Representation of Binders

Using the Bindlib (OCaml) Library



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MOTIVATIONS: BINDERS ARE (WERE?) A PAIN

We need to develop programming languages / proof assistants

This requires many technical but boring elements:

- Source code parsing (notations, unicode)
- Representation of binders (functions, quantifiers)
- Computation of dependencies, management of modules

There are many applications of binders:

- Functions, type abstraction, polymorphic types
- Quantifiers (possibly higher-order), predicates
- Pattern-matching, unification variables, metavariables

STANDARD TECHNIQUES TO DEAL WITH BINDERS

```
module DB = struct
  (* With De Bruijn indices. *)
 type term =
    Var of string (* Free variable.
     Idx of int (* Bound variable index. *)
     Abs of term (* Abstraction (function). *)
     App of term * term (* Function application. *)
end
module HOAS = struct
  (* With higher-order abstract syntax. *)
 type term =
     Var of string (* Free variable.
     Abs of (term -> term) (* Abstraction (function). *)
     App of term * term (* Function application.
end
```

BINDLIB, A HISTORY

Bindlib was designed by Christophe Raffalli in the nineties

I contributed several improvements:

- New, well-documented implementation (almost) from scratch
- Bindlib without the old (camlp4) syntax extension
- Lighter free variable management, "unbind" function, ...

Implemented systems relying on Bindlib:

- Lambdapi (new version of the Dedukti logical framework)
- PML proof system, SubML language (with subtyping)
- Pure type systems (PTS) and combinatory reduction systems (CRS)



ABSTRACT SYNTAX REPRESENTATION

SUBSTITUTION AND DESTRUCTIVE TRAVERSAL

```
val subst : ('a,'b) Bindlib.binder -> 'a -> 'b
val unbind : ('a,'b) Bindlib.binder -> 'a var * 'b
let rec eval : term * stack -> term * stack = function
  , Push(u,pi))
  (LAbs(f) , Push(t,pi)) -> eval (Bindlib.subst f t , pi
  | (MAbs(f) , pi ) -> eval (Bindlib.subst f pi, pi
  | (Name(pi,t), ) -> eval (t
                                                   , pi
  l whnf
                      -> whnf
let rec to string : term -> string = function
   TVar(x) -> Bindlib.name of x
  | LAbs(f) \rightarrow let (x,t) = Bindlib.unbind f in
               Printf.sprintf "\\%s.%s" (Bindlib.name of x) (to string t)
  | Appl(t,u) -> Printf.sprintf "(%s) %s" (to string t) (to string u)
           -> "<...>"
```

THINKING INSIDE THE BOX

```
(* There is no generic function like the following. *)
val bind var : 'a Bindlib.var -> 'b -> ('a,'b) Bindlib.binder
(* However, Bindlib provides the following function. *)
val bind var : 'a Bindlib.var -> 'b Bindlib.box
               -> ('a.'b) Bindlib.binder Bindlib.box
(* A value of the type ['a Bindlib.box] represents:
     - an element of type ['a] under construction,
     - its free variables are available for binding. *)
(* The ['a Bindlib.box] type is an applicative functor. *)
val box : 'a -> 'a Bindlib.box
val apply box : ('a -> 'b) Bindlib.box -> 'a Bindlib.box -> 'b Bindlib.box
val box var : 'a Bindlib.var -> 'a Bindlib.box
```

SMART CONSTRUCTORS AND LIFTING

```
let TVar : term Bindlib.var -> term Bindlib.box =
 fun x -> Bindlib.box var x
let LAbs : (term, term) Bindlib.binder Bindlib.box -> term Bindlib.box =
 fun b -> Bindlib.box apply (fun f -> LAbs(f)) b
let Appl : term Bindlib.box -> term Bindlib.box -> term Bindlib.box =
 fun t u -> Bindlib.box apply2 (fun t u -> Appl(t,u)) t u
let rec lift term : term -> term Bindlib.box = function
   TVar(x) -> TVar x
  LAbs(b) -> LAbs (Bindlib.box binder lift term b)
  | Appl(t,u) -> Appl (lift term t) (lift term u)
             -> failwith "Not implemented..."
```

EXAMPLES OF TERMS

```
(* Fresh (free variables). *)
let x : term Bindlib.var = Bindlib.new var (fun x -> TVar(x)) "x"
let y : term Bindlib.var = Bindlib.new var (fun x -> TVar(x)) "y"
(* Usual terms. *)
let id : term Bindlib.box =
 LAbs (Bindlib.bind var x ( TVar x))
let fst : term Bindlib.box =
 LAbs (Bindlib.bind var x ( LAbs (Bindlib.bind var y ( TVar x))))
let delta : term Bindlib.box =
 LAbs (Bindlib.bind var x ( Appl ( TVar x) ( TVar x)))
(* Unboxed term (fully constructed). *)
let omega : term = Bindlib.unbox ( Appl delta delta)
```

WORKING UNDER BINDERS

```
let rec snf : term -> term = function
  | Appl(t,u) ->
      begin
        let v = snf u in
        match snf t with
        | LAbs(b) -> snf (Bindlib.subst b v)
         h \rightarrow Appl(h,v)
      end
   LAbs(b) ->
      begin
        let (x,t) = Bindlib.unbind b in
        let v = snf t in
        Bindlib.unbox ( LAbs (Bindlib.bind var x (lift term v)))
      end
   TVar(x) \rightarrow TVar(x)
              -> failwith "not a lambda-term"
```

INTERNAL REPRESENTATION: BINDERS

```
type ('a,'b) binder =
  { b name : string (* Name of the bound variable.
  ; b bind : bool (* Indicates whether the variable occurs. *)
  ; b rank : int (* Number of remaining free variables.
                                                                  *)
  ; b mkfree : 'a var -> 'a (* Injection of variables into domain.
                                                                  *)
                                                                  *) }
  ; b value : 'a -> 'b (* Substitution function.
let subst : ('a,'b) binder -> 'a -> 'b =
  fun b v -> b.b value v
let unbind : ('a,'b) binder -> 'a var * 'b = fun b ->
   let x = new var b.b mkfree (binder name b) in
    (x, subst b (b.b mkfree x))
```

INTERNAL REPRESENTATION: VARIABLES AND BOX

```
type 'a closure = varpos -> Env.t -> 'a
type 'a box =
  | Box of 'a
  (* Element of type ['a] with no free variable. *)
  | Env of any var list * int * 'a closure
 (* Element of type ['a] with free variables stored in an environment. *)
and 'a var =
 { var_key : int (* Unique identifier.
                                                                 *)
  ; var prefix : string (* Name as a free variable (prefix). *)
 ; var suffix : int (* Integer suffix.
                                                                 *)
  ; var mkfree : 'a var -> 'a (* Function to build a term. *)
  ; mutable var box : 'a box (* Bindbox containing the variable. *) }
let box var : 'a var -> 'a box = fun x -> x.var box
```

INTERNAL REPRESENTATION: VARIABLE CREATION

```
(* type any var = Any : 'a var -> any [@@ unboxed] FIXME *)
type any var = Obj.t var
let new var closure key = fun vp -> Env.get (IMap.find key vp).index
let new var : ('a var -> 'a) -> string -> 'a var =
  fun var mkfree name ->
    let var key = fresh key () in
    let (var prefix, var suffix) = split name name in
 (* let rec x =
      { var key; var prefix; var suffix; var mkfree
      ; var box = Env([Any x], 0, new var closure var key) }
   in x *)
    let var box = Env([], 0, fun -> assert false) in
    let x = {var key; var prefix; var suffix; var mkfree; var box} in
   x.var box <- Env([0bj.magic x], 0, new var closure var key); x
```

THE (OBJ.) MAGIC OF BINDLIB

```
module Env : sig
  type t
  val create : int -> t
  val set : t -> int -> 'a -> unit
  val get : int -> t -> 'a
  (* ... *)
end = struct
  (* ... *)
  (* Safe as soon as we write/read at a fixed type for each index. *)
  let set env i e = Array.set env.tab i (Obj.repr e)
  let get i env = Obj.obj (Array.get env.tab i)
  (* ... *)
end
```

RELATED AND FUTURE WORK

Formal proof of correctness:

- Coq implementation of Bindlib (Bruno Barras)
- With axiomatized environment operations
- PML implementation of Bindlib? (Bootstrap)

Complexity analysis (Bruno Barras)

We need feedback from new users!

Thanks!

https://github.com/rlepigre/ocaml-bindlib