University of Science, VNU-HCM Faculty of Information Technology

Data Structure and Algorithm

M-way Tree BTree

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Outline

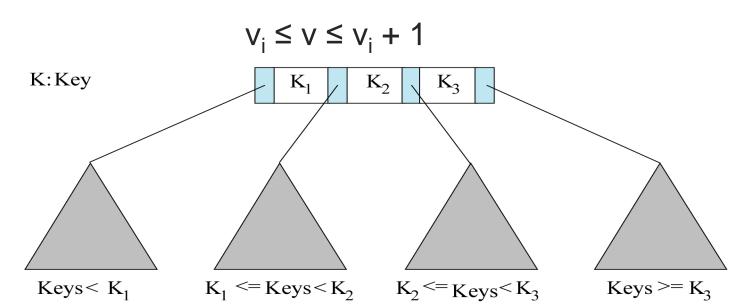
- M-way Tree
- B-Tree

About binary tree

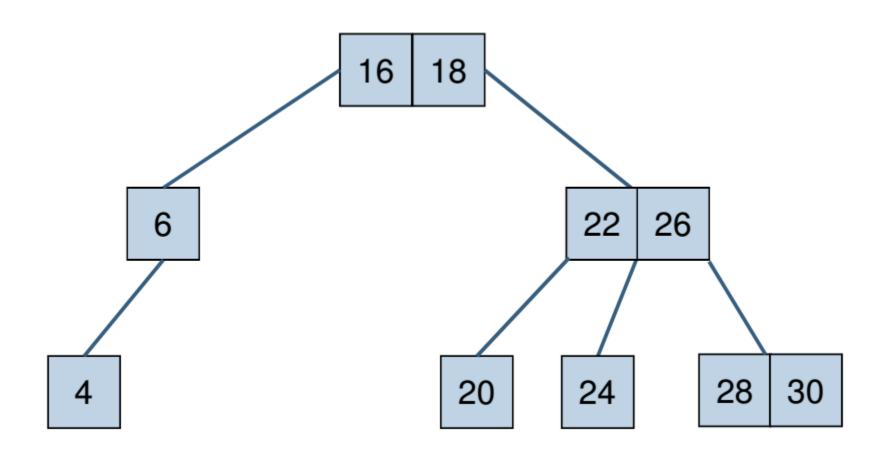
- Binary tree: each node has 1 data value and 2 branches
- Some disadvantages:
 - Each stored value need at least two more memory (left pointer, right pointer)
 - Tree can be very high if number of values is large.
 - Implement a balanced binary search tree is complex.

Multi-way Tree

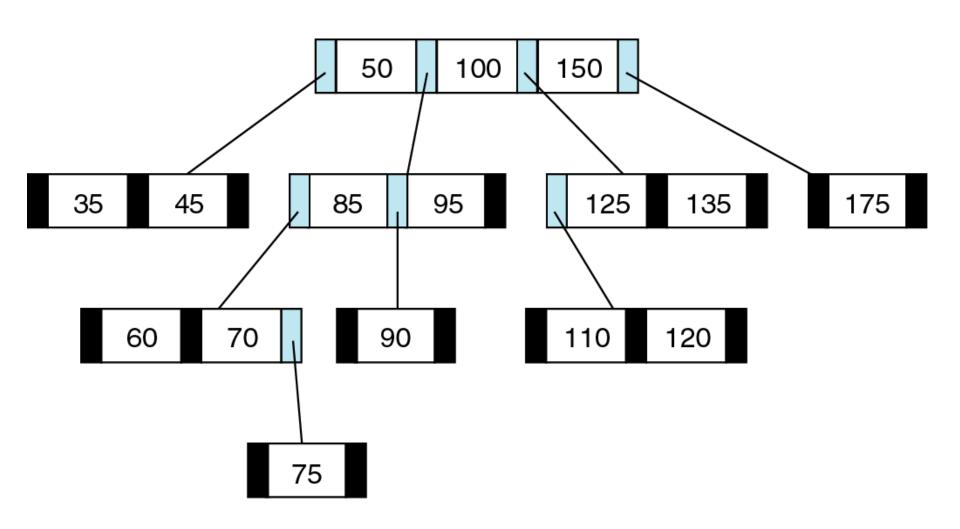
- The multi-way tree (m-way tree) is a tree:
 - Each node contains 1 to m-1 keys with distinct values
 - The keys in each node are ordered (ascending).
 - A node with k keys will have k + 1 subtree, the subtree can be empty.
 - The ith subtree $(0 \le i \le k)$ of the node contains the keys such that:



3-way tree example



4-way tree example

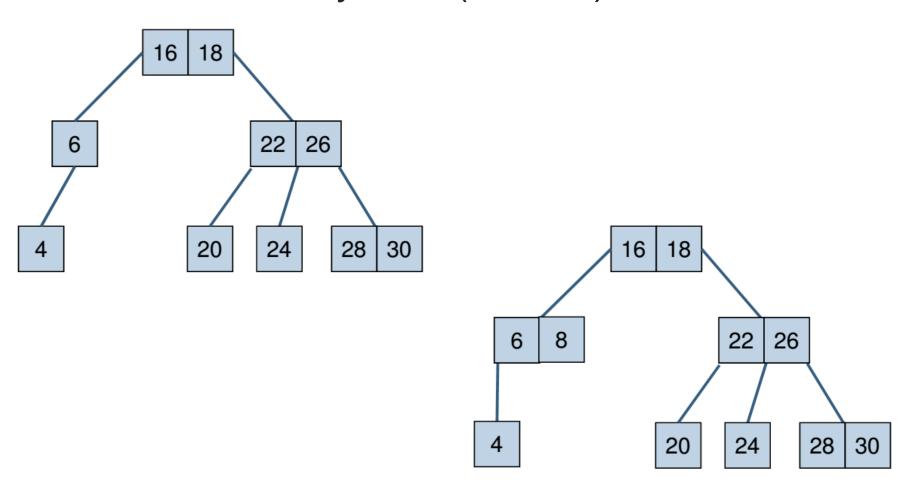


Insert new element

- Insert a key v into the tree:
 - Travese the tree until it finds an empty node
 - Case 1: If the parent node still has slot: add the key v to the parent node at this slot.
 - Case 2: If the parent node is full: create a new node and add the key v to it.

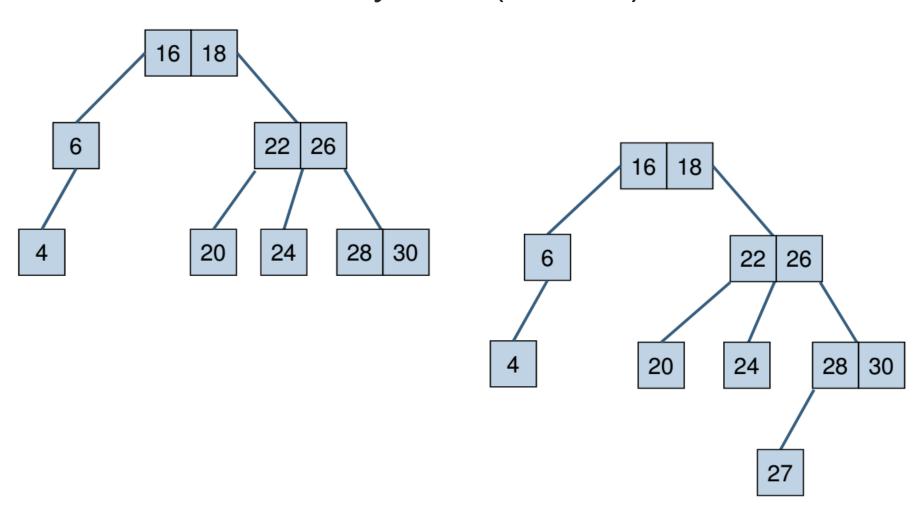
Insert new element

Insert 8 to 3-way tree (case 1)



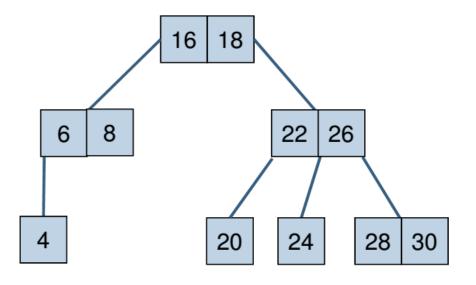
Insert new element

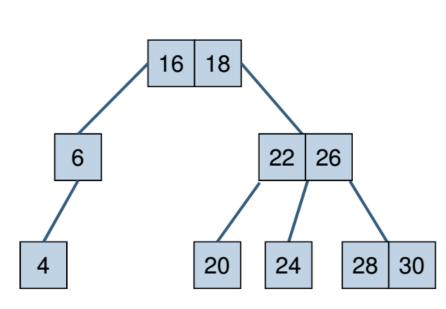
Insert 27 to 3-way tree (case 2):



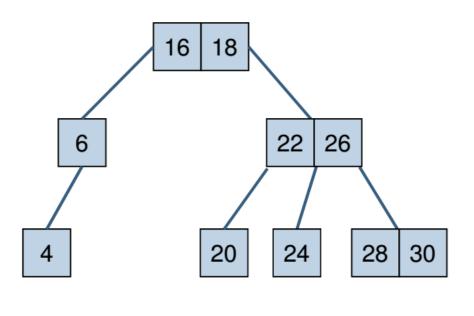
- Delete a key v from the tree:
 - Case 1: If v has no child (between 2 empty subtree) then just delete v.
 - Case 2: If v has any child, replace v with largest element in the left subtree of v or smallest element in the right subtree of v.

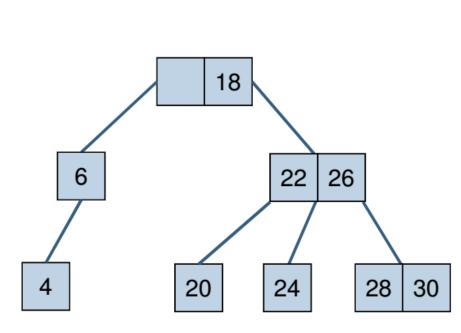
• Delete 8 (case 1)



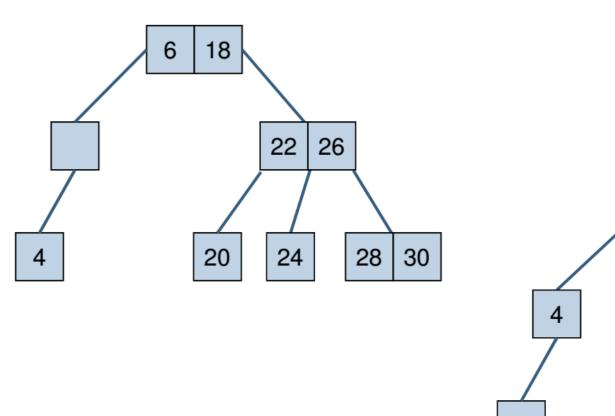


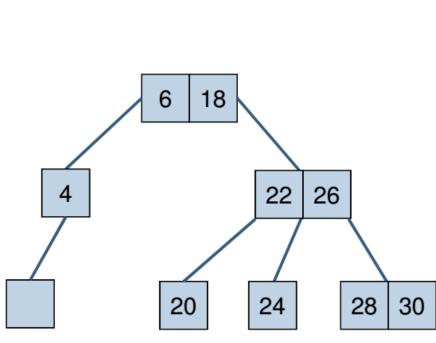
• Delete 16 (case 2):



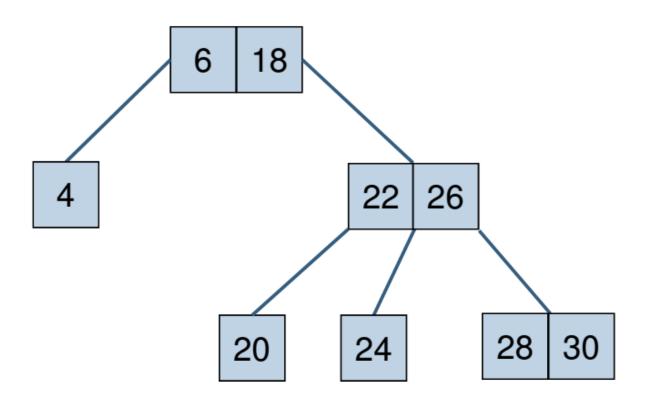


• Delete 16 (case 2):





• Delete 16 (case 2):

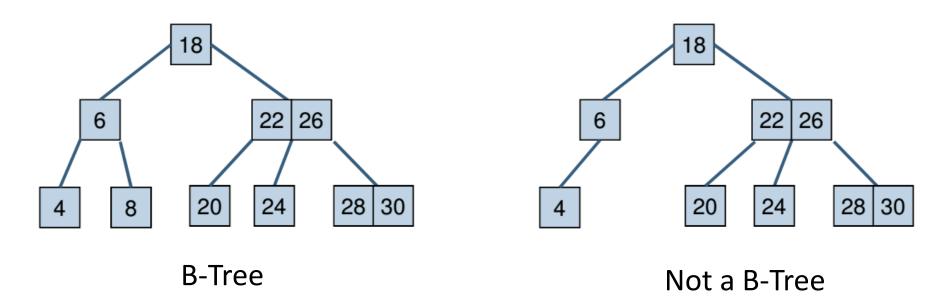


Outline

- M-way Tree
- B-Tree

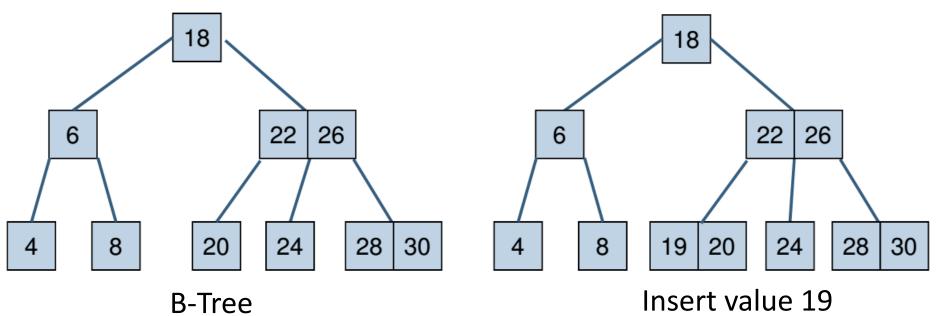
Balanced m-way Tree: B-Tree

- B-Tree is a m-way tree which satisfies:
 - Root node has at least 1 key
 - Branch nodes have at least [(m-1) / 2] +1 subtree
 - i.e. have at least [(m-1) / 2] keys
 - All empty nodes belong to the same level

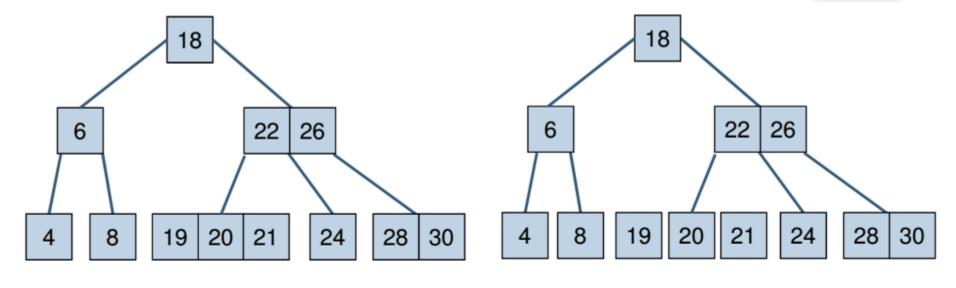


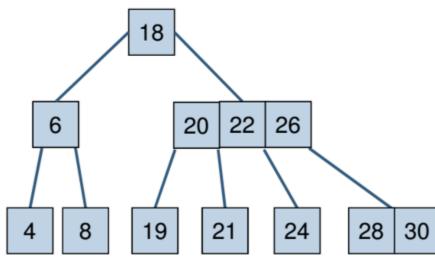
Insert new element in B-Tree

- Insert a key v to the B-tree
 - Add v to a leaf node
 - If the leaf node is full: splits the leaf node in half and moves the middle element onto the parent node.



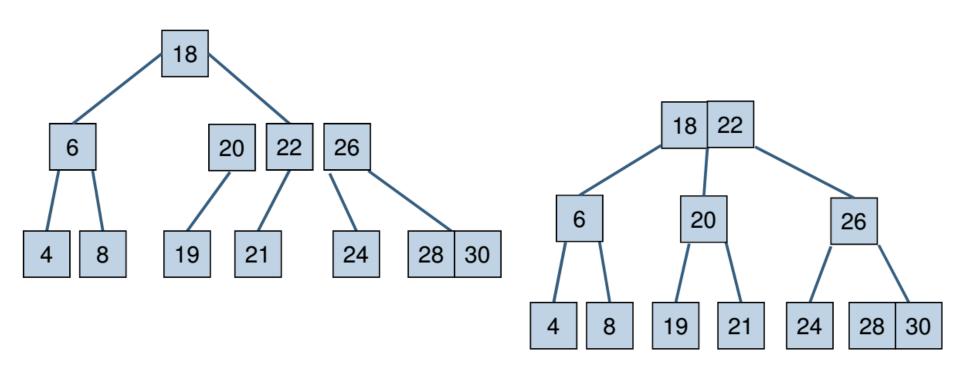
Example for 3-way tree





Adding key 21 at leaf node cause it full, so splitting the leaf node in half and moving the middle element onto the parent node.

After middle node move to parent node:

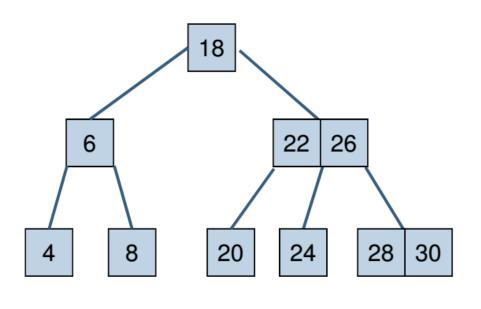


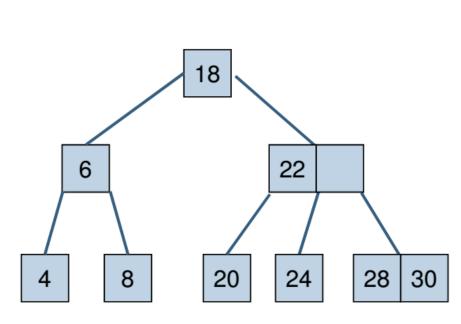
Parent node is full, so repeat same action

Delete an element from B-Tree

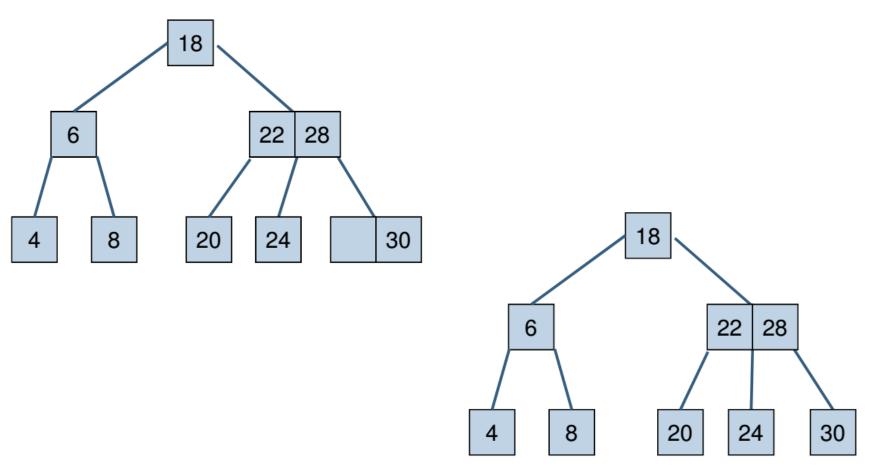
- Delete a key v from the tree
 - Do the same as a m-way tree.
 - If a node has less than [(m-1) / 2] keys
 - Borrow 1 key from the adjacent sibling node if the sibling node have enough key, or
 - Merge with an adjacent sibling node if the sibling node does not have enough key and a corresponding key from the parent node.

• Delete value 26:

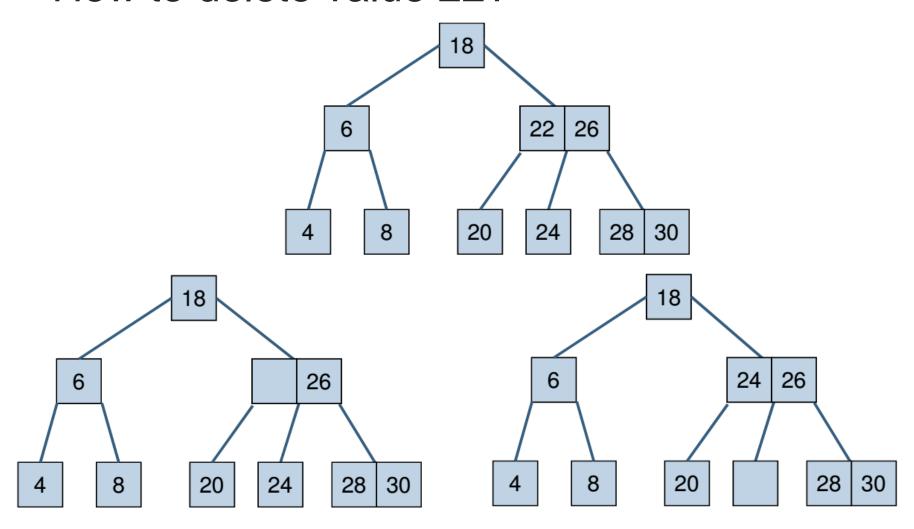




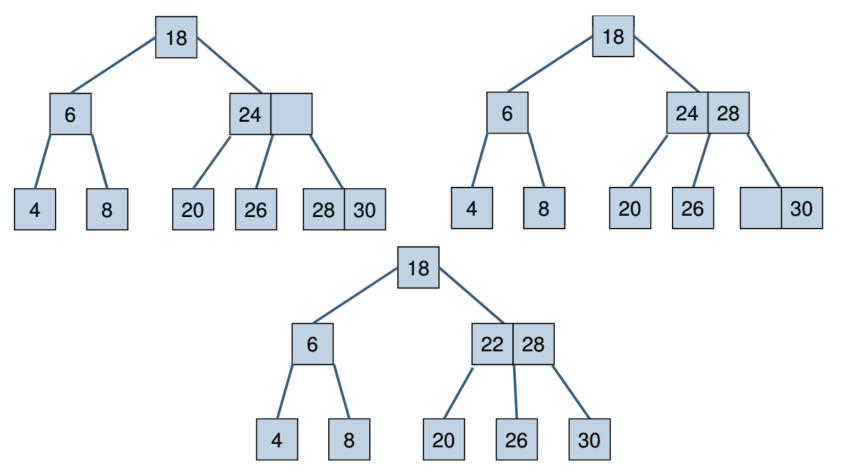
 To delete value 26, replace with the smallest value on the left subtree



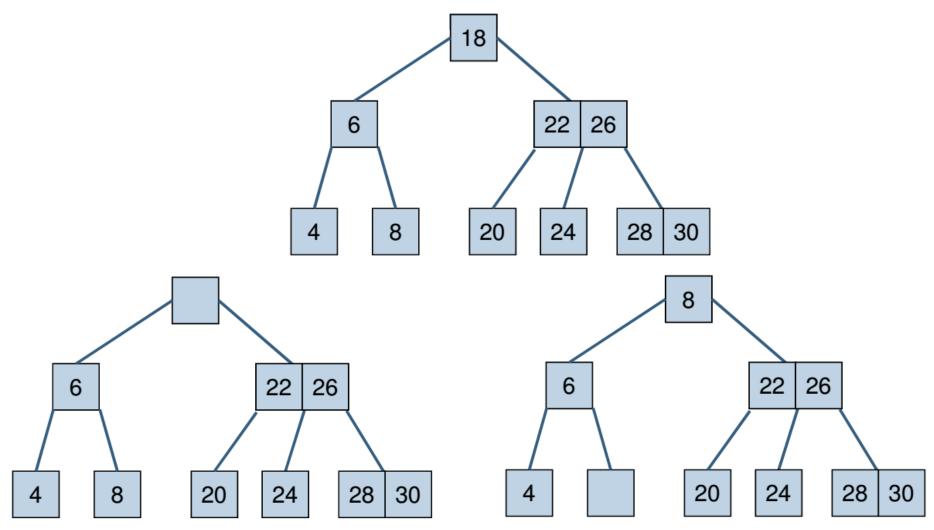
How to delete value 22?

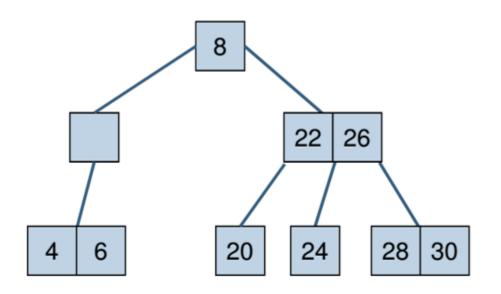


 After delete 22, a node is not enough key, so adjust it (case 1)



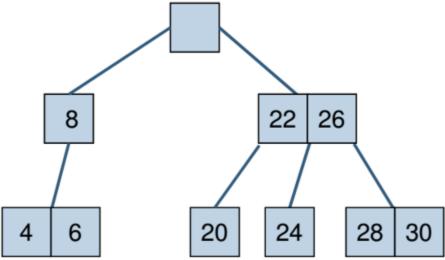
• How to delete 18?



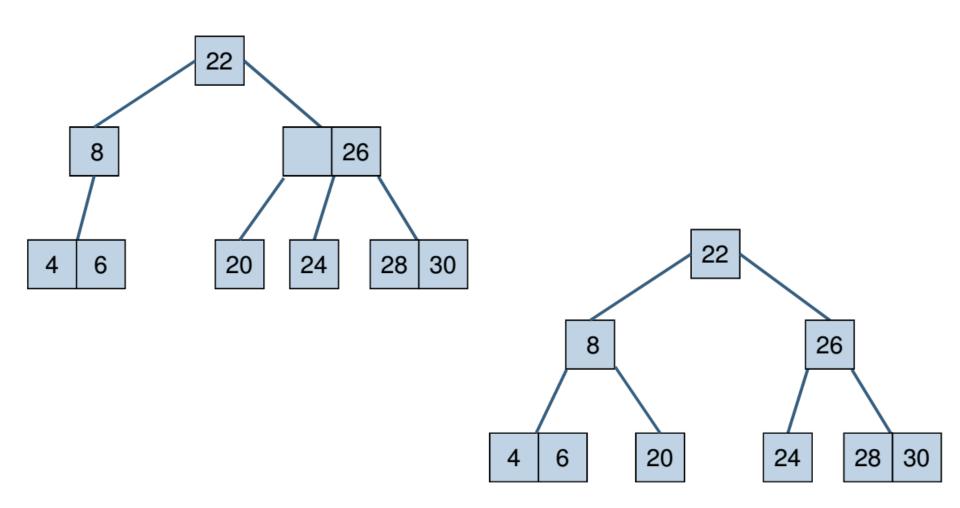


Merge with sibling sandwiching their separator taken off from their parent

Rebalance the parent because the parent has not enough key.



• Balance nodes:



Exercises

 Let's create a 5-way B-Tree with the following data in turn:

```
3, 7, 9, 23, 45, 1, 5, 14, 25, 24, 13, 11, 8, 19, 4, 31, 35, 56, 2, 6, 12.
```

Delete the following keys:

