

Propositional Logic

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Fall Semester, 2018

- In this class we'll give some attention to logical reasoning and inference.

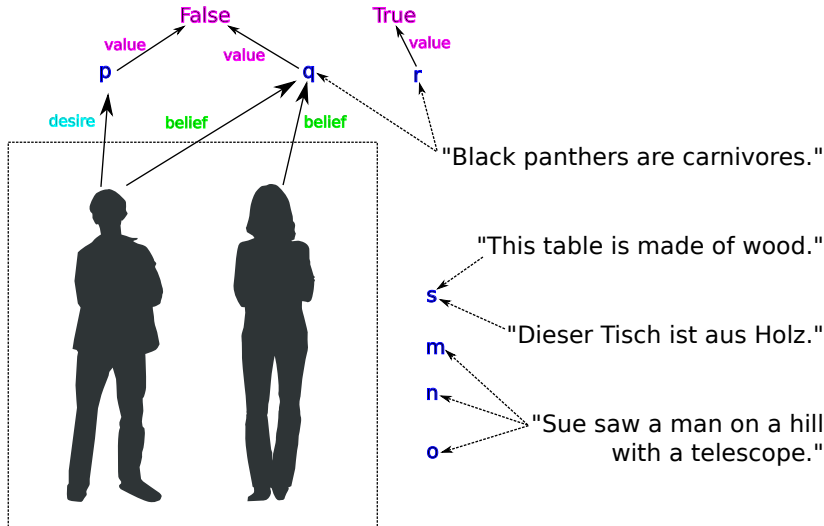
Reasoning vs. representation

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- The same rules that govern valid argumentation also enable information systems to answer interesting questions.
- Logic is also a means to describe a *domain* (some part of the world we're interested in).

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- $\{pqr, pq\bar{r}, p\bar{q}r, p\bar{q}\bar{r}, \bar{p}qr, \bar{p}q\bar{r}, \bar{p}\bar{q}r, \bar{p}\bar{q}\bar{r}\}$

Logical Operators

Table 1: 2.15 from van Benthem, et al.

Symbol	In natural language	Technical name
\neg	not	negation
\wedge	and	conjunction
\vee	or	disjunction
\rightarrow	if ... then	implication
\leftrightarrow	if and only if	equivalence

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- The sentence is mapped to a truth value via the following tables

Semantics of the operators

φ	$\neg\varphi$
0	1
1	0

Table 3: 2.18 from van Benthem, et al.

φ	ψ	$\varphi \wedge \psi$	$\varphi \vee \psi$	$\varphi \rightarrow \psi$	$\varphi \leftrightarrow \psi$
0	0	0	0	1	1
0	1	0	1	1	0
1	0	0	1	0	0
1	1	1	1	1	1

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- a = “He is currently employed abroad.”
- s = “He is entitled to social security.”

Drawing truth tables for expressions

p	q	r	$((\neg p \vee q) \rightarrow r)$				
1	1	1	0	1	1	1	1
1	1	0	0	1	1	0	0
1	0	1	0	1	0	1	1
1	0	0	0	1	0	1	0
0	1	1	1	0	1	1	1
0	1	0	1	0	1	0	0
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0	1	0	1	0	1	0	0
0	0	1	1	0	1	1	1
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Exercise 2.11

p	q	$((p \rightarrow q) \vee (q \rightarrow p))$						
1	1	1	1	1	1	1	1	1
1	0	1	0	0	1	0	1	1
0	1	0	1	1	1	1	0	0
0	0	0	1	0	1	0	1	0

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p	q	r	$(((p \vee \neg q) \wedge r) \leftrightarrow (\neg (p \wedge r) \vee q))$										
1	1	1	1	1	0	1	1	1	0	1	1	1	1
1	1	0	1	1	0	1	0	0	0	1	1	0	0
1	0	1	1	1	1	0	1	1	0	1	1	1	0
1	0	0	1	1	1	0	0	0	0	1	1	0	0
0	1	1	0	0	0	1	0	1	0	1	0	0	1
0	1	0	0	0	0	1	0	0	0	1	0	0	1
0	0	1	0	1	1	0	1	1	1	0	0	1	1
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Grammar of propositional logic

Let P be a set of proposition letters and let $p \in P$.

The following expression defines the recursive grammar for a logical expression φ in Backus–Naur Form:

$$\varphi ::= p \mid \neg \varphi \mid (\varphi \wedge \varphi) \mid (\varphi \vee \varphi) \mid (\varphi \rightarrow \varphi) \mid (\varphi \leftrightarrow \varphi)$$

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

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- $\neg p \vee q \rightarrow r$

How many correct expressions are consistent with the last one?

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

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- $(\neg p \vee (q \rightarrow r))$

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- $\neg\neg \wedge q \vee p$
- $p\neg q$

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- For " $V(\varphi) = 1$ " we also write " $V \models \varphi$ " read as " V is a model of φ " or " V satisfies φ ."
- If V doesn't satisfy φ we write " $V \not\models \varphi$ ". In other words $V(\varphi) = 0$.

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- A set of propositional logic statements is consistent if at least one state of affairs satisfies every statement in the set.

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- A set of propositional logic statements is inconsistent if no state of affairs satisfies every statement in the set.

- A conclusion is *valid* with respect to a set of premises if the conclusion is true in every situation where the premises are true (van Benthem, et al, page 2-4).

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- One can validly infer a conclusion φ from a set of premises P if the negation of φ is inconsistent with the set of statements P .

(From van Benthem, et al., chapter 2)

- ① Computing a truth value for a formula takes linear time.

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- ① Computing a truth value for a formula takes linear time.
- ② Computing a truth table for validity takes exponential time.
- ③ The problem of testing for validity in propositional logic is decidable: there exists a mechanical method that computes the answer, at least in principle.