David Eccles



INFO90002 Database Systems & Information Modelling

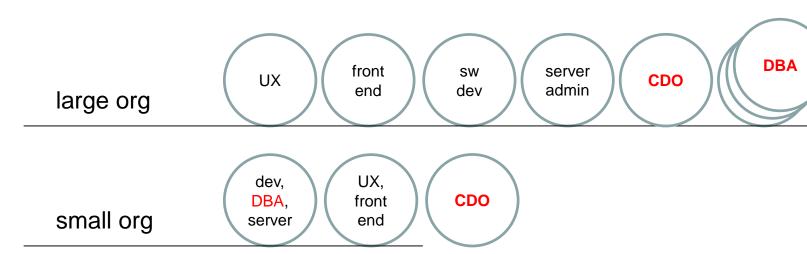
Week 8
Database Administration

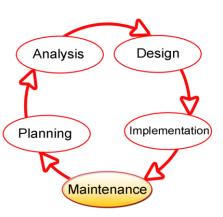
- The Database 'Administrator' role
- The Data Manager role
- Architecture Understanding the DBMS
 - concepts
- Performance improvement
 - concepts
 - common approaches e.g. indexes



The DBA role

- primarily concerned with "maintenance" / "ops" phase
- but should be consulted during all phases of development
- "Database Administrator" or "DBA"
- often framed as a "job" or a "person"
- Large companies many DBA's
- Small company developer is the DBA
- DBA role can be made redundant by Cloud-based DBMS or "database as a service" DAAS (often IAAS or PAAS)

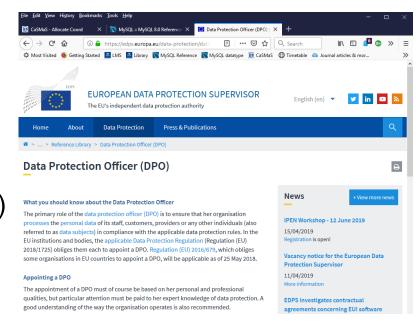






Data and Database Administration

- Data Administrator (CDO / DPO)
 - data policies, procedures and standards
 - planning
 - data conflict resolution
 - managing data
 - internal marketing & education
 - compliance
- Database Administrator (technical role)
 - analyze and design DB
 - select DBMS / tools / vendor
 - install and upgrade DBMS
 - tune DBMS performance
 - manage security, privacy, integrity
 - backup and recovery



(Hoffer et al., chapter 11)



Architecture of a Database Management System (DBMS)

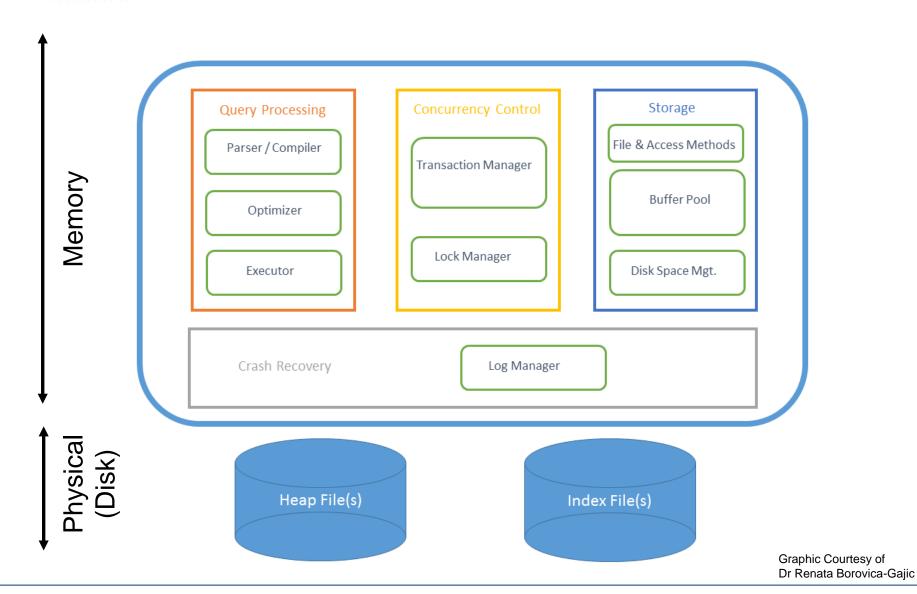


Database Systems Architecture

- A Database Management System (DBMS)
- Exists as one entity in two places
 - In Memory
 - Physically on disk
- Both places manage
 - Data (the reason we have the DBMS)
 - Performance (how it performs as it is used & grows)
 - Concurrency (manage high volumes of users)
 - Availability (assist in recovery and availability)
- One place is persistent the other transient
 - Disk representation is always present
 - Memory transient only exists when DBMS is running



DBMS Overview





Query Processing



Parsing

- Syntax is correct & can "compile"
- DBMS User Permissions
- Resources (Data, Code, be able to Record Changes/Retrieve results)

Optimizing

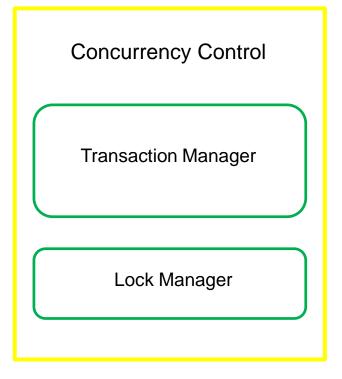
- Execution Plan & Execution Cost
- Evaluate indexes, table scans, hashing
- Eliminate worst, consider best options
- Lowest cost theoretically "best"

Execution

- Meet the ACID test,
- Atomic: All rows succeed or all fail
- Ensure resources are available
 - Data, Log changes, Memory, Cursor to do the work for the USER



Concurrency Control

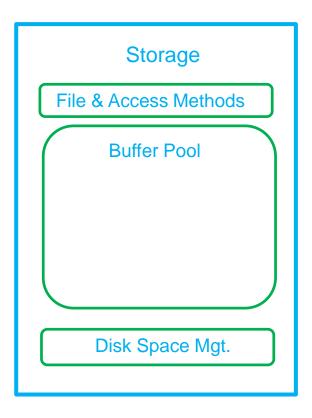


- Manages the work of the DBMS.
- Transaction Manager handles all aspects of the SQL transaction - which DBMS user wants WHAT resource
- Lock Manager is a list of what resources are locked and by which user at what level
- not only tables, indexes
 - buffers, cursor, memory addresses of resources



MELBOURNE Concurrency Control

- Essential to manage large scalable DBMS
- Enables 1,000,000s of concurrent users
- Like a Traffic Policemen controlling the flow of traffic
 - Who can do what (allowed to do what they need to do)
 - Who has to wait (queue)
 - Who can travel through the intersection concurrently
 - Usually readers of data



- File & Access Methods
 - Disk to Memory to Disk
 - Read a buffer or a block of buffers
- Buffer Pool
 - Data in memory
 - Row data
 - Index data
 - Organised
- Disk Space Management
 - How to organise growth of data on disk efficiently by writing efficiently.

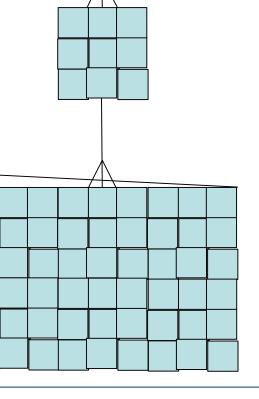


Storage – Disk Space Management

- Hierarchical Structure
- Storage of Object (tables, indexes, rollback, logs)
- Row by Row space management inefficient
- Buffer better

Free

Block Best



Row

Buffer

Block

- - Many Object Types (Tables, Indexes, Undo)
 - Each buffer contains rows, b+tree leaf etc.
 - Each buffer can have one of four status types:
 - Current
 - In use current committed version of data (row)
 - Active
 - Most recent change (may not be <u>committed</u>)
 - COMMIT (Current: DepartmentID=9)
 - Stale
 - An old version of the data
 - Aged
 - Old and about to be removed from buffer pool

Buffer Pool

DepartartmentID=8

DepartartmentID=9

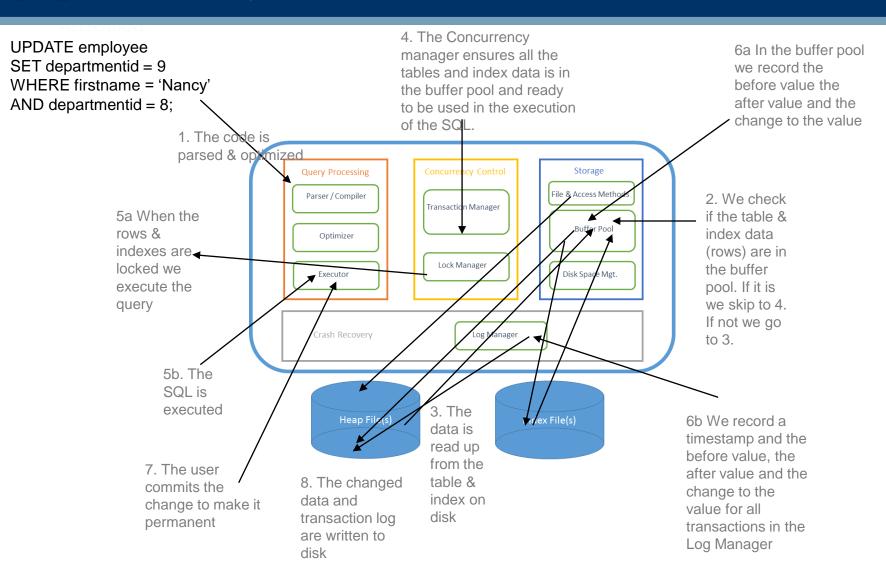
DepartmentID=8

DepartmentID=8

Crash Recovery

Log Manager

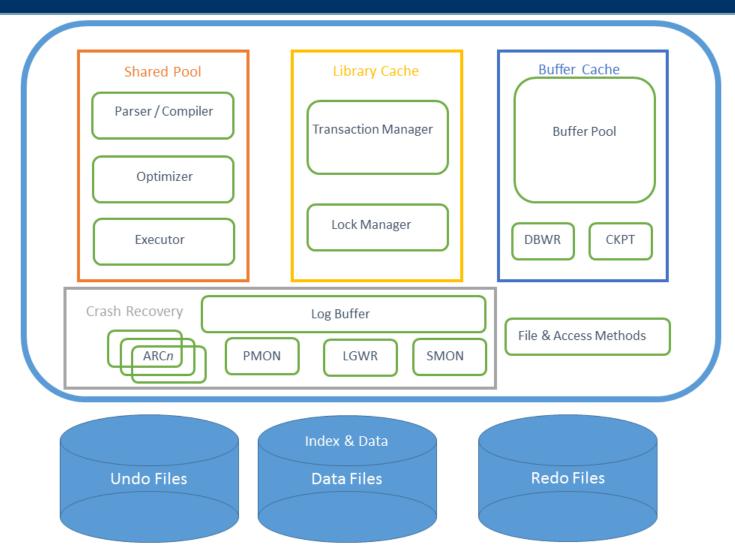
- Recovery
- Log Manager records ALL changes
 - Statement
 - Transaction
 - Statement
 - Rollback values
 - Before and After values
 - Timestamp begin
 - transaction, savepoint & commit timestamps
 - Database
 - Data Dictionary Changes



Graphic Courtesy of Dr Renata Borovica-Gajic



How Oracle* DBMS looks



^{*} This slide is not examinable

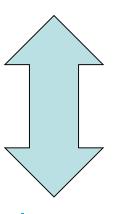


Database Performance

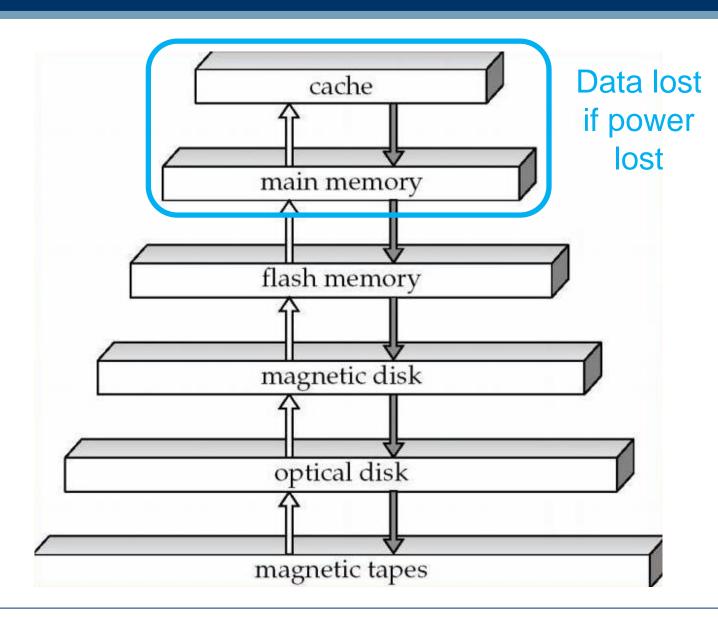


Storage media hierarchy

faster, more expensive, smaller capacity



slower, cheaper, older, bigger capacity





What affects database performance?

- Caching data in memory, e.g. data buffers
- Placement of data files across disc drives
- Fast storage such as SSD
- Database replication and server clustering
- Use of indexes to speed up searches and joins
- Good choice of data types (especially PKs)
- Good program logic (no long running CRUD)
- Good query execution plans
- Good code (no deadlocks)



Caching Data in Memory

- Data and Code found in memory
- Avoids a read
- Reads are expensive
- Goal in to minimize reads (& writes)
 - Writes are necessary (recovery logs, changed data)
- "in memory databases"
 - all code all data loaded into memory on db start & stays until shutdown

Buffer Pool (Table & Index rows)

Parser / Compiler (SQL)

Data file location & Fast Disks (SSD)

- Spread the files across the physical server
 - RAID (0, 0 + 1, 5)
- We can't avoid writes
 - Spread files across many disks
 - Avoid contention
 - (many users competing for same resource)
 - Recovery Logs (always writing)
 - faster disk
- SSD (Solid State Drives)
 - No moving parts nothing to break down
 - Faster I/O (Input & Output compared to other disk types)



Distribution & Replication

- Distributed data
 - Spreads the load
 - Data kept only where it is needed
 - Less work per physical server – faster response times



- Replicated Data
 - Spreads Load
 - Less work per physical server – faster response times



When to create indexes

- for each table, choose the columns you will index:
- column is queried frequently (used in Where clauses)
- columns that are used for joins (PK to FK)
- primary keys (automatic in most DBMS)
- foreign keys (automatic in MySQL)
- unique columns (automatic in most DBMS)
- large tables only small tables do not require indexes
 - if you frequently retrieve less than about 15% of the rows
- wide range of values (good for regular indexes).
- small range of values (good for bitmap/hash indexes).

covered already in week 6

source: Oracle® Database Application Developer's Guide

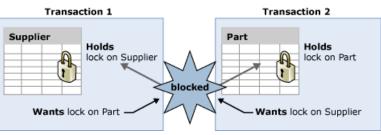
- Good Data Types
 - INTEGERS for PK FK & PFK
- Good Program Logic & Code
 - Transaction design

BEGIN TRANSACTION
SELECT
UPDATE
UPDATE
UPDATE
COMMIT

Avoid Long Complex Transactions that never commit or

savepoint

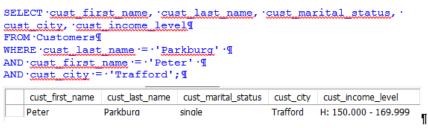
- Avoid coding deadlocks!
- Appropriate Locking strategy
 - Row b4 Buffer b4 Table b4 Database
 - Consider Lock Timeouts (if not automatic)



Deadlock Scenario

Good Execution Plans

- The best execution plan has the lowest "cost"
- Known as Cost Based Optimization (CBO)



query_block#1

12135.2 54.75K rows

Full Table Scan

Customers

Index on where condition (cust_city) improves cost:

```
CREATE · INDEX · cust city idx · ¶

ON · customers (cust city); ¶

¶

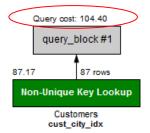
SELECT · cust first name, · cust last name, · cust marital status, · cust city, · cust income level ¶

FROM · Customers ¶

WHERE · cust last name · = · ' Parkburg' · ¶

AND · cust first name · = · ' Peter' · ¶

AND · cust city · = · ' Trafford'; ¶
```



Performance Tuning

- Collect a baseline of performance statistics
 - (EG Long running update takes 43 seconds)
 - Retrieval of all rows in the largest table takes 14 seconds
 - Snapshot memory, usage at the same time every day every week.
- Initiate a SINGLE change
 - Allocate more memory to the database

OR (NOT BOTH!)

- Spread the files over different disks
- Retest your baseline metrics
 - Did performance improve or decline
- It is the change in the value not the absolute value that is important in Database Performance Tuning

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