

Homework 8

Problem 1. Show that, for constant $p \in (0, 1)$, almost no graph in $\mathcal{G}(n, p)$ has a separating complete subgraph.

Solution. This is a simple application of the ‘almost always true of property $P_{i,j}$ ’.

Detail:

- **Separating subgraph:** Given $G = (V, E)$, and some $X \subseteq V \cup E$, we call X a separating subgraph if there exists two vertices $u, v \in V(G - X)$ such that u, v are in the same component of G , while u, v lie in two disconnected components of $G - X$ (i.e., X separates u and v).
- **Separating complete subgraph:** If the above subgraph X is also a complete graph.

Now consider a graph $G = (V, E)$ with property $\mathcal{P}_{2,1}$. We claim that a graph with property $\mathcal{P}_{2,1}$ has the following property: For any pair of vertices $u, v \in G$, there exists a pair of vertices w_1, w_2 such that

$$(w_1, u) \in E, \quad (w_1, v) \in E$$

$$(w_2, u) \in E, \quad (w_2, v) \in E$$

$$(w_1, w_2) \notin E.$$

To prove the claim: consider vertices u, v and an arbitrary vertex x . By property $\mathcal{P}_{2,1}$, there exists a vertex w_1 which is neighbor to u and v , but not to x . Now using property $\mathcal{P}_{2,1}$ again (with x replaced by w_1) it follows that there exists a vertex w_2 which is neighbor to u and v , but not to w_1 . Thus the claim holds.

Finally, consider a complete subgraph $H \subset G$ and two arbitrary vertices u and v in $G - V(H)$. By the claim above, there are two non-adjacent vertices w_1 and w_2 in G which are both neighbors of both u and v . Since H is complete, it follows that w_1 and w_2 cannot both belong to H , therefore remove H will not separate u and v . In another word, H does not separate G . The statement now follows since almost all graphs in $\mathcal{G}(n, p)$ have property $\mathcal{P}_{2,1}$ for any constant $p \in (0, 1)$. □

Problem 2. What is the expected number of trees with k vertices in $G \in \mathcal{G}(n, p)$?

Solution. By Cayley's formula and the linearity of expectation, it is $\binom{n}{k} k^{k-2} p^{k-1}$ \square

Problem 3. Show that if almost all $G \in \mathcal{G}(n, p)$ have a graph property \mathcal{P}_1 and almost all $G \in \mathcal{G}(n, p)$ have a graph property \mathcal{P}_2 , then almost all $G \in \mathcal{G}(n, p)$ have both properties.

Solution. The portion of the graphs have both properties equals 1 minus the portion of the graphs which does not have property \mathcal{P}_1 or \mathcal{P}_2 . However the portion of the graph does not have property \mathcal{P}_1 or \mathcal{P}_2 is bounded by the sum of the portion of the graphs does not have property \mathcal{P}_1 and the the portion of the graphs does not have property \mathcal{P}_2 , which both tend to 0 as n approaches ∞ . The claim in the question then follows. \square

Problem 4. (Optional)

1. Prove that the threshold for the existence of cycles in $\mathcal{G}(n, p)$ is $p = \frac{1}{n}$.
2. Search the World Wide Web to find some real world graphs in machine readable form or data bases that could automatically be converted to graphs.
 - (a) Plot the degree distribution of each graph.
 - (b) Compute the average degree of each graph.
 - (c) Count the number of connected components of each size in each graph.
 - (d) Describe what you find.
3. Create a simulation (an animation) to show the evolution of the $\mathcal{G}(n, p)$ (Erdős-Rényi) random graph as its density p is gradually increased. Observe the phase transitions for trees of increasing orders, followed by the emergence of the giant component, etc.