



# *Interactive Oracle Proofs*

A learning group for ZK and SNARK application  
development

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# Logistics: ZK Learning Group

Every month, third thursday in 2025, from 18 (CET)

One hour, presentation + short discussion

Different topics on zero knowledge proof,

- mostly from programmer and application developers perspective
- with some theory

Coordination:

- Discord channel: LF Decentralized Trust

<https://discord.com/channels/905194001349627914/1329201532628898036>

- Meetup.com: <https://www.meetup.com/lfdt-hungary/events/305634614/>

- Repo with all the contents: <https://github.com/LF-Decentralized-Trust-labs/>  
<https://github.com/Daniel-Szego/zk-learning-group>

Quizzes and small programming challenges, LFDT merchs at the end



# ***Logistics: Hunting for the SNARK***

February - Introduction, Theory : Definitions and building blocks

March - Theory : Polynomial commitments

**April** - Theory : Interactive oracle proofs

May - Programming : Circom

June - Programming : Circom

July - Programming : Noir

August - Programming : Noir

September : Applications : Off-chain transaction

October : Applications : Proving solvency

November : Applications : Rollup

December : Wrap up, Applications

*Subject to change based on community discussion*



# Agenda



- (zk)SNARK
- *Elements of building a SNARK*
- *Interactive Oracle Proof (IOP)*
- *Building blocks*
- PLONK
- *Computational trace*
- *Constructing polynomials*
- *Summary*
- *Literature and Links*
- *Q&A and discussion*

## ***(zk)SNARK - Succinct Non-interactive ARgument of Knowledge***

**Computation:** arithmetic circuit :  $C(x, w) \rightarrow F$

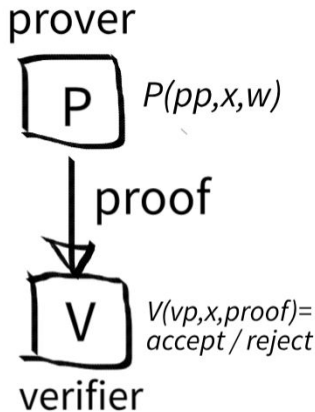
- x public input
- w private input, witness
- high level computation
- arithmetic circuit
- polynomials

**Prover** algorithm:  $P(pp, x, w) \rightarrow \text{proof}$

**Verifier** algorithm:  $V(vp, x, \text{proof}) \rightarrow \text{accept / reject}$

**Properties:**

- Succinct:
- Complete:
- Knowledge sound:
- Zero knowledge



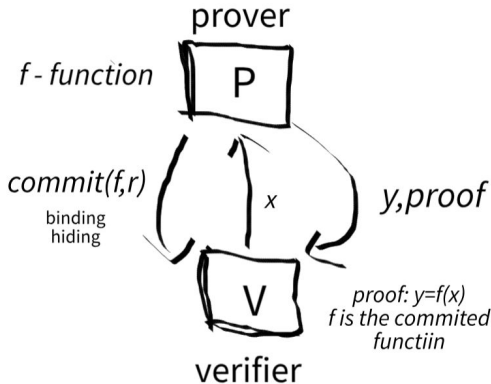
# Elements of building a SNARK

## Elements of a SNARK:

- Polynomial commitment
- Interactive Oracle Proof
- Fiat-Shamir

## Functional commitment:

- Having a family of functions  $F$ .
- Prover: choose an  $f$  function from the  $F$  function family
- Committing  $f$  function to the verifier, com commitment
- Verifier sends an  $x$  point of the function.
- Verifier sends  $y$  and proof that  $y=f(x)$  and that  $f$  is in the function family and that  $f$  is the committed function



# Interactive Oracle Proof (IOP)

Proving that  $C(x, w) = 0$

Committing  $f_0, f_1, \dots, f_N$  functions

Committed functions can be evaluated at chosen points with the help of functional commitments

Repeated steps,  $n$  times:

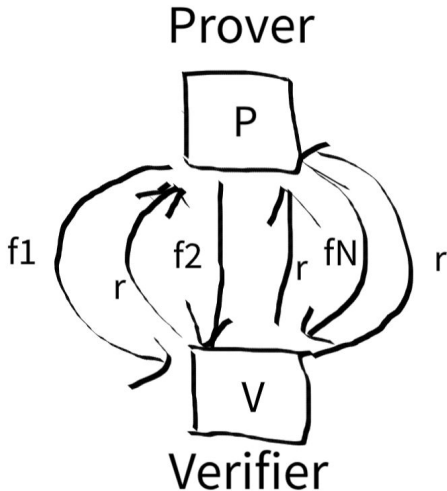
- Prover sends committed functions
- Verifier responds with random values

Final step: verification:

- functions / function commitments
- public input
- opening functions at certain points

Properties:

- complete
- knowledge sound
- zero knowledge



# Building blocks

Testing if a **polynomial is zero**:

- If a truly randomly chosen point evaluates to 0, the whole polynomial is zero with a very high probability.
- $d/p$  :  $d$  degree,  $d$  roots,  $p$  is the domain

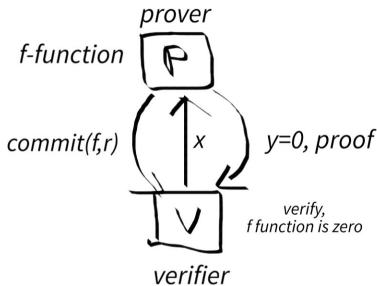
Testing if two **polynomials are equal**:

- if two polynomials  $f$  and  $g$  are equal in a randomly chosen  $r$  point  $f(r)=g(r)$ , then the two polynomials are equal with a very high probability.
- $f(r)-g(r)$  zero test

Testing if a polynomial is **zero on a set**: quotation polynomial

**Sum check**: sum of certain input values of a committed polynomial is a given value.

**Prod check**: product of certain input values of a committed polynomial is a given value.



# PLONK - constructing computational trace

Constructing a polynomial IOP

for a  $C(x, w) = 0$  circuit

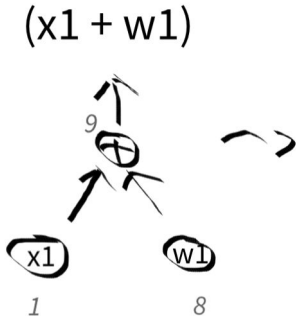
Computational trace

- table
- first row is the input
- for each gate, there is a left input, right input and output

Using computational trace instead of arithmetic circuit

Proving if the computational trace is correct

Proving if the output is zero



inputs:	1	8	
<hr/>			
Gate 0:	1	8	9
Gate 1:	left input	right input	output
Gate 2:	left input	right input	output
...			
			↑ output of the circuit

# PLONK - constructing polynomial

## Computational trace to polynomial:

Constraints to the polynomial

Encoding all inputs:

-  $P(w \text{ to the } -i) = i$  for all  $i$  input

Encoding all wires:

-  $P(w \text{ to the } 3l) = \text{left input of the } l \text{ gate}$

-  $P(w \text{ to the } 3l+1) = \text{right input of the } l \text{ gate}$

-  $P(w \text{ to the } 3l+2) = \text{output of the } l \text{ gate}$

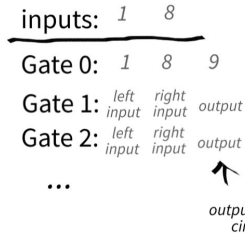
Prover can use the Fast Fourier

Transformation to construct coefficients of the polynomial

$$P(w^{-1})=1, \quad P(w^{-2})=8$$

$$P(w^0)=1, \quad P(w^1)=8, \quad P(w^2)=9$$

$$P(w^3)=\text{left}, \quad P(w^4)=\text{right}, \quad P(w^5)=\text{output}$$

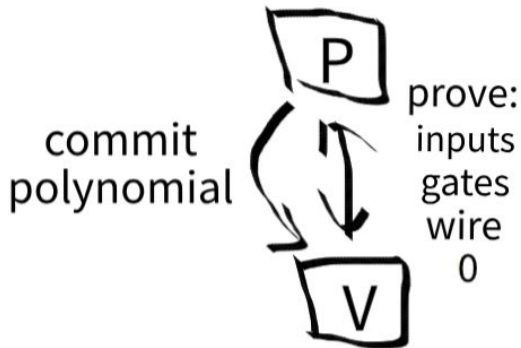


# PLONK - Proving validity

Proving validity of the created polynomial:

- All **inputs** are encoded correctly by the polynomial
- **Gates** are evaluated correctly
- **Wiring** is implemented correctly
- **Output** of the last gate is 0

Proving that the last gate is evaluated to 0, opening the polynomial at the  $P(w$  to the  $3/ + 2) =$  and testing if that is zero



# PLONK - Proving validity of gates and inputs

## Proving correct input:

Additional  $V$  polynomial for encoding all inputs

- $V(w \text{ to the } i) = i$  for all  $i$  input

Prove that the  $V$  input polynomial and the  $P$  committed polynomial agree on all input points: Zero on a set test  $P-V$

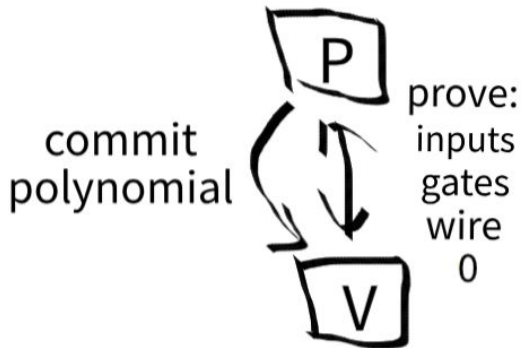
## Proving correctness of gates:

Selector polynomial  $S$ :

- $S(w \text{ to the } 3l) = 1$  if it is an addition gate
- $S(w \text{ to the } 3l) = 0$  if it is a multiplication gate

Combining with the  $P$  polynomial and using the zero set test

Independent from the input values



# PLONK - Proving validity of wiring

## Proving correct wiring:

Equalities for the wiring  
constraint

Proving wiring constraints

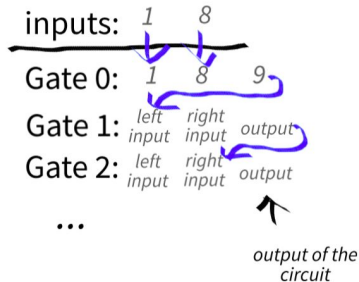
W wiring polynomial:

Rotation of coefficients

Independent from the input  
values

Can be computed at the setup

$$P(w^{-1})=P(w^0) \quad P(w^{-2})=P(w^1) \\ P(w^3)=...$$



# PLONK - Summary

Setup / preprocessing:

- compute wiring polynomial
- compute selector polynomial

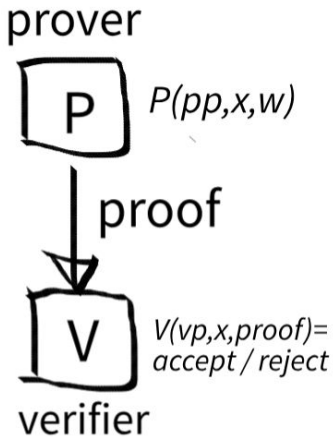
Prover:

- build  $P$  polynomial for the computation

trace and commits

- prove inputs
- prove gates
- prove wiring
- prove output is zero

Using the Fiat Shamir transformation to  
make it non-interactive



# Challenge



**Quiz:**

**Will be posted in the discord channel:**

**<https://discord.com/channels/905194001349627914/1329201532628898036>**

# Links, Resources, Literature



*A guide to Zero Knowledge Proofs (Part 2)*

[https://medium.com/@Luca\\_Franceschini/a-guide-to-zero-knowledge-proofs-part-2-7904dee9758d](https://medium.com/@Luca_Franceschini/a-guide-to-zero-knowledge-proofs-part-2-7904dee9758d)

*Interactive Oracle Proofs*

<https://www.iacr.org/archive/tcc2016b/99850156/99850156.pdf>

*What is PLONK*

[https://medium.com/@Luca\\_Franceschini/what-is-plonk-29c56f326cf6](https://medium.com/@Luca_Franceschini/what-is-plonk-29c56f326cf6)

*Plonk Interactive Oracle Proofs (IOP)*

<https://hackmd.io/@0xsachink/ByuqZfD63>

*Proofs, Arguments, and Zero-Knowledge, Chapter 4*

<https://people.cs.georgetown.edu/jthaler/ProofsArgsAndZK.pdf>



# *Happy Hunting for the SNARK :)*

## **Q & A**

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