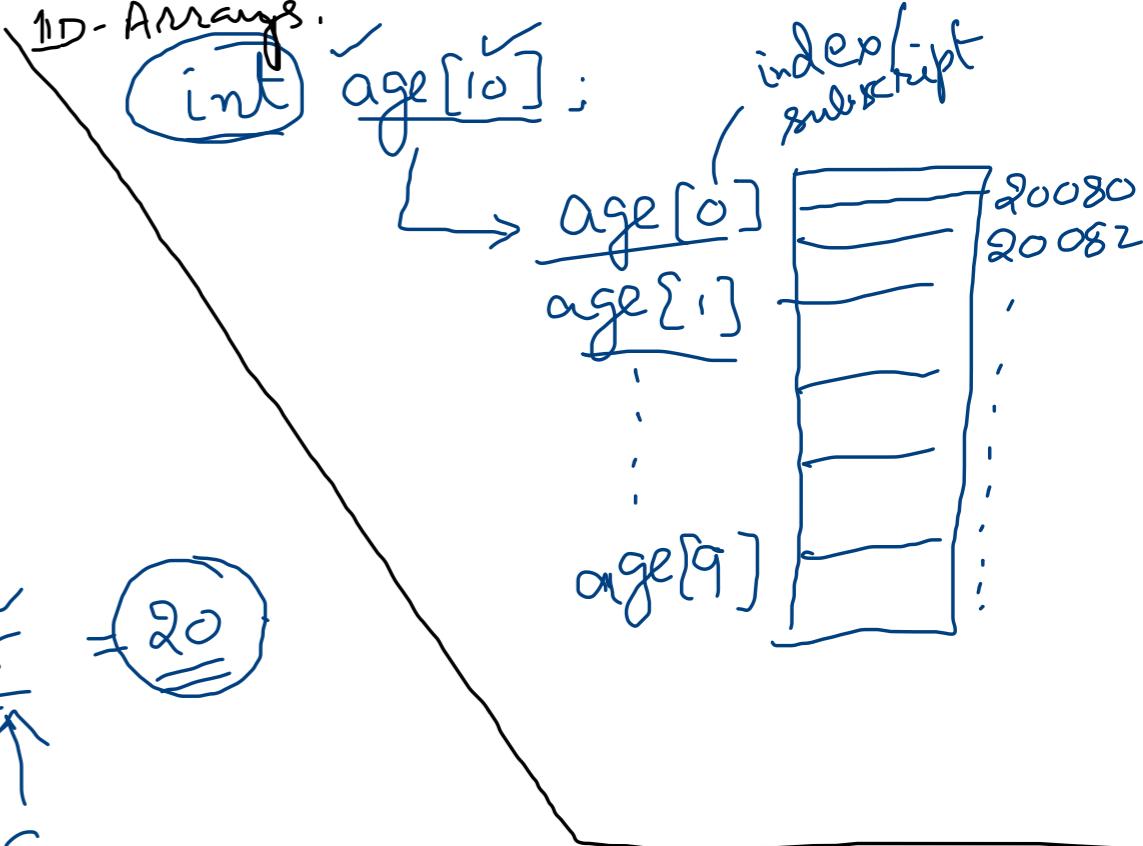
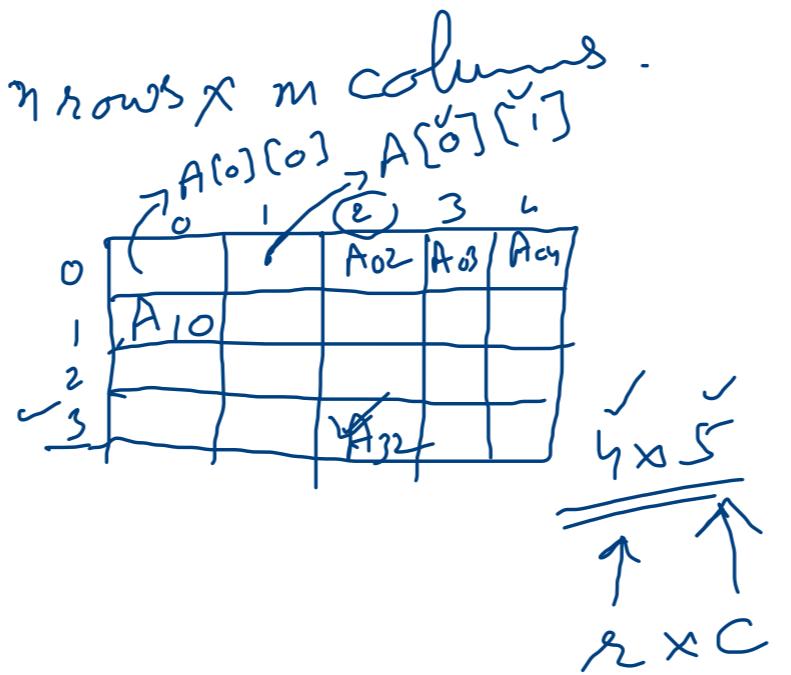
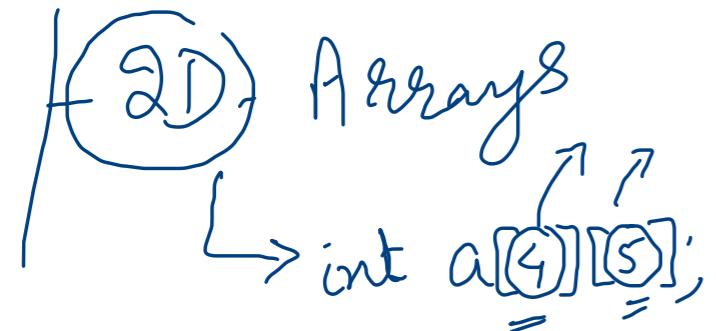
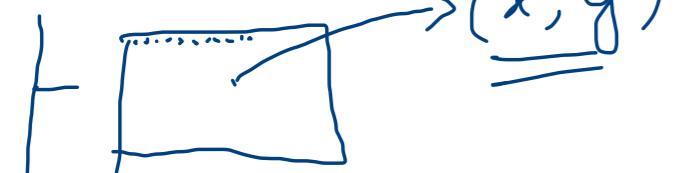


Unit - 2:

Multidimensional Arrays :



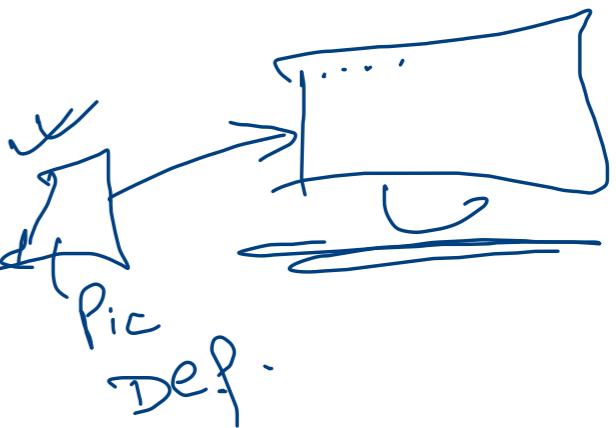
2D Arrays



Excel

- Mathematical objects \rightarrow Matrices

- Frame Buffer



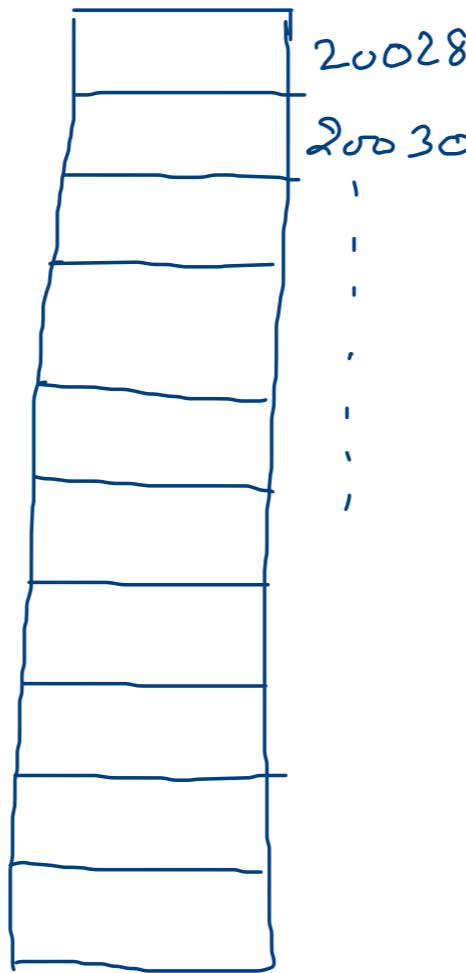
Representing 2D Arrays in Memory:

→ Row Major order .

→ Column Major order .

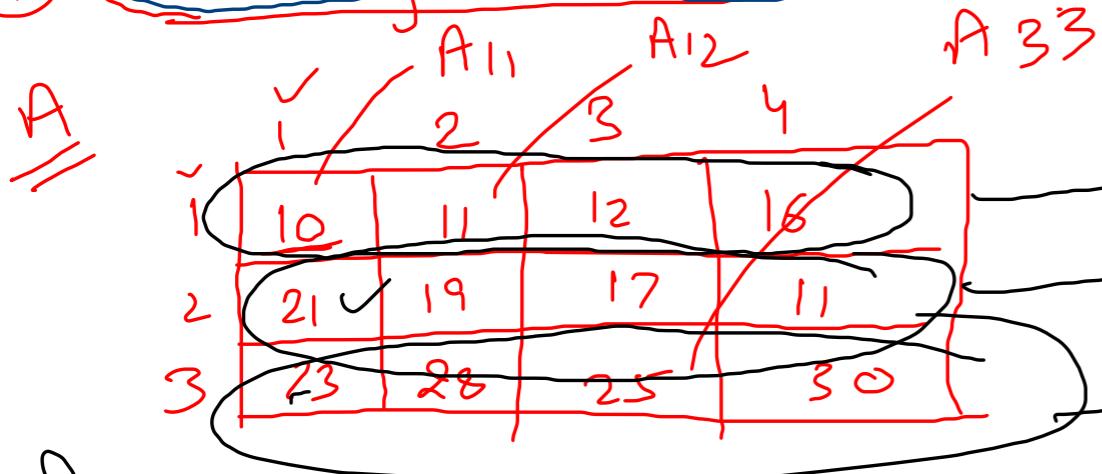
2D

a[4][5]

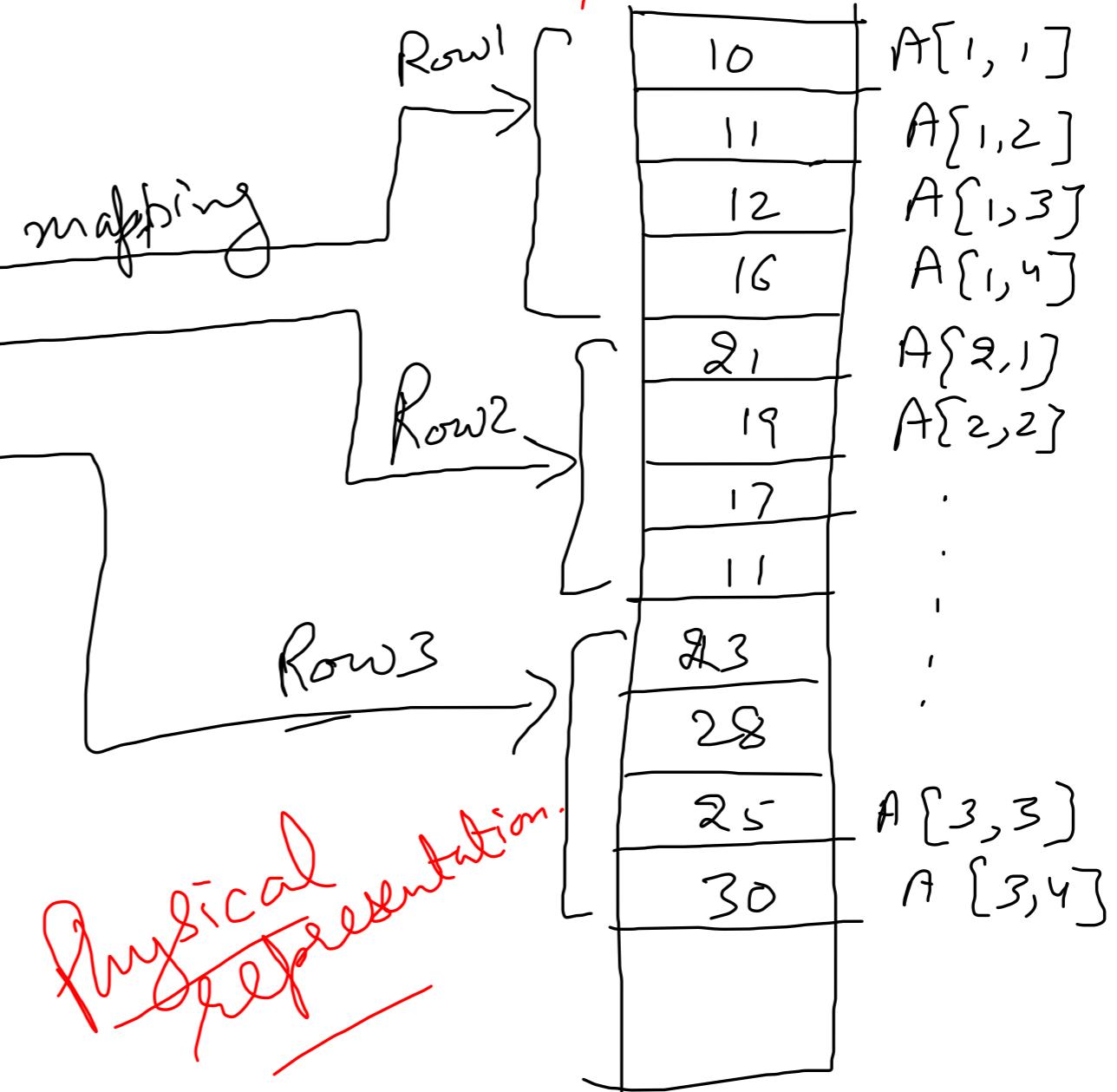


Rep. of 2D Arrays

① Row Major order:



✓ Logical representation
3x4 array A
 $r=3$
 $c=4$
Total Elements = $3 \times 4 = 12$
Used in various HLLs -
C, C++, Java, Pascal,
Modula 2, etc.
Ada,



How to calculate the address of an element of a 2D-Array ? in case of Row - Major order?

Let there be an $M \times N$ array A (ie $A[M, N]$)

Let it's base address be $\text{Base}(A)$

s = size of an element of the Array.

Q: To find address of specific element

$A[i, j]$

Multiply

$(i-1) \times \text{size of a row}$

$N \times s$

① Address of i^{th} Row : ie How many rows to skip over ie $(i-1)$

② Address of j^{th} Column:

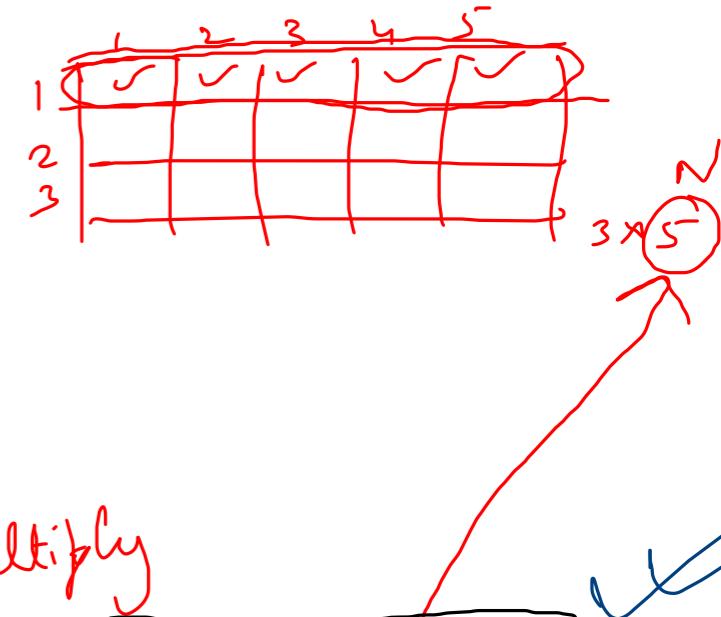
$(j-1) \times s$

\times

$(i-1) \times N \times s$

$$\text{Address of } A[i, j] = \boxed{\text{Base}(A)} + \boxed{(i-1) \times N \times s} + \boxed{(j-1) \times s}$$

$$\rightarrow = \boxed{\text{Base}(A) + s((i-1)N + (j-1))}$$



Formula:
 Address of $A[i, j] = \underline{\text{Base}(A)} + 8[(i-1)N + (j-1)]$

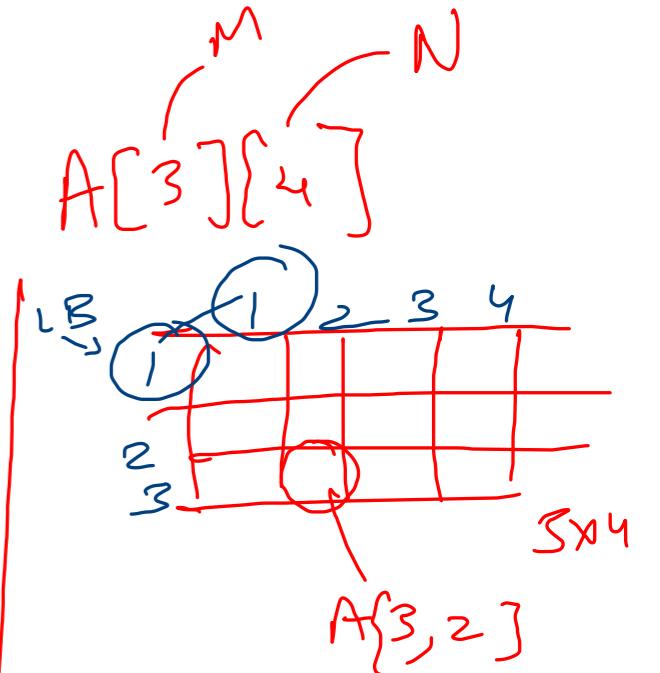
Example:

$$\begin{aligned}
 \text{Address of } A[3,2] &= 200 + 2[(3-1)4 + (2-1)] \\
 &= 200 + 2[(2)4 + 1] \\
 &= 200 + 2(8 + 1) \\
 &= 200 + 18 \\
 &= \underline{\underline{218}}
 \end{aligned}$$

$A[0,0]$

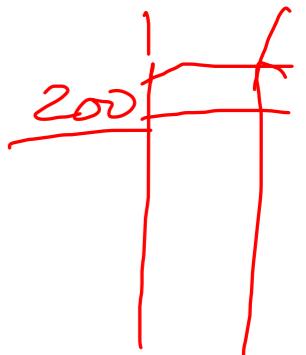
$A[10,18]$
LB UB

General Formula:
 Address of $A[i, j] = \underline{\text{Base}(A)} + 8[(i-lb)N + (j-ub)]$



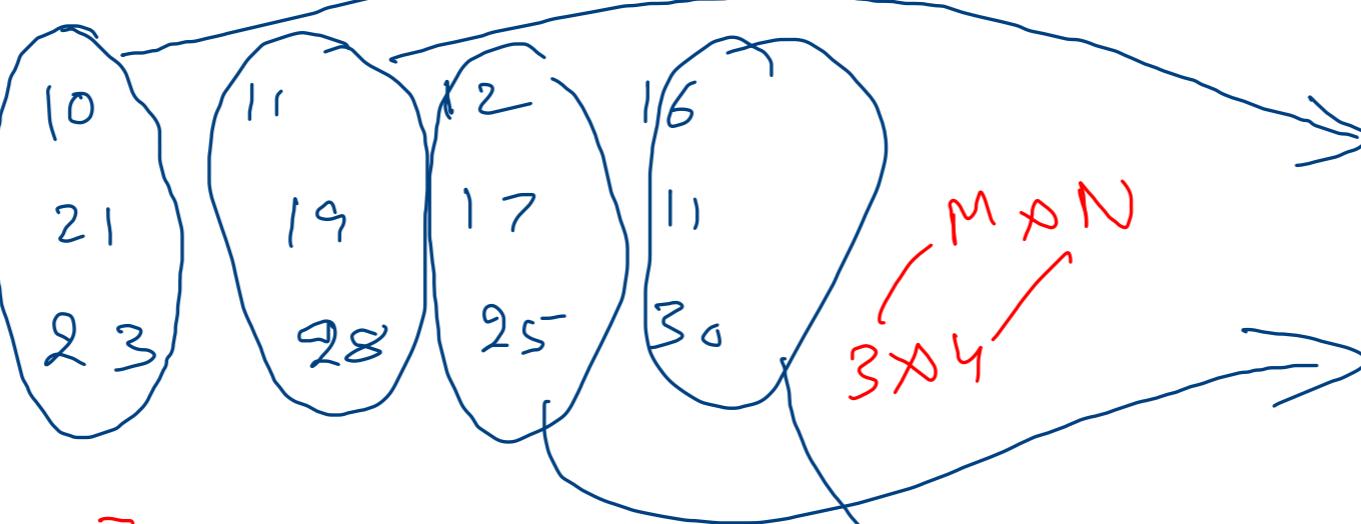
Base(A) = 200

$s = 2$ bytes



2nd Method of Representing 2D Arrays in Memory:

② Column Major Order.



Formula:

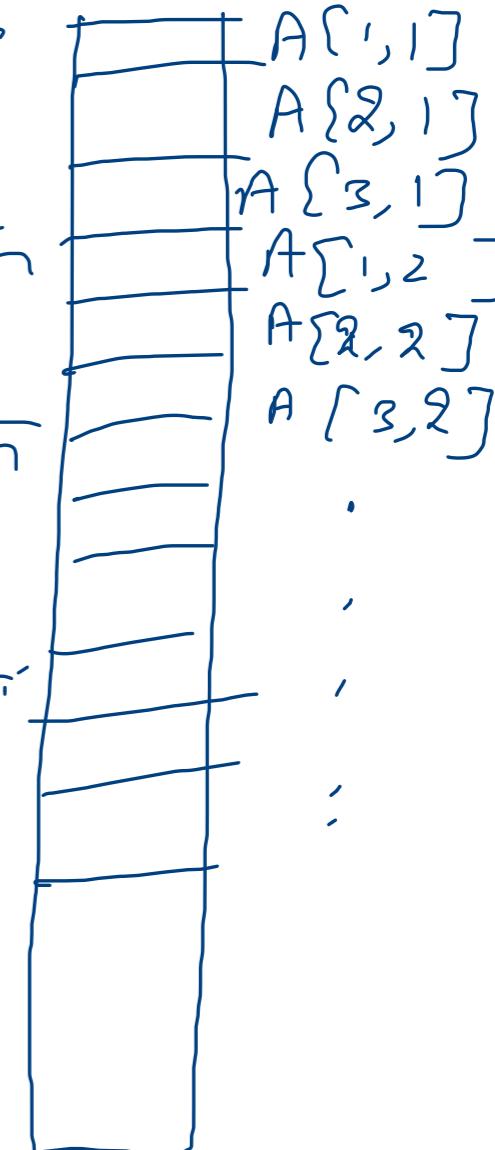
$$\text{Address of } A[i, j] = \text{Base}(A) + s[(j-1)M + (i-1)]$$

Eg: Address of $A[3, 2]$ = $200 + 2[(2-1)3 + (3-1)]$

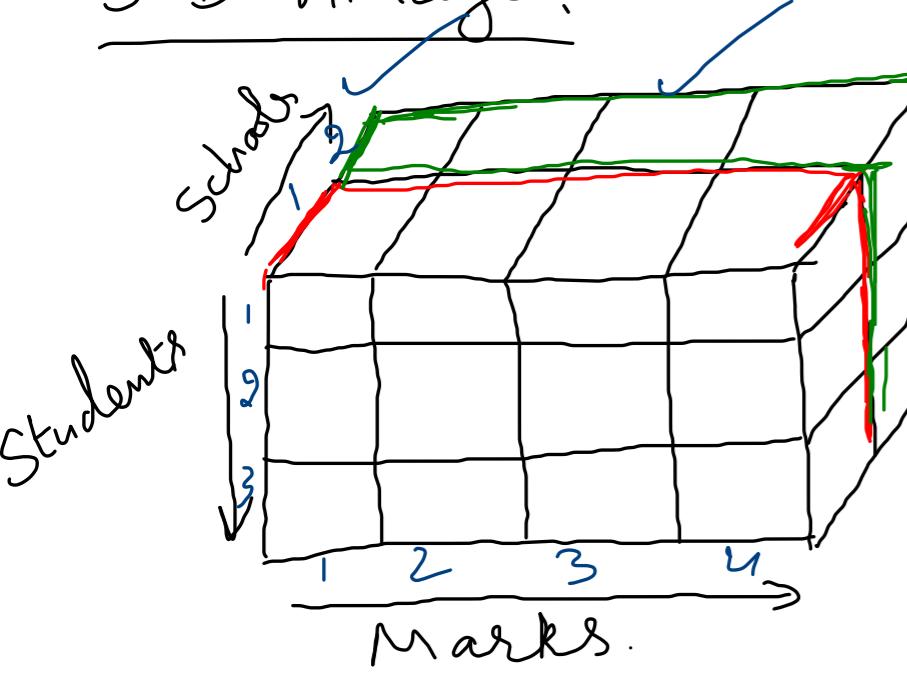
$= 200 + 2[3 + 2]$

$= 210$

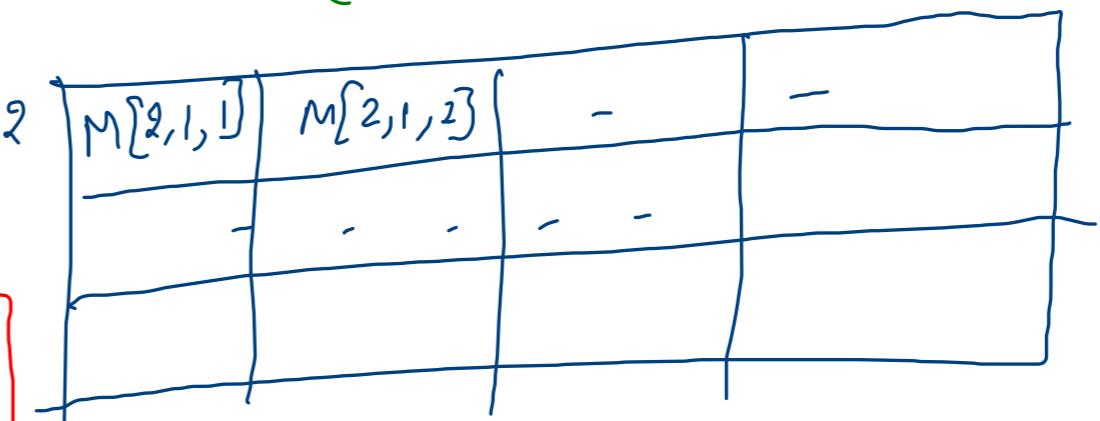
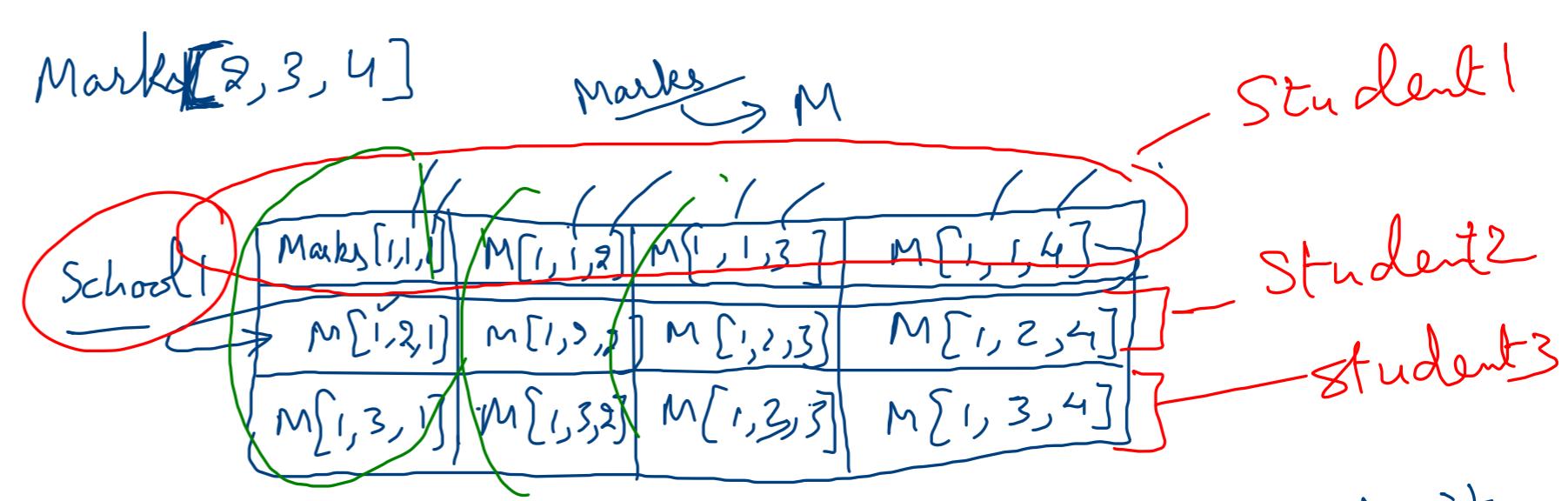
in
column
major
order



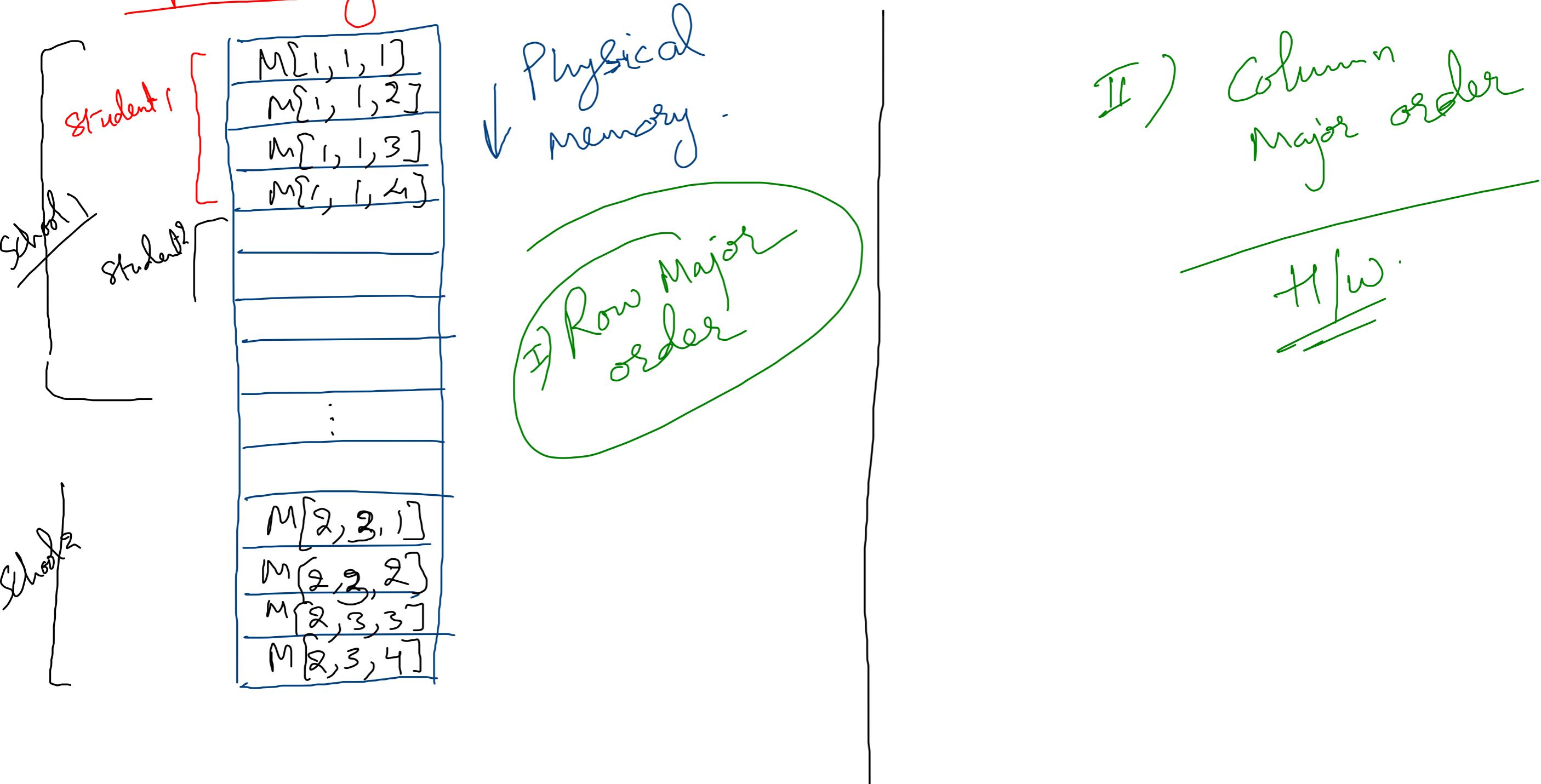
3-D Arrays:



```
for(i=1; i<=M; i++)  
    {  
        for(j=1; j<=N; j++)  
            {  
                for(k=1; k<=Q; k++)  
                    cin>> M[i][j][k];  
            }  
    }  
}
```



Representing 3D Arrays in memory:



Finding address of an element in 3D-Array.

$$1 \leq i \leq M, \quad 1 \leq j \leq N, \quad 1 \leq k \leq Q$$

$$(i-1)NQ + (j-1)Q$$

Formula:

$$\text{Address of } A[i, j, k] = \text{Base}(A) + \delta [((i-1)N + (j-1))Q + (k-1)]$$

e.g.: $A[2, 2, 3]$

$$\begin{aligned} &= 200 + 2[((2-1)3 + (2-1))4 + (3-1)] \\ &= 200 + 2[(16 + 2)] \\ &= 200 + 2(18) \\ &= 236 \end{aligned}$$

$$B(A) = 200$$

$$\delta = 2$$

$$M = 2$$

$$N = 3$$

$$Q = 4$$

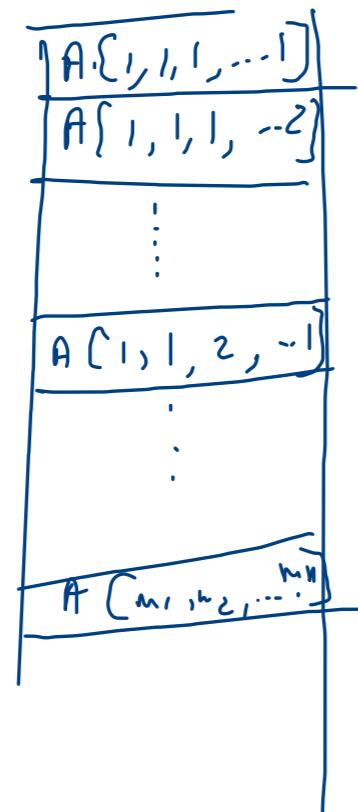
Generalization of Multidimensional Array:

↳ n -D array:

Range: $1 \leq k_1 \leq m_1, 1 \leq k_2 \leq m_2, \dots, 1 \leq k_n \leq m_n$.

Element: $A[k_1, k_2, \dots, k_n]$

Memory Rep.
(Row Major) order



Formula for finding address of an Element.

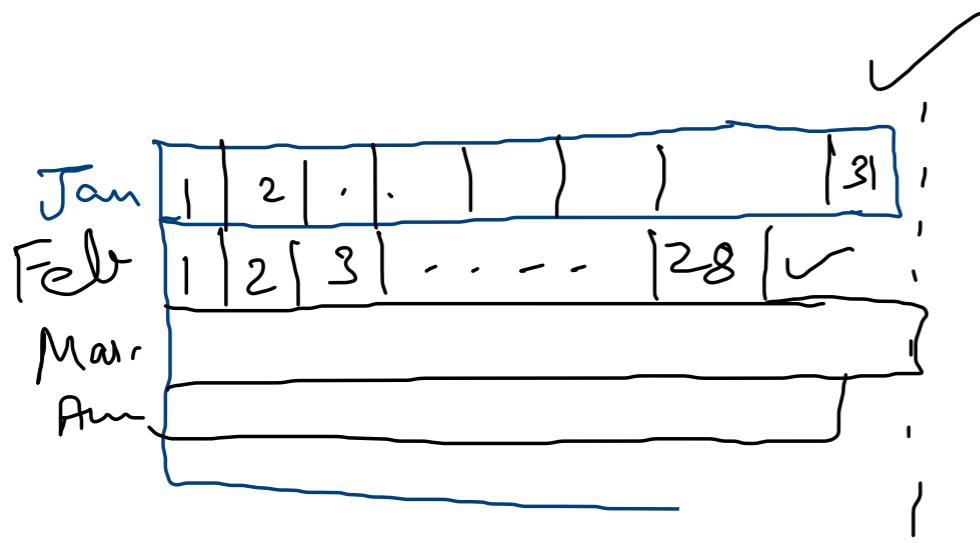
$$A[k_1, k_2, \dots, k_n] = \text{Base}(A)$$

$$+ \delta \left[(k_1 - 1)m_2 m_3 \dots m_n + (k_2 - 1)m_3 m_4 \dots m_n + \dots + (k_{n-1} - 1)m_n + (k_n - 1) \right]$$

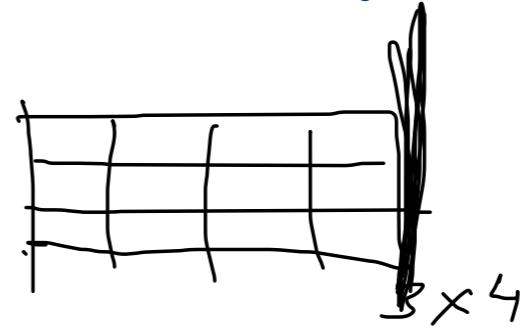
$$= \text{Base}(A) + \delta \left[\sum_{i=1}^n (p_i)(k_i - 1) \right]$$

$$p_i = \prod_{\substack{i < x \\ i \leq n}} m_x$$

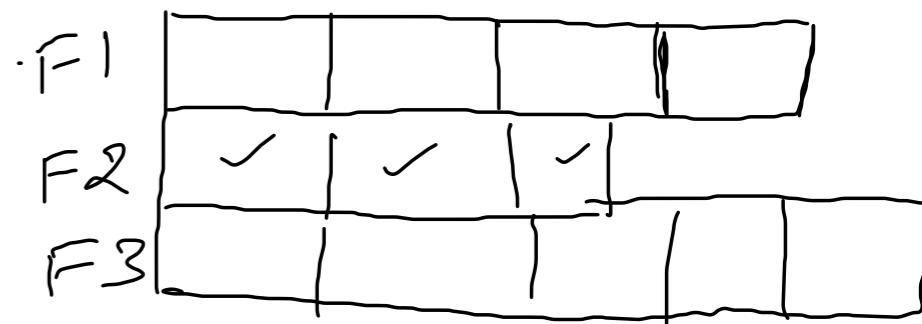
Jagged Arrays and Pointer Arrays:



~~Eg:-~~ Calander for a particular year.



An array in which no. of elements in each row need not be of same size is called a jagged array.



Storing Jagged arrays in Memory: (We must know about Pointer and array of pointers)

(P) is a var that points to an element of an array.
→ contains address of an element of an array.

PTR is an array of pointers.

