CAT3053/N Distributed Computing

Communication in Distributed Systems

Objectives

- To identify different layers of communication from protocols to Application Programming Interfaces (APIs).
- To present the various ways and widely-used models of communication in distributed systems
- To specify different paradigms of communications
- To be able to program applications with distributed objects and Internet protocols

References

- Distributed Systems: Principles and Paradigms by Tanenbaum (Chap 2, 3)
- Distributed Systems: Concepts and Design by Coulouris (Chap 3, 4, 5, 6)
- 3. Distributed and Parallel Computing by El-Rewini and Lewis (10,11)

Overview

- Protocols
- Models
 - Remote Procedure Call
 - Remote Method Invocation
 - Message Passing Interface Standard (MPI)
 - Streams
- Topology of communication
 - Client server
 - Peer to peer
 - Group

What is MPI?

- A message passing library specification
 - Message-passing model
 - Programming model which assumes each processor has its own local memory. Data are transferred by means of message passing.
 - Not a compiler specification (i.e. not a language)
 - Not a specific product
- Designed for parallel computers, clusters, and heterogeneous networks

The MPI Process

- Development began in early 1992
- Open process/Broad participation
 - IBM,Intel, TMC, Meiko, Cray, Convex, Ncube
 - PVM, p4, Express, Linda, ...
 - Laboratories, Universities, Government
- Final version of draft in May 1994
- Public (LAM/MPI, MPICH) and vendor implementations are now widely available
- de facto standard for message-passing-based application.
-Supersedes earlier libraries.

MPI-1 Services

- Supports point-to-point communication
 - communication modes blocking and non-blocking (standard, buffered, synchronise and ready)
- collective communication
 - data movement (gather/scatter)
 - reduction/collective computation
- virtual topologies conceptualise processes in an application-oriented topology (grids and graphs)
- derived datatype user-defined datatype

MPI-2 Services

- Dynamic task creation
- Client/server functions
- One-sided communication
- □ Parallel I/O

Communicators

- The notion of context and group are combined in a single communicator (object)
- All MPI communication calls take a communicator handle as a parameter, which is effectively the context in which the communication will take place
- E.g. MPI_INIT defines a communicator called MPI_COMM_WORLD for each process that calls it (default)

Communicator (cont.)

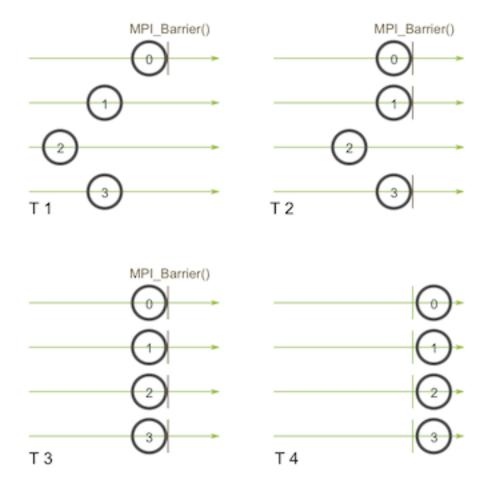
- Every communicator contains a group which is a list of processes
- The processes are ordered and numbered consecutively from 0.
- The number of each process is known as its rank
 - The rank identifies each process within the communicator
- The group of MPI_COMM_WORLD is the set of all MPI processes

Types of Collective Transfers

Barrier

- Synchronizes processors
- No data is exchanged but the barrier blocks until all processes have called the barrier routine
- Broadcast (sometimes multicast)
 - A broadcast is a one-to-many communication
 - One processor sends one message to several destinations
- Reduction
 - Often useful in a many-to-one communication

Barrier



https://mpitutorial.com/tutorials/mpi-broadcast-and-collective-communication/

What's in a Message?

- An MPI message is an array of elements of a particular MPI datatype
- All MPI messages are typed
 - The type of the contents must be specified in both the send and the receive

Basic C Datatypes in MPI

advanced datatype (derived datatype) - user construct own datatypes using MPI routines and data can be non-continuous

MPI Datatype	C datatype
MPI_CHAR	signed char
MPI_SHORT	signed short int
MPI_INT	signed int
MPI_LONG	signed long int
MPI_UNSIGNED_CHAR	unsigned char
MPI_UNSIGNED_SHORT	unsigned short int
MPI_UNSIGNED_INT	unsigned int
MPI_UNSIGNED_LONG	unsigned long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_BYTE	
MPI_PACKED	

MPI Basic

- comprises 129 functions
- a complete message passing programs can be written using 6 functions

MPI_INIT
MPI_FINALIZE
MPI_COMM_SIZE
MPI_COMM_RANK
MPI_SEND
MPI_RECV

Initiate an MPI computation.
Terminate a computation.
Determine num. of processes
Determine my process id.
Send a message.
Receive a message.

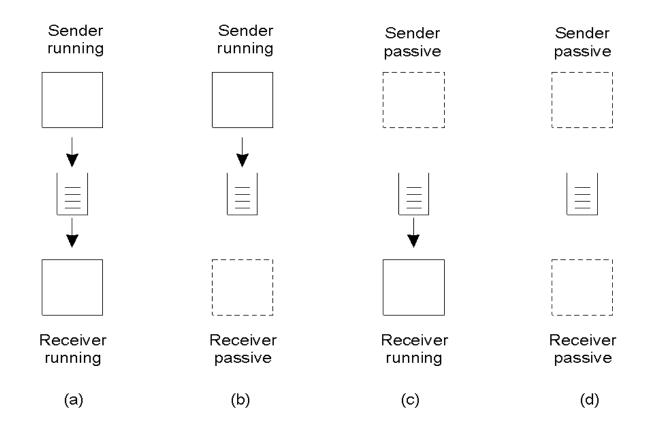
Skeleton MPI Program

```
program main
begin
    MPI_INIT ()
    MPI_COMM_SIZE (MPI_COMM_WORLD, count)
    MPI_COMM_RANK (MPI_COMM_WORLD, myid)
    print ("I am", myid, "of", count)

    MPI_FINALIZE ()
end
```

Message-Queuing Model (1)

 Four combinations for loosely-coupled communications using queues.

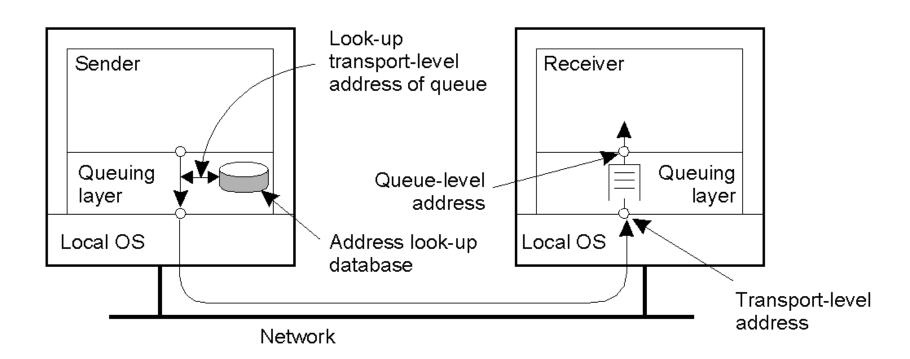


Message-Queuing Model (2)

Basic interface to a queue in a message-queuing system.

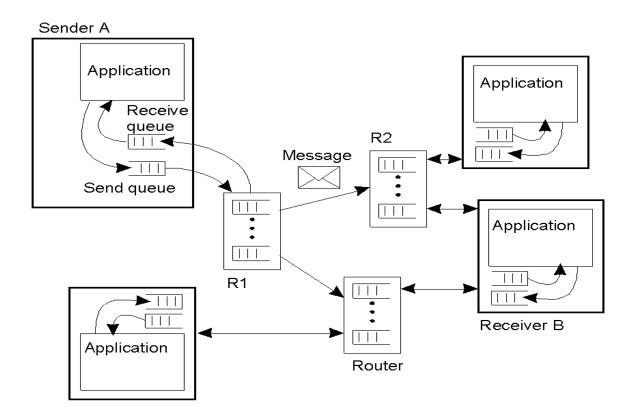
Primitive	Meaning
Put	Append a message to a specified queue
Get	Block until the specified queue is nonempty, and remove the first message
Poll	Check a specified queue for messages, and remove the first. Never block.
Notify	Install a handler to be called when a message is put into the specified queue.

General Architecture of a Message-Queuing System (1)

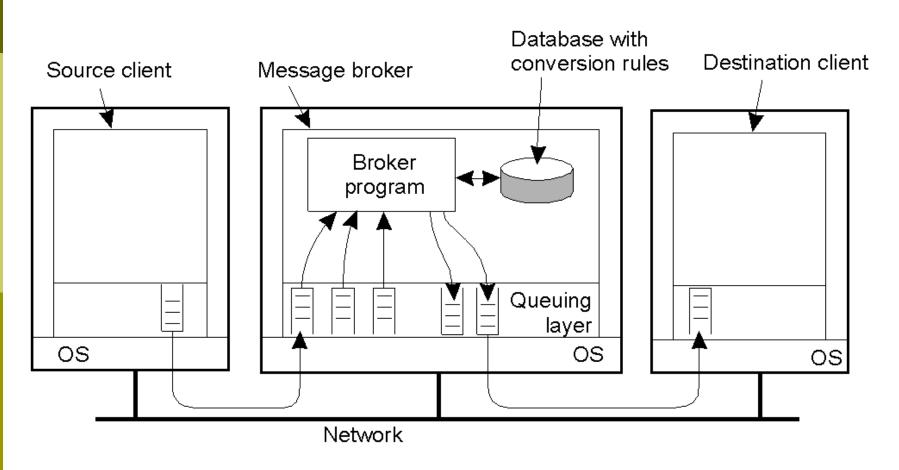


General Architecture of a Message-Queuing System (2)

The general organization of a message-queuing system with routers.



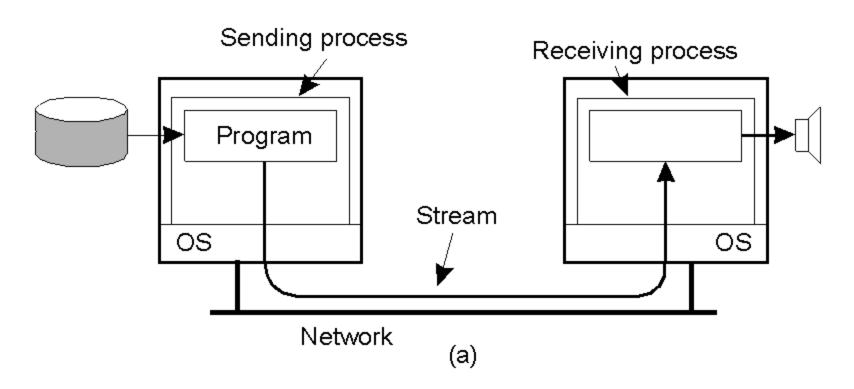
Message Brokers





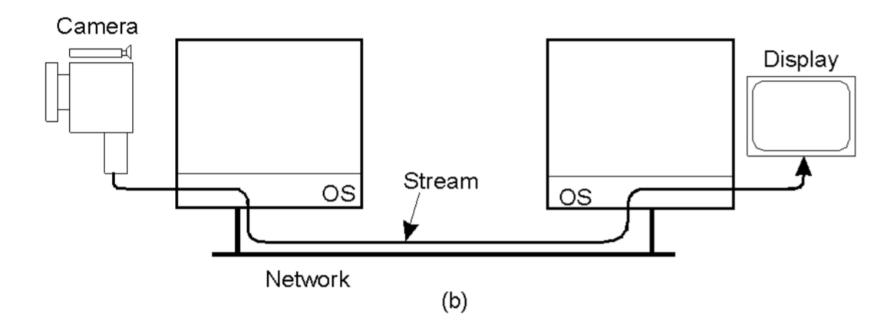
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Setting up a stream between two processes across a network.



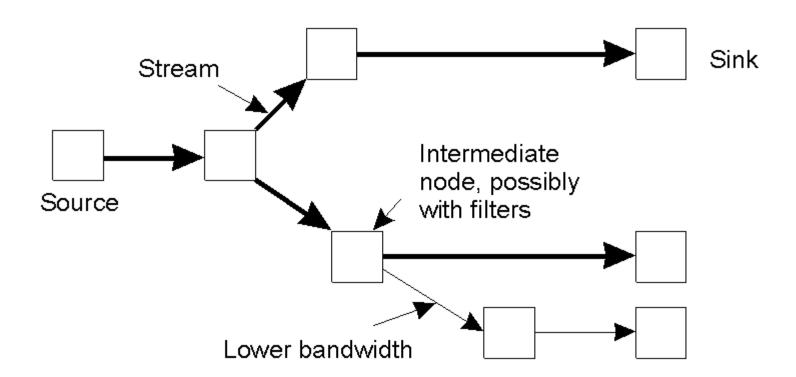
Data Stream (2)

Setting up a stream directly between two devices.



Data Stream (3)

An example of multicasting a stream to several receivers.



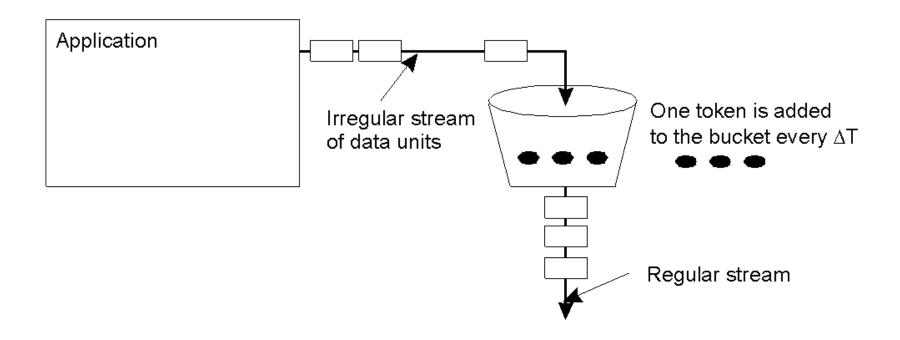
Specifying QoS (1)

A flow specification.

Characteristics of the Input	Service Required
□maximum data unit size (bytes) □Token bucket rate (bytes/sec)	□Loss sensitivity (bytes) □Loss interval (µsec)
□Toke bucket size (bytes) □Maximum transmission rate (bytes/sec)	 □Burst loss sensitivity (data units) □Minimum delay noticed (μsec) □Maximum delay variation (μsec)
	□Quality of guarantee

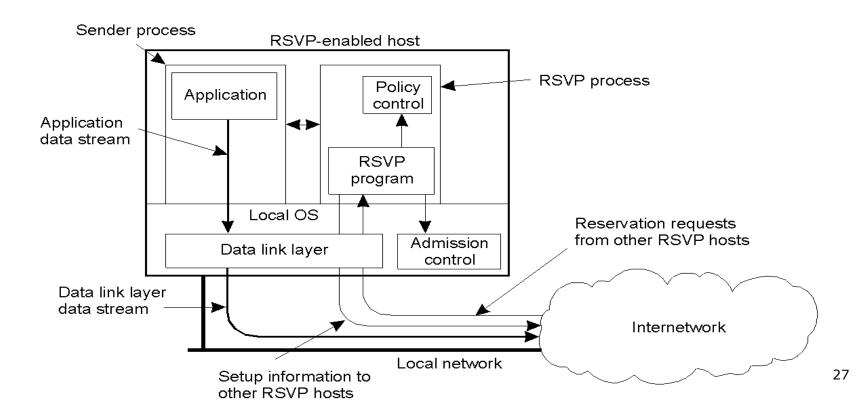
Specifying QoS (2)

□ The principle of a token bucket algorithm.



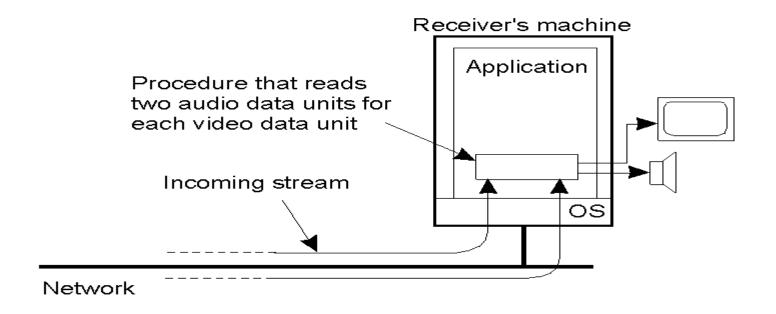
Setting Up a Stream

The basic organization of RSVP for resource reservation in a distributed system.



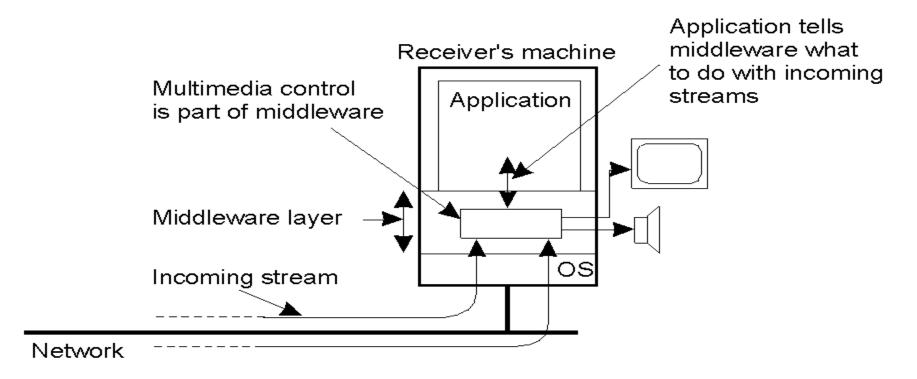
Synchronization Mechanisms (1)

The principle of explicit synchronization on the level data units.



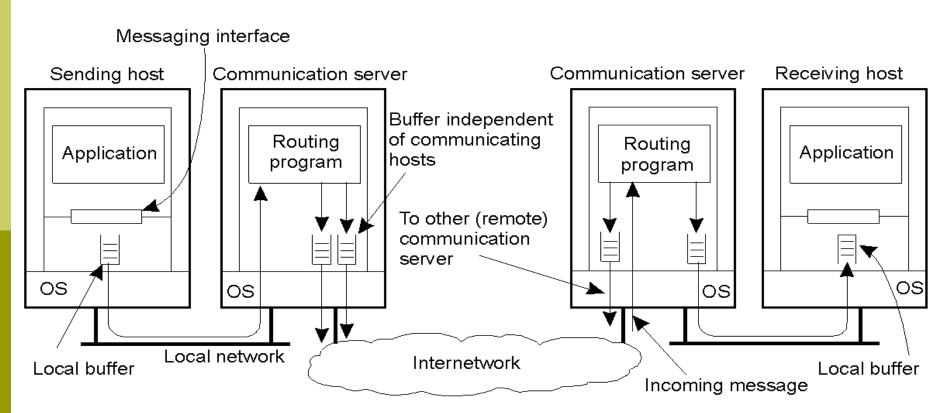
Synchronization Mechanisms (2)

The principle of synchronization as supported by high-level interfaces.



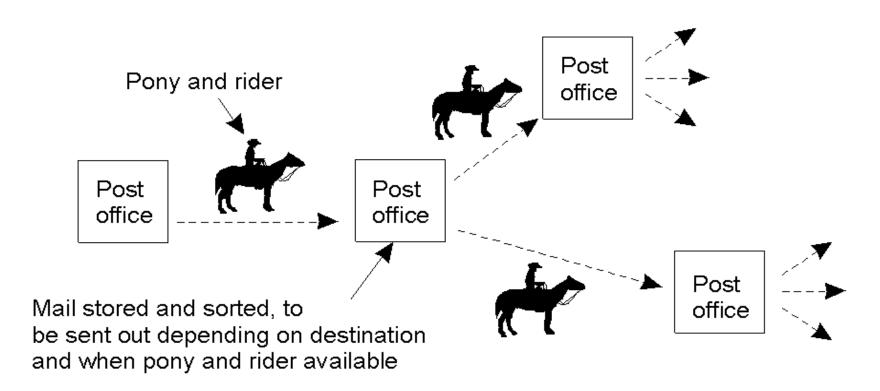
Persistence and Synchronicity in Communication (1)

General organization of a communication system in which hosts are connected through a network



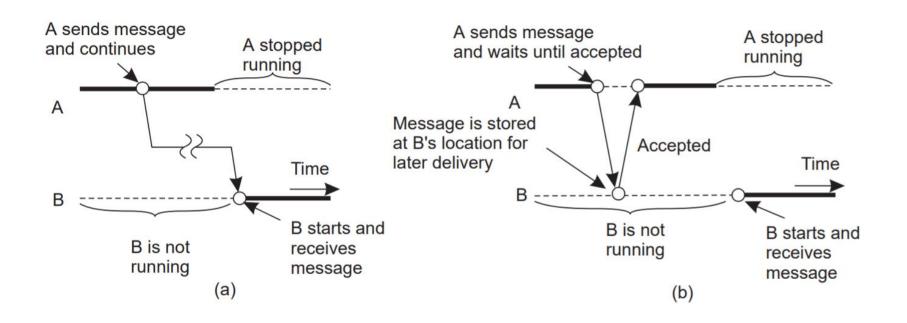
Persistence and Synchronicity in Communication (2)

Persistent communication of letters back in the days of the Pony Express.

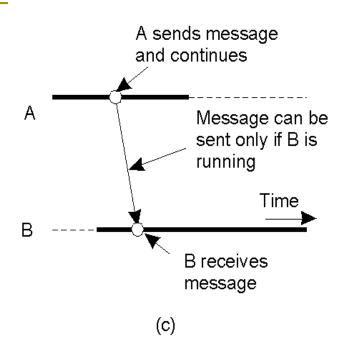


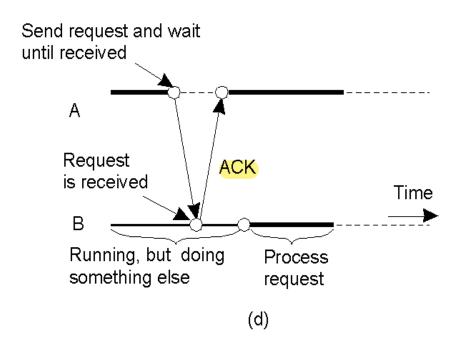
Persistence and Synchronicity in Communication (3)

- a) Persistent asynchronous communication
- b) Persistent synchronous communication



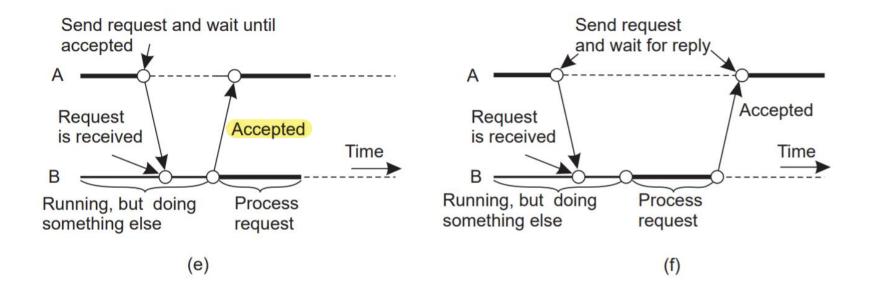
Persistence and Synchronicity in Communication (4)





- c) Transient asynchronous communication
- d) Receipt-based transient synchronous communication

Persistence and Synchronicity in Communication (5)



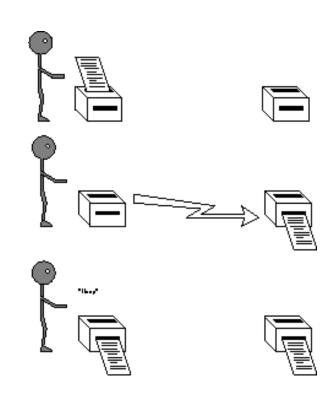
- e) Delivery-based transient synchronous communication at message delivery
- f) Response-based transient synchronous communication

Point-to-point Communication

- Always involves exactly two processes
- Example
 - *MPI_send* (buf, count, datatype, dest, tag, comm)
 - dest is identified by its rank within the communicator
 - comm identifies a group of processes and a communication context
 - MPI_Recv (buf, count, datatype, source, tag, comm, status)
- Several variations on how the sending of a message can interact with a program

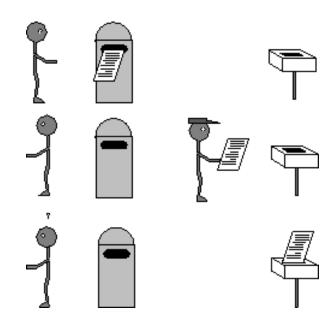
Synchronous

- A synchronous communication does not complete until the message has been received
 - A FAX or registered mail



Asynchronous

- An asynchronous communication completes as soon as the message is on the way.
 - A post card or email



Blocking and Non-blocking

- Blocking operations only return when the operation has been completed
 - Normal FAX machines
- Non-blocking operations return right away and allow the program to do other work
 - Receiving a FAX

May 23, 2018 38

Group Communication

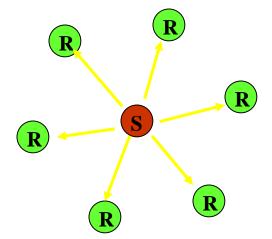
- RPC communication involves two processes (point-to-point)
- Can RPC handle one-to-many?

E.g. a client has to send message to all

servers.



Point-to-point



Group Communications

- Characteristics:
 - Dynamics create/destroy group, process join or leave and can be member to more than one group
- Multicast sends a single message from one process to each of the members of a group of process.
- Broadcast message delivered to all machines

Collective Communications

- Point-to-point communications involve pairs of processes.
- Many message passing systems provide operations which allow larger numbers of processes to participate

Thank you