Fundamentals of Computer Algorithms Homework 1 Additional Problems

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1. Prove that
$$\sum_{k=0}^{n} k = (1/2)n(n+1)$$
.

2. Prove that
$$\sum_{k=0}^{n} a^k = \frac{a^{n+1}-1}{a-1}$$
.

- 3. For each of the sums below, (1) write a python function to evaluate it and (2) evaluate it by hand, showing your work (you need not simplify).
 - (i) $\sum_{k=0}^{100} k$.
 - (ii) $\sum_{k=0}^{100} k^2$.

(iii)
$$\sum_{k=12}^{123} k^2 + k + 1$$
.

4. Evaluate
$$\sum_{0 \leq k < n^2} \sum_{j=1}^{\lfloor \sqrt{k} \rfloor} 1.$$

5. Consider the algorithm below.

```
def S(N):
    A = []
    for n in range(2,N+1):
        do_add_n = True
    for a in A:
        if n % a == 0:
             do_add_n = False
             break
    if do_add_n:
        A.append(n)
    return A
```

- (i) Implement this code in python.
- (ii) What is S(10)?
- (iii) What is S(100)?
- (iv) What does S do?
- (v) It is a famous fact that, in the outer loop, $len(A) = O(n/\ln(n))$. Use this to estimate the O-complexity of S(N). It may help in your analysis to ignore the break statement.



Hbb/

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HW1 Additional probs

Luke Dercher

1. Proof
$$\sum_{k=0}^{n} k = \frac{N(n+1)}{2}$$

let
$$S = 1 + 2 + \cdots + n$$

 $2S = (n+1) + (n+1) + \cdots + (n+1) = n(n+1)$

$$\frac{25}{2} = 5 = \frac{n(n+1)}{2}$$

$$a \cdot 5 = a + a^2 + a^3 + \dots + a^{n+1}$$

$$a \cdot 5 - 5 = (a + a^2 + \dots + a^{n+1}) - (1 + \dots + a^n)$$

$$\frac{S(a-1)}{a-1} = S = \frac{a^{n+1}-1}{a-1}$$

$$\frac{100}{k=0} k = \frac{100(100+1)}{2} = 5050$$

3.(iii)
$$\sum_{k=12}^{123} k^2 + k + 1 = \frac{n(n+1)(2n+1)}{6} + \frac{n(n+1)}{2} + n - k$$

$$= \frac{123(123+1)(2(123)+1)}{6} + \frac{123(123+1)}{2} + 123 - 12$$

$$= 635040$$

4. Eval
$$\sum_{k=1}^{2} \sum_{0 \le k \le n^{2}} \sum_{j=1}^{2} \frac{1}{||\sum_{k=1}^{2} ||\sum_{j=1}^{2} ||\sum_{j=1}^{2}$$

$$= \sum_{k=0}^{n-1} 2k^2 + k$$

$$=2\sum_{\ell=0}^{n-1}\ell^{2}+\sum_{\ell=0}^{n-1}\ell$$

$$=2\left(\frac{n-1((n-1)+1)(2(n-1)+1)}{6}\right)+\frac{n-1((n-1)+1)}{2}$$

5. (i) see +1 Wt.px (ii)...

(IV.) S peturns a list of all of the primes < N

(V.) The inner loop is ron n times. Therefore $S(N) = O(\frac{n^2}{\ln(n)})$