Semantics of Functional Programming

The Scott Model of **PCF**

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Formosan Summer School on Logic, Language, and Computation 2014

1 Scott domain model

Denotational semantics of PCF

Instead of specifying *how* a program runs, we specify *what* a program is, the *denotation* of a program. To assign a denotation to a program,

- each type σ is interpreted as some domain D_σ;
- a context $\Gamma = x_1 : \sigma_1, \dots, x_n : \sigma_n$ is interpreted as a product $\prod_{i=1}^{n} D_{\sigma_i}$ of domains;
- in particular, each term of type τ under the empty context is an element of D_τ.

In the end, we show that $(\lambda x.\, \mathsf{M})\ \mathsf{N}$ and $\lambda x.\, \mathsf{M}\ x$ have the same denotation as $\mathsf{M}[\mathsf{N}/x]$ and M respectively, and also the Compactness Theorem for the Scott domain model of **PCF**.

Interpretation of types and contexts

Define the denotation of a type inductively:

Definition 1. Every type σ in **PCF** associates with a domain D_σ as follows:

1. $D_{\mathtt{nat}} \coloneqq \mathbb{N}_\perp$, and
2. $D_{\tau\to\sigma} \coloneqq [D_\tau \to D_\sigma]$.

Define the denotation of a context inductively on its length:

Definition 2. For each context $\Gamma = x_1 : \sigma_1, \dots, x_n : \sigma_n$, the associated domain is defined as

$$D_\Gamma \coloneqq D_{\sigma_1} \times D_{\sigma_2} \times \cdots \times D_{\sigma_n}$$

and the associated domain of the empty context is $1 = \{*\}$.

Interpretation of judgements

To proceed with the denotational semantics, we further define the denotation for each judgement inductively on its derivation of the following form

- Every judgement $\Gamma \vdash \mathsf{M} : \tau$ is interpreted as a *continuous* function

$$[\![\Gamma \vdash \mathsf{M} : \tau]\!] : D_\Gamma \to D_\tau.$$

- In particular,

$$[\![\ \vdash \mathsf{M} : \tau]\!] : 1 \to D_\tau$$

is identified with an element $[\![\ \vdash \mathsf{M} : \tau]\!](*) = d$ of D_τ.

Convention

In the following context, $[\![\Gamma \vdash \mathsf{M} : \tau]\!](\vec{d})$ is written as

$$[\![\Gamma \vdash \mathsf{M} : \tau]\!]\ \vec{d}.$$

for any sequence $\vec{d} \in D_\Gamma$ if there is no danger of ambiguity.

(var) Suppose that $\Gamma \vdash \mathsf{M} : \tau$ is of the form

$$x_1 : \sigma_1,\ \dots,\ x_n : \sigma_n \vdash x_i : \sigma_i$$

derived by the rule (var). It is interpreted as the projection from D_Γ to its i-th component D_{σ_i}

$$[\![x_1 : \sigma_1,\ \dots,\ x_n : \sigma_n \vdash x_i : \sigma_i]\!](\vec{d}) \coloneqq (d_i)$$

for $i = 1, \dots, n$ where $\vec{d} = (d_1, \dots, d_n) \in D_{\sigma_1} \times \cdots \times D_{\sigma_n}$.

Note that the denotation of this judgement is equal to

$$[\![x_1 : \sigma_1, \dots, x_n : \sigma_n \vdash x_i : \sigma_i]\!] = \pi_i$$

where $\pi_i \colon D_\Gamma \to D_{\sigma_i}$ is the i-th projection and thus it is a continuous function.

(abs) Let $f \coloneqq [\![\Gamma, x : \sigma \vdash \mathsf{M} : \tau]\!]$ be the continuous function from $D_\Gamma \times D_\sigma$ to D_τ.

$$[\![\Gamma \vdash \lambda x.\, \mathsf{M} : \sigma \to \tau]\!] \coloneqq \Lambda f$$

where $\Lambda f : D_\Gamma \to [D_\sigma \to D_\tau]$ is the *curried* f. In other words

$$\left([\![\Gamma \vdash \lambda x.\, \mathsf{M} : \sigma \to \tau]\!]\ \vec{d}\right)\ d = [\![\Gamma, x : \sigma \vdash \mathsf{M} : \tau]\!]\ (\vec{d}, d).$$

(app) Define

$$\begin{aligned}&[\![\Gamma \vdash \mathsf{M}\ \mathsf{N} : \tau]\!]\ \vec{d} \\ &:= ev\left([\![\Gamma \vdash \mathsf{M} : \sigma \to \tau]\!]\ \vec{d}, [\![\Gamma \vdash \mathsf{N} : \sigma]\!]\ \vec{d}\right)\end{aligned}$$

where $ev\colon [D_1 \to D_2] \times D_1 \to D_2$ is the *evaluation map* which maps a continuous function $f\colon D_1 \to D_2$ with an element $d \in D_1$ to $f(d)$.

The cases for `zero` and `suc` M are rather obvious:

(z) `zero` is a constant, so it does not matter what the context is:

$$[\![\Gamma \vdash \mathtt{zero} : \mathtt{nat}]\!]\ \vec{d} := 0$$

i.e. a constant function.

(s) The denotation of `suc` is the successor function

$$[\![\Gamma \vdash \mathtt{suc}\ \mathsf{M} : \mathtt{nat}]\!]\ \vec{d} := (S \circ [\![\Gamma \vdash \mathsf{M} : \mathtt{nat}]\!])\ \vec{d}$$

where $S\colon \mathbb{N}_\perp \to \mathbb{N}_\perp$ is defined by

$$S(n) := \begin{cases}\perp & \text{if } n = \perp \\ n+1 & \text{if } n \in \mathbb{N}.\end{cases}$$

(Y) The denotation of `Y` is the fixpoint operation

$$[\![\Gamma \vdash \mathtt{Y}x.\,\mathsf{M} : \sigma]\!]\ \vec{d} := \mu\left([\![\Gamma, x : \sigma \vdash \mathsf{M} : \sigma]\!]\ \vec{d}\right)$$

where μ is defined previously as $\mu(f) := \bigsqcup_{i\in\mathbb{N}} f^i(\perp)$.

(ifz) The denotation of `ifz`

$$\begin{aligned}&[\![\Gamma \vdash \mathtt{ifz}(\mathsf{M}; \mathsf{M}_0; \mathsf{M}_1) : \tau]\!]\ \vec{d} \\ &:= \mathit{ifz}_\tau(n, d, f)\end{aligned}$$

where

1. $n := [\![\Gamma \vdash \mathsf{M} : \mathtt{nat}]\!]\ \vec{d}$,
2. $d := [\![\Gamma \vdash \mathsf{M}_0 : \tau]\!]\ \vec{d}$,
3. $f := [\![\Gamma, x : \mathtt{nat} \vdash \mathsf{M}_1 : \tau]\!]\ \vec{d}$,

and ifz_τ is defined by

$$\mathit{ifz}_\tau(n, x, f) := \begin{cases}\perp & \text{if } n = \perp, \\ x & \text{if } n = 0, \\ f(m) & \text{if } n = m+1.\end{cases}$$

Theorem 3. *For every judgement $\Gamma \vdash \mathsf{M} : \tau$, the associated function*

$$[\![\Gamma \vdash \mathsf{M} : \tau]\!]\colon D_\Gamma \to D_\tau$$

is Scott continuous.

Proof sketch. It is not hard to see that each case of $[\![\Gamma \vdash \mathsf{M} : \tau]\!]$ is a Scott continuous function. □

Examples

In-class exercise

2 Substitution and Compactness

Substitution Lemma

Lemma 4. *Let $\Gamma = x_1 : \sigma_1, \ldots, x_n : \sigma_n$ be a context, and $\Gamma \vdash \mathsf{M} : \tau$ a judgement. Then the following equation*

$$[\![\Delta \vdash \mathsf{M}[\vec{N}/\vec{x}] : \tau]\!]\ \vec{d} = [\![\Gamma \vdash \mathsf{M} : \tau]\!]\left([\![\Delta \vdash \mathsf{N}_1]\!]\ \vec{d}, \ldots, [\![\Delta \vdash \mathsf{N}_n]\!]\ \vec{d}\right)$$

holds for any context Δ and judgements $\Delta \vdash \mathsf{N}_i : \sigma_i$ for $i = 1, \ldots, n$.

Proof. We prove it by induction on derivations of $\Gamma \vdash \mathsf{M} : \tau$. □

Proof of Substitution Lemma

(var) Suppose that $\Gamma \vdash \mathsf{M} : \tau$ is of the form

$$x_1 : \sigma_1, \ldots, x_n : \sigma_n \vdash x_i : \sigma_i$$

for $i = 1, \ldots, n$. Then, for each family of judgements $\Delta \vdash \mathsf{N}_i : \sigma_i$, it follows that

$$\begin{aligned}&[\![\Delta \vdash x_i[\vec{n}/\vec{x}] : \sigma_i]\!] \\ &= [\![\Delta \vdash \mathsf{N}_i : \sigma_i]\!] \\ &= \pi_i\left([\![\Delta \vdash \mathsf{N}_1 : \sigma_1]\!], \ldots, [\![\Delta \vdash \mathsf{N}_n : \sigma_n]\!]\right)\end{aligned}$$

where $[\![x_1 : \sigma_1, \ldots, x_n : \sigma_n \vdash x_i : \sigma_i]\!] = \pi_i$.

Corollary 5 (Application)**.** *For every judgement $\Gamma, x : \sigma \vdash \mathsf{M} : \tau$ and $\Gamma \vdash \mathsf{N} : \sigma$, we have*

$$[\![\Gamma \vdash (\lambda x.\,\mathsf{M})\ \mathsf{N} : \tau]\!] = [\![\Gamma \vdash \mathsf{M}[\mathsf{N}/x] : \tau]\!].$$

Observe that $\vec{d} = ([\![\Gamma \vdash x_1 : \sigma_1]\!]\ \vec{d}, \ldots, [\![\Gamma \vdash x_n : \sigma_n]\!]\ \vec{d})$ for any context $\Gamma = x_1 : \sigma_1, \ldots, x_n : \sigma_n$. Then, this corollary is a series of simple facts:

$$\begin{aligned}&[\![\Gamma \vdash (\lambda x.\,\mathsf{M})\ \mathsf{N} : \tau]\!]\ \vec{d} \\ &= ev\left([\![\Gamma \vdash (\lambda x.\,\mathsf{M}) : \sigma \to \tau]\!]\ \vec{d}, [\![\Gamma \vdash \mathsf{N} : \sigma]\!]\ \vec{d}\right) \\ &= ev\left([\![\Gamma, x : \sigma \vdash \mathsf{M} : \tau]\!]\ \vec{d}, [\![\Gamma \vdash \mathsf{N} : \sigma]\!]\ \vec{d}\right) \\ &= [\![\Gamma, x : \sigma \vdash \mathsf{M} : \tau]\!]\ (\vec{d}, [\![\Gamma \vdash \mathsf{N} : \sigma\ \vec{d}]\!]) \\ &= [\![\Gamma \vdash \mathsf{M}[\vec{x}, \mathsf{N}/\vec{x}, x] : \tau]\!]\ \vec{d} \\ &= [\![\Gamma \vdash \mathsf{M}[\mathsf{N}/x] : \tau]\!]\ \vec{d}\end{aligned}$$

Lemma 6 (Weakening)**.** *Let* $\Gamma \vdash \mathsf{M} : \tau$ *be a judgement. Then the following*

$$[\![\Gamma \vdash \mathsf{M} : \tau]\!] = [\![\Gamma, x : \sigma \vdash \mathsf{M} : \tau]\!]$$

holds for any variable $x : \sigma$ *not in* Γ*.*

It follows from Substitution Lemma. (*Why?*)

Corollary 7 (Extensionality)**.** *Let* $\Gamma \vdash \mathsf{M} : \sigma \to \tau$ *be a judgement. Then,*

$$[\![\Gamma \vdash \lambda x.\, \mathsf{M}\; x : \sigma \to \tau]\!] = [\![\Gamma \vdash \mathsf{M} : \sigma \to \tau]\!]$$

if x *is not a variable in* Γ*.*

For every sequence $\vec{d} \in D_\Gamma$ and $d \in D_\sigma$, we have

$$\begin{aligned}
&\left([\![\Gamma \vdash \lambda x.\, \mathsf{M}\; x : \sigma \to \tau]\!]\; \vec{d}\right)\; d \\
&= [\![\Gamma, x : \sigma \vdash \mathsf{M}\; x : \tau]\!](\vec{d}, d) \\
&= ev\left([\![\Gamma, x : \sigma \vdash \mathsf{M} : \sigma \to \tau]\!](\vec{d}, d), [\![\Gamma, x : \sigma \vdash x : \sigma]\!](\vec{d}, d)\right) \\
&= ev\left([\![\Gamma, x : \sigma \vdash \mathsf{M} : \sigma \to \tau]\!](\vec{d}, d), d\right) \\
&= \left([\![\Gamma, x : \sigma \vdash \mathsf{M} : \sigma \to \tau]\!](\vec{d}, d)\right) d \\
&= \left([\![\Gamma \vdash \mathsf{M} : \sigma \to \tau]\!]\; \vec{d}\right)\; d.
\end{aligned}$$

Compactness

Define $\mathtt{Y}^i x.\, \mathsf{M}$ inductively for each $i \in \mathbb{N}$ by

1. $\mathtt{Y}^0 x.\, \mathsf{M} := \mathtt{Y} x.\, x$ and
2. $\mathtt{Y}^{n+1} x.\, \mathsf{M} := \mathsf{M}[\mathtt{Y}^n x.\, \mathsf{M}/x]$.

Theorem 8. *For every judgement* $\Gamma, x : \sigma \vdash \mathsf{M} : \sigma$*, we have*

$$[\![\Gamma \vdash \mathtt{Y} x.\, \mathsf{M} : \sigma]\!] = \bigsqcup_{i \in \mathbb{N}} [\![\Gamma \vdash \mathtt{Y}^i x.\, \mathsf{M} : \sigma]\!].$$

To show this theorem, it suffices to show the following

$$[\![\, \vdash \mathtt{Y}^i x.\, \mathsf{M} : \sigma]\!] = [\![x : \sigma \vdash \mathsf{M} : \sigma]\!]^i(\bot)$$

for $i \in \mathbb{N}$. (Why?)

For $n = 0$ we show that $[\![\, \vdash \mathtt{Y}^0 x.\, \mathsf{M} : \sigma]\!] = \bot_{D_\sigma} \in D_\sigma$.

By definition, $\mathtt{Y}^0 x.\, \mathsf{M} : \sigma$ is equal to $\mathtt{Y} x.\, x$, so

$$\begin{aligned}
[\![\, \vdash \mathtt{Y} x.\, x : \sigma]\!] = \mu(id) &= \bigsqcup_{i \in \mathbb{N}} id^i(\bot) \\
&= \bigsqcup \bot = \bot
\end{aligned}$$

For $i = n + 1$ it suffices to show that

$$[\![\, \vdash \mathtt{Y}^{n+1} x.\, \mathsf{M} : \sigma]\!] = [\![x : \sigma \vdash \mathsf{M} : \sigma]\!]\left([\![\, \vdash \mathtt{Y}^n x.\, \mathsf{M} : \sigma]\!]\right),$$

so the statement follows by the induction hypothesis.

By definition, $\mathtt{Y}^{n+1} x.\, \mathsf{M}$ is equal to $\mathsf{M}[\mathtt{Y}^n x.\, \mathsf{M}/x]$, and by Substitution Lemma we have

$$[\![\, \vdash \mathsf{M}[\mathtt{Y}^n x.\, \mathsf{M}/x]]\!] = [\![x : \sigma \vdash \mathsf{M} : \tau]\!]\left([\![\, \vdash \mathtt{Y}^n x.\, \mathsf{M}]\!]\right).$$

Examples

In-class exercise