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From: Matthias Gerstner <mgerstner@...e.de>

To: oss-security@...ts.openwall.com

Subject: Security Issues in the spice-vdagentd daemon

Hello list,

please find below a security report regarding the spice-vdagentd [1].

Attached to this mail is a tarball containing the final set of patches that upstream developed to address the issues. These patches are by now already published in the upstream git repository. Furthermore attached is a tarball containing scripts and source code to reproduce some of the issues discussed in the report below. CVEs have been assigned by upstream for the individual issues, please find them in the respective sections of the report.

[1]: [https://gitlab.freedesktop.org/spice/linux/vd\\_agent](https://gitlab.freedesktop.org/spice/linux/vd_agent)

#### # 1) Introduction

The SUSE security team got a request to review the 'spice-vdagentd' daemon. It has been brought to our attention that this daemon can start indirectly via udev ('70-spice-vdagentd.rules') even if the service was not explicitly activated by the Administrator. Thus it might introduce additional security attack surface in default SUSE installations.

The 'spice-vdagentd' is typically used in Qemu virtual machines to provide additional features like:

- a shared clipboard between host and guest system to allow seamless copy/paste
- file transfers from the host to the guest machine
- sharing information about the virtual display size and monitor configuration
- synchronizing audio device volume between host and guest

'spice-vdagentd' is running as \*root\* as a socket activated systemd service.

The following security report is based on the spice-vdagent version 0.20.0, which was the latest upstream release available at the time of its writing.

#### # 2) The 'vdagentd' Communication Channels

This section gives a short introduction about the communication channels used by 'vdagentd' for readers that are not familiar with its design.

##### ## a) The com.redhat.spice.0 Virtio Serial Device

In Qemu virtual machine instances, that have the spice protocol enabled, a virtual serial device will be emulated, typically enabled by passing switches like these to qemu:

...

```
-spice disable-ticketing,unix,addr=/path/to/spice_socket, -device virtio-serial-pci -device virtserialport,chardev=spicechannel0,name=com.redhat.spice.0 -chardev spicevmc,id=spicechannel0,name=vdagent
```

Within the virtual machine 'spice-vdagentd' talks to this serial device ('/dev/virtio-ports/com.redhat.spice.0') to exchange information with the host machine and implement the spice features. The SPICE protocol is used on this communication channel.

This serial device is of lesser importance in the context of this security report.

##### ## b) The spice-vdagentd UNIX Domain Socket

The 'spice-vdagentd' creates a listening UNIX domain socket of type 'SOCK\_STREAM' in '/run/spice-vdagentd/spice-vdagent-sock'. This socket path is accessible to all users in the system (file mode 0666). It is used by the counterpart of 'spice-vdagentd', the 'spice-vdagent', which runs in the context of a graphical user session with user privileges. This role will also be called the "user agent" for the rest of this document.

The 'spice-vdagent' e.g. informs the 'spice-vdagentd' about changes in clipboard contents in the graphical VM session, which in turn forwards this information via the serial device to the host machine. Similarly requests from the host machine to e.g. initiate a file transfer will be forwarded from 'spice-vdagentd' to the 'spice-vdagent', which will then store the file data in the context of the user session.

Multiple user agents can connect to 'spice-vdagentd' at the same time. At most one user agent may be present in the currently active session, however. Only a user agent running in the currently active session can access most of the features of 'spice-vdagentd'. 'spice-vdagentd' determines the session a user agent is part of by querying systemd facilities (e.g. 'sd\_seat\_get\_active()', 'sd\_pid\_get\_session()'). The user agent (if any) associated with the currently active session will also be the target of any forwarded requests received from the host on the virtual serial device.

#### # 3) Security Issues

This report is accompanied by two Python scripts ('vdagent.py', 'create\_vdagent\_conns.py') and a C++ source file ('socket\_pid\_atFack.cxx') which can be used to reproduce the findings that follow. These reproducers will be mentioned in the individual sections. All reproducers need to be executed in the context of a host system running a SPICE enabled Qemu virtual machine (see section 2.a for the required Qemu switches, or use something like livbird and its Qemu backend with appropriate settings).

##### ## a) CVE-2020-25650: Memory DoS via Arbitrary Entries in 'active\_xfers' Hash Table

The 'spice-vdagentd' maintains a hash map named 'active\_xfers' that maps 'task\_ids' to UNIX domain socket connections they belong to. These 'task\_ids' refer to ongoing file transfers from the host to the virtual machine.

An arbitrary client connected to 'spice-vdagentd' via a UNIX domain socket can trigger an entry into this hash map, without the requirement that the client is associated with the currently active graphical session (function 'do\_agent\_file\_xfer\_status', specifically 'vdagentd.c:1025'). There is no limit on the maximum amount of file transfers ongoing in parallel and there are no timeouts applied for a file transfer to be finished.

Therefore any unprivileged local user with access to the '/run/spice-vdagentd/spice-vdagent-sock' socket path can perform a memory denial-of-service by entering a large amount of entries into this hash map.

##### ### Impact

- an unprivileged local attacker inside a virtual machine can cause

large memory allocation in a daemon running as `*root*`. These memory allocations will persist as long as the UNIX domain socket connection remains intact. After a longer while, depending on the available system memory and swap space, the system might enter an out of memory situation, causing a denial-of-service for `'spice-vdagentd'` or even other processes in the system (the Linux kernel OOM killer is known to punish innocent victims at times).

- after a large amount of entries have been made by an attacker this way and once the UNIX domain socket connection is terminated, the `'spice-vdagentd'` will remove all the hash map entries belonging to this connection (function `'agent_disconnect'`, specifically `'vdagentd.c:955'`). Each "cancelled" file transfer will be logged in the system (`'vdagentd.c:910'`, `'vdagentd.c:342'`). In my tests this also caused a high load in `'systemd-journald'` and high disk space consumption from the logs stored in `'/var/log'`.

### ### Fixes

- Patch 02 changes the handling of entries into the hash map. User agents will no longer be able to trigger arbitrary entries into it.
- Patch 03 limits the maximum number of transfers allowed to be active at the same time.

### ### Reproducer

Using the supplied Python program the DoS can be performed this way:

```
'''
user $ ./vdagent.py --xfer-status=DoS
'''
```

The program will infinitely create new entries in the `'active_xfers'` hash map up until the maximum 64-bit (for `'x86_64'` architectures) integer value. Memory consumption will grow rather slowly but reliably. An exponential growth allocation scheme seems to be in place, because the observed memory consumption (as seen in `'top'`) only happens in jumps.

### ## b) CVE-2020-25651: Possible File Transfer DoS and Information Leak via `'active_xfers'` Hash Map

The same basic problem as described in section 3.a can lead to a file transfer information leak. The file transfer protocol roughly works like this:

- The host will send a `'VD AGENT FILE XFER START'` message that is forwarded to the user agent (function `'do_client_file_xfer()'`, specifically `'vdagentd.c:376'`). This message contains a `'task_id'` that identifies the file transfer process in future messages.
- The `'spice-vdagent'` will check free disk space and allocate a file of the expected size in the file system. If all checks pass then it will reply with a `'VDAGENTD FILE XFER STATUS CAN SEND DATA'` message, which causes `'spice-vdagentd'` to associate the client connection with the ongoing file transfer.
- The host will now start sending out chunks of the file data with `'VDAGENTD FILE XFER DATA'` messages (processed in function `'do_client_file_xfer()'`, specifically `'vdagentd.c:386'`). `'spice-vdagentd'` will forward each chunk to the client connection stored in the `'active_xfers'` hash map.

The host application (tested with `'remote-viewer'` from the `virt-viewer` package [2]) chooses an incrementally growing `'task_id'` for file exchanges which starts counting at 1. Thus the `'task_id'` is predictable. Since any unauthenticated local client can replace the mapping of `'task_id'` to client connection by its own client connection, there is a possibility for an attacker to obtain parts of the transferred file data.

The attacker needs to win a race condition here, because it needs to hit the time window after the legitimate client sends out the `'VDAGENTD FILE XFER STATUS CAN SEND DATA'` message and before the host starts sending out file chunks via `'VDAGENTD FILE XFER DATA'`. If the attacker sends his own `'VDAGENTD FILE XFER STATUS CAN SEND DATA'` using the correct `'task_id'` during this time window, then he can obtain the complete file. At least for large file exchanges bigger parts of the file are feasible to be obtained, even when the initial parts of the file are transferred to the legitimate client.

The more difficult part for an attacker will be to identify when such a file transfer will take place. The reproducer shows the basic attack technique.

[2]: <https://gitlab.com/virt-viewer/virt-viewer>

### ### Impact

File data from the host system can end up in full or in parts in the client connection of an illegitimate local user in the VM system. Exploitability will be difficult if there is not a suitable side channel with information about file transfers going on.

In any case active file transfers from other users can also be interrupted (DoS aspect).

### ### Fixes

Additionally to the fixes in section 3.a:

- Patch 08 cleans all entries from the hash map when the active user agent disconnects.
- Patch 09 prevents existing entries in the hash map to be replaced.

### ### Reproducer

This reproducer has the precondition that the victim's home directory is world-readable (or at least readable by the attacker). The script will wait (using the `'inotify'` API) for a new file to be created in the home directory (assuming that this is the initiation of a file transfer).

For reproducing perform the following steps:

1. Open a fresh `'remote-viewer'` connection for the test VM. This will reset the `'task_id'` used for file transfers to start to 1.
2. Run the reproducer script in an attacker context as follows:

```
# replace this path by the actual victim's home directory
attacker$ VICTIM=/home/victim
attacker$ ./vdagent.py --send-xfer-status 1 --wait-for-file-create $VICTIM
```

3. Log in as the `'victim'` user in a graphical session in the VM using the `'remote-viewer'` connection started in step 1.
4. Drag-and-drop a larger file (I tested using an 8 megabyte text file) into the `'remote-viewer'` window to initiate file transfer.

You should see that the attacker's `'vdagent.py'` script outputs large parts of the file transfer to the terminal. The original file transfer will indicate an error condition without specific details about the problem.

Remember that for each subsequent attack attempt you will need to increment the `'task_id'` (`'--send-xfer-status'` parameter) or open a fresh `'remote-viewer'` window for the VM, to reset the `'task_id'` used on the host side.

### ## c) CVE-2020-25652: Possibility to Exhaust File Descriptors in `'vdagentd'`

`'spice-vdagentd'` does not apply a limit to the amount of client connections that can be established via the UNIX domain socket in `'/run/spice-vdagentd/spice-vdagent-sock'`. Also existing connections aren't subject to a timeout or any kind of preconditions for them to stay alive. Thus it is easy to exhaust the file descriptor limit for

the 'spice-vdagentd' process (typically 1024 file descriptors by default, this limit is also imposed by system calls like 'select()').

Any local user in the virtual machine can open around ~1020 connections to 'spice-vdagentd' and simply keep them open without transmitting any data. The 'spice-vdagentd' will then become unable to open further connections for legitimate clients or perform other tasks (like opening the serial device, see section 2.a, or invoking systemd library calls that require opening files).

### ### Impact

By exhausting file descriptors in 'spice-vdagentd' the following effects can be achieved:

- The attack can prevent legitimate 'spice-vdagent' instances from connecting to the 'spice-vdagentd'. SPICE features won't be available to affected sessions.
- The attack can cause 'vdagentd' to exit on error conditions if tuned carefully. For example, an attacker can exhaust all file descriptors in 'spice-vdagentd' except for one and then wait for a legitimate client from an active session to connect. This connection attempt will succeed, but the subsequent attempt to open the serial device (see section 2.a) will fail, and 'spice-vdagentd' will exit. This will then also cause the involved 'spice-vdagent' to exit, because the connection to the system daemon is lost.
- 'spice-vdagentd' will enter a 100 % CPU load infinite loop, because it tries to 'accept()' the new connection, which is impossible, but also doesn't close the listening socket or abort execution.
- This attack vector makes security issue 3.d better exploitable, which will be explained there in more detail.

### ### Fixes

- Patch 04 implements an upper limit of client connections accepted by the 'spice-vdagentd'.
- Patch 07 implements a limit for client connections established from the same session.

### ### Reproducer

Using the supplied script 'create\_vdagent\_conns.py' the maximum number of connections that 'spice-vdagentd' can process, can be established. The result will be that no new user agents can register with the system daemon. A 100 % CPU loop will occur in 'spice-vdagentd'.

### ## d) CVE-2020-25653: UNIX Domain Socket Peer PID Retrieved via 'SO\_PEERCREID' is Subject to Race Condition

One major security property of 'spice-vdagentd' is that it only allows those clients access to most of the SPICE features (like clipboard, file transfer) that are currently in an active session according to systemd (see also section 2.b). It is possible for arbitrary local users (like 'nobody') to connect to 'spice-vdagentd' but these connections should not be able to interact with the host machine, because they don't belong to the active session.

The session check is performed after a new UNIX domain socket connection is established in 'agent\_connect()' in 'vdagentd.c:937'. The check basically relies on these two source code lines:

```
...
pid = vdagent_connection_get_peer_pid(VDAGENT_CONNECTION(conn), &err);
agent_data->session = session_info_session_for_pid(session_info, pid);
...
```

The peer's PID is obtained via glib's 'g\_socket\_get\_credentials' which boils down to the 'SO\_PEERCREID' socket option that is supported for UNIX domain sockets (see 'man 7 socket', 'man 7 unix', 'struct ucred'). The man page says about this:

> The returned credentials are those that were in effect at the time of the call  
> to connect(2) or socketpair(2).

This means that there is a race condition between the point in time when a client performs the 'connect()' call to establish a connection with 'spice-vdagentd' and the time 'spice-vdagentd' retrieves and checks the PID in its 'agent\_connect()' function. The PID in question can already have been replaced by an unrelated process. Therefore the session that 'spice-vdagentd' associates with this PID might be a different one than the actual peer process belonged to, when the 'connect()' system call was performed.

An attack to exploit the race condition requires the following steps:

1. an attacker can inherit a UNIX domain socket file descriptor to a child process that performs the 'connect()' to 'spice-vdagentd' and exits immediately again, thereby freeing the PID (let's call it the malicious PID) in the system as soon as the parent process performs a 'wait()' on the exited child process. This malicious PID will now be associated in the kernel with the 'SO\_PEERCREID' data returned for the connected UNIX domain socket.
2. now the attacker needs to perform a PID cycle in the system (i.e. create many useless child processes to cause the maximum PID - typically 32768 - to be reached in the system and new processes get assigned small PIDs again). When the PIDs assigned by the kernel are getting close to the malicious PID, the attacker needs to stop creating child processes and wait for unrelated processes from other users to come into existence.
3. Once the malicious PID gets reassigned to an unrelated process and the 'agent\_connect()' function runs in 'spice-vdagentd', it will retrieve wrong session information for the existing connection. If the malicious PID gets reassigned to a process running in the active session, then the connection that the attacker uses will get access to the SPICE features and can communicate with the host, although the attacker would otherwise not have sufficient privileges to do so.

The described race condition is very hard to hit under normal circumstances, because step 2., the PID cycle, is taking a long time and the 'agent\_connect()' function in 'spice-vdagentd' is very likely to run before an unrelated process gets reassigned the malicious PID in question. When combined with the file descriptor exhaustion security issue described in section 3.c, however, then this attack will become way more feasible.

This combined attack works like follows:

- Exhaust all file descriptors in 'spice-vdagentd' as described in section 3.c.
- Perform the attack steps 1. and 2. as described previously. What happens now is that the attacker's UNIX domain socket 'connect()' will succeed, because on kernel level this is still possible. 'spice-vdagentd' won't be able to 'accept()' this connection, though, because no more file descriptors are available to do so. The connection remains pending on the listening socket, however.
- Now for step 3., once the attacker notices that the malicious PID got assigned to an unrelated process, he can stop the file descriptor exhaustion put into place previously, thus making it possible for 'spice-vdagentd' to 'accept()' the malicious connection pending in the kernel. Only now will the 'agent\_connect()' function run, and it will more reliably determine the wrong session for the connection.

### ### Impact

1. A compromised local account with little privileges inside the virtual machine like 'nobody' can try to become the 'active agent' for 'spice-vdagentd' for the graphical session of a legitimate local user. If successful then the attacker can access the host's clipboard content or send malicious clipboard content to the host. The attacker can also retrieve file data from the host (compare section 2.b) or send invalid screen resolution and display

- information to the host.
- The combined attack using the file descriptor exhaustion and the 'SO\_PEERCREC' race condition is still not 100 % reliable but it can be repeated many times to increase chances of success. The only unpredictable ingredient is victim child processes appearing that get assigned the desired malicious PID and stay around for long enough for 'spice-vdagentd' to pick up the wrong session information.
  - If the victim's graphical session already runs a legitimate 'spice-vdagent' then a successful attack will trigger an information leak protection logic in 'vdagentd.c:874'. This has the effect of a denial-of-service, because neither the attacker nor the legitimate user will be able to use the SPICE features anymore.
  - If the victim's graphical session is not running a 'spice-vdagent' then the attacker can achieve all the effects described in 1.
  - If 3. applies (the victim's is already running 'spice-vdagent') then the attacker could try to crash the currently running 'spice-vdagentd' (see section 3.c for a possible attack vector). systemd should then restart the 'spice-vdagentd' while the victim's 'spice-vdagent' should exit but not be restarted. After this situation 4) applies.

### Fixes

- Patches 05 and 06 change the session check logic mainly by also taking into account the connected client's UID. If the UID of the determined session and the client don't match then the connection will be terminated.

### Reproducer

This attack is quite complex and requires the combination of all supplied programs to succeed. All programs must be placed in the same directory, because they interact with each other. Perform the following steps:

- For reduced complexity it is a precondition that an active graphical session already exists in the VM, but no 'spice-vdagent' is running in it. Therefore log in the "victim" user in a graphical session, kill any 'spice-vdagent' that got started automatically and keep the session open.
- In the context of an "attacker" user run 'create\_vdagent\_conns.py':

```
attacker$ ./create_vdagent_conns.py
Established 1015 Active connections to vdagentd.
Waiting for Ctrl-C. Or press ENTER to close 5 sockets.
```

Wait until the program displays the message seen above. Keep the program running in this state.

- Still in the context of the attacker compile and run the C++ program:

```
attacker$ g++ -O2 socket_pid_attack.cxx -o socket_pid_attack
attacker$ ./socket_pid_attack
Target UDS PID = ???
Cycled to PID ....
[...]
Closing in to target_pid ????: Got child PID ???
Now waiting for ??? to get reassigned.
```

This program creates a UNIX domain socket connected to 'spice-vdagentd' but the 'connect()' call is performed in a short-living child process, thus freeing the PID the kernel stored for 'SO\_PEERCREC'. The program then waits for the malicious PID to be assigned to an unrelated process before it continues. Keep the program running in this state.

- In the context of the "victim" user create a number of long running new child processes in the graphical session. For example something like this:

```
victim$ for i in `seq 20`; do sleep 1h & done
```

This simulates new child processes being created in the victim's context to replace the malicious PID the attacker is waiting for.

- If successful then you should see in the program from step 3. that something happened:

```
PID ??? now exists, but can't read exe: Permission denied
Running './vdagent.py --socket-fd 3 '
Using existing connected socket file descriptor
```

If this did not work right away then you need to repeat steps 3) and 4) (rather quickly) until it succeeds.

In the case of success, the 'socket\_pid\_attack' program will have replaced itself by the 'vdagent.py' Python program, which needs to be present in the same directory. The Python program will block, trying to receive data on the connected UNIX domain socket, because 'spice-vdagentd' can't accept the connection. Now terminate the program still running from step 1. via Ctrl-C to free up the blocked file descriptors in 'spice-vdagentd'. The malicious connection should now be accepted by 'spice-vdagentd' and the attacker should have become the active agent for the victim's graphical session. The additional output of the program from step 3. should look similar to this:

```
...
MessageType.VERSION arg1 = 0 arg2 = 0 bytes = 7
b'302e32302e3000'
sending resolution of 1024 768
MessageType.GRAPHICS_DEVICE_INFO arg1 = 0 arg2 = 0 bytes = 4
b'00000000'
MessageType.CLIENT_DISCONNECTED arg1 = 0 arg2 = 0 bytes = 0
MessageType.GRAPHICS_DEVICE_INFO arg1 = 0 arg2 = 0 bytes = 4
b'00000000'
MessageType.GRAPHICS_DEVICE_INFO arg1 = 0 arg2 = 0 bytes = 4
b'00000000'
MessageType.AUDIO_VOLUME_SYNC arg1 = 0 arg2 = 0 bytes = 7
b'01eb0200000000'
MessageType.AUDIO_VOLUME_SYNC arg1 = 0 arg2 = 0 bytes = 7
b'00eb0200000000'
MessageType.CLIPBOARD_RELEASE arg1 = 1 arg2 = 0 bytes = 0
sending clipboard grab request for ClipboardType.SELECTION_PRIMARY
...
```

If the attacker's connection is not considered to be part of the active session then the 'AUDIO\_VOLUME\_SYNC' and 'GRAPHICS\_DEVICE\_INFO' messages will not be received from the 'spice-vdagentd'. This can happen if some other user in the system received the malicious PID or if the PID was not in existence for long enough. If this is the case repeat steps 3. to 5. until it succeeds.

- If the previous steps all succeeded then you should be able to see for example the clipboard content from the host when a SPICE capable viewer is used to connect to the virtual machine. The attacker can also send back "malicious clipboard content" to the host.

# 4) Other Suggestions

## a) fchmod() instead of chmod() for UNIX domain socket

I made a minor suggestion about the 'chmod()' call in 'vdagentd.c:1223'. If possible this should be replaced by a call to 'fchmod()', if the file descriptor can be obtained from the glib functions. If custom, unsafe paths are used for the UNIX domain socket then symlink attacks might become possible through the use of 'chmod()'.

This was addressed by upstream in Patch 01.

# 5) Timeline

2020-09-18: I privately sent a version of this report and the reproducers to the upstream developer(s).

2020-10-16: Upstream shared their first version of the patches with me,

I reviewed them and we discussed some details.  
2020-10-26: I was asked to send this prenotification to the  
linux-distros mailing list to prepare the publication of the  
issues.

# 6) Further References

- SUSE bugzilla tracker bug [3]

[3]: [https://bugzilla.suse.com/show\\_bug.cgi?id=1173749](https://bugzilla.suse.com/show_bug.cgi?id=1173749)

Cheers

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