## Talos Vulnerability Report

TALOS-2020-1118

## Microsoft Azure Sphere AF\_AZSPIO socket memory corruption vulnerability

JULY 31, 2020

CVE NUMBER

CVE-2020-16970

Summary

A memory corruption vulnerability exists in the AF\_AZSPIO socket functionality of Microsoft Azure Sphere 20.05. A sequence of socket operations can cause a double-free and out-of-bounds read in the kernel. An attacker can write a shellcode to trigger this vulnerability.

Tested Versions

Microsoft Azure Sphere 20.05

Product URLs

https://azure.microsoft.com/en-us/services/azure-sphere/

CVSSv3 Score

8.1 - CVSS:3.0/AV:L/AC:H/PR:N/UI:N/S:C/C:H/I:H/A:H

CWE

CWE-415 - Double Free

Details

Microsoft's Azure Sphere is a platform for the development of internet-of-things applications. It features a custom SoC that consists of a set of cores that run both high-level and real-time applications, enforces security and manages encryption (among other functions). The high-level applications execute on a custom Linux-based OS, with several modifications to make it smaller and more secure, specifically for IoT applications.

The Azure Sphere has a procedure that utilizes special AF\_AZSPI0 sockets to communicate between high-level and real-time applications: https://docs.microsoft.com/en-us/azure-sphere/app-development/high-level-inter-app

As mentioned, in order to allow this connection to occur in the kernel, one must edit their high-level application's manifest to include: "AllowedApplicationConnections": [
"005180BC-402F-4CB3-A662-72937DBCDE47"], where by the value is the component\_id of the high-level app.

But regardless of if this entry is in the application manifest or not, a high-level app can still utilize an AF\_AZSPIO socket via the normal linux socket api:

```
#define AF_AZSPIO 43
int azspio_sock = socket(AF_AZSPIO, 0x80002, 0x0);
```

Just like any other socket, the creation is pretty simple, it's more the binding and sending addresses that's atypical:

The notes above in the comments must be followed in order for any bind, send, recvmsg, sendmsg, to work, and this can be done with the below code:

```
struct sockaddr_azspio sockaddr;
sockaddr.sa_family = AF_AZSPIO;
sockaddr.sa_port = 0x0;

memset(&sockaddr.sa_component_id,0x0,sizeof(sockaddr.sa_component_id));
// make sa_component_id match app's component id.
```

There's also a little complication on converting a component\_id from string form ("ComponentId": "b0000000-0000-0000-0000-000000000000") to the matching binary form, the first three groups of the component\_id are in little endian, whereas the last two are big endian. To illustrate, if I have a uint8\_t[16] component\_id, each index in that char array would correspond to the string as such:

```
03 02 01 00 - 05 04 - 07 06 - 08 09 - 0a 0b 0c 0d 0e 0f
```

```
memset(&sockaddr.sa_component_id,0x0,sizeof(sockaddr.sa_component_id));
sockaddr.sa_component_id[3] = 0xb0;
```

And likewise, then doing memcpy(ssockaddr.sa\_component\_id[6],"\xde\xad\xbe\xef",4); would result in a component id of b0000000-0000-adde-beef-000000000000. Moving on, this component\_id acts essentially as the IP address of the socket (with regards to other apps), and now that we have a valid sa\_component\_id, we can bind our socket:

```
ret = bind(azspio_sock, &sockaddr, sizeof(sockaddr));
```

This bind results in a few different things, but most importantly we hit the following code:

```
// kernel/net/azspio/azspio.c
if (!azspio_id_matches(component_id, &current_component_id)) {
     return -EACCES;
  .:∠apped 68
return 0;
}
   rc = azspio_socket_assign(ipc);
                                                               // [2]
     return rc;
  /* unbind previous, if any */
if (!zapped) // [3]
    azspio_socket_remove(ipc);
  [...]
  if (!rc) {
      sock_reset_flag(sk, SOCK_ZAPPED);
                                                              //[4]
  return rc;
}
```

At [1], the azspio\_socket\_assign function adds ipc->item into a static linked list (@azspio\_all\_sockets) via the kernel's list api (list\_add):

```
/* Add socket to global socket list. */
static int azspio_socket_assign(struct azspio_sock *ipc)
{
    mutex_lock(&azspio_socket_lock);
    list_add(&ipc->item, &azspio_all_sockets);
    mutex_unlock(&azspio_socket_lock);
    sock_hold(&ipc->sk);
    return 0;
}
```

Continuing on in \_\_azspio\_bind:

At [2], we see that if the SOCK\_ZAPPED flag is not set for our socket, the socket immediately gets removed, however since this is a socket that has not been bound before, this is not the case. The only place that an AF\_AZSPIO socket can have its SOCK\_ZAPPED flag unset is at [3], after it has successfully bound to a socket.

Now, something interesting to note about \_\_azspio\_bind further up:

At [1], [2], and [3], we can see the process of verifying the component\_id in a couple ways, the most interesting being [3], since if zapped is false (i.e. we have already bound the socket), as long as the component\_id we are binding with matches our application's component\_id, we continue on. Thus at [4], we can actually add the same socket into the list twice, and since zapped is set, we also immediately remove it from this list. Let us now look at such a scenario in memory, first with the declarations and structures in the kernel:

```
struct list_head {
    struct list_head *next, *prev;
};

#define LIST_HEAD_INIT(name) { &(name), &(name) }

#define LIST_HEAD(name) \
    struct list_head name = LIST_HEAD_INIT(name)
```

Thus, static LIST\_HEAD(azspio\_all\_sockets); results in:

```
[^~^]> x/2wx 6azspio_all_sockets
0xc0620458 <azspio_all_sockets>: 0xc0620458 0xc0620458
```

Moving on, after the first azspio\_socket\_assign:

We can see both the head [1] and our new socket [2] having prev and next point to the other as normal. The interesting part happens on the second add:

```
Run till exit from #0 azspio_socket_assign (ipc=<optimized out>) at net/azspio/azspio.c:515
__azspio_bind (sock=<optimized out>, addr=<optimized out>, zapped=0) at net/azspio/azspio.c:562
562 in net/azspio/azspio.c

[~.~]> x/2wx 6azspio_all_sockets
0xc0620458 <azspio_all_sockets>: 0xc18d11d8 0xc18d1d8 // [1]

[0.0]> x/2wx 0xc18d11d8
0xc18d11d8: 0xc18d11d8 0xc0620458 // [2]
```

Our initial list head has not changed, however sock->item->next now points to itself [2]. When the first list\_del happens, immediately after, further corruption happens:

```
Run till exit from #0 azspio_socket_remove (ipc=0xc18d1000) at net/azspio.c:524
0xc0360786 in __azspio_bind (sock=<optimized out>, addr=<optimized out>, zapped=0) at net/azspio/azspio.c:563
563 in net/azspio/azspio.c

[0.0]> x/2wx & & azspio_all_sockets
0xc0620458 < azspio_all_sockets>: 0xc18d11d8
0xc18d11d8

[>_>]> x/2wx & 0xc18d11d8
0xc18d11d8: 0x00000100 0x00000200
```

Due to the corrupted list, the \_\_list\_del at [1] actually does nothing, and then our socket's list members get poisoned with 0x200 and 0x100. This leaves us in an interesting position since, if azspio socket remove gets called again, we will end up dereferencing 0x100 and 0x200.

Assuming we close our socket now and destroy our last reference to the socket file descriptor, azspio\_release gets called, which will cause another azspio\_socket\_remove to happen, resulting in a crash, and it's worth noting that this also occurs if the process exits. Assuming that one could map the NULL page, this would more than likely result in a privilege escalation, but until a secondary vulnerability is found to map the NULL page, this vulnerability would only cause a denial-of-service.

Crash Information

Timeline

2020-07-02 - Vendor Disclosure 2020-07-31 - Public Release

CREDI

Discovered by Lilith >\_>, Claudio Bozzato and Dave McDaniel of Cisco Talos.

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