

CMP2020M Artificial Intelligence

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CM2020M Artificial Intelligence



- Introduction
- Logic Programming(**1**/3)
- Knowledge Representations
- Games AI & Search(/3)
- Planning I
- Planning II
- Probabilistic AI

My Lectures



1. Propositional Calculus
2. Predicate Calculus
3. Prolog Programming
4. Knowledge Representation

Propositional Calculus



- What is Logic Programming
- The Propositional Calculus
 - Syntax
 - Semantics
 - Proofs
- Next week: Predicate calculus and Unification

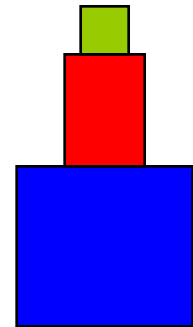
What is Logic Programming?



- In logic programming, you present facts and rules to **infer** new facts by just ask questions.
- When you asked a question, the run time system searches through the database of facts and rules to determine (by **logical deduction**) the answer.

Example

- Given:
 - "The red block is above the blue block"
 - "The green block is above the red block"
- Infer?
 - "The green block is above the blue block"
 - "The blocks form a tower"



Logic consists of



- A language
 - tells us how to build up sentences (i.e., *syntax*), and what that sentences mean (i.e., *semantics*)
- An inference procedure
 - which tells us which sentences are valid inference from given sentences

Propositional Logic

- The **syntax** of propositional logic consists of
 - the propositional symbols:
 - P, Q, R, S, \dots
 - and connectives:
 - $\wedge, \vee, \neg, \rightarrow, \equiv$
- The **semantics (interpretation)** is assigning a truth value (T or F) to each sentence.

Examples of Propositional Logic sentences

- P means "It is hot."
- Q means "It is humid."
- R means "It is raining."
- $(P \wedge Q) \rightarrow R$
"If it is hot and humid, then it is raining"
- $Q \rightarrow P$
"If it is humid, then it is hot"

Propositional calculus semantics

P	$\neg P$
T	F
F	T

Truth table for the operator \neg Negation (not)

Propositional calculus semantics

P	Q	$P \wedge Q$
T	T	T
T	F	F
F	T	F
F	F	F

Truth table for the operator \wedge Conjunction (and)

Propositional calculus semantics

P	Q	$P \vee Q$
T	T	T
T	F	T
F	T	T
F	F	F

Truth table for the operator \vee Disjunction (or)

Propositional calculus semantics

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

Truth table for the operator \rightarrow Implication

Propositional calculus semantics

P	Q	$P \equiv Q$
T	T	T
T	F	F
F	T	F
F	F	T

**Truth table for the operator $\leftrightarrow \Leftrightarrow \equiv$ Equivalence
(if and only if)**

Truth table is demonstrating the equivalence of $\neg P \vee Q$ and $P \rightarrow Q$

P	Q	$\neg P$	$\neg P \vee Q$	$P \rightarrow Q$	$(\neg P \vee Q) \equiv (P \rightarrow Q)$

Truth table is demonstrating the equivalence of $\neg P \vee Q$ and $P \rightarrow Q$

P	Q	$\neg P$	$\neg P \vee Q$	$P \rightarrow Q$	$(\neg P \vee Q) \equiv (P \rightarrow Q)$
T	T				
T	F				
F	T				
F	F				

Truth table is demonstrating the equivalence of $\neg P \vee Q$ and $P \rightarrow Q$

P	Q	$\neg P$	$\neg P \vee Q$	$P \rightarrow Q$	$(\neg P \vee Q) \equiv (P \rightarrow Q)$
T	T	F			
T	F	F			
F	T	T			
F	F	T			

Truth table is demonstrating the equivalence of $\neg P \vee Q$ and $P \rightarrow Q$

P	Q	$\neg P$	$\neg P \vee Q$	$P \rightarrow Q$	$(\neg P \vee Q) \equiv (P \rightarrow Q)$
T	T	F	T		
T	F	F	F		
F	T	T	T		
F	F	T	T		

Truth table is demonstrating the equivalence of $\neg P \vee Q$ and $P \rightarrow Q$

P	Q	$\neg P$	$\neg P \vee Q$	$P \rightarrow Q$	$(\neg P \vee Q) \equiv (P \rightarrow Q)$
T	T	F	T	T	
T	F	F	F	F	
F	T	T	T	T	
F	F	T	T	T	

Truth table is demonstrating the equivalence of $\neg P \vee Q$ and $P \rightarrow Q$

P	Q	$\neg P$	$\neg P \vee Q$	$P \rightarrow Q$	$(\neg P \vee Q) \equiv (P \rightarrow Q)$
T	T	F	T	T	T
T	F	F	F	F	T
F	T	T	T	T	T
F	F	T	T	T	T

Tautology and Contradiction

- A **tautology** is a formula which is "always true".
 - It is **true** for every assignment of truth values to its simple components.
- The opposite of a tautology is a **contradiction**, a formula which is "always false".
 - It is **false** for every assignment of truth values to its simple components.

Exercise 1

- Use a truth table to demonstrate the equivalence of $\neg(P \vee Q)$ and $(\neg P \wedge \neg Q)$

P	Q	$P \vee Q$	$\neg(P \vee Q)$	$\neg P$	$\neg Q$	$\neg P \wedge \neg Q$	$\neg(P \vee Q) \equiv (\neg P \wedge \neg Q)$
T	T						
T	F						
F	T						
F	F						

Exercise 1 : Answer

- Use a truth table to demonstrate the equivalence of $\neg(P \vee Q)$ and $(\neg P \wedge \neg Q)$

P	Q	$P \vee Q$	$\neg(P \vee Q)$	$\neg P$	$\neg Q$	$\neg P \wedge \neg Q$	$\neg(P \vee Q) \equiv (\neg P \wedge \neg Q)$
T	T	T	F				
T	F	T	F				
F	T	T	F				
F	F	F	T				

Exercise 1 : Answer

- Use a truth table to demonstrate the equivalence of $\neg(P \vee Q)$ and $(\neg P \wedge \neg Q)$

P	Q	$P \vee Q$	$\neg(P \vee Q)$	$\neg P$	$\neg Q$	$\neg P \wedge \neg Q$	$\neg(P \vee Q) \equiv (\neg P \wedge \neg Q)$
T	T	T	F	F	F		
T	F	T	F	F	T		
F	T	T	F	T	F		
F	F	F	T	T	T		

Exercise 1 : Answer

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P	Q	$P \vee Q$	$\neg(P \vee Q)$	$\neg P$	$\neg Q$	$\neg P \wedge \neg Q$	$\neg(P \vee Q) \equiv (\neg P \wedge \neg Q)$
T	T	T	F	F	F	F	
T	F	T	F	F	T	F	
F	T	T	F	T	F	F	
F	F	F	T	T	T	T	

Exercise 1 : Answer

- Use a truth table to demonstrate the equivalence of $\neg(P \vee Q)$ and $(\neg P \wedge \neg Q)$

P	Q	$P \vee Q$	$\neg(P \vee Q)$	$\neg P$	$\neg Q$	$\neg P \wedge \neg Q$	$\neg(P \vee Q) \equiv (\neg P \wedge \neg Q)$
T	T	T	F	F	F	F	T
T	F	T	F	F	T	F	T
F	T	T	F	T	F	F	T
F	F	F	T	T	T	T	T

The $\neg(P \vee Q) \equiv (\neg P \wedge \neg Q)$ is a tautology sentence

Proving things



- A **proof** is a sequence of sentences, where each sentence is either a premise or a sentence derived from earlier sentences in the proof by one of the inference rules.
- The last sentence is the **theorem** (also called goal or query) that we want to **prove**.

Exercise 2



- Symbolise the following propositions. The meaning of symbols chosen for statements also needs to be given:
- Terry likes Science-fiction and going to the Gym.
- Lectures are valuable if and only if they are structured carefully
- Today is Thursday or Yesterday was Wednesday.
- If I do not work hard then I will not pass my AI exam.

Exercise 2: Answer

- Terry likes Science-fiction and going to the Gym.
 - T : Terry likes Science fiction, G : Terry likes going to the gym
 - $T \wedge G$
- Lectures are valuable if and only if they are structured carefully
 - L : Lectures are valuable, S : Lectures are structured carefully
 - $L \equiv S$
- Today is Thursday or Yesterday was Wednesday.
 - T : Today is Thursday, Y : Yesterday was Wednesday
 - $T \vee Y$
- If I do not work hard then I will not pass my AI exam.
 - W : I work hard, A : I will pass my AI exam
 - $\neg W \rightarrow \neg A$

Proofs in Propositional Calculus

- If it is sunny today, then the sun shines on the screen. If the sun shines on the screen, the blinds are brought down. The blinds are not down. (*these are the given premises*)
 - Is it sunny today? (*this is the goal*)
 - P: It is sunny today.
 - Q: The sun shines on the screen.
 - R: The blinds are down.
 - Given: $P \rightarrow Q, Q \rightarrow R, \neg R$
 - Question: P
-
- Premise
- goal

Proof using a truth table

Variables			Given			Trial Conclusions	
P	Q	R	$P \rightarrow Q$	$Q \rightarrow R$	$\neg R$	P	$\neg P$
T	T	T					
T	T	F					
T	F	T					
T	F	F					
F	T	T					
F	T	F					
F	F	T					
F	F	F					

If it is sunny today, then the sun shines on the screen. If the sun shines on the screen, the blinds are brought down. The blinds are not down.

Is it sunny today? (*this is the goal*)

Proof using a truth table

Variables			Given			Trial Conclusions	
P	Q	R	$P \rightarrow Q$	$Q \rightarrow R$	$\neg R$	P	$\neg P$
T	T	T	T	T	F	T	F
T	T	F	T	F	T	T	F
T	F	T	F	T	F	T	F
T	F	F	F	T	T	T	F
F	T	T	T	T	F	F	T
F	T	F	T	F	T	F	T
F	F	T	T	T	F	F	T
F	F	F	T	T	T	F	T

Given:

$P \rightarrow Q$

$Q \rightarrow R$

$\neg R$

Question: P

Given: $P \rightarrow Q$, $Q \rightarrow R$, $\neg R$
Question: P

Proof using a truth table

Variables			Given			Trial Conclusions	
P	Q	R	$P \rightarrow Q$	$Q \rightarrow R$	$\neg R$	P	$\neg P$
T	T	T	T	T	F	T	F
T	T	F	T	F	T	T	F
T	F	T	F	T	F	T	F
T	F	F	F	T	T	T	F
F	T	T	T	T	F	F	T
F	T	F	T	F	T	F	T
F	F	T	T	T	F	F	T
Answer:	F	F	T	T	T	F	T

It is not sunny today

Proof procedure



- The problem with proof using truth tables is that the number of rows required grows very quickly as the number of propositional variables increases (2^n).
- A **proof procedure** is another method of proving statements using **inference rules**.

Inference rules

- *Modus ponens*

If P and $P \rightarrow Q$ are true, then infer Q .

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

Inference rules

- *Modus ponens*

If P and $P \rightarrow Q$ are true, then infer Q .

- *Modus tollens*

If $P \rightarrow Q$ and $\neg Q$ are true, then infer $\neg P$.

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

Inference rules

- *Modus ponens*

If P and $P \rightarrow Q$ are true, then infer Q .

- *Modus tollens*

If $P \rightarrow Q$ and $\neg Q$ are true, then infer $\neg P$.

- *And elimination*

If $P \wedge Q$ is true, then infer both P and Q are true

P	Q	$P \wedge Q$
T	T	T
T	F	F
F	T	F
F	F	F

Inference rules

- *Modus ponens*

If P and $P \rightarrow Q$ are true, then infer Q .

- *Modus tollens*

If $P \rightarrow Q$ and $\neg Q$ are true, then infer $\neg P$.

- *And elimination*

If $P \wedge Q$ is true, then infer both P and Q are true

- *And introduction*

If both P and Q are true, then infer $P \wedge Q$

P	Q	$P \wedge Q$
T	T	T
T	F	F
F	T	F
F	F	F

The Rules of Inference

Rule of inference	Tautology	Name
$\begin{array}{l} p \rightarrow q \\ \underline{p} \\ \hline \therefore q \end{array}$	$[p \wedge (p \rightarrow q)] \rightarrow q$	Modus ponens
$\begin{array}{l} \neg q \\ \underline{p \rightarrow q} \\ \hline \therefore \neg p \end{array}$	$[\neg q \wedge (p \rightarrow q)] \rightarrow \neg p$	Modus tollens
$\begin{array}{l} p \rightarrow q \\ \underline{q \rightarrow r} \\ \hline \therefore p \rightarrow r \end{array}$	$[(p \rightarrow q) \wedge (q \rightarrow r)] \rightarrow (p \rightarrow r)$	Hypothetical syllogism
$\begin{array}{l} p \vee q \\ \underline{\neg p} \\ \hline \therefore q \end{array}$	$((p \vee q) \wedge \neg p) \rightarrow q$	Disjunctive syllogism
$\begin{array}{l} \underline{p} \\ \hline \therefore p \vee q \end{array}$	$p \rightarrow (p \vee q)$	Addition
$\begin{array}{l} \underline{p \wedge q} \\ \hline \therefore p \end{array}$	$(p \wedge q) \rightarrow p$	Simplification
$\begin{array}{l} p \\ \underline{q} \\ \hline \therefore p \wedge q \end{array}$	$((p) \wedge (q)) \rightarrow (p \wedge q)$	Conjunction
$\begin{array}{l} p \vee q \\ \underline{\neg p \vee r} \\ \hline \therefore q \vee r \end{array}$	$[(p \vee q) \wedge (\neg p \vee r)] \rightarrow (q \vee r)$	Resolution

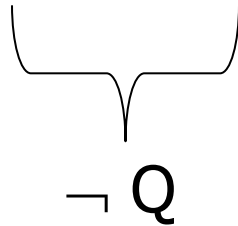
Proof using inference rules

- P: It is sunny today.
- Q: The sun shines on the screen.
- R: The blinds are down.
- Given: $P \rightarrow Q$, $Q \rightarrow R$, $\neg R$
- Question: P

Modus ponens: If P and $P \rightarrow Q$ are true, then infer Q .
Modus tollens: If $P \rightarrow Q$ and $\neg Q$ are true, then infer $\neg P$.
And elimination: If $P \wedge Q$ is true, then infer both P and Q are true
And introduction: If both P and Q are true, then infer $P \wedge Q$

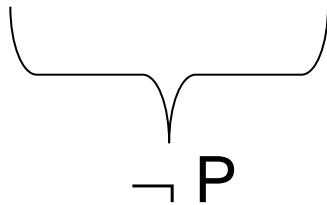
Proof using inference rules

- $P \rightarrow Q, Q \rightarrow R, \neg R$



Given

Modus Tollens



Modus Tollens

So P is false

It is not sunny today

Modus ponens: If P and $P \rightarrow Q$ are true, then infer Q .

Modus tollens: If $P \rightarrow Q$ and $\neg Q$ are true, then infer $\neg P$.

And elimination: If $P \wedge Q$ is true, then infer both P and Q are true

And introduction: If both P and Q are true, then infer $P \wedge Q$

Example 2:

- Consider the following statements for the “weather problem”.

- | | | | |
|----|---|---------|-------------------------------------|
| 1. | Humid | Premise | “It is humid” |
| 2. | Humid \rightarrow Hot | Premise | “If it is humid, it is hot” |
| 3. | (Hot \wedge Humid) \rightarrow Rain | Premise | “If it’s hot & humid, it’s raining” |

Is it raining?

Goal

- | | | | |
|----|--------------------|-------------------------|-----------------------|
| 4. | Hot | Modus Ponens (1, 2) | “It is hot” |
| 5. | Hot \wedge Humid | And Introduction (1, 4) | “It is hot and humid” |
| 6. | Rain | Modus Ponens (3, 5) | “It is raining” |

Exercise 3

- [illegible]

Exercise 3: Answer

- Consider the following statements:
- IF food is older than 7 days THEN it is unsafe to eat. The food is 8 days old.
- Is it save to eat the food? why?
 - F: food is older than 7 days.
 - E: food is safe to eat
 - given:
 - $F \rightarrow \neg E$
 - F
 - Infer
 - $\neg E$
 - The food is unsafe to eat
- What inference rule has been used?
 - Modus Ponens : If P and $P \rightarrow Q$ are true, then infer Q.

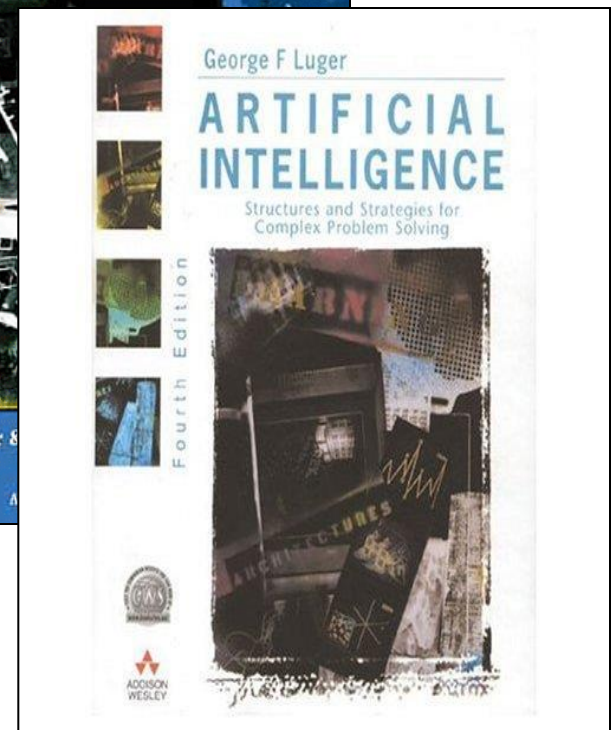
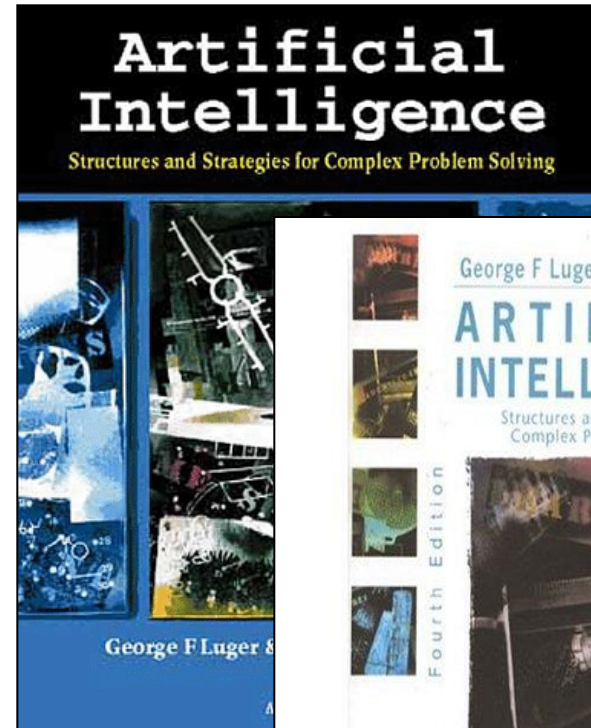
Next lecture: Predicate Calculus



- Predicate Calculus
- Unification
- Substitution

A little reading

- Luger, G.F. and Stubblefield, W.A., **Artificial Intelligence – Structures and Strategies for Complex Problem Solving**, (Addison-Wesley, 1998).
- Chapter 2 (3rd edition) 'The Predicate Calculus'
- take care with alternative editions – chapter numbers differ, and later editions are by Luger only



Thank you for listening!

