

First Intuition of Category Theory

wilberchao 2019/2/26

Agenda

- Definition of Category Theory
- Functor
- Natural Transformation

Definition of Category Theory

- Primitive
 - Object
 - Morphism
- Properties
 - Composition
 - Identity
 - Associativity

Object

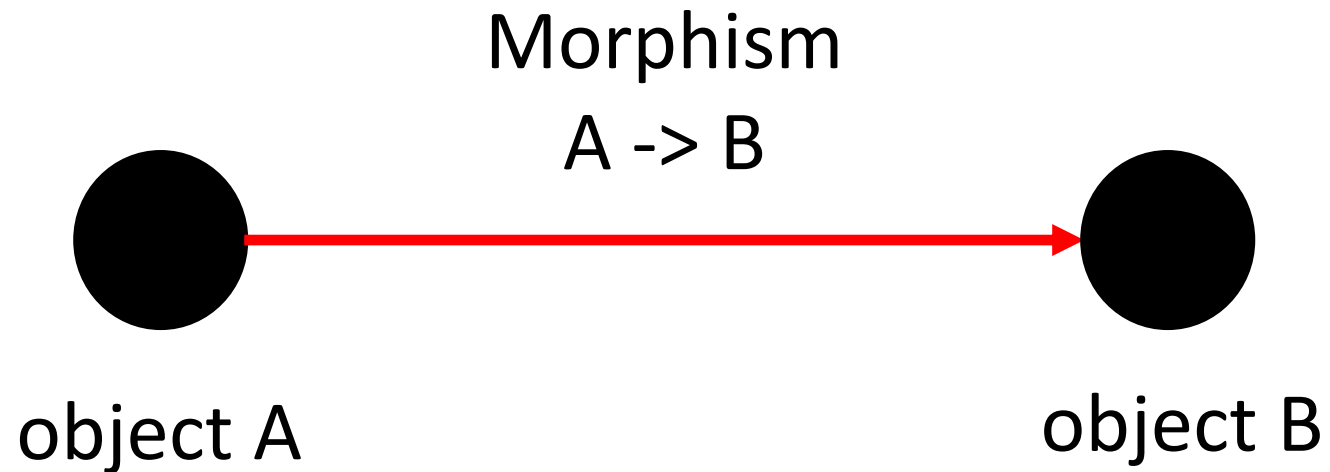
- Not Object in Object-Oriented programming
- Not Class in Object-Oriented programming
- Just a point
- Primitive
- No properties

Morphism

- Primitive
- No properties
- Just an arrow
 - Identifying start and end

Object - Morphism

- The object is used to identify the start and end of the morphism
- The morphism shows the relationship between objects



Object – Morphism

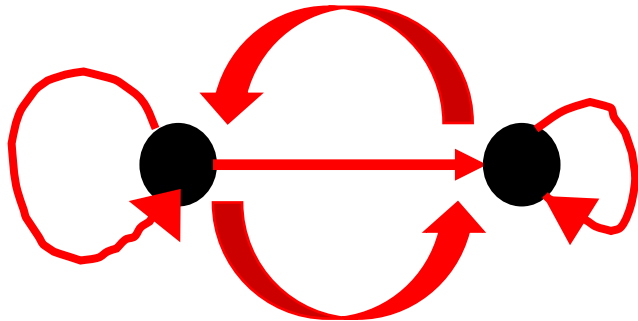
only one morphism



no morphism



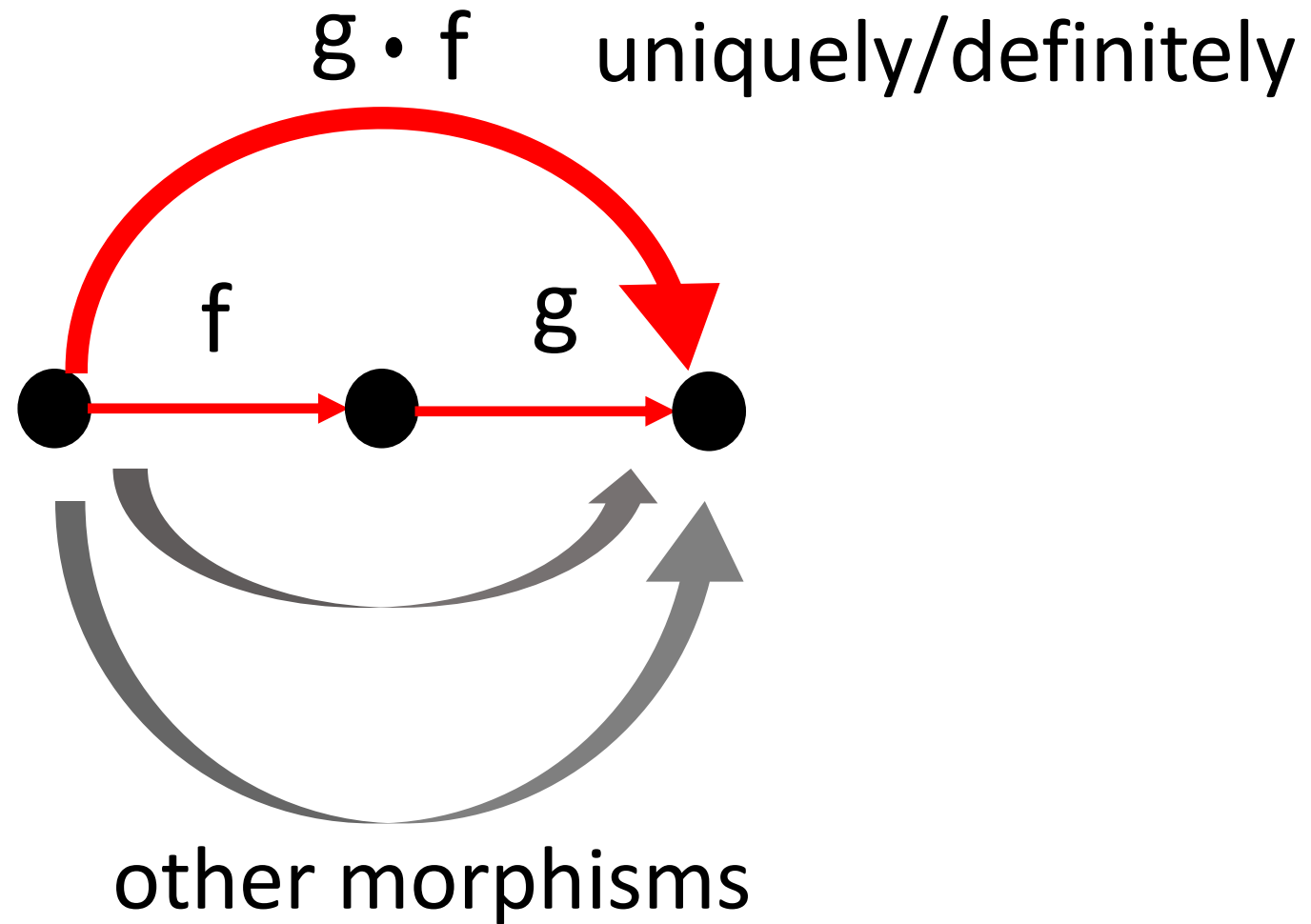
many morphisms



infinite morphisms



Composition

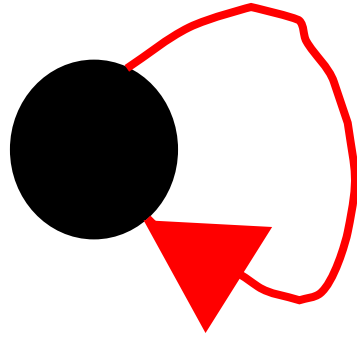


Composition

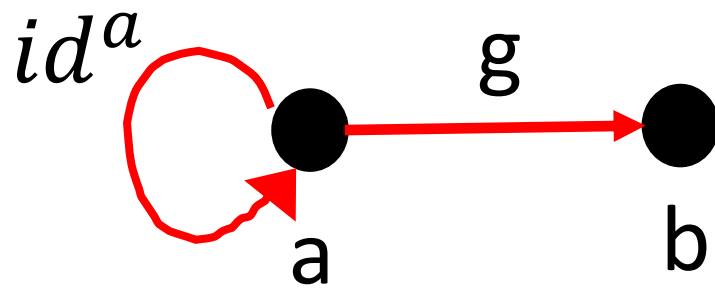
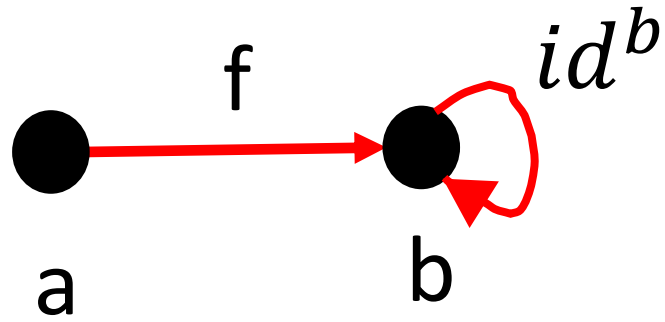
- Objects, morphisms don't have structures
- Composition has structures
 - Encoding everything into compositions

Identity

Every object has one identity morphism at least



Identity



$$id^b \circ f = f$$

left identity



$$g \circ id^a = g$$

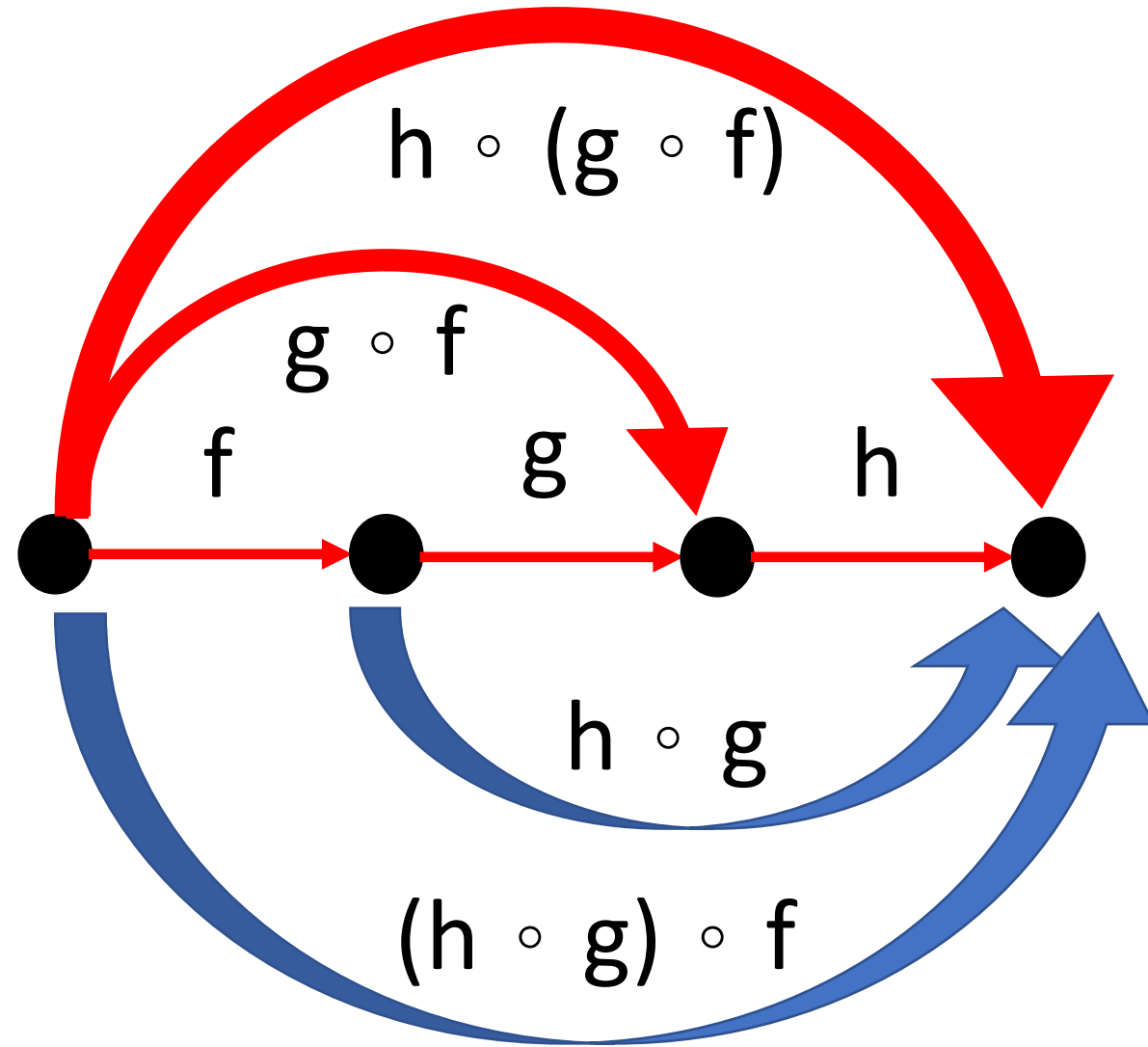
right identity

Associativity

$$h \circ (g \circ f)$$



$$(h \circ g) \circ f$$



Set Category

- Object: Set
- Morphism: functions between sets

Scala Category

- Object: Type
- Morphism: Function

```
scala> val f:Int => String = a => a.toString
f: Int => String = $$Lambda$1079/671384775@5a6f6cac

scala> val g: String => Long = s => 0L
g: String => Long = $$Lambda$1080/202968316@77ce8bc5

scala> f andThen g
res2: Int => Long = scala.Function1$$Lambda$1078/1861236708@335cdd2

scala> g compose f
res3: Int => Long = scala.Function1$$Lambda$1218/860285190@443ac5b8
```

Why Referential transparency

- Making functions produce a value rather than effects
 - `val f: JDBC => IO[Read Mysql]`
 - `val g: JDBC => Future[Read Mysql]`
- We can compose functions which produce value
- We can't compose functions which produce effect
 - Effects disobey the associativity of composition

Functor ?

Functor is a type class that abstracts over type constructors that can be `map`'ed over. Examples of such type constructors are `List`, `Option`, and `Future`.

```
trait Functor[F[_]] {  
  def map[A, B](fa: F[A])(f: A => B): F[B]  
}  
  
// Example implementation for Option  
implicit val functorForOption: Functor[Option] = new Functor[Option] {  
  def map[A, B](fa: Option[A])(f: A => B): Option[B] = fa match {  
    case None    => None  
    case Some(a) => Some(f(a))  
  }  
}
```

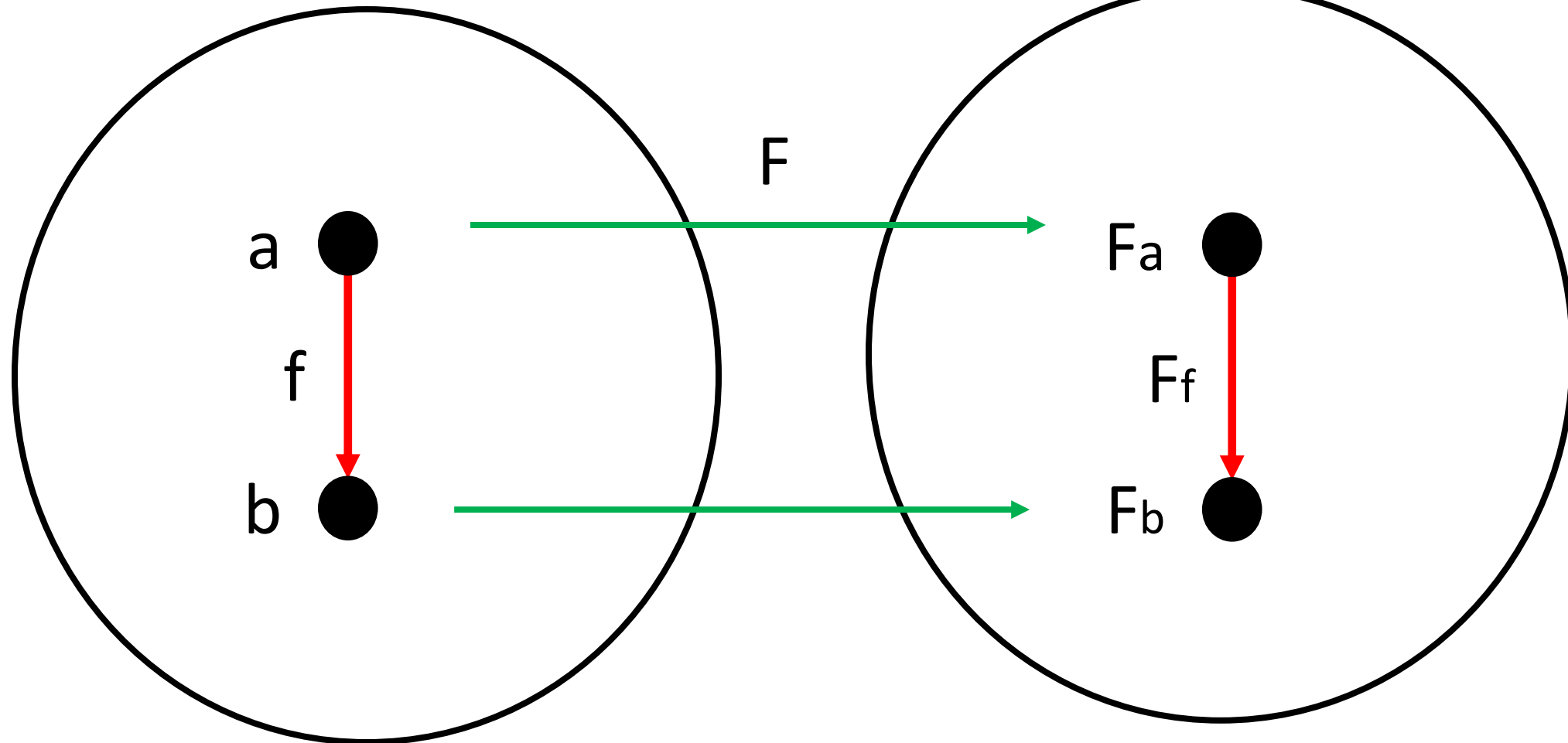

Functor

- Mappings between categories
 - Object
 - Morphism
- Preserving Structure

Functor Mapping

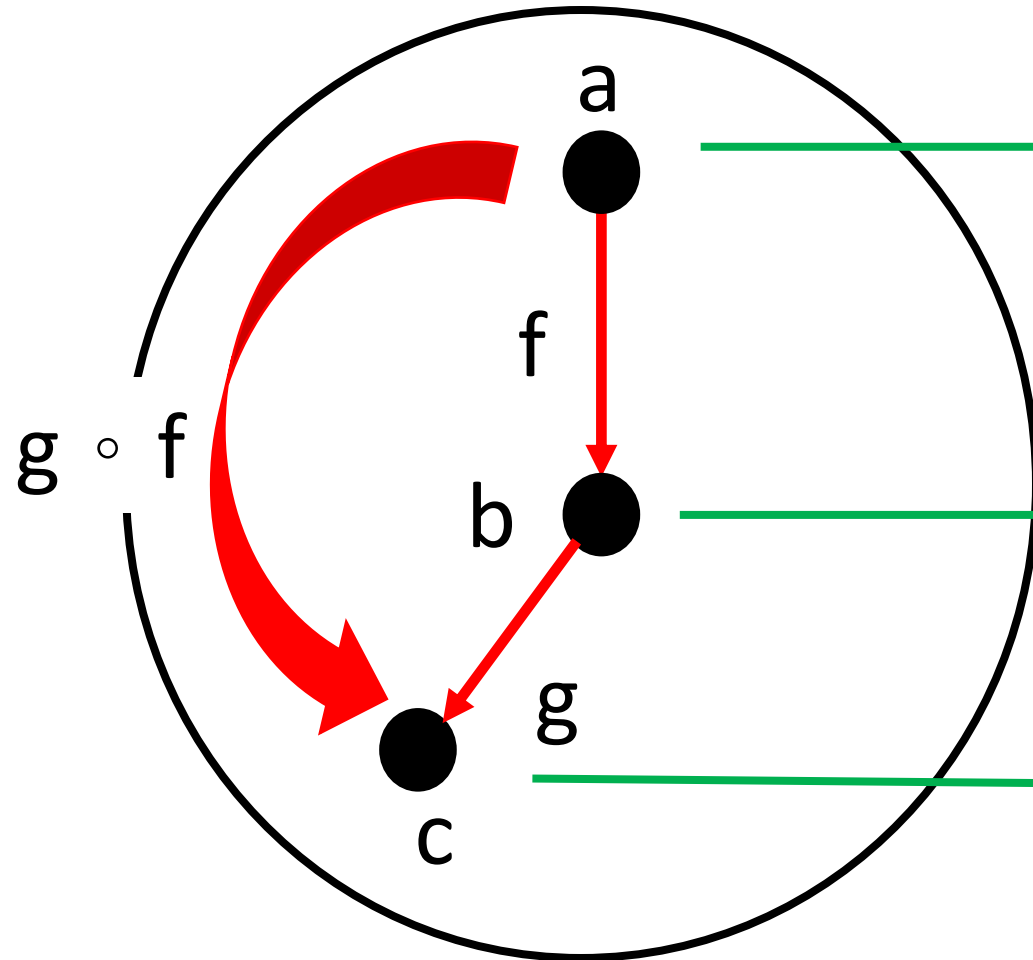
Category C

Category D



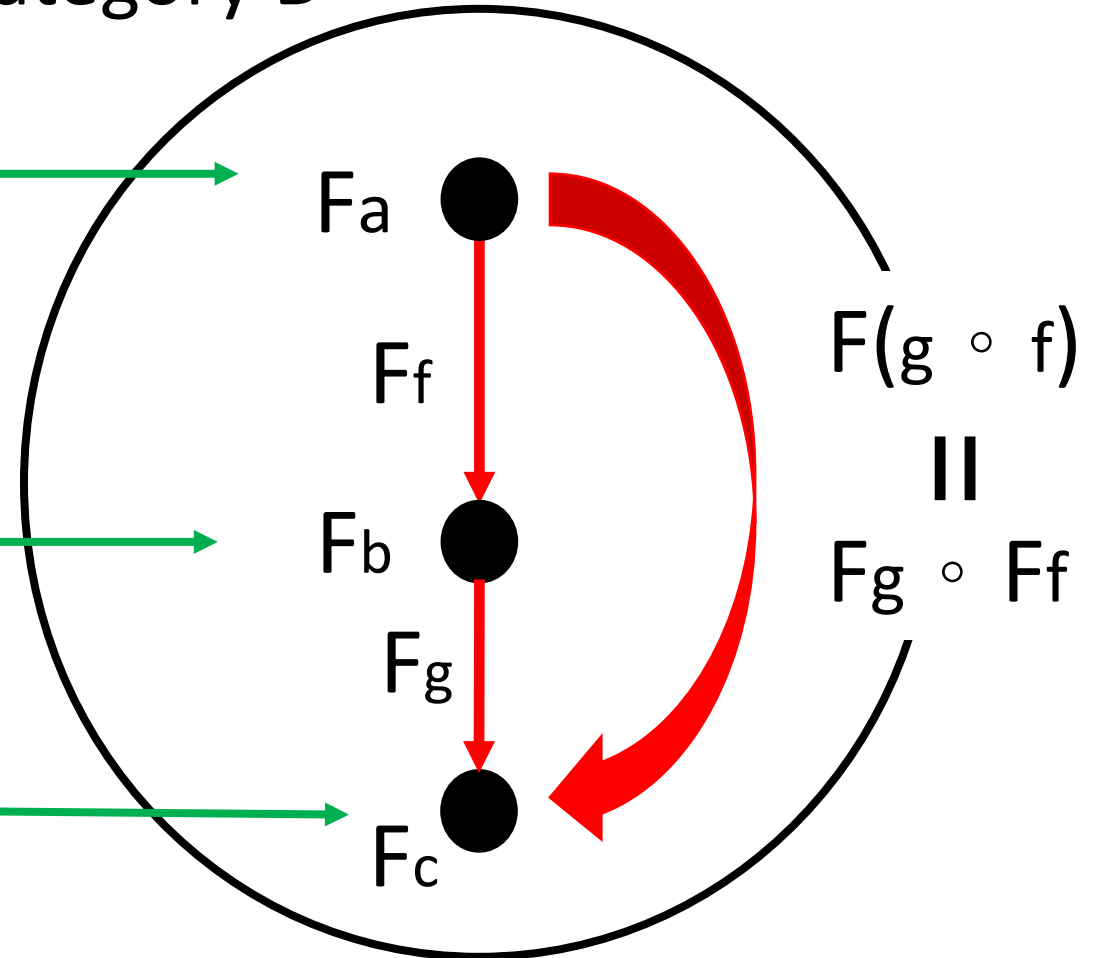
Functor Preserving Structure

Category C

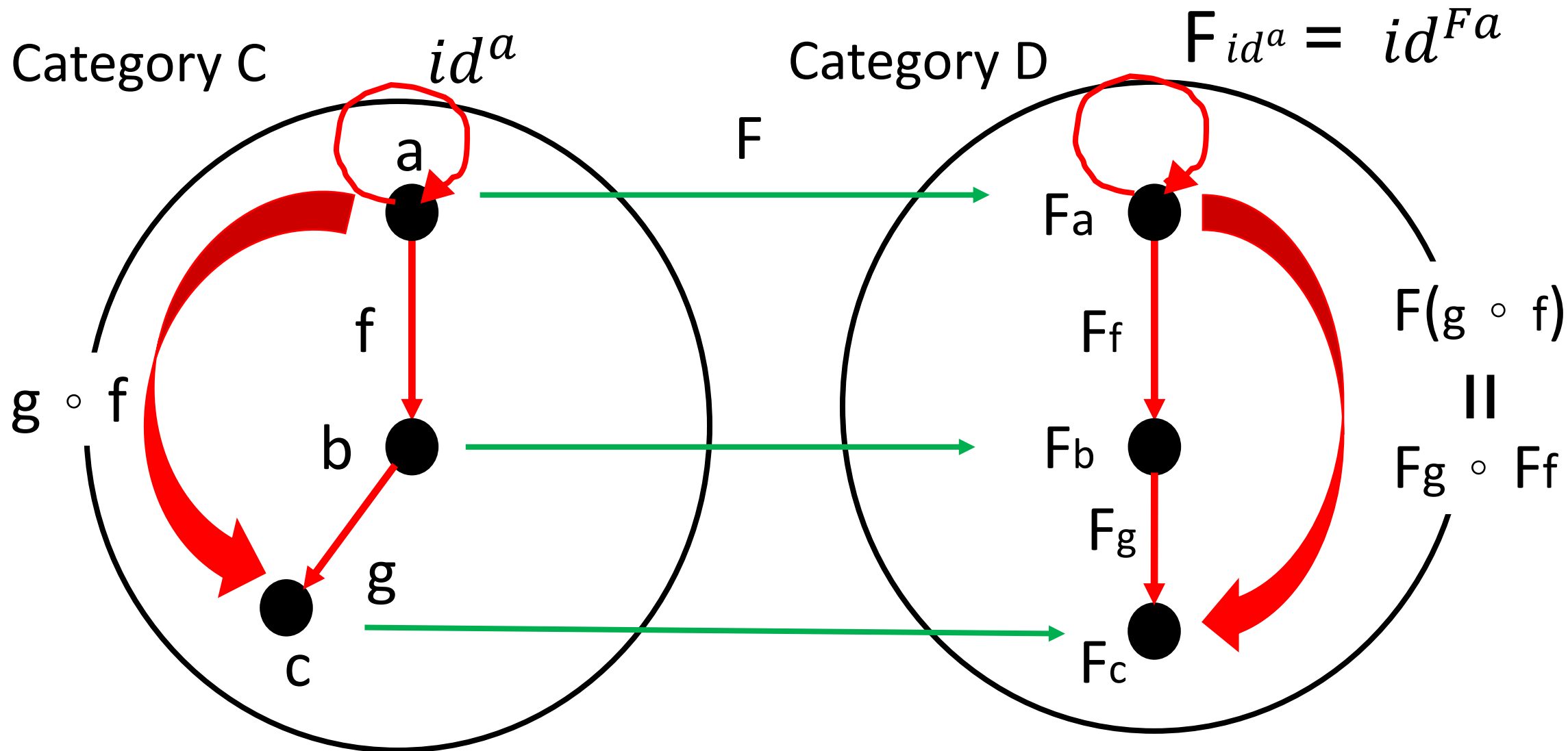


Category D

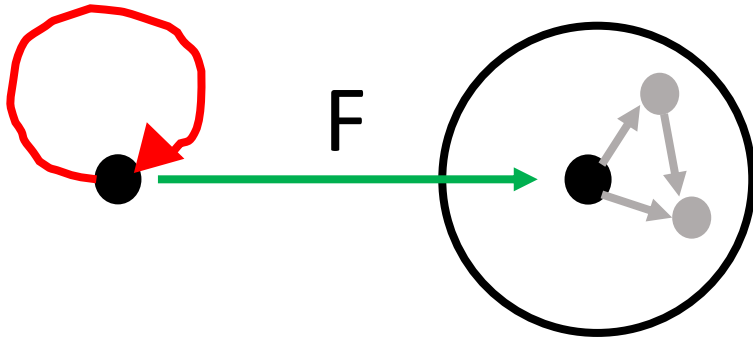
F



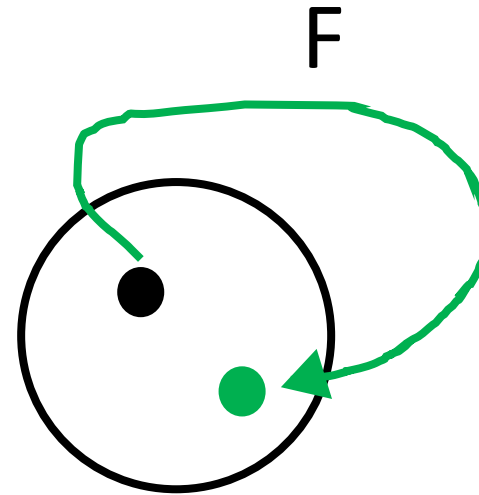
Functor Identity



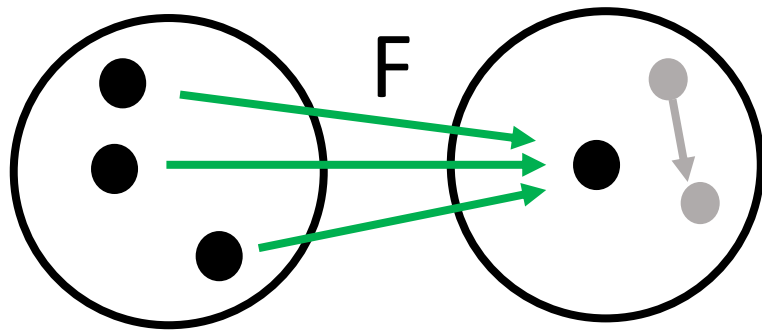
Functors



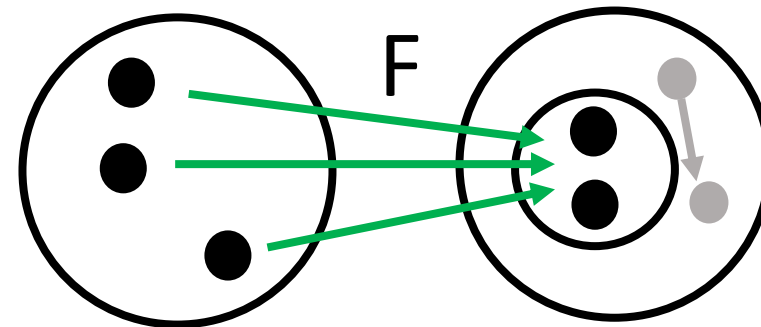
Functor maps a single object category to another point in another category



Functor maps a category to itself

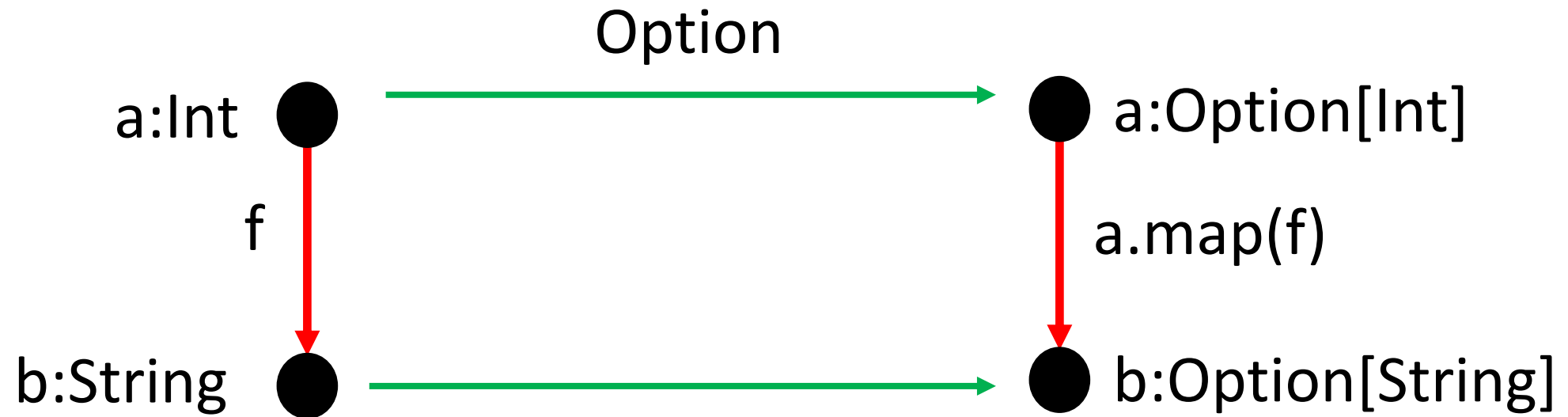


Functor maps whole category to an object in another category



Functor maps whole category to an subset of another category

Functor in Scala



Functor in Cats

```
trait Functor[F[_]] {  
  def map[A, B](fa: F[A])(f: A => B): F[B]  
}  
  
// Example implementation for Option  
implicit val functorForOption: Functor[Option] = new Functor[Option] {  
  def map[A, B](fa: Option[A])(f: A => B): Option[B] = fa match {  
    case None    => None  
    case Some(a) => Some(f(a))  
  }  
}
```

Functor law

A `Functor` instance must obey two laws:

- Composition: Mapping with `f` and then again with `g` is the same as mapping once with the composition of `f` and `g`
 - `fa.map(f).map(g) = fa.map(f.andThen(g))`
- Identity: Mapping with the identity function is a no-op
 - `fa.map(x => x) = fa`

Functor Law

- Functor of `Option[_]`
- Identity
 - $\text{map Option}(a) \text{ Id}_a = \text{Id}_a \text{ Option}(a)$
- Composition
 - $\text{map}(\text{Option}(a))(g \circ f) = (\text{map } g \circ \text{map } f)(\text{Option}(a))$

Functor Container

- Option[A], either None or Some(A)
 - You can just put a function into `map`
- List[A], either 0 or many A in a list
 - You can just put a function into `map`
- Future[A], either happened or pending
 - You can just put a function into `map`

Natural Transformation ?

Natural Transformation

I think we now have enough ammunition on our hands to tackle naturality. Let's skip to the middle of the book, section 7.4.

A natural transformation is a morphism of functors. That is right: for fix categories \mathbf{C} and \mathbf{D} , we can regard the functors $\mathbf{C} \Rightarrow \mathbf{D}$ as the object of a new category, and the arrows between these objects are what we are going to call natural transformations.

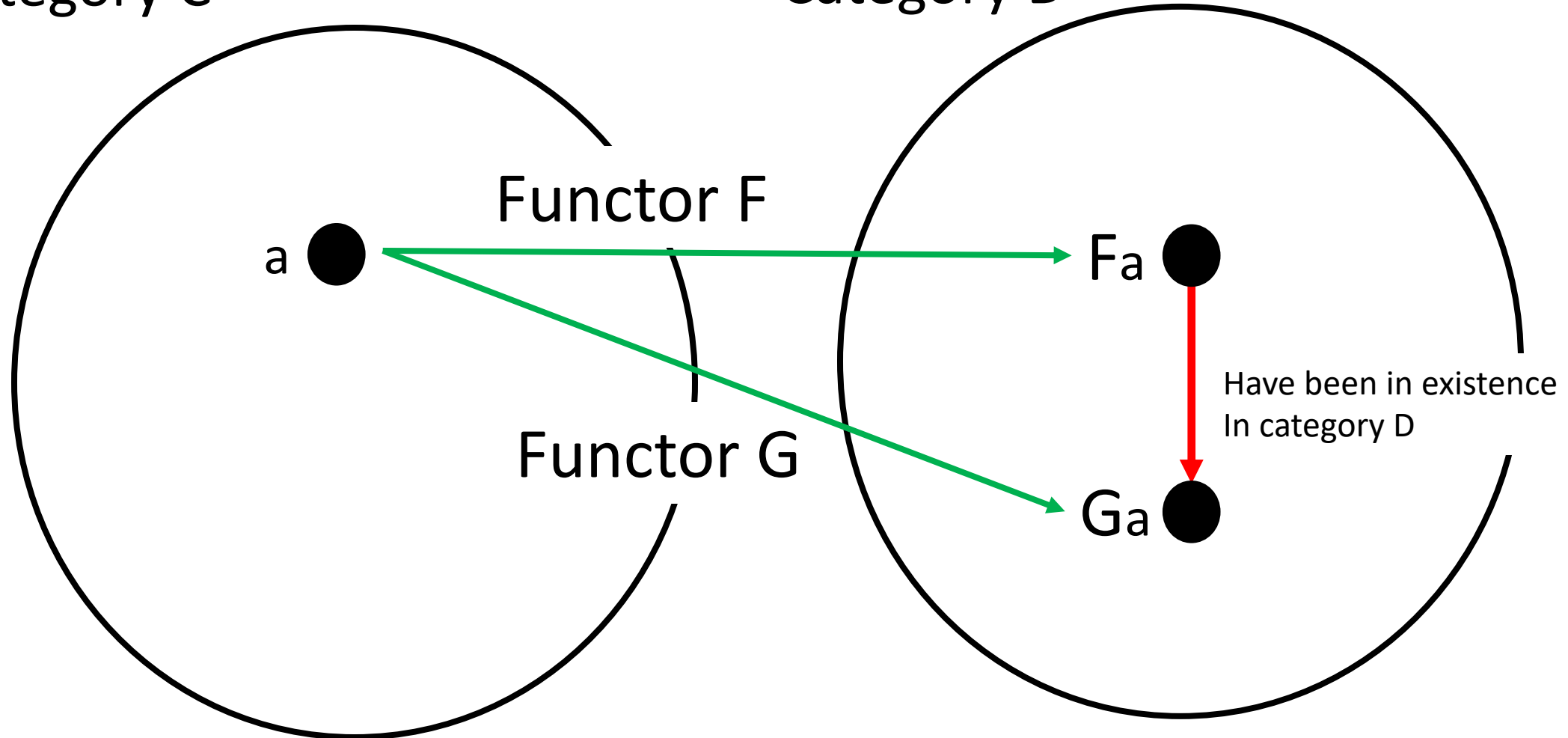
Natural Transformation

- Mapping between Functors
- Preserving structure

Natural Transformation

Category C

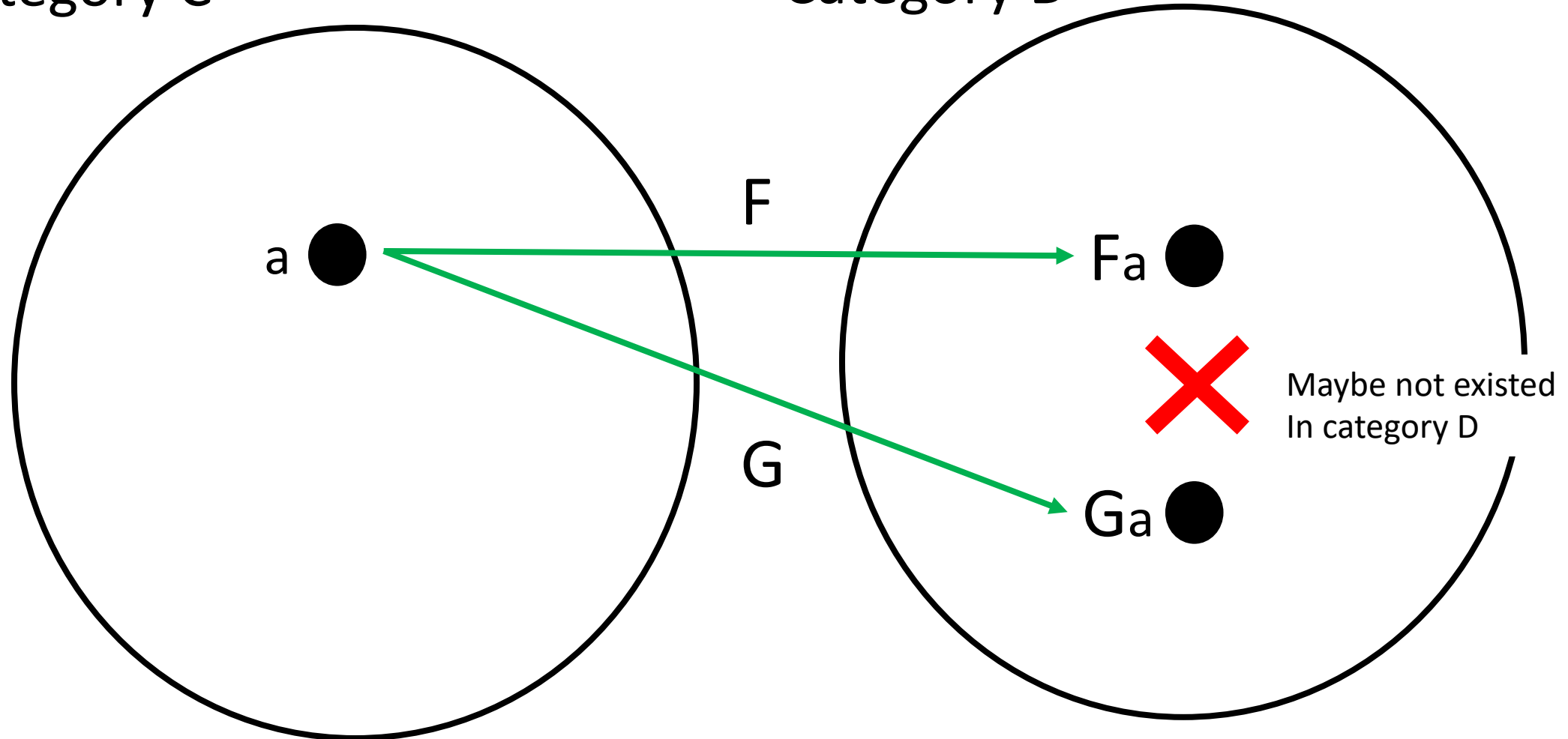
Category D



Natural Transformation

Category C

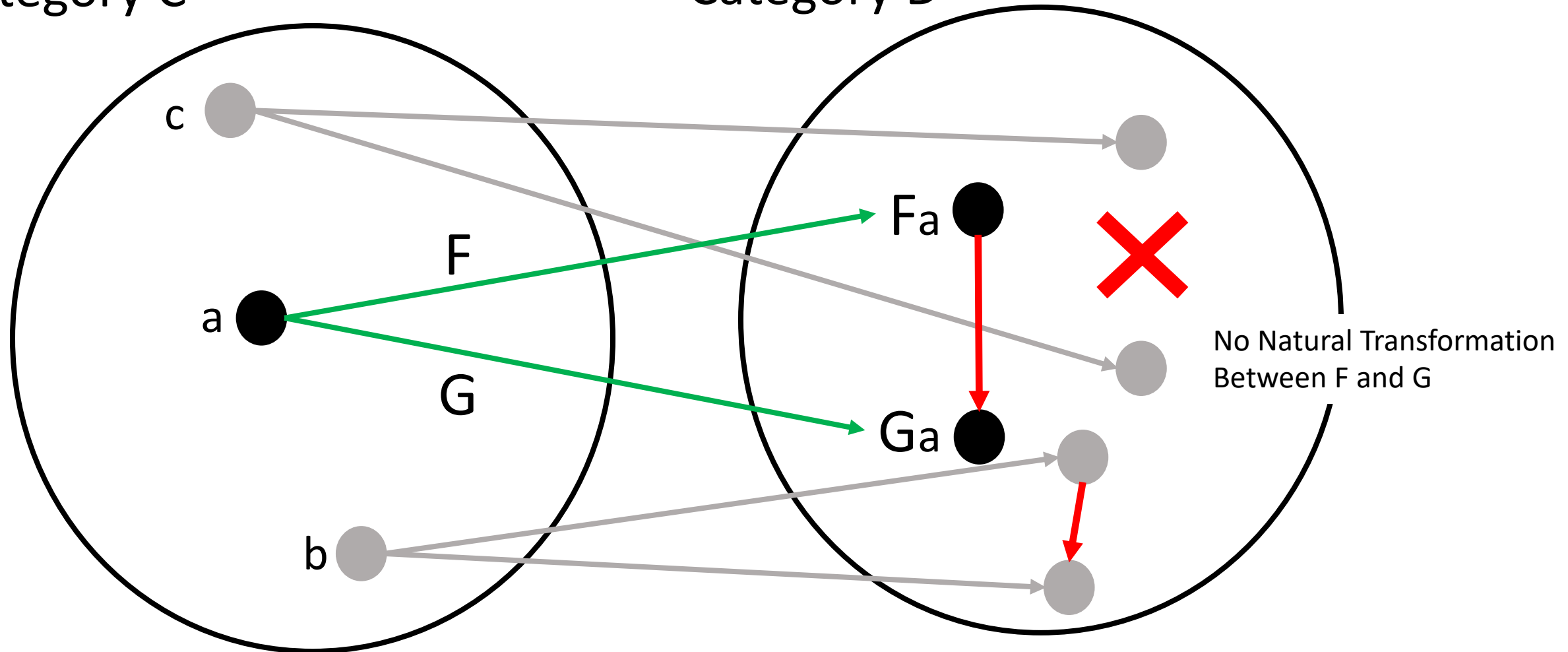
Category D



Natural Transformation

Category C

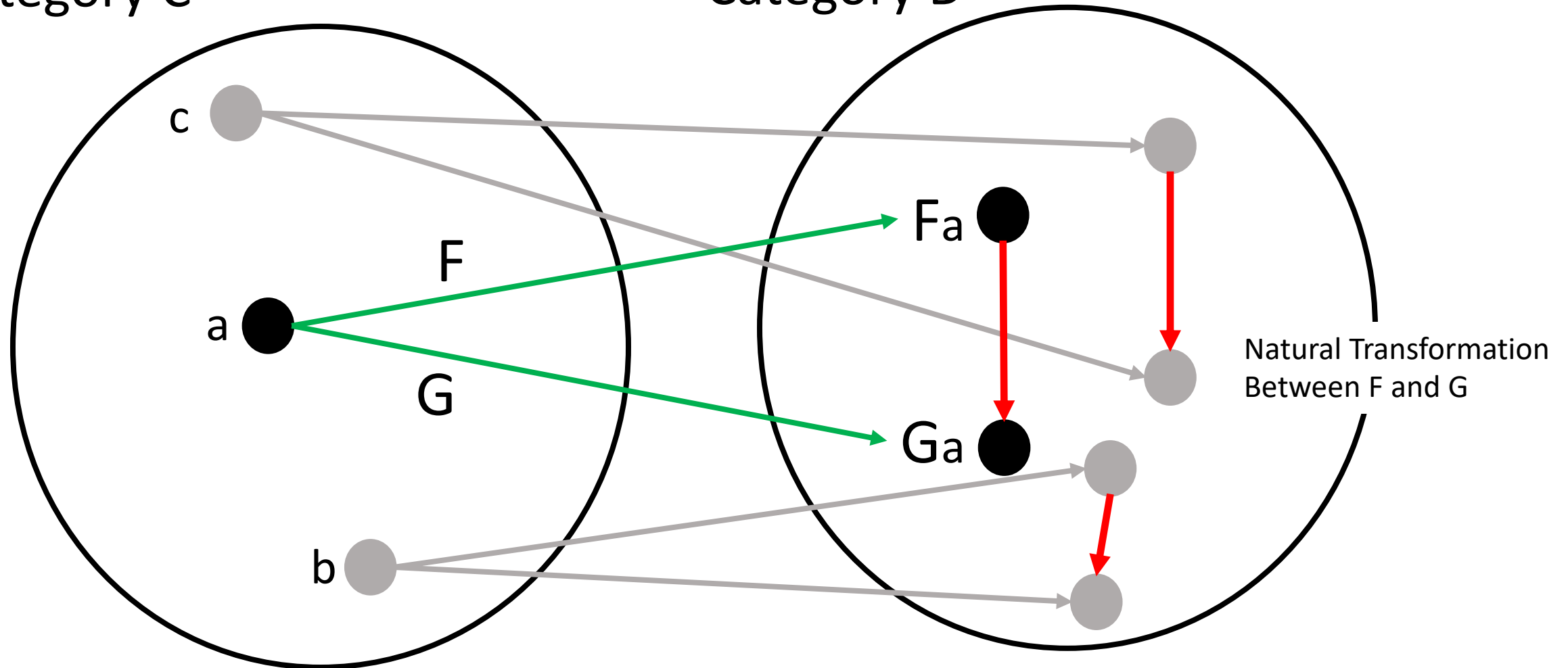
Category D



Natural Transformation

Category C

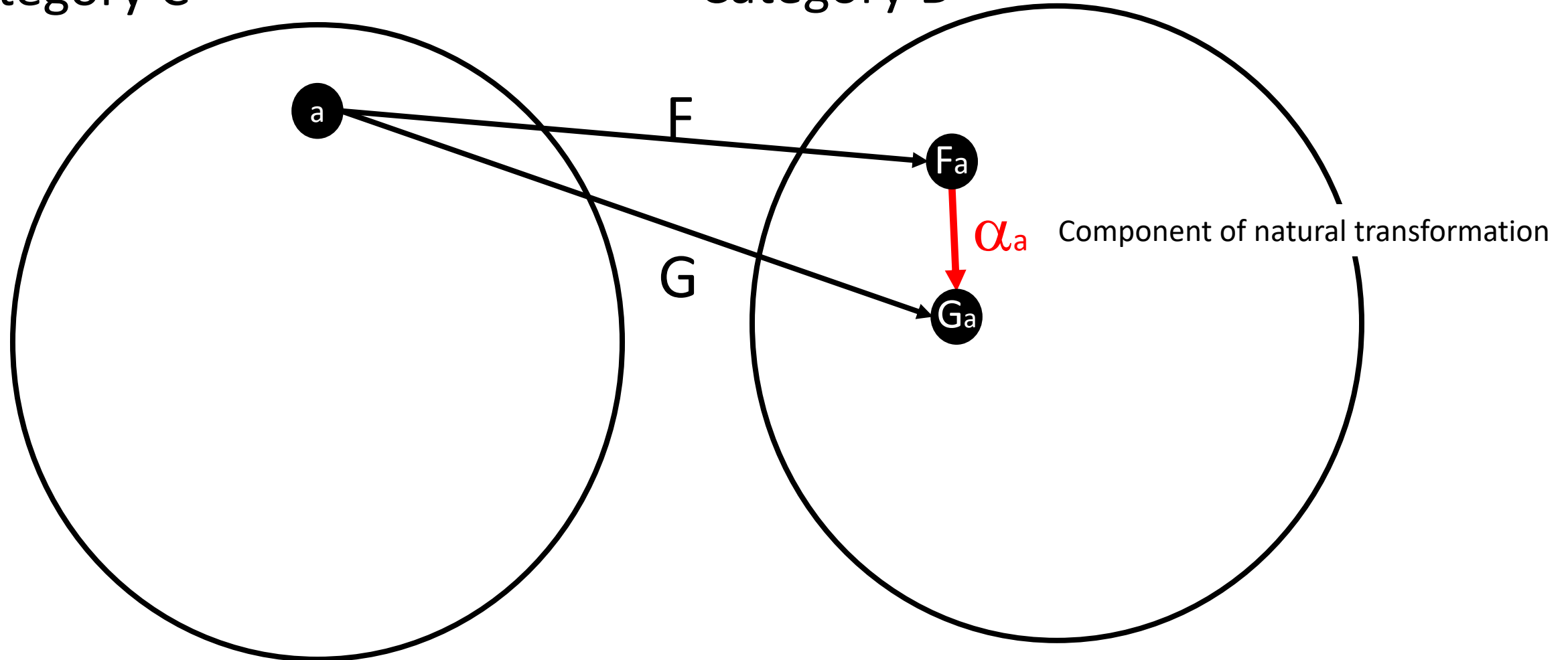
Category D



Natural Transformation – Preserving Structure

Category C

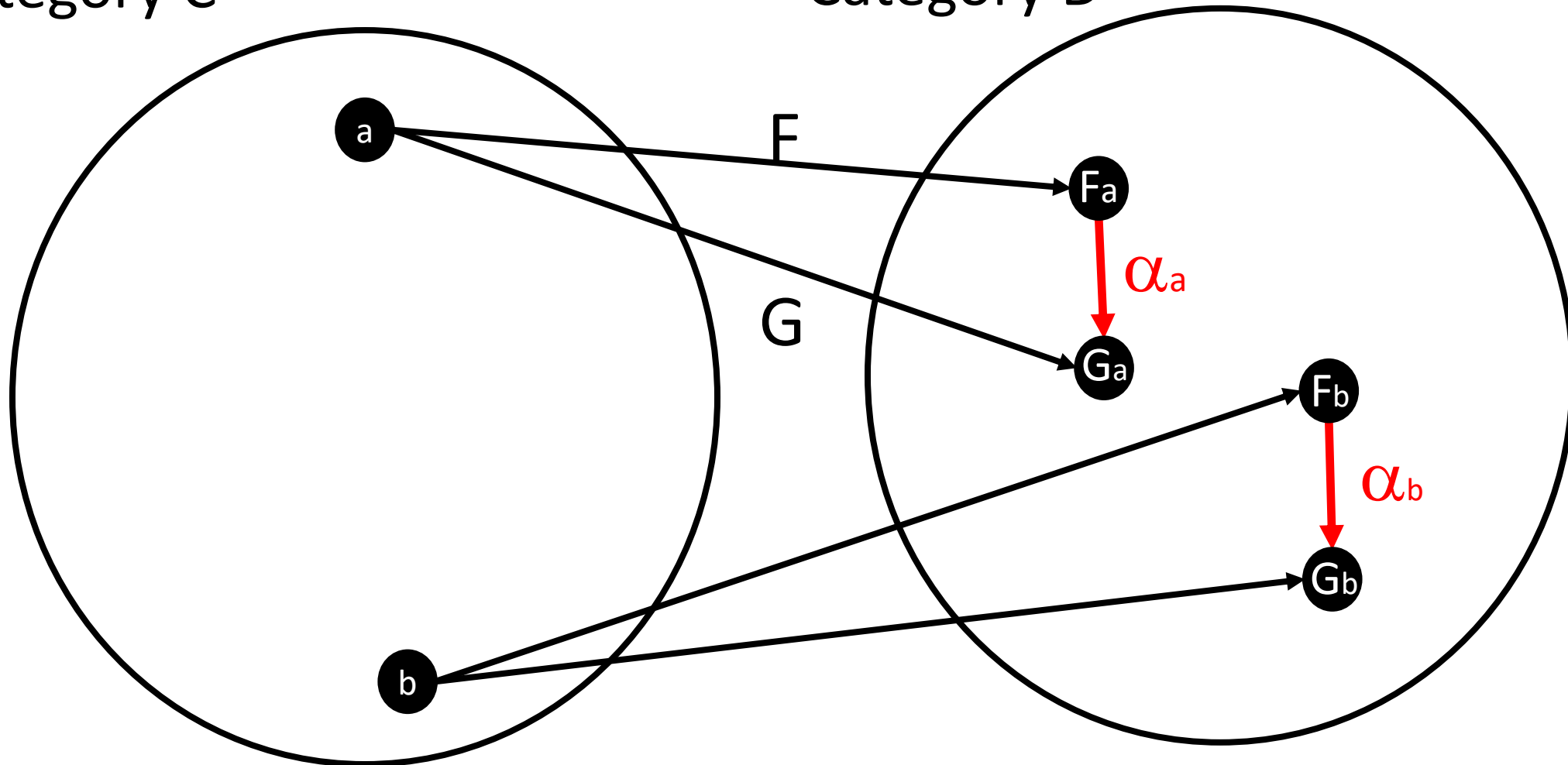
Category D



Natural Transformation – Preserving Structure

Category C

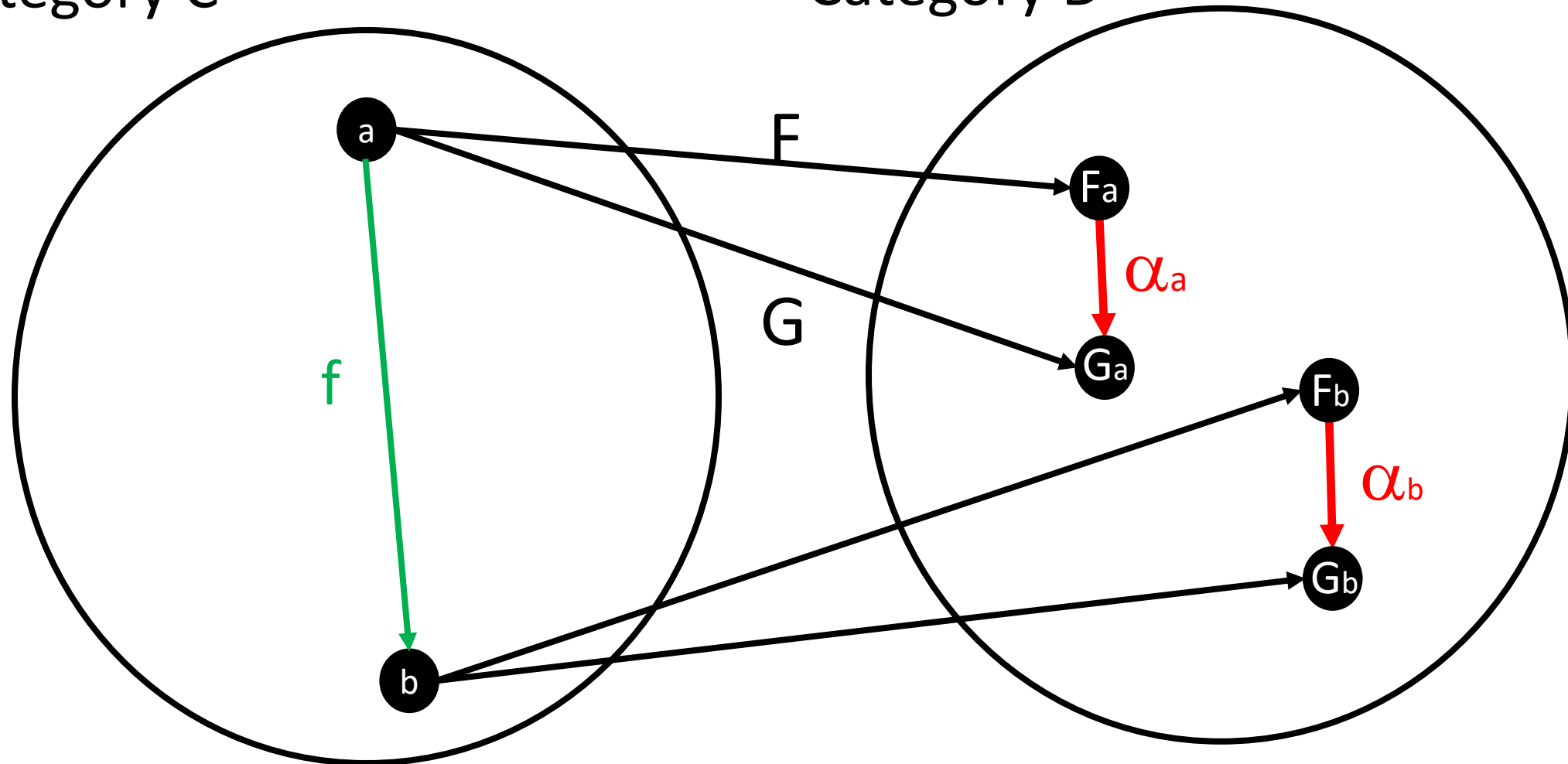
Category D



Natural Transformation – Preserving Structure

Category C

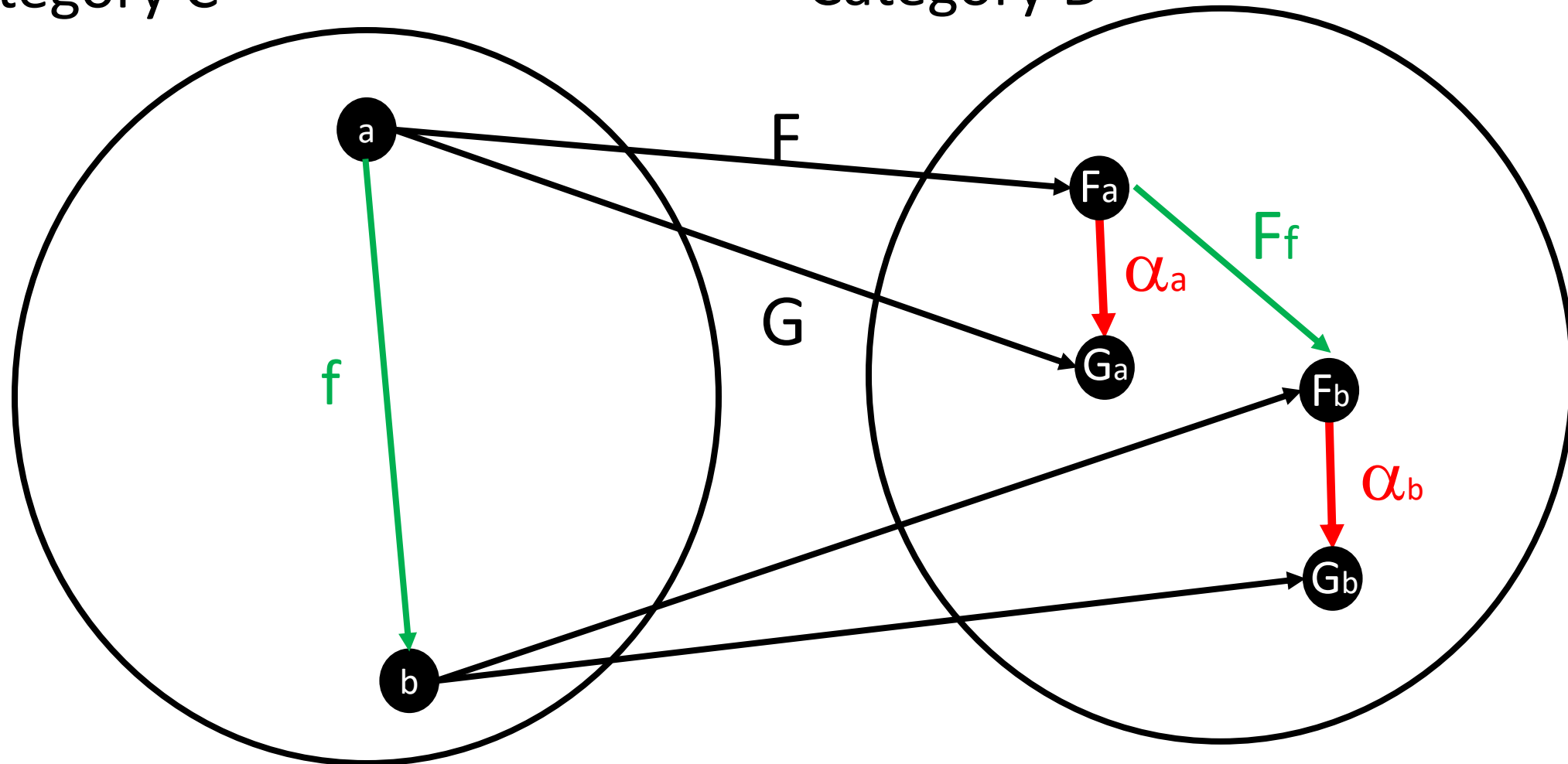
Category D



Natural Transformation – Preserving Structure

Category C

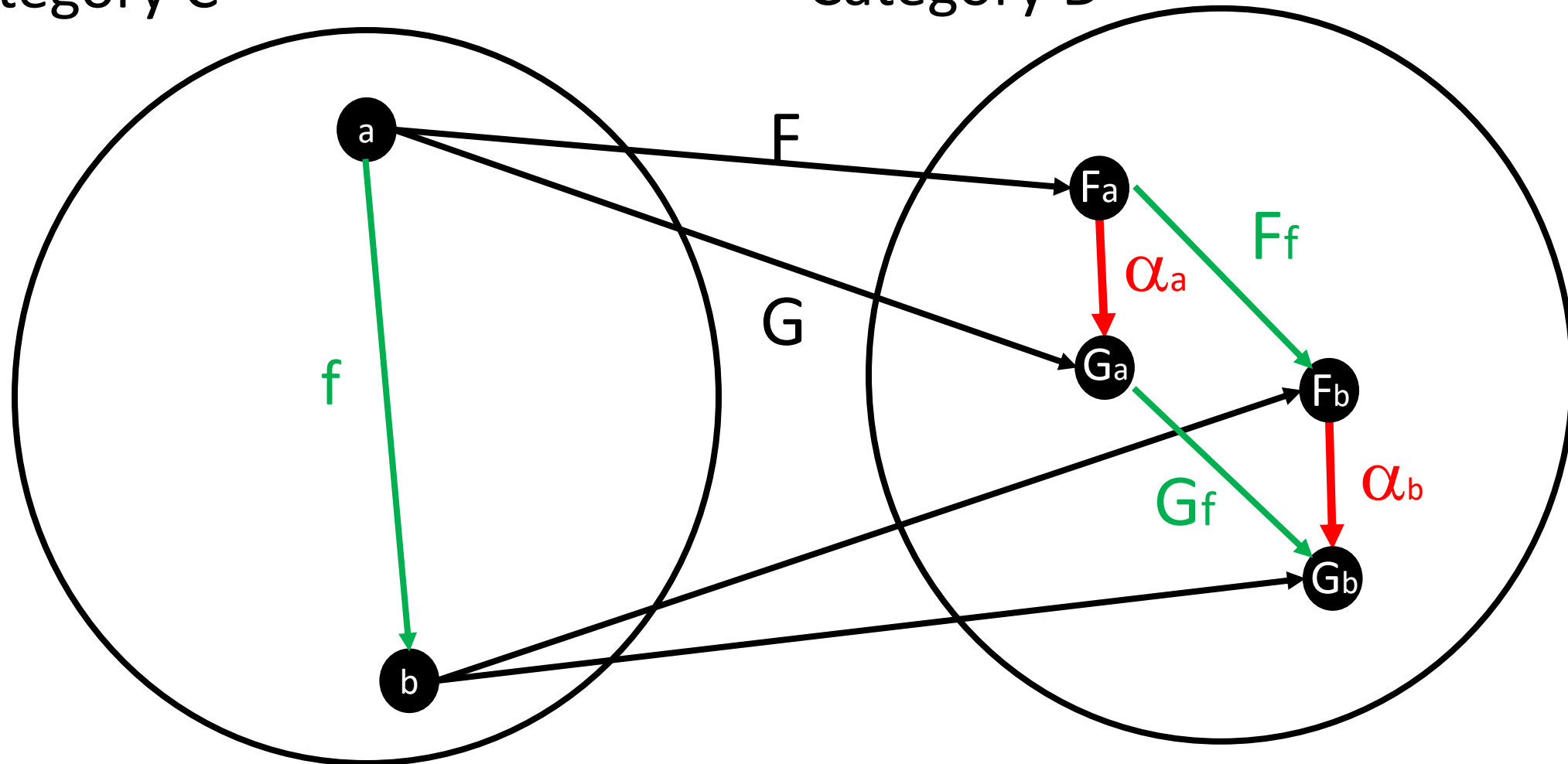
Category D



Natural Transformation – Preserving Structure

Category C

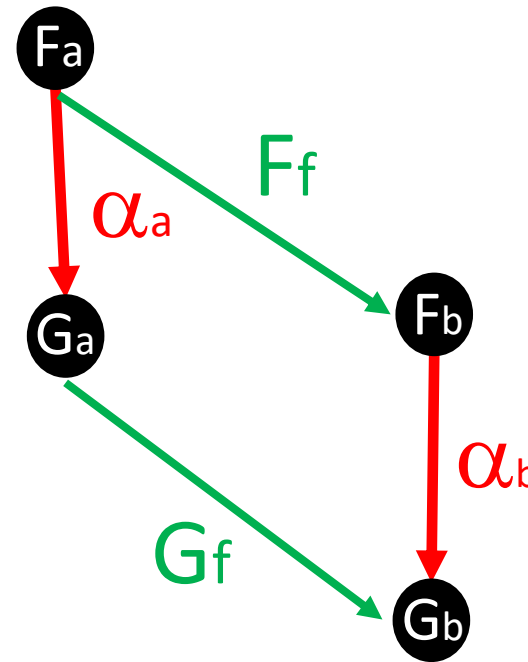
Category D



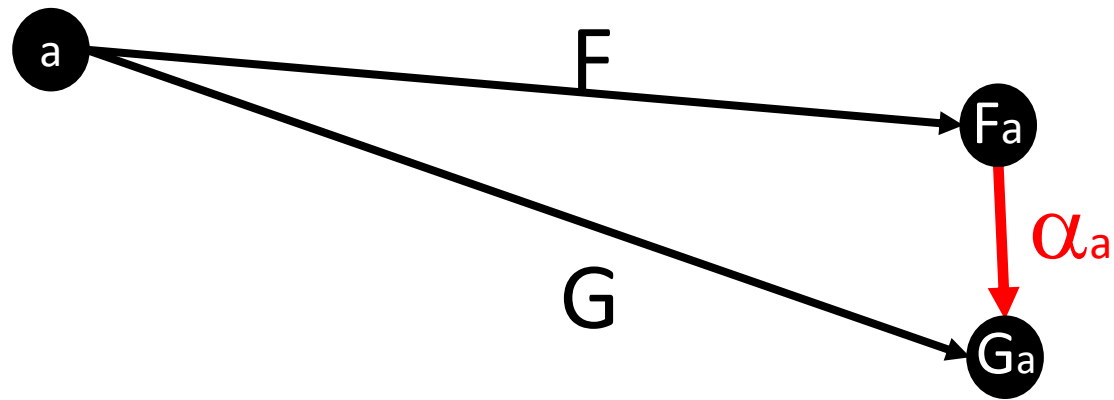
Natural Transformation – Preserving Structure

Relationship between F_f and G_f
is Naturality Square

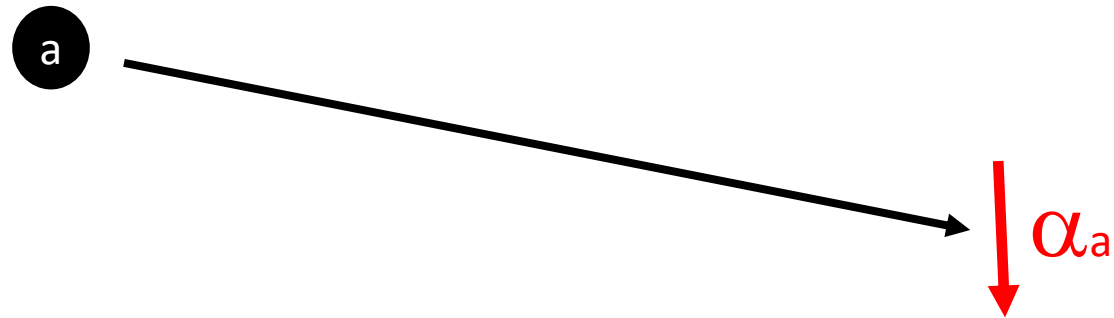
$$\alpha_b \circ F_f = G_f \circ \alpha_a$$



Natural Transformation

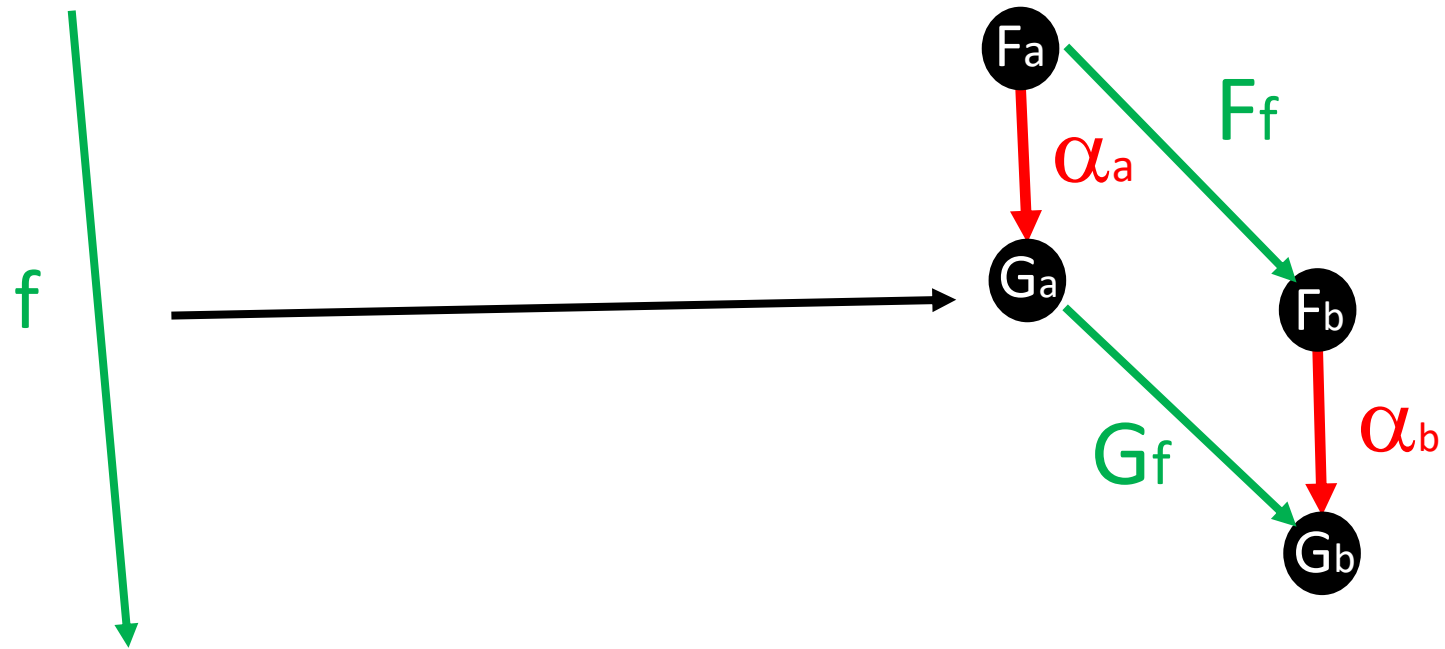


Natural Transformation



Mapping object a to morphism α_a

Natural Transformation

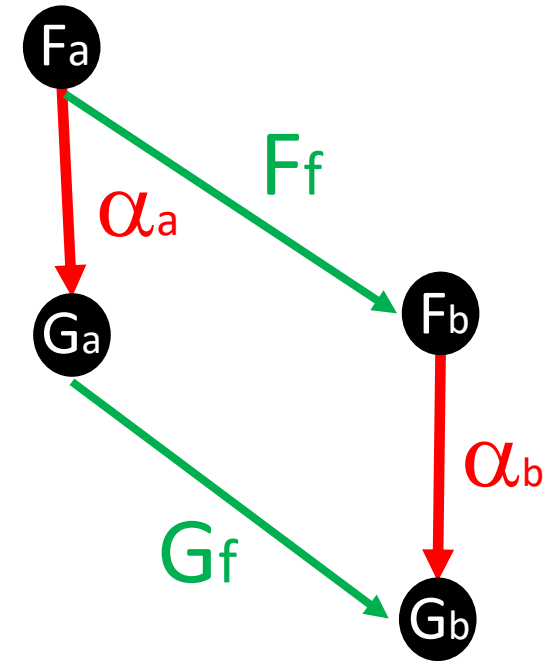


Mapping morphism f to a commuting diagram

Natural Transformation-Polymorphic function

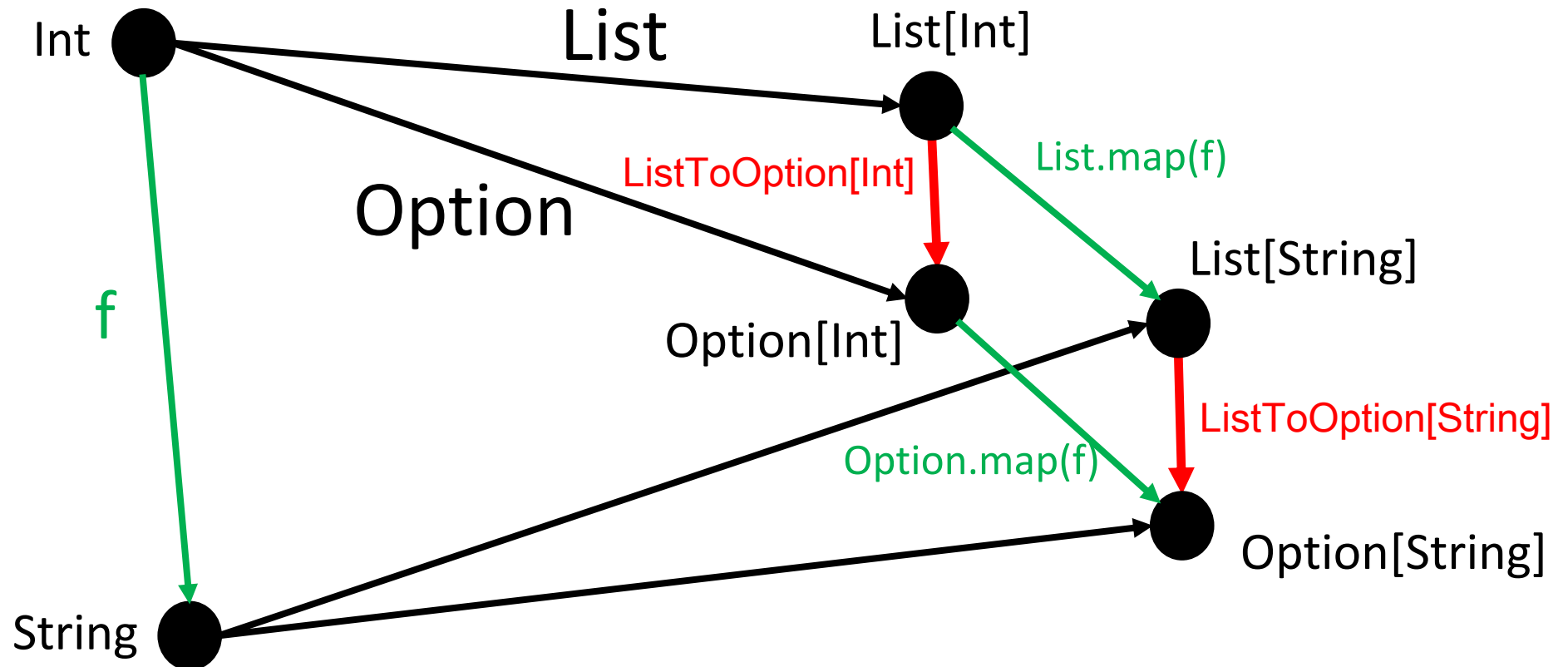
α : forall a, $F_a \Rightarrow G_a$

$$\alpha_b \circ \text{map}_F f = \text{map}_G f \circ \alpha_a$$



Natural Transformation-Parametric function

```
val ListToOption[A]: List[A] => Option[A] = list => list.headOption
```



Functor vs Natural Transformation

- Functor as a container
 - *map* modifies the content without reshaping the container
- Natural Transformation
 - To fulfill that mapping all objects from Functor F to Functor F
 - The content object is too polymorphic to modify it
 - Reshaping/transforming one container into another container without modifying content
 - Polymorphic/parametric functions

Reference

- Bartosz Milewski – Category theory for programmer
 - [Category Theory 1.2: What is a category?](#)
 - [Category Theory 6.1: Functors](#)
 - [Category Theory 9.1: Natural transformations](#)
- [Category Theory for Programmers: The Preface](#)

Thank you for your attention