

Foolproof Proof-writing

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Introduction: High School Lied to You

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Who remembers writing proofs like this?

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Prove the identity

$$\cot(x) + \tan(x) = \cos(x) \csc(x) (\sin^2(x) \sec^2(x) + \sin^2(x) \csc^2(x)).$$

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$$= \cot(x) (\tan^2(x) + \cos^2(x) + \sin^2(x)) \quad (3)$$

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Can you spot any mistakes in this proof?

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Of course you can't!

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Of course you can't!

Because this isn't a good proof.

What is a good proof?

Any ideas?

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1. It has to be *clear*.

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Here are some of ours:

1. It has to be *clear*.
2. It has to have good *structure*.
3. It has to *flow*.

Outline

1. Structure
2. Clarity
3. Flow
4. One-on-One Feedback

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1. **Structure**
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Structure: Proofs as Essays

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- ▶ Write a first draft, using complete sentences.
- ▶ Proofread. (Literally)

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Example

~~It will be proved via contradiction...~~

We now prove via contradiction...

Structure: Sentence Structure

- ▶ Simple sentence structure is generally easier to read.
- ▶ Don't worry about sounding a little formulaic.
- ▶ Use the active voice.
- ▶ Try to only justify one thing per sentence.

Structure: Overall Structure

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- ▶ Some proof types have structure that you can use to your advantage!
 - ▶ Induction
 - ▶ Element Method
 - ▶ Bijections
 - ▶ Bidirectional Proofs (If and Only If)

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- ▶ Avoid using lists inside a proof.

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The `description` environment looks nice though!

Injectivity Proof of the injectivity of f would go here. It nicely aligns the paragraphs within the proof.

Surjectivity Proof of the surjectivity of f would go here.

Example Proof 1: Problem Statement

Consider the function $f : \mathbb{Z} \rightarrow \mathbb{E}$, $f(x) = 2x$.

Prove that f is a bijection.

Example Proof 1: Rough Draft

Proof.

It is necessary to show that f is surjective and injective, or that $f(x) \neq f(y) \implies x \neq y \forall x, y \in \mathbb{Z}$ and that $\forall y \in \mathbb{E}, \exists x \in \mathbb{Z}$ where $f(x) = y$. For any $y \in \mathbb{E}$ that you can think of, by definition of an even number, $y = 2x$ for some $x \in \mathbb{Z}$, since every even number can be divided by 2, no matter what. And if $f(x) \neq f(y)$, then $2x \neq 2y$ which would suggest that $x \neq y$. \square

Example Proof 1: Polished

Proof.

To prove that f is a bijection, we must show injectivity and surjectivity.

Injectivity Suppose we have $x, y \in \mathbb{Z}$ such that $f(x) \neq f(y)$.
Then $2x \neq 2y$, which means $x \neq y$, as needed.

Surjectivity Consider an arbitrary $y \in \mathbb{E}$. By definition of an even number, $y = 2x$ for some $x \in \mathbb{Z}$, as needed.

Thus, f is a bijection.



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3. Flow
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Clarity: Keeping the Reader Informed

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Example

To prove a function is odd, we must show...

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Example

To prove a function is odd, we must show...

In order to prove that R is an equivalence relation, we need...

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Example

Thus, we have...

But we recall from earlier that...

Combining this with our result from case 1...

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...by the Fundamental Theorem of Arithmetic.

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Example

...by the Fundamental Theorem of Arithmetic.

By definition of... (Sparingly!)

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- ▶ Conclusion: What did you just do?

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...thus we have reached a contradiction.

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Example

...thus we have reached a contradiction.

Since we have proven $P(1)$ and have shown $P(k)$ implies $P(k + 1)$, we have shown $P(n)$ for all $n \in \mathbb{Z}^+$.

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$\forall x \in S$

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 - ▶ Don't replace a single word with a single symbol, just like you wouldn't write " $3 + \text{four}$ ".
 - ▶ Similarly, don't write " $\text{for an element } \in S$ ". Be consistent within a given context.
- ▶ Short notation tips.

Example Proof 2: Problem Statement

Prove that there are infinitely many primes.

Example Proof 2: Rough Draft

Proof.

What if there were only finitely many primes? p_1, p_2 , through p_n is the finite list of all these primes.

$$Q = p_1 p_2 \cdots p_n + 1$$

If Q is prime, then Q is greater than $p_i = Q$ is not \in the list of primes. $\Rightarrow \Leftarrow$. If Q is not prime then $p_i \mid Q$ and p_i divides $p_1 p_2 \cdots p_n$. p_i doesn't divide 1. $Q - p_1 p_2 \cdots p_n = 1$. $\Rightarrow \Leftarrow$ □

Example Proof 2: Polished

Proof.

Assume for the sake of contradiction that there are finitely many primes. Let $P = \{p_1, p_2, \dots, p_n\}$ be the set of all primes. Now, let us consider $Q = p_1 p_2 \cdots p_n + 1$. We aim to show that Q can be neither prime nor composite. We consider the two cases:

Prime Suppose Q is prime. But $Q > p_i \forall i$, meaning that $Q \notin P$. This contradicts our definition of P .

Composite Suppose Q is not prime; by the Fundamental Theorem of Arithmetic, Q can be factored into primes. Consider p_i , one of these prime factors. Since $p_i \mid Q$ and $p_i \mid p_1 p_2 \cdots p_n$, we know that $p_i \mid (Q - p_1 p_2 \cdots p_n)$. But $Q - p_1 p_2 \cdots p_n = 1$, meaning that $p_i \mid 1$. This is a contradiction.

Thus, we have proven that there cannot be finitely many primes.



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Example

We are given that B_1, \dots, B_k partitions U **into distinct blocks such that every element in U is in some block.**

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- ▶ Exception: Recalling an earlier proven point or citing a sub-result out of context.

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Example

...it is a bijection. Because it is surjective...

Recall that R is an equivalence relation. By the transitivity of R ...

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- ▶ It suffices to show...

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In order to prove...
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- ▶ ...as needed.

Flow: Using Meaningful Transitions

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- ▶ We need to show...
In order to prove...
- ▶ It suffices to show...
- ▶ ...as needed.
- ▶ Suppose...

Flow: Using Meaningful Transitions

- ▶ Hence, thus, therefore.
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In order to prove...
- ▶ It suffices to show...
- ▶ ...as needed.
- ▶ Suppose...
- ▶ Let x ...

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In order to prove...
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- ▶ Consider...

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- ▶ Suppose...
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- ▶ Consider...
- ▶ Recall...

Flow: Using Meaningful Transitions

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- ▶ Consider...
- ▶ Recall...
- ▶ In particular...

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- ▶ Without loss of generality (wlog)

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- ▶ Consider...
- ▶ Recall...
- ▶ In particular...
- ▶ Without loss of generality (wlog)
- ▶ ~~Clearly, obviously, trivially~~

Example Proof 3: Problem Statement

Consider the following relation on the set of integers:

$\forall a, b \in \mathbb{Z}, (a, b) \in R$ if and only if a and b have the same remainder when divided by 3.

Prove that R is transitive.

Example Proof 3: Rough Draft

Proof.

We know that dividing integers by integers will yield integer remainders, by properties of division. So let r_a be the remainder when you divide a by 3. Similarly for r_b and r_c with b, c .

Definition of transitivity:

$$(a, b), (b, c) \in R \implies (a, c) \in R \quad \forall a, b, c \in \mathbb{Z}$$

so we need this to be true to show transitivity. (e.g.

$$(1, 2), (2, 3) \in R \implies (1, 3) \in R.)$$

Notice $(a, b) \in R \implies r_a = r_b$ and $(b, c) \in R \implies r_b = r_c$ so $r_a = r_c$.

So R is transitive because $(a, c) \in R$ for all $(a, b), (b, c) \in R$. □

Example Proof 3: Polished

Proof.

For transitivity to hold, we need

$$(a, b), (b, c) \in R \implies (a, c) \in R \quad \forall a, b, c \in \mathbb{Z}.$$

Let r_a , r_b , and r_c be the remainders when you divide a , b , and c by 3, respectively. Since $(a, b) \in R$, we know that $r_a = r_b$. Since $(b, c) \in R$, we know that $r_b = r_c$. Thus, by the transitivity of equality, we have $r_a = r_c$. By definition of the relation R , $(a, c) \in R$, as needed.

Thus, we have shown that R is transitive. □

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