

CS170–Spring 2019 — Homework 2 Solutions

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1 Study Group

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2 Asymptotic Complexity Comparisons

(a) $3 \leq 7 \leq 2 \leq 5 \leq 4 \leq 9 \leq 8 \leq 6 \leq 1$

(b) (i) $\log_3 n = \Theta(\log_4 n)$

Proof:

$$\log_3 n = \log_{(4^{\log_4 3})} n = \left(\frac{1}{\log_4 3}\right) \log_4 n = \Theta(\log_4 n)$$

(ii) $n \log(n^4) = O(n^2 \log(n^3))$

Proof:

$$\lim_{n \rightarrow +\infty} \frac{n^2 \log(n^3)}{n \log(n^4)} = \lim_{n \rightarrow +\infty} \frac{3n^2 \log n}{4n \log n} = \lim_{n \rightarrow +\infty} \frac{3}{4} n = \infty$$

(iii) $\sqrt{n} = \Omega((\log n)^3)$

Proof: (Use L'Hôpital's rule)

$$\begin{aligned} \lim_{n \rightarrow +\infty} \frac{\sqrt{n}}{(\log n)^3} &= \lim_{n \rightarrow +\infty} \frac{\frac{1}{2\sqrt{n}}}{3(\log n)^2 \frac{1}{n}} = \lim_{n \rightarrow +\infty} \frac{\sqrt{n}}{(\log n)^2} \\ &= \lim_{n \rightarrow +\infty} \frac{\frac{1}{2\sqrt{n}}}{2(\log n) \frac{1}{n}} = \lim_{n \rightarrow +\infty} \frac{\sqrt{n}}{\log n} \\ &= \lim_{n \rightarrow +\infty} \frac{\frac{1}{2\sqrt{n}}}{\frac{1}{n}} = \lim_{n \rightarrow +\infty} \sqrt{n} = \infty \end{aligned}$$

(iv) $2^n = \Theta(2^{n+1})$

Proof:

$$2^{n+1} = 2 * 2^n = \Theta(2^n)$$

(v) $n = \Omega((\log n)^{\log \log n})$

Proof: (Use L'Hôpital's rule)

Take the logarithm of both functions

$$\log(\log n)^{\log \log n} = (\log \log n)^2$$

Therefore, it is equivalent to compare $(\log \log n)^2$ and $\log n$

$$\begin{aligned} \lim_{n \rightarrow +\infty} \frac{(\log \log n)^2}{\log n} &= \lim_{n \rightarrow +\infty} \frac{2(\log \log n) \frac{1}{\log n} \frac{1}{n}}{\frac{1}{n}} = \lim_{n \rightarrow +\infty} \frac{2(\log \log n)}{\log n} \\ &= \lim_{n \rightarrow +\infty} \frac{\frac{1}{\log n} \frac{1}{n}}{\frac{1}{n}} = \lim_{n \rightarrow +\infty} \frac{1}{\log n} = 0 \end{aligned}$$

(vi) $n + \log n = \Theta(n + (\log n)^2)$

Proof: (Use L'Hôpital's rule)

$$\begin{aligned} \lim_{n \rightarrow +\infty} \frac{n + \log n}{n + (\log n)^2} &= \lim_{n \rightarrow +\infty} \frac{1 + \frac{1}{n}}{1 + 2(\log n) \frac{1}{n}} = \lim_{n \rightarrow +\infty} \frac{n + 1}{n + 2 \log n} \\ &= \lim_{n \rightarrow +\infty} \frac{1}{1 + \frac{2}{n}} = 1 \end{aligned}$$

(vii) $\log(n!) = O(n \log n)$

Proof: Since $n \log n = \log(n^n)$, it's equivalent to compare $n!$ and n^n . Obviously, $n! = O(n^n)$.

3 In Between Functions

4 Bit Counter

n	# Total Flips
1	1
2	4
3	11
4	26

There're two stages when counting from 0 to $2^n - 1$ for a n-bit counter.

- Stage1 : when most-significant bit is 0
- Stage2 : when most-significant bit is 1

Suppose $f(n)$ represents the total number of flips need for n-bit long counter. Both in stage 1 and 2, $f(n-1)$ flips is needed to flip all bits to 1 except the most-significant bit. When switching between stage 1 and 2, another n flips is introduced. 1 flip to switch the most-significant bit from 0 to 1. $n-1$ flips needed to switch all other bits to 0. Therefore, we have the following recursive formula:

$$f(n) = 2f(n-1) + n$$

To solve this formula, we keep using it recursively, providing $f(1) = 1$ as base case:

$$\begin{aligned}
 f(n) &= 2f(n-1) + n \\
 &= 4f(n-2) + 2(n-1) + n \\
 &= 8f(n-3) + 4(n-2) + 2(n-1) + n \\
 &= 2^k f(n-k) + 2^{k-1}(n-k+1) + \dots + 2^2(n-2) + 2^1(n-1) + 2^0(n-0) \\
 &= 2^{n-1}f(1) + 2^{n-2}(2) + \dots + 2^2(n-2) + 2^1(n-1) + 2^0(n-0) \\
 &= 2^{n-1}(1) + 2^{n-2}(2) + \dots + 2^2(n-2) + 2^1(n-1) + 2^0(n-0)
 \end{aligned} \tag{1}$$

Then we multiply it by 2:

$$2f(n) = 2^n(1) + 2^{n-1}(2) + \dots + 2^3(n-2) + 2^2(n-1) + 2^1(n-0) \tag{2}$$

By (2) - (1), we have

$$\begin{aligned}
 2f(n) - f(n) &= 2^n + 2^{n-1} + \dots + 2^1 - n \\
 f(n) &= \frac{2(1-2^n)}{1-2} - n \\
 &= 2^{n+1} - n - 2 \\
 &= \Theta(2^n)
 \end{aligned}$$

5 Recurrence Relations

Master theorem:

Suppose $a, b, c \in \mathbb{R}^+$, $b > 1$ and $T[1] = 1$

Given

$$T[n] = aT\left[\frac{n}{b}\right] + O(n^c),$$

We have

$$T[n] = \begin{cases} O(n^{\log_b a}) & c < \log_b a \\ O(n^c \log_2 n) & c = \log_b a \\ O(n^c) & c > \log_b a \end{cases}$$

(a) a

6 Hadamard matrices

(a)

$$H_0 = [1]$$

$$H_1 = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$H_2 = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix}$$

(b)

$$H_2 v = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ -1 \\ -1 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 4 \end{bmatrix}$$

(c)

$$\begin{aligned} u_1 &= H_1(v_1 + v_2) \\ &= \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \left(\begin{bmatrix} 1 \\ -1 \end{bmatrix} + \begin{bmatrix} -1 \\ 1 \end{bmatrix} \right) \\ &= \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \end{bmatrix} \\ &= \begin{bmatrix} 0 \\ 0 \end{bmatrix} \end{aligned}$$

$$\begin{aligned} u_2 &= H_1(v_1 - v_2) \\ &= \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \left(\begin{bmatrix} 1 \\ -1 \end{bmatrix} - \begin{bmatrix} -1 \\ 1 \end{bmatrix} \right) \\ &= \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 2 \\ -2 \end{bmatrix} \\ &= \begin{bmatrix} 0 \\ 4 \end{bmatrix} \end{aligned}$$

$$u = \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 4 \end{bmatrix}$$

u is identical to $H_2 v$

(d)

$$\begin{aligned} H_k v &= \begin{bmatrix} H_{k-1} & H_{k-1} \\ H_{k-1} & -H_{k-1} \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \end{bmatrix} \\ &= \begin{bmatrix} H_{k-1}v_1 + H_{k-1}v_2 & H_{k-1}v_1 + H_{k-1}v_2 \\ H_{k-1}v_1 + H_{k-1}v_2 & -H_{k-1}v_1 - H_{k-1}v_2 \end{bmatrix} \\ &= \begin{bmatrix} H_{k-1}(v_1 + v_2) & H_{k-1}(v_1 + v_2) \\ H_{k-1}(v_1 + v_2) & -H_{k-1}(v_1 + v_2) \end{bmatrix} \end{aligned}$$