

Propositional Logic

Dave Dubin

Fall Semester, 2017

“Bene disserere est finis logices”

- Most logic books (including your assigned reading) introduce propositional logic as a tool or system for analyzing the validity and soundness of reasoning.

“Bene disserere est finis logices”

- Most logic books (including your assigned reading) introduce propositional logic as a tool or system for analyzing the validity and soundness of reasoning.
- van Benthem, et al. contrast arguments like this invalid one:

“Bene disserere est finis logices”

- Most logic books (including your assigned reading) introduce propositional logic as a tool or system for analyzing the validity and soundness of reasoning.
- van Benthem, et al. contrast arguments like this invalid one:

“Bene disserere est finis logices”

- Most logic books (including your assigned reading) introduce propositional logic as a tool or system for analyzing the validity and soundness of reasoning.
- van Benthem, et al. contrast arguments like this invalid one:

(2.4) If you take my medication, you will get better
 But you are not taking my medication.

 So, you will not get better. \therefore

“Bene disserere est finis logices”

- Most logic books (including your assigned reading) introduce propositional logic as a tool or system for analyzing the validity and soundness of reasoning.
- van Benthem, et al. contrast arguments like this invalid one:

(2.4) If you take my medication, you will get better
 But you are not taking my medication.

 So, you will not get better. \therefore

- ... with this valid argument:

(2.6) If you take my medication, you will get better
 But you are not getting better.

 So, you have not taken my medication. \therefore

Reasoning vs. representation

- In this class we'll give some attention to logical reasoning and inference.

Reasoning vs. representation

- In this class we'll give some attention to logical reasoning and inference.
- The same rules that govern valid argumentation also enable information systems to answer interesting questions.

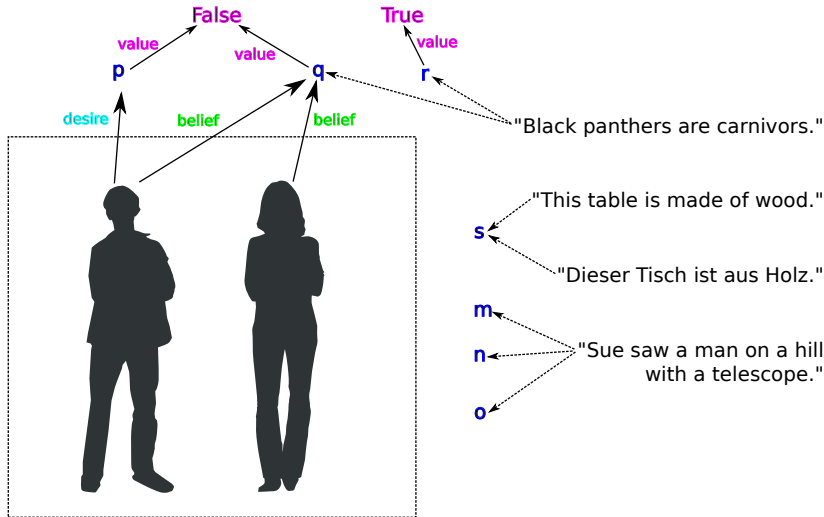
Reasoning vs. representation

- In this class we'll give some attention to logical reasoning and inference.
- The same rules that govern valid argumentation also enable information systems to answer interesting questions.
- But it takes an entire semester course in logic before you start to get good at logical inference.

Reasoning vs. representation

- In this class we'll give some attention to logical reasoning and inference.
- The same rules that govern valid argumentation also enable information systems to answer interesting questions.
- But it takes an entire semester course in logic before you start to get good at logical inference.
- This semester we'll focus mainly on logic as a means to describe a *domain* (some part of the world we're interested in).

Propositions, statements, and sentences



Propositions, statements, and sentences

- Propositions are the bearers of truth values: the kinds of things that can be true or false.

Propositions, statements, and sentences

- Propositions are the bearers of truth values: the kinds of things that can be true or false.
- States of affairs are the parts of reality responsible for making propositions true or false.

- We represent propositions with lower case letters (typically p , q , and r).

Propositional logic

- We represent propositions with lower case letters (typically p , q , and r).
- A set of proposition letters generates a set of states of affairs: ways the world might be.

- We represent propositions with lower case letters (typically p , q , and r).
- A set of proposition letters generates a set of states of affairs: ways the world might be.
- n proposition letters generate 2^n states of affairs.

- We represent propositions with lower case letters (typically p , q , and r).
- A set of proposition letters generates a set of states of affairs: ways the world might be.
- n proposition letters generate 2^n states of affairs.
- $\{pqr, pq\bar{r}, p\bar{q}r, p\bar{q}\bar{r}, \bar{p}qr, \bar{p}q\bar{r}, \bar{p}\bar{q}r, \bar{p}\bar{q}\bar{r}\}$

Table 1: 2.15 from van Benthem, et al.

Symbol	In natural language	Technical name
\neg	not	negation
\wedge	and	conjunction
\vee	or	disjunction
\rightarrow	if ... then	implication
\leftrightarrow	if and only if	equivalence

- A sentence constructed from proposition letters and operators is true or false in each state of affairs.

- A sentence constructed from proposition letters and operators is true or false in each state of affairs.
- Consider, for example: $((\neg p \vee q) \rightarrow r)$

- A sentence constructed from proposition letters and operators is true or false in each state of affairs.
- Consider, for example: $((\neg p \vee q) \rightarrow r)$
- The sentence is mapped to a truth value via the following tables

Semantics of the operators

φ	$\neg\varphi$
0	1
1	0

Table 3: 2.18 from van Benthem, et al.

φ	ψ	$\varphi \wedge \psi$	$\varphi \vee \psi$	$\varphi \rightarrow \psi$	$\varphi \leftrightarrow \psi$
0	0	0	0	1	1
0	1	0	1	1	0
1	0	0	1	0	0
1	1	1	1	1	1

Exercise 2.6

- “I will only go to school if I get a cookie now.”

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: $s =$ “I will go to school” and $c =$ “I get a cookie.”

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: $s =$ “I will go to school” and $c =$ “I get a cookie.”
- John and Mary are running

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: $s =$ “I will go to school” and $c =$ “I get a cookie.”
- John and Mary are running
- $(j \wedge m)$

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: s = “I will go to school” and c = “I get a cookie.”
- John and Mary are running
- $(j \wedge m)$
- A foreign national is entitled to social security if he has legal employment or if he has had such less than three years ago, unless he is currently also employed abroad.

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: s = “I will go to school” and c = “I get a cookie.”
- John and Mary are running
- $(j \wedge m)$
- A foreign national is entitled to social security if he has legal employment or if he has had such less than three years ago, unless he is currently also employed abroad.
- $((e \vee l) \wedge \neg a) \rightarrow s$

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: s = “I will go to school” and c = “I get a cookie.”
- John and Mary are running
- $(j \wedge m)$
- A foreign national is entitled to social security if he has legal employment or if he has had such less than three years ago, unless he is currently also employed abroad.
- $((e \vee l) \wedge \neg a) \rightarrow s)$
- e = “A foreign national has legal employment.”

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: s = “I will go to school” and c = “I get a cookie.”
- John and Mary are running
- $(j \wedge m)$
- A foreign national is entitled to social security if he has legal employment or if he has had such less than three years ago, unless he is currently also employed abroad.
- $((e \vee l) \wedge \neg a) \rightarrow s$
- e = “A foreign national has legal employment.”
- l = “He has had legal employment less than three years ago.”

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: s = “I will go to school” and c = “I get a cookie.”
- John and Mary are running
- $(j \wedge m)$
- A foreign national is entitled to social security if he has legal employment or if he has had such less than three years ago, unless he is currently also employed abroad.
- $((e \vee l) \wedge \neg a) \rightarrow s)$
- e = “A foreign national has legal employment.”
- l = “He has had legal employment less than three years ago.”
- a = “He is currently employed abroad.”

Exercise 2.6

- “I will only go to school if I get a cookie now.”
- Compare with “I will go to school if I get a cookie”
- The latter is cookie implies school, but the former is school implies cookie.
- You want $(s \rightarrow c)$
- Gloss: s = “I will go to school” and c = “I get a cookie.”
- John and Mary are running
- $(j \wedge m)$
- A foreign national is entitled to social security if he has legal employment or if he has had such less than three years ago, unless he is currently also employed abroad.
- $((e \vee l) \wedge \neg a) \rightarrow s$
- e = “A foreign national has legal employment.”
- l = “He has had legal employment less than three years ago.”
- a = “He is currently employed abroad.”
- s = “He is entitled to social security.”

Drawing truth tables for expressions

p	q	r	$((\neg p \vee q) \rightarrow r)$				
1	1	1	0	1	1	1	1
1	1	0	0	1	1	0	0
1	0	1	0	1	0	1	1
1	0	0	0	1	0	1	0
0	1	1	1	0	1	1	1
0	1	0	1	0	1	0	0
0	0	1	1	0	1	1	1
0	0	0	1	0	1	0	0

Drawing truth tables for expressions

p	q	r	$((\neg p \vee q) \rightarrow r)$				
1	1	1	0	1	1	1	1
1	1	0	0	1	1	0	0
1	0	1	0	1	0	0	1
1	0	0	0	1	0	0	1
0	1	1	1	0	1	1	1
0	1	0	1	0	1	0	0
0	0	1	1	0	1	0	1
0	0	0	1	0	1	0	0

Drawing truth tables for expressions

p	q	r	$((\neg p \vee q) \rightarrow r)$				
1	1	1	0	1	1	1	1
1	1	0	0	1	1	0	0
1	0	1	0	1	0	1	1
1	0	0	0	1	0	1	0
0	1	1	1	0	1	1	1
0	1	0	1	0	1	0	0
0	0	1	1	0	0	1	1
0	0	0	1	0	0	0	0

Drawing truth tables for expressions

p	q	r	$((\neg p \vee q) \rightarrow r)$				
1	1	1	0	1	1	1	1
1	1	0	0	1	1	0	0
1	0	1	0	0	0	1	1
1	0	0	0	0	0	1	0
0	1	1	1	0	1	1	1
0	1	0	1	0	1	0	0
0	0	1	1	0	0	1	1
0	0	0	1	0	0	0	0

Exercise 2.11

p	q	$((p \rightarrow q) \vee (q \rightarrow p))$						
1	1	1	1	1	1	1	1	1
1	0	1	0	0	1	0	1	1
0	1	0	1	1	1	1	0	0
0	0	0	1	0	1	0	1	0

Exercise 2.11

p	q	r	$(((p \vee \neg q) \wedge r) \leftrightarrow (\neg (p \wedge r) \vee q))$										
1	1	1	1	1	0	1	1	1	0	1	1	1	1
1	1	0	1	1	0	1	0	0	0	1	1	0	0
1	0	1	1	1	1	0	1	1	0	1	1	1	0
1	0	0	1	1	1	0	0	0	0	1	1	0	0
0	1	1	0	0	0	1	0	1	0	1	0	0	1
0	1	0	0	0	0	1	0	0	0	1	0	0	1
0	0	1	0	1	1	0	1	1	1	0	0	1	1
0	0	0	0	1	1	0	0	0	0	1	0	0	0

Grammar of propositional logic

Let P be a set of proposition letters and let $p \in P$.

The following expression defines the recursive grammar for a logical expression φ in Backus–Naur Form:

$$\varphi ::= p \mid \neg\varphi \mid (\varphi \wedge \varphi) \mid (\varphi \vee \varphi) \mid (\varphi \rightarrow \varphi) \mid (\varphi \leftrightarrow \varphi)$$

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

- $\neg \vee p$

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

- $\neg \vee p$
- $\vee)p\neg$

How many correct expressions are consistent with the last one?

Syntactically conforming expressions

Let $P = \{o, q, r, s\}$

Examples of grammatically conforming expressions include:

- r
- $\neg q$
- $(s \leftrightarrow o)$
- $(\neg(s \leftrightarrow \neg\neg\neg o) \rightarrow (q \wedge q))$

Grammatically *incorrect* expressions would include:

- $\neg \vee p$
- $\vee)p\neg$
- $\neg p \vee q \rightarrow r$

How many correct expressions are consistent with the last one?

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$
- $\neg((p \vee q) \rightarrow r)$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$
- $\neg((p \vee q) \rightarrow r)$
- $(\neg p \vee (q \rightarrow r))$

These conforming expressions are all consistent with $\neg p \vee q \rightarrow r$

- $((\neg p \vee q) \rightarrow r)$
- $(\neg(p \vee q) \rightarrow r)$
- $\neg((p \vee q) \rightarrow r)$
- $(\neg p \vee (q \rightarrow r))$
- $\neg(p \vee (q \rightarrow r))$

Exercise 2.7

Which of the following are formulas in propositional logic?

- $p \rightarrow \neg q$

Exercise 2.7

Which of the following are formulas in propositional logic?

- $p \rightarrow \neg q$
- $\neg\neg \wedge q \vee p$

Exercise 2.7

Which of the following are formulas in propositional logic?

- $p \rightarrow \neg q$
- $\neg\neg \wedge q \vee p$
- $p\neg q$

- Semantics is the relationship of a language to the part of the world that we're modeling.

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.
- “ $V(\varphi) = 1$ ” means the formula (or sentence) φ is true in the state of affairs represented by the function V . “ $V(\varphi) = 0$ ” means that φ is false in the state of affairs represented by the function V .

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.
- “ $V(\varphi) = 1$ ” means the formula (or sentence) φ is true in the state of affairs represented by the function V . “ $V(\varphi) = 0$ ” means that φ is false in the state of affairs represented by the function V .
- For “ $V(\varphi) = 1$ ” we also write “ $V \models \varphi$ ” read as “ V is a model of φ ” or “ V satisfies φ .”

- Semantics is the relationship of a language to the part of the world that we're modeling.
- Valuations are functions from expressions to truth values.
- “ $V(\varphi) = 1$ ” means the formula (or sentence) φ is true in the state of affairs represented by the function V . “ $V(\varphi) = 0$ ” means that φ is false in the state of affairs represented by the function V .
- For “ $V(\varphi) = 1$ ” we also write “ $V \models \varphi$ ” read as “ V is a model of φ ” or “ V satisfies φ .”
- If V doesn't satisfy φ we write “ $V \not\models \varphi$ ”. In other words $V(\varphi) = 0$.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.
- Examples:

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.
- Examples:
 - ① $(q \vee \neg q)$ is logically true.

Logical truth and logical falsity.

- A statement φ is logically true if it is true in every state of affairs generated by its propositional variables.
- A statement φ is logically false if it is false in every state of affairs generated by its propositional variables.
- If a statement φ is neither logically true or logically false then it is contingent.
- Examples:
 - ① $(q \vee \neg q)$ is logically true.
 - ② $(q \wedge \neg q)$ is logically false.

- A set of propositional logic statements is consistent if at least one state of affairs satisfies every statement in the set.

- A set of propositional logic statements is consistent if at least one state of affairs satisfies every statement in the set.
- A set of propositional logic statements is inconsistent if no state of affairs satisfies every statement in the set.

- A conclusion is *valid* with respect to a set of premises if the conclusion is true in every situation where the premises are true (van Benthem, et al, page 2-4).

- A conclusion is *valid* with respect to a set of premises if the conclusion is true in every situation where the premises are true (van Benthem, et al, page 2-4).
- One can validly infer a conclusion φ from a set of premises P if the negation of φ is inconsistent with the set of statements P .

(From van Benthem, et al., chapter 2)

- 1 Computing a truth value for a formula takes linear time.

(From van Benthem, et al., chapter 2)

- ① Computing a truth value for a formula takes linear time.
- ② Computing a truth table for validity takes exponential time.

(From van Benthem, et al., chapter 2)

- ① Computing a truth value for a formula takes linear time.
- ② Computing a truth table for validity takes exponential time.
- ③ The problem of testing for validity in propositional logic is decidable: there exists a mechanical method that computes the answer, at least in principle.