

Integrity Constraints & Transactions in SQL

Queries and Transactions

- ❑ Queries: requests to the DBMS to retrieve data from the database
- ❑ Updates: requests to the DBMS to insert, delete or modify existing data
- ❑ Transactions: logical grouping of query and update requests to perform a task
 - Logical unit of work (like a function/subroutine)

Execution Abstraction

- ❑ A **transaction** is a **logical unit of work** in DBMSs
 - It is the execution of a **program segment** that performs some function or task by accessing shared data (e.g., a db)
 - logical grouping of query and update requests needed to perform a task
- ❑ Examples:
 - banking transaction
 - Deposit, withdraw, transfer \$
 - (airline reservation
 - reserve a seat on a flight
 - inventory transaction
 - Receive, Ship, Update

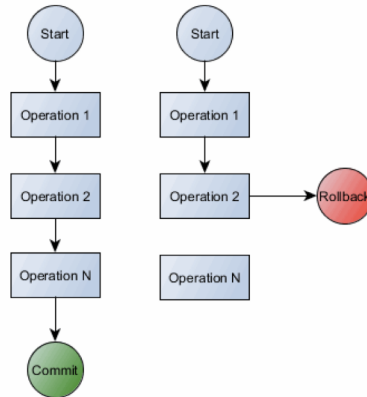


Transaction's ACID Properties

- ❑ **Atomicity** (alias failure atomicity)
Either all the operations associated with a transaction happen or none of them happens
- ❑ **Consistency Preservation**
A transaction is a correct program segment. It satisfies the integrity constraints on the database at the transaction's boundaries
- ❑ **Isolation** (alias concurrency atomicity / serializability)
Transactions are independent, the result of the execution of concurrent transactions is the same as if transactions were executed serially, one after the other
- ❑ **Durability** (alias persistence / permanence)
The effects of completed transactions become permanent surviving any subsequent failures

SQL TRANSACTIONS

- ❑ Start:
 - SET TRANSACTION
- ❑ COMMIT ;
- ❑ ROLLBACK;



SQL TRANSACTIONS

- ❑ Basic transaction statements:
 - SET TRANSACTION READ WRITE NAME <name>;
(SQL1: DECLARE TRANSACTION READ WRITE;)
 - SET TRANSACTION READ ONLY NAME <name>;
(SQL1: DECLARE TRANSACTION READ ONLY;)
 - COMMIT ;
 - ROLLBACK;
- ❑ ROLLBACK default action

Transaction Atomicity

- ❑ What do we expect with Atomicity?
 - “All or nothing”
- ❑ Consider a transaction:


```

set transaction read write name 'test';
insert into Student values (23, 'John', 'CS');
insert into Dept values ('CS', 501);
Commit;
            
```
- ❑ What happens if the first insert fails, e.g., due to a referential constraint violation?
 - Is the new tuple inserted into Department? No?
- ❑ If no error happens at commit time, the second insert is still committed!!!

Modes of Constraints Enforcement

- ❑ **NOT DEFERRABLE** or **IMMEDIATE**
 - Evaluation is performed at input time
 - By default constraints are created as NON DEFERRABLE
 - It *cannot* be changed during execution
- ❑ **DEFERRED**
 - Constraints are not evaluated until commit time
- ❑ **DEFERRABLE**
 - It can be changed within a transaction to be DEFERRED using SET CONSTRAINT
- ❑ Modes can be specified when a table is created.
 - INITIALLY IMMEDIATE: constraint validation to happen immediate
 - INITIALLY DEFERRED: constraint validation to defer until commit

Specifying Initial Eval. Mode in Tables

```
❑ CREATE TABLE SECTION
( SectNo sectno_dom,
  Name section_dom,
  HeadSSN ssn_dom,
  Budget budget_dom,
  CONSTRAINT section_PK
    PRIMARY KEY (SectNo) DEFERRABLE,
  CONSTRAINT section_FK
    FOREIGN KEY (HeadSSN) REFERENCES LIBRARIAN(SSN)
      INITIALLY DEFERRED DEFERRABLE,
  CONSTRAINT section_name_UN UNIQUE (Name)
    DEFERRABLE INITIALLY IMMEDIATE
);
```

Changing Constraint Evaluation Mode

- ❑ It is permitted only for deferrable constraints
- ❑ Setting the constraint validation mode within a transaction
 - set mode of all deferrable constraints

```
SET CONSTRAINT ALL IMMEDIATE;
```

```
SET CONSTRAINT ALL DEFERRED;
```
 - set mode of specific deferrable constraints (list)

```
SET CONSTRAINT section_budget_IC1 IMMEDIATE;
```

```
SET CONSTRAINT section_budget_IC1 DEFERRED;
```

Specifying Transaction Atomicity

- ❑ Errors at commit time: only when **deferred constraints** are violated
 - Constraints can be deferred if specified as **deferrable** in the table schema, and
 - deferred in the scope of the transaction
- ❑ E.g., *assume the constraints are deferrable*

```
set transaction read write name 'test';
set constraints all deferred;
insert into Student values (23, 'John', 'CS');
insert into Dept values ('CS', 501);
Commit;
```
- ❑ No constraint violation of the first insert is detected at *commit time* → the whole transaction is committed

Specifying Transaction Atomicity (2)

- ❑ E.g. 2, *assume the constraints are deferrable and assume SID 23 exists in that Database*

```
set transaction read write name 'test';
set constraints all deferred;
insert into Student values (23, 'John', 'CS');
insert into Dept values ('CS', 501);
Commit;
```
- ❑ The constraint violation of the first insert is detected at *commit time* → the whole transaction is rollback

Integrity Constraints in SQL

Structural Constraints

- ❑ Constraints (on Attributes):
 - NOT NULL
 - DEFAULT value
 - without the DEFAULT-clause, the default value is NULL
 - PRIMARY KEY (attribute-list)
 - UNIQUE (attribute list)
 - allows the specification of alternative key
 - FOREIGN KEY (key) REFERENCES table (key)

Referential Triggered Actions

- ❑ Actions if a Referential Integrity constraint is violated
 - SET NULL
 - CASCADE (propagate action)
 - SET DEFAULT
- ❑ Qualify actions by the triggering condition:
 - ON DELETE
 - ON UPDATE
- ❑ Note: Oracle does not support ON UPDATE & SET DEFAULT

Create Table with RI Trigger Actions

```
CREATE TABLE LIBRARIAN          /* or Micro_db.LIBRARIAN */
(
  Name    name_dom,
  SSN     ssn_dom,
  Section INTEGER,
  Address address_dom,
  Gender  gender_dom,
  Birthday DATE,
  Salary  DEC(8,2),

  CONSTRAINT librarian_PK PRIMARY KEY (SSN) DEFERRABLE,
  CONSTRAINT librarian_FK
    FOREIGN KEY (Section) REFERENCES SECTION (SNO)
    On Delete SET DEFAULT On Update CASCADE
    DEFERRABLE
);
```

Semantic Integrity Constraints

- ❑ A constraint is expressed as a *Predicate*, a condition similar to the one at the WHERE-clause of a query
- ❑ Three DDL constructs
 - Checks
 - Assertions
 - Triggers

Check Constraints

- ❑ CHECK *prohibits* an operation on a table that would violate the constraint. It is a *local* constraint.
- ❑ **CREATE TABLE** SECTION
(SectNo sectno_dom,
Name section_dom,
HeadSSN ssn_dom,
Budget budget_dom,
CONSTRAINT section_PK
PRIMARY KEY (SectNo) DEFERRABLE,
CONSTRAINT section_FK
FOREIGN KEY (HeadSSN) REFERENCES LIBRARIAN(SSN)
DEFERRABLE,

Check Constraints...

```
CONSTRAINT section_budget_IC1
CHECK ((Budget >= 0) AND (Budget IS NOT NULL))
DEFERRABLE,
CONSTRAINT section_budget_IC2
CHECK (NOT EXISTS
      (SELECT * FROM SECTION WHERE budget <
       (SELECT SUM (Salary) FROM LIBRARIAN)))
DEFERRABLE,
CONSTRAINT Head_Lib_IC3
CHECK (HeadSSN <> ALL (SELECT SSN FROM Retiree))
DEFERRABLE
);
```

Assertions

- ❑ Similar to CHECK but they are **global** constraints
CREATE OR REPLACE ASSERTION <assertion_name>
CHECK <Predicate> [Mode of Evaluation];
 - **Predicate** usually involves EXISTS and NOT EXISTS
- ❑ E.g., **CREATE OR REPLACE ASSERTION** budget_constraint
CHECK (NOT EXISTS
 (SELECT * FROM SECTION WHERE budget <
 (SELECT SUM (Salary) FROM LIBRARIAN));
❑ Dropping an assertion...
DROP ASSERTION budget_constraint;

VQuery
that
violates
IC

Triggers

- ❑ A trigger consists of 3 parts:
 1. Event(s),
 2. Condition, and
 3. Action
- ❑ E.g., Notify the Dean whenever the number of students in any major exceeds 1800

Triggers vs. Assertions

- ❑ Assertion
 - Condition must be true for each database state
 - DBMS rejects operations that violate such condition
- ❑ Trigger
 - DBMS takes a certain **action** when condition is true
 - Action could be: stored procedure, SQL statements, Rollback, etc.

My First Trigger

- Notify the Dean when the # of students in any major exceeds 1800

CREATE TRIGGER Major_Limit

Event(s)

Condition

Action

Example

- Notify the Dean when the # of students in any major exceeds 1800

CREATE TRIGGER Major_Limit

Event(s)

```
WHEN ( EXISTS (
    SELECT Major_Code, COUNT (*)
    FROM Student
    GROUP BY Major_Code
    HAVING COUNT (*) > 1800 ) )
```

Action

Example

- Notify the Dean when the # of students in any major exceeds 1800

CREATE TRIGGER Major_Limit

Event(s)

```
WHEN( EXISTS (
    SELECT Major_Code, COUNT (*)
    FROM Student
    GROUP BY Major_Code
    HAVING COUNT (*) > 1800 ))
```

```
CALL email_dean(Major_code);
```

Example

- Notify the Dean when the # of students in any major exceeds 1800

CREATE TRIGGER Major_Limit

AFTER INSERT OR UPDATE OF Major_Code
ON Student

```
WHEN( EXISTS (
    SELECT Major_Code, COUNT (*)
    FROM Student
    GROUP BY Major_Code
    HAVING COUNT (*) > 1800 ))
```

```
CALL email_dean(Major_code);
```

Example: Assertions Vs. Triggers

- ❑ **CREATE OR REPLACE ASSERTION** budget_constraint
CHECK (NOT EXISTS
(SELECT * FROM SECTION WHERE budget <
(SELECT SUM (Salary) FROM LIBRARIAN)));
- ❑ **CREATE OR REPLACE TRIGGER** budget_constraint_trigger
after INSERT, UPDATE of Salary
ON LIBRARIAN
WHEN (EXISTS (SELECT * FROM SECTION WHERE budget <
(SELECT SUM (Salary) FROM LIBRARIAN))

ROLLBACK;

Triggers (SQL99)

- ❑ **CREATE or REPLACE TRIGGER** <trigger-name>
 <time events> ON <list-of-tables>
 REFERENCING { NEW | OLD } AS <user-name>
 [FOR EACH { ROW | STATEMENT }]
 [WHEN (<Predicate>)]
 <action>
- ❑ time: **before** or **after**
- ❑ events: **Insert, Delete, Update** [of <list of attributes>]
- ❑ **NEW & OLD** refer to new & old (existing) tuples/table respectively
- ❑ The REFERENCING clause assigns aliases to NEW and OLD
- ❑ action: Stored procedure or
 BEGIN ATOMIC {<SQL procedural statements>} **END**

Oracle Example: Statement Trigger

- ❑ Statement-level trigger fires once by the triggering statement & defined on a single table
- ❑ **No WHEN-clause** in the definition of statement trigger
- ❑ **CREATE OR REPLACE TRIGGER** Audit_Updater
AFTER
INSERT OR DELETE OR UPDATE
ON STUDENTS
BEGIN
INSERT INTO AUDIT_Table VALUES ('STUDENT', sysdate);
END;
/
- ❑ The end slash ("/") installs and activates the trigger

Oracle Example: Row-Level Trigger

- ❑ Row- or tuple-level trigger fires once for each row affected by the triggering statement
- ❑ In Oracle, triggers are defined on a single table.
CREATE OR REPLACE TRIGGER trigger_deans_list
AFTER INSERT ON STUDENTS
REFERENCING NEW AS newRow
FOR EACH ROW
WHEN (newRow.QPA > 3.5)
BEGIN
INSERT INTO DL VALUES (:newRow.SID, :newRow.QPA);
END;
/
- ❑ Scope Rules: In the trigger body, NEW and OLD are global “variables” and must be preceded by a colon (“:”), but in the WHEN clause (triggering condition), they do not have a preceding colon!

Mutating Trigger

- ❑ Recursive call of triggers is not permitted
- ❑ Table read in a trigger it cannot be updated
- CREATE OR REPLACE TRIGGER** my_bad_auto_sid
AFTER INSERT ON STUDENTS
FOR EACH ROW
BEGIN
SELECT MAX(SID) +1 **INTO** : NEW.SID
FROM Students;
END;
/
- ❑ ERROR at line 1:
ORA-04084: cannot change NEW values for this trigger type

Mutating Trigger

- ❑ **Before** does not acquire locks
- CREATE OR REPLACE TRIGGER** my_auto_sid
BEFORE INSERT ON STUDENTS
FOR EACH ROW
BEGIN
SELECT MAX(SID) +1 **INTO** : NEW.SID
FROM Students;
END;
/
- ❑ Trigger created.
- ❑ **INTO**: the tuple assignment operator in PL/SQL

Auto-increment in Oracle

- ❑ It is achieved with a trigger
- ❑ Two special tables **dual** & **sequence**
`CREATE SEQUENCE oracle_example
NOCYCLE MAXVALUE 99999 START With 1;`

- ❑ dual is a table created by Oracle

SQL> describe dual;

Name	Null?	Type
DUMMY		VARCHAR2(1)

SQL> select * from dual;

DUMMY
X

Auto-increment in Oracle ...

- ❑ Trigger on the table with auto-increment

```
CREATE OR REPLACE TRIGGER
auto_increment_example
BEFORE INSERT ON STUDENTS
FOR EACH ROW
DECLARE next_id integer
BEGIN
  SELECT oracle_example.NEXTVAL INTO next_id
  FROM dual;

  :new.SID := next_id;
END;
/
```

Enable & Disable Triggers

- ❑ Enable/Disable All Triggers:
`ALTER TABLE <table_name> ENABLE ALL TRIGGERS;`
`ALTER TABLE <table_name> DISABLE ALL TRIGGERS;`
 - E.g., `ALTER TABLE Librarian DISABLE ALL TRIGGERS;`
- ❑ Enable/Disable Individual Trigger
`ALTER TRIGGER <trigger_name> ENABLE | DISABLE;`
 - E.g., `ALTER TRIGGER librarian_salary_trigger DISABLE;`
- ❑ Drop A Trigger
`DROP TRIGGER <trigger-name>;`
 - E.g., `DROP TRIGGER librarian_salary_trigger;`

Dropping a Trigger & Final note on IC

- ❑ Assertions and Checks Vs. Triggers
 - Assertions and Checks support the declarative approach of supporting Integrity Constraints
 - Triggers combine the declarative and procedural approach of implementing integrity constraints