



Interactive Oracle Proofs

A learning group for ZK and SNARK application development

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Logistics: ZK Learning Group

Every month, third thursday in 2025, from 18 (CET)

One hour, presentation + short discussion

Different topics on zero knowledge proof,

- mostly from programmer and application developers perspective
- with some theory

Coordination:

- Discord channel: LF Decentralized Trust

<https://discord.com/channels/905194001349627914/1329201532628898036>

- Meetup.com: <https://www.meetup.com/lfdt-hungary/events/305634614/>

- Repo with all the contents:<https://github.com/LF-Decentralized-Trust-labs/>

<https://github.com/Daniel-Szego/zk-leraning-group>

Quizzes and small programming challenges, LFDT merchs at the end



Logistics: Hunting for the SNARK

February - Introduction, Theory : Definitions and building
blocks

March - Theory : Polynomial commitments

April - Theory : Interactive oracle proofs

May - Programming : Circom

June - Programming : Circom

July - Programming : Noir

August - Programming : Noir

September : Applications : Off-chain transaction

October : Applications : Proving solvency

November : Applications : Rollup

December : Wrap up, Applications





Agenda

- (zk)SNARK
- *Elements of building a SNARK*
- *Interactive Oracle Proof (IOP)*
- *Building blocks*
- PLOK
- *Computational trace*
- *Constructing polynomials*
- *Summary*
- *Literature and Links*
- *Q&A and discussion*

(zk)SNARK - Succinct Non-interactive ARgument of Knowledge

Computation: arithmetic circuit : $C(x, w) \rightarrow F$

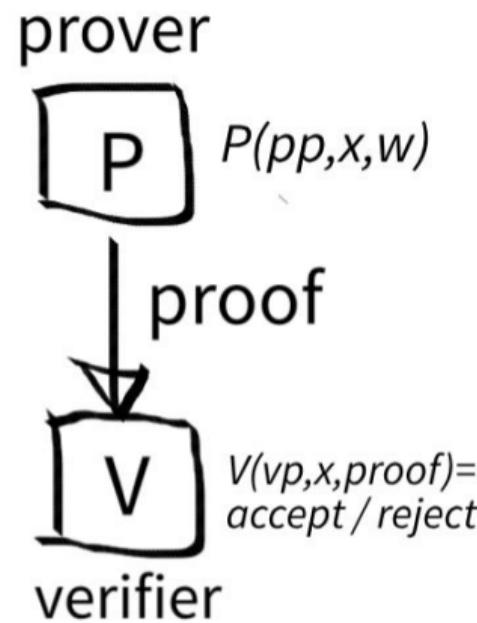
- x public input
- w private input, witness
- high level computation
- arithmetic circuit
- polynomials

Prover algorithm: $P(pp, x, w) \rightarrow proof$

Verifier algorithm: $V(vp, x, proof) \rightarrow accept / reject$

Properties:

- *Succinct*:
- *Complete*:
- *Knowledge sound*:
- *Zero knowledge*



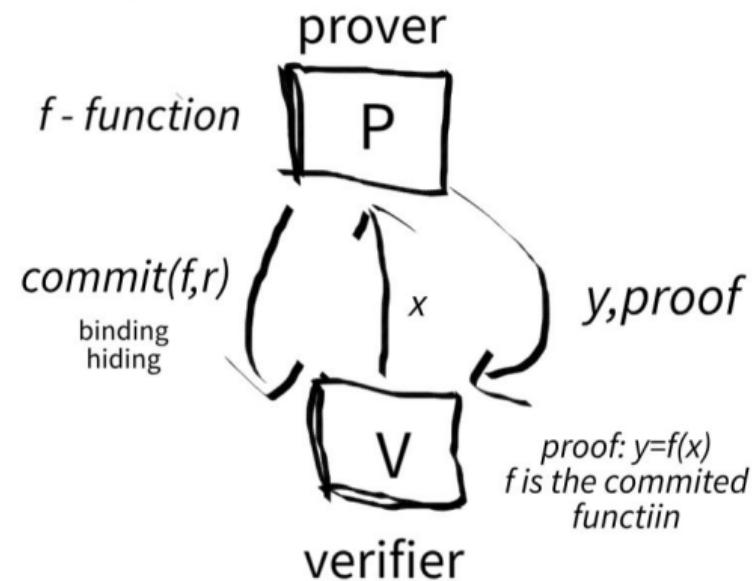
Elements of building a SNARK

Elements of a SNARK:

- Polynomial commitment
- Interactive Oracle Proof
- Fiat-Shamir

Functional commitment:

- Having a family of functions F.
- Prover: choose an f function from the F function family
- Committing f function to the verifier, com commitment
- Verifier sends an x point of the function.
- Verifier sends y and proof that $y=f(x)$ and that f is in the function family and that f is the committed function



Interactive Oracle Proof (IOP)

Proving that $C(x, w) = 0$

Committing f_0, f_1, \dots, f_N functions

Committed functions can be evaluated at chosen points with the help of functional commitments

Repeated steps, n times:

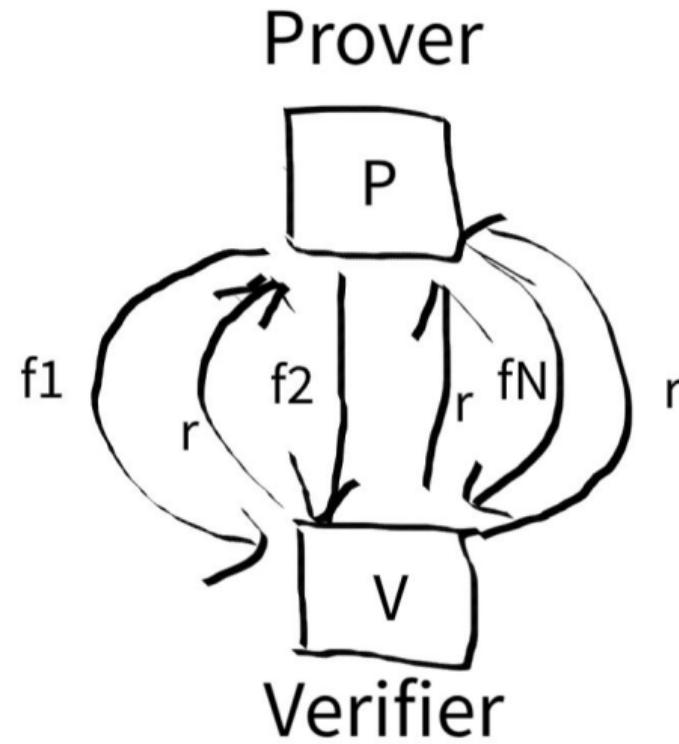
- Prover sends committed functions
- Verifier responds with random values

Final step: verification:

- functions / function commitments
- public input
- opening functions at certain points

Properties:

- complete
- knowledge sound
- zero knowledge



Building blocks

Testing if a **polynomial is zero**:

- If a truly randomly chosen point evaluates to 0, the whole polynomial is zero with a very high probability.
- d/p : d degree, d roots, p is the domain

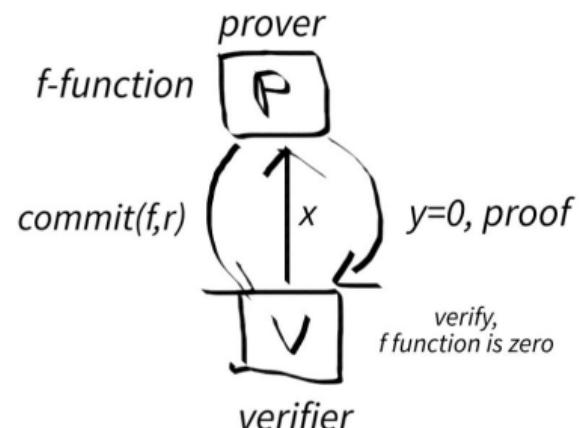
Testing if two **polynomials are equal**:

- if two polynomials f and g are equal in a randomly chosen r point $f(r)=g(r)$, then the two polynomials are equal with a very high probability.
- $f(r)-g(r)$ zero test

Testing if a polynomial is **zero on a set**: quotation polynomial

Sum check: sum of certain input values of a committed polynomial is a given value.

Prod check: product of certain input values of a committed polynomial is a given value.



PLONK - constructing computational trace

Constructing a polynomial IOP

for a $C(x, w) = 0$ circuit

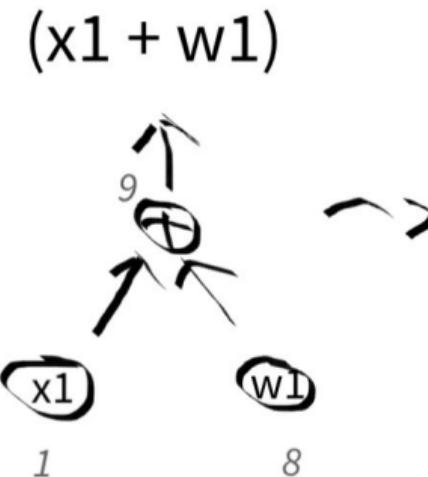
Computational trace

- table
- first row is the input
- for each gate, there is a left input, right input and output

Using computational trace
instead of arithmetic circuit

Proving if the computational
trace is correct

Proving if the output is zero



inputs:			1	8	
Gate 0:			1	8	9
Gate 1:			left input	right input	output
Gate 2:			left input	right input	output
...					
					↑ <i>output of the circuit</i>

PLONK - constructing polynomial

Computational trace to polynomial:

Constraints to the polynomial

Encoding all inputs:

- $P(w \text{ to the } -i) = i$ for all i input

Encoding all wires:

- $P(w \text{ to the } 3l) = \text{left input of the } l \text{ gate}$
- $P(w \text{ to the } 3l+1) = \text{right input of the } l \text{ gate}$
- $P(w \text{ to the } 3l+2) = \text{output of the } l \text{ gate}$

Prover can use the Fast Fourier

Transformation to construct coefficients of
the polynomial

$$P(w^{-1})=1, \quad P(w^2)=8$$

$$P(w^0)=1, \quad P(w^1)=8, \quad P(w^2)=9$$

$$P(w^3)=\text{left}, \quad P(w^4)=\text{right}, \quad P(w^5)=\text{output}$$

inputs: 1 8

Gate 0: 1 8 9

Gate 1: left right output
 input input output

Gate 2: left right output
 input input output

...



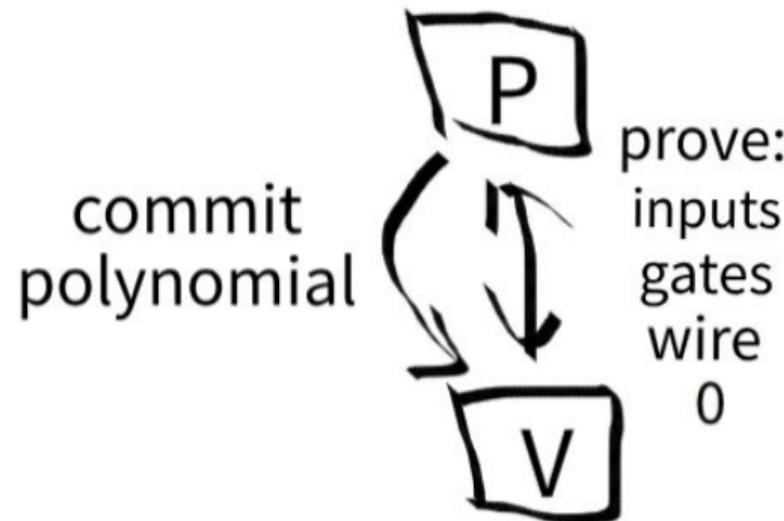
output of the
circuit

PLONK - Proving validity

Proving validity of the created polynomial:

- All **inputs** are encoded correctly by the polynomial
- **Gates** are evaluated correctly
- **Wiring** is implemented correctly
- **Output** of the last gate is **0**

Proving that the last gate is evaluated to 0,
opening the polynomial at the $P(w \text{ to the } 3l+2) =$ and testing if that is zero



PLONK - Proving validity of gates and inputs

Proving correct input:

Additional V polynomial for encoding all inputs

- $V(w \text{ to the } -i) = i$ for all i input

Prove that the V input polynomial and the P committed polynomial agree on all input points: Zero on a set test $P-V$

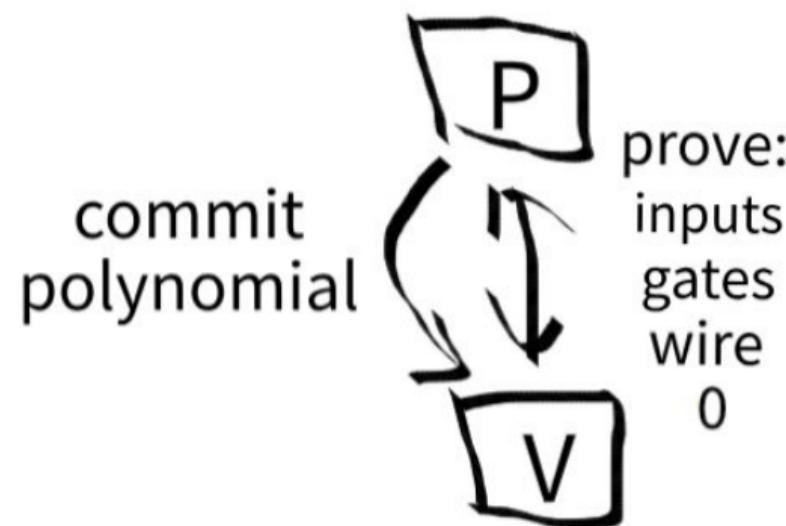
Proving correctness of gates:

Selector polynomial S :

- $S(w \text{ to the } 3l) = 1$ if it is an addition gate
- $S(w \text{ to the } 3l) = 0$ if it is a multiplication gate

Combining with the P polynomial and using the zero set test

Independent from the input values



PLONK - Proving validity of wiring

Proving correct wiring:

Equalities for the wiring
constraint

Proving wiring constraints

W wiring polynomial:

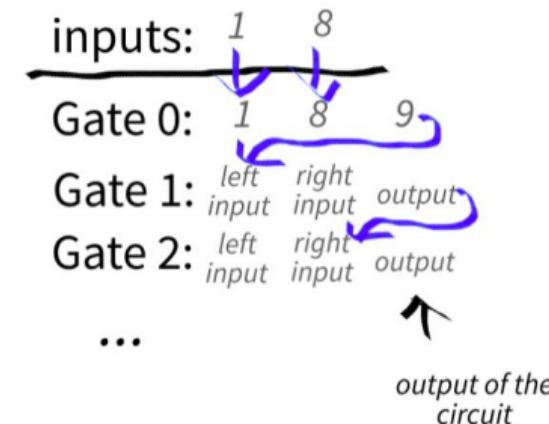
Rotation of coefficients

Independent from the input
values

Can be computed at the setup

$$P(w^{-1}) = P(w^0) \quad P(w^{-2}) = P(w^1)$$

$$P(w^3) = \dots$$



PLONK - Summary

Setup / preprocessing:

- compute wiring polynomial
- compute selector polynomial

Prover:

- build P polynomial for the computation
- trace and commits
- prove inputs
- prove gates
- prove wiring
- prove output is zero

Using the Fiat Shamir transformation to make it non-interactive

prover



$P(pp, x, w)$

proof



$V(vp, x, proof) =$
accept / reject

verifier

Challenge



Quiz:

Will be posted in the discord channel:

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Links, Resources, Literature

A guide to Zero Knowledge Proofs (Part 2)

https://medium.com/@Luca_Franceschini/a-guide-to-zero-knowledge-proofs-part-2-7904dee9758d

Interactive Oracle Proofs

<https://www.iacr.org/archive/tcc2016b/99850156/99850156.pdf>

What is PLONK

https://medium.com/@Luca_Franceschini/what-is-plonk-29c56f326cf6

Plonk Interactive Oracle Proofs (IOP)

<https://hackmd.io/@0xsachink/ByuqZfD63>

Proofs, Arguments, and Zero-Knowledge, Chapter 4

<https://people.cs.georgetown.edu/jthaler/ProofsArgsAndZK.pdf>

Happy Hunting for the SNARK :)

Q & A

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