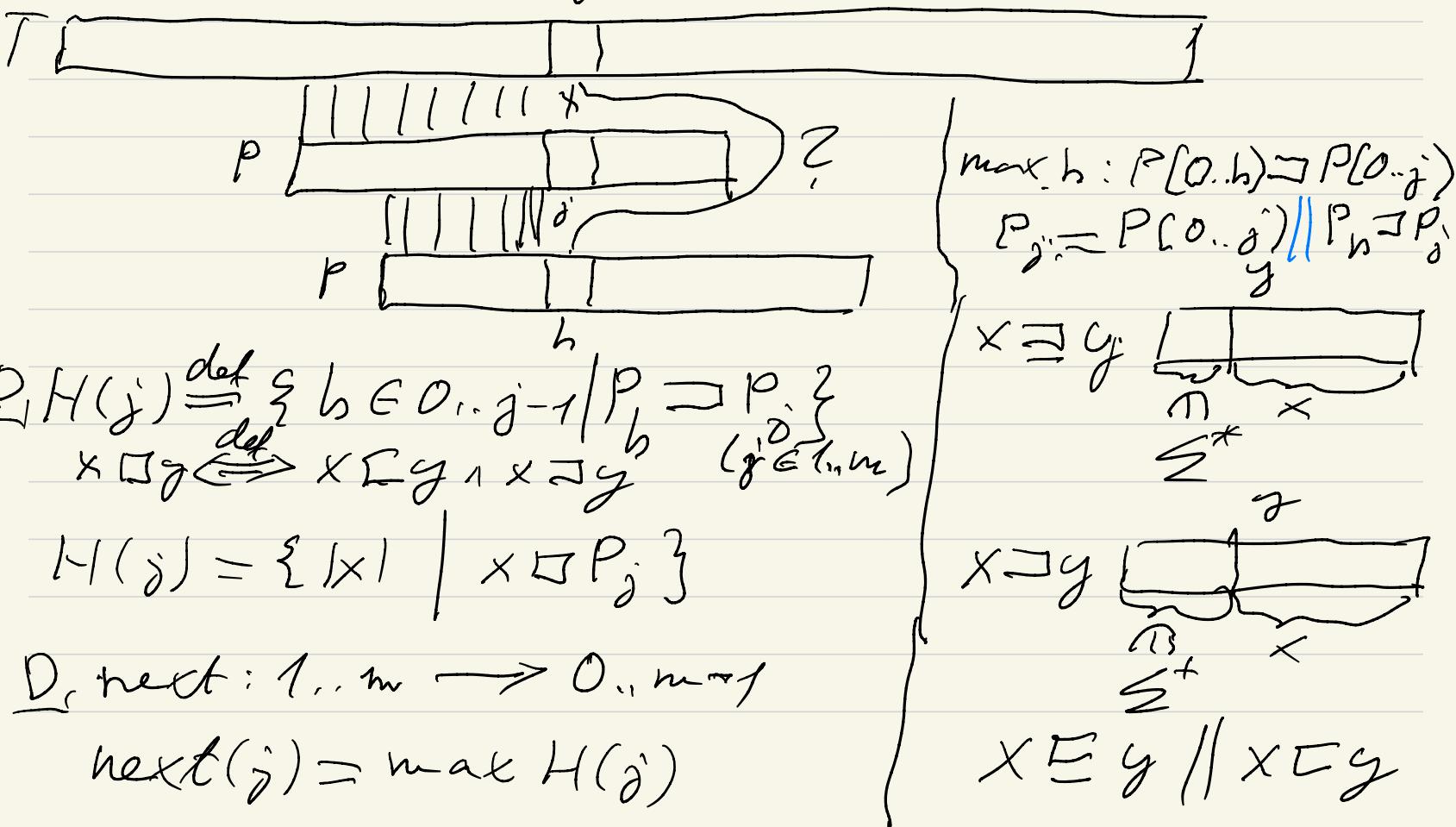
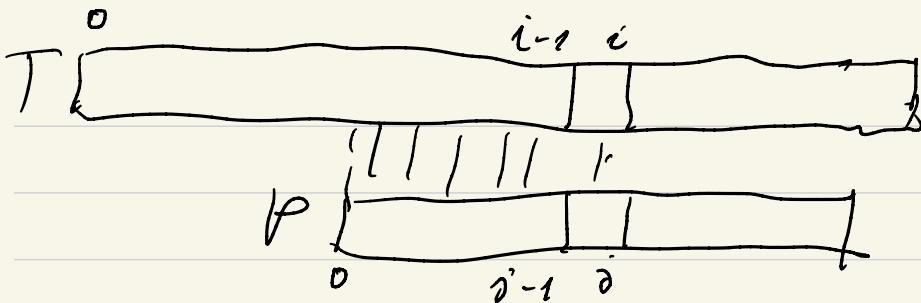


KMP algorithm





$$P_{j+1} \equiv T_{i+j}$$

↑
↓

$$P_j \equiv T_i \times P[j] = T[i]$$

Tac

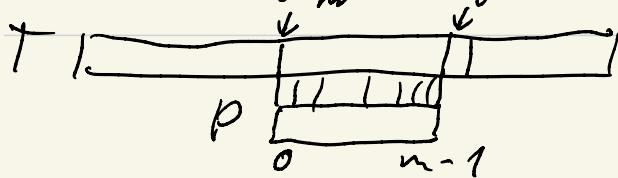
$$P_{j+1} \supseteq T_{i+j} \Leftrightarrow P_j \supseteq T_{i+1} P[j] = T[i]$$

$\text{next}(j) \in 0..j-1 \quad (j \in 1..m)$

$\text{next}(j+1) \leq \text{next}(j) + 1 \quad (j \in 1..m-1)$

$$P[j-1] = \boxed{\begin{matrix} B & A & B & A & B & B & A & B \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 \end{matrix}}$$

$$\text{next}(j) = \boxed{\begin{matrix} 0 & 0 & 1 & 2 & 3 & 1 & 2 & 3 \\ i-m & i & & & & & & \end{matrix}}$$



KMP($T: \Sigma^n$; $P: \Sigma^m$; $S: \mathbb{N} \{ \}$)

$\text{next}[1:N[m]]$; $\text{init}(\text{next}, P)$

$S := \{\}$; $i := j := 0$

$i < n$

$P[j] = T[i]$

$i++$; $j++$

$j = 0$

$j = m$

$S := S \cup \{i-m\}$ | $i++$

$j := \text{next}[j]$

$\exists i \forall P_j \geq T_i$

$1 \leq j \leq i \leq n$

$1 \leq j \leq m$

$t(i, j) := 2i - j \in [0..2n]$

Initially: $t(i, 0) = 0$

with each iteration: $t(i, j)$ strictly increases

\Rightarrow KMP: max. $2n$ iteration

KMP: max. n iteration
(initially: $i=0$ & finally: $i=n$,
 i increases max. by 1 in each iteration)

$mT(n, m), MT(n, m) \in \Theta(n) + \Theta(m)$ $mT_{init}(n), MT_{init}(m) \in \Theta(n)$

$$n \leq n$$

 $mT(n, m), MT(n, m) \in \Theta(n)$

(KMP($T : \Sigma[n] ; P : \Sigma[m] ; S : \mathbb{N}^{\{ \}}$))

$next/1 : \mathbb{N}[m] ; init(next, P)$

$S := \{ \} ; i := j := 0$

$i < n$

$P[j] = T[i]$

$i++ ; j++$

$j = 0$

$j = m$

$S := S \cup \{i - m\}$

SKIP

$j := next[m]$

$i++$

$j := next[j]$

$$\text{next}(1) = 0 \quad (1 \leq j < m)$$

$$\begin{aligned}
 \text{next}(j+1) &= \max H(j+1) = \\
 &= \max \{ h \in 0..j \mid P_h \supseteq P_{j+1} \} = \quad // i := h-1 \\
 &= \max \{ i+1 \in 0..j \mid P_{j+1} \supseteq P_{i+1} \} = \\
 &= \max \{ i+1 \in 1..j \mid P_{i+1} \supseteq P_{j+1} \} \cup \{ 0 \} = \\
 &= \max \{ i+1 \in 1..j \mid \underbrace{P_i \supseteq P_j}_{\text{if } P[i] = P[j]} \wedge P[i] = P[j] \} \cup \{ 0 \}.
 \end{aligned}$$

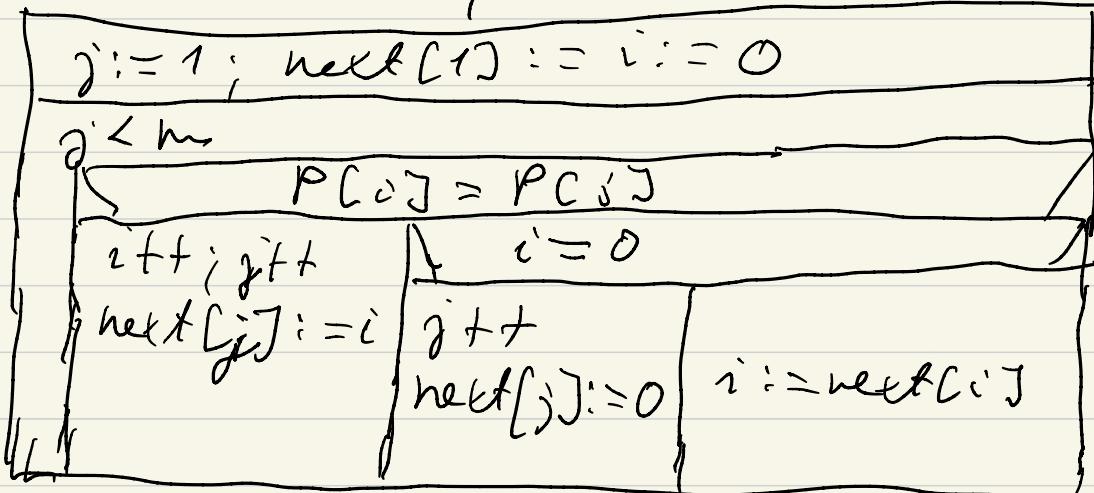
$$h(j) := \{ i \in 0..j-1 \mid \begin{array}{l} i \in H(j) \\ P[i] = P[j] \end{array} \}$$

$$h(j) \subseteq H(j)$$

$$\text{next}(j+1) = \begin{cases} \max h(j) + 1 & \text{if } h(j) \neq \emptyset \\ 0 & \text{if } h(j) = \emptyset \end{cases}$$

$$\begin{aligned}
 \boxed{\max_{\ell+1} H(j)} &= \text{next}(\max_\ell H(j)) \\
 (\max_\ell H(j)) &= \text{next}(j)
 \end{aligned}$$

init(next[1:N[m]; P:Σ[m])



Invariant

$$\text{next}[1..j] = \text{next}[1..j']$$

$$1 \leq i < j \leq m$$

$$\wedge P_i \supseteq P_j \quad (i \in H(s))$$

To prove: $mT(m), MT_{\underset{\text{init}}{\text{init}}}(m) \in \Theta(m)$

See in the printed lecture notes !!!

init(next[1:N[m]]; P:Σ[m])

$j := 1; \text{next}(1) := i := 0$
 $j < m$
 $P[0] = P[i]$
 $i++;$
 $\text{next}[j] := i$
 $i = 0$
 $j++$
 $\text{next}[j] := 0$
 $i := \text{next}[i]$

$PL[i-1] =$	B	A	B	A	B	B	A	B
$j^* =$	1	2	3	4	5	6	7	8
$nextC[j] =$	0	0	1	2	3	1	2	3