

Propositional Logic

Dave Dubin and Jodi Schneider

October 10, 2016

Concepts from earlier in the semester

- Propositions are the bearers of truth values: the kinds of things that can be true or false.

Concepts from earlier in the semester

- Propositions are the bearers of truth values: the kinds of things that can be true or false.
- States of affairs are the parts of reality responsible for making propositions true or false.

- We represent propositions with lower case letters (typically p , q , and r).

- We represent propositions with lower case letters (typically p , q , and r).
- A set of proposition letters generates a set of states of affairs: ways the world might be.

- We represent propositions with lower case letters (typically p , q , and r).
- A set of proposition letters generates a set of states of affairs: ways the world might be.
- n proposition letters generate 2^n states of affairs.

- We represent propositions with lower case letters (typically p , q , and r).
- A set of proposition letters generates a set of states of affairs: ways the world might be.
- n proposition letters generate 2^n states of affairs.
- $\{pqr, pq\bar{r}, p\bar{q}r, p\bar{q}\bar{r}, \bar{p}qr, \bar{p}q\bar{r}, \bar{p}\bar{q}r, \bar{p}\bar{q}\bar{r}\}$

Logical Operators

Table 1: 2.15 from van Benthem, et al.

| Symbol | In natural language | Technical name |
|-------------------|---------------------|----------------|
| \neg | not | negation |
| \wedge | and | conjunction |
| \vee | or | disjunction |
| \rightarrow | if ... then | implication |
| \leftrightarrow | if and only if | equivalence |

- A sentence constructed from proposition letters and operators is true or false in each state of affairs.

- A sentence constructed from proposition letters and operators is true or false in each state of affairs.
- Consider, for example: $(\neg p \vee q) \rightarrow r$

- A sentence constructed from proposition letters and operators is true or false in each state of affairs.
- Consider, for example: $(\neg p \vee q) \rightarrow r$
- The sentence is mapped to a truth value via the following tables

Semantics of the operators

| φ | $\neg\varphi$ |
|-----------|---------------|
| 0 | 1 |
| 1 | 0 |

Table 3: 2.18 from van Benthem, et al.

| φ | ψ | $\varphi \wedge \psi$ | $\varphi \vee \psi$ | $\varphi \rightarrow \psi$ | $\varphi \leftrightarrow \psi$ |
|-----------|--------|-----------------------|---------------------|----------------------------|--------------------------------|
| 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 | 1 | 1 |

Exercise 2.3

You are given the information that p -or- q and $(\text{not-}p)$ -or- r . What can you conclude about q and r ? What is the strongest valid conclusion you can draw? (A statement is stronger than another statement if it rules out more possibilities.)

Drawing truth tables for expressions

| p | q | r | $((\sim p \vee q) \rightarrow r)$ | | | | |
|---|---|---|-----------------------------------|---|---|---|---|
| 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |

Drawing truth tables for expressions

| p | q | r | $((\sim p \vee q) \rightarrow r)$ | | | | |
|---|---|---|--------------------------------------|---|---|---|---|
| 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |

Drawing truth tables for expressions

| p | q | r | $((\sim p \vee q) \rightarrow r)$ | | | | |
|---|---|---|-----------------------------------|---|---|---|---|
| 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |

Drawing truth tables for expressions

| p | q | r | $((\sim p \vee q) \rightarrow r)$ | | | | |
|---|---|---|-----------------------------------|---|---|---|---|
| 1 | 1 | 1 | 0 | 1 | 1 | 1 | 1 |
| 1 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 1 | 0 | 1 | 0 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 0 | 0 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |

Grammar of propositional logic

Let P be a set of proposition letters and let $p \in P$.

The following expression defines the recursive grammar for a logical expression φ in Backus–Naur Form:

$$\varphi ::= p \mid \neg\varphi \mid (\varphi \wedge \varphi) \mid (\varphi \vee \varphi) \mid (\varphi \rightarrow \varphi) \mid (\varphi \leftrightarrow \varphi)$$

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

- $\neg \vee p$

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

- $\neg \vee p$
- $\vee)p\neg$

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

- $\neg \vee p$
- $\vee)p\neg$
- $\neg p \vee q \rightarrow r$

How many correct expressions are consistent with the last one?

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$
- $\neg((p \vee q) \rightarrow r)$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$
- $\neg((p \vee q) \rightarrow r)$
- $(\neg p \vee (q \rightarrow r))$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$
- $\neg((p \vee q) \rightarrow r)$
- $(\neg p \vee (q \rightarrow r))$
- $\neg(p \vee (q \rightarrow r))$

- Semantics is the relationship of a language to the part of the world that we're modeling.

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.
- “ $V(\varphi) = 1$ ” means the formula (or sentence) φ is true in the state of affairs represented by the function V . “ $V(\varphi) = 0$ ” means that φ is false in the state of affairs represented by the function V .

More on semantics

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.
- “ $V(\varphi) = 1$ ” means the formula (or sentence) φ is true in the state of affairs represented by the function V . “ $V(\varphi) = 0$ ” means that φ is false in the state of affairs represented by the function V .
- For “ $V(\varphi) = 1$ ” we also write “ $V \models \varphi$ ” read as “ V is a model of φ ” or “ V satisfies φ .”

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.
- “ $V(\varphi) = 1$ ” means the formula (or sentence) φ is true in the state of affairs represented by the function V . “ $V(\varphi) = 0$ ” means that φ is false in the state of affairs represented by the function V .
- For “ $V(\varphi) = 1$ ” we also write “ $V \models \varphi$ ” read as “ V is a model of φ ” or “ V satisfies φ .”
- If V doesn't satisfy φ we write “ $V \not\models \varphi$ ”. In other words $V(\varphi) = 0$.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.
- Examples:

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.
- Examples:
 - ① $(q \vee \neg q)$ is logically true.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.
- Examples:
 - ① $(q \vee \neg q)$ is logically true.
 - ② $(q \wedge \neg q)$ is logically false.

- A set of propositional logic statements is consistent if at least one state of affairs satisfies every statement in the set.

- A set of propositional logic statements is consistent if at least one state of affairs satisfies every statement in the set.
- A set of propositional logic statements is inconsistent if no state of affairs satisfies every statement in the set.

- A conclusion is *valid* with respect to a set of premises if the conclusion is true in every situation where the premises are true (van Benthem, et al, page 2-4).

- A conclusion is *valid* with respect to a set of premises if the conclusion is true in every situation where the premises are true (van Benthem, et al, page 2-4).
- One can validly infer a conclusion φ from a set of premises P if the negation of φ is inconsistent with the set of statements P .

(From van Benthem, et al., chapter 2)

- ① Computing a truth value for a formula takes linear time.

(From van Benthem, et al., chapter 2)

- 1 Computing a truth value for a formula takes linear time.
- 2 Computing a truth table for validity takes exponential time.

(From van Benthem, et al., chapter 2)

- ① Computing a truth value for a formula takes linear time.
- ② Computing a truth table for validity takes exponential time.
- ③ The problem of testing for validity in propositional logic is decidable: there exists a mechanical method that computes the answer, at least in principle.