

(*) Building an 8-bit Breadboard Computer

by Ben Eater

(<https://youtube.com/playlist?list=PLowKtXNTBypGqImE405J2565dvjafg1HU>)

^ SR Latch

(<https://youtu.be/KM0DdEaY5sY>)

Let the initial STABLE state be $S = 0$, $R = 0$, $Q = 0$ and $Q' = 1$.



Let the PROPAGATION delays of the upper and lower NOR gates be 2 units and 3 units, respectively.

So, the EFFECTS of the changed inputs will be visible only AFTER the propagation delays.

In order to draw the TIMING diagrams, for every logic gate, look at its inputs at time $(t - T)$, where T is its propagation delay.

For eg., if the propagation delay of a logic gate is 3 units, then its OUTPUT at $t = 3.5$ units (say) depends upon its INPUTS at $t = 0.5$ units.

Now, let S get changed to 1 at $t = 0$.

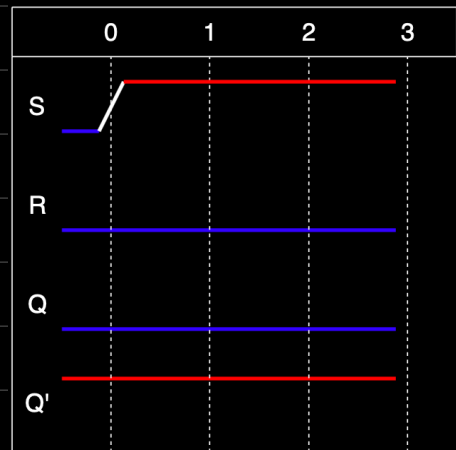
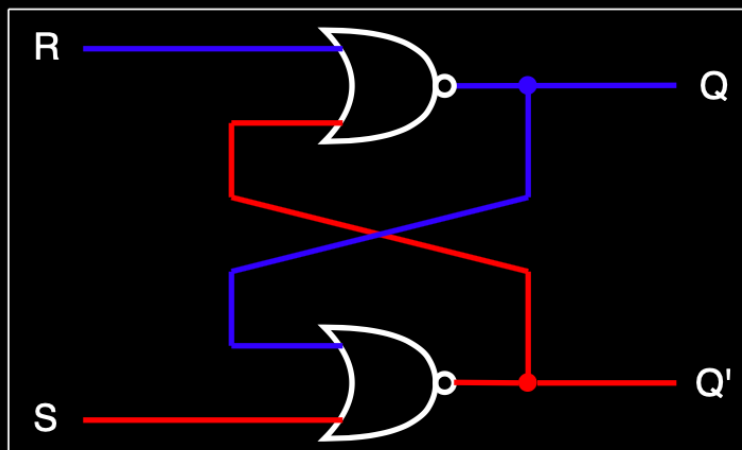
From $t = 0$ to $t = 1$:-



From $t = 1$ to $t = 2$:-



From $t = 2$ to $t = 3$:-



From $t = 3$ to $t = 4$:-



From $t = 4$ to $t = 5$:-



From $t = 5$ to $t = 6$:-



Now, the latch will stay STABLE in this SET state, even if S gets changed back to 0.

Let S get changed back to 0 at $t = 6$.

From $t = 6$ to $t = 7$:-



From $t = 7$ to $t = 8$:-



Hence, the latch will stay STABLE in this SET state, as long as S was maintained at 1 for a LONG ENOUGH time, i.e. if S gets changed back to 0 AFTER $t = 5$, then the latch will result in a STABLE SET state.

Now, let S get changed back to 0 at $t = 4$, instead of at $t = 6$.

From $t = 4$ to $t = 5$:-



From $t = 5$ to $t = 6$:-



From $t = 6$ to $t = 7$:-



From $t = 7$ to $t = 8$:-



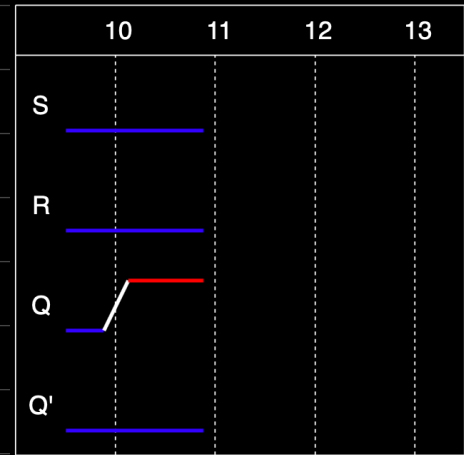
From $t = 8$ to $t = 9$:-



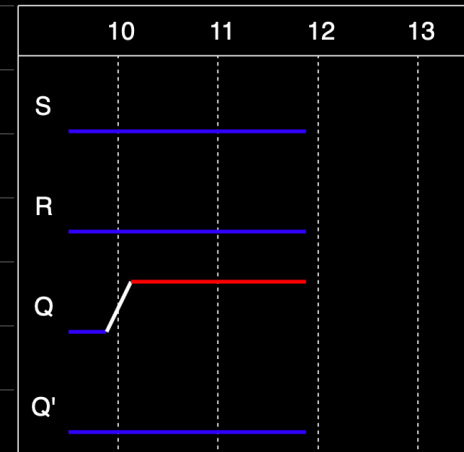
From $t = 9$ to $t = 10$:-



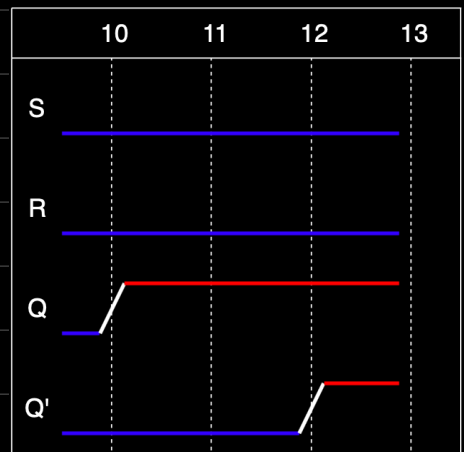
From $t = 10$ to $t = 11$:-



From $t = 11$ to $t = 12$:-



From $t = 12$ to $t = 13$:-



The latch will keep oscillating like this FOREVER.

Previously, when the latch was being SET, it temporarily switched to an INVALID state, i.e. $Q = Q' = 0$, but then it became STABLE in the SET state.

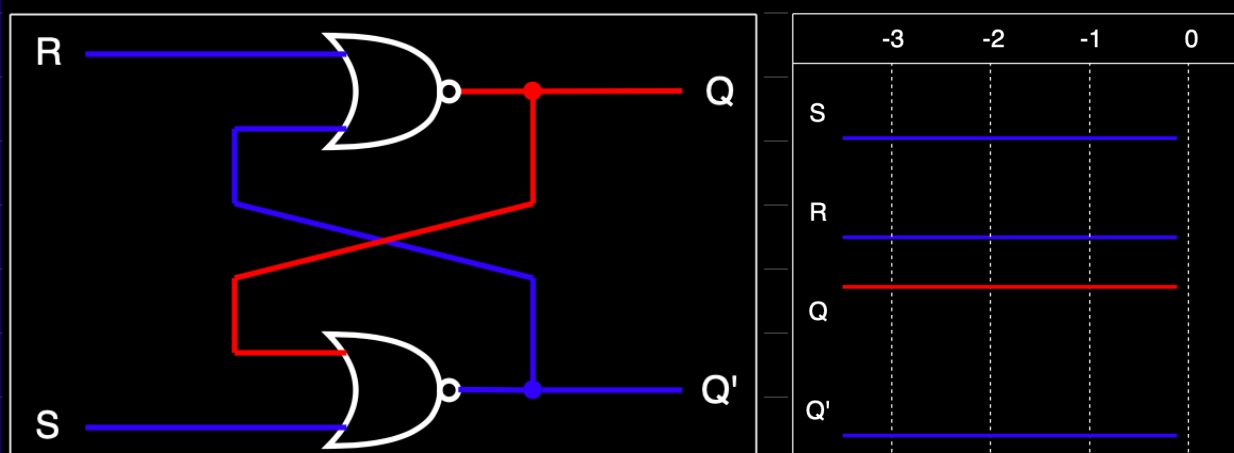
But, in this case, the latch will NEVER become STABLE in ANY state.

Therefore, if S is NOT maintained at 1 for a LONG ENOUGH time, then problems like this may occur.

It should be noted that the above analysis corresponds to an IDEAL version of logic gates. In real-world circuits, for EDGE cases like this, logic gates may NOT work as expected. For eg., even though ideally the latch should keep oscillating FOREVER, in real-world circuits, the latch MAY or MAY NOT keep oscillating due to the time it takes to turn on and off transistors, to charge internal capacitors, etc.

In any case, if used CORRECTLY, i.e. by keeping S/R high for a LONG ENOUGH time when setting/resetting the latch and by NOT keeping S & R high at the same time, then the latch will work AS EXPECTED in ideal as well as real-world circuits.

Now, Let the initial STABLE state be $S = 0$, $R = 0$, $Q = 1$ and $Q' = 0$.



Let R get changed to 1 at $t = 0$.

From $t = 0$ to $t = 1$:-



From $t = 1$ to $t = 2$:-



From $t = 2$ to $t = 3$:-



From $t = 3$ to $t = 4$:-



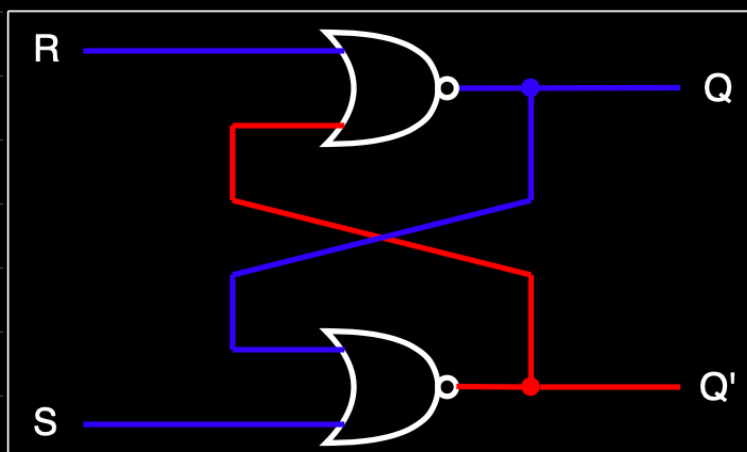
	3	4	5	6
S	low	low	low	low
R	high	high	low	low
Q	low	low	low	low
Q'	high	high	high	high

From $t = 4$ to $t = 5$:-



	3	4	5	6
S	low	low	low	low
R	high	high	low	low
Q	low	low	low	low
Q'	high	high	high	high

From $t = 5$ to $t = 6$:-



	3	4	5	6
S	low	low	low	low
R	high	high	low	low
Q	low	low	low	low
Q'	high	high	high	high

Now, the latch will stay STABLE in this RESET state, even if R gets changed back to 0, as long as R was maintained at 1 for a LONG ENOUGH time, i.e. if R gets changed back to 0 AFTER $t = 5$, then the latch will result in a STABLE RESET state.

Therefore, if R is NOT maintained at 1 for a LONG ENOUGH time, then similar problems may occur like when S is NOT maintained at 1 for a LONG ENOUGH time when SETTING the latch.

Now, Let the initial STABLE state be $S = 1$, $R = 1$, $Q = 0$ and $Q' = 0$ (INVALID).



Let S & R both get changed to 0 at $t = 0$.

From $t = 0$ to $t = 1$:-



