

Logic of Computer Science

Lecture 1: Foundations & Set Theory Review

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Learning Objectives

By the end of this session, you will be able to:

① **Recall basic set operations.**

- Form unions, intersections, set-differences, complements.
- Interpret notation $A \subseteq B$, $x \in A$, $A \cup B$, etc.

② **Reinforce proof techniques on set identities.**

- Practise direct proofs and contrapositives.
- Canonical identity example: $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$.

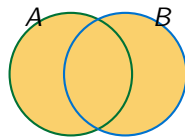
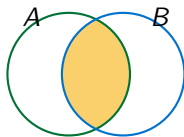
Key Definitions

1. Set, Element, Subset

- **Set** A : collection of distinct objects.
- **Element** $x \in A$: x belongs to A .
- **Subset** $A \subseteq B$: $\forall x(x \in A \Rightarrow x \in B)$.
- Empty set \emptyset .

2. Operations

- Union $A \cup B = \{x : x \in A \vee x \in B\}$.
- Intersection
 $A \cap B = \{x : x \in A \wedge x \in B\}$.
- Difference $A - B = \{x : x \in A \wedge x \notin B\}$.
- Complement $\bar{A} = U - A$.



Intersection vs. Union visual

Relations & Functions

Relation $R \subseteq A \times B$

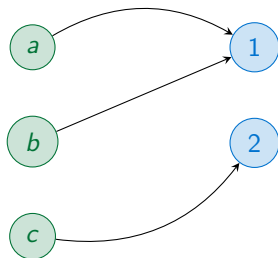
$$(a, b) \in R \Leftrightarrow a R b$$

$$\text{Domain}(R) = \{a \in A \mid \exists b (a, b) \in R\}$$

Function $f: A \rightarrow B$

Each $a \in A$ has exactly one $b \in B$ with $f(a) = b$.

- **Injective:** distinct inputs \Rightarrow distinct outputs.
- **Surjective:** every $y \in B$ hit by some x .
- **Bijective:** both injective and surjective.



Example:
surjective but not injective.

Proof Techniques Refresher

Set Equality $X = Y$

Show $X \subseteq Y$ and $Y \subseteq X$.

Direct Proof

Assume P ; derive Q .

Contrapositive

Prove $\neg Q \Rightarrow \neg P$ instead of $P \Rightarrow Q$.

Tip: identify P and Q clearly before choosing a method.

Example Proof — Distributive Law

Proposition

$$A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$$

\subseteq Let $x \in A \cap (B \cup C)$. Then $x \in A$ and $(x \in B \vee x \in C)$.

- If $x \in B \Rightarrow x \in A \cap B$.
- If $x \in C \Rightarrow x \in A \cap C$.

Hence $x \in (A \cap B) \cup (A \cap C)$.

\supseteq Let $y \in (A \cap B) \cup (A \cap C)$.

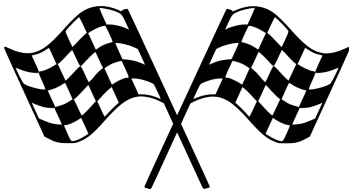
- If $y \in A \cap B \Rightarrow y \in A$ and $y \in B \Rightarrow y \in B \cup C$.
- Similar for $A \cap C$.

Thus $y \in A \cap (B \cup C)$.



In-Class Activity — Set-Proof Race

- 1 Split into **small groups**
- 2 Each group receives **one set identity** to prove, e.g.
 - $A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$
 - $\overline{A \cup B} = \overline{A} \cap \overline{B}$
 - $A \Delta B = (A \cup B) - (A \cap B)$
- 3 12 minutes to craft proofs & choose a presenter.
- 4 Rapid-fire presentations (2 minutes each).



Summary & Takeaways

Key Points





- Set theory is foundational for logic and programming semantics.
- Precise **definitions and notation** underpin logical reasoning.
- Proof techniques build rigor and form the basis for formal logic systems.
- These concepts set the stage for propositional and predicate logic.

Up Next — Week 2

Truth tables, logical connectives, more complex reasoning.

Home Practice: Prove $\overline{A \cap B} = \overline{A} \cup \overline{B}$.

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Thank You!