

# CS 1550

Week 4 – Synchronization with xv6

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#### CS 1550 – Announcement

- Project 1
  - put struct cs1550\_sem declaration inside a header file names sem.h
  - add **#include** "**sem.h**" into the test cases
  - The sem.h file should be in the same folder as the test case file when compiling
  - Two more test cases are posted:
    - trafficsim-mutex.c
      - gcc -m32 -lm -o trafficsim-mutex
    - trafficsim-strict-order.c

## Keep in mind the different qemu

- qemu with xv6 (Labs) Refer to Lab 1 if needed!
- qemu-x86 (Project 1)







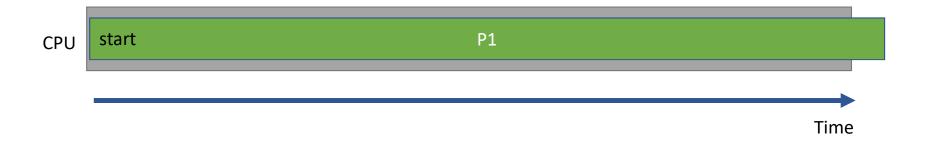
• OS chooses another processes to execute once the first finishes



• OS chooses another processes to execute once the first finishes

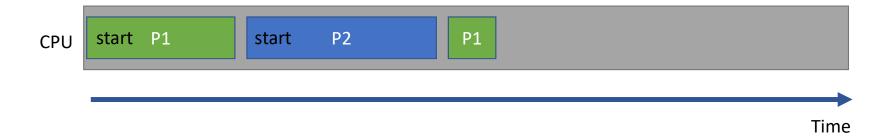


• What if P1 is a big process?



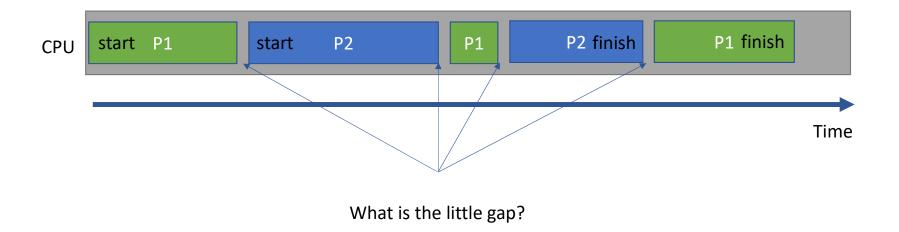


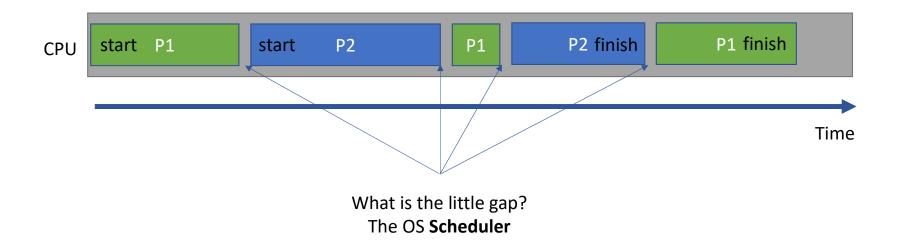








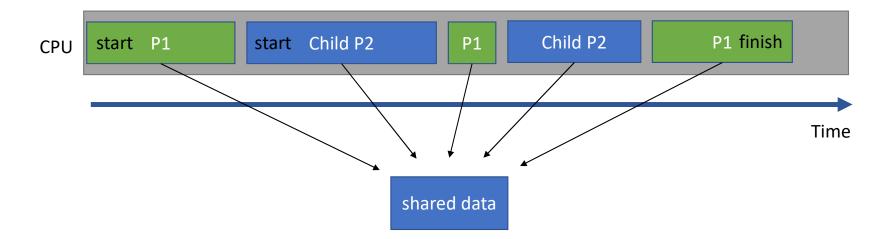




• What happens in Parent-Child Process scenario?



- What happens in Parent-Child Process scenario?
- How to keep integrity/correctness on race conditions?



```
struct list {
  int data;
  struct list *next;
};
```

```
struct list {
  int data;
  struct list *next;
};
struct list *list = 0;
```

```
struct list {
   int data;
   struct list *next;
};

struct list *list = 0;

void
insert(int data) {
    struct list *l;
    l = malloc(sizeof *l);
    l->data = data;
    l->next = list;
    list = l;
}
```

```
struct list {
         int data;
         struct list *next;
       };
                                               CPU
                                                      P1
       struct list *list = 0;
       void
       insert(int data) {
              struct list *1;
                                               P1 stops here the
                                               OS switches to P2
              1 = malloc(sizeof *1);
P1 stopped
              1->data = data;
              1->next = list;
              list = 1;
```

```
struct list {
         int data;
         struct list *next;
      };
                                                           CP2
                                            CPU
                                                   P1
      struct list *list = 0;
      void
       insert(int data) {
             struct list *1;
             l = malloc(sizeof *1);
             l->data = data;
P1 stopped
             1->next = list;
             list = 1;
```

CP2

P1

CPU

P1

```
struct list {
   int data;
   struct list *next;
};

struct list *list = 0;

void
   insert(int data) {
      struct list *l;
      l = malloc(sizeof *l);
      l->data = data;

CP2 stopped l->next = list;
      list = l;
}
```

- Sharing CPU among processes
- Ensuring data integrity/correctness
- Ensure that a critical section of your code is only executed by one process

```
struct list *list = 0;
struct lock listlock;

void
insert(int data)
{
    struct list *l;

    acquire(&listlock);
    l = malloc(sizeof *l);
    l->data = data;
    l->next = list;
    list = l;
    release(&listlock);
}
```



```
struct list *list = 0;
struct lock listlock;

void
insert(int data)
{
    struct list *l;

    acquire(&listlock);
    l = malloc(sizeof *l);
    l->data = data;
    l->next = list;
    list = l;
    release(&listlock);
}
```



```
struct list *list = 0;
struct lock listlock;
void
insert(int data)
                                       CPU
                                                       CP2
                                              P1
                                                                P1
      struct list *1;
      acquire(&listlock);
      1 = malloc(sizeof *1);
                                              P1 gets locks the lock
      1->data = data;
      1->next = list;
      list = 1;
      release(&listlock);
```

```
struct list *list = 0;
struct lock listlock;
void
insert(int data)
                                       CPU
                                                       CP2
                                              P1
                                                                P1
      struct list *1;
      acquire(&listlock);
      l = malloc(sizeof *1);
                                              P1 gets locks the lock
      l->data = data;✓
      1->next = list;
      list = 1;
      release(&listlock);
```

```
struct list *list = 0;
   struct lock listlock;
  void
   insert(int data)
                                          CPU
                                                          CP2
                                                 P1
                                                                   P1
         struct list *1; *
         acquire(&listlock);
         l = malloc(sizeof *l);
                                               When the OS schedule CP2
P1 stopped
        1->data = data;
         1->next = list;
         list = 1;
         release(&listlock);
```

```
struct list *list = 0;
struct lock listlock;

void
insert(int data)
{
    struct list *l;

    acquire(&listlock);
    l = malloc(sizeof *l);

P1 stopped l->data = data;
    l->next = list;
    list = l;
    release(&listlock);
}
```



It will try to get the lock but won't.

```
struct list *list = 0;
struct lock listlock;

void
insert(int data)
{
    struct list *l;

    acquire(&listlock);
    l = malloc(sizeof *l);

P1stopped l->data = data;
    l->next = list;
    list = l;
    release(&listlock);
}
```



It will try to get the lock but won't.

It will be constantly try to get it (in a loop). Until the OS switches back to P1

```
struct list *list = 0;
    struct lock listlock;
    void
    insert(int data)
                                               CPU
                                                                CP2
                                                      P1
                                                                         P1
           struct list *1;
           acquire(&listlock);
CP2 stopped
           1 = malloc(sizeof *1);
           1->data = data;
                                                       P1 release the lock P2 will finally be
           1->next = list;
                                                       able to execute, once scheduled
           list = 1;
           release(&listlock);
```

```
struct list *list = 0;
    struct lock listlock;
    void
    insert(int data)
                                               CPU
                                                                CP2
                                                      P1
                                                                         P1
                                                                                CP2
           struct list *1;
CP2 proceeds acquire(&listlock);
           1 = malloc(sizeof *1);
           1->data = data;
                                                       P1 release the lock P2 will finally be
           1->next = list;
                                                       able to execute, once scheduled
           list = 1;
           release(&listlock);
```

#### SpinLock

```
Void
acquire (struct spinlock *lk)
{
    for(;;) {
        if(!lk->locked) {
            lk->locked = 1;
            break;
        }
    }
}
```

- Keep spinning until find lock is released
- But we can have the same issue as before
- We need to check and lock atomically

- Xv6 relies on a special 386 hardware instruction, xchg
- Atomically check and change a register value
  - xchg(&lk->locked, 1)

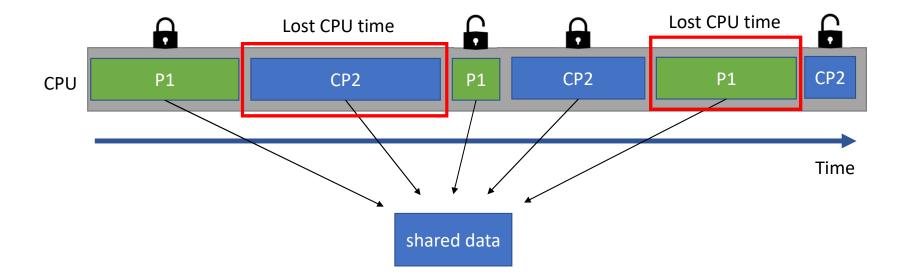
- Swap a word in memory with the contents of a register
- In acquire function:
  - Repeats xchg instruction in a loop
  - Each round atomically read lock and set the lock to 1

```
void
acquire(struct spinlock *lk)
{
    pushcli(); // disable interrupts to
avoid deadlock.
    ...

// The xchg is atomic.
    while(xchg(&lk->locked, 1) != 0);
    ...

// Record info about lock acquisition for debugging.
    lk->cpu = mycpu();
    getcallerpcs(&lk, lk->pcs);
}
```

- But the we have another issue
  - Busy waiting



- SpinLock
  - Busy waiting
  - Useful for short critical sections
    - E.g. increment a counter, access an array element, etc.
  - Not useful, when the period of wait is unpredictable or will take a long time
    - E.g. read page from disk

- Sleep Locks
  - For code need to hold a lock for a long time (read/write to disk)
- Avoids the schedule of "spin locked" processes

- Sleep Locks
  - For code need to hold a lock for a long time (read/write to disk)
- Avoids the schedule of "spin locked" processes

```
void
acquiresleep(struct sleeplock *lk)
{
    acquire(&lk->lk);
    while (lk->locked) {
        sleep(lk, &lk->lk);
    }
    lk->locked = 1;
    lk->pid = myproc()->pid;
    release(&lk->lk);
}

release(&lk->lk);
}
```

- Sleep Locks
  - For code need to hold a lock for a long time (read/write to disk)
- Avoids the schedule of "spin locked" processes

```
void
acquiresleep(struct sleeplock *lk)
{
    acquire(&lk->lk);
    while (lk->locked) {
        sleep(lk, &lk->lk);
        lk->locked = 0;
        lk->pid = 0;
        wakeup(lk);
        lk->pid = myproc()->pid;
        release(&lk->lk);
}

    release(&lk->lk);
}
```

- Put one process to sleep waiting for event
- Mark current process as sleeping
- Call sched() to release the processor

```
void
sleep(void *chan, struct spinlock *lk)
{
   struct proc *p = myproc();
    ...
   p->state = SLEEPING;
   sched();
   ...
}
```

```
void
                                            sleep(void *chan, struct spinlock *lk)
                                              struct proc *p = myproc();
                                              if(p == 0)
            Sanity checks
                                                panic("sleep");

    Must be a current process

    Must have been passed a lock

                                              if(lk == 0)
                                                panic("sleep without lk");
                                              if(lk != &ptable.lock) {
                                                acquire(&ptable.lock);
Hold the ptable.lock, it is safe to release lk
                                              \rightarrow release(lk);
                                              p->chan = chan;
                                              p->state = SLEEPING;
                                              sched();
                                              p->chan = 0
                                              if(lk != &ptable.lock) {
                                                release (&ptable.lock);
                                                acquire(lk);
```

- Wake up process when event happened
- Mark a waiting process as runnable

```
static void
wakeup(void *chan)
{
    acquire(&ptable.lock);
    wakeup1(chan);
    release(&ptable.lock);
}
```

- Who needs to be a syscall?
  - SpinLocks
  - SleepLocks