Assumption Tracking for Optimistic Optimizations

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```
for (k = 1; k \le M; k++) {
  mc[k] = mpp[k - 1] + tpmm[k - 1];
  if ((sc = ip[k - 1] + tpim[k - 1]) > mc[k]) mc[k] = sc;
  if ((sc = dpp[k - 1] + tpdm[k - 1]) > mc[k]) mc[k] = sc;
  if ((sc = xmb + bp[k]) > mc[k]) mc[k] = sc;
  mc[k] += ms[k];
  if (mc[k] < -INFTY) mc[k] = -INFTY;</pre>
  dc[k] = dc[k - 1] + tpdd[k - 1];
  if ((sc = mc[k - 1] + tpmd[k - 1]) > dc[k]) dc[k] = sc;
  if (dc[k] < -INFTY) dc[k] = -INFTY;
  if (k < M) {
    ic[k] = mpp[k] + tpmi[k];
    if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
    ic[k] += is[k]:
    if (ic[k] < -INFTY) ic[k] = -INFTY;</pre>
```



```
#pragma clang loop vectorize(enable)
for (k = 1; k \le M; k++) {
  mc[k] = mpp[k - 1] + tpmm[k - 1];
  if ((sc = ip[k - 1] + tpim[k - 1]) > mc[k]) mc[k] = sc;
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```

+ up to 30% speedup



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  if (k < M) {
    ic[k] = mpp[k] + tpmi[k];
    if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
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    if ((sc = ip[k] + tpii[k]) > ic[k]) ic[k] = sc;
    ic[k] += is[k]:
    if (ic[k] < -INFTY) ic[k] = -INFTY;</pre>
```

+ up to 50% speedup





1 vectorized loop \implies + up to 30% speedup



```
1 vectorized loop \implies + up to 30% speedup 2 vectorized loops \implies + up to 50% speedup
```



```
1 vectorized loop \implies + up to 30% speedup 2 vectorized loops \implies + up to 50% speedup possible aliasing \implies - runtime alias checks
```



```
1 vectorized loop ⇒ + up to 30% speedup
2 vectorized loops ⇒ + up to 50% speedup
possible aliasing ⇒ - runtime alias checks
possible dependences ⇒ - static dependence analysis
```







```
float BlkSchlsEqEuroNoDiv(float sptprice, float strike, float rate,
                          float volatility, float time, int otype) {
    float xD1, xD2, xDen, d1, d2, FutureValueX, NofXd1, NofXd2, NegNofXd1,
          NegNofXd2, Price;
    xD1 = rate + volatility * volatility; * 0.5;
    xD1 = xD1 * time:
    xD1 = xD1 + \log(\text{ sptprice / strike });
    xDen = volatility * sqrt(time);
    xD1 = xD1 / xDen:
    xD2 = xD1 - xDen:
    d1 = xD1:
    d2 = xD2;
    NofXd1 = CNDF(d1):
    NofXd2 = CNDF(d2):
    FutureValueX = strike * ( exp( -(rate)*(time) ) );
    if (otype == 0) {
        Price = (sptprice * NofXd1) - (FutureValueX * NofXd2):
    } else {
        NegNofXd1 = (1.0 - NofXd1):
        NegNofXd2 = (1.0 - NofXd2);
        Price = (FutureValueX * NegNofXd2) - (sptprice * NegNofXd1);
    return Price;
}
```





+ 2.9× speedup for manual parallelization on a guad-core i7



- + 2.9× speedup for manual parallelization on a quad-core i7
- + 2.8× speedup for automatic parallelization on a guad-core i7



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- + 2.8× speedup for automatic parallelization on a quad-core i7
- Possible aliasing arrays



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- Possible execution of non-pure calls



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- Possible execution of dead-iterations (0 <= j < NUM_RUNS 1)



- + 2.9× speedup for manual parallelization on a quad-core i7
- + 2.8× speedup for automatic parallelization on a quad-core i7
- + 6.5× speedup for sequential execution (native input)
- Possible aliasing arrays
- Possible execution of non-pure calls
- Possible execution of dead-iterations (0 <= j < NUM_RUNS 1)





```
void compute_rhs() {
  int i, j, k, m;
  double rho_inv, uijk, up1, um1, vijk, vp1, vm1, wijk, wp1, wm1;
  if (timeron) timer start(t rhs);
  for (k = 0; k \le grid_points[2]-1; k++) {
    for (j = 0; j <= grid_points[1]-1; j++) {
      for (i = 0; i <= grid_points[0]-1; i++) {
        rho_{inv} = 1.0/u[k][j][i][0];
        rho_i[k][j][i] = rho_inv;
        us[k][j][i] = u[k][j][i][1] * rho_inv;
        vs[k][j][i] = u[k][j][i][2] * rho_inv;
        ws[k][j][i] = u[k][j][i][3] * rho_inv;
        square[k][j][i] = 0.5* (
            u[k][j][i][1]*u[k][j][i][1] +
            u[k][i][i][2]*u[k][i][i][2] +
            u[k][j][i][3]*u[k][j][i][3] ) * rho_inv;
        qs[k][j][i] = square[k][j][i] * rho_inv;
    }
```



```
for (k = 0; k \le grid_points[2]-1; k++) {
  for (j = 0; j <= grid_points[1]-1; j++) {</pre>
   for (i = 0; i <= grid_points[0]-1; i++) {</pre>
      for (m = 0; m < 5; m++) {
        rhs[k][j][i][m] = forcing[k][j][i][m];
if (timeron) timer_start(t_rhsx);
for (k = 1; k \le grid_points[2]-2; k++) {
  for (j = 1; j <= grid_points[1]-2; j++) {</pre>
    for (i = 1; i <= grid_points[0]-2; i++) {</pre>
      uijk = us[k][j][i];
      up1 = us[k][j][i+1];
      um1 = us[k][j][i-1];
      rhs[k][j][i][0] = rhs[k][j][i][0] + dx1tx1 *
        (u[k][j][i+1][0] - 2.0*u[k][j][i][0] +
          u[k][j][i-1][0]) -
        tx2 * (u[k][j][i+1][1] - u[k][j][i-1][1]);
```



```
rhs[k][j][i][1] = rhs[k][j][i][1] + dx2tx1 *
  (u[k][j][i+1][1] - 2.0*u[k][j][i][1] +
  u[k][j][i-1][1]) +
  xxcon2*con43 * (up1 - 2.0*uijk + um1) -
  tx2 * (u[k][j][i+1][1]*up1 -
      u[k][j][i-1][1]*um1 +
      (u[k][j][i+1][4] - square[k][j][i+1]-
       u[k][j][i-1][4] + square[k][j][i-1]) * c2);
rhs[k][j][i][2] = rhs[k][j][i][2] + dx3tx1 *
  (u[k][i][i+1][2] - 2.0*u[k][i][i][2] +
  u[k][j][i-1][2]) +
  xxcon2 * (vs[k][j][i+1] - 2.0*vs[k][j][i] +
      vs[k][i][i-1]) -
  tx2 * (u[k][j][i+1][2]*up1 - u[k][j][i-1][2]*um1);
rhs[k][j][i][3] = rhs[k][j][i][3] + dx4tx1 *
  (u[k][j][i+1][3] - 2.0*u[k][j][i][3] +
  u[k][i][i-1][3]) +
  xxcon2 * (ws[k][j][i+1] - 2.0*ws[k][j][i] +
      ws[k][j][i-1]) -
  tx2 * (u[k][j][i+1][3]*up1 - u[k][j][i-1][3]*um1);
  /* \approx 300 more lines of similar code */
```



```
for (k = 0; k <= grid_points[2]-1; k++)
  for (j = 0; j <= grid_points[1]-1; j++)
    for (i = 0; i <= grid_points[0]-1; i++)
        for (m = 0; m < 5; m++)
            rhs[k][j][i][m] = forcing[k][j][i][m];

if (timeron) timer_start(t_rhsx);

for (k = 1; k <= grid_points[2]-2; k++) {
    for (j = 1; j <= grid_points[1]-2; j++) {
        for (i = 1; i <= grid_points[0]-2; i++) {
            /* ... */</pre>
```

^aSanvam and Yew. PLDI 15



```
for (k = 0; k <= grid_points[2]-1; k++)
  for (j = 0; j <= grid_points[1]-1; j++)
    for (i = 0; i <= grid_points[0]-1; i++)
        for (m = 0; m < 5; m++)
            rhs[k][j][i][m] = forcing[k][j][i][m];

if (timeron) timer_start(t_rhsx);

for (k = 1; k <= grid_points[2]-2; k++) {
    for (j = 1; j <= grid_points[1]-2; j++) {
        for (i = 1; i <= grid_points[0]-2; i++) {
            /* ... */</pre>
```

+ 6× speedup for 8 threads/cores ^a

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```
for (k = 0; k <= grid_points[2]-1; k++)
  for (j = 0; j <= grid_points[1]-1; j++)
    for (i = 0; i <= grid_points[0]-1; i++)
        for (m = 0; m < 5; m++)
            rhs[k][j][i][m] = forcing[k][j][i][m];

if (timeron) timer_start(t_rhsx);

for (k = 1; k <= grid_points[2]-2; k++) {
    for (j = 1; j <= grid_points[1]-2; j++) {
        for (i = 1; i <= grid_points[0]-2; i++) {
            /* ... */</pre>
```

- 6× speedup for 8 threads/cores ^a
- Possible variant loop bounds

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```
for (k = 0; k <= grid_points[2]-1; k++)
  for (j = 0; j <= grid_points[1]-1; j++)
    for (i = 0; i <= grid_points[0]-1; i++)
        for (m = 0; m < 5; m++)
            rhs[k][j][i][m] = forcing[k][j][i][m];

if (timeron) timer_start(t_rhsx);

for (k = 1; k <= grid_points[2]-2; k++) {
    for (j = 1; j <= grid_points[1]-2; j++) {
        for (i = 1; i <= grid_points[0]-2; i++) {
            /* ... */</pre>
```

- + 6× speedup for 8 threads/cores ^a
- Possible variant loop bounds
- Possible out-of-bound accesses

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```
for (k = 0; k <= grid_points[2]-1; k++)
  for (j = 0; j <= grid_points[1]-1; j++)
    for (i = 0; i <= grid_points[0]-1; i++)
        for (m = 0; m < 5; m++)
            rhs[k][j][i][m] = forcing[k][j][i][m];

if (timeron) timer_start(t_rhsx);

for (k = 1; k <= grid_points[2]-2; k++) {
    for (j = 1; j <= grid_points[1]-2; j++) {
        for (i = 1; i <= grid_points[0]-2; i++) {
            /* ... */</pre>
```

- + 6× speedup for 8 threads/cores ^a
- Possible variant loop bounds
- Possible out-of-bound accesses
- Possible execution of non-pure calls

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```
for (k = 0; k <= grid_points[2]-1; k++)
  for (j = 0; j <= grid_points[1]-1; j++)
    for (i = 0; i <= grid_points[0]-1; i++)
        for (m = 0; m < 5; m++)
            rhs[k][j][i][m] = forcing[k][j][i][m];

if (timeron) timer_start(t_rhsx);

for (k = 1; k <= grid_points[2]-2; k++) {
    for (j = 1; j <= grid_points[1]-2; j++) {
        for (i = 1; i <= grid_points[0]-2; i++) {
            /* ... */</pre>
```

- + 6× speedup for 8 threads/cores ^a
- Possible variant loop bounds
- Possible out-of-bound accesses
- Possible execution of non-pure calls
- Possible integer under/overflows complicate loop bounds

^aSanvam and Yew. PLDI 15





Be Optimistic!

Programs might be nasty but programmers are not.

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Optimistic Assumptions & Speculative Versioning



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Optimistic Assumptions

1 Take optimistic assumptions to (better) optimize loops



Optimistic Assumptions & Speculative Versioning

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- Take optimistic assumptions to (better) optimize loops
- Derive simple runtime conditions that imply these assumptions



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- Take optimistic assumptions to (better) optimize loops
- Derive simple runtime conditions that imply these assumptions
- 3 Version the code based on the assumptions made and conditions derived.

When static information is insufficient

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Optimistic Assumptions & Speculative Versioning

Optimistic Assumptions

- Take optimistic assumptions to (better) optimize loops
- Derive simple runtime conditions that imply these assumptions
- 3 Version the code based on the assumptions made and conditions derived.

Speculative Versioning

```
if (/* Runtime Conditions */)
  /* Optimized Loop Nest */
else
  /* Original Loop Nest */
```

When static information is insufficient

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Runtime Conditions

Runtime Conditions

- Fast to derive (compile time)
- Fast to verify (runtime)
- High probability to be true

The polyhedral loop optimizer in LLVM



The polyhedral loop optimizer in LLVM



Polyhedral Optimizer

- Loop Nest Optimizer
- Precise compute model (affine constraints)
- Combines many classical loop optimizations
 - ► Tiling, Interchange, Fusion, Fission, ...
- Examples: Pluto, ppcg

The polyhedral loop optimizer in LLVM



Polly — Advantages

- Automatic & Semi-automatic Mode
- Robust & Widely Applicable
- Embedded in LLVM
 - Source & Target Independent
 - Information flow between Polly and other passes



Possibly Invariant Loads

```
void loop_bounds(int *size0, int *size1) {
   for (int i = 0; i < *size0; i++)
     for (int j = 0; j < *size1; j++)
        ...
}</pre>
```

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Possibly Invariant Loads

```
void loop_bounds(int *size0, int *size1) {
  int size0val = *size0;
  int size1val = *size1;

for (int i = 0; i < size0val; i++)
  for (int j = 0; j < size1val; j++)
    ...
}</pre>
```

Hoist invariant loads

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Possibly Invariant Loads

```
void loop_bounds(int *size0, int *size1) {
  int size0val = *size0;
  int size1val = 0;

  if (size0val > 0)
     size1val = *size1;

  for (int i = 0; i < size0val; i++)
     for (int j = 0; j < size1val; j++)
        ...
}</pre>
```

- Hoist invariant loads
- Keep conditions for conditionally executed loads

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Possibly Invariant Loads

```
void loop_bounds(int *size0, int *size1) {
  int size0val = *size0;
  int size1val = 0;

  if (size0val > 0)
     size1val = *size1;

  for (int i = 0; i < size0val; i++)
     for (int j = 0; j < size1val; j++)
     ...
}</pre>
```

- Hoist invariant loads
- Keep conditions for conditionally executed loads
- Powerful in combination with alias checks





	Today in LLVM	Polly
Target	Inner Loops	(Unstructured) Loops Nests



	Today in LLVM	Polly
Target	Inner Loops	(Unstructured) Loops Nests
Approach	Seq. of Specialized Passes	Combined Modeling



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Analysis	Heuristic-based Dependences Instruction Cost Model	Precise Flow-Dependences Optimistic Assumptions Memory Access Analysis



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- (A) Applicability/Correctness
 - No Alias Assumption¹
 - No Wrapping Assumption²
 - 3 Finite Loops Assumption²
 - 4 Array In-bounds Assumption²
 - 5 Valid Multidimensional View Assumption (Delinearization)³
 - 6 Possibly Invariant Loads
- (B) Optimizations
 - Array In-bounds Check Hoisting²
 - Parametric Dependence Distances⁴

¹Joint work Fabrice Rastello (INRIA Grenoble) & others. [OOPSLA'15]

²Joint work with Tobias Grosser (ETH)

³Tobias Grosser & Sebastian Pop (Samsung) [ICS'15]

⁴Joint work with Zino Benaissa (Qualcomm)



■ Polly Mainline

- Improved optimization choices
- Profile guided optimization
- More powerful checks (e.g., generate inspector loops)
- ▶ User feedback, register tiling, user provided assumptions
- ▶ ...



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- ▶ ..
- Integrate Polly into LLVM mainline



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Independent Projects using Polly

- Heterogeneous Compute (OpenCL)
- High-level Synthesis
- Dynamic Compilation (JITs)



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Independent Projects using Polly

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Thank You.

Incomplete feature list



Loops & Conditions

Styles

for/while/do/goto

Arbitrary Presburger Expressions

```
for (i = 0; i < 22 && i < mod(i + b, 13); i += 2)
if (5 * i + b <= 13 || 12 > b)
```

- Multiple back-edges/exit-edges & unstructured control flow break; continue; goto; switch
- Data-dependent control flow

```
if (B[i]) A[i] = A[i] / B[i];
```

Accesses

- Multi-dimensionality: A[] [n] [m] / A[] [10] [100]
- ► Non-affine: A[i * j]
- Scalar: x = A[i]; ...; B[i] = x;

User Annotations

- __builtin_assume(M > 8)
- __builtin_assume(__polly_profitable == 1)



```
rhs.c:47:3: remark: SCoP begins here. [-Rpass-analysis=polly-scops]
  for (k = 0; k <= grid_points[2]-1; k++) {
    /* ... */
    rhs.c:418:16: remark: SCoP ends here. [-Rpass-analysis=polly-scops]
    if (timeron) timer_stop(t_rhs);</pre>
```







```
rhs.c:419:1: remark: No-overflows assumption: [grid_points, grid_points', grid_points'', timeron] -> {: (grid_points >= 3 and grid_points' >= 3 and grid_points' >= 3 and grid_points' >= 3 and grid_points' >= -2147483643 or (grid_points >= 3 and grid_points' >= -2147483643 and grid_points'' >= -2147483646) or (grid_points <= 2 and grid_points >= -2147483643 and grid_points' >= 3 and grid_points'' >= -2147483646) or (grid_points <= 2 and grid_points >= -2147483644 and grid_points' <= 2 and grid_points' >= -2147483646) or (grid_points' = -2147483644 and grid_points >= 3 and grid_points'' <= 2 and grid_points'' >= -2147483646) or (grid_points = -2147483644 and grid_points' >= 3 and grid_points'' <= 2 and grid_points'' >= -2147483646) }
```

```
_builtin_assume(grid_points[0] >= -2147483643 && grid_points[1] >= -2147483643 && grid_points[2] >= -2147483643);
```



```
rhs.c:50:23: remark: Possibly aliasing pointer, use restrict keyword.
    [-Rpass-analysis=polly-scops]
        rho_inv = 1.0/u[k][j][i][0];

rhs.c:56:13: remark: Possibly aliasing pointer, use restrict keyword.
    [-Rpass-analysis=polly-scops]
        u[k][j][i][1]*u[k][j][i][1] +
```



Array In-bounds Assumptions

```
for (int i = 0; i < N; i++)
    for (int j = 0; j < M; j++)
        A[j][i] += A[2*j+1][i];
}</pre>
```



Array In-bounds Assumptions



Array In-bounds Assumptions

Out-of-bound accesses introduce multiple addresses for one memory location (e.g., &A[1][0] == &A[0][128])





```
void mem_copy(int N, float *A, float *B) {
   if (&A[0] >= &B[N+5] ||

    #pramga vectorize
   for (i = 0; i < N; i++)
        A[i] = B[i+5];
} else {
    /* original code */
}
</pre>
```



No Alias Assumptions

```
void mem_copy(int N, float *A, float *B) {
   if (&A[0] >= &B[N+5] || &A[N] <= &B[5]) {
     #pramga vectorize
     for (i = 0; i < N; i++)
        A[i] = B[i+5];
} else {
        /* original code */
}
</pre>
```

Compare minimal/maximal accesses to possible aliasing arrays



No Alias Assumptions

```
void evn_odd(int N, int *Evn, int *Odd, int *A, int *B) {
  if
    for (int i = 0: i < N: i += 2)
      if (N % 2)
        Odd[i/2] = A[i+1] - B[i+1];
      else
        Evn[i/2] = A[i] + B[i];
  } else {
    /* original code */
```

Compare minimal/maximal accesses to possible aliasing arrays



No Alias Assumptions

```
void evn_odd(int N, int *Evn, int *Odd, int *A, int *B) {
  if
    for (int i = 0: i < N: i += 2)
      if (N % 2)
        Odd[i/2] = A[i+1] - B[i+1];
      else
        Evn[i/2] = A[i] + B[i];
  } else {
    /* original code */
```

- Compare minimal/maximal accesses to possible aliasing arrays
- Do not compare accesses to read-only arrays



No Alias Assumptions

```
void evn_odd(int N, int *Evn, int *Odd, int *A, int *B) {
  if
    for (int i = 0: i < N: i += 2)
      if (N % 2)
        Odd[i/2] = A[i+1] - B[i+1];
      else
        Evn[i/2] = A[i] + B[i];
  } else {
    /* original code */
```

- Compare minimal/maximal accesses to possible aliasing arrays
- Do not compare accesses to read-only arrays
- Use the iteration domain of the accesses



No Alias Assumptions

```
void evn odd(int N. int *Evn. int *Odd. int *A. int *B) {
  if (N\%2 ? ((\&B[N+1] \le \&Odd[0] | | \&Odd[(N+1)/2] \le \&B[1]) \&\&
                 (&A[N+1] \le &Odd[0] | &Odd[(N+1)/2] \le &A[1])
          : ((\&B[N] \le \&Evn[0] | | \&Evn[(N+1)/2] \le \&B[0]) \&\&
                 (&A[N] \le &Evn[O] \mid &Evn[(N+1)/2] \le &A[O]))) {
    for (int i = 0: i < N: i += 2)
      if (N % 2)
        Odd[i/2] = A[i+1] - B[i+1];
      else
        Evn[i/2] = A[i] + B[i];
  } else {
    /* original code */
```

- Compare minimal/maximal accesses to possible aliasing arrays
- Do not compare accesses to read-only arrays
- Use the iteration domain of the accesses





```
void mem_shift(unsigned char N, float *A) {
  if (N <= 128) {

    #pramga vectorize
    for (unsigned char i = 0; i < N; i++)
        A[i] = A[N + i];

} else {
    /* original code */
}
}</pre>
```

- Finite bit width can cause integer expressions to "wrap around"
- Wrapping causes multiple addresses for one memory location



$$\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1} \equiv_{p} (\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1}) \mod 2^k$$



$$\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1} \equiv_p (\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1}) \mod 2^k$$

$$\underline{i} \in [0, N-1] \land N \in [0, 2^8]$$

$$(N + \underline{i}) \equiv_p (N + \underline{i}) \mod 2^8$$

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$$\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1} \equiv_{p} (\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1}) \mod 2^{k}$$

$$i \in [0, N-1] \land N \in [0, 2^{8}]$$

$$(N + i) \equiv_{p} (N + i) \mod 2^{8}$$

$$\Longrightarrow (N + i) \leq_{p} 255$$



$$\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1} \equiv_{p} (\underline{i} * \underline{c_0} + \underline{p} * \underline{c_1}) \mod 2^k$$

$$i \in [0, N-1] \land N \in [0, 2^8]$$

$$(N + i) \equiv_{p} (N + i) \mod 2^8$$

$$\Longrightarrow (N + i) \leq_{p} 255$$

$$\Longrightarrow N \leq 128$$





Finite Loops Assumption

```
void mem_shift(unsigned N, float *A) {
  if (N % 2 == 0) {
    #pramga vectorize
    for (unsigned i = 0; i != N; i+=2)
        A[i+4] = A[i];
} else {
    /* original code */
}
```

- Allows to provide other LLVM passes real loop bounds
- Infinite loops create unbounded optimization problems







Valid Multidimensional View Assumption

- Define multi-dimensional view of a linearized (one-dimensional) array
- Derive conditions that accesses are in-bounds for each dimension





```
struct SafeArray { int Size, int *Array };
inline void set(SafeArray A, int idx, int val) {
  if (idx < 0 || A.Size <= idx)
      throw OutOfBounds;
  A.Array[idx] = val;
}

void set_safe_array(int N, SafeArray A) {
  for (int i = 0; i < N; i++)
      for (int j = 0; j < i/2; j++)
      set(A, i+j, 1); /* Throws out-of-bounds */
}</pre>
```



Array In-bounds Check Hoisting



Array In-bounds Check Hoisting



Array In-bounds Check Hoisting

Hoist in-bounds access conditions out of the loop nest



Check Hoisting

```
void copy(int N, double A[N][N], double B[N][N]) {
    if (DebugLevel <= 5) {</pre>
      #pragma parallel
      for (int i = 0; i < N; i++)
        #pragma simd
        for (int j = 0; j < N; j++)
          A[i][i] = B[i][i];
    } else {
      for (int i = 0; i < N; i++) {
        for (int j = 0; j < N; j++)
          A[i][i] = B[i][i];
        if (DebugLevel > 5)
          printf("Columnu\%ducopied\n", i)
      }
```





Parametric Dependence Distances

```
void vectorize(int N, double *A) {
   if (c >= 4) {

    #pragma vectorize width(4)
   for (int i = c; i < N+c; i++)
        A[i-c] += A[i];

} else {
    /* original code */
}
</pre>
```

Assume large enough dependence distance, e.g., for vectorization





- Modern off-the-shelf processors are complex and powerful
 - low overhead vector units
 - multiple cores and levels of cache



- Modern off-the-shelf processors are complex and powerful
 - low overhead vector units
 - multiple cores and levels of cache
- Programs do not exploit this power
 - not cache aware
 - written for single threaded execution



```
for (i = 0; i < nx; i++) {
  for (j = 0; j < ny; j++) {
    q[i] = q[i] + A[i][j] * p[j];
    s[j] = s[j] + r[i] * A[i][j];
}</pre>
```



```
if ((nv >= 1)) {
 ub1 = floord((nx + -1), 256);
#pragma omp parallel for private(c2, c3, c4, c5, c6) firstprivate(ub1)
 for (c1 = 0; c1 \le ub1; c1++)
   for (c2 = 0; c2 \le floord((nv + -1), 256); c2++)
     for (c3 = (8 * c1); c3 \le min(floord((nx + -1), 32), ((8 * c1) + 7)); c3++)
        for (c4 = (8 * c2); c4 \le min(floord((nv + -1), 32), ((8 * c2) + 7)); c4++)
          for (c5 = (32 * c4); c5 \le min(((32 * c4) + 31), (nv + -1)); c5++)
#pragma ivdep
#pragma vector always
#praama simd
            for (c6 = (32 * c3); c6 \le min(((32 * c3) + 31), (nx + -1)); c6++)
              q[c6]=q[c6]+A[c6][c5]*p[c5];
}
if ((nx >= 1)) {
 ub1 = floord((ny + -1), 256);
#pragma omp parallel for private(c2, c3, c4, c5, c6) firstprivate(ub1)
 for (c1 = 0; c1 <= ub1; c1++)
   for (c2 = 0; c2 \le floord((nx + -1), 256); c2++)
      for (c3 = (8 * c1); c3 \le min(floord((ny + -1), 32), ((8 * c1) + 7)); c3++)
        for (c4 = (8 * c2); c4 \le min(floord((nx + -1), 32), ((8 * c2) + 7)); c4++)
         for (c5 = (32 * c4); c5 \le min(((32 * c4) + 31), (nx + -1)); c5++)
#pragma ivdep
#pragma vector always
#pragma simd
            for (c6 = (32 * c3); c6 \le min(((32 * c3) + 31), (ny + -1)); c6++)
              s[c6]=s[c6]+r[c5]*A[c5][c6]:
}
```



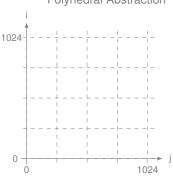
Input Loop Nest

```
for (i = 0; i <= 1024; i++) {
  for (j = 0; j <= 1024; j++) {
    s[j] = s[j] + r[i] * A[i][j];
    q[i] = q[i] + A[i][j] * p[j];
  }
}</pre>
```



Input Loop Nest

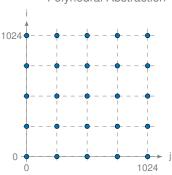
```
for (i = 0; i <= 1024; i++) {
  for (j = 0; j <= 1024; j++) {
    s[j] = s[j] + r[i] * A[i][j];
    q[i] = q[i] + A[i][j] * p[j];
}</pre>
```





Input Loop Nest

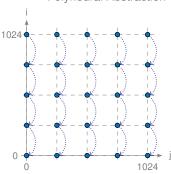
```
for (i = 0; i <= 1024; i++) {
  for (j = 0; j <= 1024; j++) {
    s[j] = s[j] + r[i] * A[i][j];
    q[i] = q[i] + A[i][j] * p[j];
}</pre>
```





Input Loop Nest

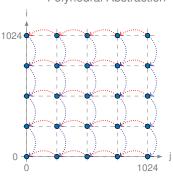
```
for (i = 0; i <= 1024; i++) {
  for (j = 0; j <= 1024; j++) {
    s[j] = s[j] + r[i] * A[i][j];
    q[i] = q[i] + A[i][j] * p[j];
}</pre>
```



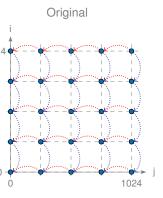


Input Loop Nest

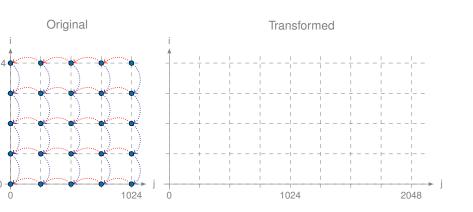
```
for (i = 0; i <= 1024; i++) {
  for (j = 0; j <= 1024; j++) {
    s[j] = s[j] + r[i] * A[i][j];
    q[i] = q[i] + A[i][j] * p[j];
}</pre>
```



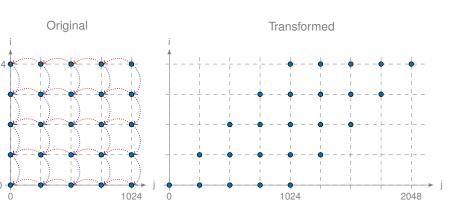




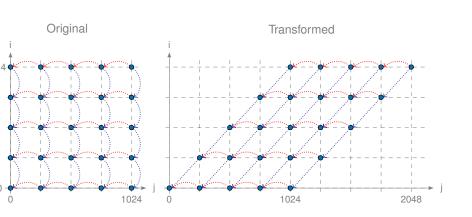




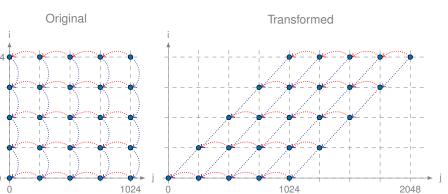




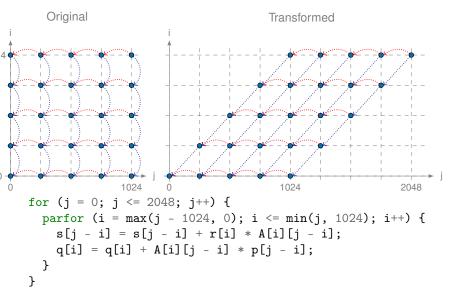














```
for (i = 0; i <= 1024; i++) {
  for (j = 0; j <= 1024; j++) {
    s[j] = s[j] + r[i] * A[i][j];
    q[i] = q[i] + A[i][j] * p[j];
}</pre>
```

Input Loop Nest



```
for (i = 0; i <= 1024; i++) {
  for (j = 0; j <= 1024; j++) {
    s[j] = s[j] + r[i] * A[i][j];
    q[i] = q[i] + A[i][j] * p[j];
  }
}</pre>
```

Static Analysis

Input Loop Nest





```
for (i = 0; i <= 1024; i++) { for (j = 0; j <= 1024; j++) { s(j) = s(j) + r(i) * A(i)(j); q(i) = q(i) + A(i)(j) * p(j); } } } Input Loop Nest
```

Static Analysis

Polyhedral Abstraction Scheduler Schedule



```
for (i = 0; i <= 1024; i++) {
                                                               for (j = 0; j <= 2048; j++) {
      for (j = 0; j <= 1024; j++) {
                                                                parfor (i = max(j - 1024, 0); i <= min(j, 1024); i++) {
        s[j] = s[j] + r[i] * A[i][j];
                                                                  s[j-i] = s[j-i] + r[i] * A[i][j-i];
        q[i] = q[i] + A[i][i] * p[i];
                                                                  q[i] = q[i] + A[i][i - i] * p[i - i];
   Input Loop Nest
                                                                  Output Loop Nest
               Static Analysis
                                                                               Code Generation
                                           Scheduler
Polyhedral Abstraction
                                                                       Schedule
```