



Polynomial Commitments

A learning group for ZK and SNARK application development

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Logistics: ZK Learning Group

Every month, third thursday in 2025, from 18 (CET)

One hour, presentation + short discussion

Different topics on zero knowledge proof,

- mostly from programmer and application developers perspective
- with some theory

Coordination:

- Discord channel: LF Decentralized Trust

<https://discord.com/channels/905194001349627914/1329201532628898036>

- Meetup.com: <https://www.meetup.com/lfdt-hungary/events/305634614/>

- Repo with all the contents:<https://github.com/LF-Decentralized-Trust-labs/>

<https://github.com/Daniel-Szego/zk-leraning-group>

Quizzes and small programming challenges, LFDT merchs at the end



Logistics: Hunting for the SNARK

February - Introduction, Theory : Definitions and building blocks

March - Theory : Polynomial commitments

April - Theory : Interactive oracle proofs

May - Programming : Circom

June - Programming : Circom

July - Programming : Noir

August - Programming : Noir

September : Applications : Off-chain transaction

October : Applications : Proving solvency

November : Applications : Rollup

December : Wrap up, Applications

Subject to change based on community discussion





Agenda

- zkSNARK
- *Development flow*
- *Elements of building a SNARK*
- Commitments
- Functional commitments
- Types of functional commitment
- KZG10
- Protocols with functional commitment
- Fiat-Shamir transformation
- Quiz, literature, links and Q and A

(zk)SNARK - Succinct Non-interactive ARgument of Knowledge

Computation: arithmetic circuit : $C(x, w) \rightarrow F$

- x public input
- w private input, witness

Prover algorithm: $P(pp, x, w) \rightarrow \text{proof}$

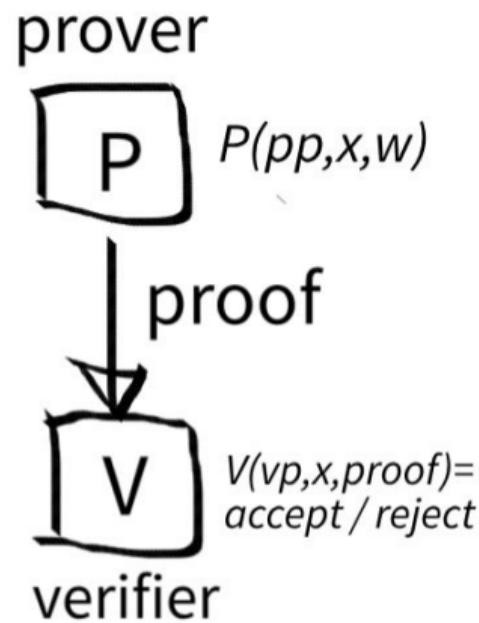
Verifier algorithm: $V(vp, x, \text{proof}) \rightarrow \text{accept / reject}$

Succinct: small proof in size, easy / fast to verify

Complete: if there is a proof, it is accepted by the prover by very high probability

Knowledge sound: If the verifier accepts, then the prover “knows” a w witness (secret input), that satisfies the computation. (malicious prover can not convince the verifier)

Zero knowledge (optional): the proof reveals nothing from the witness (secret input)



Development flow

High level language:

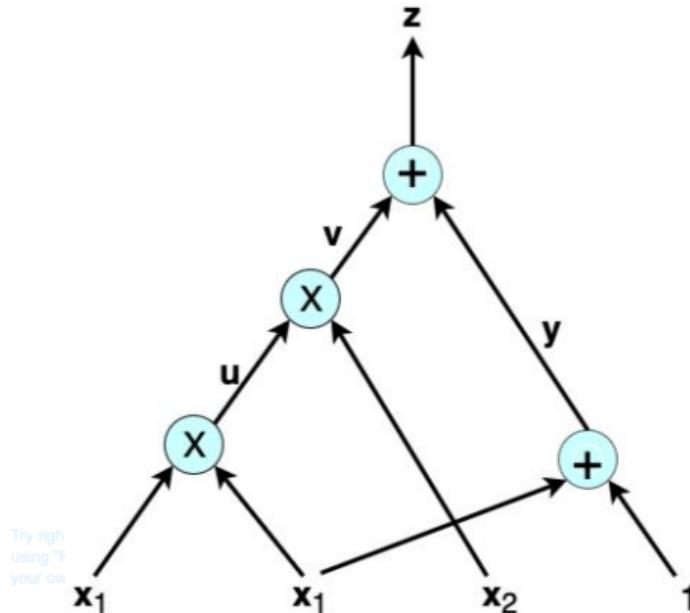
- Likie Cirom, Noir

Low level abstraction:

- Arithmetic circuit
- DAG: directed acyclic graphs
- Inputs, gates, constants, connections (wires)
- Nodes are labelled with “+” addition and “x” multiplication

Polynomials:

- R1CS, rank one constraint
- $a \times b = g$
- a, b and g are affine combination of variables
- examples: $w_1 \times (w_2 - w_3 + 1) = x_1$

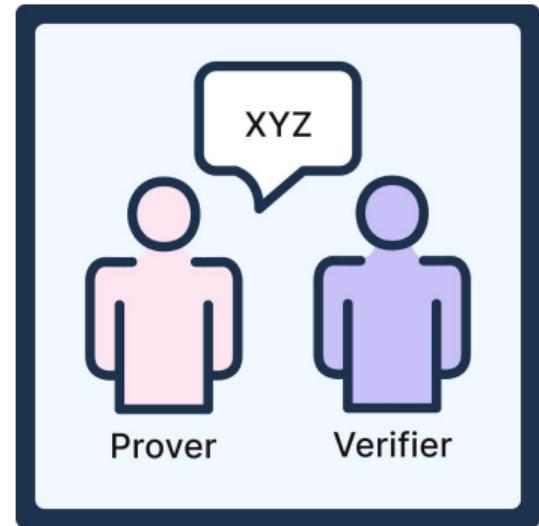


Elements of building a SNARK

Polynomial commitment scheme: get succinct interactive argument: Bind and commit a polynomial.

Interactive Oracle Proof (IOP): R1CS or circuit satisfiability. Evaluating (committed) polynomials in one or multiply points in multiply rounds, extend polynomial commitment to general circuit satisfiability.

Fiat-Shamir: transforming a interactive proofs to non-interactive:



Commitments

Commit to a chosen value (or chosen statement), keeping it hidden to others, reveal the committed value later:

1. the **commit** phase during which a value is chosen and committed to
2. the reveal phase during which the value is revealed by the sender, then the receiver **verifies** its authenticity

Properties:

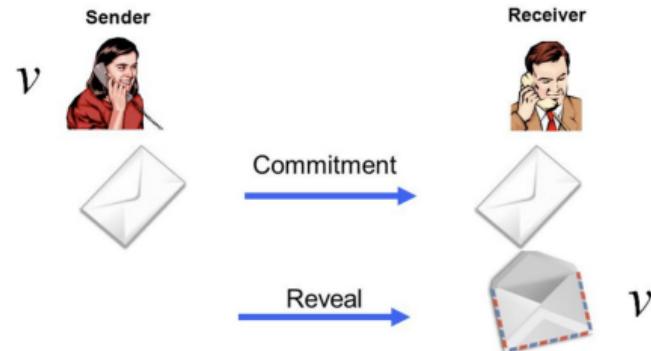
- **hiding**: the committed information can not be revealed
- **binding**: the prover can not fake the committed message

Example: commitment with hash function (hash commitment)

- commit: commitment=sha(message)
- verify: send message, verify if commitment=sha(message)

Commitment Scheme

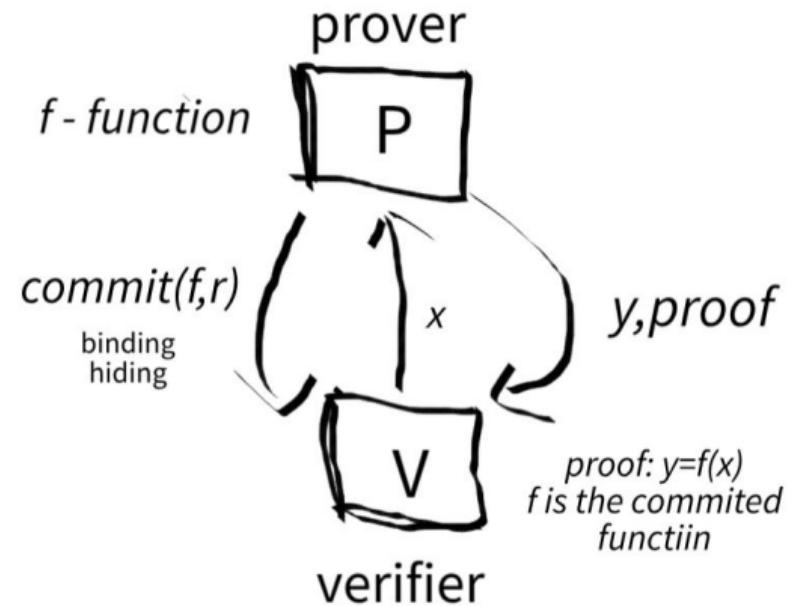
The “digital analogue” of sealed envelopes.



Functional commitment

Committing a function and opening at in a point:

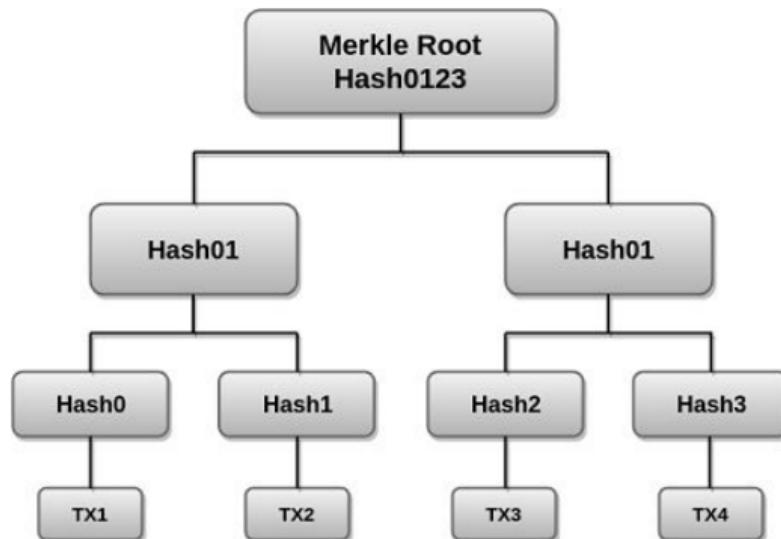
- Having a family of functions F .
- Prover: choose an f function from the F function family
- Committing f function to the verifier, com commitment
- Verifier sends an x point of the function.
- Verifier sends y and proof that $y=f(x)$ and that f is in the function family and that f is the committed function



Types of functional commitment

Different types of functional commitments:

- Polynomial commitment: univariate polynomial (at most degree d).
 - Bilinear groups: KZG 10
 - Hash functions: FRI
 - Elliptic curves: Bulletproof
- Multilinear commitment: several variable but the degree is at most one.
- Vector commitments: committing a vector and opening it at a certain point / index.
E.g. merkle tree.
- Inner product commitment



KZG10 system (Kate-Zaverucha-Goldberg)

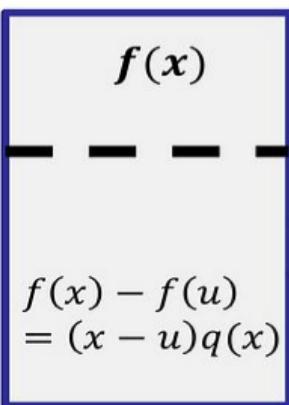
KZG polynomial commitment

Univariate polynomials of degree $\leq d$

Prover

$$gp = (g, g^\tau, g^{\tau^2}, \dots, g^{\tau^d})$$

Verifier

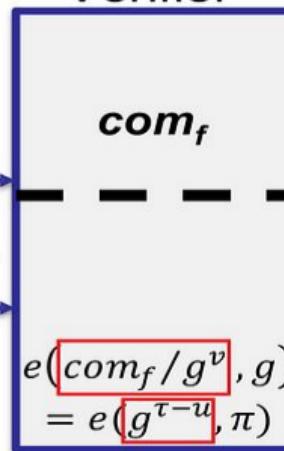


$$com_f = g^{f(\tau)}$$

u

$$f(x) - f(u) = (x - u)q(x)$$

$$v, \text{ proof } \pi = g^{q(\tau)}$$



Protocols with functional commitment

Testing if a **polynomial is zero**:

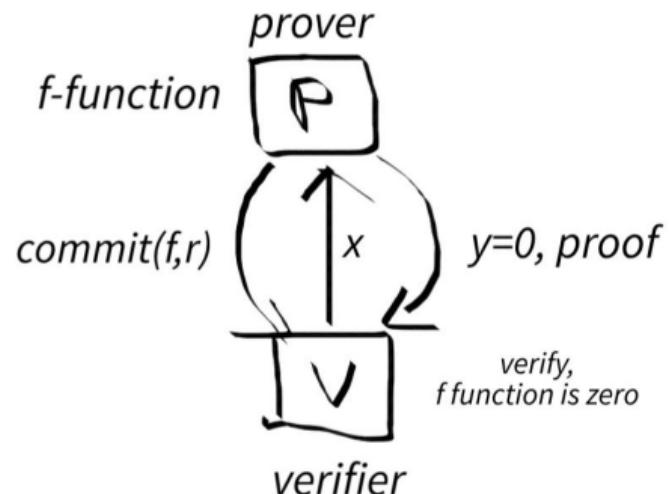
- If a truly randomly chosen point evaluates to 0, the whole polynomial is zero with a very high probability.
- d/p : d degree, d roots, p is the domain

Testing if two **polynomials are equal**:

- if two polynomials f and g are equal in a randomly chosen r point $f(r)=g(r)$, then the two polynomials are equal with a very high probability.
- $f(r)-g(r)$ zero test

Sum check: sum of certain input values of a committed polynomial is a given value.

Prod check: product of certain input values of a committed polynomial is a given value.



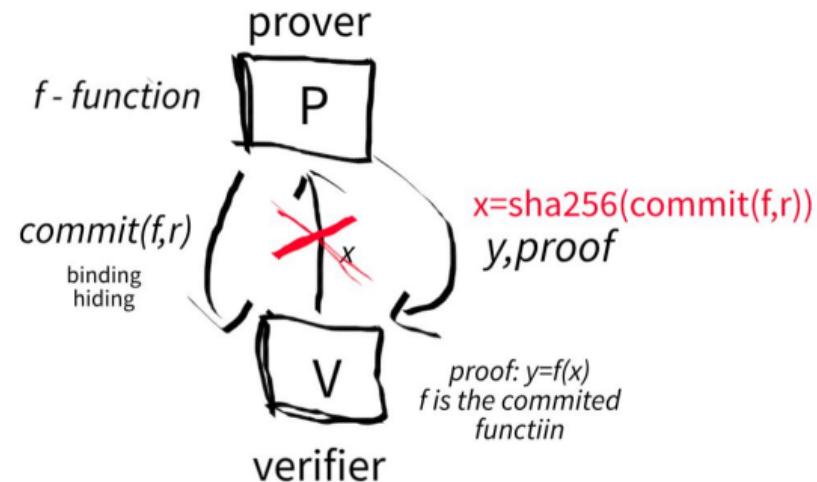
Fiat-Shamir transformation

Transforming interactive certain, random challenge, proofs to non-interactive.
(flip coin protocols)

Proof schema:

- The prover first generates some random value, the commitment, and sends it to the verifier.
- The verifier responds with a challenge value generated uniformly at random.
- The prover computes the final proof based on both the commitment and challenge.

The prover can compute this random value themselves by using a random function, such as a cryptographic hash function.



Challenge

Quiz:

Will be posted in the discord channel:

<https://discord.com/channels/905194001349627914/1329201532628898036>

Links, Resources, Literature

Commitment schemes:

<https://medium.com/@icostan/commitment-schemes-8b523d48aa1e>

Functional Commitments: ZK under a different lens

<https://geometry.xyz/notebook/functional-commitments-zk-under-a-different-lens>

Explaining KZG Commitment with Code Walkthrough

<https://kaijuneer.medium.com/explaining-kzg-commitment-with-code-walkthrough-216638a620c9>

Exploring Elliptic Curve Pairings

<https://medium.com/@VitalikButerin/exploring-elliptic-curve-pairings-c73c1864e627>

Fiat-Shamir transformation

<https://www.zkdocs.com/docs/zkdocs/protocol-primitives/fiat-shamir/>

Happy Hunting for the SNARK :)

Q & A

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Merkle tree - naive polynomial commitment

