

## Talos Vulnerability Report

TALOS-2020-0984

### CoTURN HTTP Server POST-parsing information leak vulnerability

FEBRUARY 18, 2020

CVE NUMBER

CVE-2020-6061

#### Summary

An exploitable heap out-of-bounds read vulnerability exists in the way CoTURN 4.5.1.1 web server parses POST requests. A specially crafted HTTP POST request can lead to information leaks and other misbehavior. An attacker needs to send an HTTPS request to trigger this vulnerability.

#### Tested Versions

CoTURN 4.5.1.1

#### Product URLs

<https://github.com/coturn/coturn>

#### CVSSv3 Score

7.0 - CVSS:3.0/AV:N/AC:H/PR:N/UI:N/S:U/C:L/I:L/A:H

#### CWE

CWE-125: Out-of-bounds Read

#### Details

CoTURN is a TURN server implementation. A TURN Server is a VoIP media traffic NAT traversal server and gateway. CoTURN can be used as a general-purpose network traffic TURN server and gateway.

For administration purposes, it includes a web server. Code responsible for parsing POST request body variables contains a bug that can lead to out of bounds memory access.

When preparing to parse the POST request body, the following code is executed:

```
static struct headers_list * post_parse(char *data, size_t data_len)
{
    while((*data=='\r')||(*data=='\n')) ++data;           [1]
    char *post_data = (char*)calloc(data_len + 1, sizeof(char));
    memcpy(post_data, data, data_len);                    [2]
```

To start, at [1] newline and carriage return characters are skipped in order to get to the start of POST data. However, while data pointer is incremented, the data\_len isn't decremented. Then, at [2], the memcpy call will copy data from incremented data pointer into newly allocated post\_data memory buffer using the unchanged data\_len. This results in bytes beyond the end of original data buffer being accessed.

A POST request of following form can be used to trigger this issue:

```
"POST /logon HTTP/1.1\r\nContent-Length: 32717\r\n\r\n" + "\x0d"*33000 + "u\r\n\r\n\r\n"
```

Using extra \r or \n characters, we can control how much the data pointer gets incremented at [1]. Also, the content length header controls the allocation size. By aligning those two, we can have the while loop at [1] skip till the actual end of the data buffer which would result in a large out of bounds access at [2].

Depending on the memory layout, this could potentially result in further memory corruption, access to sensitive information from other requests and other unforeseen consequences.

#### Timeline

2020-02-11 - Vendor Disclosure

2020-02-17 - Vendor patched

2020-02-18 - Public Release

#### CREDIT

Discovered by Aleksandar Nikolic of Cisco Talos.

