CS170–Spring 2019 — Homework 2 Solutions

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1. Study Group

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2. Asymptotic Complexity Comparisons

- (a) $3 \le 7 \le 2 \le 5 \le 4 \le 9 \le 8 \le 6 \le 1$
- (b) (i) $\log_3 n = \Theta(\log_4 n)$ Proof:

$$\log_3 n = \log_{(4^{\log_4 3})} n = (\frac{1}{\log_4 3}) \log_4 n = \Theta(\log_4 n)$$

(ii) $n\log(n^4) = O(n^2\log(n^3))$

Proof:

$$\lim_{n \to +\infty} \frac{n^2 \log(n^3)}{n \log(n^4)} = \lim_{n \to +\infty} \frac{3n^2 \log n}{4n \log n} = \lim_{n \to +\infty} \frac{3}{4}n = \infty$$

(iii) $\sqrt{n} = \Omega((\log n)^3)$

Proof: (Use L'Hôpital's rule)

$$\lim_{n \to +\infty} \frac{\sqrt{n}}{(\log n)^3} = \lim_{n \to +\infty} \frac{\frac{1}{2\sqrt{n}}}{3(\log n)^2 \frac{1}{n}} = \lim_{n \to +\infty} \frac{\sqrt{n}}{(\log n)^2}$$
$$= \lim_{n \to +\infty} \frac{\frac{1}{2\sqrt{n}}}{2(\log n) \frac{1}{n}} = \lim_{n \to +\infty} \frac{\sqrt{n}}{\log n}$$
$$= \lim_{n \to +\infty} \frac{\frac{1}{2\sqrt{n}}}{\frac{1}{n}} = \lim_{n \to +\infty} \sqrt{n} = \infty$$

(iv) $2^n = \Theta(2^{n+1})$ Proof:

$$2^{n+1} = 2 * 2^n = \Theta(2^n)$$

(v) $n = \Omega((\log n)^{\log \log n})$

Proof:(Use L'Hôpital's rule)

Take the logarithm of both functions

$$\log(\log n)^{\log\log n} = (\log\log n)^2$$

Therefore, it is equivalent to compare $(\log\log n)^2$ and $\log n$

$$\lim_{n \to +\infty} \frac{(\log \log n)^2}{\log n} = \lim_{n \to +\infty} \frac{2(\log \log n) \frac{1}{\log n} \frac{1}{n}}{\frac{1}{n}} = \lim_{n \to +\infty} \frac{2(\log \log n)}{\log n}$$
$$= \lim_{n \to +\infty} \frac{\frac{1}{\log n} \frac{1}{n}}{\frac{1}{n}} = \lim_{n \to +\infty} \frac{1}{\log n} = 0$$

(vi) $n + \log n = \Theta(n + (\log n)^2)$

Proof:(Use L'Hôpital's rule)

$$\lim_{n \to +\infty} \frac{n + \log n}{n + (\log n)^2} = \lim_{n \to +\infty} \frac{1 + \frac{1}{n}}{1 + 2(\log n)\frac{1}{n}} = \lim_{n \to +\infty} \frac{n + 1}{n + 2\log n}$$
$$= \lim_{n \to +\infty} \frac{1}{1 + \frac{2}{n}} = 1$$

(vii) $\log(n!) = O(n \log n)$

Proof: Since $n \log n = \log(n^n)$, it's equivalent to compare n! and n^n . Obviously, $n! = O(n^n)$.

3. Understanding Academic Dishonesty

- (a) Not OK
- (b) Not OK
- (c) Not OK
- (d) OK