

1. Prove that language  $L = \{0^n \mid n \text{ is a power of } 2\}$  is not regular.

证明思路：泵引理，反证法。取  $s = 0^{2^N}$ ，则  $2^N < |xy^2z| < 2^N + N < 2^N + 2^N = 2^{(N+1)}$

2. If  $L$  is a language, and  $a$  is a symbol, then  $L/a$ , the quotient of  $L$  and  $a$ , is the set of strings  $w$  such that  $wa$  is in  $L$ . For example, if  $L = \{a, aab, baa\}$ , then  $L/a = \{\varepsilon, ba\}$ . Prove that if  $L$  is regular, so is  $L/a$ . Hint: Start with a DFA for  $L$  and consider the set of accepting states.

证明：令  $L = L(M)$ , 其中  $M = (Q, \Sigma, \delta, q_0, F)$

构造  $M' = (Q, \Sigma, \delta, q_0, F')$ , 其中  $F' = \{q \mid \delta(q, a) \in F\}, q \in Q, a \in \Sigma$

证明  $L(M') = L/a$ ,  $\because \forall w \in L(M')$  即  $\hat{\delta}(q_0, w) \in F'$  即  $\delta(\hat{\delta}(q_0, w), a) \in F \therefore w \in L/a$

又  $\because \forall w \in L/a$  有  $wa \in L$  即  $\hat{\delta}(q_0, wa) \in F$  即  $\delta(\hat{\delta}(q_0, w), a) \in F$  即  $\hat{\delta}(q_0, w) \in F' \therefore w \in L(M')$

Design context-free grammars for the following languages:

3. The set  $\{a^i b^j c^k \mid i \neq j \text{ or } j \neq k\}$ , that is, the set of strings of  $a$ 's followed by  $b$ 's followed by  $c$ 's, such that there are either a different number of  $a$ 's and  $b$ 's or a different number of  $b$ 's and  $c$ 's, or both.

$S \rightarrow A_1 C \mid A_2 C \mid AB_1 \mid AB_2$

$A_1 \rightarrow aA_1b \mid aA_1 \mid a$

$A_2 \rightarrow aA_2b \mid A_2b \mid b$

$C \rightarrow Cc \mid \varepsilon$

$B_1 \rightarrow bB_1c \mid bB_1 \mid b$

$B_2 \rightarrow bB_2c \mid B_2c \mid c$

$A \rightarrow Aa \mid \varepsilon$

(注意:  $Cc \mid \varepsilon$  若为  $Cc \mid c$  则不能产生  $a, c$  同时为 0 个, 或  $b, c$ )

4. The set of all strings over  $\{0, 1\}$  with twice as many 0's as 1's.

$S \rightarrow S0S0S1S \mid S0S1S0S \mid S1S0S0S \mid \varepsilon$

5. The set of all strings over  $\{a, b\}$  that are **not** of the form  $ww$ , for some string  $w$ . Explain how your grammar works. You needn't prove it's correctness formally.

如果串长为奇数, 显然不是  $ww$  形式 (对应下面文法中的  $A$  或  $B$ )。而对于长度为偶数 ( $2n$ ) 的串, 至少存在一对儿距离为  $n$  (串长度的一半) 的两字符不相同。为了能够产生两个不相同字符的距离刚好是整个长度的一半, 使用两个变元  $A$  和  $B$  分别产生基数长的串, 然后合并即可。对  $A$  或  $B$ , 在为产生串时, 如果增加了两字符间的字符数, 那么也要增加两字符外的字符数。如  $aaaabbbb = \underline{aaa}ab\underline{bbb}$  或  $aabaaa = \underline{aab}aa\underline{a}$ 。

$S \rightarrow A \mid B \mid AB \mid BA$

$A \rightarrow XAX \mid a$

$B \rightarrow XBX \mid b$

$X \rightarrow a \mid b$