

소프트웨어학과 201620914 이휘진 HW-1

① $\neg(A \vee B) \Leftrightarrow \neg A \wedge \neg B$

A	B	$\neg A$	$\neg B$	$A \vee B$	$\neg(A \vee B)$	$\neg A \wedge \neg B$	$\neg(A \vee B) \Leftrightarrow \neg A \wedge \neg B$
t	t	f	f	t	f	f	t
t	f	f	t	t	f	f	t
f	t	t	f	t	f	f	t
f	f	t	t	f	t	t	t

$$\textcircled{2} (A \vee B) \wedge (\neg B \vee C) \Rightarrow (A \vee C) \quad \Rightarrow \neg((A \vee B) \wedge (\neg B \vee C)) \vee (A \vee C)$$

A	B	C	$\neg B$	$A \vee B$	$\neg B \vee C$	$(A \vee B) \wedge (\neg B \vee C)$	$\neg((A \vee B) \wedge (\neg B \vee C))$	$A \vee C$	$(A \vee B) \wedge (\neg B \vee C) \Rightarrow (A \vee C)$
t	t	t	f	t	t	t	f	t	t
f	f	f	t	f	t	f	t	f	t
f	f	t	t	f	t	f	t	f	t
f	t	t	f	t	t	t	f	t	t
f	t	f	f	t	f	f	t	t	t
t	t	f	f	t	f	f	t	t	t
t	f	f	t	t	t	t	f	t	t
t	f	t	t	t	t	t	f	t	t

B. Transform the following statement into conjunctive normal form. (2

Points)

$$A \wedge (A \Rightarrow B) \Rightarrow B$$

$$\begin{aligned} \text{CNF)} \quad & A \wedge (\neg A \vee B) \Rightarrow B \\ & (A \wedge \neg A) \vee (A \wedge B) \\ & (A \wedge B) \Rightarrow B \\ & \begin{aligned} & \neg(A \wedge B) \vee B \\ & = (\neg A \vee \neg B) \vee B \\ & = (\neg A \vee B) \vee (\neg B \vee B) \\ & = (\neg A \vee B) \leftarrow \text{disjunction} \\ & = \neg(A \wedge \neg B) \\ & \quad \text{answer} \end{aligned} \end{aligned}$$

C. Consider the following knowledge base (KB).

KB = {A, B, C, D, E, (A ∧ B ∧ C ⇒ F), (C ∧ F ∧ E ⇒ H)}

Show $(KB \models H)$ by SLD resolution-refutation. (4 Points)

Claim. $KB \models H \Rightarrow KB \wedge \neg H$

Res(6, 1) : $(B \wedge C \Rightarrow F)_8$

Res(8, 2) : $(C \Rightarrow F)_9$

Res(9, 3) : $(F)_{10}$

Res(10, 11) : $(C \wedge E \Rightarrow H)_{11}$

Res(11, 3) : $(E \Rightarrow H)_{12}$

Res(12, 5) : $(H)_{13}$

+ $(\neg H)_{14}$

Res(14, 13) : $()_{15}$

→ From there we derive the empty clause using resolution
Thus the claim has been proved.