

# Remote Method Invocation

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# Figure 5.1 Middleware layers

This chapter (and Chapter 6)

**Applications** 

Remote invocation, indirect communication

Underlying interprocess communication primitives:

Sockets, message passing, multicast support, overlay networks

**UDP** and TCP

Middleware layers

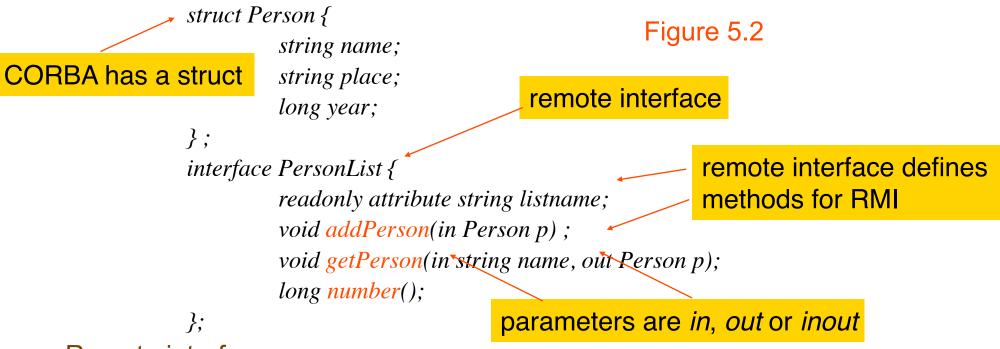
## Middleware programming models

- Distributed objects and remote object invocation is the model explained:
  - illustrated by Java RMI
- CORBA:
  - it provides remote object invocation between a client program written in one language and a server program written in another language
  - our book uses Java CORBA to illustrate the use of CORBA
  - another language commonly used in CORBA is C++
- Other programming models
  - remote event notification
  - remote SQL access
  - distributed transaction processing

#### External data representation

- It masks the differences due to different computer hardware and software.
- CORBA CDR
  - only defined in CORBA 2.0 in 1998, before that, each implementation of CORBA had an external data representation, but they could not generally work with one another. That is:
    - the heterogeneity of hardware was masked
    - but not the heterogeneity due to different programmers (until CORBA 2)
  - CORBA CDR represents simple and constructed data types (sequence, string, array, struct, enum and union)
    - note that it does not deal with objects
  - it requires an IDL specification of data to be serialised
- Java object serialisation
  - represents both objects and primitive data values
  - it uses reflection to serialise and deserialise objects
     – it does not need an IDL specification of the objects

## CORBA IDL example



- Remote interface:
  - specifies the methods of an object available for remote invocation
  - an interface definition language (or IDL) is used to specify remote interfaces. E.g. the above in CORBA IDL.
  - Java RMI would have a class for Person, but CORBA has a struct

Figure 5.3 Remote and local method invocations

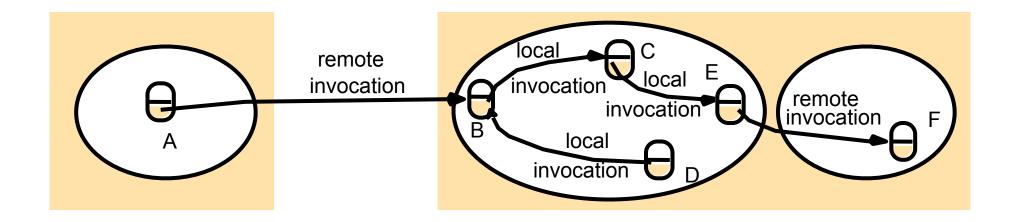


Figure 5.4 A remote object and its remote interface

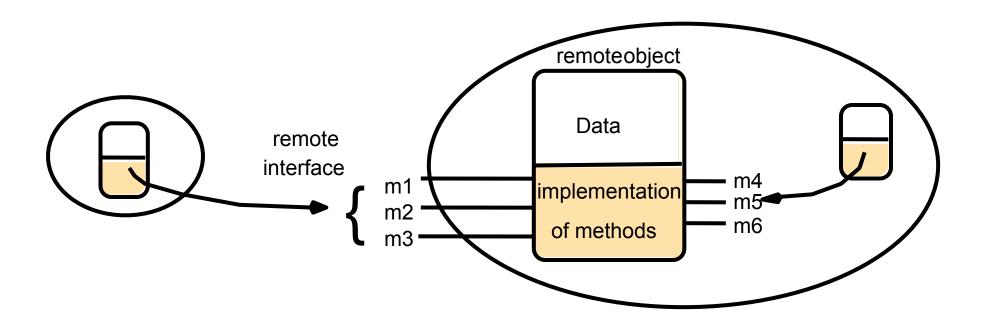


Figure 5.14 Instantiation of remote objects

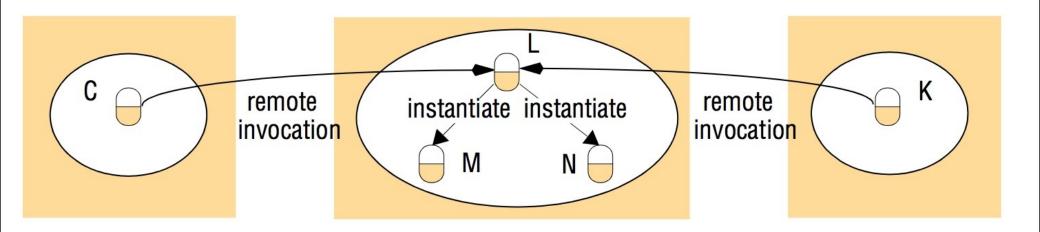


Figure 5.5 Invocation semantics

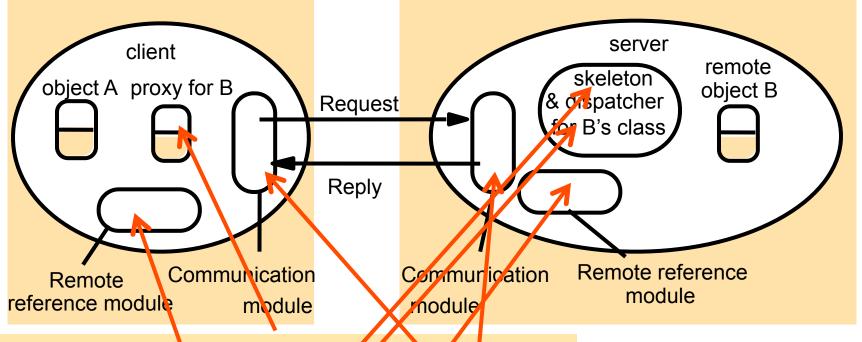
Fault tolerance measures			Invocation semantics
Retransmit request message	Duplicate filtering	Re-execute procedure or retransmit reply	
No	Not applicable	Not applicable	Maybe
Yes	No	Re-execute procedure	At-least-once
Yes	Yes	Retransmit reply	At-most-once

#### Invocation semantics: failure model

- Maybe, At-least-once and At-most-once can suffer from crash failures when the server containing the remote object fails.
- Maybe if no reply, the client does not know if method was executed or not
  - omission failures if the invocation or result message is lost
- At-least-once the client gets a result (and the method was executed at least once) or an exception (no result)
  - arbitrary failures. If the invocation message is retransmitted, the remote object may execute the method more than once, possibly causing wrong values to be stored or returned.
  - if idempotent operations are used, arbitrary failures will not occur
- At-most-once the client gets a result (and the method was executed exactly once) or an exception (instead of a result, in which case, the method was executed once or not at all)

#### The architecture of remote method invocation

## Figure 5.6



Proxy - makes RMI transparent to carries out Request-remote interface. Marshals request results. Forwards request.

translates between local and remote object references and creates remote object nvokes references. Uses remote object table

RMI software - between application level objects and communication and remote reference modules

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Figure 5.7
Role of client and server stub procedures in RPC

