

ER-to-Relational Mapping

- There is almost a one-to-one correspondence between the ER constructs and the relational ones.
- The two major distinctions are:
 - In a relational schema, relationships are represented *implicitly* through primary and foreign keys of participating entities.
 - In a relational schema, columns of relations cannot be multi-valued or composite.
 - Composite attributes are replaced with their simple component ones, and
 - multi-valued attributes are stored in a separate relation.

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Mapping Algorithm

- 1. For each strong entity type E.
 - Create a new table.
 - Include as its columns, all the simple attributes and simple components of the composite attributes of E.
 - Identify the primary key and the alternate keys.

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Mapping Algorithm...Weak Entities

- 2. For each weak entity W that is associated with only one 1:1 identifying owner relationship.
 - Identify the table T of the owner entity type.
 - Include as columns of T, all the simple attributes and simple components of the composite attributes of W.
- 3. For each weak entity W that is associated with a 1:N or M:N identifying relationship, or participates in more than one relationship.
 - Create a new table T.
 - Include as its columns, all the simple attributes and simple components of the composite attributes of W.
 - Form its primary key by including as a foreign key in R, the primary key of its owner entity.

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Mapping Algorithm...Binary Relationships

- 4. For each binary 1:1 relationship type R.
 - Identify the tables S and T of the participating entity types.
 - Choose S (preferably the one with total participation).
 - Include as foreign key in S, the primary key of T.
 - Include as Columns of S, all the simple attributes and simple components of the composite attributes of R.
- 5. For each binary 1:N relationship type R.
 - Identify the table S (at the *N-side*) and T of the participating entities.
 - Include as a foreign key in S, the primary key of T.
 - Include as columns of S, all the simple attributes and simple components of composite attributes of R.

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Mapping Algorithm...Any Relationship

- 6. For each N-ary relationship (including binary N:M relationship) type R.
 - Create a new table T.
 - Include as columns of T, all the simple attributes and simple components of composite attributes of R.
 - Include as foreign keys, the primary keys of the participating (strong or weak) entity types.
 - Specify as the primary key of T, the list of foreign keys.

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Mapping Algorithm...Multi-valued Attributes

- 7. For each multi-valued attribute A.
 - Create a new table T.
 - Include as columns of T, the simple attribute or simple components of the attribute A.
 - Include as foreign key, the primary key of the entity or relationship type that has A.
 - Specify as the primary key of T, the foreign key and the columns corresponding to A.

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EER Mapping Algorithm

- 8. For each specialization with disjoint subclasses.
 - Create a new table T_i for each subclass S_i.
 - Include as columns of T_{i,} the simple attributes and simple component attributes of the superclass.
 - Include as columns of T_i, the simple attributes and simple component attributes specific to S_i.
 - Identify the primary key.

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EER Mapping Algorithm...

- 9. For each specialization with overlapping subclasses.
 - Create a new table O for the superclass.
 - Include as columns of O, the simple attributes and the simple component attributes of the superclass.
 - Identify its primary key and alternate keys.
 - Create a new table T_i for each subclass S_i
 - Include as columns of T_i, the simple attributes and simple component attributes specific to S_i.
 - Include as a foreign key in T_i (to be part of the primary key of T_i), the primary key of O.

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