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#### PROBLEM A: EDIT DISTANCE

The edit distance between two words—sometimes also called the *Levenshtein* distance—is the *minimum* number of letter insertions, letter deletions, and letter substitutions required to transform one word into another.

For example, the edit distance between FOOD and MONEY is at most four:

FOOD → MOOD → MON\_D → MONED → MONEY

Given two strings, find the edit distance between them.

INPUT:

Line 1: the first string, A

Line 2: the second string, B

OUTPUT:

Edit distance between the two strings

- 1) We are transforming string A to string B. Assume that string A[0] .. A[i-1] have been transformed to be identical to B[0] .. B[j-1], and the consideration now is on A[i] and B[j].

The table below lists all possible scenarios at state (i, j) and edit operations that can be performed.

What is the consequential state for each combination of condition and operation ?

condition	edit operation	next state to consider
A[i] == B[j]	None	(i+1, j+1)
A[i] != B[j]	Insert B[j] in front of A[i]	(i, j+1)
A[i] != B[j]	Delete A[i]	(i+1, j)
A[i] != B[j]	Change A[i] to B[j]	(i+1, j+1)

- 2) What is the beginning state? (i = 0, j = 0)
- 3) If A runs out, but B has not yet, in other words, i == len(A), but j < len(B), what is the additional edit distance required to complete the transformation? len(B) - j
- 4) If B runs out, but A has not yet, what is the additional edit distance required to complete the transformation? len(A) - j
- 5) Use the concepts obtained from step 1 to 4 above in write a recursive brute-force solution for this problem. The zipped test case file is downloadable from Class Materials.
- 6) Given that a string can be up to 1000 letters long, improve the brute-force solution so that the program will finish in no more than 2.5 seconds (CPU processing time).

PROBLEM B: Dynamic Programming for Minimum Coin Change  
 (Targeted time to finish : not longer than 1 hour)

INPUT:

Line 1 : the list of coin denominator

Line 2 : the amount of change

OUTPUT: The minimum number of coins required for the change

EXAMPLE

INPUT	OUTPUT
1 3 4 5 7	2
1 2 5 10 13 3377	260

The following code is a memoized minimum coin change function.

```
mm = [-1] * (V+1)

def mincoin(v):
    global coin, mm

    if mm[v] == -1:
        if v == 0:
            mm[v] = 0
        else:
            minc = 10000000000
            for c in coin:
                if c <= v:
                    minc = min(minc, 1 + mincoin(v-c))
            mm[v] = minc
    return mm[v]
```

- Given that  $v_1 \geq v_2$ ,
  - which recursive call, to  $\text{mincoin}(v_1)$  or to  $\text{mincoin}(v_2)$ , is made first?  
**mincoin(v1) is made first**
  - which recursive function,  $\text{mincoin}(v_1)$  or  $\text{mincoin}(v_2)$ , returns first?  
**mincoin(v2) return first**
  - which  $\text{mm}$ 's entry,  $\text{mm}[v_1]$  or  $\text{mm}[v_2]$ , obtains its final value first?  
**mm[v2] get value first**
- Develop a **non-recursive** minimum coin change solution i.e. does not utilize recursive function, by iterating through  $\text{mm}$ 's indices with an appropriate sequence, computing value of corresponding  $\text{mm}$ 's entry along the way.