



# ***Polynomial Commitments***

A learning group for ZK and SNARK application development

***Daniel Szego***

In: <https://www.linkedin.com/in/daniel-szego/>



# Logistics: ZK Learning Group

Every month, third thursday in 2025, from 18 (CET)

One hour, presentation + short discussion

Different topics on zero knowledge proof,

- mostly from programmer and application developers perspective
- with some theory

Coordination:

- Discord channel: LF Decentralized Trust

<https://discord.com/channels/905194001349627914/1329201532628898036>

- Meetup.com: <https://www.meetup.com/lfdt-hungary/events/305634614/>

- Repo with all the contents: <https://github.com/LF-Decentralized-Trust-labs/>  
<https://github.com/Daniel-Szego/zk-learning-group>

Quizzes and small programming challenges, LFDT merchs at the end



# Logistics: Hunting for the SNARK

February - Introduction, Theory : Definitions and building blocks

**March** - Theory : Polynomial commitments

April - Theory : Interactive oracle proofs

May - Programming : Circom

June - Programming : Circom

July - Programming : Noir

August - Programming : Noir

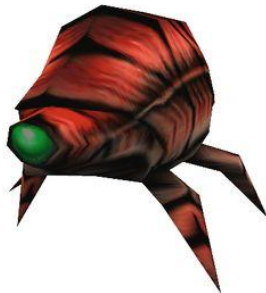
September : Applications : Off-chain transaction

October : Applications : Proving solvency

November : Applications : Rollup

December : Wrap up, Applications

*Subject to change based on community discussion ....*



# Agenda



- zkSNARK
- *Development flow*
- *Elements of building a SNARK*
- Commitments
- Functional commitments
- Types of functional commitment
- KZG10
- Protocols with functional commitment
- Fiat-Shamir transformation
- Quiz, literature, links and Q and A

## (zk)SNARK - Succinct Non-interactive ARgument of Knowledge

**Computation:** arithmetic circuit :  $C(x, w) \rightarrow F$

-  $x$  public input

-  $w$  private input, witness

**Prover** algorithm:  $P(pp, x, w) \rightarrow \text{proof}$

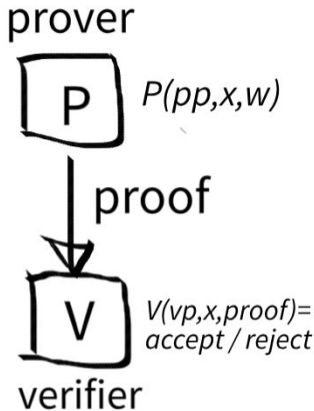
**Verifier** algorithm:  $V(vp, x, \text{proof}) \rightarrow \text{accept} / \text{reject}$

**Succinct:** small proof in size, easy / fast to verify

**Complete:** if there is a proof, it is accepted by the prover by very high probability

**Knowledge sound:** If the verifier accepts, then the prover “knows” a  $w$  witness (secret input), that satisfies the computation. (malicious prover can not convince the verifier)

**Zero knowledge (optional):** the proof reveals nothing from the witness (secret input)



# Development flow

## High level language:

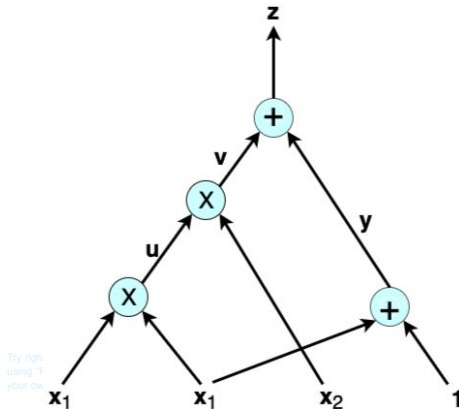
- Like Cirom, Noir

## Low level abstraction:

- Arithmetic circuit
- DAG: directed acyclic graphs
- Inputs, gates, constants, connections (wires)
- Nodes are labelled with “+” addition and “x” multiplication

## Polynomials:

- R1CS, rank one constraint
- $a \times b = g$
- $a, b$  and  $g$  are affine combination of variables
- examples:  $w_1 \times (w_2 - w_3 + 1) = x_1$

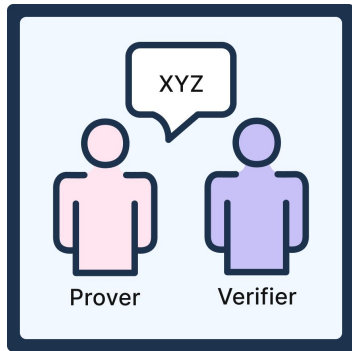


# Elements of building a SNARK

**Polynomial commitment** scheme: get succinct interactive argument: Bind and commit a polynomial.

**Interactive Oracle Proof** (IOP): R1CS or circuit satisfiability. Evaluating (committed) polynomials in one or multiply points in multiply rounds, extend polynomial commitment to general circuit satisfiability.

**Fiat-Shamir**: transforming a interactive proofs to non-interactive:



# Commitments

Commit to a chosen value (or chosen statement), keeping it hidden to others, reveal the committed value later:

1. the **commit** phase during which a value is chosen and committed to
2. the reveal phase during which the value is revealed by the sender, then the receiver **verifies** its authenticity

Properties:

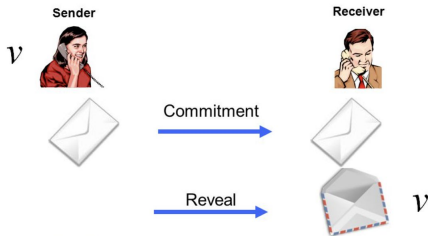
- **hiding**: the committed information can not be revealed
- **binding**: the prover can not fake the committed message

Example: commitment with hash function (hash commitment)

- commit:  $\text{commitment} = \text{sha}(\text{message})$
- verify: send message, verify if  $\text{commitment} = \text{sha}(\text{message})$

## Commitment Scheme

The “digital analogue” of sealed envelopes.

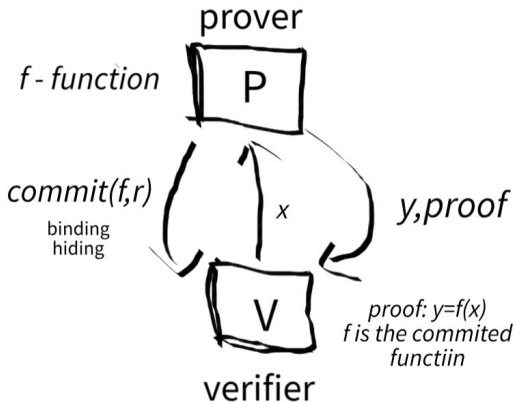




# Functional commitment

Committing a function and opening at in a point:

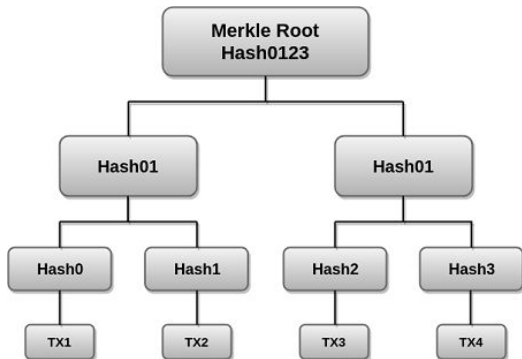
- Having a family of functions  $F$ .
- Prover: choose an  $f$  function from the  $F$  function family
- Committing  $f$  function to the verifier,  $com$  commitment
- Verifier sends an  $x$  point of the function.
- Verifier sends  $y$  and proof that  $y=f(x)$  and that  $f$  is in the function family and that  $f$  is the committed function



# Types of functional commitment

Different types of functional commitments:

- Polynomial commitment: univariate polynomial (at most degree  $d$ ).
  - Bilinear groups: KZG 10
  - Hash functions: FRI
  - Elliptic curves: Bulletproof
- Multilinear commitment: several variable but the degree is at most one.
- Vector commitments: committing a vector and opening it at a certain point / index.  
E.g. merkle tree.
- Inner product commitment



# KZG10 system (Kate-Zaverucha-Goldberg)

## KZG polynomial commitment

Univariate polynomials of degree  $\leq d$

$$gp = (g, g^\tau, g^{\tau^2}, \dots, g^{\tau^d})$$

Prover

$$f(x)$$

$$f(x) - f(u) = (x - u)q(x)$$

$$com_f = g^{f(\tau)}$$

$$u$$

$$v, \text{ proof } \pi = g^{q(\tau)}$$

Verifier

$$com_f$$

$$e(com_f / g^v, g) = e(g^{\tau-u}, \pi)$$

# Protocols with functional commitment

Testing if a **polynomial is zero**:

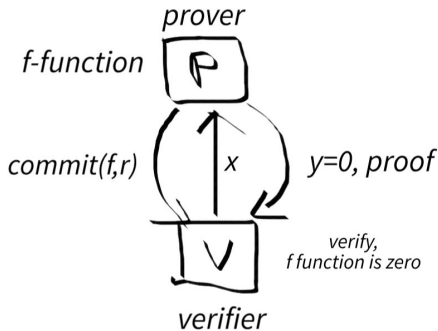
- If a truly randomly chosen point evaluates to 0, the whole polynomial is zero with a very high probability.
- $d/p$  :  $d$  degree,  $d$  roots,  $p$  is the domain

Testing if two **polynomials are equal**:

- if two polynomials  $f$  and  $g$  are equal in a randomly chosen  $r$  point  $f(r)=g(r)$ , then the two polynomials are equal with a very high probability.
- $f(r)-g(r)$  zero test

**Sum check**: sum of certain input values of a committed polynomial is a given value.

**Prod check**: product of certain input values of a committed polynomial is a given value.



# Fiat-Shamir transformation

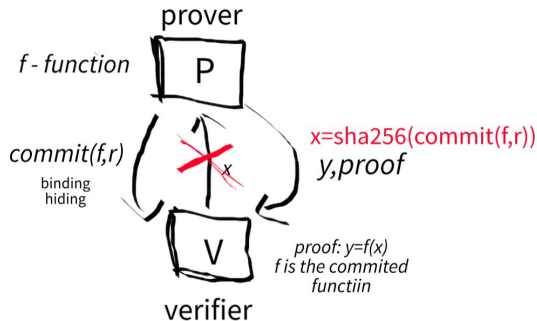
Transforming interactive certain, random challenge, proofs to non-interactive.

**(flip coin protocols)**

Proof schema:

- The prover first generates some random value, the commitment, and sends it to the verifier.
- The verifier responds with a challenge value generated uniformly at random.
- The prover computes the final proof based on both the commitment and challenge.

The prover can compute this random value themselves by using a random function, such as a cryptographic hash function.



# Challenge



## Quiz:

Will be posted in the discord channel:

<https://discord.com/channels/905194001349627914/1329201532628898036>

# Links, Resources, Literature



*Commitment schemes:*

<https://medium.com/@icostan/commitment-schemes-8b523d48aa1e>

*Functional Commitments: ZK under a different lens*

<https://geometry.xyz/notebook/functional-commitments-zk-under-a-different-lens>

*Explaining KZG Commitment with Code Walkthrough*

<https://kaijuneer.medium.com/explaining-kzg-commitment-with-code-walkthrough-216638a620c9>

*Exploring Elliptic Curve Pairings*

<https://medium.com/@VitalikButerin/exploring-elliptic-curve-pairings-c73c1864e627>

*Fiat-Shamir transformation*

<https://www.zkdocs.com/docs/zkdocs/protocol-primitives/flat-shamir/>



# *Happy Hunting for the SNARK :)*

## **Q & A**

Daniel Szego

In: <https://www.linkedin.com/in/daniel-szego/>





# Merkle tree - naive polynomial commitment

