

# Propositional Logic

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- States of affairs are the parts of reality responsible for making propositions true or false.

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- $\{pqr, pq\bar{r}, p\bar{q}r, p\bar{q}\bar{r}, \bar{p}qr, \bar{p}q\bar{r}, \bar{p}\bar{q}r, \bar{p}\bar{q}\bar{r}\}$

Table 1: 2.15 from van Benthem, et al.

Symbol	In natural language	Technical name
$\neg$	not	negation
$\wedge$	and	conjunction
$\vee$	or	disjunction
$\rightarrow$	if ... then	implication
$\leftrightarrow$	if and only if	equivalence



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- The sentence is mapped to a truth value via the following tables

# Semantics of the operators

$\phi$	$\neg\phi$
0	1
1	0

Table 3: 2.18 from van Benthem, et al.

$\phi$	$\psi$	$\phi \wedge \psi$	$\phi \vee \psi$	$\phi \rightarrow \psi$	$\phi \leftrightarrow \psi$
0	0	0	0	1	1
0	1	0	1	1	0
1	0	0	1	0	0
1	1	1	1	1	1

## Exercise 2.3

You are given the information that  $p$ -or- $q$  and  $(\text{not-}p)$ -or- $r$ . What can you conclude about  $q$  and  $r$ ? What is the strongest valid conclusion you can draw? (A statement is stronger than another statement if it rules out more possibilities.)

# Drawing truth tables for expressions

p	q	r	$((\sim p \vee q) \rightarrow r)$				
1	1	1	0	1	1	1	1
1	1	0	0	1	1	0	0
1	0	1	0	1	0	1	1
1	0	0	0	1	0	1	0
0	1	1	1	0	1	1	1
0	1	0	1	0	1	0	0
0	0	1	1	0	1	1	1
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p	q	r	$(( \sim p \vee q ) \rightarrow r )$				
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1	1	0	0	1	1	0	0
1	0	1	0	1	0	1	1
1	0	0	0	1	0	1	0
0	1	1	1	0	1	1	1
0	1	0	1	0	1	0	0
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# Drawing truth tables for expressions

$(\sim p \vee q) \rightarrow r$

---

0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1

# Drawing truth tables for expressions

$(\sim p \vee q) \rightarrow r$

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10	0	0
10	0	1
10	1	0
10	1	1
01	0	0
01	0	1
01	1	0
01	1	1

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$(\sim p \vee q) \rightarrow r$

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10 1 0      0

10 1 0      1

10 1 1      0

10 1 1      1

01 0 0      0

01 0 0      1

01 1 1      0

01 1 1      1

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$(\sim p \vee q) \rightarrow r$

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10 1 0 0 0

10 1 0 1 1

10 1 1 0 0

10 1 1 1 1

01 0 0 1 0

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# Grammar of propositional logic

Let  $P$  be a set of proposition letters and let  $p \in P$ .

The following expression defines the recursive grammar for a logical expression  $\phi$  in Backus–Naur Form:

$$\phi ::= p \mid \neg\phi \mid (\phi \wedge \phi) \mid (\phi \vee \phi) \mid (\phi \rightarrow \phi) \mid (\phi \leftrightarrow \phi)$$

# Syntactically conforming expressions

Let  $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- $r$

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

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- $\vee)p\neg$
- $\neg p \vee q \rightarrow r$

How many correct expressions are consistent with the last one?

These conforming expressions are all consistent with  $\neg p \vee q \rightarrow r$

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- “ $V(\phi) = 1$ ” means the formula (or sentence)  $\phi$  is true in the state of affairs represented by the function  $V$ . “ $V(\phi) = 0$ ” means that  $\phi$  is false in the state of affairs represented by the function  $V$ .

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- For “ $V(\phi) = 1$ ” we also write “ $V \models \phi$ ” read as “ $V$  is a model of  $\phi$ ” or “ $V$  satisfies  $\phi$ .”
- If  $V$  doesn't satisfy  $\phi$  we write “ $V \not\models \phi$ ”. In other words  $V(\phi) = 0$ .

# Logical truth and logical falsity.

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- If a statement  $\phi$  is neither logically true or logically false then it is contingent.
- Examples:
  - ①  $(q \vee \neg q)$  is logically true.
  - ②  $(q \wedge \neg q)$  is logically false.

- A set of propositional logic statements is consistent if at least one state of affairs satisfies every statement in the set.

- A set of propositional logic statements is consistent if at least one state of affairs satisfies every statement in the set.
- A set of propositional logic statements is inconsistent if no state of affairs satisfies every statement in the set.

- A conclusion is *valid* with respect to a set of premises if the conclusion is true in every situation where the premises are true (van Benthem, et al, page 2-4).



- A conclusion is *valid* with respect to a set of premises if the conclusion is true in every situation where the premises are true (van Benthem, et al, page 2-4).
- One can validly infer a conclusion  $\phi$  from a set of premises  $P$  if the negation of  $\phi$  is inconsistent with the set of statements  $P$ .

(From van Benthem, et al., chapter 2)

- ① Computing a truth value for a formula takes linear time.

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- ① Computing a truth value for a formula takes linear time.
- ② Computing a truth table for validity takes exponential time.
- ③ The problem of testing for validity in propositional logic is decidable: there exists a mechanical method that computes the answer, at least in principle.