Chapter 02 The Physical Layer

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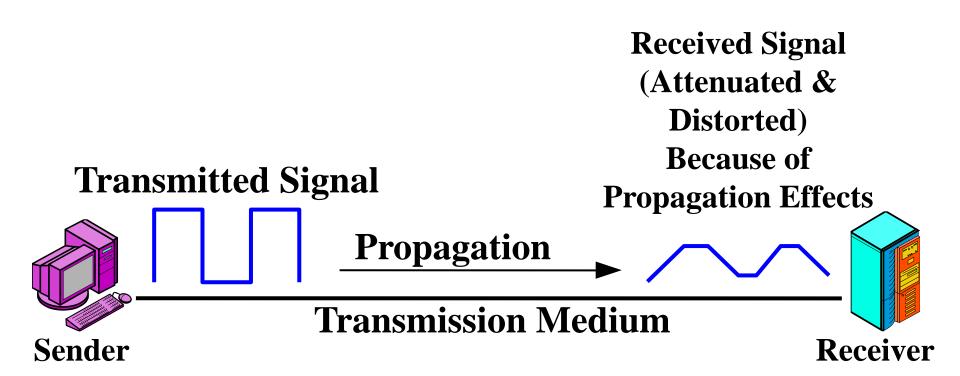**Beijing Institute of Technology**

Key Points

物理层	功能与作用	熟练掌握
	接口特性	熟练掌握
通信基础	基本概念	熟练掌握
	奈奎斯特定理，香农定理	熟练掌握
	编码与调制	熟练掌握
	多路复用	熟练掌握
传输介质	种类，特点	掌握
物理层设备	中继器	掌握
	集线器	掌握

Chapter 2: Roadmap

Physical Layer

Transmitted Signal

Propagation

Transmission Medium

Received Signal (Attenuated & Distorted) Because of Propagation Effects

Sender

Receiver

Physical Layer

- **The lowest layer in any network architecture model**

- **Concerned with the transparent transmission of "raw" bits across a communications medium**

- **Deals with the physical characteristics (mechanical, electrical, functional, procedural) of data transmission and communication.**

Physical Layer

- **Responsible for:**
  - **providing basic signaling (control, data)**
  - **signal modulation**
  - **encoding/decoding**
  - **activate/deactivate physical medium (PM)**
  - **bit-timing (clocking)**
  - **mapping between different formats**
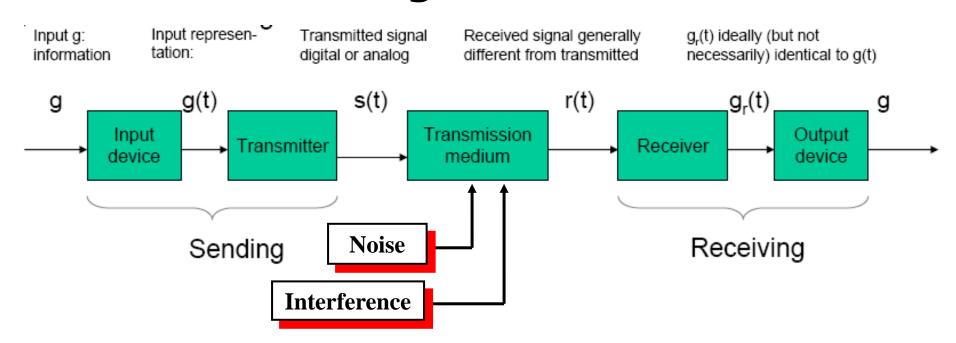
Chapter 2: Roadmap

The Basis for Data Communication

- **What is communication?**

  The transmission of information from one point to another through a succession of certain processes.
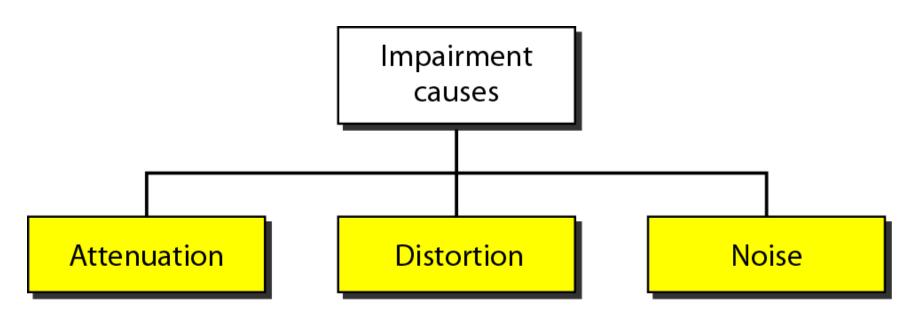
The Basis for Data Communication

- ## Important transformations:
  - ☐ **Source coding**
  - ☐ **Channel coding**

| Input g: information | Input represen-tation: | Transmitted signal digital or analog | Received signal generally different from transmitted | $g_r(t)$ ideally (but not necessarily) identical to g(t) |

g   g(t)   s(t)   r(t) $g_r(t)$ g

| Input device | Transmitter | Transmission medium | Receiver | Output device |

Sending

Noise
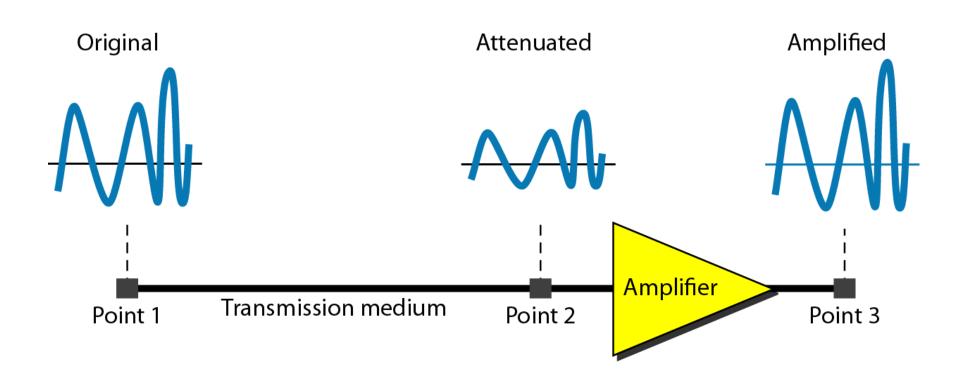
Interference

Receiving

The basic model of a data comm. system
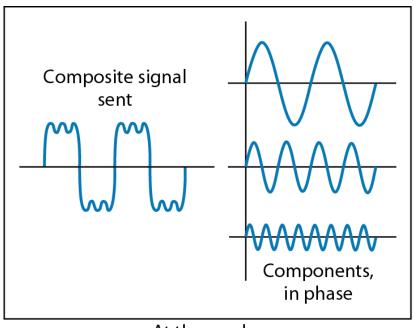
Transmission Impairment

- **Signals travel through transmission media, which are not perfect.**
- **Three causes of impairment are attenuation, distortion, and noise.**
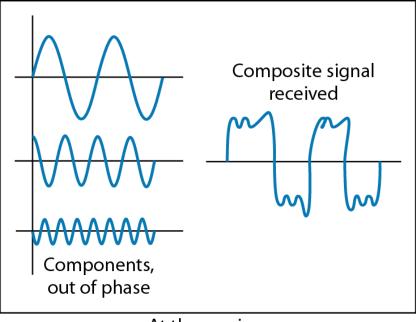
Attenuation: loss of energy

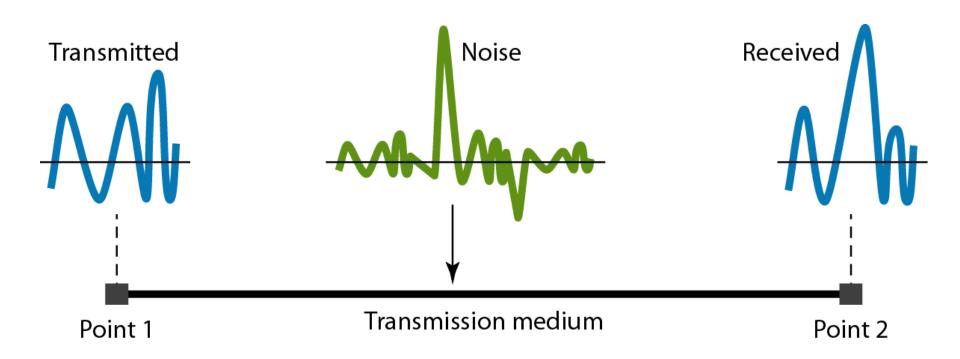# Distortion

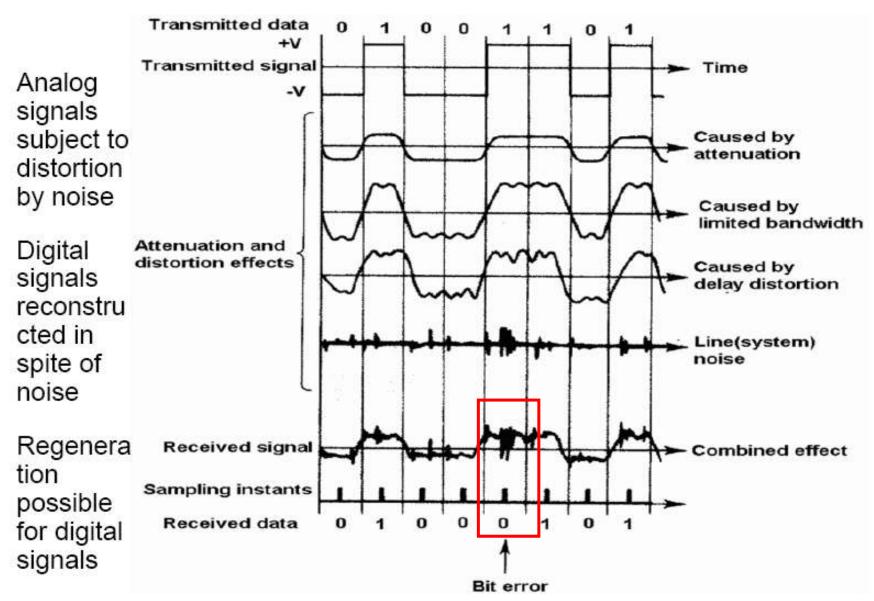At the sender   At the receiver

This is because each frequency signal has its own propagation speed through a medium

Noise

Transmission of Complex Signals

Analog signals subject to distortion by noise

Digital signals reconstructed in spite of noise

Regeneration possible for digital signals

Fourier Analysis

傅里叶，J.-B.-J.

法国数学家。**1768年3月21日**生于奥塞尔，**1830年5月16日**卒于巴黎。**1795**年曾在巴黎综合工科学校任讲师。 **1798**年随拿破仑远征埃及，当过埃及学院的秘书。**1801**年回法国，又任伊泽尔地区的行政长官。**1817**年傅里叶被选为科学院院士，并于**1822**年成为科学院的终身秘书。**1827**年又当选为法兰西学院院士。

Fourier Analysis

- **A periodic function with period *T* (and frequency *f* = 1/*T*), *g*(*t*) can be written as:**

$$g(t) = \frac{1}{2}c + \sum_{n=1}^{\infty} a_n \sin(2\pi nft) + \sum_{n=1}^{\infty} b_n \cos(2\pi nft)$$

$$c = \frac{2}{T}\int_0^T g(t)\,dt$$

$$a_n = \frac{2}{T}\int_0^T g(t)\sin(2\pi nft)\,dt$$

$$b_n = \frac{2}{T}\int_0^T g(t)\cos(2\pi nft)\,dt$$

***f* = 1/*T* is the fundamental frequency**

Fourier Analysis

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$$g(t) = \frac{1}{2}c + \sum_{n=1}^{\infty} a_n \sin(2\pi nft) + \sum_{n=1}^{\infty} b_n \cos(2\pi nft)$$

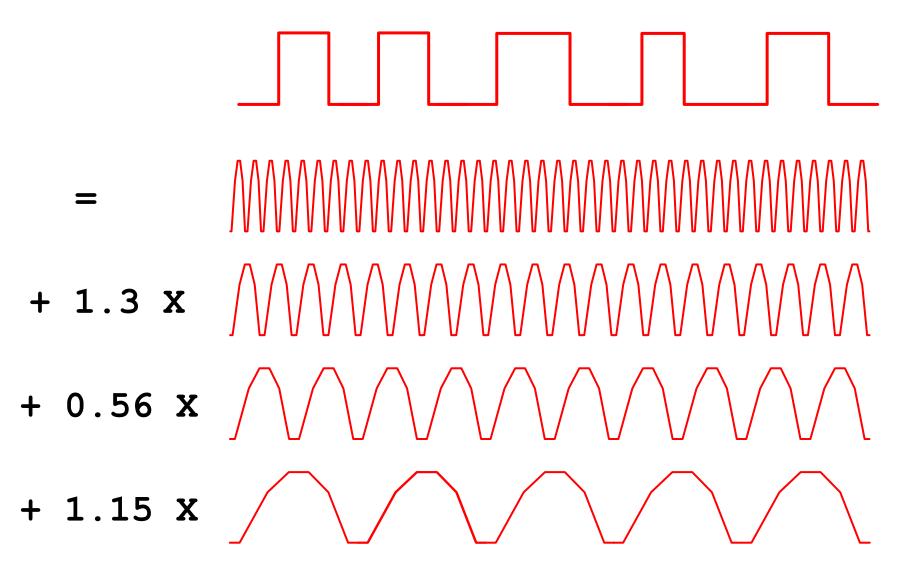$$\int_{-\pi}^{\pi} \cos nx\,dx = 0 \qquad (n = 1, 2, 3, \ldots)$$

$$\int_{-\pi}^{\pi} \sin nx\,dx = 0 \qquad (n = 1, 2, 3, \ldots)$$

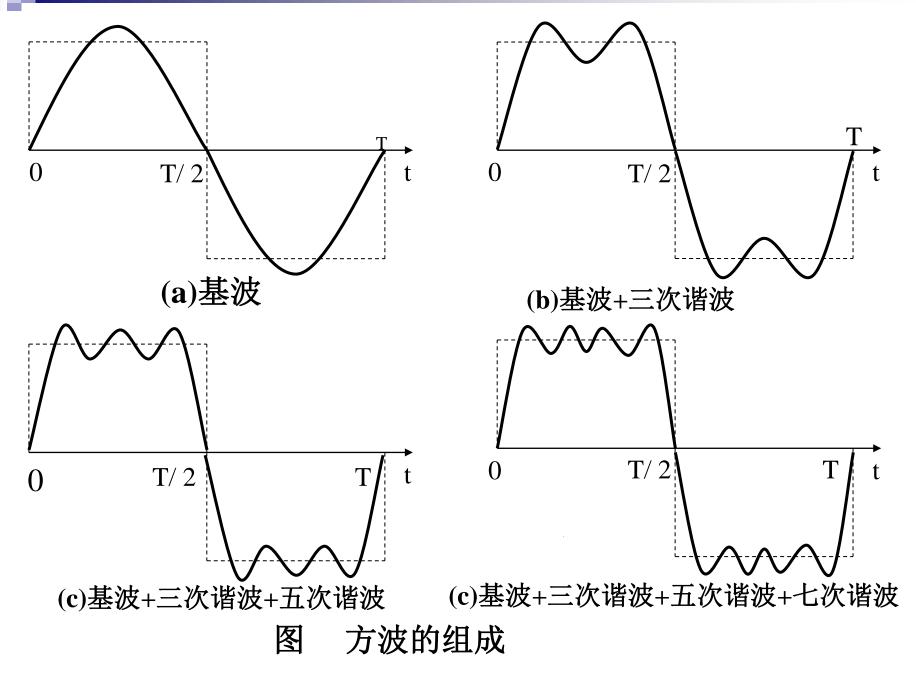$$\int_{-\pi}^{\pi} \sin kx \cdot \cos nx\,dx = 0 \qquad (k, n = 1, 2, 3, \ldots)$$

$$\int_{-\pi}^{\pi} \cos kx \cdot \cos nx\,dx = 0 \qquad (k, n = 1, 2, 3, \ldots ; k \neq n)$$

$$\int_{-\pi}^{\pi} \sin kx \cdot \sin nx\,dx = 0 \qquad (k, n = 1, 2, 3, \ldots ; k \neq n)$$
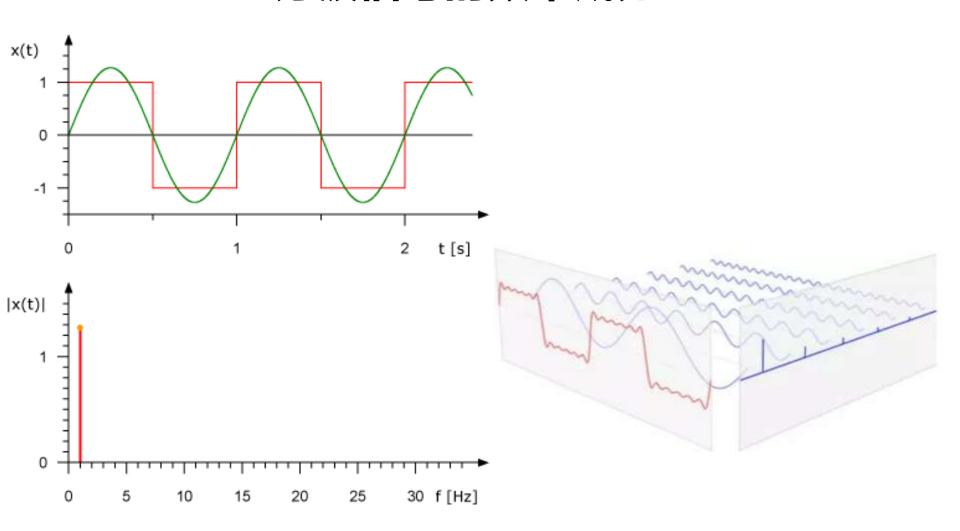
f = 1/

funda

Signal = Sum of Waves

=

+ 1.3 x

+ 0.56 x

+ 1.15 x

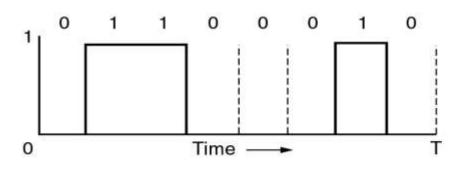(a)基波
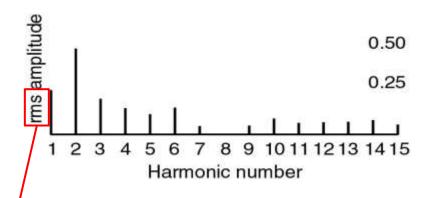
(b)基波+三次谐波

(c)基波+三次谐波+五次谐波

(c)基波+三次谐波+五次谐波+七次谐波

图　方波的组成

方波信号的频率成分

Bandwidth-Limited Signals

■ **Example: 01100010, 8 bit for ASCII "b"**

a_n =1/πn[cos(πn/4) -cos(3πn/4) + cos(6πn/4) - cos(7πn/4)]

b_n = 1/πn[sin(3πn/4) -sin(πn/4) + sin(7πn/4) - sin(6πn/4)]

c_n = 3/4
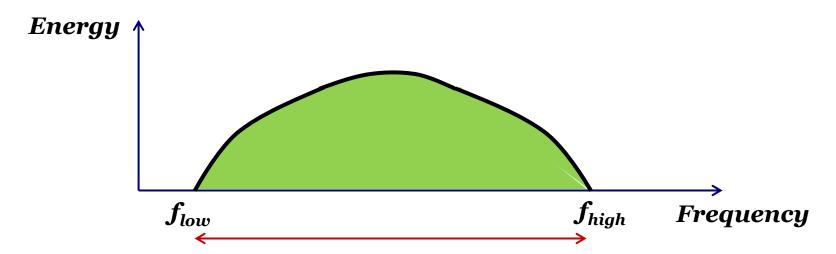
The root mean square amplitude is $\sqrt{a_n^2 + b_n^2}$

与对应频率处的能量成正比
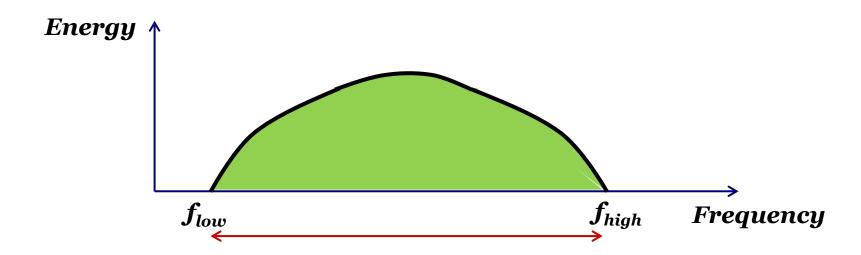
Bandwidth-Limited Signals

- **In theory, signal spectrum spreads from −∞ to ∞**

- **Practically, most of signal energy (~95%) in the spectrum is contained in a finite range of frequencies.**

A bandwidth-limited signal:
a signal who has a finite spectrum.

Bandwidth-Limited Signals
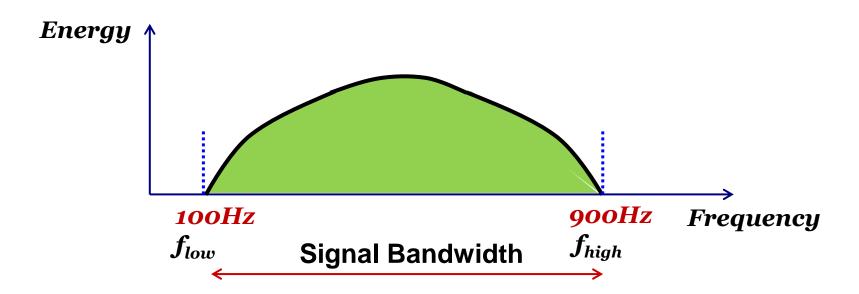
It means:
- **The quality of transmission is frequency dependent.**
- **Not all parts of the spectral components of a signal get through the channel as you would expect.**

Signal Bandwidth (Hz)
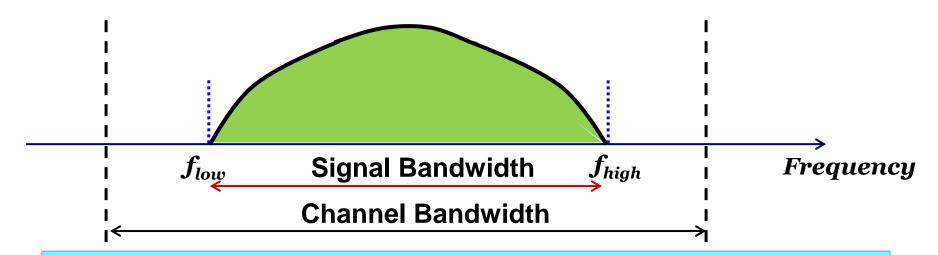
- **The difference between upper and lower frequency limits of the signal.**

$$B = f_{high} - f_{low} = 900 - 100 = 800 Hz$$

Chanel Bandwidth (Hz)

- **Chanel bandwidth** is the range of frequencies that the channel can carry.

The channel bandwidth diagram showing f_{low} and f_{high} with Signal Bandwidth and Channel Bandwidth, versus Frequency.

The channel bandwidth should always be **greater** than the bandwidth of the signal to be transmitted else loss of information takes place.

Chanel Bandwidth (Hz)

- **Varied forms of transmission media have different bandwidths.**

Type of the channel	Frequency range (Approx.)
Twisted pair	1MHz – 600 MHz
Coaxial cable	0 – 750 MHz
Microwave	1 GHz-30 GHz
Satellite	1 GHz – 40 GHz
Fibre optics	180 THz – 330 THz

Baud-rate and Bit-rate

- **baud rate (Signal/Symbol rate):**

  The number of signal or symbol (one of several voltage, frequency, or phase changes) changes per second.

  B=1/T **(T is the frequency period of signal or symbol)**

- **bit rate (Data rate):**

  The number of bits transmitted per second.

  b/s or bps (bits per second)

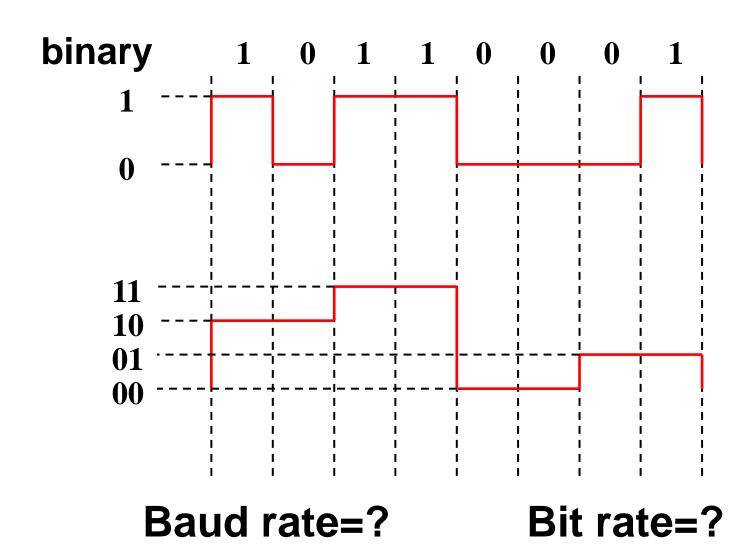# Relationship between baud-rate and bit-rate

Bit rate = Baud rate x the number of bit per baud.

■ $S = B * \log_2 v$

S is bit rate,

B is baud rate,

v is number of signal values.

■ **Ex:**

B=9600, v=2, s=9600 bps;

B=2400, v=16, s=9600 bps

Relationship between baud-rate and bit-rate

binary   1   0   1   1   0   0   0   1

1

0

11
10
01
00

Baud rate=?   Bit rate=?

Difference Between Bit Rate And Baud Rate

	Bit Rate	Baud Rate
1.	Bit rate is defined as the transmission of number of bits per second.	Baud rate is defined as the number of signal units per second.
2.	Bit rate is also defined as per second travel number of bits.	Baud rate is also defined as per second number of changes in signal.
3.	Bit rate emphasized on computer efficiency.	While baud rate emphasized on data transmission.
4.	The formula of Bit Rate is:= baud rate x the number of bit per baud	The formula of Baud Rate is:= bit rate / the number of bit per baud
5.	Bit rate is not used to decide the requirement of bandwidth for transmission of signal.	While baud rate is used to decide the requirement of bandwidth for transmission of signal.

Relationship between bandwidth and data rate

- **In theory, bandwidth is related to data rate by:**
  - **Nyquist formula**
  - **Shannon formula**

Nyquist formula

- **For a noiseless channel, the theoretical maximum bit rate is**

$$\text{Max. baud rate} = 2H \text{ (baud)}$$
$$\text{Max. bit rate} = 2H \log_2 V \text{ (bps)}$$

(where H is the bandwidth, V is the number of signal values)

Bandwidth is a fixed quantity, so it cannot be changed. Hence, the data rate is directly proportional to the number of signal levels.

Nyquist proved that if an arbitrary signal has been run through a low-pass filter of bandwidth, the filtered signal can be completely reconstructed by making only 2*Bandwidth samples per second.

Nyquist formula Examples

- **1, Consider a noiseless channel with a bandwidth of 3000 Hz transmitting a signal with four signal levels. What can be the maximum bit rate?**

  Bit Rate = 2 * 3000 * $\log_2(4)$ = 12000bps

- **2, We need to send 256 kbps over a noiseless channel with a bandwidth of 20 kHz. How many signal levels do we need?**

  $256000 = 2 * 20000 * \log_2(V)$

  $\log_2(V) = 6.4$

  $V = 2^{6.4} = 84.45$ levels

Shannon formula

- **For a noisy channel with a signal-to-noise ratio S/N, the theoretical maximum bit rate is**

$$\text{Max. bit rate} = H \log_2 (1+S/N) \text{ (bps)}$$

(where H is the bandwidth)

Bandwidth is a fixed quantity, so it cannot be changed. Hence, the channel capacity is directly proportional to the power of the signal, as SNR = (Power of signal) / (power of noise).

Shannon formula

- **The signal-to-noise ratio (S/N) is usually expressed in decibels (dB) given by the formula:**

$$10*\log_{10}(S/N)$$

- **S/N = 10, $10*\log_{10}(10) = 10$ dB**
- **S/N = 1000, $10*\log_{10}(1000) = 30$ dB**

Shannon formula Examples

- **1, We have a channel with a 2-MHz bandwidth. Assume that SNR = 36dB. and the channel bandwidth is 2 MHz. What is the theoretical channel capacity?**

$$S/N = 10^{36/10} = 10^{3.6} = 3981$$

$$C = 2 * 10^6 * \log_2(1 + 3981) = 24\text{Mbps}$$

Shannon formula Examples

- **2, We have a channel with a 1-MHz bandwidth. The SNR for this channel is 63. What are the appropriate bit rate and signal level?**

$$C = 1 * 10^6 * \log_2(1 + 63) = 6\text{Mbps}$$

- **Then we use the Nyquist formula to find the number of signal levels. If we choose something lower, 4 Mbps, for example.**

$$4\text{Mbps} = 2 * 1 \text{ MHz} * \log_2(V)$$

$$V = 4$$

Shannon formula Examples

- **The Shannon capacity gives us the upper limit;**
- **the Nyquist formula tells us how many signal levels we need.**

Difference between Bandwidth and Data Rate

Bandwidth	Data Rate
It is the number of bits per second that a link can send or receive.	It is the speed of data transmission.
Normally it is measured in Hz, bps, Mbps or Gbps.	It is normally measured in Mbps or MBps.
It refers to the maximum data transmission speed.	It refers to the actual data transmission speed.
It is the potential of the data that is to be transferred in a specific period of time.	It is the amount of data transmitted during a specified time period over a network.
It is physical layer property in OSI model.	While it is common in all layers.
It shows the capacity of the channel.	It shows the present speed of data transmission.
It does not depend on properties of sender or receiver.	While it gets affected by sender or receiver.

Chapter 2: Roadmap

Copper Wires

- **Twisted pair**
  - ☐ **Shielded Twisted Pair (STP)**
  - ☐ **Unshielded Twisted Pair (UTP)**

Categories of Unshielded Twisted Pair

The EIA/TIA (Electronic Industry Association/Telecommunication Industry Association) has established standards of UTP.

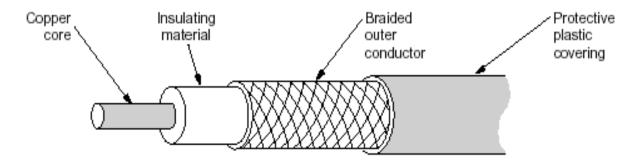| Type	Use
Category 1	1MHz Voice Only (Telephone Wire)
Category 2	Data to 4 MHz (4Mbps LocalTalk)
Category 3	Data to 16 MHz (10Mbps Ethernet)
Category 4	Data to 20 MHz (16Mbps Token Ring)
Category 5	Data to 100 MHz (100Mbps)
Category 5e	Data to 100 MHz (1Gbps)
Category 6	Data to 250 MHz (10Gbps)
Category 7	Data to 600 MHz (10Gbps)

Copper Wires

- ## Coaxial Cable

Copper core — Insulating material — Braided outer conductor — Protective plastic covering
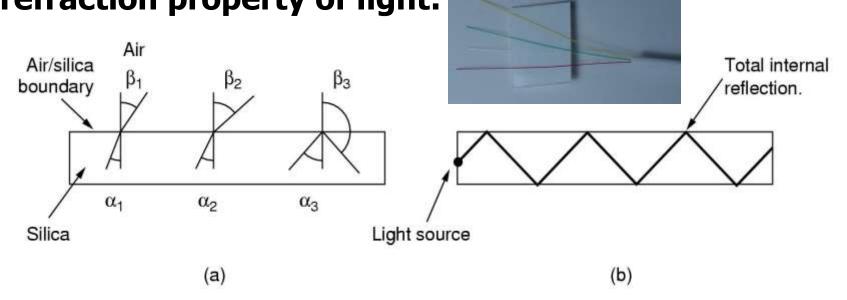
Baseband Coax

- ☐ **50-ohm cable for digital transmission**
- ☐ **10Base-2, BNC, Thin-LAN, 185m/per segment**
- ☐ **10Base-5, AUI, Thick-LAN, 500m/ per segment**
- ☐ **At most 5 segments, up to 945m/2500m.**

- ## Broadband Coax

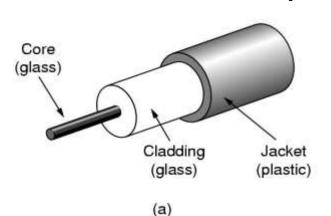- ☐ **75-ohm cable for analog transmission, like cable TV.**
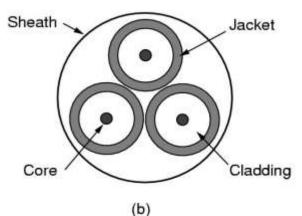
Fiber Optics

- **Principle:** optical signals that are passed through optical fiber. Principal working is based on the refraction property of light:

(a) Three examples of a light ray from inside a silica fiber impinging on the air/silica boundary at different angles.

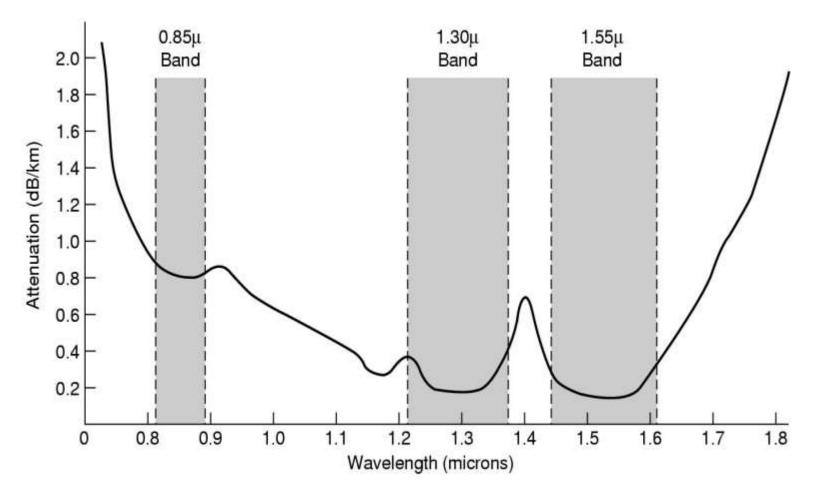(b) Light trapped by total internal reflection.

As it turns out, attenuation is extremely well in optical fiber. This means that they can be used for long distances. In addition, the bandwidth is enormous.
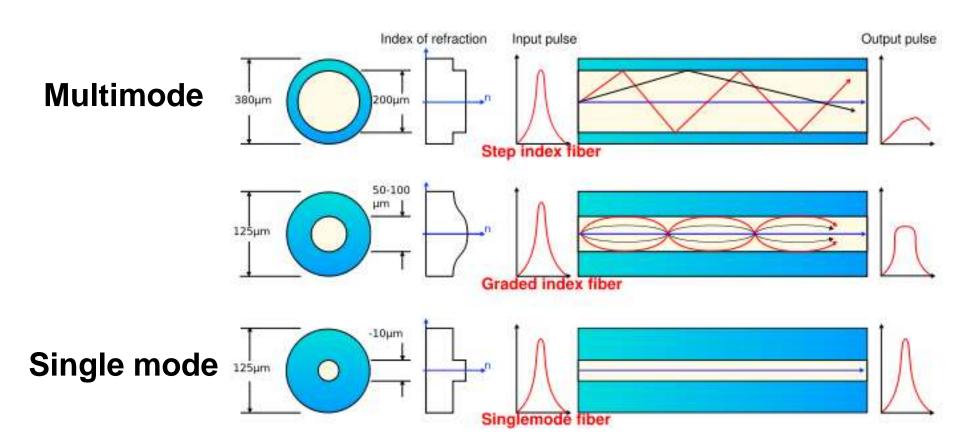

Item	LED	Semiconductor Laser
Data rate	Low	High
Mode	Multimode	Multimode or single mode
Distance	Short	Long
Lifetime	Long life	Short life
Temperature sensitivity	Minor	Substantial
Cost	Low cost	Expensive

- **Attenuation in decibels = $10\log_{10}$(transmitted power/received power)**

Attenuation of light through fiber in the infrared region.

Fiber Types

Multimode

Single mode

The electromagnetic spectrum

- **Wireless transmissions travel at the speed of light (c), uses a frequency (f) which has a wavelength (l). The relation is that:**
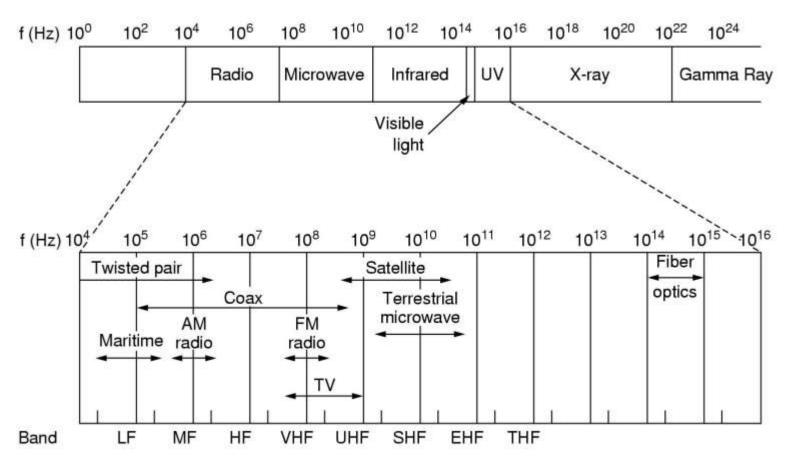
$$c = \lambda * f$$

c is a constant, approximately 3 x 10^8m/sec, that is theoretical in vacuum.

In practice, c is about 2 x 10^8m/sec either in fiber or copper. That is **200m/µsec**. (IMPORTANT!)

The larger the wavelength is, the longer the distance it can travel without attenuation. Also, the dispersion of higher frequencies is much lower.

The electromagnetic spectrum and its uses for communication

Microwave Transmission

- **Waves travel in straight lines and can be narrowly focused above 100MHz, concentrating energy into a small beam using dish-like antenna.**

- **Repeaters are needed periodically, for 100m high towers, repeaters can be spaced 80km apart.**

- **Multipath-fading is weather and frequency dependent.**

- **Frequency range: 2.4GHz-2.484GHz, no FCC (federal communication commission) licensing needed.**
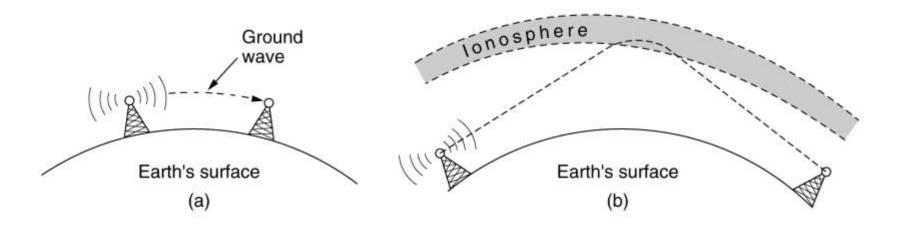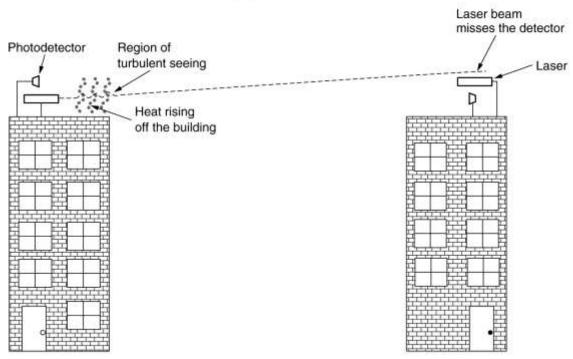
Fig. 2-12. (a) In the VLF, VF, and MF bands, radio waves follow the curvature of the earth. (b) In the HF they bounce off the ionosphere.

- **Microwave transmission is also popular and is good for long distances, as long as it's directed.**

- **Problem is the density in the spectrum, requiring higher frequency ranges (which are hard for unguided transmissions)**

Infrared and Millimeter Waves Light Wave Transmission
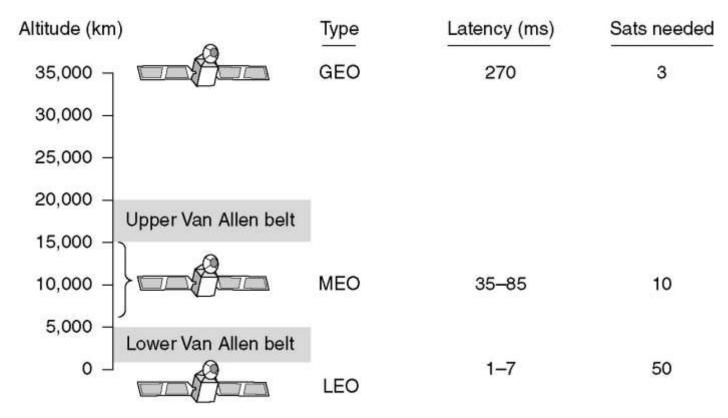

Convection currents can interfere with laser communication systems. A bidirectional system, with two lasers, is pictured here.

Communication Satellites

- **Geostationary Satellites(同步卫星)**
- **Medium-Earth Orbit Satellites**
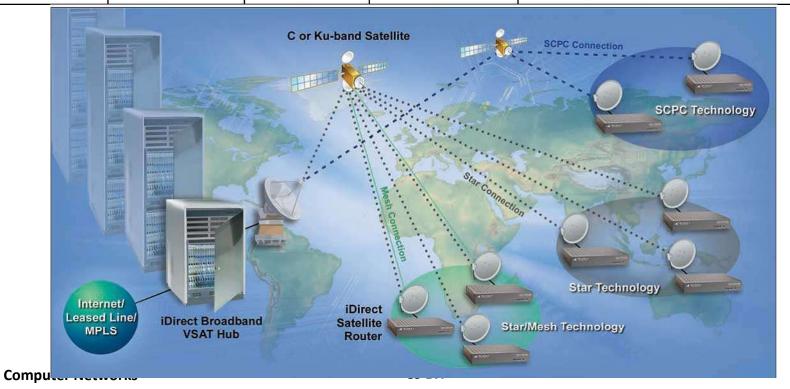- **Low-Earth Orbit Satellites**
- **Satellites versus Fiber**

Communication Satellites

Altitude (km)		Type	Latency (ms)	Sats needed
35,000		GEO	270	3
30,000				
25,000				
20,000				
	Upper Van Allen belt			
15,000				
10,000		MEO	35–85	10
5,000				
	Lower Van Allen belt			
0		LEO	1–7	50

- **Communication satellites and some of their properties, including altitude above the earth, round-trip delay time and number of satellites needed for global coverage.**

The principal satellite bands

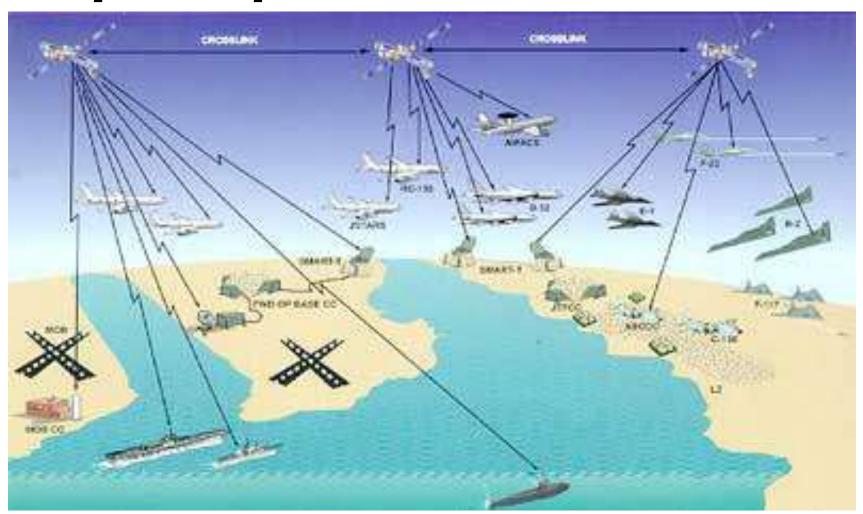| Band	Downlink	Uplink	Bandwidth	Problems
L	1.5 GHz	1.6 GHz	15 MHz	Low bandwidth; crowded
S	1.9 GHz	2.2 GHz	70 MHz	Low bandwidth; crowded
C	4.0 GHz	6.0 GHz	500 MHz	Terrestrial interference
Ku	11 GHz	14 GHz	500 MHz	Rain
Ka	20 GHz	30 GHz	3500 MHz	Rain, equipment cost

The principal satellite bands

VSATs using a hub
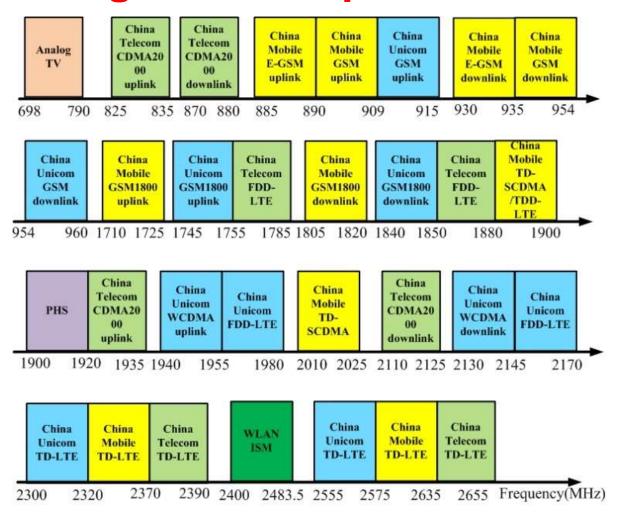**VSAT: very small aperture terminals**

The principal satellite bands

美国" 军事星系统"工作示意图

Radio spectrum Allocation

Licencing of the frequencies

Radio spectrum Allocation

ISM means Industrial, Scientific and Medical frequency band.

ISM band for unlicensed use

Frequency range		Bandwidth	Center frequency
6.765 MHz	6.795 MHz	30 kHz	6.780 MHz
13.553 MHz	13.567 MHz	14 kHz	13.560 MHz
26.957 MHz	27.283 MHz	326 kHz	27.120 MHz
40.660 MHz	40.700 MHz	40 kHz	40.680 MHz
433.050 MHz	434.790 MHz	1.84 MHz	433.920 MHz
902.000 MHz	928.000 MHz	26 MHz	915.000 MHz
2.400 GHz	2.500 GHz	100 MHz	2.450 GHz
5.725 GHz	5.875 GHz	150 MHz	5.800 GHz
24.000 GHz	24.250 GHz	250 MHz	24.125 GHz
61.000 GHz	61.500 GHz	500 MHz	61.250 GHz
122.000 GHz	123.000 GHz	1 GHz	122.500 GHz
244.000 GHz	246.000 GHz	2 GHz	245.000 GHz

Respect laws of your country regarding EMI and the maximum TX power allowed per band

15

Chapter 2: Roadmap

Data and Signals

Note

To be transmitted, data must be transformed to electromagnetic signals.

Analog and Binary Data

Data can be analog or digital

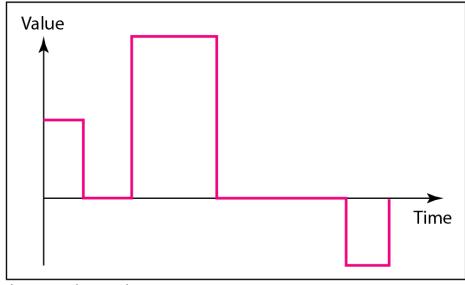| Analog Data	Binary Data
	11010110000111100101
Smoothly changing among an infinite number of states (loudness levels, etc.)	**Two states:** One state represents 1 The other state represents 0

Analog and Binary Signals

Signals can be analog or digital.

Analog signals can have an infinite number of values in a range

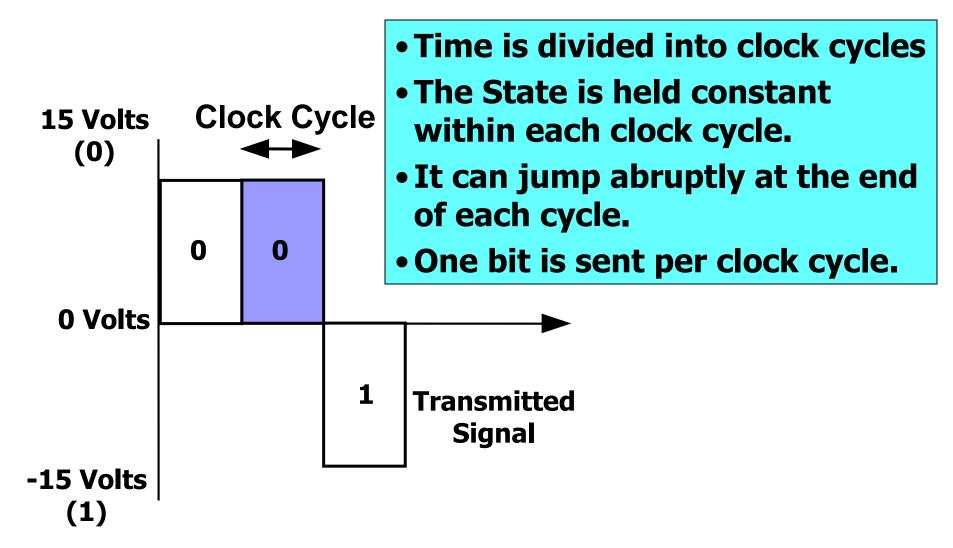**Digital signals can have only a limited number of values.**

a. Analog signal

b. Digital signal

Binary Data and Binary Signal

15 Volts (0)

Clock Cycle

- **Time is divided into clock cycles**
- **The State is held constant within each clock cycle.**
- **It can jump abruptly at the end of each cycle.**
- **One bit is sent per clock cycle.**

0 Volts

0   0

1   **Transmitted Signal**

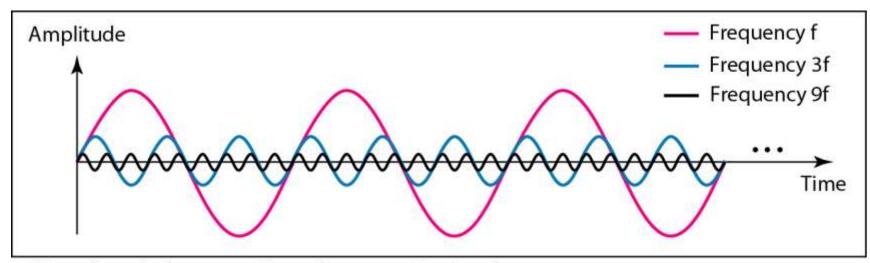**-15 Volts (1)**

Analog and Binary Signals

Note

In data communications, we commonly use periodic analog signals and nonperiodic digital signals.

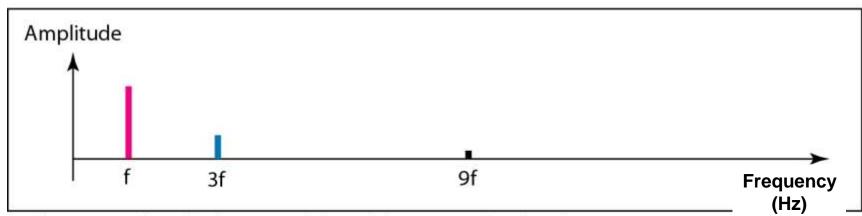*Periodic analog signal is used as data carrier (such as AM/FM radio)*

Periodic analog signals

- **Periodic analog signals can be classified as simple or composite.**

- **A simple periodic analog signal, a sine wave, cannot be decomposed into simpler signals.**

- **A composite periodic analog signal is composed of multiple sine waves.**
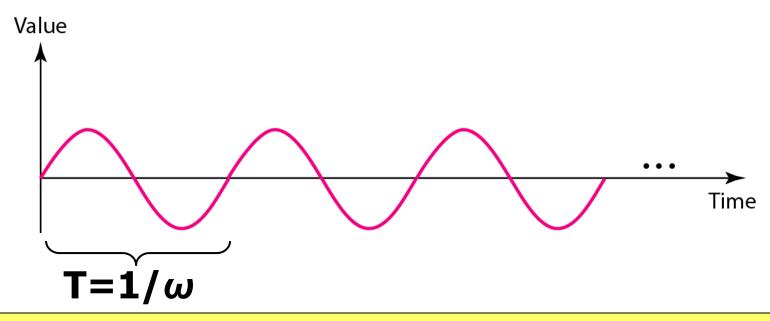
Periodic analog signals

a. Time-domain decomposition of a composite signal

b. Frequency-domain decomposition of the composite signal
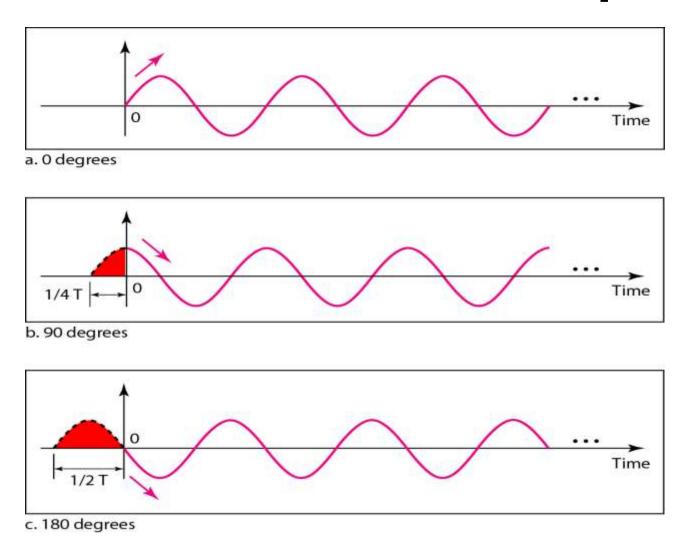
A sine wave

$$A\,sin\,(\omega t\,+\varphi)$$

A : amplitude = maximum strength of signal (V)

ω **: frequency = rate of change of signal or**
cycles per second (Hz)

period = time for one repetition (T)

φ **: phase = relative position in time**

3 sine waves with different phases

a. 0 degrees

b. 90 degrees

c. 180 degrees

Phase unit: degree (360°) or radians (2π rad)
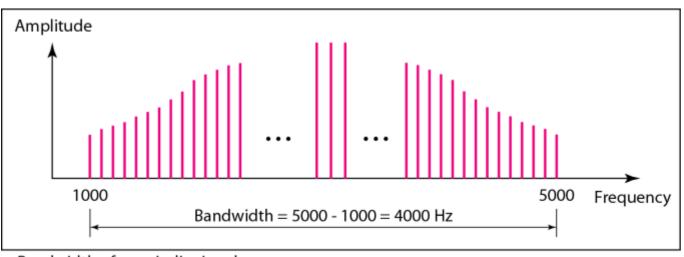
Periodic analog signals

Note

A **single**-frequency sine wave **is not** useful in data communications; we need to send a **composite** signal, a signal made of many simple sine waves.

Note

According to Fourier analysis, any composite signal is a combination of simple sine waves with different frequencies, amplitudes, and phases.
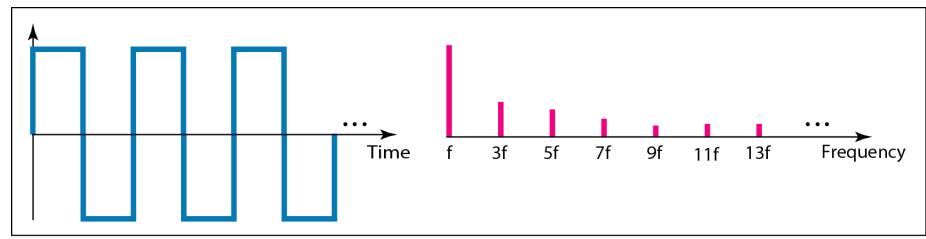
Bandwidth

a. Bandwidth of a periodic signal

The **bandwidth** of a composite signal is the difference between the highest and the lowest frequencies contained in that signal.
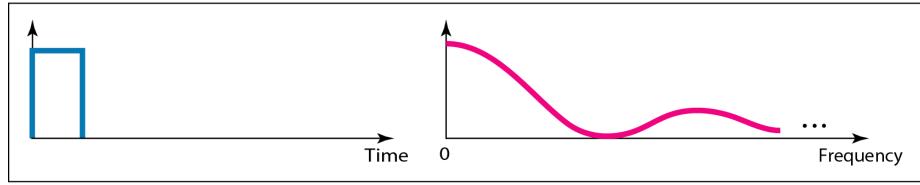
Digital Signals

- **In addition to being represented by an analog signal, information can also be represented by a digital signal.**

- For example, a 1 can be encoded as a positive voltage and a 0 as zero voltage.

- **A digital signal can have more than two levels. In this case, we can send more than 1 bit for each level.**

Digital Signals

a. Time and frequency domains of periodic digital signal
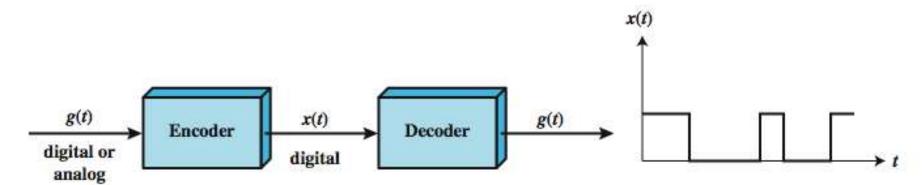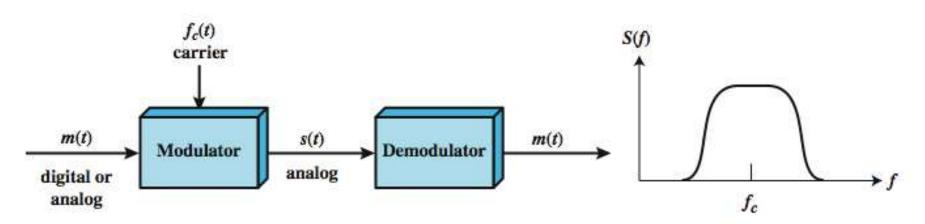
Periodic digital signal

b. Time and frequency domains of nonperiodic digital signal

Nonperiodic digital signal

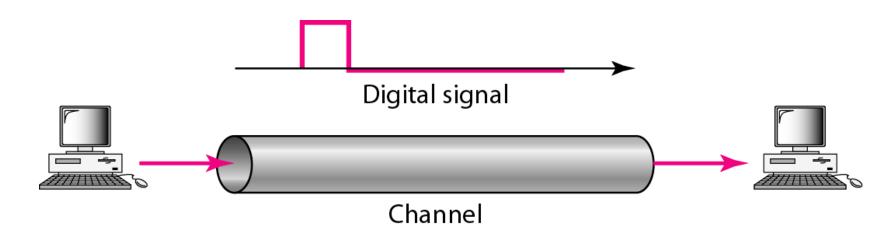# Encoding and Modulation Techniques

(a) Encoding onto a digital signal

(b) Modulation onto an analog signal

Encoding and Modulation Techniques

- **There are two types of Digital Data Transmission:**

  1) Base-Band data transmission
  - Uses low frequency carrier signal to transmit the data

  2) Band-Pass data transmission
  - Uses high frequency carrier signal to transmit the data

Baseband Transmission

Digital signal

Channel
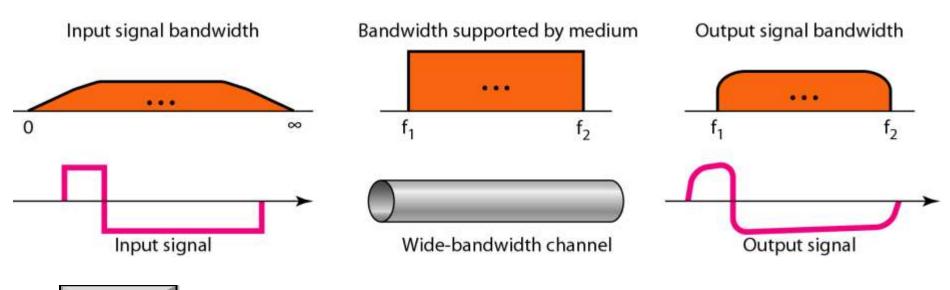
The baseband: a type of signal which has a frequency range near to zero.

Baseband Transmission: sending digital signal without changing digital to an analog signal.

Example: 10Base5, 100BaseT, 1000BaseT

Baseband Transmission

- **Baseband channel (or lowpass channel)**

  a communication channel that can transfer frequencies that are very near zero.

- **Baseband bandwidth**

  equal to the highest frequency of a signal or system, or an upper bound on such frequencies

Baseband transcription using a **dedicated** medium

Input signal bandwidth

Bandwidth supported by medium

Output signal bandwidth

Input signal

Wide-bandwidth channel

Output signal

Note

Baseband transmission of a digital signal that preserves the shape of the digital signal is possible only if we have a low-pass channel with an infinite or very wide bandwidth.

Baseband Transmission

Note

In baseband transmission, the required bandwidth is proportional to the bit rate; if we need to send bits faster, we need more bandwidth.

Baseband Transmission

Table *Bandwidth requirements for different bit rate*

Bit Rate	Harmonics 1	Harmonics 1,3	Harmonics 1,3,5
1 kbps	B=500 Hz	B=1.5 KHz	B=2.5 KHz
10 kbps	B=5k Hz	B=15 kHz	B=25 kHz
100 kbps	B=50 kHz	B=150 kHz	B=250 kHz

Baseband Transmission

What is the required bandwidth of a low-pass channel if we need to send 1 Mbps by using baseband transmission?

Solution

The answer depends on the accuracy desired.

a. The minimum bandwidth, is B = bit rate /2, or 500 kHz.

b. A better solution is to use the first and the third harmonics with B = 3 ✕ 500 kHz = 1.5 MHz.

c. Still a better solution is to use the first, third, and fifth harmonics with B = 5 ✕ 500 kHz = 2.5 MHz.
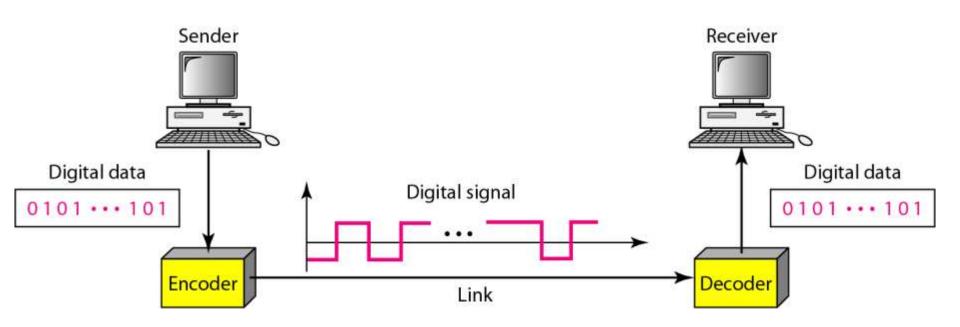
Baseband Transmission

We have a low-pass channel with bandwidth 100 kHz. What is the maximum bit rate of this channel?

Solution

Assume binary data are transmitted.
The bit rate is 2 times the available bandwidth, or 200 kbps.

Digital Data, Digital Signal Encoding

- **Digital Data:**

  - **Sequence of bits to transmit**

- **Physical signal:**

  - **Physical value e.g. voltage, changing in discrete time epochs.**

- **Encoding (channel coding):**

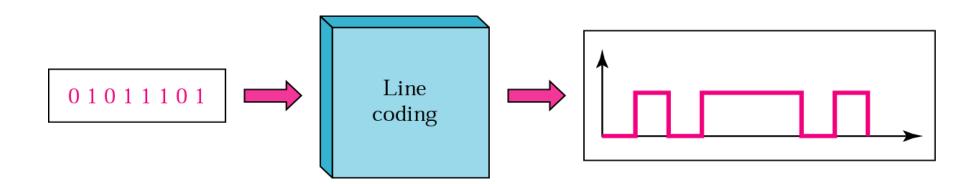 - **Mapping of bit sequences to signals. This can be done in many ways...**

Digital Data, Digital Signal Encoding

Digital Data, Digital Signal Encoding

- **Interpreting Signals**
  - **Need to know**
    - **Timing of bits - when they start and end**
    - **Signal levels**
  - **Factors affecting successful interpreting of signals**
    - **Signal to noise ratio**
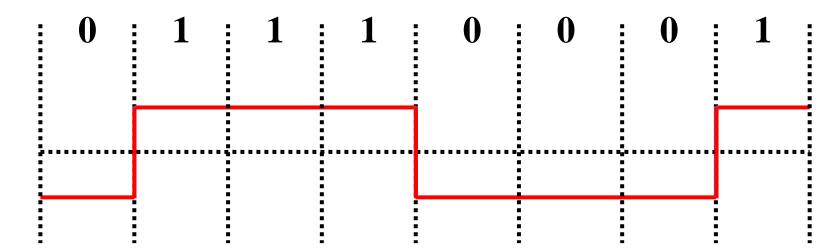   - **Data rate**
    - **Bandwidth**
- **Issues**
  - **Bit timing**
  - **Recovery from signal**
  - **Noise immunity**

Digital Data, Digital Signal Encoding

Line coding is the process of converting binary data, a sequence of bits, to a digital signal.

Digital Data, Digital Signal Encoding
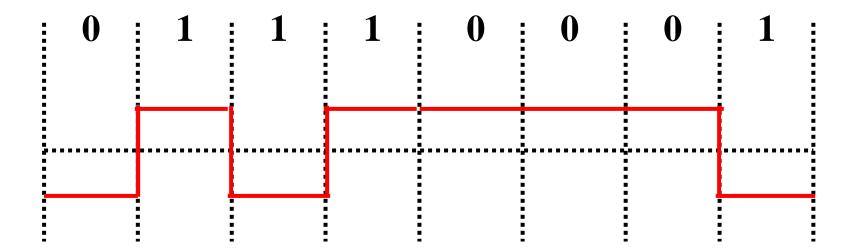
- **Encoding Schemes**
  - **Nonreturn to Zero-Level (NRZ-L)**
  - **Nonreturn to Zero Inverted (NRZI)**
  - **Manchester**
  - **Differential Manchester**
  - **……**

Digital Data, Digital Signal Encoding

- **NRZ-L (Nonreturn-to-Zero-Level)**
  - **0 = low level**
  - **1 = high level**
  - **Or exactly the other way round…☺**

0	1	1	1	0	0	0	1

问题：接收方如何知道有几个"1"和"0"？

Digital Data, Digital Signal Encoding

- ## NRZI (Nonreturn-to-Zero-Inverted) - *polarity not important!!!*
  - □ **0 = no transition at beginning of interval**
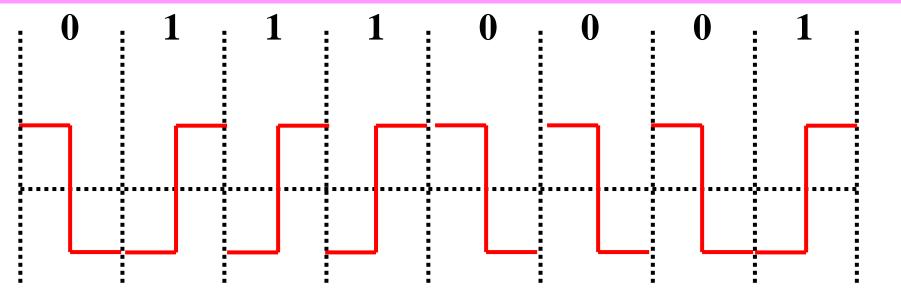 - □ **1 = transition at beginning of interval**

| 0 | 1 | 1 | 1 | 0 | 0 | 0 | 1 |

问题：接收方如何知道有几个 "0"？
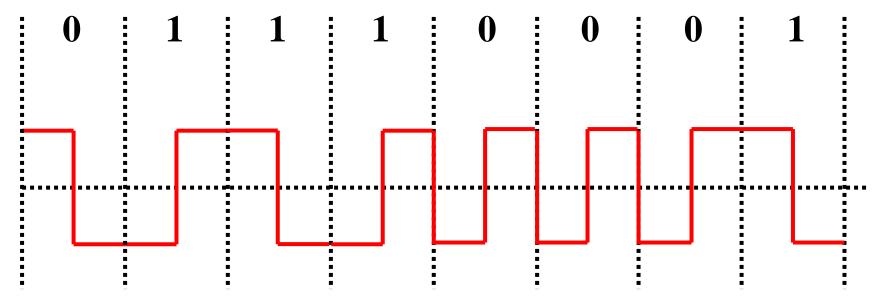
Digital Data, Digital Signal Encoding

■ **Manchester:** *Solves the problem of long constant values…*

- ❑ **0 = transition high ➔ low in middle of interval**
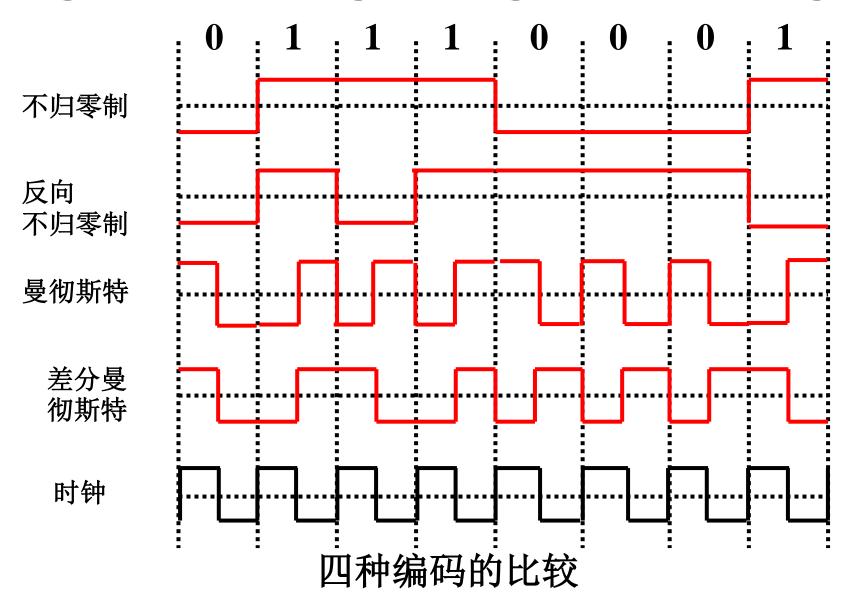- ❑ **1 = transition low ➔ high in middle of interval**

Note: Bit rate lower than baud rate!!

The direction of the mid-bit transition represents the digital data.

$$0 \quad 1 \quad 1 \quad 1 \quad 0 \quad 0 \quad 0 \quad 1$$

Digital Data, Digital Signal Encoding

- **Differential Manchester:**
  - **Always a transition in middle of interval**
    - **mid-bit transition is ONLY for clocking.**
  - **0 = transition at beginning of interval**
  - **1 = no transition at beginning of interval**

$$0 \quad 1 \quad 1 \quad 1 \quad 0 \quad 0 \quad 0 \quad 1$$

Digital Data, Digital Signal Encoding

四种编码的比较

NRZ Pros and Cons

- **Pros**
  - **Easy to engineer**
  - **Make good use of bandwidth**
- **Cons**
  - **dc component**
  - **Lack of synchronization capability**
- **Used for magnetic recording**
- **Not often used for signal transmission**

Manchester Pros and Cons

- **Pros**
  - **Synchronization on mid bit transition (self clocking)**
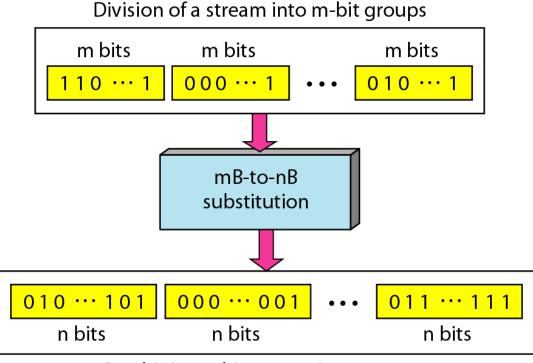 - **No dc component**
  - **Error detection**
- **Con**
  - **At least one transition per bit time and possibly two**
  - **Maximum modulation rate is twice NRZ**
  - **Requires more bandwidth**
  - **Bandwidth inefficient: 50%**

Block coding

- **Block coding is normally referred to as mB/nB coding;**

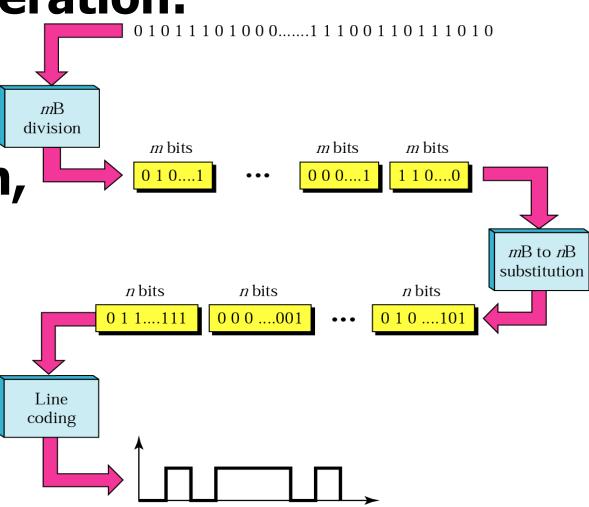- **it replaces each m-bit group with an n-bit group.**

Division of a stream into m-bit groups

m bits	m bits		m bits
1 1 0 ⋯ 1	0 0 0 ⋯ 1	⋯	0 1 0 ⋯ 1

mB-to-nB substitution

0 1 0 ⋯ 1 0 1	0 0 0 ⋯ 0 0 1	⋯	0 1 1 ⋯ 1 1 1
n bits	n bits		n bits

Combining n-bit groups into a stream

Block coding

- **Stages of operation:**
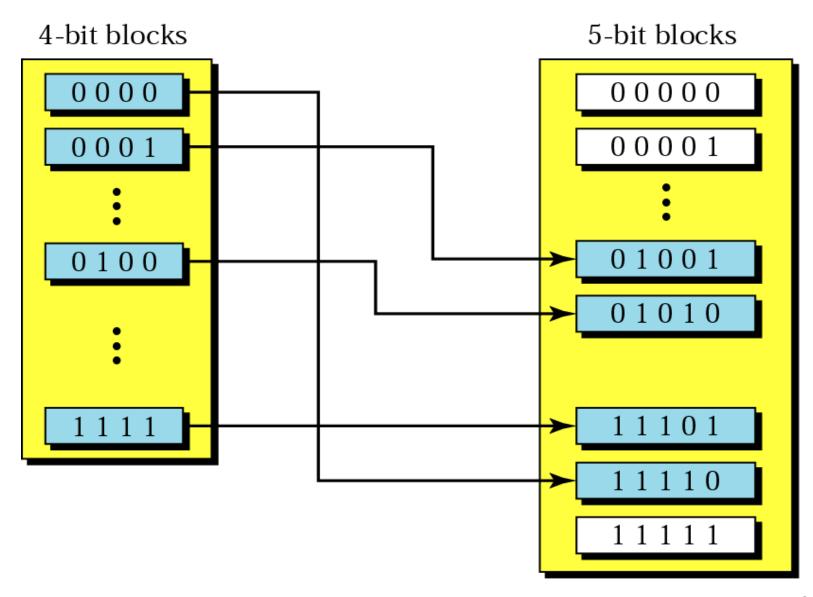
  - **Division,**
  - **Substitution,**
  - **Line Coding**

0 1 0 1 1 1 0 1 0 0 0.......1 1 1 0 0 1 1 0 1 1 1 0 1 0

mB division

m bits — m bits — m bits

0 1 0....1   ...   0 0 0....1   1 1 0....0

mB to nB substitution

n bits n bits n bits

0 1 1....111   0 0 0 001   ...   0 1 0 101

Line coding

4B/5B Block Coding

4-bit blocks

| 0 0 0 0 |
| 0 0 0 1 |
| ⋮ |
| 0 1 0 0 |
| ⋮ |
| 1 1 1 1 |

5-bit blocks

| 0 0 0 0 0 |
| 0 0 0 0 1 |
| ⋮ |
| 0 1 0 0 1 |
| 0 1 0 1 0 |
| 1 1 1 0 1 |
| 1 1 1 1 0 |
| 1 1 1 1 1 |

4B/5B Mapping Codes (1/2)
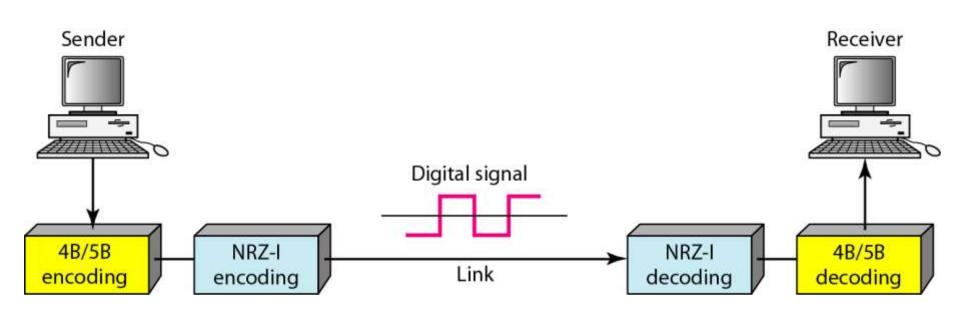
Data	Code	Data	Code
0000	11110	1000	10010
0001	01001	1001	10011
0010	10100	1010	10110
0011	10101	1011	10111
0100	01010	1100	11010
0101	01011	1101	11011
0110	01110	1110	11100
0111	01111	1111	11101

- **The selection of the 5-bit code is such that each code contains no more than one leading 0 and no more than two trailing 0s.**
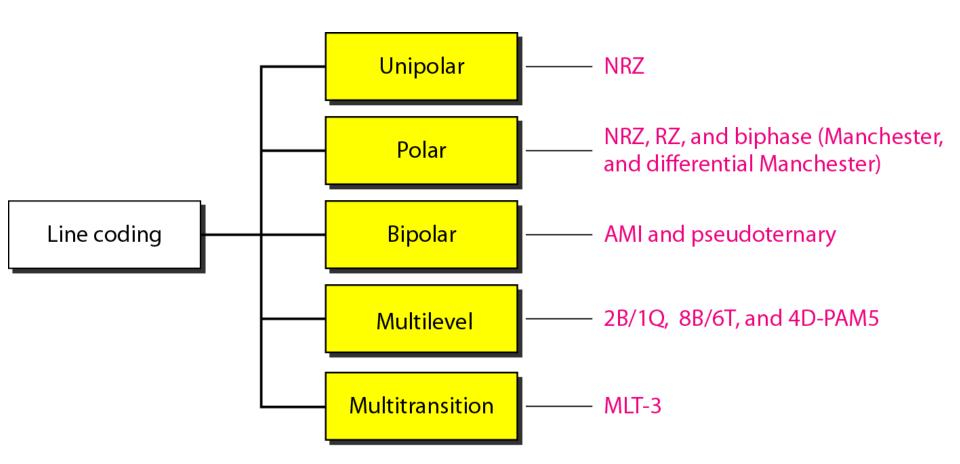
4B/5B Mapping Codes (2/2)

Data	Code
Q (Quiet)	00000
I (Idle)	11111
H (Halt)	00100
J (start delimiter)	11000
K (start delimiter)	10001
T (end delimiter)	01101
S (Set)	11001
R (Reset)	00111

4B/5B Block Coding
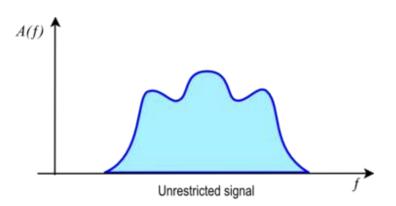

Using block coding 4B/5B with NRZ-I line coding scheme
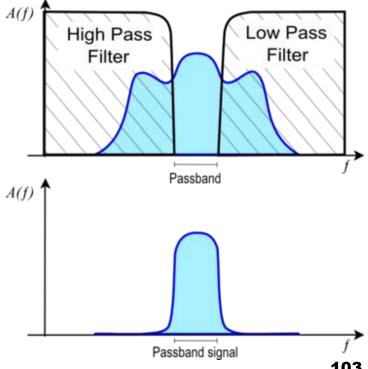
Line coding schemes

Pass-band Transmission

- **Passband transmission is the transmission after shifting the baseband frequencies to some higher frequency range using modulation.**
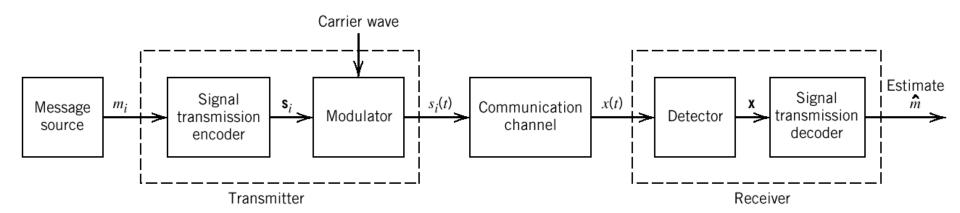
Pass-band Transmission

- **a passband (通带) is the portion of the frequency spectrum that is transmitted by some filtering device.**
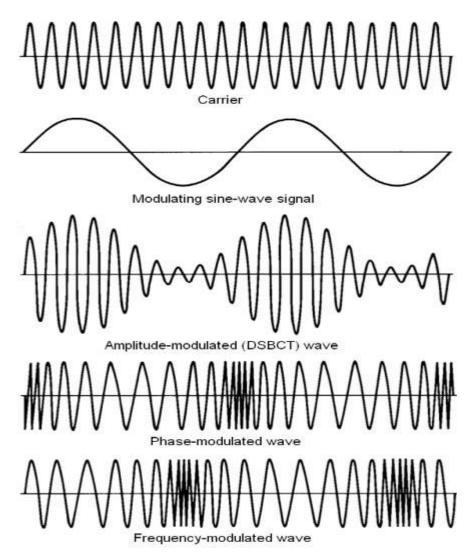
Only specific frequency bands are used.

Unrestricted signal

High Pass Filter   Low Pass Filter

Passband

Passband signal

Pass-band Transmission

- **Passband different from the "natural spectrum of the signal".**

- **The "useful" signal has to be "shifted" where desired...**

- **Called "pass-band transmission"**

- **We do this, using modulation ...**
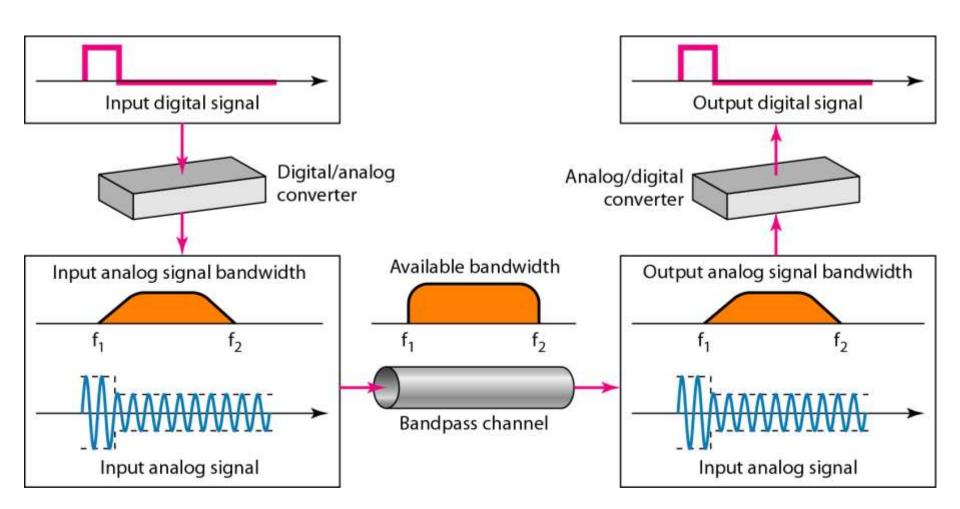
Pass-band Transmission

Functional model of pass-band data transmission system.

Pass-band Transmission
(Analog Data and Analog Signals!)

- **We shift the frequency spectrum elsewhere...**

- **The different kinds of modulation can be combined**

Carrier

Modulating sine-wave signal

Amplitude-modulated (DSBCT) wave

Phase-modulated wave

Frequency-modulated wave

Pass-band Transmission
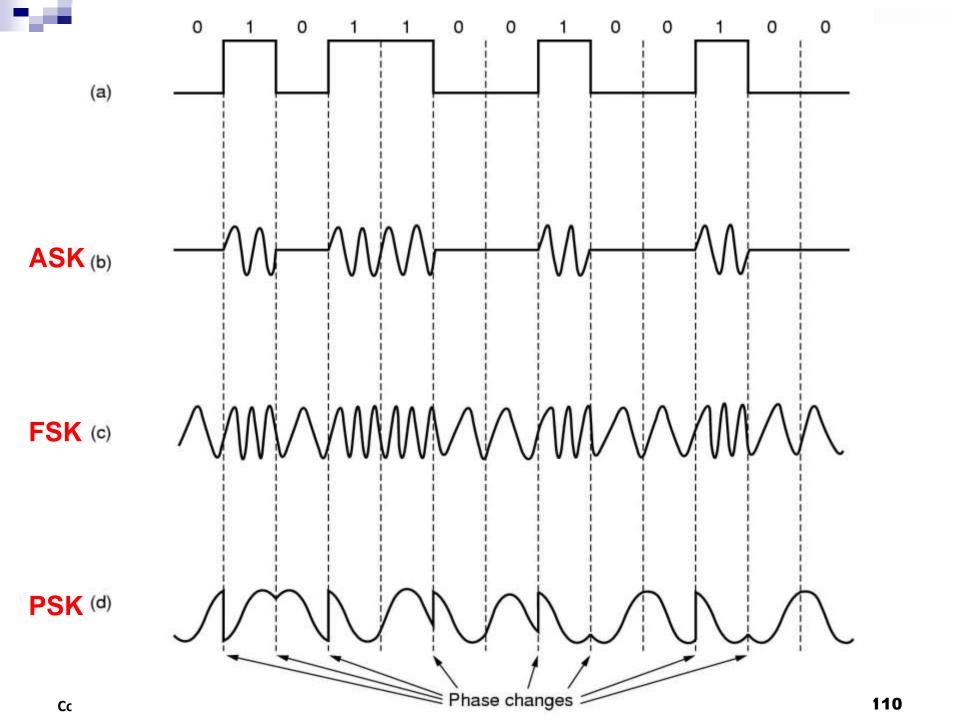## (Digital Data and Analog Signals!)

Modulation of a digital signal for transmission on a bandpass channel

Pass-band Transmission (Digital Data and Analog Signals!)

- **Digital data is encoded by modulating one of the three characteristics of the carrier:**
  - **amplitude,**
  - **frequency, or**
  - **Phase, or**
  - **some combination of these.**

Fundamental digital modulation methods

- **ASK (amplitude-shift keying):** a finite number of amplitudes are used.

- **FSK (frequency-shift keying):** a finite number of frequencies are used.

- **PSK (phase-shift keying):** a finite number of phases are used.

- **QAM (quadrature amplitude modulation):** a finite number of at least two phases and at least two amplitudes are used.

ASK (b)

FSK (c)

PSK (d)

Phase changes

Increasing Transmission Rates

- **Observation: An important issue is to use low-baud modems for high transmission rates, by increasing the number of signal values—according to Nyquist, Bit rate =2 * bandwidth * $\log_2 V$.**

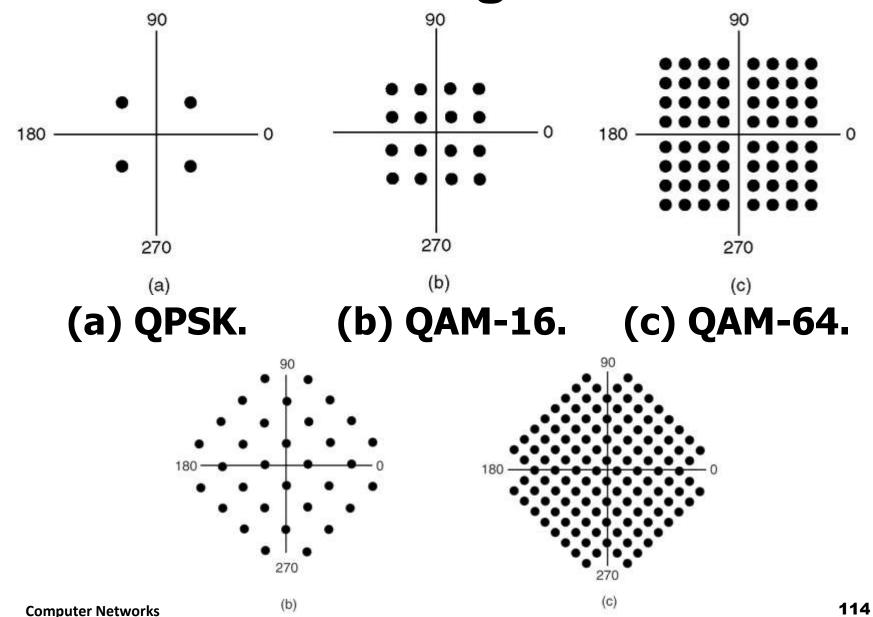- **The key issue: how to present more bit information for each sample.**

Modems

- **An acronym for modulator-demodulator**
- **Uses a constant-frequency signal known as a carrier signal**
- **Converts a series of binary voltage pulses into an analog signal by modulating the carrier signal**
- **The receiving modem translates the analog signal back into digital data**

Modems

- **All advanced modems use a *combination of modulation techniques* to transmit multiple bits per baud.**

- **Multiple amplitude and multiple phase shifts are combined to transmit several bits per symbol.**

- **Modems actually use Quadrature Amplitude Modulation (QAM).**

- **These concepts are explained using constellation points where a point determines a specific amplitude and phase.**
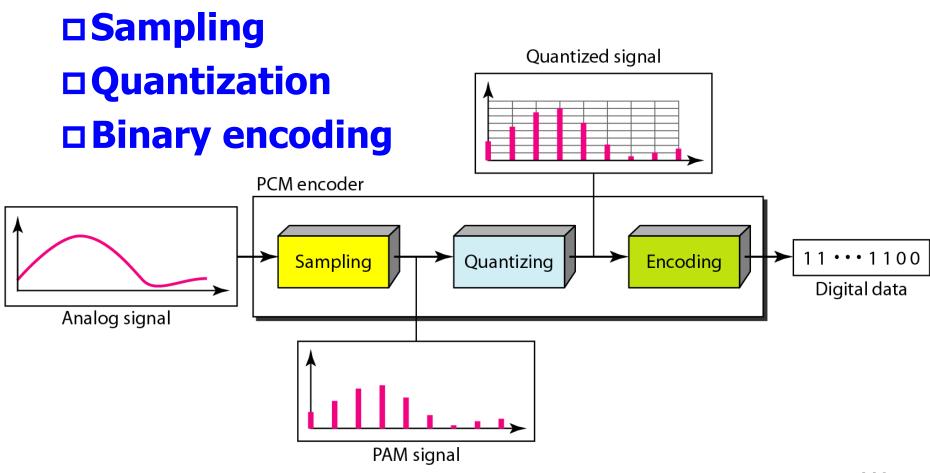
Constellation Diagrams

(a) QPSK.   (b) QAM-16.   (c) QAM-64.

Analog-to-Digital Conversion

■ **A digital signal is superior to an analog signal because it is more robust to noise and can easily be recovered, corrected and amplified.**

■ **For this reason, the tendency today is to change an analog signal to digital data. In this section we describe two techniques:**

  ❑ **Pulse Code Modulation (PCM) and**
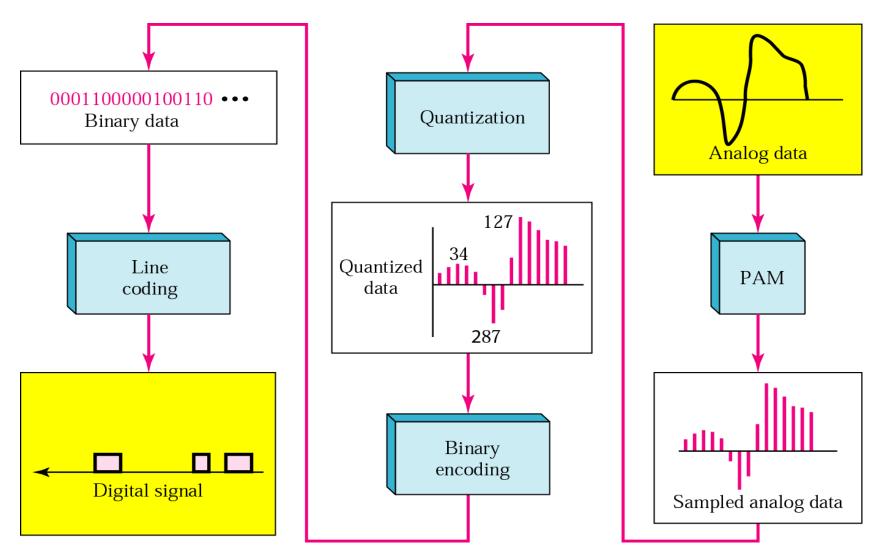
  ❑ **Delta Modulation (DM)**

PCM

- **PCM consists of three steps to digitize an analog signal:**
  - **Sampling**
  - **Quantization**
  - **Binary encoding**

Quantized signal

PCM encoder

Sampling → Quantizing → Encoding → 11 · · · 1100

Analog signal

Digital data

PAM signal
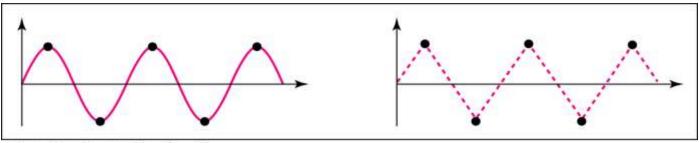
PCM

From analog signal to PCM digital code

Sampling

- **Analog signal is sampled every T_S secs.**

- **T_s is referred to as the sampling interval.**

- **$f_s = 1/T_s$ is called the sampling rate or sampling frequency.**
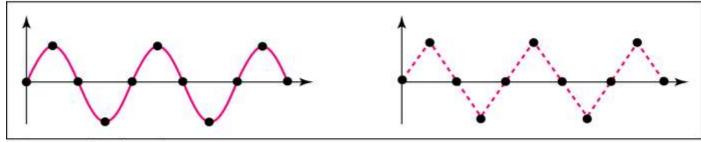
Note

**According to the Nyquist theorem, the sampling rate must be
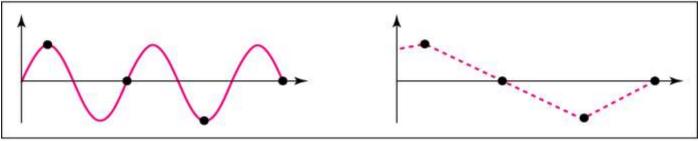at least 2 times the highest frequency contained in the signal.**

Example

a. Nyquist rate sampling: $f_s = 2f$

b. Oversampling: $f_s = 4f$

c. Undersampling: $f_s = f$

Recovery of a sampled sine wave for different sampling rates

Quantization

- **Sampling results in a series of pulses of varying amplitude values ranging between two limits: a min and a max.**

- **The amplitude values are infinite between the two limits.**

- **We need to map the *infinite* amplitude values onto a finite set of known values.**

- **This is achieved by dividing the distance between min and max into L zones, each of height Δ.**

$$\Delta = (max - min)/L$$

Quantization Levels

- **The midpoint of each zone is assigned a value from 0 to L-1 (resulting in L values)**

- **Each sample falling in a zone is then approximated to the value of the midpoint.**

Quantization Zones

- **Assume we have a voltage signal with amplitutes V_{min}=-20V and V_{max}=+20V.**

- **We want to use L=8 quantization levels.**

- **Zone width $\triangle$ = (20 - -20)/8 = 5**

- **The 8 zones are: -20 to -15, -15 to -10, -10 to -5, -5 to 0, 0 to +5, +5 to +10, +10 to +15, +15 to +20**

- **The midpoints are: -17.5, -12.5, -7.5, -2.5, 2.5, 7.5, 12.5, 17.5**

Assigning Codes to Zones

- **Each zone is then assigned a binary code.**
- **The number of bits required to encode the zones, or the number of bits per sample as it is commonly referred to, is obtained as follows:**

$$n_b = \log_2 L$$

- **Given our example, $n_b = 3$**
- **The 8 zone (or level) codes are therefore: 000, 001, 010, 011, 100, 101, 110, and 111**
- **Assigning codes to zones:**
  - **000 will refer to zone -20 to -15**
  - **001 to zone -15 to -10, etc.**

Bit rate and bandwidth requirements of PCM

- **The bit rate of a PCM signal can be calculated form the number of bits per sample x the sampling rate**

$$\textbf{Bit rate} = n_b \textbf{ x } f_s$$

- **The bandwidth required to transmit this signal depends on the type of line encoding used.**

Example

We want to digitize the human voice. What is the bit rate, assuming 8 bits per sample?

Solution

The human voice normally contains frequencies from 0 to 4000 Hz. So the sampling rate and bit rate are calculated as follows:

$$\text{Sampling rate} = 4000 \times 2 = 8000 \text{ samples/s}$$
$$\text{Bit rate} = 8000 \times 8 = 64{,}000 \text{ bps} = 64 \text{ kbps}$$
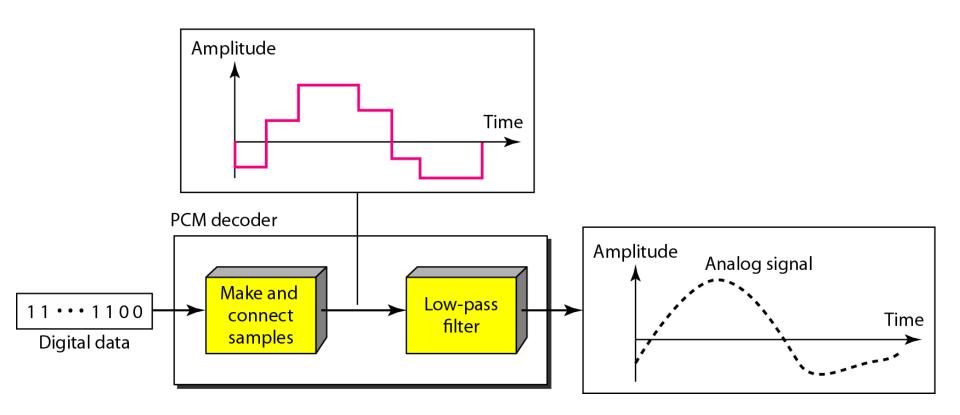
Example

- **A digitized signal will always need more bandwidth than the original analog signal.**

- **We have a low-pass analog signal of 4 kHz.**
  - **If we send the analog signal, we need a channel with a minimum bandwidth of 4 kHz.**
  - **If we digitize the signal and send 8 bits per sample, we need a channel with a minimum bandwidth of**

    8×4 kHz $= 32$ kHz.

PCM Decoder

- **To recover an analog signal from a digitized signal we follow the following steps:**

  - **We use a hold circuit that holds the amplitude value of a pulse till the next pulse arrives.**

  - **We pass this signal through a low pass filter with a cutoff frequency that is equal to the highest frequency in the pre-sampled signal.**

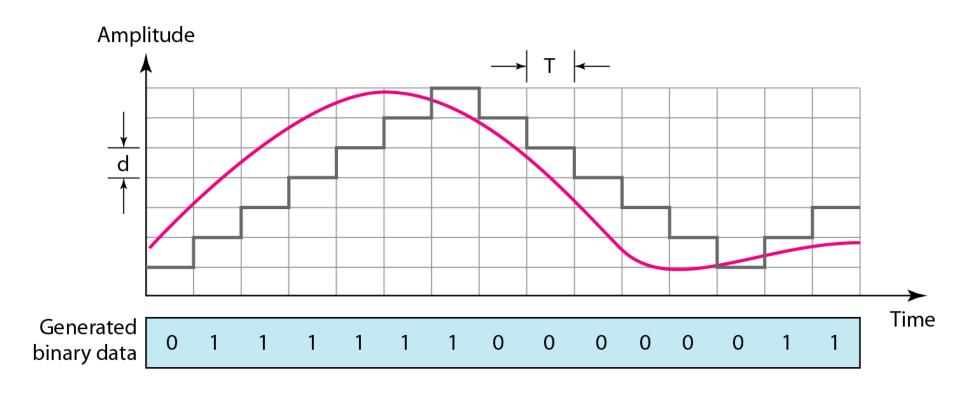- **The higher the value of L, the less distorted a signal is recovered.**

PCM Decoder

Delta Modulation

- **This scheme sends only the difference between pulses, if the pulse at time t_{n+1} is higher in amplitude value than the pulse at time t_n, then a single bit, say a "1", is used to indicate the positive value.**

- **If the pulse is lower in value, resulting in a negative value, a "0" is used.**
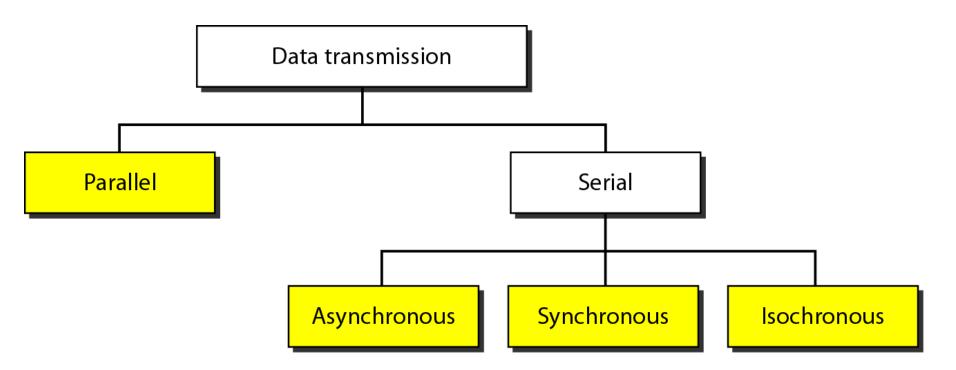
Delta Modulation

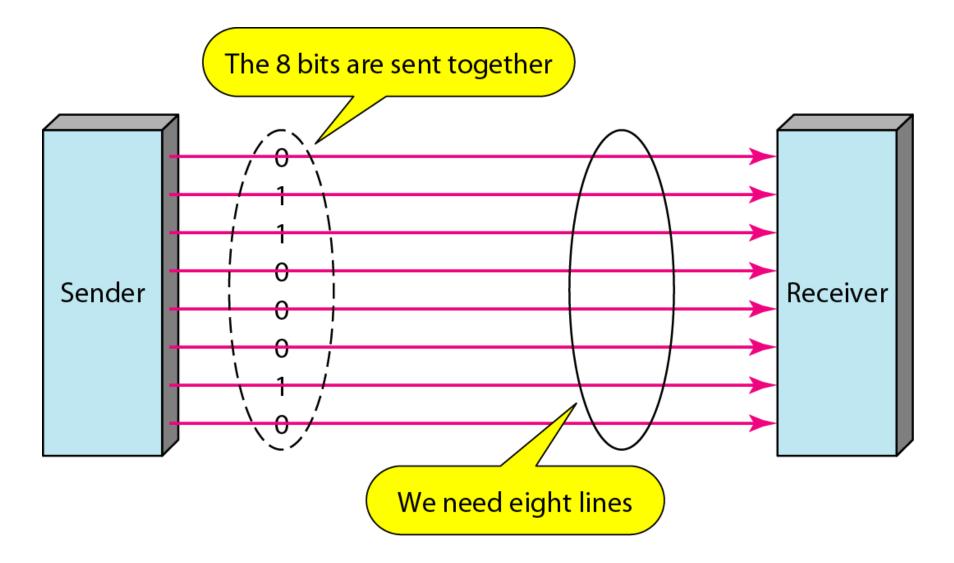*The process of delta modulation*

Delta PCM (DPCM)

- **Instead of using one bit to indicate positive and negative differences, we can use more bits -> quantization of the difference.**

- **Each bit code is used to represent the value of the difference.**

- **The more bits the more levels -> the higher the accuracy.**
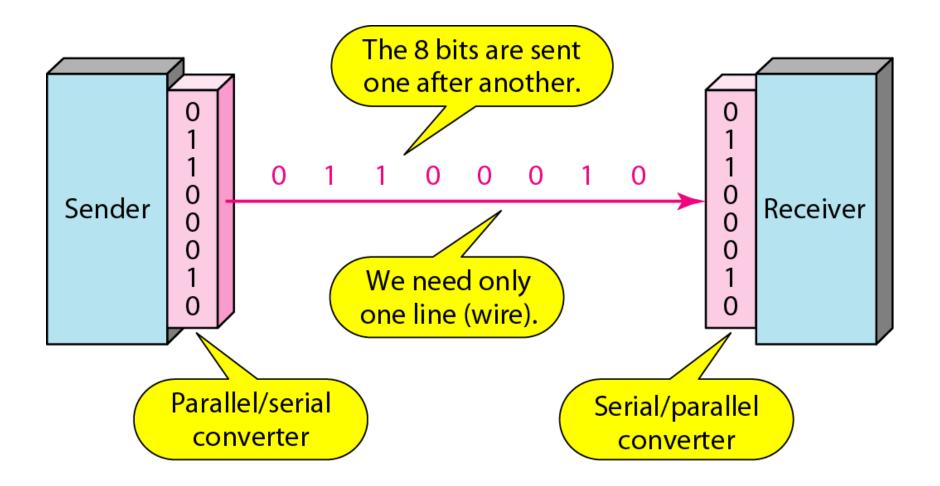
Transmission Modes

- **The transmission of binary data across a link can be accomplished in either parallel or serial mode.**

Parallel transmission

The 8 bits are sent together
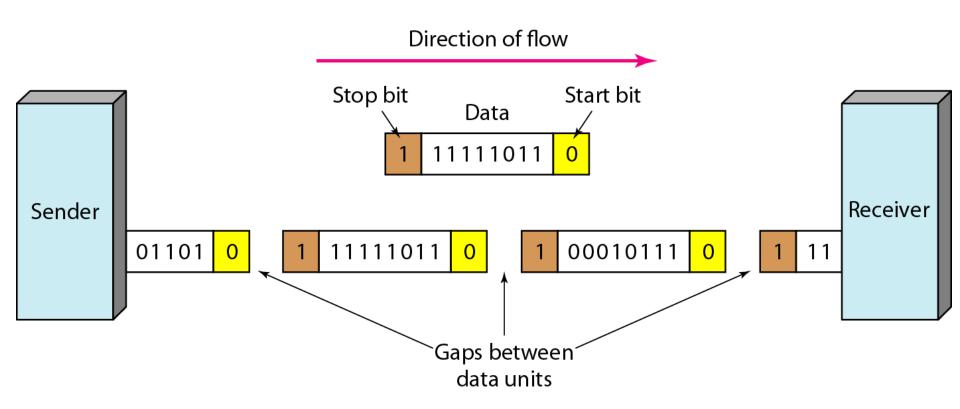
We need eight lines

Serial transmission

Asynchronous transmission

- **In asynchronous transmission, we send 1 start bit (0) at the beginning and 1 or more stop bits (1s) at the end of each byte. There may be a gap between each byte.**

Note

**Asynchronous here means
"asynchronous at the byte level,"
but the bits are still synchronized;
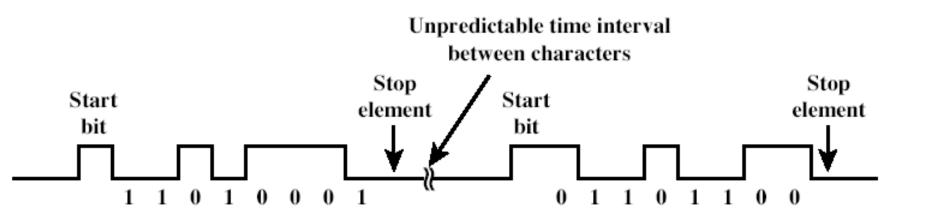their durations are the same.**

Asynchronous transmission
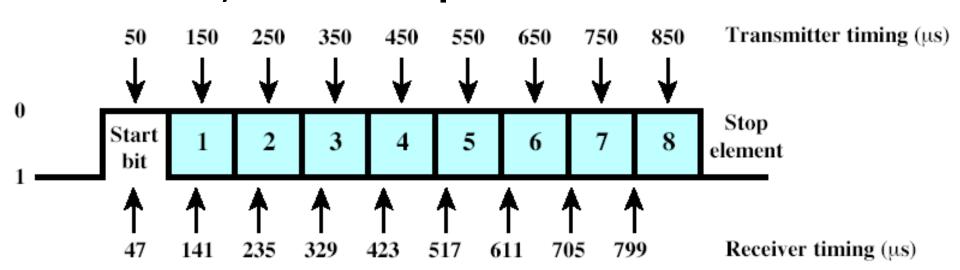
Asynchronous transmission

Idle state of line

odd or even parity or unused

Remain idle or next start bit

5 to 8 data bits

1 to 2 bit times

0

1

Start bit

P bit

Stop element

(a) Character format

Unpredictable time interval between characters

Start bit

Stop element

Start bit

Stop element

1 1 0 1 0 0 0 1

0 1 1 0 1 1 0 0

(b) 8-bit asynchronous character stream

❑ **Example: The figure below shows the effects of a timing error of sufficient magnitude to cause error in reception. In this example, we assume a data rate of 10Kbps; therefore each bit is 100µs duration. Assume that the receiver is fast by 6%, or 6µs per bit time. Thus, the receiver samples the incoming character every 94µs. As we can see, the last sample is erroneous.**
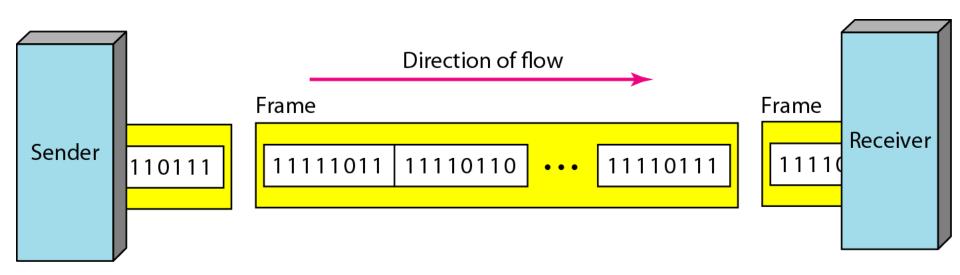

(c) Effect of timing error

Asynchronous Transmission- Behavior

- ❑ **simple**
- ❑ **cheap**
- ❑ **overhead of 2 or 3 bits per char (~20%)**
- ❑ **good for data with large gaps (keyboard)**

Synchronous transmission

- **In synchronous transmission, we send bits one after another without start or stop bits or gaps.**

- **It is the responsibility of the receiver to group the bits. The bits are usually sent as bytes and many bytes are grouped in a frame.**

- **A frame is identified with a start and an end byte.**

Synchronous transmission

Direction of flow

Sender 110111

Frame
| 11111011 | 11110110 | ••• | 11110111 |

Frame
11110... Receiver

Isochronous

- **In isochronous transmission we cannot have uneven gaps between frames.**

- **Transmission of bits is fixed with equal gaps.**
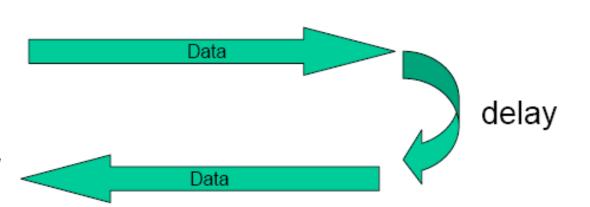
Simplex *vs.* Half Duplex *vs.* Full Duplex

Simplex data flow

Data →

One direction, e.g. Television

Half-Duplex data flow

Data →

delay

Either direction, but only one way at a time, e.g. police radio

← Data

Full-Duplex data flow

Data →

no delay

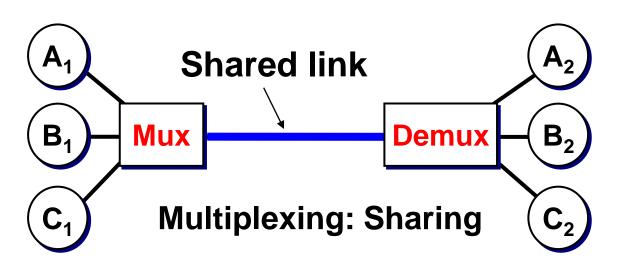Both directions at the same time e.g. telephone
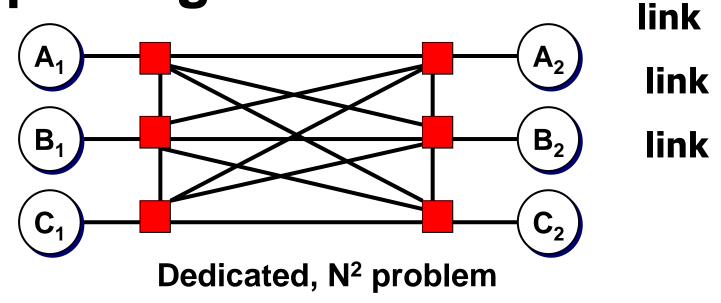
← Data

Chapter 2: Roadmap

Bandwidth utilization is the wise use of available bandwidth to achieve specific goals.

Efficiency can be achieved by multiplexing; privacy and anti-jamming can be achieved by spreading.
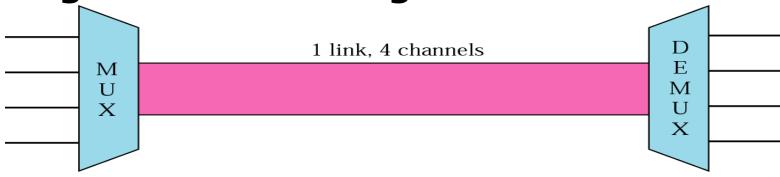
Multiplexing

link

link

link

A_1 A_2

B_1 B_2

C_1 C_2

Dedicated, N^2 problem

A_1 **Shared link** A_2

B_1 **Mux** **Demux** B_2

C_1 **Multiplexing: Sharing** C_2

Multiplexing

- ## Multiplexing

  - **A set of techniques that allows the simultaneous transmission of multiple signals across a single data link.**
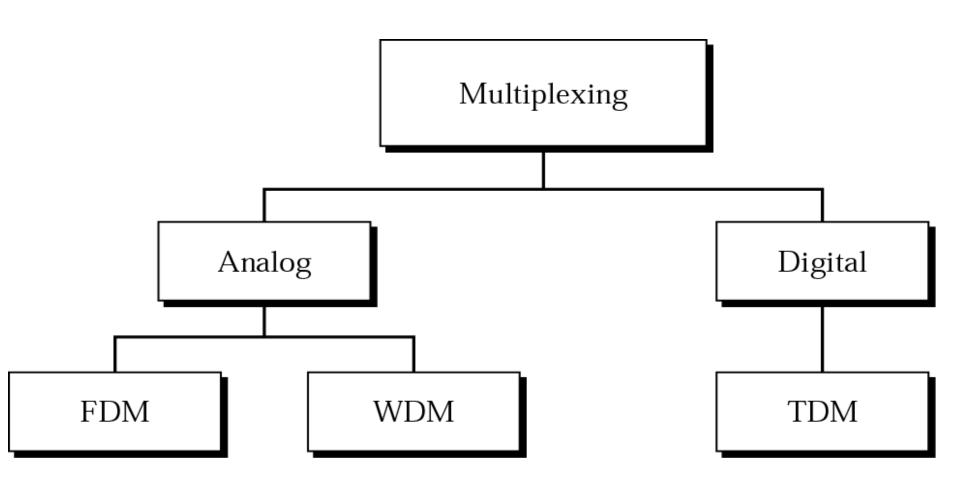

Multiplexer (MUX) :

    Combines multiple streams into a single stream (many to one).
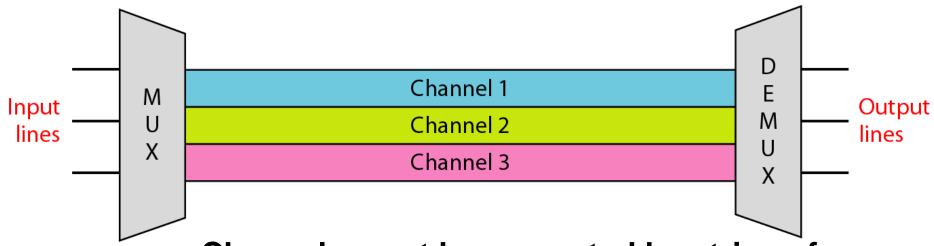
Demultiplexer (DEMUX) :

    Separates the stream back into its component transmission (one to many) and directs them to their correct lines.

Category of Multiplexing

FDM : Frequency-division multiplexing

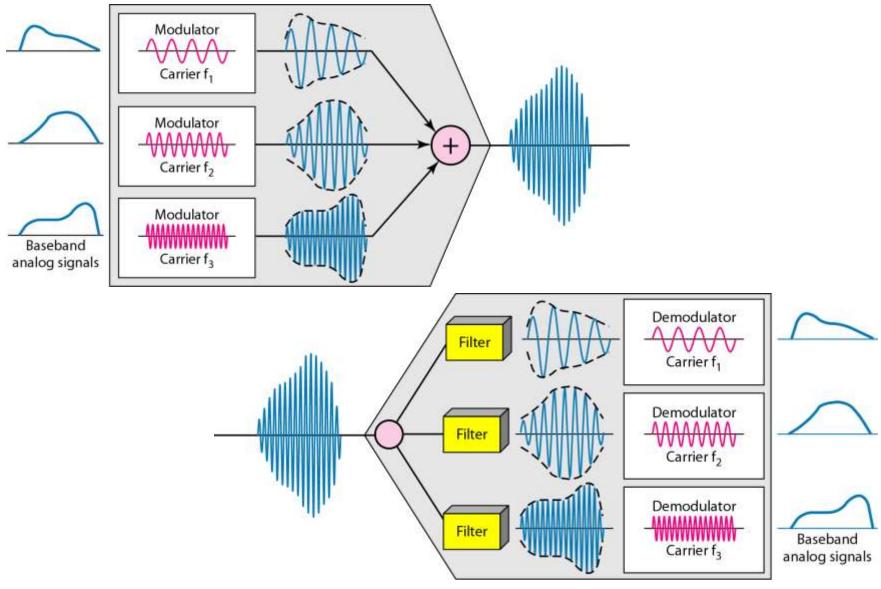## Assigns different analog frequencies to each connected device

Input lines — MUX — Channel 1 / Channel 2 / Channel 3 — DEMUX — Output lines

Channels must be separated by strips of unused bandwidth - *guard bandwidth*

Note

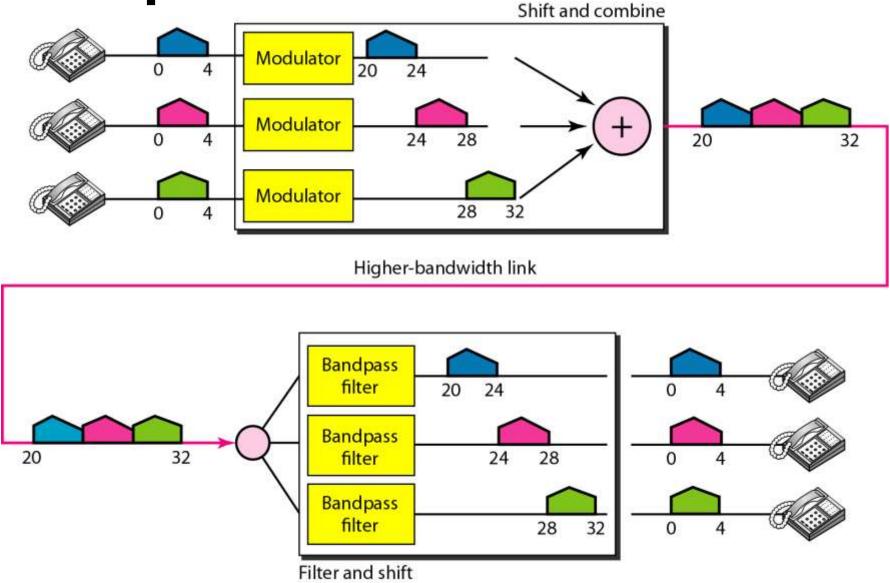**FDM is an analog multiplexing technique that combines analog signals.**

FDM process

Example 1

Assume that a voice channel occupies a bandwidth of 4 kHz. We need to combine three voice channels into a link with a bandwidth of 12 kHz, from 20 to 32 kHz. Show the configuration, using the frequency domain. Assume there are no guard bands.
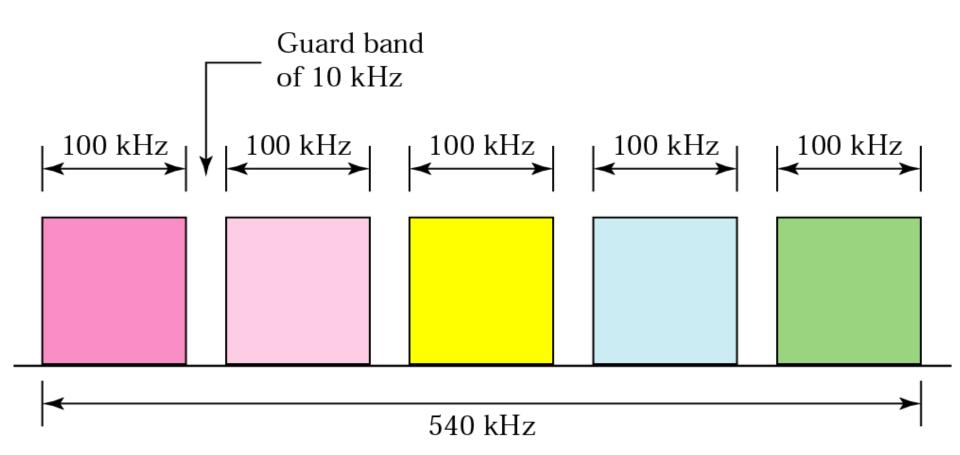
Example 1

Solution

We shift (modulate) each of the three voice channels to a different bandwidth. We use the 20- to 24-kHz bandwidth for the first channel, the 24- to 28-kHz bandwidth for the second channel, and the 28- to 32-kHz bandwidth for the third one. Then we combine them as shown in Figure.
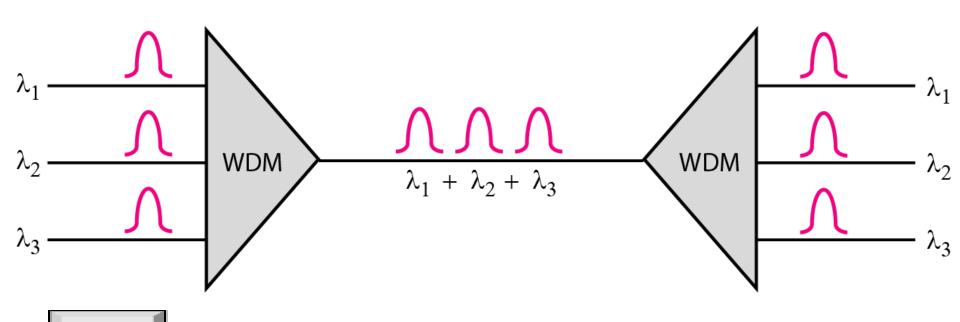
Example 1

Example 2

Five channels, each with a 100-kHz bandwidth, are to be multiplexed together. What is the minimum bandwidth of the link if there is a need for a guard band of 10 kHz between the channels to prevent interference?

Solution

For five channels, we need at least four guard bands. This means that the required bandwidth is at least $5 \times 100 + 4 \times 10 = 540\ kHz$

Example 2

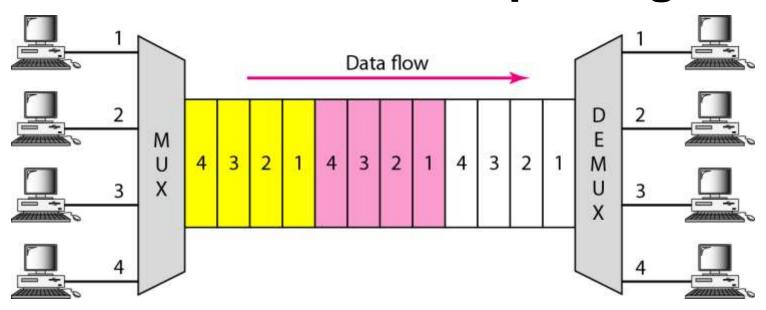# *Wavelength-division multiplexing*

Note

WDM is an analog multiplexing technique to combine optical signals.

TDM : Time Division Multiplexing

Sharing signal is accomplished by dividing available transmission time on a medium among users.

The whole bandwidth is used all the time, but alternatively by different channels!

TDM : Time Division Multiplexing

Note

TDM is a digital multiplexing technique for combining several low-rate channels into one high-rate one.

- **Time division multiplexing comes in two basic forms:**
  - **Synchronous time division multiplexing**
  - **Statistical, or asynchronous time division multiplexing**

TDM : Time Division Multiplexing
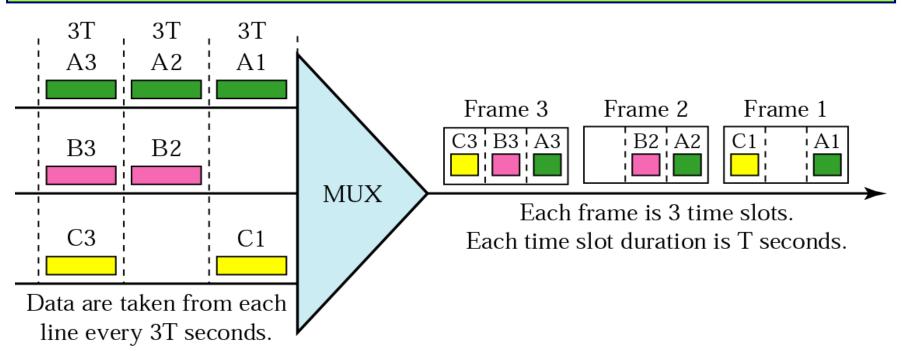
- ## Time Slots

  - Each terminal/host given a "slice" of time (time slot)

- ## Frames

  - In TDM, a frame consists of one complete cycle of time slots, with one slot dedicated to each sending device.

Synchronous time division multiplexing

the multiplexer allocates exactly the same time slot to each device at all times, whether or not a device has anything to transmit.

3T   3T   3T
A3   A2   A1

B3   B2

C3   C1

Data are taken from each line every 3T seconds.

MUX

Frame 3   Frame 2   Frame 1
C3 B3 A3   B2 A2   C1   A1

Each frame is 3 time slots.
Each time slot duration is T seconds.

Synchronous time division multiplexing

Note

In synchronous TDM, Data rate of medium exceeds data rate of digital signal to be transmitted.

Note

In synchronous TDM, the data rate of the link is n times faster, and the unit duration is n times shorter.

Example 4

Four 1-Kbps connections are multiplexed together. A unit is 1 bit. Find (1) the duration of 1 bit before multiplexing, (2) the transmission rate of the link, (3) the duration of a time slot, and (4) the duration of a frame?

Solution

1. The duration of 1 bit is 1/1 Kbps, or 0.001 s (1 ms).
2. The rate of the link is 4 Kbps.
3. The duration of each time slot 1/4 ms or 250 ms.
4. The duration of a frame 1 ms.

Solution in detail

$$DataRate_{link} = 4 \times 1kbps = 4kbps = 4000bps$$

$$BitDuration_{link} = \frac{1bit}{4000bps} = \frac{1bit}{4000bit/\sec ond}$$

$$= 0.25ms/bit = 250\mu s/bit$$
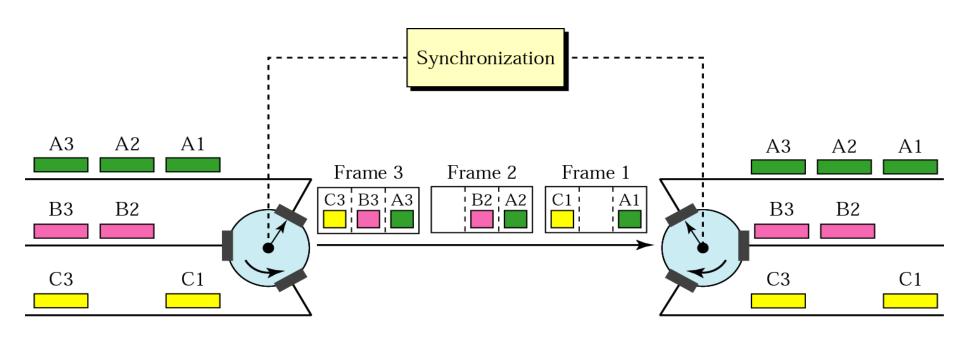
$$TimeSlotDuration_{link} = BitDuration \times UnitSize$$

$$= 250\mu s/bit \times 1bit/TimeSlot = 250\mu s/TimeSlot$$

$$FrameDuration$$

$$= TimeSlotDuration \times ChannelNumber$$

$$= 250\mu s/TimeSlot * 4TimeSlot/Frame$$

$$= 1ms/Frame$$
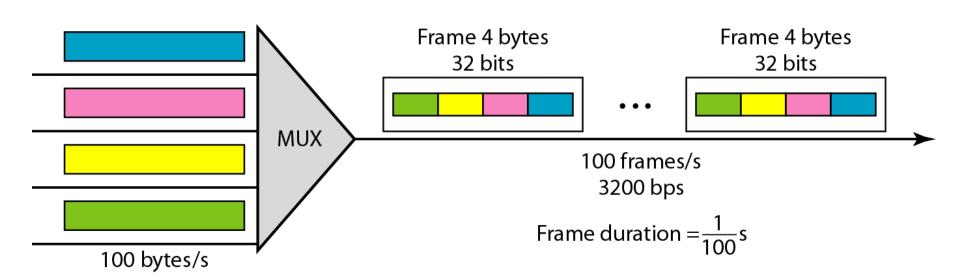
Interleaving of data segments

Example 5

Four channels are multiplexed using TDM. If each channel sends 100 bytes /s and we multiplex 1 byte per channel, show the frame traveling on the link, the size of the frame, the duration of a frame, the frame rate, and the bit rate for the link.
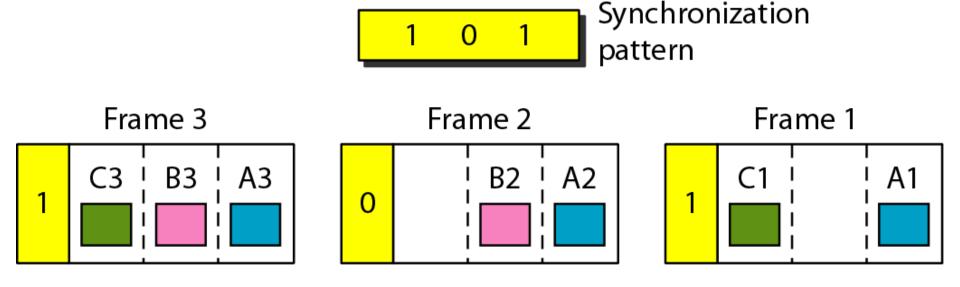
Example 5

Solution

Each frame carries 1 byte from each channel; the size of each frame, therefore, is 4 bytes, or 32 bits. Because each channel is sending 100 bytes/s and a frame carries 1 byte from each channel, the frame rate must be 100 frames per second. The bit rate is 100 $\times$ 32, or 3200 bps.

SYNCHRONIZING

- **One or more *Framing bit (s)* is (are) added to each frame for synchronization between the multiplexer and demultiplxer**

- **If 1 framing bit per frame, framing bits are alternating between 0 and 1**

Synchronization pattern

Frame 3 — 1 | C3 | B3 | A3

Frame 2 — 0 | | B2 | A2

Frame 1 — 1 | C1 | | A1

Example 6

We have four sources, each creating 250 characters per second. If the interleaved unit is a character and 1 synchronizing bit is added to each frame, find (1) the data rate of each source, (2) the duration of each character in each source, (3) the frame rate, (4) the duration of each frame, (5) the number of bits in each frame, and (6) the data rate of the link.

Example 6

1. The data rate of each source is $250 \times 8 = 2000$ bps $= 2$ Kbps.

2. The duration of a character is $1/250$ s, or 4 ms.

3. The link needs to send 250 frames per second to keep the transmission rate of each source.

4. The duration of each frame is $1/250$ s, or 4 ms.

5. Each frame is $4 \times 8 + 1 = 33$ bits.

6. The data rate of the link is 250×33, or 8250 bps.

Example 6

Solution in detail

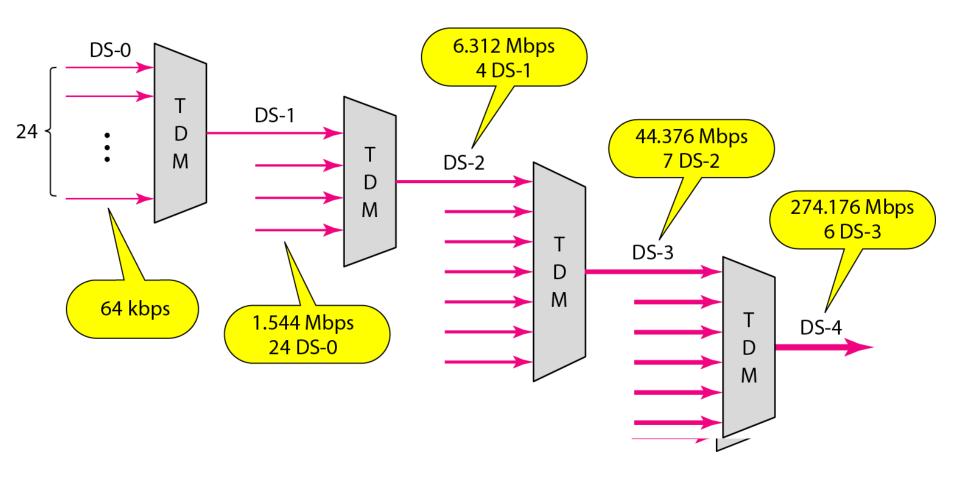$$FrameSize = ChannelNumber \times UnitSize + Fra\min gBits$$

$$= 4timeslot / frame \times 1character / timeslot + 1bit / frame$$

$$= 33bits / frame$$

$$FrameRate = 250 frame / \sec ond$$

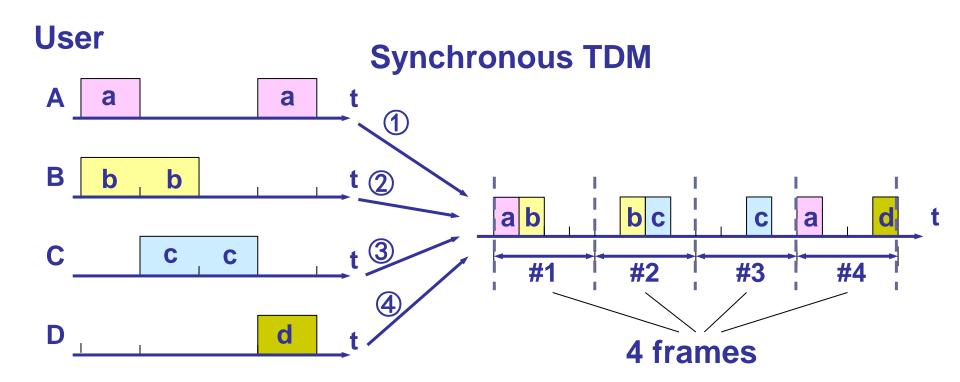$$DataRate = FrameRate \times FrameSize$$

$$= 250 frame / \sec ond \times 33bits / frame$$

$$= 8250bit / \sec ond$$

TDM hierarchy
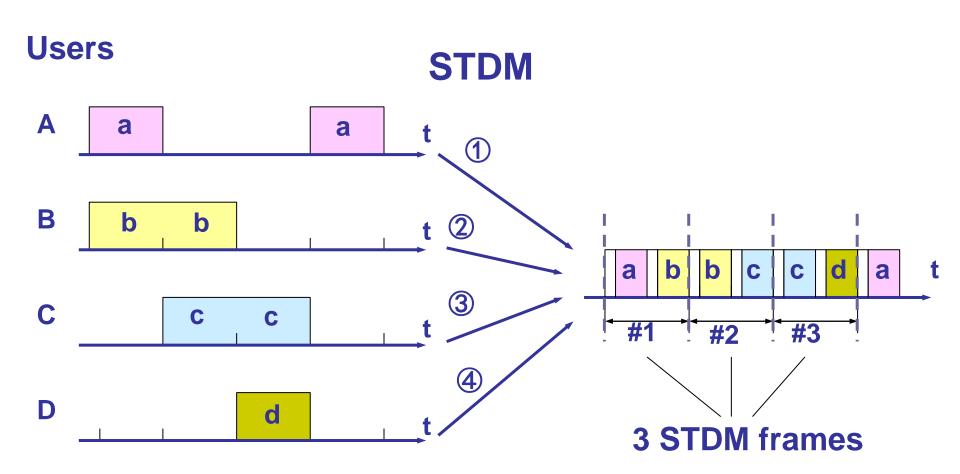
Asynchronous TDM

- **Synchronous TDM does not guarantee that the full capacity of a link is used.**

User

Synchronous TDM

A a   a t ①

B b b t ②

C c c t ③

D d t ④

a b b c c a d t

#1 #2 #3 #4

4 frames

Asynchronous TDM

- **Asynchronous TDM, or statistical TDM (STDM), is designed to avoid this type of waste.**

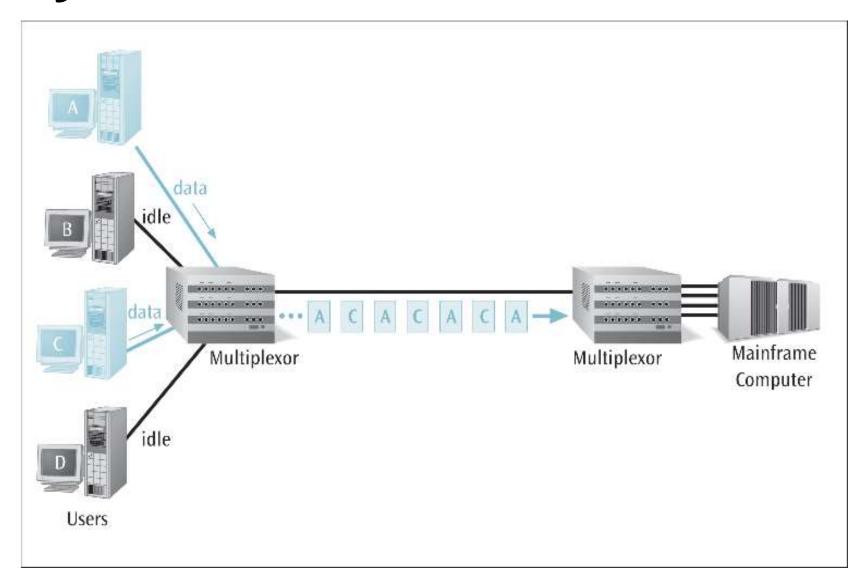- **Statistical TDM allocates time slots dynamically based on demand.**

STDM

Users

STDM

A

| a | | a |

t ①

B

| b | b |

t ②

C

| c | c |

t ③

D

| d |

t ④

| a | b | b | c | c | d | a | t

#1 #2 #3

3 STDM frames
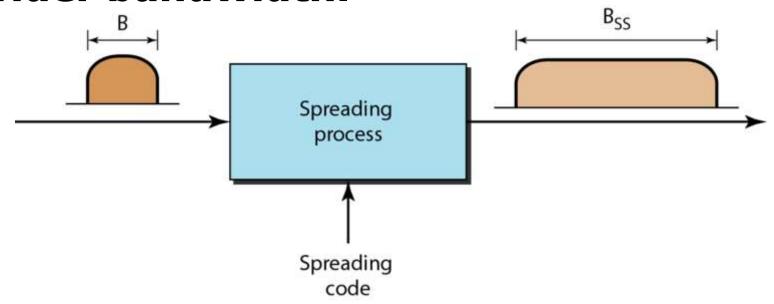
Asynchronous TDM

- **Statistical multiplexor transmits only the data from active workstations**

  - **If a workstation is not active, no space is wasted on the multiplexed stream**

- **A statistical multiplexor**

  - **Accepts incoming data streams**

  - **Creates a frame containing only the data to be transmitted**
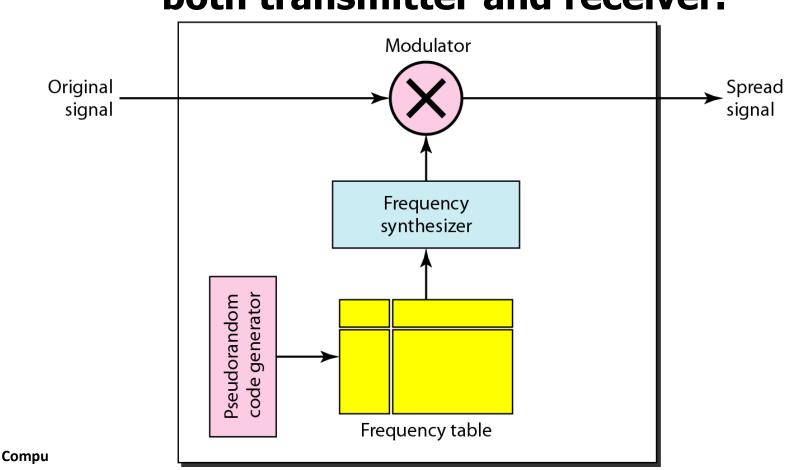
Asynchronous TDM

Spread Spectrum

- **Methods by which a signal generated with a particular bandwidth is deliberately spread in the frequency domain, resulting in a signal with a wider bandwidth.**

Spread Spectrum

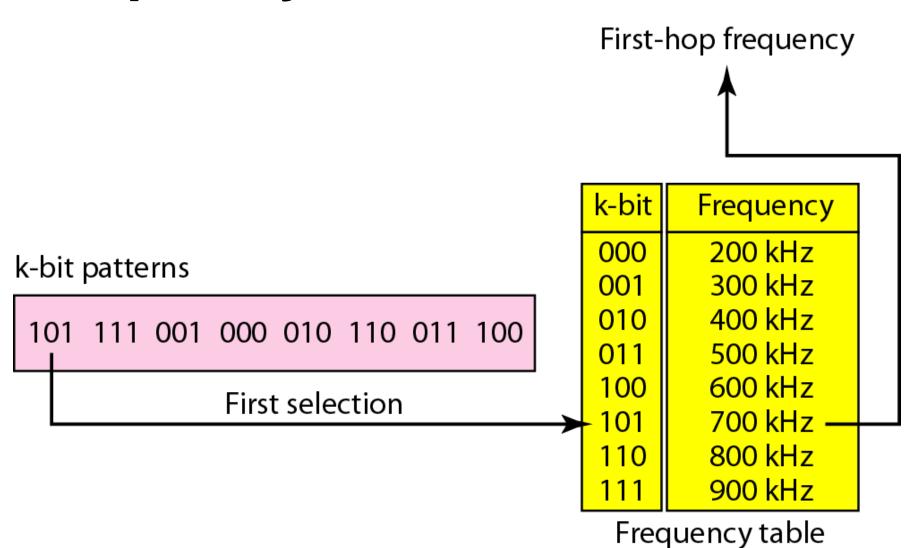- **Spread Spectrum Techniques**
  - **Frequency-hopping spread spectrum (FHSS)**
  - **direct-sequence spread spectrum (DSSS)**
  - **time-hopping spread spectrum (THSS)**
  - **chirp spread spectrum (CSS)**
  - **The combinations of these techniques**

FHSS transmitting radio signals by rapidly switching a carrier among many frequency channels (HOPS), using a pseudorandom sequence known to both transmitter and receiver.

Modulator

Original signal

Spread signal

Frequency synthesizer

Pseudorandom code generator

Frequency table

Frequency selection in FHSS

First-hop frequency

k-bit patterns

| 101 | 111 | 001 | 000 | 010 | 110 | 011 | 100 |

First selection

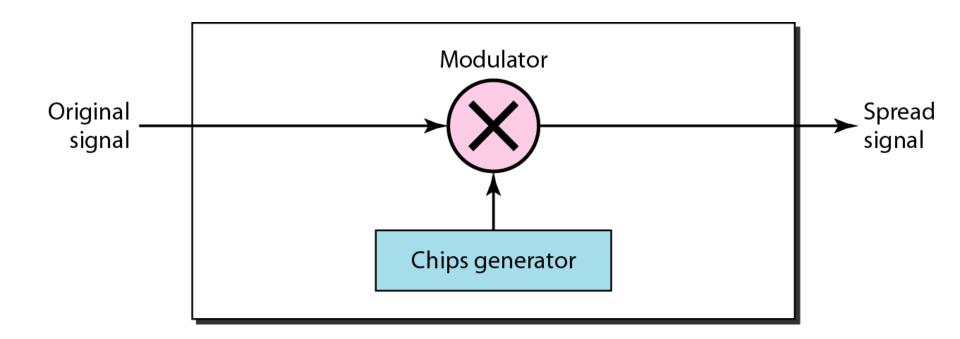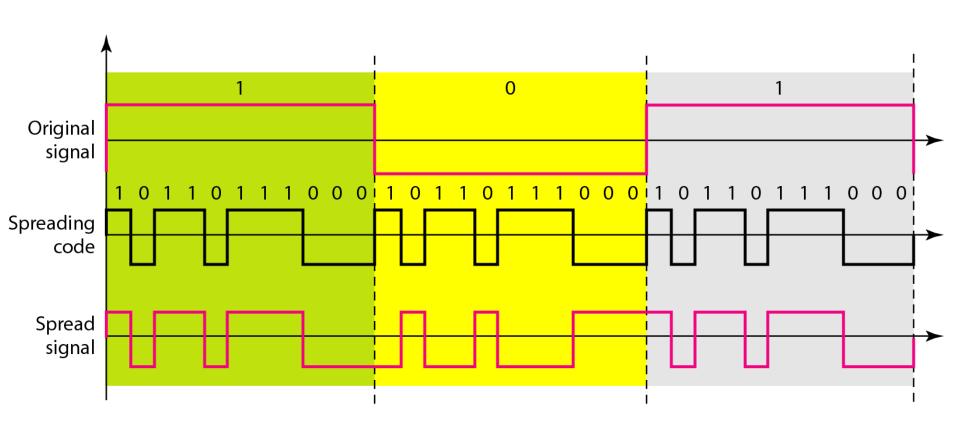| k-bit	Frequency
000	200 kHz
001	300 kHz
010	400 kHz
011	500 kHz
100	600 kHz
101	700 kHz
110	800 kHz
111	900 kHz

Frequency table

FHSS

- **The original 802.11 FHSS standard supports 1 and 2 Mbps data rate.**

- **FHSS uses the 2.402 – 2.480 GHz frequency range in the ISM band.**

  - **It splits the band into *79* non-overlapping channels with each channel *1 MHz* wide.**

  - **FHSS hops between channels at a minimum rate of 2.5 times per second. Each hop must cover at least 6 MHz.**

  - **The hopping channels for the US and Europe are shown below.**

Set	Hopping Pattern
1	{0,3,6,9,12,15,18,21,24,27,30,33,36,39,42,45,48,51,54,57,60,63,66,69,72,75}
2	{1,4,7,10,13,16,19,22,25,28,31,34,37,40,43,46,49,52,55,58,61,64,67,70,73,76}
3	{2,5,8,11,14,17,20,23,26,29,32,35,38,41,44,47,50,53,56,59,62,65,68,71,72,77}

DSSS spreads the signal by combining information *bits* (data signal) with a higher data rate bit sequence - *pseudorandom number (PN) -* called a *Chipping Code*
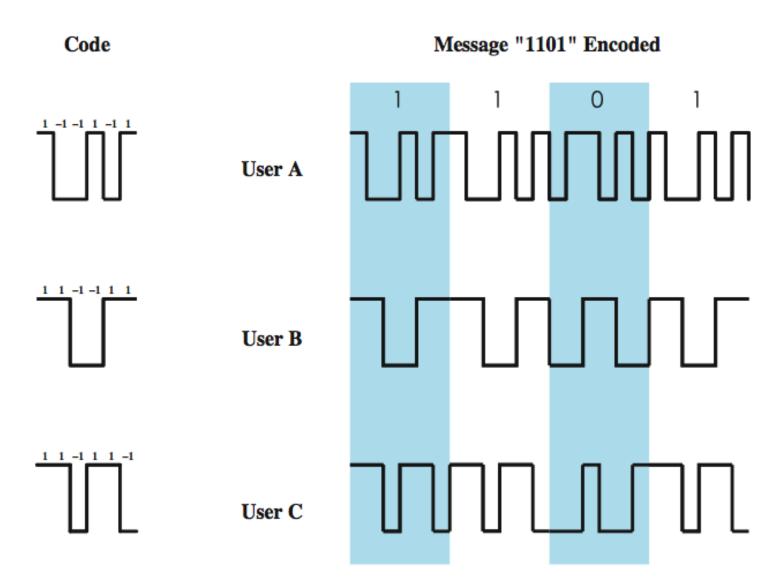
DSSS example

Code Division Multiple Access (CDMA)

- **Basic Principles of CDMA**
  - D = rate of data signal (bit data rate)
  - Break each bit into *k chips*
    - Chips are a user-specific fixed pattern (user code, chipping code)
  - Chip data rate of new channel = *kD*
- **With CDMA, the receiver can sort out transmission from the desired sender, even when there may be other users broadcasting in the same cell.**

CDMA Example

码片序列(chip sequence)

- 每个站被指派一个唯一的 *m* bit 码片序列。
  - 如发送比特 1，则发送自己的 *m* bit 码片序列。
  - 如发送比特 0，则发送该码片序列的二进制反码。
- 例如，S 站的 8 bit 码片序列是 00011011。
  - 发送比特 1 时，就发送序列 00011011,
  - 发送比特 0 时，就发送序列 11100100。
- S 站的码片序列：(–1 –1 –1 +1 +1 –1 +1 +1)

每个站分配的码片序列不仅必须各不相同，并且还必须互相正交(orthogonal)。
在实用的系统中是使用伪随机码序列。

码片序列的正交关系

- 令向量 S 表示站 S 的码片向量，令 T 表示其他任何站的码片向量。

- 两个不同站的码片序列正交，就是向量 S 和 T 的规格化内积(inner product)都是 0：

$$\mathbf{S} \bullet \mathbf{T} \equiv \frac{1}{m} \sum_{i=1}^{m} S_i T_i = 0$$
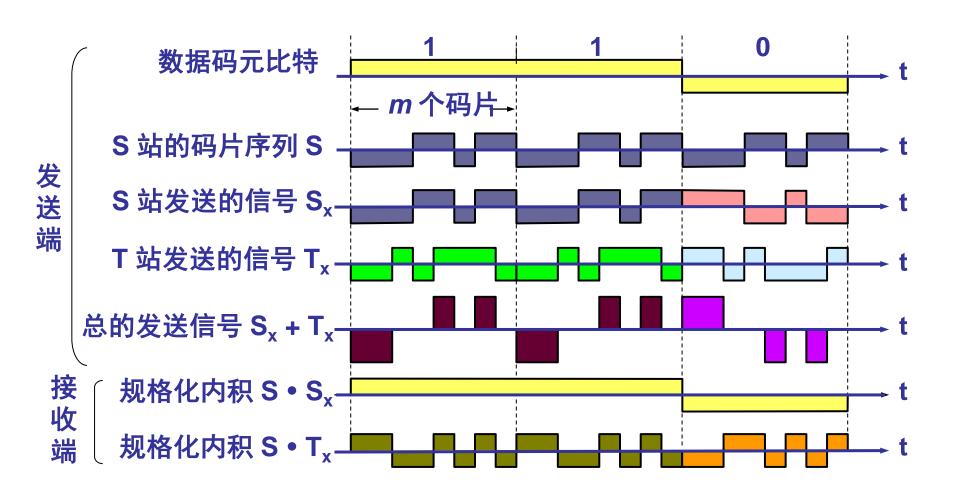
正交关系的另一个重要特性

- **任何一个码片向量和该码片向量自己的规格化内积都是1 。**

$$\mathbf{S} \bullet \mathbf{S} = \frac{1}{m} \sum_{i=1}^{m} S_i S_i = \frac{1}{m} \sum_{i=1}^{m} S_i^2 = \frac{1}{m} \sum_{i=1}^{m} (\pm 1)^2 = 1$$

- **一个码片向量和该码片反码的向量的规格化内积值是 –1 。**

CDMA 的工作原理

CDMA – Code Division Multiple Access

A: 0 0 0 1 1 0 1 1
B: 0 0 1 0 1 1 1 0
C: 0 1 0 1 1 1 0 0
D: 0 1 0 0 0 0 1 0

(a)

A: (−1 −1 −1 +1 +1 −1 +1 +1)
B: (−1 −1 +1 −1 +1 +1 +1 −1)
C: (−1 +1 −1 +1 +1 +1 −1 −1)
D: (−1 +1 −1 −1 −1 −1 +1 −1)

(b)

$$正交：\ S \bullet T = \frac{1}{m}\sum_{i=1}^{m} S_i T_i = 0$$

$$S \bullet S = 1; \ S \bullet \overline{S} = -1$$

Six examples:

− − 1 −	C	$S_1 = (−1 +1 −1 +1 +1 +1 −1 −1)$
− 1 1 −	B + C	$S_2 = (−2 \ 0 \ 0 \ 0 +2 +2 \ 0 −2)$
1 0 − −	A + B	$S_3 = (\ 0 \ 0 −2 +2 \ 0 −2 \ 0 +2)$
1 0 1 −	A + B + C	$S_4 = (−1 +1 −3 +3 +1 −1 −1 +1)$
1 1 1 1	A + B + C + D	$S_5 = (−4 \ 0 −2 \ 0 +2 \ 0 +2 −2)$
1 1 0 1	A + B + C̄ + D	$S_6 = (−2 −2 \ 0 −2 \ 0 −2 +4 \ 0)$

(c)
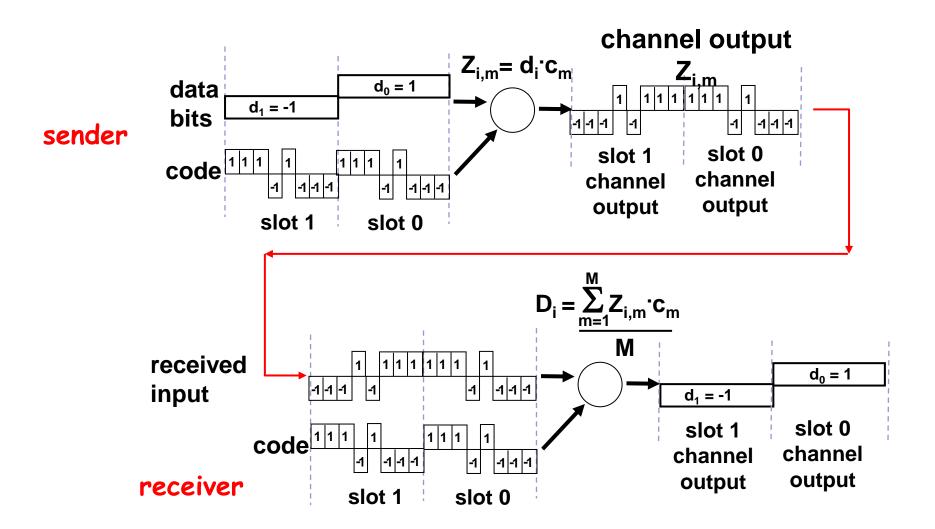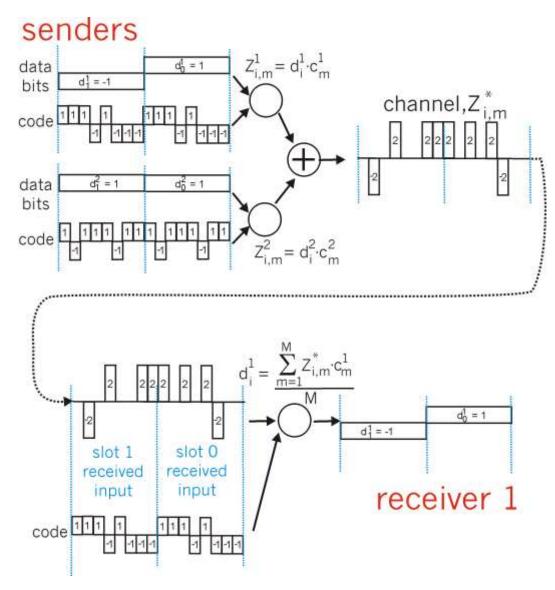
$S_1 \bullet C = (1 +1 +1 +1 +1 +1 +1 +1)/8 = 1$
$S_2 \bullet C = (2 +0 +0 +0 +2 +2 +0 +2)/8 = 1$
$S_3 \bullet C = (0 +0 +2 +2 +0 −2 +0 −2)/8 = 0$
$S_4 \bullet C = (1 +1 +3 +3 +1 −1 +1 −1)/8 = 1$
$S_5 \bullet C = (4 +0 +2 +0 +2 +0 −2 +2)/8 = 1$
$S_6 \bullet C = (2 −2 +0 −2 +0 −2 −4 +0)/8 = −1$

(d)

(a) Binary chip sequences for four stations
(b) Bipolar chip sequences
(c) Six examples of transmissions
(d) Recovery of station C's signal

CDMA – Code Division Multiple Access

CDMA – Code Division Multiple Access

RS-232-C: An Example Protocol of Physical Layer

- 规范了计算机与调制解调器间的一个串行物理接口标准。

- 电子工业协会**EIA（Electronic Industries Association)1969**年制定的标准。

RS-232-C: An Example Protocol of Physical Layer

- ## **Mechanical**

  - **25芯或9芯D型连接器，DTE侧为插针，DCE侧为插孔。**
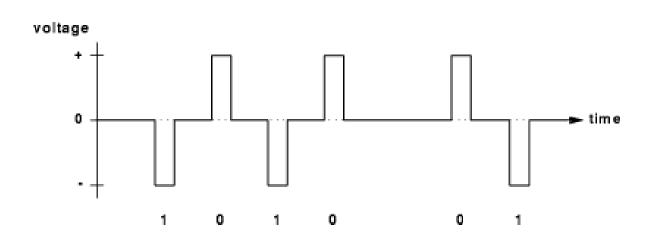
RS-232-C: An Example Protocol of Physical Layer

- **Electrical: Negative logic**
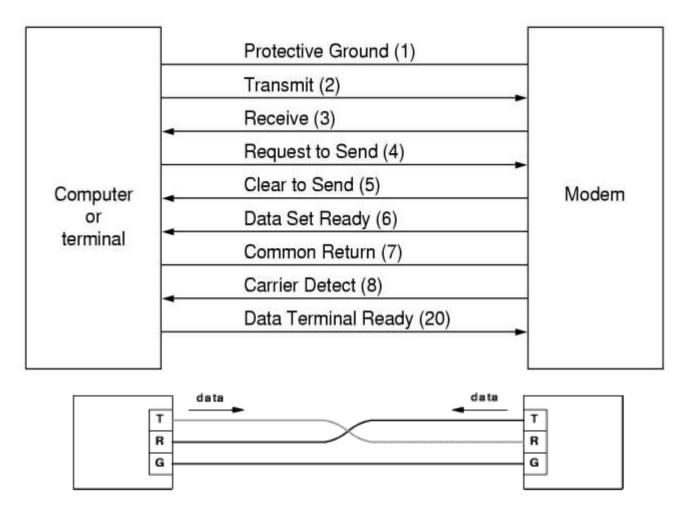
  "0": +3v ~ +15v, we use +12v

  "1": -3v ~ -15v, we use -12v

  -3v ~ +3v: not defined, illegitimate

RS-232-C: An Example Protocol of Physical Layer
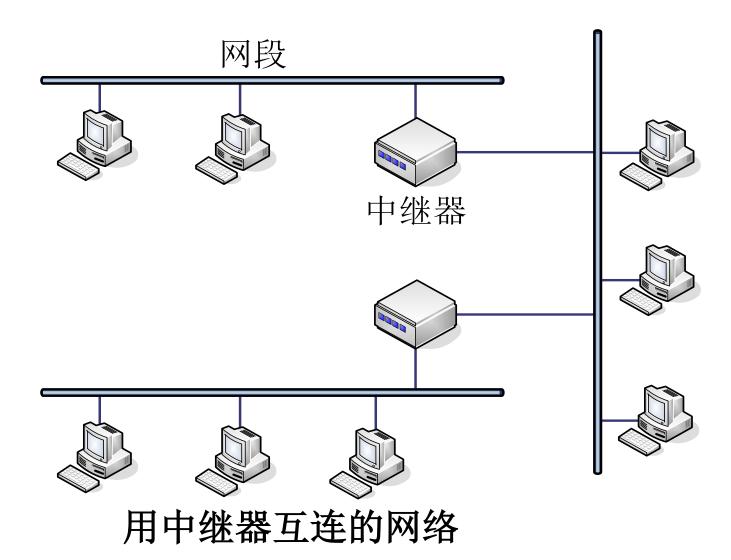
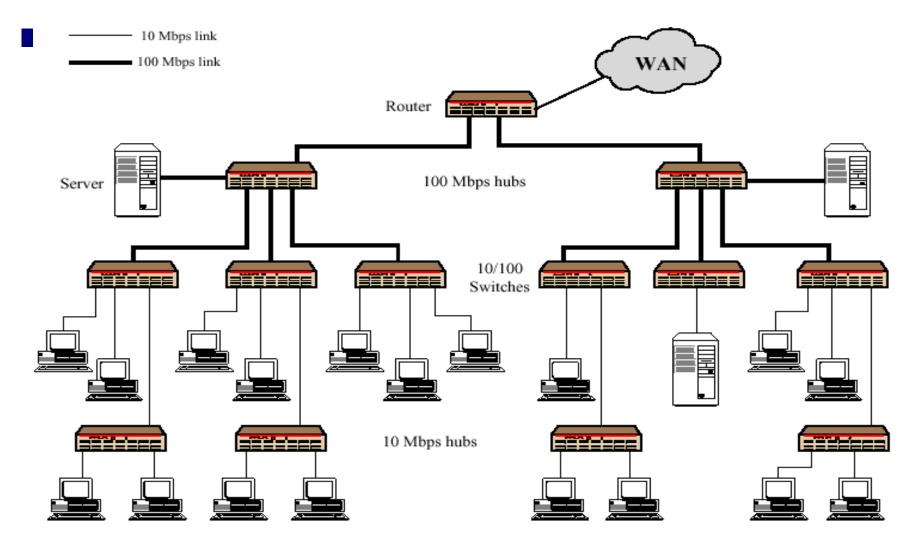## ■ **Functional**

Physical-layer Device

- **NIC (**网卡**)**
  - 典型实现物理层和数据链路层功能
- **Repeater (**中继器**)**
  - 实现信号的再生、放大和转发
  - 主要用于延迟电缆长度，扩展网段距离
  - 可以互连不同的传输介质
  - 不能隔离网段（即广播会通过中继器）
  - 互连的网络在逻辑上是一个网络（即一个物理网络系统）。（中继器对站点透明）
- **Hub (**集线器**)：** 多端口中继器

Physical-layer Device

网段

中继器

用中继器互连的网络

Physical-layer Device

- 10 Mbps link
- 100 Mbps link

WAN

Router

Server

100 Mbps hubs

10/100 Switches

10 Mbps hubs

用集线器构建的网络