

### Laboratory practice No. 4: Hash tables and trees

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### 3) Practice for final project defense presentation

**3.1** We chose a hash table as data structure because the time complexity was the priority and the hash table insertion time is  $O(1)$  and finding a required bee is also done in  $O(1)$  time as well as the deletion, so given the condition we thought that the hash table was the most optimum choice in time terms.

#### 3.4

**2.1**  $O(\log n)$

#### 3.5

**2.1 n:** nodes quantity that the tree has

### 4) Practice for midterms

#### 4.1

**4.1.1 b)** the chains that start with the same letter collide

**4.1.2 d)**  $O(1)$

#### 4.2

**4.2.1** It returns the nearest common ancestor node

**4.2.2**  $O(n)$

**4.2.3** The tree can be balanced into an AVL type

#### 4.3

**4.3.1** return true.

**4.3.2** Is  $O(n+m)$ .

#### 4.4

**4.4.1 c)**  $T(n)=2.T(n/2)+C$ , que es  $O(n)$

**4.4.2 a)**  $O(n)$

**ESTRUCTURA DE DATOS 1**  
**Código ST0245**

**4.4.3 d)** Wilkenson, Joaquina, Eustaquia, Florinda, Eustaquio, Jovín, Sufranio, Piolina, Wilberta, Piolín, Usnavy.

**4.4.4 c)** Change the order in 03, 04 y 05 lines for 03, 05, 04

**4.5**

**a)**  $p \neq \text{null}$

**b)**  $\text{toInsert} > p$

**4.6**

**4.6.1 d)** 4

**4.6.2** return new Nodo(suma);

**4.6.3**  $\text{raiz.hijos.size()} == 0$

**4.7**

**4.7.1 a)** 0, 2, 1, 7, 5, 10, 13, 11, 9, 4

**4.7.2 b)** 2

**4.7.3 d)**  $O(n)$

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