
int insert(int x, const Point &p, int type = 1) {
 96
         pivot = type;
 97
         if (x == 0) {
 98
             tree[++size].reset(p);
 99
             return size;
100
         tree[x].rectangle.add(p);
         if (compare(p, tree[x].seperator)) {
102
103
             tree[x].child[0] = insert(tree[x].child[0], p,
                \hookrightarrow type ^ 1);
104
         } else {
             tree[x].child[1] = insert(tree[x].child[1], p,
105
                \hookrightarrow type ^ 1);
106
107
         return x;
    ۱,
108
109
110 // For minimum distance
    // For maximum: 下面递归 query 时 0, 1 换顺序;< and
111
       \hookrightarrow >;min and max
    void query(int x, const Point &p, std::pair<long long,</pre>
       \hookrightarrow int> &answer, int type = 1) {
113
         pivot = type;
         if (x == 0 || tree[x].rectangle.dist(p) >
114

    answer.first) {
             return:
115
116
117
         answer = std::min(answer,
118
                  std::make_pair(dist(tree[x].seperator, p),

    tree[x].seperator.id));
119
         if (compare(p, tree[x].seperator)) {
             query(tree[x].child[0], p, answer, type ^ 1);
120
121
             query(tree[x].child[1], p, answer, type ^ 1);
122
         } else {
123
             query(tree[x].child[1], p, answer, type ^ 1);
124
             query(tree[x].child[0], p, answer, type ^ 1);
125
126
    }
127
    std::priority_queue<std::pair<long long, int> > answer;
128
129
    void query(int x, const Point &p, int k, int type = 1) {
130
         pivot = type;
131
         if (x == 0 || (int)answer.size() == k &&
132
            \hookrightarrow tree[x].rectangle.dist(p) > answer.top().first) {
133
134
         }
         answer.push(std::make_pair(dist(tree[x].seperator, p),
135

    tree[x].seperator.id));
         if ((int)answer.size() > k) {
136
137
             answer.pop();
138
139
         if (compare(p, tree[x].seperator)) {
             query(tree[x].child[0], p, k, type ^{1});
140
             query(tree[x].child[1], p, k, type ^ 1);
141
142
         } else {
143
             query(tree[x].child[1], p, k, type ^ 1);
             query(tree[x].child[0], p, k, type ^ 1);
144
145
146
    | }
```

5.3 Treap

```
1 struct Node{
2
     int mn, key, size, tag;
     bool rev:
3
     Node* ch[2]:
5
     Node(int mn, int key, int size): mn(mn), key(key),
        \hookrightarrow size(size), rev(0), tag(0){}
     void downtag();
6
7
     Node* update(){
       mn = min(ch[0] \rightarrow mn, min(key, ch[1] \rightarrow mn));
8
9
       size = ch[0] \rightarrow size + 1 + ch[1] \rightarrow size;
```

```
10
        return this;
11
12
   };
13
   typedef pair<Node*, Node*> Pair;
   Node *null. *root:
14
   void Node::downtag(){
15
     if(rev){
        for(int i = 0; i < 2; i++)
17
18
          if(ch[i] != null){
19
            ch[i] -> rev ^= 1;
            swap(ch[i] -> ch[0], ch[i] -> ch[1]);
20
21
22
       rev = 0;
23
     }
24
      if(tag){
        for(int i = 0; i < 2; i++)
25
          if(ch[i] != null){
26
27
            ch[i] -> key += tag;
28
            ch[i] -> mn += tag;
29
            ch[i] -> tag += tag;
          }
30
31
        tag = 0;
32
33
34
   int r(){
35
     static int s = 3023192386;
36
     return (s += (s << 3) + 1) & (^{\circ}0u >> 1);
37
   bool random(int x, int y){
38
39
     return r() \% (x + y) < x;
40
   Node* merge(Node *p, Node *q){
41
42
     if(p == null) return q;
     if(q == null) return p;
43
     p -> downtag();
44
45
     q -> downtag();
46
     if(random(p -> size, q -> size)){
        p -> ch[1] = merge(p -> ch[1], q);
47
        return p -> update();
48
49
     }else{
50
        q -> ch[0] = merge(p, q -> ch[0]);
51
        return q -> update();
52
53
   }
54
   Pair split(Node *x, int n){
55
     if(x == null) return make_pair(null, null);
     x -> downtag();
56
57
     if(n \le x \rightarrow ch[0] \rightarrow size){
58
        Pair ret = split(x -> ch[0], n);
59
        x \rightarrow ch[0] = ret.second;
60
        return make_pair(ret.first, x -> update());
61
     Pair ret = split(x \rightarrow ch[1], n - x \rightarrow ch[0] \rightarrow size -
62
         \hookrightarrow 1):
63
     x \rightarrow ch[1] = ret.first;
     return make_pair(x -> update(), ret.second);
64
65
66
   pair<Node*, Pair> get_segment(int 1, int r){
67
     Pair ret = split(root, 1 - 1);
68
     return make_pair(ret.first, split(ret.second, r - 1 +
         \hookrightarrow 1));
69
70
   int main(){
71
     null = new Node(INF, INF, 0);
72
     null \rightarrow ch[0] = null \rightarrow ch[1] = null;
73
     root = null;
74
```

5.4 Splay

```
template<class T>void checkmin(T &x,T y) {
```

```
3
    if(y < x) x = y;
                                                                        76
4 }
                                                                        77
                                                                               now->pushdown();
5
   struct Node
                                                                        78
                                                                               int t = (now->c[0] ? now->c[0]->size : 0) + 1;
6 | {
                                                                        79
7
     Node *c[2], *fa;
                                                                        80
8
     int size, rev;
                                                                               if(t == k) break;
                                                                        81
9
                                                                        82
10
     LL val, add, min;
                                                                        83
                                                                               if(t > k) now = now->c[0];
11
                                                                        84
                                                                               else now = now->c[1], k -= t;
     Node *init(LL v)
                                                                        85
12
13
                                                                        86
14
       val = min = v;
                                                                        87
                                                                             return now:
15
       add = rev = 0;
                                                                        88
16
       c[0] = c[1] = fa = NULL;
                                                                        89
                                                                           void rotate(Node *x,Node* &k)
17
       size = 1:
                                                                        90
                                                                             Node *y = x->fa, *z = y->fa;
18
                                                                        91
19
                                                                        92
       return this:
    }
                                                                             if(y != k) z -> c[z -> c[1] == y] = x;
20
                                                                        93
21
     void rvs()
                                                                             else k = x;
22
                                                                        95
       std::swap(c[0], c[1]);
23
                                                                        96
                                                                             x->fa = z:
       rev ^= 1;
                                                                        97
24
                                                                             int i = (y->c[1] == x);
25
     }
                                                                        98
26
     void inc(LL x)
                                                                        99
27
                                                                       100
                                                                             setc(y, i, x->c[i ^ 1]);
28
       val += x;
                                                                       101
                                                                             setc(x, i ^ 1, y);
29
       add += x;
                                                                       102
       min += x;
                                                                       103
                                                                             y->update(), x->update();
30
     }
                                                                           }
31
                                                                       104
     void pushdown()
                                                                           void spaly(Node *x,Node* &k)
32
                                                                       105
33
     {
                                                                       106
34
       if(rev)
                                                                       107
                                                                             static Node *st[maxs];
35
                                                                       108
                                                                             int top = 0;
         if(c[0]) c[0]->rvs();
36
                                                                       109
                                                                             Node *y, *z;
         if(c[1]) c[1]->rvs();
                                                                       110
37
38
         rev = 0;
                                                                       111
                                                                             y = x;
                                                                             while(y != k) st[++top] = y, y = y->fa;
39
                                                                       112
40
       if (add)
                                                                       113
                                                                             st[++top] = y;
41
                                                                       114
         if(c[0]) c[0]->inc(add);
                                                                             while(top) st[top]->pushdown(), top--;
42
                                                                       115
         if(c[1]) c[1]->inc(add);
43
                                                                       116
44
         add = 0;
                                                                       117
                                                                             while(x != k)
45
                                                                       118
     }
                                                                               y = x-fa, z = y-fa;
46
                                                                       119
47
     void update()
                                                                       120
                                                                               if(y != k)
48
                                                                       121
49
       min = val;
                                                                       122
50
       if(c[0]) checkmin(min, c[0]->min);
                                                                       123
                                                                                 if((y == z-c[1]) ^ (x == y-c[1])) rotate(x, k);
51
       if(c[1]) checkmin(min, c[1]->min);
                                                                       124
                                                                                  else rotate(y, k);
52
                                                                       125
53
       size = 1;
                                                                       126
54
       if(c[0]) size += c[0]->size;
                                                                       127
                                                                               rotate(x, k);
       if(c[1]) size += c[1]->size;
55
                                                                       128
                                                                       129
56
                                                                           Node *subtree(int 1,int r)
57
                                                                       130
                                                                       131
   } *root;
59
                                                                       132
                                                                             assert((++1) <= (++r));
   | Node* newnode(LL x)
                                                                             spaly(find(1 - 1), root);
60
                                                                       133
61 {
                                                                       134
                                                                             spaly(find(r + 1), root->c[1]);
62
     static Node pool[maxs], *p = pool;
                                                                       135
63
                                                                       136
                                                                             return root->c[1]->c[0];
64
     return (++p)->init(x);
                                                                       137
65 }
                                                                       138
                                                                           void ins(int pos,int v)
66
   void setc(Node *x,int t,Node *y)
                                                                       139
67 | ₹
                                                                       140
                                                                             pos++:
     x->c[t] = y;
                                                                             spaly(find(pos), root);
68
                                                                       141
69
     if(y) y->fa = x;
                                                                       142
                                                                             spaly(find(pos + 1), root->c[1]);
70 }
                                                                             setc(root->c[1], 0, newnode(v));
                                                                       143
71 Node *find(int k)
                                                                       144
                                                                             root->c[1]->update();
72 | {
                                                                       145
                                                                             root->update();
73
                                                                       146
     Node *now = root:
74
                                                                       147
                                                                           void del(int pos)
75
     while(true)
                                                                       148
                                                                           {
```

```
149
      pos++:
      spaly(find(pos - 1), root);
150
151
      spaly(find(pos + 1), root->c[1]);
      root->c[1]->c[0] = NULL;
152
      root->c[1]->update();
153
     root->update();
154
155 }
   void init()
156
157 {
158
     root = newnode(0);
159
      setc(root, 1, newnode(0));
160
      root->update();
161 }
```

5.5 Link cut Tree

```
1 inline void reverse(int x) {
2
     tr[x].rev ^= 1; swap(tr[x].c[0], tr[x].c[1]);
3 }
4
5
   inline void rotate(int x, int k) {
     int y = tr[x].fa, z = tr[y].fa;
6
       tr[x].fa = z; tr[z].c[tr[z].c[1] == y] = x;
7
8
       tr[tr[x].c[k ^ 1]].fa = y; tr[y].c[k] = tr[x].c[k ^
9
       tr[x].c[k ^ 1] = y; tr[y].fa = x;
10 }
11
12 | inline void splay(int x, int w) {
    int z = x; pushdown(x);
13
     while (tr[x].fa != w) {
15
      int y = tr[x].fa; z = tr[y].fa;
16
       if (z == w) {
         pushdown(z = y); pushdown(x);
17
18
         rotate(x, tr[y].c[1] == x);
19
         update(y); update(x);
20
       } else {
21
         pushdown(z); pushdown(y); pushdown(x);
22
         int t1 = tr[y].c[1] == x, t2 = tr[z].c[1] == y;
         if (t1 == t2) rotate(y, t2), rotate(x, t1);
23
24
         else rotate(x, t1), rotate(x, t2);
25
         update(z); update(y); update(x);
26
27
     }
28
     update(x);
     if (x != z) par[x] = par[z], par[z] = 0;
29
30
  }
31
32
   inline void access(int x) {
33
     for (int y = 0; x; y = x, x = par[x]) {
34
       splay(x, 0);
35
       if (tr[x].c[1]) par[tr[x].c[1]] = x, tr[tr[x].c[1]].fa
          \Rightarrow = 0:
36
       tr[x].c[1] = y; par[y] = 0; tr[y].fa = x; update(x);
37
38
  }
39
40
   inline void makeroot(int x) {
41
     access(x); splay(x, 0); reverse(x);
42
   }
43
44
   inline void link(int x, int y) {
45
     makeroot(x); par[x] = y;
46
  ۱,
47
48 inline void cut(int x, int y) {
49
    access(x); splay(y, 0);
50
     if (par[y] != x) swap(x, y), access(x), splay(y, 0);
51
     par[y] = 0;
52 }
53
   inline void split(int x, int y) { // x will be the root
     \hookrightarrow of the tree
```

```
makeroot(y); access(x); splay(x, 0);
56
```

5.6 树上莫队

```
void dfs(int u)
 1
2
 3
     dep[u] = dep[fa[u][0]] + 1;
 4
     for(int i = 1; i < logn; i++)</pre>
       fa[u][i] = fa[fa[u][i - 1]][i - 1];
5
6
 7
     stk.push(u);
8
     for(int i = 0; i < vec[u].size(); i++)</pre>
9
10
       int v = vec[u][i];
11
       if(v == fa[u][0]) continue;
12
13
14
       fa[v][0] = u, dfs(v);
15
16
        size[u] += size[v];
17
       if(size[u] >= bufsize)
18
19
20
          ++bcnt;
21
22
          while(stk.top() != u)
23
24
            block[stk.top()] = bcnt;
25
            stk.pop();
26
27
28
          size[u] = 0;
29
30
     }
31
32
     size[u]++;
33
34
   void prework()
35
36
     dfs(1);
37
38
     ++bcnt;
39
     while(!stk.empty())
40
41
       block[stk.top()] = bcnt;
42
        stk.pop();
43
44
45
   void rev(int u)
46
47
     now -= (cnt[val[u]] > 0);
48
49
     if(used[u])
50
     {
        cnt[val[u]]--;
51
52
        used[u] = false;
     7
53
     else
54
55
56
       cnt[val[u]]++;
       used[u] = true;
57
58
59
     now += (cnt[val[u]] > 0);
60
61
62
   void move(int &x,int y,int z)
63
64
     int fwd = y;
65
     rev(getlca(x, z));
66
     rev(getlca(y, z));
67
```

```
68
69
     while(x != y)
70
       if(dep[x] < dep[y]) std::swap(x, y);</pre>
71
72
       rev(x), x = fa[x][0];
73
74
75
76
     x = fwd:
77
  |}
78
   void solve()
79
  1
80
     std::sort(query + 1, query + m + 1);
81
     int L = 1, R = 1;
82
     rev(1);
83
84
     for(int i = 1; i <= m; i++)</pre>
85
86
       int 1 = query[i].u;
87
88
       int r = query[i].v;
89
       move(L, 1, R);
90
91
       move(R, r, L);
92
93
       ans[query[i].t] = now;
94
95 }
```

5.7 CDQ 分治

```
1
2
   {
3
     int x, y, z, idx;
4
5
     friend bool operator == (const Node &a,const Node &b)
6
7
       return a.x == b.x && a.y == b.y && a.z == b.z;
8
     }
9
     friend bool operator < (const Node &a,const Node &b)
10
11
       return a.y < b.y;</pre>
12
13
14
   } triple[maxn];
15
16
  | bool cmpx(const Node &a,const Node &b)
17 | {
     if(a.x != b.x) return a.x < b.x;</pre>
18
19
    if(a.y != b.y) return a.y < b.y;</pre>
     return a.z < b.z;</pre>
20
21
  | }
22
23
   void solve(int l,int r)
24 | {
25
     if(1 == r) return;
26
27
     int mid = (1 + r) >> 1;
28
     solve(1, mid);
29
30
31
     static std::pair<Node,int> Lt[maxn], Rt[maxn];
     int Ls = 0, Rs = 0;
32
33
34
     for(int i = 1; i <= mid; i++)
      Lt[++Ls] = std::make_pair(triple[i], i);
35
     for(int i = mid + 1; i <= r; i++)
36
37
      Rt[++Rs] = std::make_pair(triple[i], i);
38
39
     int pos = 1;
40
     std::sort(Lt + 1, Lt + Ls + 1);
41
     std::sort(Rt + 1, Rt + Rs + 1);
42
```

```
43
44
     backup.clear();
45
     for(int i = 1; i <= Rs; i++)</pre>
46
        while(pos <= Ls && !(Rt[i].first < Lt[pos].first))</pre>
47
48
          insert(Lt[pos].first.z, 1);
49
50
         pos++;
51
52
53
54
        f[Rt[i].second] += query(Rt[i].first.z);
55
56
     for(int i = 0; i < backup.size(); i++) pre[backup[i]] =</pre>
57
58
59
     solve(mid + 1, r);
60
```

5.8 整体二分

```
void solve(int l,int r,std::vector<int> q)
2
   {
3
     if(1 == r || q.empty())
 4
5
       for(int i = 0; i < q.size(); i++)</pre>
6
 7
          ans[q[i]] = 1;
8
       }
9
10
     else
11
       int mid = (1 + r) >> 1;
12
13
14
       backup.clear();
15
16
        for(int i = 1; i <= mid; i++)</pre>
17
18
          Event e = event[i];
19
          if(e.1 <= e.r)
20
21
22
            add(e.1, e.v);
23
            add(e.r + 1, -e.v);
24
          }
25
          else
26
          {
27
            add(1, e.v);
            add(e.r + 1, -e.v);
28
29
            add(e.1, e.v);
30
          }
        }
31
32
33
        std::vector<int> qL, qR;
34
35
        for(int i = 0; i < q.size(); i++)</pre>
36
37
          LL val = 0;
38
39
          for(int j = 0; j < vec[q[i]].size(); j++)</pre>
40
41
            val += count(vec[q[i]][j]);
42
            if(val >= p[q[i]]) break;
43
44
45
          if(cnt[q[i]] + val >= p[q[i]])
46
47
48
            qL.push_back(q[i]);
          }
49
50
          else
```

```
51
52
            cnt[q[i]] += val;
53
            qR.push_back(q[i]);
54
55
56
        for(int i = 0; i < backup.size(); i++) sum[backup[i]]</pre>
57
58
        solve(1, mid, qL);
59
        solve(mid + 1, r, qR);
60
61
   }
```

6. 图论

6.1 2-SAT tarjan

```
template<class TAT>void checkmin(TAT &x,TAT y)
1
2
   {
3
     if(y < x) x = y;
  }
4
5
   void tarjan(int u)
6 4
7
     dfn[u] = low[u] = ++dt;
     flag[u] = true;
8
9
     stk.push(u);
10
     for(int i = 0; i < vec[u].size(); i++)</pre>
11
12
13
       int v = vec[u][i];
14
15
       if(!dfn[v])
16
17
         tarjan(v);
18
         checkmin(low[u], low[v]);
19
20
       else if(flag[v])
21
22
         checkmin(low[u], dfn[v]);
23
       }
     }
24
25
     if(low[u] == dfn[u])
26
27
       ++bcnt:
28
       while(stk.top() != u)
29
30
         block[stk.top()] = bcnt;
31
         flag[stk.top()] = false;
32
33
         stk.pop();
34
35
       block[u] = bcnt;
36
       flag[u] = false;
37
38
       stk.pop();
39
40
  }
41
   bool solve()
42
43
       for(int i = 1; i <= 2 * n; i++)
44
         if(!dfn[i]) tarjan(i);
45
46
       bool ans = true;
47
       for(int i = 1; i <= n; i++)
48
49
         if(block[2 * i] == block[2 * i - 1])
50
51
            ans = false;
52
            break;
53
54
55
       return ans;
56 }
```

6.2 KM

```
struct KM {
     // Truly O(n^3)
2
     // 邻接矩阵,不能连的边设为 -INF, 求最小权匹配时
        →边权取负, 但不能连的还是 -INF, 使用时先对 1
        → -> n 调用 hungary() , 再 get_ans() 求值
 4
     int w[N][N]:
5
     int lx[N], ly[N], match[N], way[N], slack[N];
6
     bool used[N];
 7
     void init() {
8
       for (int i = 1; i <= n; i++) {
9
         match[i] = 0;
10
         lx[i] = 0;
11
         ly[i] = 0;
         way[i] = 0;
12
13
       }
14
15
     void hungary(int x) {
16
       match[0] = x;
       int j0 = 0;
17
       for (int j = 0; j \le n; j++) {
18
19
         slack[j] = INF;
20
         used[j] = false;
21
22
       do {
23
24
         used[j0] = true;
         int i0 = match[j0], delta = INF, j1 = 0;
25
         for (int j = 1; j \le n; j++) {
26
27
           if (used[j] == false) {
28
             int cur = -w[i0][j] - 1x[i0] - 1y[j];
29
             if (cur < slack[j]) {</pre>
30
               slack[j] = cur;
31
               way[j] = j0;
32
33
             if (slack[j] < delta) {</pre>
34
               delta = slack[j];
35
               j1 = j;
36
37
           }
38
         }
         for (int j = 0; j \le n; j++) {
39
40
           if (used[j]) {
41
             lx[match[j]] += delta;
42
             ly[j] -= delta;
43
44
           else slack[j] -= delta;
         }
45
46
         j0 = j1;
       } while (match[j0] != 0);
47
48
49
       do {
        int j1 = way[j0];
50
51
         match[j0] = match[j1];
52
         j0 = j1;
53
       } while (j0);
54
     }
55
56
     int get_ans() {
57
       int sum = 0;
58
       for(int i = 1; i <= n; i++) {
         if (w[match[i]][i] == -INF); // 无解
59
60
         if (match[i] > 0) sum += w[match[i]][i];
       }
61
62
       return sum;
    }
63
   } km;
```

6.3 点双连通分量

```
const bool BCC_VERTEX = 0, BCC_EDGE = 1;
2
   struct BCC { // N = NO + MO. Remember to call
      \hookrightarrow init(&raw_graph).
3
     Graph *g, forest; // g is raw graph ptr.
4
     int dfn[N], DFN, low[N];
5
     int stack[N], top;
                            // Where edge i is expanded to in
6
     int expand to[N]:

→ expaned graph.

7
     // \ensuremath{\text{Vertex i expaned to i.}}
     int compress_to[N]; \ \ // Where vertex i is compressed to.
8
9
     bool vertex_type[N], cut[N], compress_cut[N], branch[M];
10
     //std::vector<int> BCC_component[N]; // Cut vertex
        \hookrightarrow belongs to none.
11
     __inline void init(Graph *raw_graph) {
12
       g = raw_graph;
     }
13
     void DFS(int u, int pe) {
14
15
       dfn[u] = low[u] = ++DFN; cut[u] = false;
16
       if (!~g->adj[u]) {
17
         cut[u] = 1;
         compress_to[u] = forest.new_node();
18
19
         compress_cut[compress_to[u]] = 1;
20
21
       for (int e = g->adj[u]; ~e; e = g->nxt[e]) {
22
         int v = g \rightarrow v[e];
23
         if ((e ^ pe) > 1 && dfn[v] > 0 && dfn[v] < dfn[u]) {
24
            stack[top++] = e;
25
           low[u] = std::min(low[u], dfn[v]);
26
27
         else if (!dfn[v]) {
28
            stack[top++] = e; branch[e] = 1;
            DFS(v, e);
29
           low[u] = std::min(low[v], low[u]);
30
            if (low[v] >= dfn[u]) {
31
32
              if (!cut[u]) {
33
                cut[u] = 1;
34
                compress_to[u] = forest.new_node();
35
                compress_cut[compress_to[u]] = 1;
36
              }
37
              int cc = forest.new_node();
              forest.bi_ins(compress_to[u], cc);
38
              compress_cut[cc] = 0;
39
              //BCC_component[cc].clear();
40
41
              do {
42
                int cur_e = stack[--top];
43
                compress_to[expand_to[cur_e]] = cc;
                compress_to[expand_to[cur_e^1]] = cc;
44
45
                if (branch[cur_e]) {
                  int v = g->v[cur_e];
46
                  if (cut[v])
47
48
                    forest.bi_ins(cc, compress_to[v]);
49
                  else {
                    //BCC component[cc].push back(v):
50
51
                    compress_to[v] = cc;
52
53
54
              } while (stack[top] != e);
55
56
         }
57
       }
     }
58
     void solve() {
59
60
       forest.init(g->base);
61
       int n = g - n;
       for (int i = 0; i < g > e; i + +) {
62
63
         expand_to[i] = g->new_node();
65
       memset(branch, 0, sizeof(*branch) * g->e);
66
       memset(dfn + g->base, 0, sizeof(*dfn) * n); DFN = 0;
       for (int i = 0; i < n; i++)
67
         if (!dfn[i + g->base]) {
68
```

```
69         top = 0;
70         DFS(i + g->base, -1);
71         }
72     }
73     } bcc;
74
75     bcc.init(&raw_graph);
76     bcc.solve();
77     // Do something with bcc.forest ...
```

6.4 边双连通分量

```
struct BCC {
 2
     Graph *g, forest;
     int dfn[N], low[N], stack[N], tot[N], belong[N], vis[N],
 3

→ top, dfs_clock;

     // tot[] is the size of each BCC, belong[] is the BCC
 4
         \hookrightarrow that each node belongs to
 5
     pair<int, int > ori[M]; // bridge in raw_graph(raw node)
     bool is_bridge[M];
 7
      __inline void init(Graph *raw_graph) {
 8
        g = raw_graph;
9
        memset(is_bridge, false, sizeof(*is_bridge) * g -> e);
10
        memset(vis + g \rightarrow base, 0, sizeof(*vis) * g \rightarrow n);
11
12
      void tarjan(int u, int from) {
        dfn[u] = low[u] = ++dfs_clock; vis[u] = 1;
13
           \hookrightarrow stack[++top] = u;
        for (int p = g -> adj[u]; ~p; p = g -> nxt[p]) {
14
         if ((p ^ 1) == from) continue;
15
          int v = g -> v[p];
16
17
          if (vis[v]) {
18
            if (vis[v] == 1) low[u] = min(low[u], dfn[v]);
19
          } else {
20
            tarjan(v, p);
            low[u] = min(low[u], low[v]);
21
22
            if (low[v] > dfn[u]) is_bridge[p / 2] = true;
23
        }
24
        if (dfn[u] != low[u]) return;
25
        tot[forest.new_node()] = 0;
26
27
          belong[stack[top]] = forest.n;
28
29
          vis[stack[top]] = 2;
          tot[forest.n]++;
30
31
          --top;
       } while (stack[top + 1] != u);
32
33
34
     void solve() {
35
        forest.init(g -> base);
        int n = g \rightarrow n;
36
37
        for (int i = 0; i < n; ++i)
         if (!vis[i + g -> base]) {
38
            top = dfs_clock = 0;
39
40
            tarjan(i + g \rightarrow base, -1);
41
42
        for (int i = 0; i < g -> e / 2; ++i)
43
          if (is_bridge[i]) {
44
            int e = forest.e;
45
            forest.bi_ins(belong[g -> v[i * 2]], belong[g ->
               \hookrightarrow v[i * 2 + 1]], g \rightarrow w[i * 2]);
46
            ori[e] = make_pair(g -> v[i * 2 + 1], g -> v[i *
47
            ori[e + 1] = make_pair(g \rightarrow v[i * 2], g \rightarrow v[i * 2
               \hookrightarrow + 1]);
48
          }
49
     }
   } bcc;
```

6.5 最小树形图

```
const int MAXN,INF;// INF >= sum( W_ij )
   int from[MAXN + 10][MAXN * 2 + 10],n,m,edge[MAXN +
 2
       \hookrightarrow 10] [MAXN * 2 + 10];
   int sel[MAXN * 2 + 10], fa[MAXN * 2 + 10], vis[MAXN * 2 +
      → 101:
   int getfa(int x){if(x == fa[x]) return x; return fa[x] =
      \hookrightarrow getfa(fa[x]):
   void liuzhu(){ // 1-base: root is 1, answer = (sel[i], i)
 5
      \hookrightarrow for i in [2..n]
     fa[1] = 1;
 6
      for(int i = 2; i <= n; ++i){
 8
        sel[i] = 1; fa[i] = i;
        for(int j = 1; j <= n; ++j) if(fa[j] != i)</pre>
g
10
          if(from[j][i] = i, edge[sel[i]][i] > edge[j][i])
              \hookrightarrow sel[i] = j;
11
     }
12
      int limit = n;
13
      while(1){
        int prelimit = limit; memset(vis, 0, sizeof(vis));
14
           \hookrightarrow vis[1] = 1;
        for(int i = 2; i <= prelimit; ++i) if(fa[i] == i &&</pre>
15
           \hookrightarrow !vis[i]){
16
          int j = i; while(!vis[j]) vis[j] = i, j =
              \hookrightarrow getfa(sel[j]);
          if(j == 1 || vis[j] != i) continue; vector<int> C;
17
              \hookrightarrow int k = j;
18
          do C.push_back(k), k = getfa(sel[k]); while(k != j);
19
          ++limit;
20
          for(int i = 1; i <= n; ++i){
21
             edge[i][limit] = INF, from[i][limit] = limit;
22
          fa[limit] = vis[limit] = limit;
23
24
          for(int i = 0; i < int(C.size()); ++i){</pre>
             int x = C[i], fa[x] = limit;
25
             for(int j = 1; j \le n; ++j)
26
               if(edge[j][x] != INF && edge[j][limit] >
27
                  \hookrightarrow \texttt{edge[j][x] - edge[sel[x]][x])} \{
                  edge[j][limit] = edge[j][x] - edge[sel[x]][x];
28
29
                  from[j][limit] = x;
30
               }
          }
31
32
          for(int j=1;j<=n;++j) if(getfa(j)==limit)</pre>
              \hookrightarrow \text{edge[j][limit]} = \text{INF};
          sel[limit] = 1;
33
34
          for(int j = 1; j \le n; ++j)
35
             if(edge[sel[limit]][limit] > edge[j][limit])
                \hookrightarrow sel[limit] = j;
36
        if(prelimit == limit) break;
38
      for(int i = limit; i > 1; --i) sel[from[sel[i]][i]] =
39
         \hookrightarrow sel[i];
40 }
```

6.6 带花树

```
vector<int> link[maxn];
2 int n,match[maxn],Queue[maxn],head,tail;
3 int pred[maxn],base[maxn],start,finish,newbase;
4 | bool InQueue [maxn], InBlossom [maxn];
5 | void push(int u){ Queue[tail++]=u;InQueue[u]=true; }
  int pop(){ return Queue[head++]; }
  int FindCommonAncestor(int u,int v){
    bool InPath[maxn]:
8
Q
    for(int i=0;i<n;i++) InPath[i]=0;</pre>
10
    while(true){ u=base[u];InPath[u]=true;if(u==start)
       while(true){ v=base[v];if(InPath[v])
11
       12
    return v:
13 }
```

```
14 void ResetTrace(int u){
16
              while(base[u]!=newbase){
17
                   v=match[u]:
                    InBlossom[base[u]]=InBlossom[base[v]]=true;
18
19
                    u=pred[v]:
                    if(base[u]!=newbase) pred[u]=v;
20
21
22
        }
23
        void BlossomContract(int u,int v){
24
              newbase=FindCommonAncestor(u.v):
25
              for (int i=0;i<n;i++)</pre>
26
              InBlossom[i]=0:
27
              ResetTrace(u);ResetTrace(v);
28
              if(base[u]!=newbase) pred[u]=v;
29
              if(base[v]!=newbase) pred[v]=u;
              for(int i=0:i<n:++i)</pre>
30
31
              if(InBlossom[base[i]]){
                   base[i]=newbase:
33
                    if(!InQueue[i]) push(i);
34
              }
35
        bool FindAugmentingPath(int u){
36
37
              bool found=false;
38
              for(int i=0;i<n;++i) pred[i]=-1,base[i]=i;</pre>
39
              for (int i=0;i<n;i++) InQueue[i]=0;</pre>
40
              start=u;finish=-1; head=tail=0; push(start);
41
              while(head<tail){</pre>
42
                   int u=pop();
43
                   for(int i=link[u].size()-1;i>=0;i--){
                         int v=link[u][i];
45
                         if (base[u]!=base[v]&&match[u]!=v)
46
                               \label{eq:continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous_continuous
47
                                    BlossomContract(u,v);
48
                               else if(pred[v]==-1){
49
                                    pred[v]=u;
50
                                     if(match[v]>=0) push(match[v]);
                                     else{ finish=v; return true; }
51
52
53
54
55
             return found;
        void AugmentPath(){
58
              int u=finish,v,w;
59
              while(u>=0){

    v=pred[u]; w=match[v]; match[v]=u; match[u]=v; u=w; }

60
61
         void FindMaxMatching(){
              for(int i=0;i<n;++i) match[i]=-1;</pre>
63
              for(int i=0;i<n;++i) if(match[i]==-1)</pre>
                      \hookrightarrow \texttt{if}(\texttt{FindAugmentingPath(i)}) \  \, \texttt{AugmentPath()};
64
```

6.7 支配树

```
vector<int> prec[N], succ[N];
   vector<int> ord;
 3
   int stamp, vis[N];
   int num[N]:
   int fa[N];
   void dfs(int u) {
     vis[u] = stamp;
 7
8
     num[u] = ord.size();
9
     ord.push_back(u);
10
     for (int i = 0; i < (int)succ[u].size(); ++i) {</pre>
11
       int v = succ[u][i];
       if (vis[v] != stamp) {
12
13
         fa[v] = u;
14
         dfs(v);
       }
15
     }
16
```

```
17 }
18
   int fs[N], mins[N], dom[N], sem[N];
19
   int find(int u) {
    if (u != fs[u]) {
20
       int v = fs[u]:
21
       fs[u] = find(fs[u]);
22
       if (mins[v] != -1 && num[sem[mins[v]]] <</pre>
          \hookrightarrow num[sem[mins[u]]]) {
24
         mins[u] = mins[v];
25
       }
     }
26
27
     return fs[u]:
28
   void merge(int u, int v) { fs[u] = v; }
29
30
   vector<int> buf[N]:
31 | int buf2[N]:
  void mark(int source) {
32
33
    ord.clear();
34
    ++stamp;
     dfs(source);
35
     for (int i = 0; i < (int)ord.size(); ++i) {</pre>
36
      int u = ord[i];
37
       fs[u] = u, mins[u] = -1, buf2[u] = -1;
38
39
40
     for (int i = (int) ord.size() - 1; i > 0; --i) {
41
       int u = ord[i], p = fa[u];
42
       sem[u] = p;
       for (int j = 0; j < (int)prec[u].size(); ++j) {</pre>
43
         int v = prec[u][j];
44
         if (use[v] != stamp) continue;
45
46
         if (num[v] > num[u]) {
            find(v); v = sem[mins[v]];
47
48
         }
49
         if (num[v] < num[sem[u]]) {</pre>
            sem[u] = v:
50
51
52
       buf[sem[u]].push_back(u);
53
       mins[u] = u;
54
55
       merge(u, p);
56
       while (buf[p].size()) {
         int v = buf[p].back();
57
58
         buf[p].pop_back();
59
         find(v);
60
         if (sem[v] == sem[mins[v]]) {
61
           dom[v] = sem[v];
         } else {
62
63
            buf2[v] = mins[v];
64
65
       }
66
     }
     dom[ord[0]] = ord[0];
67
     for (int i = 0; i < (int)ord.size(); ++i) {</pre>
68
       int u = ord[i];
69
70
       if (~buf2[u]) {
          dom[u] = dom[buf2[u]];
71
72
73
     }
74 | }
```

6.8 无向图最小割

```
int cost[maxn] [maxn], seq[maxn], len[maxn], n, m, pop, ans;
   bool used[maxn];
 3
   void Init(){
 4
    int i,j,a,b,c;
5
     for(i=0;i<n;i++) for(j=0;j<n;j++) cost[i][j]=0;
6
     for(i=0;i<m;i++){</pre>
       scanf("%d %d %d",&a,&b,&c); cost[a][b]+=c;
 7
           \hookrightarrow cost[b][a]+=c;
     }
 8
     pop=n; for(i=0;i<n;i++) seq[i]=i;
9
10 }
```

```
11 void Work(){
12
     ans=inf; int i,j,k,l,mm,sum,pk;
13
     while(pop > 1){
       for(i=1;i<pop;i++) used[seq[i]]=0; used[seq[0]]=1;</pre>
14
15
       for(i=1;i<pop;i++) len[seq[i]]=cost[seq[0]][seq[i]];</pre>
       pk=0; mm=-inf; k=-1;
16
       for(i=1;i<pop;i++) if(len[seq[i]] > mm){
           18
       for(i=1;i<pop;i++){</pre>
19
         used[seq[l=k]]=1;
20
         if(i==pop-2) pk=k;
         if(i==pop-1) break;
21
         mm=-inf;
22
23
         for(j=1;j<pop;j++) if(!used[seq[j]])</pre>
24
            if((len[seq[j]]+=cost[seq[1]][seq[j]]) > mm)
25
              mm=len[seq[j]], k=j;
       }
26
27
       sum=0;
28
       for(i=0;i<pop;i++) if(i != k)</pre>
          \hookrightarrow sum+=cost[seq[k]][seq[i]];
29
       ans=min(ans,sum);
30
       for(i=0;i<pop;i++)</pre>
         cost[seq[k]][seq[i]]=cost[seq[i]][seq[k]]+=cost[seq[pk]][s
31
32
       seq[pk]=seq[--pop];
33
34
     printf("%d\n",ans);
35
```

6.9 最大团搜索

```
const int N = 1000 + 7;
2
   vector<vector<bool> > adj;
3
   class MaxClique {
4
       const vector<vector<bool> > adj;
       const int n:
       vector<int> result, cur_res;
       vector<vector<int> > color_set;
8
       const double t_limit; // MAGIC
9
     int para, level;
     vector<pair<int, int> > steps;
10
11
   public:
12
       class Vertex {
13
       public:
14
           int i, d;
15
           Vertex(int i, int d = 0) : i(i), d(d) {}
16
       }:
17
       void reorder(vector<Vertex> &p) {
18
           for (auto &u : p) {
19
               u.d = 0;
20
               for (auto v : p) u.d += adj[v.i][u.i];
21
           }
22
           sort(p.begin(), p.end(), [&](const Vertex &a,
              23
     // reuse p[i].d to denote the maximum possible clique
24
        \hookrightarrow for first i vertices.
25
       void init_color(vector<Vertex> &p) {
26
           int maxd = p[0].d;
27
           for (int i = 0; i < p.size(); i++) p[i].d = min(i,</pre>
              \hookrightarrow maxd) + 1:
28
29
       bool bridge(const vector<int> &s, int x) {
           for (auto v : s) if (adj[v][x]) return true;
30
31
           return false;
32
       }
33
     // approximate estimate the p[i].d
     // Do not care about first mink color class (For better
34
        \hookrightarrow result, we must get some vertex in some color class
        → larger than mink )
       void color_sort(vector<Vertex> &cur) {
35
```

```
int totc = 0, ptr = 0, mink =
36
               \hookrightarrow 0):
            for (int i = 0; i < cur.size(); i++) {</pre>
37
                int x = cur[i].i, k = 0;
38
                while (k < totc && bridge(color_set[k], x))</pre>
39
40
                if (k == totc) color_set[totc++].clear();
41
                color_set[k].push_back(x);
42
                if (k < mink) cur[ptr++].i = x;</pre>
           }
43
           if (ptr) cur[ptr - 1].d = 0;
44
45
            for (int i = mink; i < totc; i ++) {</pre>
                for (auto v : color_set[i]) {
46
47
                    cur[ptr++] = Vertex(v, i + 1);
                7
48
           }
49
50
       void expand(vector<Vertex> &cur) {
51
       steps[level].second = steps[level].second -
52

    steps[level].first + steps[level - 1].first;

       steps[level].first = steps[level - 1].second;
53
           while (cur.size()) {
54
                if (cur_res.size() + cur.back().d <=</pre>
55

    result.size()) return ;

56
                int x = cur.back().i;
57
                cur_res.push_back(x); cur.pop_back();
                vector<Vertex> remain;
58
                for (auto v : cur) {
59
60
                    if (adj[v.i][x]) remain.push_back(v.i);
                }
61
62
                if (remain.size() == 0) {
63
                    if (cur_res.size() > result.size()) result
                       } else {
64
65
            // Magic ballance.
            if (1. * steps[level].second / ++para < t_limit)</pre>
66

    reorder(remain);
67
                    color_sort(remain);
            steps[level++].second++;
68
                    expand(remain);
69
70
           level--;
71
                }
72
                cur_res.pop_back();
73
       }
74
75
   public:
76
       MaxClique(const vector<vector<bool> > &_adj, int n,
          \hookrightarrow double tt = 0.025) : adj(_adj), n(n), t_limit(tt)
77
           result.clear();
78
           cur_res.clear();
79
           color_set.resize(n);
       steps.resize(n + 1):
80
81
       fill(steps.begin(), steps.end(), make_pair(0, 0));
82
       level = 1;
83
       para = 0;
84
85
       vector<int> solve() {
86
           vector<Vertex> p;
87
           for (int i = 0; i < n; i++)
               \hookrightarrow p.push_back(Vertex(i));
88
            reorder(p);
89
            init_color(p);
90
            expand(p);
91
           return result:
92
93 };
```

6.10 斯坦纳树

```
void SPFA(int *dist)
{
```

```
static int line[maxn + 5];
 3
 4
        static bool hash[maxn + 5];
 5
        int f = 0, r = 0;
 6
 7
        for(int i = 1; i <= N; i++)
 8
            if(dist[i] < inf)</pre>
 9
            {
10
                 line[r] = i;
11
                 hash[i] = true;
12
                 r = (r + 1) \% (N + 1);
13
14
15
        while(f != r)
16
            int t = line[f];
17
            hash[t] = false;
18
            f = (f + 1) \% (N + 1);
19
20
            for(int i = head[t]; i ; i = edge[i].next)
22
                 int v = edge[i].v, dt = dist[t] + edge[i].w;
23
24
                 if(dt < dist[v])</pre>
25
26
27
                     dist[v] = dt;
28
29
                     if(!hash[v])
30
                          if(dist[v] < dist[line[f]])</pre>
31
32
                              f = (f + N) \% (N + 1);
33
34
                              line[f] = v;
35
                          }
36
                          else
37
38
                              line[r] = v;
39
                              r = (r + 1) \% (N + 1);
40
41
                          hash[v] = true;
42
                     }
43
44
                 }
45
            }
46
47
   }
48
   void solve()
49
   {
50
        for(int i = 1; i <= S; i++)
51
52
            for(int j = 1; j <= N; j++)</pre>
                 for(int k = (i - 1) & i; k; k = (k - 1) & i)
53
54
                     G[i][j] = std::min(G[i][j], G[k][j] + G[k]
                        \hookrightarrow ^ i][j]);
55
            SPFA(G[i]);
56
57
        }
58
```

6.11 虚树

```
bool cmp(const int lhs,const int rhs)
2
3
     return dfn[lhs] < dfn[rhs];</pre>
4
  }
5
   void build()
6
7
     std::sort(h + 1, h + 1 + m, cmp);
8
9
     int top = 0;
10
     for (int i = 1; i <= m; i++)
11
12
```

```
13
        if (!top) father[st[++top] = h[i]] = 0;
14
15
        {
            int p = h[i], lca = LCA(h[i],st[top]);
16
17
            while(d[st[top]] > d[lca])
18
19
20
                if (d[st[top - 1]] <= d[lca])</pre>
21
                     father[st[top]] = lca;
22
23
                top--;
            }
24
25
26
            if (st[top] != lca)
27
                t[++tot] = lca;
28
                father[lca] = st[top];
29
30
                st[++top] = lca;
31
32
            father[p] = lca;
33
            st[++top] = p;
34
35
       }
36
     }
37
   }
```

6.12 点分治

```
template<class TAT>void checkmax(TAT &x,TAT y)
2
  {
3
     if(x < y) x = y;
4
   template<class TAT>void checkmin(TAT &x,TAT y)
5
6
7
     if(y < x) x = y;
8
  ۱,
9
   void getsize(int u,int fa)
10 {
11
     size[u] = 1;
12
     smax[u] = 0;
13
     for(int i = 0; i < G[u].size(); i++)</pre>
14
15
       int v = G[u][i];
16
17
18
       if(v == fa || ban[v]) continue;
19
       getsize(v, u);
20
21
22
       size[u] += size[v];
        checkmax(smax[u], size[v]);
23
24
25
  |}
   int getroot(int u,int ts,int fa)
26
27
   {
28
     checkmax(smax[u], ts - size[u]);
29
30
     int res = u;
31
32
     for(int i = 0; i < G[u].size(); i++)</pre>
33
       int v = G[u][i];
34
35
36
       if(v == fa || ban[v]) continue;
37
       int w = getroot(v, ts, u);
38
39
40
       if(smax[w] < smax[res]) res = w;</pre>
41
42
43
     return res;
44 }
45 void solve()
```

```
46 {
47
     static int line[maxn];
48
     static std::vector<int> vec;
     int f = 0, r = 0;
49
50
51
     line[r++] = 1;
52
53
     while(f != r)
54
55
       int u = line[f++];
56
57
        getsize(u, 0);
58
        u = getroot(u, size[u], 0);
59
        ban[u] = true;
60
61
        vec.clear();
62
63
        for(int i = 0; i < G[u].size(); i++)</pre>
          if(!ban[G[u][i]]) vec.push_back(G[u][i]);
65
66
67
68
        do something you like...
69
70
71
72
        for(int i = 0; i < vec.size(); i++)</pre>
          line[r++] = vec[i];
73
74
75
   }
```

6.13 最小割最大流

```
bool BFS()
 1
2
   {
3
        for(int i = 1; i <= ind; i++) dep[i] = 0;</pre>
 4
 5
        dep[S] = 1, line.push(S);
 6
 7
        while(!line.empty())
8
9
            int now = line.front();
10
            line.pop();
11
12
            for(int i = head[now], p; i ; i = edge[i].next)
13
                if(edge[i].cap && !dep[p = edge[i].v])
                     dep[p] = dep[now] + 1, line.push(p);
14
       }
15
16
        if(dep[T])
17
18
19
            for(int i = 1; i <= ind; i++)
                cur[i] = head[i];
20
21
            return true;
22
       }
23
24
            return false;
25
26
   int DFS(int a,int flow)
27
        if(a == T) return flow;
28
29
30
       int ret = 0;
31
        for(int &i = cur[a], p; i ; i = edge[i].next)
32
33
            if(dep[p = edge[i].v] == dep[a] + 1 &&
               \hookrightarrow \texttt{edge[i].cap})
34
35
                 int ff = DFS(p, std::min(flow, edge[i].cap));
36
                flow -= ff, edge[i].cap -= ff;
37
                ret += ff, edge[i ^ 1].cap += ff;
38
```

```
39
40
                 if(!flow) break;
41
42
43
            return ret;
44 }
   int solve()
45
46
   {
47
        int totflow = 0;
48
        while(BFS())
49
50
        {
51
            totflow += DFS(S, INF);
52
53
54
        return totflow;
55 }
```

6.14 最小费用流

```
bool SPFA()
 2
   {
        static int line[maxv];
3
        static bool hash[maxv];
 4
        register int f = 0, r = 0;
5
 6
     for(int i = 1; i <= ind; i++)</pre>
 7
8
      {
9
          dist[i] = inf;
          from[i] = 0;
10
11
      }
12
13
        dist[S] = 0, line[r] = S, r = (r + 1) \% maxv;
14
        hash[S] = true;
15
        while(f != r)
16
17
        {
            int x = line[f];
18
19
20
            line[f] = 0, f = (f + 1) % maxv;
21
            hash[x] = false;
22
            for(int i = head[x]; i; i = edge[i].next)
23
                 if(edge[i].cap)
24
25
26
                      int v = edge[i].v;
27
                      int w = dist[x] + edge[i].cost;
28
                      if(w < dist[v])</pre>
29
30
                      {
                          dist[v] = w;
31
32
                          from[v] = i;
33
                          if(!hash[v])
34
35
                          {
                               if(f != r && dist[v] <=</pre>
36

    dist[line[f]])

                                    f = (f - 1 + maxv) \% maxv,
37
                                       \hookrightarrow line[f] = v;
                               else
38
                                    line[r] = v, r = (r + 1) %
39
                                       \hookrightarrow maxv;
40
41
                               hash[v] = true;
                          }
42
                     }
43
                 }
44
45
46
47
        return from[T];
48 }
49
50 int back(int x,int flow)
```

```
52
     if(from[x])
53
       flow = back(edge[from[x] ^ 1].v, std::min(flow,
54
           \hookrightarrow edge[from[x]].cap));
55
        edge[from[x]].cap -= flow;
56
57
        edge[from[x] ^ 1].cap += flow;
58
59
60
     return flow;
61
62
   int solve()
63
        int mincost = 0, maxflow = 0;
64
65
        while(SPFA())
66
67
            int flow = back(T, inf);
68
69
            mincost += dist[T] * flow;
70
71
            maxflow += flow;
72
        }
73
74
        return mincost;
75
```

6.15 zkw 费用流

```
int S, T, totFlow, totCost;
 1
3
   int dis[N], slack[N], visit[N];
   int modlable () {
5
6
       int delta = INF;
 7
        for (int i = 1; i <= T; i++) {
            if (!visit[i] && slack[i] < delta) delta =</pre>
               \hookrightarrow slack[i];
9
            slack[i] = INF;
       }
10
11
       if (delta == INF) return 1;
12
        for (int i = 1; i <= T; i++)
13
            if (visit[i]) dis[i] += delta;
14
        return 0;
15
16
   int dfs (int x, int flow) {
17
       if (x == T) {
18
19
            totFlow += flow;
20
            totCost += flow * (dis[S] - dis[T]);
21
            return flow;
22
       }
23
       visit[x] = 1;
24
       int left = flow;
25
        for (int i = e.last[x]; ~i; i = e.succ[i])
            if (e.cap[i] > 0 && !visit[e.other[i]]) {
26
27
                int y = e.other[i];
28
                if (dis[y] + e.cost[i] == dis[x]) {
                    int delta = dfs (y, min (left, e.cap[i]));
29
                    e.cap[i] -= delta;
30
                    e.cap[i ^ 1] += delta;
31
32
                    left -= delta;
33
                     if (!left) { visit[x] = 0; return flow; }
                } else {
34
                    slack[y] = min (slack[y], dis[y] +
35
                       \hookrightarrow e.cost[i] - dis[x]);
36
37
            }
38
        return flow - left;
39
40
41 pair <int, int> minCost () {
```

```
totFlow = 0; totCost = 0;
42
43
       fill (dis + 1, dis + T + 1, 0);
44
       do {
45
                fill (visit + 1, visit + T + 1, 0);
46
           } while (dfs (S, INF));
47
       } while (!modlable ());
48
49
       return make_pair (totFlow, totCost);
50
  }
```

6.16 最小割树

```
#include<iostream>
   #include<cstdio>
 3
   #include<cstdlib>
   #include<cstring>
5
   #include<algorithm>
   #include<queue>
6
   #define inf 0x3f3f3f3f
 7
8
   #define N 155
   using namespace std;
10
11
   int
      \hookrightarrow \texttt{cnt,n,m,dis[N],last[N],a[N],tmp[N],ans[N][N],s,t,mark[N];}
   struct edge{int to,c,next;}e[N*200];
12
   queue <int> q;
13
15
   void addedge(int u,int v,int c)
16
   {
17
           \hookrightarrow \texttt{e[++cnt].to=v;e[cnt].c=c;e[cnt].next=last[u];last[\rlap/\mu]=cntg}
18
           \hookrightarrow e[++cnt].to=u;e[cnt].c=c;e[cnt].next=last[v];last[v]=cnt1
19
   }
20
21
   bool bfs()
22
   {
23
        memset(dis,0,sizeof(dis));
24
        dis[s]=2;
25
        while (!q.empty()) q.pop();
26
        q.push(s);
        while (!q.empty())
28
        {
29
            int u=q.front();
30
            q.pop();
31
            for (int i=last[u];i;i=e[i].next)
32
                 if (e[i].c&&!dis[e[i].to])
33
                     dis[e[i].to]=dis[u]+1;
34
35
                     if (e[i].to==t) return 1;
36
                     q.push(e[i].to);
37
38
39
        return 0;
40
   }
41
   int dfs(int x,int maxf)
42
43
        if (x==t||!maxf) return maxf;
44
        int ret=0:
45
        for (int i=last[x];i;i=e[i].next)
46
47
            if (e[i].c&&dis[e[i].to]==dis[x]+1)
48
49
                 int f=dfs(e[i].to,min(e[i].c,maxf-ret));
50
                 e[i].c-=f;
                 e[i^1].c+=f;
51
                 ret+=f:
52
53
                 if (ret==maxf) break;
55
        if (!ret) dis[x]=0;
56
        return ret;
   ۱,
57
58
```

```
60
 61
         mark[x]=1;
 62
         for (int i=last[x];i;i=e[i].next)
 63
              if (e[i].c&&!mark[e[i].to]) dfs(e[i].to);
 64
 65
 66
    void solve(int l,int r)
 67
    {
 68
         if (l==r) return;
 69
         s=a[1];t=a[r];
 70
         for (int i=2;i<=cnt;i+=2)</pre>
 71
              e[i].c=e[i^1].c=(e[i].c+e[i^1].c)/2;
 72
         int flow=0;
 73
         while (bfs()) flow+=dfs(s.inf):
 74
         memset(mark, 0, sizeof(mark));
 75
         dfs(s):
 76
         for (int i=1;i<=n;i++)</pre>
 77
             if (mark[i])
                  for (int j=1;j<=n;j++)
 78
                      if (!mark[j])
 79
 80
                              \rightarrow ans[i][j]=ans[j][i]=min(ans[i][j],flow)
 81
         int i=1, j=r;
         for (int k=1;k<=r;k++)</pre>
 83
              if (mark[a[k]]) tmp[i++]=a[k];
 84
              else tmp[j--]=a[k];
 85
         for (int k=1;k<=r;k++)</pre>
             a[k]=tmp[k];
 86
 87
         solve(1,i-1);
         solve(j+1,r);
    int main()
 92
    {
 93
         int cas;
 94
         scanf("%d",&cas);
 95
         while (cas--)
 96
 97
             scanf("%d%d",&n,&m);
 98
             cnt=1:
 99
             for (int i=1;i<=n;i++)</pre>
                  a[i]=i:
              memset(last,0,sizeof(last));
102
             memset(ans,inf,sizeof(ans));
103
             for (int i=1;i<=m;i++)</pre>
104
              {
105
                  int x,y,z;
                  scanf("%d%d%d",&x,&y,&z);
107
                  addedge(x,y,z);
108
             }
109
              solve(1,n);
110
             int q;
             scanf("%d",&q);
111
112
              for (int i=1;i<=q;i++)</pre>
113
114
                  int x,tot=0;
115
                  scanf("%d",&x);
116
                  for (int i=1;i<n;i++)</pre>
117
                       for (int j=i+1;j<=n;j++)</pre>
118
                            if (ans[i][j]<=x) tot++;</pre>
119
                  printf("%d\n",tot);
120
             }
121
              cout << end1;
122
123
         return 0:
124
```

void dfs(int x)

6.17 上下界网络流建图

B(u,v) 表示边 (u,v) 流量的下界,C(u,v) 表示边 (u,v) 流量的上界,F(u,v) 表示边 (u,v) 的流量。设

$$G(u,v) = F(u,v) - B(u,v)$$
,显然有
$$0 \le G(u,v) \le C(u,v) - B(u,v)$$

6.17.1 无源汇的上下界可行流

建立超级源点 S^* 和超级汇点 T^* ,对于原图每条边 (u,v) 在新网络中连如下三条边: $S^* \to v$,容量为 B(u,v); $u \to T^*$,容量为 B(u,v); $u \to v$,容量为 C(u,v) - B(u,v)。 最后求新网络的最大流,判断从超级源点 S^* 出发的边是否都满流即可,边 (u,v) 的最终解中的实际流量为 G(u,v) + B(u,v)。

6.17.2 有源汇的上下界可行流

从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的边。按照**无源汇的上下界可行流**一样做即可,流量即为 $T\to S$ 边上的流量。

6.17.3 有源汇的上下界最大流

- 1. 在**有源汇的上下界可行流**中,从汇点 T 到源点 S 的 边改为连一条上界为 ∞ ,下届为 x 的边。x 满足二分性质,找到最大的 x 使得新网络存在**无源汇的上下界可行流**即为原图的最大流。
- 2. 从汇点 T 到源点 S 连一条上界为 ∞ ,下界为 0 的 边,变成无源汇的网络。按照**无源汇的上下界可行流** 的方法,建立超级源点 S^* 和超级汇点 T^* ,求一遍 $S^* \to T^*$ 的最大流,再将从汇点 T 到源点 S 的这条边拆掉,求一次 $S \to T$ 的最大流即可。

6.17.4 有源汇的上下界最小流

- 1. 在**有源汇的上下界可行流**中,从汇点 T 到源点 S 的 边改为连一条上界为 x,下界为 0 的边。x 满足二分性质,找到最小的 x 使得新网络存在**无源汇的上下界可行流**即为原图的最小流。
- 2. 按照无源汇的上下界可行流的方法,建立超级源点 S^* 与超级汇点 T^* ,求一遍 $S^* \to T^*$ 的最大流,但是注意这一次不加上汇点 T 到源点 S 的这条边,即不使之改为无源汇的网络去求解。求完后,再加上那条汇点 T 到源点 S 上界 ∞ 的边。因为这条边下界为 0,所以 S^* , T^* 无影响,再直接求一次 $S^* \to T^*$ 的最大流。若超级源点 S^* 出发的边全部满流,则 $T \to S$ 边上的流量即为原图的最小流,否则无解。

7. 其他

7.1 Dancing Links

7.1.1 精确覆盖

```
#pragma comment(linker, "/STACK:1024000000,1024000000")
   #define maxn 1000005
2
3
   using namespace std:
   int head.sz:
   int U[maxn],D[maxn],L[maxn],R[maxn];//上下左右链表指针
5
   int H[maxn],ROW[maxn],C[maxn],S[maxn],O[maxn];
7
   void remove(int c) {
8
       L[R[c]]=L[c];
9
       R[L[c]]=R[c];
10
       for(int i=D[c]; i!=c; i=D[i])
           for(int j=R[i]; j!=i; j=R[j]) {
               U[D[j]]=U[j];
12
13
               D[U[j]]=D[j];
               --S[C[j]];
14
15
16
  }
   void resume(int c) {
17
18
       for(int i=U[c]; i!=c; i=U[i]) {
19
           for(int j=L[i]; j!=i; j=L[j]) {
20
               ++S[C[i]]:
21
               U[D[j]]=j;
```

```
D[U[j]]=j;
23
24
       L[R[c]]=c;
25
       R[L[c]]=c;
26
27
   }
   void init(int m) {
28
       head=0;//头指针为 0
29
        for(int i=0; i<=m; i++) {</pre>
31
            D[i]=i;//建立双向十字链表
32
33
            L[i]=i-1;
34
            R[i]=i+1;
35
            S[i]=0:
36
37
       R[m]=0;
        L[0]=m;
38
39
       S[0] = INF+1:
40
        sz=m+1:
41
       memset(H.O.sizeof(H)):
42
43
   void insert(int i, int j) {
       if(H[i]) {
45
            L[sz] = L[H[i]];
            R[sz] = H[i];
46
            L[R[sz]] = sz;
47
48
            R[L[sz]] = sz;
49
50
        else {
51
            L[sz] = sz;
52
            R[sz] = sz:
            H[i] = sz;
53
54
55
       U[sz] = U[j];
       D[sz] = j;
56
57
       U[D[sz]] = sz;
58
       D[U[sz]] = sz;
       C[sz] = j;
59
60
        ROW[sz] = i;
61
        ++S[j];
62
   }
63
64
   bool dfs(int k,int len) {
65
        if(R[head] == head) {
            sort(0,0+len*len);
66
            int p=0;
            for(int i=0; i<len; i++) {</pre>
69
                for(int j=0; j<len; j++) {
70
                     int num=0[p++];
71
                     num=num-(i*len+j)*len;
72
                     printf("%d",num);
73
74
                puts("");
75
            }
76
            return true;
77
        }
78
        int s=INF.c:
79
        for (int t=R[head]; t!=head; t=R[t])
            if (S[t] <s) s=S[t],c=t;</pre>
81
        remove(c);
82
        for(int i=D[c]; i!=c; i=D[i]) {
83
            0[k]=ROW[i];
84
            for(int j=R[i]; j!=i; j=R[j])
85
                remove(C[j]);
            if(dfs(k+1,len))
87
                return true;
88
            for(int j=L[i]; j!=i; j=L[j])
89
                resume(C[j]);
90
91
        resume(c);
        return false;
93
```

7.1.2 重复覆盖

```
1
  int h()
 2
 3
        int i,j,k,count=0;
 4
        bool visit[N];
 5
        memset(visit,0,sizeof(visit));
        for(i=R[0];i;i=R[i])
 6
 7
 8
            if(visit[i]) continue;
 9
            count++:
10
            visit[i]=1:
            for(j=D[i];j!=i;j=D[j])
11
12
13
                 for(k=R[j];k!=j;k=R[k])
14
                     visit[C[k]]=1;
15
16
        }
17
        return count;
18
   | }
19
   void Dance(int k)
20
21
        int i,j,c,Min,ans;
22
        ans=h();
        if(k+ans>K || k+ans>=ak) return;
23
        if(!R[0])
24
25
26
            if(k<ak) ak=k;</pre>
27
            return;
28
       }
29
        for(Min=N,i=R[0];i;i=R[i])
            if(S[i]<Min) Min=S[i],c=i;</pre>
30
        for(i=D[c];i!=c;i=D[i])
31
32
33
            remove(i);
34
            for(j=R[i];j!=i;j=R[j])
35
                remove(j);
36
            Dance(k+1):
            for(j=L[i];j!=i;j=L[j])
37
38
                resume(j);
39
            resume(i);
40
        }
41
        return;
42
   ۱,
```

7.2 蔡勒公式

```
int zeller(int y,int m,int d) {
   if (m<=2) y--,m+=12; int c=y/100; y%=100;
   int w=((c>>2)-(c<<1)+y+(y>>2)+(13*(m+1)/5)+d-1)%7;
   if (w<0) w+=7; return(w);
}</pre>
```

7.3 五边形数定理

$$p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k-1} p(n - \frac{k(3k-1)}{2})$$

```
1 #include<iostream>
  #include<cstdio>
3 using namespace std;
4 #define LL __int64
  const int N=100005;
   const int MOD=1000000007;
   LL dp[N],fi[N];
8 LL five(LL x) { return (3*x*x-x)/2; }
   //五边形数
   void wbxs(){
10
11
       dp[0]=1;
       int t=1000; //其实可以等于 sqrt(N)
12
       for(int i=-t;i<=t;++i)</pre>
13
          fi[i+t]=five(i):
14
       for(int i=1;i<=100000;++i){
15
```

```
int flag=1;
17
            for(int j=1;;++j){
18
                LL a=fi[j+t],b=fi[-j+t];
                 if(a>i && b>i) break;
19
                 if(a<=i) dp[i]=(dp[i]+dp[i-a]*flag+MOD)%MOD;</pre>
20
                 if(b<=i) dp[i]=(dp[i]+dp[i-b]*flag+MOD)%MOD;</pre>
21
22
                 flag*=-1;
23
            }
24
        }
25
26
   int main(){
27
        wbxs():
28
        int T,n;
        scanf("%d",&T);
29
        while(T--){
30
            scanf("%d",&n);
31
            printf("%I64d\n",dp[n]);
32
        }
33
34
        return 0:
35
   }
```

7.4 凸包闵可夫斯基和

```
1  // cv[0..1] 为两个顺时针凸包,其中起点等于终点,求

→ 出的闵可夫斯基和不一定是严格凸包

2  int i[2] = {0, 0}, len[2] = {(int)cv[0].size() - 1,

→ (int)cv[1].size() - 1};

3  vector<P> mnk;

4  mnk.push_back(cv[0][0] + cv[1][0]);

do {

6  int d((cv[0][i[0] + 1] - cv[0][i[0]]) * (cv[1][i[1] + 1]

→ - cv[1][i[1]]) >= 0);

7  mnk.push_back(cv[d][i[d] + 1] - cv[d][i[d]] +

→ mnk.back());

8  i[d] = (i[d] + 1) % len[d];

9 } while(i[0] || i[1]);
```

8. 技巧

8.1 STL 归还空间

```
template <typename T>
__inline void clear(T& container) {
  container.clear(); // 或者删除了一堆元素
  T(container).swap(container);
}
```

8.2 大整数取模

8.3 读入优化

```
1 // getchar() 读入优化 << 关同步 cin << 此优化
2 // 用 isdigit() 会小幅变慢
3 // 返回 false 表示读到文件尾
4 namespace Reader {
5 const int L = (1 << 15) + 5;
6 char buffer[L], *S, *T;
7 __inline bool getchar(char &ch) {
8 if (S == T) {
9 T = (S = buffer) + fread(buffer, 1, L, stdin);
10 if (S == T) {
```

```
ch = EOF;
11
12
            return false;
13
            }
14
        ch = *S++:
15
        return true:
16
17
        __inline bool getint(int &x) {
18
19
        char ch; bool neg = 0;
        for (; getchar(ch) && (ch < '0' || ch > '9'); ) neg ^=
20
          \hookrightarrow ch == '-';
        if (ch == EOF) return false;
21
22
        x = ch - '0':
        for (; getchar(ch), ch >= '0' && ch <= '9'; )</pre>
23
24
         x = x * 10 + ch - '0';
25
        if (neg) x = -x;
26
        return true:
27
28
  1
```

8.4 二次随机法

```
#include <random>

int main() {
    std::mt19937 g(seed); // std::mt19937_64
    std::cout << g() << std::endl;
}</pre>
```

8.5 vimrc

```
set ruler
  set number
3 set smartindent
  set autoindent
  set tabstop=4
   set softtabstop=4
   set shiftwidth=4
   set hlsearch
   set incsearch
10 set autoread
11 set backspace=2
12
  set mouse=a
13
   syntax on
15
16
  nmap <C-A> ggVG
   vmap <C-C> "+y
17
18
   filetype plugin indent on
19
20
21
   autocmd FileType cpp set cindent
   autocmd FileType cpp map <F9> :!g++ % -o %< -g -std=c++11
     \hookrightarrow -Wall -Wextra -Wconversion && size %< <CR>
   autocmd FileType cpp map <C-F9> :!g++ % -o %< -std=c++11
     \hookrightarrow -02 && size %< <CR>
   autocmd FileType cpp map <F8> :!time ./%< < %<.in <CR>
   autocmd FileType cpp map <F5> :!time ./%< <CR>
  map <F3> :vnew %<.in <CR>
27
28 map <F4> :!gedit % <CR>
```

8.6 控制 cout 输出实数精度

```
std::cout << std::fixed << std::setprecision(5);</pre>
```

8.7 让 make 支持 c++11

```
export CXXFLAGS='-std=c++11 -Wall'
```

8.8 tuple 相关

8.9 汇编技巧

```
1 03优化
2 #define __ __attribute__ ((optimize("-03")))
3 #define _ _ __inline __attribute__ ((__gnu_inline__, __ __always_inline__, __artificial__))
4
5 汇编开栈
6 #pragma comment(linker, "/STACK:256000000")
7
8 int __size = 256 << 20;
9 char* __p__ = (char *) malloc(__size__) + __size__;
10
11 int main() {
    __asm__("mov1 %0, %%esp\n" :: "r"(__p__));
12    return 0;
14 }
```

9. 提示

9.1 线性规划转对偶

```
\begin{array}{l} \text{maximize } \mathbf{c}^T \mathbf{x} \\ \text{subject to } \mathbf{A} \mathbf{x} \leq \mathbf{b}, \mathbf{x} \geq 0 \\ \end{array} \\ \Longleftrightarrow \begin{array}{l} \text{minimize } \mathbf{y}^T \mathbf{b} \\ \text{subject to } \mathbf{y}^T \mathbf{A} \geq \mathbf{c}^T, \mathbf{y} \geq 0 \end{array}
```

9.2 NTT 素数及其原根

Prime	Primitive root	
1053818881	7	
1051721729	6	
1045430273	3	
1012924417	5	
1007681537	3	

9.3 积分表

9.3.1
$$ax^2 + bx + c(a > 0)$$

1.
$$\int \frac{\mathrm{d}x}{ax^2 + bx + c} = \begin{cases} \frac{2}{\sqrt{4ac - b^2}} \arctan \frac{2ax + b}{\sqrt{4ac - b^2}} + C & (b^2 < 4ac) \\ \frac{1}{\sqrt{b^2 - 4ac}} \ln \left| \frac{2ax + b - \sqrt{b^2 - 4ac}}{2ax + b + \sqrt{b^2 - 4ac}} \right| + C & (b^2 > 4ac) \end{cases}$$

2.
$$\int \frac{x}{ax^2 + bx + c} dx = \frac{1}{2a} \ln|ax^2 + bx + c| - \frac{b}{2a} \int \frac{dx}{ax^2 + bx + c}$$

9.3.2 $\sqrt{\pm ax^2 + bx + c}(a > 0)$

1.
$$\int \frac{dx}{\sqrt{ax^2 + bx + c}} = \frac{1}{\sqrt{a}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

2.
$$\int \sqrt{ax^2 + bx + c} dx = \frac{2ax + b}{4a} \sqrt{ax^2 + bx + c} + \frac{4ac - b^2}{8\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

3.
$$\int \frac{x}{\sqrt{ax^2 + bx + c}} dx = \frac{1}{a} \sqrt{ax^2 + bx + c} - \frac{b}{2\sqrt{a^3}} \ln|2ax + b + 2\sqrt{a}\sqrt{ax^2 + bx + c}| + C$$

4.
$$\int \frac{\mathrm{d}x}{\sqrt{c+bx-ax^2}} = -\frac{1}{\sqrt{a}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

5.
$$\int \sqrt{c + bx - ax^2} dx = \frac{2ax - b}{4a} \sqrt{c + bx - ax^2} + \frac{b^2 + 4ac}{8\sqrt{a^3}} \arcsin \frac{2ax - b}{\sqrt{b^2 + 4ac}} + C$$

6.
$$\int \frac{x}{\sqrt{c+bx-ax^2}} dx = -\frac{1}{a} \sqrt{c+bx-ax^2} + \frac{b}{2\sqrt{a^3}} \arcsin \frac{2ax-b}{\sqrt{b^2+4ac}} + C$$

9.3.3 $\sqrt{\pm \frac{x-a}{x-b}}$ 或 $\sqrt{(x-a)(x-b)}$

1. $\int \frac{\mathrm{d}x}{\sqrt{(x-a)(b-x)}} = 2\arcsin\sqrt{\frac{x-a}{b-x}} + C \ (a < b)$

$$\int \sqrt{(x-a)(b-x)} dx = \frac{2x-a-b}{4} \sqrt{(x-a)(b-x)} + \frac{(b-a)^2}{4} \arcsin \sqrt{\frac{x-a}{b-x}} + C, (a < b) \quad (1)$$

9.3.4 三角函数的积分

1.
$$\int \tan x dx = -\ln|\cos x| + C$$

2.
$$\int \cot x dx = \ln|\sin x| + C$$

3.
$$\int \sec x dx = \ln \left| \tan \left(\frac{\pi}{4} + \frac{x}{2} \right) \right| + C = \ln \left| \sec x + \tan x \right| + C$$

4.
$$\int \csc x dx = \ln \left| \tan \frac{x}{2} \right| + C = \ln \left| \csc x - \cot x \right| + C$$

5.
$$\int \sec^2 x dx = \tan x + C$$

6.
$$\int \csc^2 x dx = -\cot x + C$$

7.
$$\int \sec x \tan x dx = \sec x + C$$

8.
$$\int \csc x \cot x dx = -\csc x + C$$

9.
$$\int \sin^2 x dx = \frac{x}{2} - \frac{1}{4} \sin 2x + C$$

10.
$$\int \cos^2 x dx = \frac{x}{2} + \frac{1}{4} \sin 2x + C$$

11.
$$\int \sin^n x dx = -\frac{1}{n} \sin^{n-1} x \cos x + \frac{n-1}{n} \int \sin^{n-2} x dx$$

12.
$$\int \cos^n x dx = \frac{1}{n} \cos^{n-1} x \sin x + \frac{n-1}{n} \int \cos^{n-2} x dx$$

13.
$$\int \frac{\mathrm{d}x}{\sin^n x} = -\frac{1}{n-1} \frac{\cos x}{\sin^{n-1} x} + \frac{n-2}{n-1} \int \frac{\mathrm{d}x}{\sin^{n-2} x}$$

14.
$$\int \frac{dx}{\cos^n x} = \frac{1}{n-1} \frac{\sin x}{\cos^{n-1} x} + \frac{n-2}{n-1} \int \frac{dx}{\cos^{n-2} x}$$

$$\int \cos^{m} x \sin^{n} x dx$$

$$= \frac{1}{m+n} \cos^{m-1} x \sin^{n+1} x + \frac{m-1}{m+n} \int \cos^{m-2} x \sin^{n} x dx$$

$$= -\frac{1}{m+n} \cos^{m+1} x \sin^{n-1} x + \frac{n-1}{m+1} \int \cos^{m} x \sin^{n-2} x dx$$

16.
$$\int \frac{\mathrm{d}x}{a+b\sin x} = \begin{cases} \frac{2}{\sqrt{a^2-b^2}} \arctan \frac{a\tan \frac{x}{2}+b}{\sqrt{a^2-b^2}} + C & (a^2 > b^2) \\ \frac{1}{\sqrt{b^2-a^2}} \ln \left| \frac{a\tan \frac{x}{2}+b-\sqrt{b^2-a^2}}{a\tan \frac{x}{2}+b+\sqrt{b^2-a^2}} \right| + C & (a^2 < b^2) \end{cases}$$

17.
$$\int \frac{dx}{a+b\cos x} = \begin{cases} \frac{2}{a+b} \sqrt{\frac{a+b}{a-b}} \arctan\left(\sqrt{\frac{a-b}{a+b}} \tan\frac{x}{2}\right) + C & (a^2 > b^2) & 1. \int \ln x dx = x \ln x - x + C \\ \frac{1}{a+b} \sqrt{\frac{a+b}{a-b}} \ln\left|\frac{\tan\frac{x}{2} + \sqrt{\frac{a+b}{b-a}}}{\tan\frac{x}{2} - \sqrt{\frac{a+b}{b-a}}}\right| + C & (a^2 < b^2) & 2. \int \frac{dx}{x \ln x} = \ln\left|\ln x\right| + C \\ 3. \int x^n \ln x dx = \frac{1}{x} x^{n+1} \ln x dx = \frac{1}{x$$

18.
$$\int \frac{\mathrm{d}x}{a^2 \cos^2 x + b^2 \sin^2 x} = \frac{1}{ab} \arctan\left(\frac{b}{a} \tan x\right) + C$$

19.
$$\int \frac{dx}{a^2 \cos^2 x - b^2 \sin^2 x} = \frac{1}{2ab} \ln \left| \frac{b \tan x + a}{b \tan x - a} \right| + C$$

- 20. $\int x \sin ax dx = \frac{1}{a^2} \sin ax \frac{1}{a} x \cos ax + C$
- 21. $\int x^2 \sin ax dx = -\frac{1}{a}x^2 \cos ax + \frac{2}{a^2}x \sin ax + \frac{2}{a^3}\cos ax + C$
- 22. $\int x \cos ax dx = \frac{1}{a^2} \cos ax + \frac{1}{a} x \sin ax + C$
- 23. $\int x^2 \cos ax dx = \frac{1}{a}x^2 \sin ax + \frac{2}{a^2}x \cos ax \frac{2}{a^3}\sin ax + C$

9.3.5 反三角函数的积分 (其中 a > 0)

- 1. $\int \arcsin \frac{x}{a} dx = x \arcsin \frac{x}{a} + \sqrt{a^2 x^2} + C$
- 2. $\int x \arcsin \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arcsin \frac{x}{a} + \frac{x}{4} \sqrt{x^2 x^2} + C$
- 3. $\int x^2 \arcsin \frac{x}{a} dx = \frac{x^3}{2} \arcsin \frac{x}{a} + \frac{1}{9}(x^2 + 2a^2)\sqrt{a^2 x^2} + C$
- 4. $\int \arccos \frac{x}{a} dx = x \ \arccos \frac{x}{a} \sqrt{a^2 x^2} + C$
- 5. $\int x \arccos \frac{x}{a} dx = (\frac{x^2}{2} \frac{a^2}{4}) \arccos \frac{x}{a} \frac{x}{4} \sqrt{a^2 x^2} + C$
- 6. $\int x^2 \arccos \frac{x}{a} dx = \frac{x^3}{3} \arccos \frac{x}{a} \frac{1}{9}(x^2 + 2a^2)\sqrt{a^2 x^2} + C$
- 7. $\int \arctan \frac{x}{a} dx = x \arctan \frac{x}{a} \frac{a}{2} \ln(a^2 + x^2) + C$
- 8. $\int x \arctan \frac{x}{a} dx = \frac{1}{2}(a^2 + x^2) \arctan \frac{x}{a} \frac{a}{2}x + C$
- 9. $\int x^2 \arctan \frac{x}{a} dx = \frac{x^3}{3} \arctan \frac{x}{a} \frac{a}{6}x^2 + \frac{a^3}{6} \ln(a^2 + x^2) + C$

9.3.6 指数函数的积分

- 1. $\int a^x dx = \frac{1}{\ln a} a^x + C$
- 2. $\int e^{ax} dx = \frac{1}{2} a^{ax} + C$
- 3. $\int xe^{ax} dx = \frac{1}{a^2}(ax-1)a^{ax} + C$
- 4. $\int x^n e^{ax} dx = \frac{1}{a} x^n e^{ax} \frac{n}{a} \int x^{n-1} e^{ax} dx$
- 5. $\int xa^x dx = \frac{x}{\ln a}a^x \frac{1}{(\ln a)^2}a^x + C$
- 6. $\int x^n a^x dx = \frac{1}{\ln a} x^n a^x \frac{n}{\ln a} \int x^{n-1} a^x dx$
- 7. $\int e^{ax} \sin bx dx = \frac{1}{a^2 + b^2} e^{ax} (a \sin bx b \cos bx) + C$
- 8. $\int e^{ax} \cos bx dx = \frac{1}{a^2 + b^2} e^{ax} (b \sin bx + a \cos bx) + C$
- 9. $\int e^{ax} \sin^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \sin^{n-1} bx (a \sin bx)$ $nb\cos bx$) + $\frac{n(n-1)b^2}{a^2+b^2n^2}\int e^{ax}\sin^{n-2}bxdx$
- 10. $\int e^{ax} \cos^n bx dx = \frac{1}{a^2 + b^2 n^2} e^{ax} \cos^{n-1} bx (a \cos bx + nb \sin bx) + \frac{n(n-1)b^2}{a^2 + b^2 n^2} \int e^{ax} \cos^{n-2} bx dx$

9.3.7 对数函数的积分

- 3. $\int x^n \ln x dx = \frac{1}{n+1} x^{n+1} (\ln x \frac{1}{n+1}) + C$
- 4. $\int (\ln x)^n dx = x(\ln x)^n n \int (\ln x)^{n-1} dx$
- 5. $\int x^m (\ln x)^n dx = \frac{1}{m+1} x^{m+1} (\ln x)^n \frac{n}{m+1} \int x^m (\ln x)^{n-1} dx$