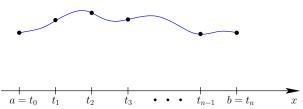
Arne Morten Kvarving

Department of Mathematical Sciences Norwegian University of Science and Technology

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- We are given a "large" dataset, i.e. a function sampled in many points.
- We want to find an approximation in-between these points.
- Until now we have seen one way to do this, namely high order interpolation - we express the solution over the whole domain as one polynomial of degree N for N + 1 data points.



• Let us consider the function

$$f(x) = \frac{1}{1+x^2}.$$

Known as Runge's example.

 While what we illustrate with this function is valid in general, this particular function is constructed to really amplify the problem.

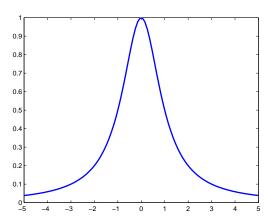


Figure: Runge's example plotted on a grid with 100 equidistantly spaced grid points.

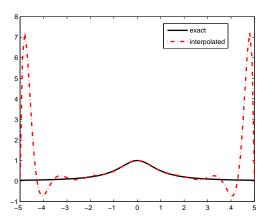
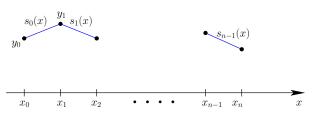


Figure: Runge's example interpolated using a 15th order polynomial based on equidistant sample points.

- It turns out that high order interpolation using a global polynomial often exhibit these oscillations hence it is "dangerous" to use (in particular on equidistant grids).
- Another strategy is to use piecewise interpolation. For instance, piecewise linear interpolation.



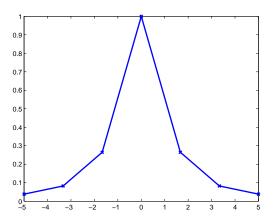
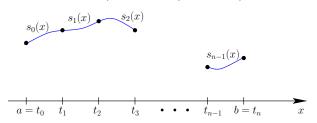


Figure: Runge's example interpolated using piecewise linear interpolation. We have used 7 points to interpolate the function in order to ensure that we can actually see the discontinuities on the plot.

A better strategy - spline interpolation

- We would like to avoid the Runge phenomenon for large datasets ⇒ we cannot do higher order interpolation.
- The solution to this is using piecewise polynomial interpolation.
- However piecewise linear is not a good choice as the regularity of the solution is only C^0 .
- These desires lead to splines and spline interpolation.



Splines - definition

A function S(x) is a spline of degree k on [a, b] if

- $S \in C^{k-1}[a, b]$.
- $a = t_0 < t_1 < \cdots < t_n = b$ and

$$S(x) = \begin{cases} S_0(x), & t_0 \le x \le t_1 \\ S_1(x), & t_1 \le x \le t_2 \\ \vdots & \vdots \\ S_{n-1}(x), & t_{n-1} \le x \le t_n \end{cases}$$

where $S_i(x) \in \mathbb{P}^k$.

Cubic spline

$$S(x) = \begin{cases} S_0(x) = & a_0x^3 + b_0x^2 + c_0x + d_0, & t_0 \le x \le t_1 \\ \vdots \\ S_{n-1}(x) = & a_{n-1}x^3 + b_{n-1}x^2 + c_{n-1}x + d_{n-1}, \ t_{n-1} \le x \le t_n. \end{cases}$$

which satisfies

$$S(x) \in C^{2}[t_{0}, t_{n}]: S'_{i-1}(x_{i}) = S'_{i}(x_{i}) S''_{i-1}(x_{i}) = S''_{i}(x_{i}) S''_{i-1}(x_{i}) = S''_{i}(x_{i})$$
, $i = 1, 2, \dots, n-1$.

Cubic spline - interpolation

Given $(x_i, y_i)_{i=0}^n$. Task: Find S(x) such that it is a cubic spline interpolant.

- The requirement that it is to be a cubic spline gives us 3(n-1) equations.
- In addition we require that

$$S(x_i) = y_i, \qquad i = 0, \cdots, n$$

which gives n+1 equations.

- This means we have 4n-2 equations in total.
- We have 4n degrees of freedom $(a_i, b_i, c_i, d_i)_{i=0}^{n-1}$.
- Thus we have 2 degrees of freedom left.

Cubic spline - interpolation

We can use these to define different subtypes of cubic splines:

- $S''(t_0) = S''(t_n) = 0$ natural cubic spline.
- $S'(t_0), S'(t_n)$ given clamped cubic spline.

•

$$S_{0}'''\left(t_{1}
ight)=S_{1}'''\left(t_{1}
ight) \\ S_{n-2}''\left(t_{n-1}
ight)=S_{n-1}'''\left(t_{n-1}
ight)
ight\}$$
 - Not a knot condition (MATLAB)

Task: Find S(x) such that it is a natural cubic spline.

- Let $t_i = x_i, i = 0, \dots, n$.
- Let $z_i = S''(x_i)$, $i = 0, \dots, n$. This means the condition that it is a natural cubic spline is simply expressed as $z_0 = z_n = 0$.
- Now, since S(x) is a third order polynomial we know that S''(x) is a linear spline which interpolates (t_i, z_i) .
- Hence one strategy is to first construct the linear spline interpolant S''(x), and then integrate that twice to obtain S(x).

The linear spline is simply expressed as

$$S_i''(x) = z_i \frac{x - t_{i+1}}{t_i - t_{i+1}} + z_{i+1} \frac{x - t_i}{t_{i+1} - t_i}.$$

• We introduce $h_i = t_{i+1} - t_i, i = 0, \dots, n$ which leads to

$$S''(x) = z_{i+1} \frac{x - t_i}{h_i} + z_i \frac{t_{i+1} - x}{h_i}.$$

We now integrate twice

$$S_i(x) = \frac{z_{i+1}}{6h_i} (x - t_i)^3 + \frac{z_i}{6h_i} (t_{i+1} - x)^3 + C_i (x - t_i) + D_i (t_{i+1} - x).$$

• Interpolation gives:

$$S_i(t_i) = y_i \Rightarrow \frac{z_i}{6}h_i^2 + D_ih_i = y_i, i = 0, \cdots, n.$$

• Continuity yields:

$$S_i(t_{i+1}) = y_{i+1} \Rightarrow \frac{z_{i+1}}{6}h_i^2 + C_ih_i = y_{i+1}.$$

 We insert these expressions to find the following form of the system

$$S_{i}(x) = \frac{z_{i+1}}{6h_{i}} (x - t_{i})^{3} + \frac{z_{i}}{6h_{i}} (t_{i+1} - x)^{3} + \left(\frac{y_{i+1}}{h_{i}} - \frac{z_{i+1}}{6}h_{i}\right) (x - t_{i}) + \left(\frac{y_{i}}{h_{i}} - \frac{h_{i}}{6}z_{i}\right) (t_{i+1} - x).$$

• We then take the derivative.

The derivative reads

$$S'_{i}(x) = \frac{z_{i+1}}{2h_{i}} (x - t_{i})^{2} - \frac{z_{i}}{2h_{i}} (t_{i+1} - x)^{2} + \underbrace{\frac{1}{h_{i}} (y_{i+1} - y_{i})}_{b_{i}} - \frac{h_{i}}{6} (z_{i+1} - z_{i}).$$

• In our abscissas this gives

$$S'_{i}(t_{i}) = -\frac{1}{2}z_{i}h_{i} + b_{i} - \frac{h_{i}}{6}z_{i+1} + \frac{1}{6}h_{i}z_{i}$$

$$S'_{i}(t_{i+1}) = \frac{z_{i+1}}{2}h_{i} + b_{i} - \frac{h_{i}}{6}z_{i+1} + \frac{1}{6}h_{i}z_{i}$$

$$S_{i-1}(t_{i}) = \frac{1}{3}z_{i}h_{i+1} + \frac{1}{6}h_{i-1}z_{i-1} + b_{i-1}$$

$$S'_{i}(t_{i}) = S_{i-1}(t_{i}) \Rightarrow$$

$$6(b_{i} - b_{i-1}) = h_{i-1}z_{-1} + 2(h_{i-1} + h_{i})z_{i} + h_{i}z_{i+1}.$$

Natural cubic splines - algorithm

This means that we can find our solution using the following procedure:

• First do some precalculations

$$h_{i} = t_{i+1} - t_{i}, i = 0, \dots, n-1$$

$$b_{i} = \frac{1}{h_{i}} (y_{i+1} - y_{i}), i = 0, \dots, n-1$$

$$v_{i} = 2 (h_{i-1} + h_{i}), i = 1, \dots, n-1$$

$$u_{i} = 6 (b_{i} - b_{i-1}), i = 1, \dots, n-1$$

$$z_{0} = z_{n} = 0$$

Natural cubic splines - algorithm

• Then solve the tridiagonal system

Natural cubic splines - example

Given the dataset.

i	0	1	2	3
X _i	0.9	1.3	1.9	2.1
Уi	1.3	1.5	1.85	2.1
$h_i = x_{i+1} - x_i$	0.4	0.6	0.2	
$b_i = \frac{1}{h_i} \left(y_{i+1} - y_i \right)$	0.5	0.5833	1.25	
$v_i = 2(h_{i-1} + h_i)$		2.0	1.6	
$u_i = 6\left(b_i - b_{i-1}\right)$		0.5	4	

• The linear system reads

$$\begin{bmatrix} 2.0 & 0.4 \\ 0.4 & 1.6 \end{bmatrix} \begin{bmatrix} z_1 \\ z_2 \end{bmatrix} = \begin{bmatrix} 0.5 \\ 0.4 \end{bmatrix}$$

Natural cubic splines - example

• We find $z_0 = 0.5$, $z_1 = 0.125$. This gives us our spline functions

$$S_0(x) = 0.208 (x - 0.9)^3 + 3.78 (x - 0.9) + 3.25 (1.3 - x)$$

$$S_1(x) = 0.035 (x - 1.3)^3 + 0.139 (1.9 - x)^3 + 0.664 - 0.62x$$

$$S_2(x) = 0.104 (x - 1.9)^3 + 10.5 (x - 1.9) + 9.25 (2.1 - x)$$

· Assume we are given a general tridiagonal system

$$\begin{bmatrix} d_1 & c_1 & & & & \\ a_1 & d_2 & c_2 & & & \\ & \ddots & \ddots & \ddots & \\ & & & c_{n-1} \\ & & & a_{n-1} & d_n \end{bmatrix}, \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_{n-1} \\ b_n \end{bmatrix}.$$

First elimination (second row) yields

$$\begin{bmatrix} d_1 & c_1 & & & & \\ 0 & \tilde{d}_2 & c_2 & & & \\ & \ddots & \ddots & \ddots & \\ & & & c_{n-1} \\ & & & a_{n-1} & d_n \end{bmatrix}, \begin{bmatrix} b_1 \\ \tilde{b}_2 \\ \vdots \\ b_{n-1} \\ b_n \end{bmatrix}, \qquad \tilde{d}_2 = d_2 - \frac{a_1}{d_1} c_1 \\ \tilde{b}_2 = b_2 - \frac{a_1}{d_1} b_1$$

This means that the elimination stage is

for
$$i=2,\cdots,n$$

$$m=a_{i-1}/d_{i-1}$$

$$\tilde{d}_i=d_i-mc_{i-1}$$

$$\tilde{b}_i=b_i-mb_{i-1}$$
 end

And the backward substitution reads

$$x_n = ilde{b}_n/d_n$$
 for $i=n-1,\cdots,1$ $x_i = \left(ilde{b}_i - c_i x_{i+1}
ight)/ ilde{d}_i$ end

where $\tilde{b}_1 = b_1$.



- This will work out fine as long as $\tilde{d}_i \neq 0$.
- Assume that $|d_i| > |a_{i-1}| + |c_i|$ i.e. diagonal dominance.
- For the eliminated system diagonal domiance means that

$$|\tilde{d}_i| < |c_i|.$$

 We now want to show that diagonal dominance of the original system implies that the eliminated system is also diagonal dominant.

• We now assume that $|\tilde{d}_{i-1}| > |c_{i-1}|$. This is obviously satisfied for $\tilde{d}_1 = d_1$.

$$|\tilde{d}_{i}| = |d_{i} - \frac{a_{i-1}}{\tilde{d}_{i-1}}c_{i-1}| \ge |d_{i}| - \frac{|a_{i-1}|}{|\tilde{d}_{i-1}|}|c_{i-1}|$$

$$> |a_{i-1} - |c_{i}| - |a_{i-1}| = |c_{i}|.$$

• Hence the diagonal domiance is preserved which means that $\tilde{d}_i \neq 0$. The algorithm produces a unique solution.

Why cubic splines?

- Now to motivate why we use cubic splines.
- First, let us introduce a measure for the smoothness of a function:

$$\mu(f) = \int_{a}^{b} (f''(x))^{2} dx.$$
 (1)

• We then have the following theorem

Theorem

Given interpolation data $(t_i, y_i)_{i=0}^n$. Among all functions $f \in C^2[a, b]$ which interpolates (t_i, y_i) , the natural cubic spline is the smoothest, where smoothness is measured through (1).

Why cubic splines?

• We need to prove that

$$\mu(f) \ge \mu(S) \,\forall \, f \in C^2[a,b].$$

Introduce

$$g(x) = S(x) - f(x),$$
 $g(x) \in C^{2}[a, b]$
 $g(t_{i}) = 0, i = 0, \dots, n.$

• Inserting this yields

$$\mu(f) = \int_{a}^{b} (S''(x) - g''(x))^{2} dx$$
$$= \mu(S) + \mu(g) - 2 \int_{a}^{b} S''(x)g''(x) dx$$

Now since $\mu(g) > 0$, we have proved our result if we can show that

$$\int_a^b S''(x)g''(x)\,\mathrm{d}x=0.$$



Why cubic splines?

We have that

$$\int_a^b S''(x)g''(x) \, \mathrm{d} x = \left. g'(x)S''(x) \right|_a^b - \int_a^b g'(x)S'''(x) \, \mathrm{d} x$$

First part on the right hand side is zero since $z_0 = z_n = 0$. Second part we split in an integral over each subdomain

$$-\int_{a}^{b} g'(x)S'''(x) dx = -\sum_{i=0}^{n-1} \int_{t^{i}}^{t^{i+1}} g'(x)S'''(x) dx$$
$$= -\sum_{i=0}^{n-1} 6a_{i} \int_{t_{i}}^{t_{i+1}} g'(x) dx$$
$$= -\sum_{i=0}^{n-1} 6a_{i} g(x)|_{t_{i}}^{t_{i+1}} = 0.$$

Cubic spline result

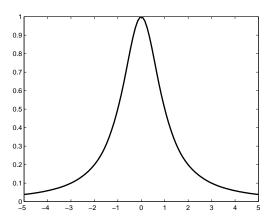


Figure: Runge's example interpolated using cubic spline interpolation based on 15 equidistant samples.