Talos Vulnerability Report

TALOS-2020-0993

Accusoft ImageGear JPEG jpegread precision code execution vulnerability

FEBRUARY 10, 2020

CVE NUMBER

CVE-2020-6069

Summary

An exploitable out-of-bounds write vulnerability exists in the igcore19d.dll JPEG jpegread precision parser of the Accusoft ImageGear 19.5.0 library. A specially crafted JPEG file can cause an out-of-bounds write, resulting in a remote code execution. An attacker needs to provide a malformed file to the victim to trigger the vulnerability.

Tested Versions

Accusoft ImageGear 19.5.0

Product URLs

https://www.accusoft.com/products/imagegear/overview/

CVSSv3 Score

9.8 - CVSS:3.0/AV:N/AC:L/PR:N/UI:N/S:U/C:H/I:H/A:H

CWE

CWE-787: Out-of-bounds Write

Details

The ImageGear library is a document imaging developer toolkit providing all kinds of functionality related to image conversion, creation, editing, annotation, etc. It supports more than 100 formats, including many image formats, DICOM, PDF, Microsoft Office and others.

There is a vulnerability in the JPEG raster image parser. A specially crafted JPEG file can lead to an out-of-bounds write resulting in remote code execution.

If we try to load a malformed JPEG file via the IG_load_file function we end up in the following situation:

```
(397c.4b88): Access violation - code c0000005 (first chance)
First chance exceptions are reported before any exception handling.
This exception may be expected and handled.
eax=ffff8887f ebx=000000000 ex-ffffe880 edx-0c0667004 esi=00000fff edi=ffff887f
eip=5b9d4669 esp=0100f084 ebp=0100f09c iopl=0 nv up ei pl zr na pe nc
cs=0023 ss=002b ds=002b es-002b fs=0053 gs=002b efl=00010246
igCore1pd/1Eq mip_lape_set+0x/02409:
5b9d4669 66895afc mov word ptr [edx-4],bx ds:002b:0c067000=????
```

Checking the buffer capacity pointed to by edx, we can see:

The calculated address 0c067000 is outside of the allocation for this buffer space.

Going back to the allocation, we land in the following place:

The getMemSize function at line 4 calculates the returned mem_size as follows:

Where the variable/fields are located in the file at offsets:

```
0xA2 : precision : SOFx[0]->precision = 7
0xA5 : SOFx[0]->x_image = 0xDE
0xA7 : SOFx[0]->nr_comp = 3
```

Next, a fixed value of 24 is added to returned value and we end up with 0x2b4 allocated bytes.

Let use take a look now at the vulnerable function:

```
Line 1 signed __int16 *__cdecl sub_5B9D45C0(int a1, int loop_limit, signed __int16 *a3, signed __int16 *a4, signed __int16 *a5, int a6, struct_a7 *a7, int a8)
Line 2 {
Line 3 s
Line 4
              signed __int16 *result; // eax
int v9; // edi
int v10; // esi
int v11; // kr00_4
int v12; // ecx
line 5
Line 6
Line 7
              int v13; // ecx
int v14; // ebx
bool v15; // zf
int v16; // [esp+0h] [ebp-Ch]
int v17; // [esp+4h] [ebp-8h]
int v18; // [esp+8h] [ebp-4h]
Line 8
Line 9
Line 10
Line 11
Line 12
Line 13
Line 14
Line 15
               _WORD *buffer; // [esp+28h] [ebp+1Ch]
               result = 0;
Line 16
               v16 = 0;
v17 = 0;
v18 = 0;
Line 17
Line 18
Line 19
              if ( loop_limit > 0 ) {
                        buffer = (\_WORD *)(a6 + 4);
Line 22
Line 23
Line 24
                        result = a5;
do
Line 25
                           v9 = *a3 + 2048;
v10 = *a4;
v11 = 5742 * *result;
Line 26
Line 27
Line 28
                           V11 = 9/4 * *result;

V12 = 90 * (2925 * *result + 1410 * v10) / 4096;

V13 = V0 * 7258 * v10 / 4096;

V14 = V0 * V11 / 4096;

if ( v9 * v11 / 4096 >= 0 )
Line 29
Line 30
Line 31
Line 32
Line 33
                                   if ( v14 > 4095 )
LOWORD(v14) = 4095;
Line 34
Line 35
Line 36
Line 37
Line 38
                           else
{
Line 39
Line 40
Line 41
                                    LOWORD(v14) = 0;
                           *(buffer - 2) = v14;
if ( v12 >= 0 )
Line 42
Line 43
Line 44
                                    if ( v12 > 4095 )
Line 45
                                       LOWORD(v12) = 4095:
Line 46
Line 47
                           else
Line 48
Line 49
Line 50
                           {
                                    LOWORD(v12) = 0;
Line 51
Line 52
Line 53
                           *(buffer - 1) = v12;
if ( v13 >= 0 )
{
                                    if ( v13 > 4095 )
LOWORD(v13) = 4095;
Line 54
Line 55
Line 56
Line 57
Line 58
                                   IOWORD(v13) = 0:
Line 59
Line 60
                           *buffer = v13;
v18 += a7->dword34;
Line 61
Line 62
                           if ( v18 == a8 )
Line 63
Line 64
Line 65
                                     ++a3;
                                    v18 = 0;
Line 66
Line 67
                           v17 += a7->dword84;
line 68
                           if ( v17 == a8 )
Line 69
Line 70
                                    ++a4;
v17 = 0;
Line 71
Line 72
Line 73
Line 74
                           v16 += a7->dwordD4;
Line 75
Line 76
Line 77
                           if ( v16 == a8 )
                                    v16 = 0;
                                    result = a5 + 1;
++a5;
Line 78
Line 79
Line 81
                           else
Line 82
Line 83
                                    result = a5;
Line 84
                           v15 = loop_limit-- == 1;
buffer += 3;
Line 85
Line 86
Line 87
                        while ( !v15 );
Line 88
Line 89
              return result;
line 90
Line 91 }
```

The loop_limit variable is equal to 0xA5 : SOFx[0]->x_image = 0xDE. Each loop cycle, the buffer pointer is increased by 3 * WORD = 6 at line 86.

So for buffer with capacity 0x2B4 bytes this loop is able to execute a maximum of 0x2B4 / 6 = 00000073 times but loop_limit equals 0xDE and is decrement by 1 during each cycle, which leads to an out-of-bounds write and memory corruption.

As we can see an attacker controls all presented variables just by proper file content manipulation.

Increasing the loop count via the loop_limit variable an attacker can cause an out-of-bounds write leading to memory corruption, which can result in remote code execution.

Timeline

2020-01-27 - Vendor Disclosure 2020-02-10 - Public Release

CREDIT

Discovered by Emmanuel Tacheau and a member of Cisco Talos.

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